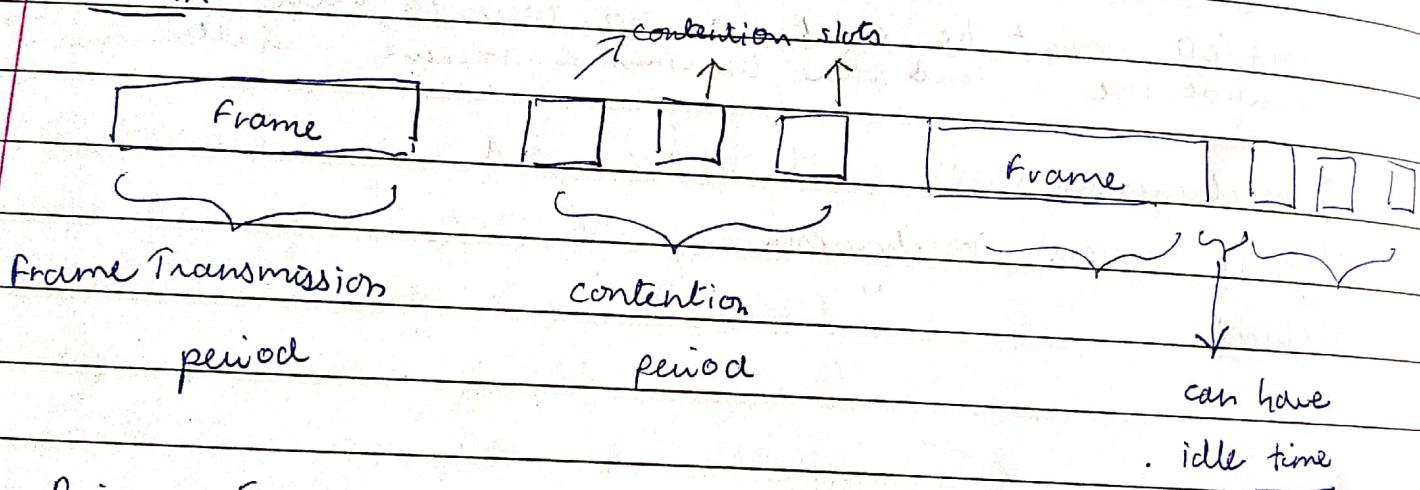


- 13/20 → ALOHA
- CSMA
- csma/cd - wired network
csma/ca - wireless network
- Contention free protocols
- Limited Contention Protocols
- Wavelength Division Multiplexing.

CSMA



Binary Exponential back off algo.

0, 1, 2, 3

0 - - - 7

0 - - - 15

↑

Max 10 attempts

0 — 1023

If can't transmit within this, gives up transmission.

ALOHA - no probability in time slots

CSMA - probability
time slots waiting time
is random

Contention Free Protocols

Polling

Polling Method

Select Method.

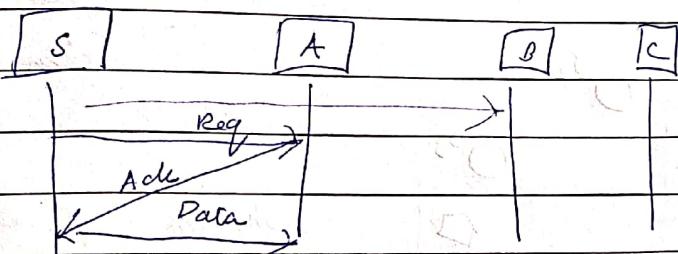
Signals are connected in daisy chain 

server wants to send to client (selecting a client to send to) - Select method.

Client wants to send to server, server must

select client - Polling method

Select Method:

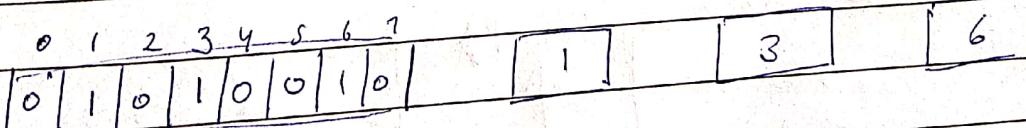


Polling Method :



Bit Map Method

e.g.



