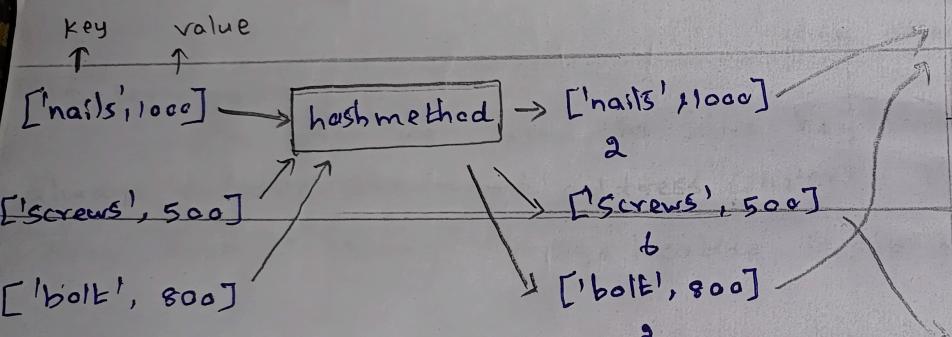


HashTable.

built in hash table python dictionary.
To demonstrate hash table Create a hardware
little store...



0	
1	
2	[['nails', 1000], ['bolt', 800]]
3	
4	
5	
6	[['screws', 500]]

Here we ran through a key through the hash method and we got our address space (0 to 6) using that address as a index we can store the key value pairs.

hash itself.

i) It's nonway, ii) Deterministic

ii) It's one way.

'key' → hash → address
function

we use key to get address through the
hash method we can't reverse that.
using address - through hash can't get the key

iii) Deterministic.

→ when we pass the same key through hash function, it will
always produce the same address (index). This ensures the given key will
always map the same storage location in the table.

Hashable - collisions.

[[milk', 100],
[nuts', 120]]

→ when we put a multiple
key value pairs in same address

It's called collision.

Technique (few deal) with collisions.

1. put them in the same address

→ Separate chaining

2. $\llbracket \text{key}, \text{value} \rrbracket, \llbracket \text{key}, \text{value} \rrbracket \rrbracket$

2. if we already have a value in that address we need to go down until find a empty address and put the key value.

Linear probing
(form of open addressing)

3. another way to do separate chaining,
instead having a nested loop we can create
a linked list.



constructor.

class HashTable:

def __init__(self, size=7):

self.data-map = [None] * size

def __hash__(self, key):

my-hash = 0

for letter in key:

my-hash = (my-hash + ord(letter) * 23)

% len(self.data-map)

return my-hash

def print_table(self):

for i, value in enumerate(self.data-map):

print(f'{i}:{value}')

→ 0: None

1: None

2: None

3: None

default size.

represents,

e	None
1	None
2	None
3	None
4	None
5	None
6	None

hash method.

ord(letter) → short for ordinal it will

give a ascii number for each char.

23 → prime number. helps distribute

values more evenly across table.

~.7 → when we divide any integer

using 7 the remainder always 0 and it

included and between .

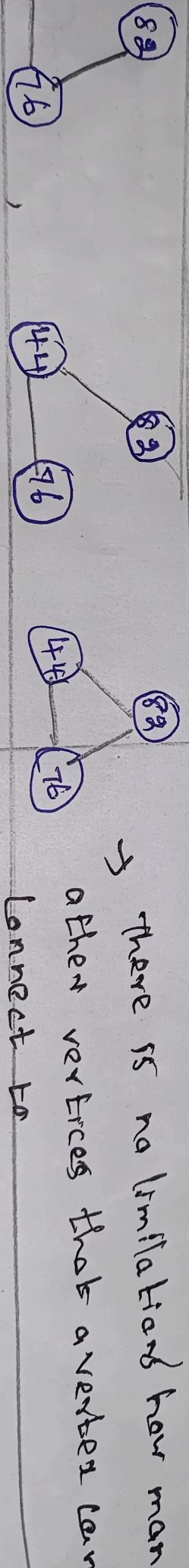
That exactly the address space.

Graph.

linked list are the form of a tree, Tree is a form of a graph.
therefore, a linked list is a form of a graph with the limitations that
they can only point to one other node.