Contention slot - all devices indicate willingness to transmit by changing $0 \rightarrow 1$
 works on reservation

8 devices, all initially have 0.

To indicate that a particular dev. wants to transmit
 Transmit in Round-Robin fashion
 other dev. have to wait for 3 frame slot time.

Date: _____

D	1	2	3	4	5	6	7
1	0	1	0	1	1	0	1

1	3	6
---	---	---

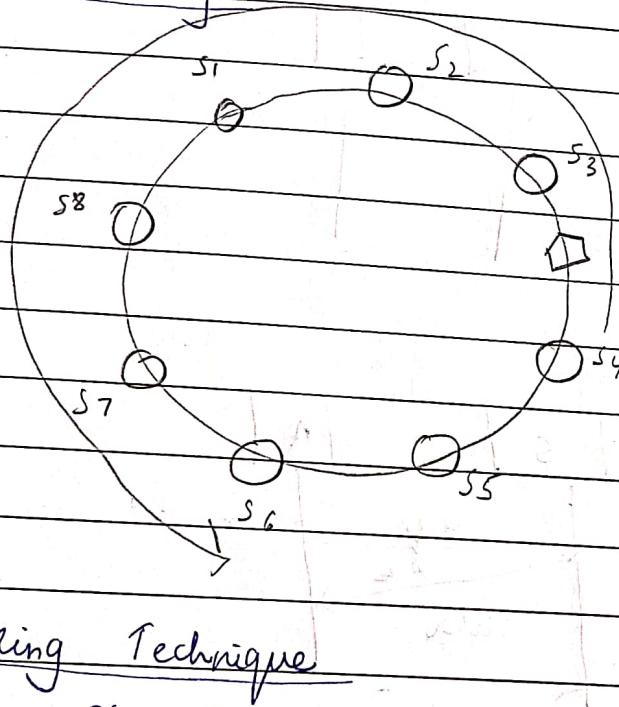
0	0	1	1	0	1	0	0	1
---	---	---	---	---	---	---	---	---

Reservation of slots is done by Bit Map

Very good for heavy load as starvation does not occur.
Not good for light load

Also commonly called Reservation Method.

Token Passing



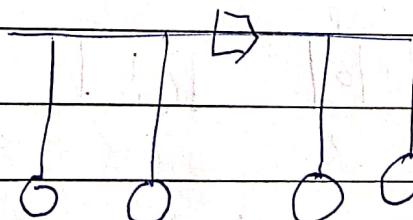
- Token Ring
- Token Bus

whoever captures token can transmit.

Token Ring Technique

Disadv: Starvation might occur as some might capture token each time.

Token Bus Technique

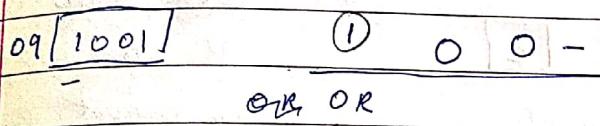
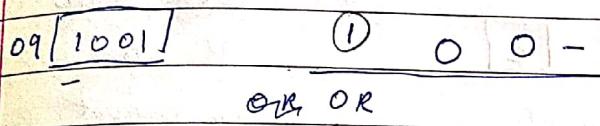
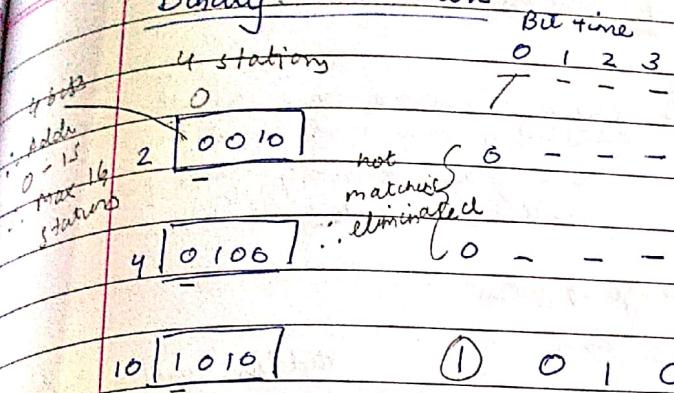


Whoever wants to transmit, captures the token.

If want delay & throughput Max - ALOHA
Contention Free

Theory +
Date: problems
for exam

Birdy Countdown



OR OR

① 0 1 0 ← address can transmit i.e. station 10

∴ Priority based ∵ largest always given.

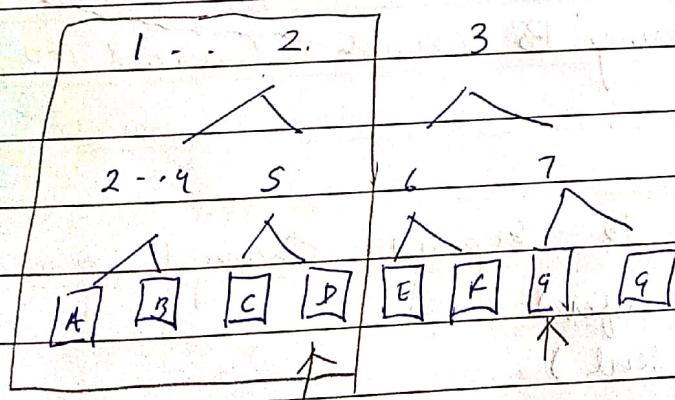
Limited Contention

Combines best features of Contention free & contention based

Adaptive Tree Walking Protocol - Divide & Conquer

Level 0 1

Systems are connected / arranged in a hierarchy.



If D & G want to transmit → collision
for next time slot only []
can compete if only D is there
→ go through

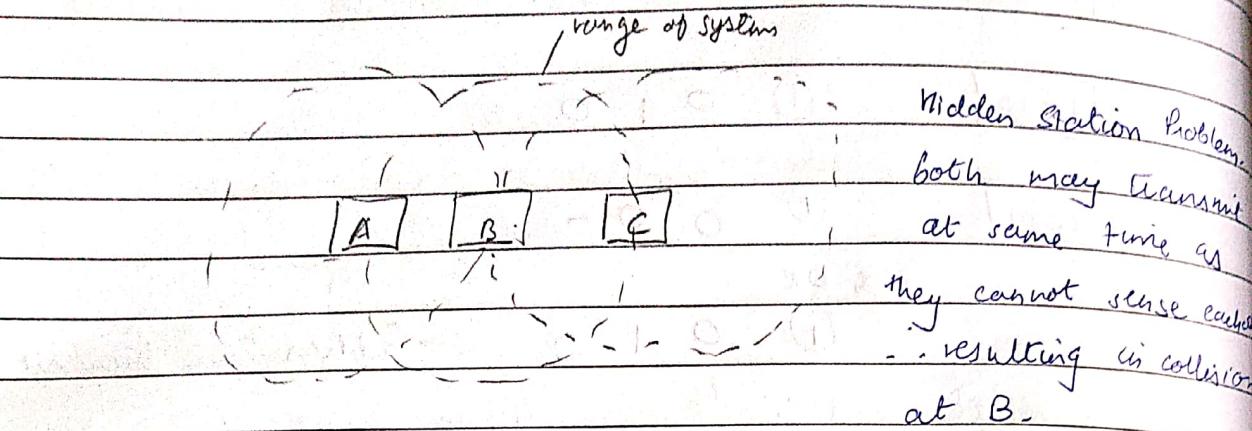
Light load situation - At time slot 0 everyone can transmit
if only 1, not no collision → go through.
If collision, they can compete for next time slot.

Contention Based Protocols Of Wireless Network.

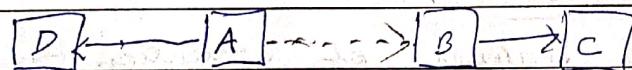
CSMA/CD cannot transmit & receive at same time as at time of transmission your own signals are strong hear you, receiving signals from other will be weak.