(44) ← vertex (or) proper way to say vertex here).
node

(44) → plural is vertices between the vertices we have
a edge or connection (proper to say edge)



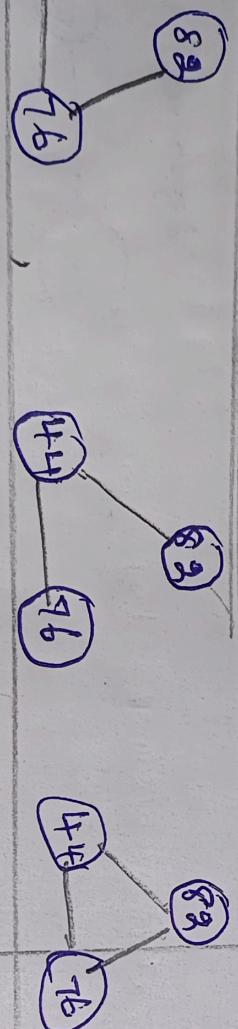
→ there is no limitation how many
other vertices that a vertex can
connect to

Graph.

linked list are the form of a tree, Tree is a form of a graph.
Therefore, a linked list is a form of a graph with the limitations that
they can only point to one other node.

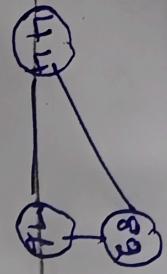
(44) ← vertex (or) (proper way to say vertex here).
node

(44) (76) → plural is vertices between the vertices we have
a edge or connection (proper to say edge)



→ there is no limitation how many
other vertices that a vertex can
connect to

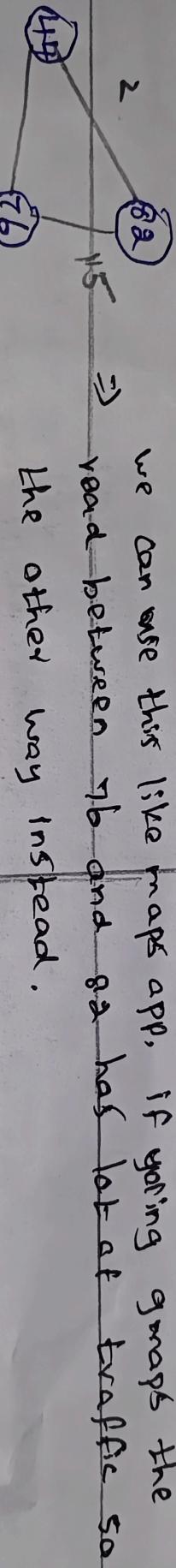
upper left one says hours of classes (less or more at a day).



assume we need to go - 76 to 82.
⇒ 11 76 to 82 → 1 (correct route).

ii) 76 to 44 and 44 to 82 → 2.

But graph, we can have not always but can have a weighted edges.

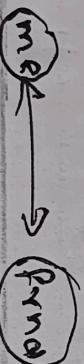


we can use this like maps app, if youing ignores the road between 76 and 82 has lot of traffic so the other way instead.

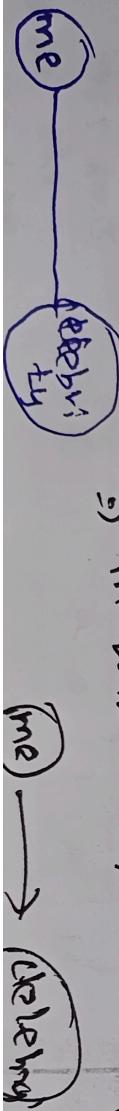
iii) ⇒ The edges can be weighted or non weighted.

another concept.

⇒ in fb we both friend follows each other (bidirectional relationship)

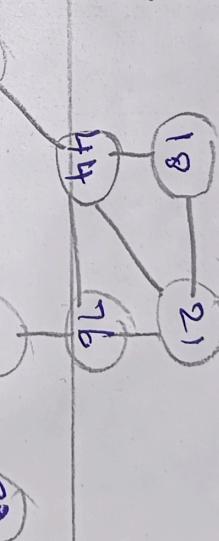


⇒ in twitter i follow a person, but they dont (unidirectional).



So the edges can be weighted (or nonweighted) they can be directional or bidirectional.

When we think about graph we can look something like this



Adjacency matrix.

2 way to locally represent graph.

1) is a. adjacency list:

This axis is the element C has an edge with.



Shamefully looking at A has 4 edges and e. so put the edge (e.g.) 1 in the box now B has edges with A and C and ...