CSMA/CA used instead

MACAW is commonly used



Exposed Terminal Problem



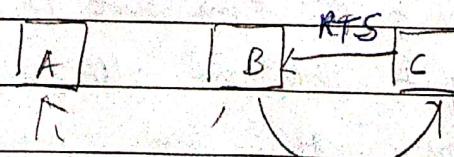
B is in the range of A, whatever A is transmitting to D, it reaches B also bc of broadcasting

Even though B was free to transmit, B assumes that transmission is occurring. B cannot transmit to C
∴ CSMA/CA

MACA Protocol solⁿ to Exposed Terminal Prob.

CTS (Request to transfer)

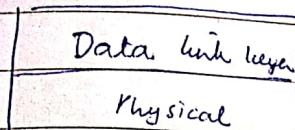
CTS (Clear to send)



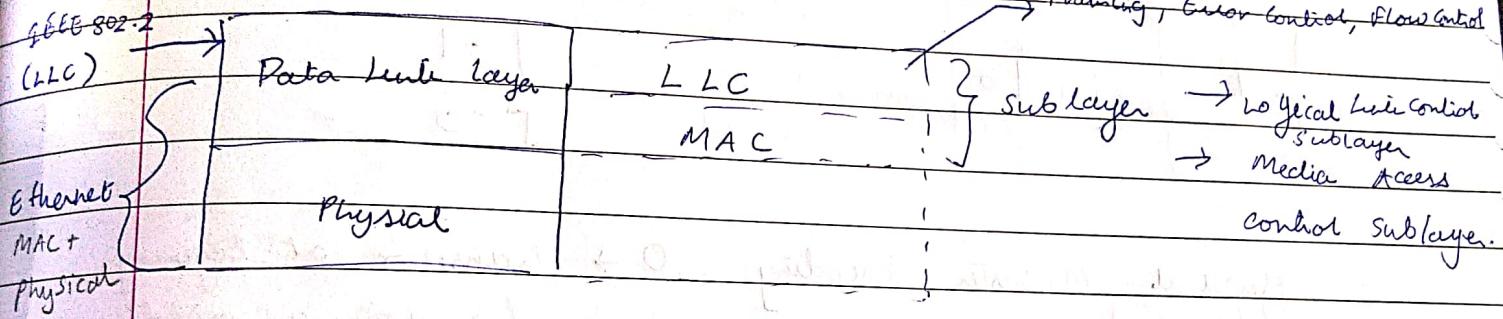
CTS means - If B is ready to transmit it gives back to A too in range CTS A cannot transfer to B

ETHERNET (LAN)

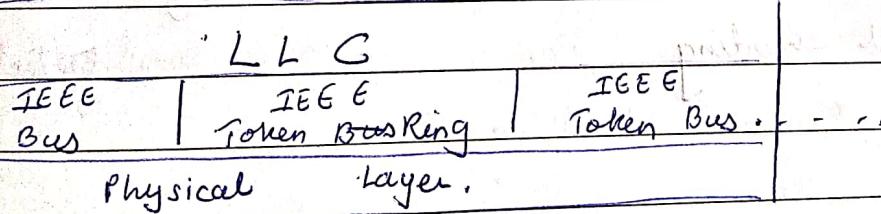
Layered Architecture



These 2 layers are sufficient for local network communication between 1 device to another device.



Ethernet follows IEEE std. → IEEE 802.3 ref to
combinatⁿ of physical layer +
part of DLL → MAC



→ Topology .

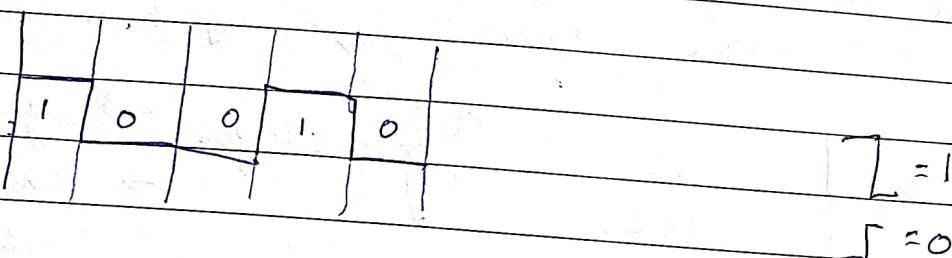
Sublayer of Physical layer.

Data encoding

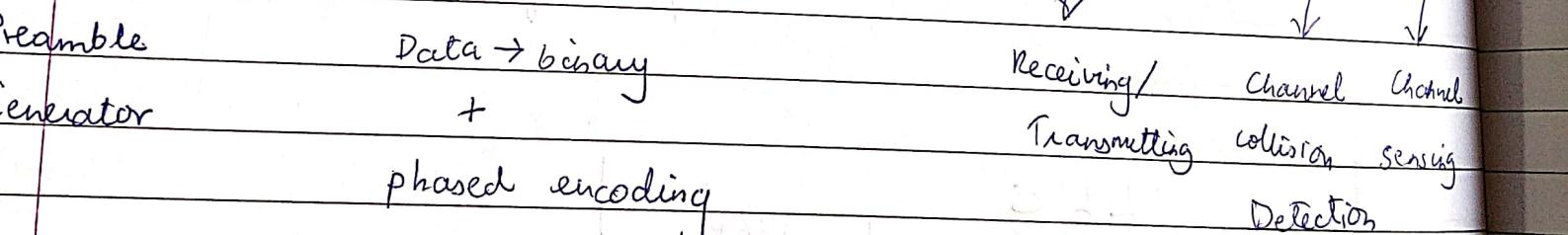
(Manchester encoding) & Differential Manchester

If ethernet is working at speed = 10MB

$$\therefore \text{Band rate} = 2 \times 10 = 20$$



Differential Manchester Encoding : 0 → transition at beginning
 1 → transition in middle



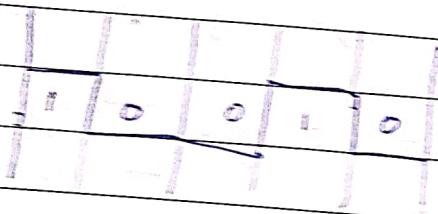
Sublayer of Physical layer.

↓
Barker encoding

(Polarized encoding) & Differential NRZ

↓ Channel coding
Fiber is working at speed = 10Mbps

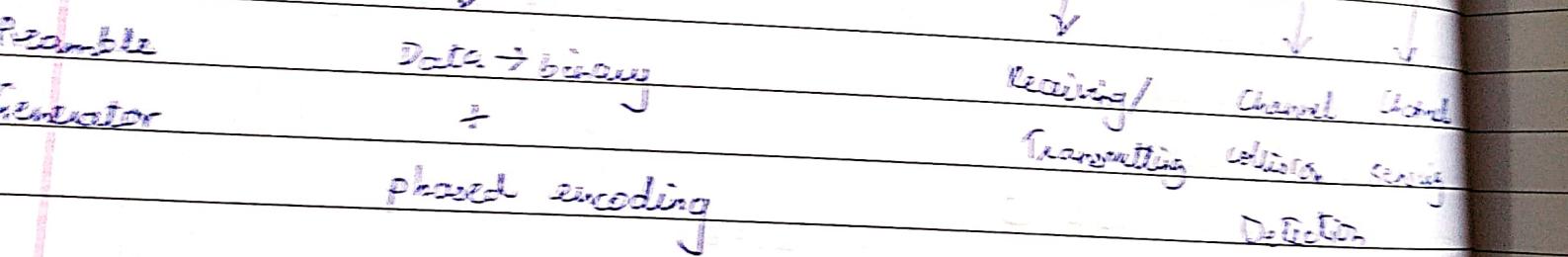
$$\therefore \text{Band width} = 2 \times 10^9 = 20$$



$$L=1$$

$$T=0$$

Differential Manchester Encoding : 0 → transition at beginning
 1 → transition in middle.