A could point to A & B and point to B ...
so we always gonna have the 45 degree line of zeros like

	A	B	C	D	E
A	0	1	0	0	0
B	1	0	1	0	0
C	0	1	0	1	0
D	0	0	1	0	1
E	1	0	1	0	0

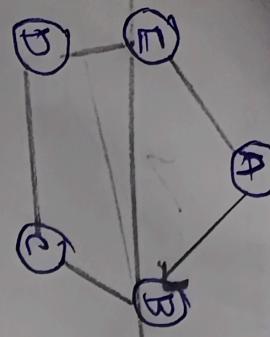
And if we have bidirectional matrix like here always have mirror image on each side of the diagonal.

This axis represent the vertex.

20 Aug

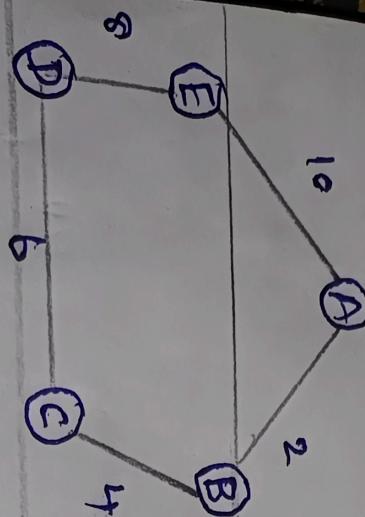
Graph Theory (Discrete Mathematics) or Probability & Statistics

here now it's not Bidirectional
A pointing B ; but B not point



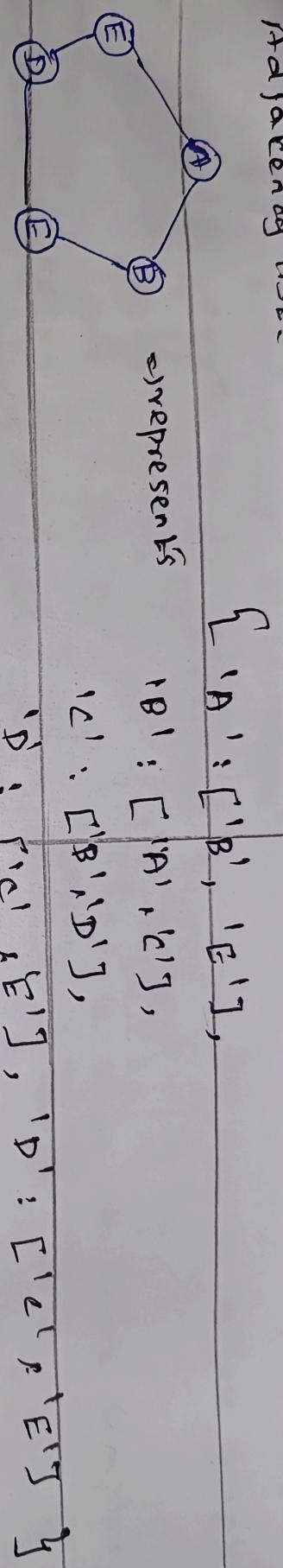
A.

weighted.



	A	B	C	D	E
A	0	2	0	0	10
B	2	0	4	0	0
C	4	0	6	0	8
D	0	6	0	8	0
E	10	0	0	1	0

Adjacency list.



Big O graph.

Adding item.

i) Adjacency matrix $\Rightarrow O(N^2)$

ii) Adjacency list $\Rightarrow O(1)$

Adding Edge : i) AM $\Rightarrow O(1)$

ii) AL $\Rightarrow O(1)$

removing edge : removing edge : number of edges.

i) AM $\Rightarrow O(1)$

ii) AL $\Rightarrow O(1) \Rightarrow O(n)$

remove vertex,

i) AM $\Rightarrow O(N^2)$

ii) AL $\Rightarrow O(M + |E|)$

matrix method not only storing problem for small graph assume a so its very inefficient in storage perspective

in list method we dont have to store zero its more simpler and more effective.

Add vertex.

{'A': []}

A

⇒

```
class Graph:  
    def __init__(self):  
        self.adj_list = {}  
    def print_graph(self):  
        for vertex in self.adj_list:  
            print(vertex, ':', self.adj_list[vertex])  
  
    def add_vertex(self, vertex):  
        if vertex not in self.adj_list.keys():  
            self.adj_list[vertex] = []  
        return True
```

also storing a lot not large million of vertices graph

```
my_graph = Graph()  
my_graph.add_vertex('A')  
my_graph.print_graph()
```

A: []