Ethernet diff

- length of segment
- length of frame
- topology

Page No.:

Date:

YOUNV

10 Mb standard

works on speed of
802.3

802.3i

802.3a

802.3j

802.3ak all → 200 Gbps | 400G Optical
(2017)

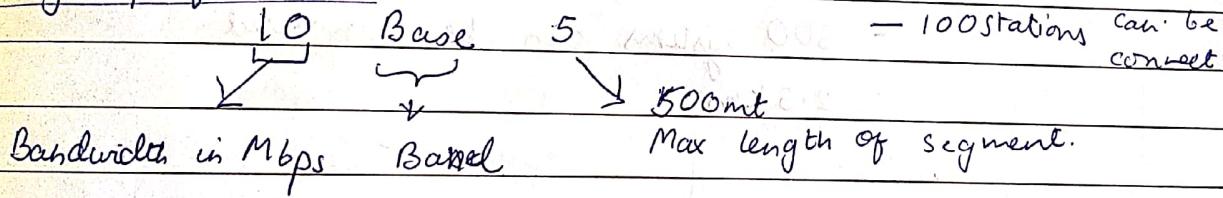
10 Base-T

10 Base-5

10 Base-2

10 Base-T

Design specifications



(Base Band: transmitting to only 1 freq

∴ Local area)

Co-axial cables are used. - commonly called thick co-axial cables. Thick Internet.

Restrictions: no of stations that can be connected

: length of tap

: dist b/w 2 stations

10 Base 2 - Bus topology
Thin co-axial cable

- 10 Mbps

- 185 mts customers can be connected (max).

Ethernet cliff

- length of segment
- length of frame
- Topology

10Mb standard

works on speed of

802.3

802.3i

802.3a

802.3j

802.3ak all → 200 Gbps / 400G Optical
(2017)

Fibre Cable

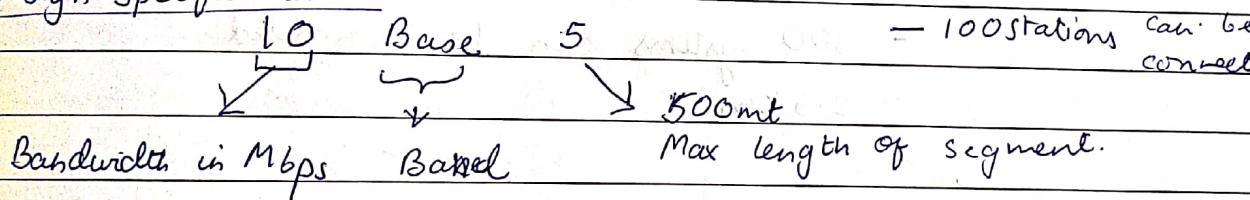
10 Base-T

10 Base-5

10 Base-2

10 Base-F

Design specifications



(Base Band: transmitting to only 1 freq

∴ Local area)

Co-axial cables are used. - commonly called Thick co-axial
cables ∵ Thick Internet.

Restrictions: no of stations that can be connected

: length of tap

: dist b/w 2 stations

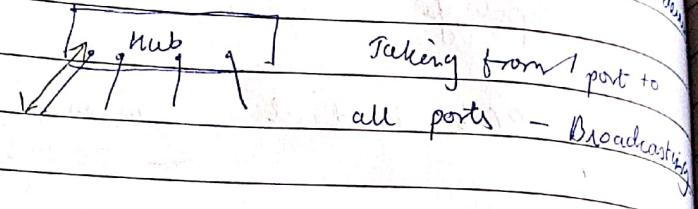
10 Base 2 - Bus topology
Thin co-axial cable

- 10 Mb/s

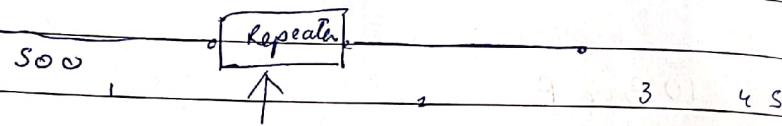
- 185 mts. items can be connected (max).

10 Base T

- 10 Mbps Base band signal Twisted pair cable
- Bus or star topology
- Hub is used (passive device used to connect devices)
- 100m

10 Base 5 - 5 - 4 - 3 Rule

To connect 500m cable



can use upto 4 repeaters to connect 5 links & only 3 links should be populated

- 300 systems can be connected.
- 2.5 km

10 Base F - star topology / Bus topology

- Fibre Optic cable
- Dist upto 2000m
- 1024 systems

Standard Ethernet - 10 Mbps

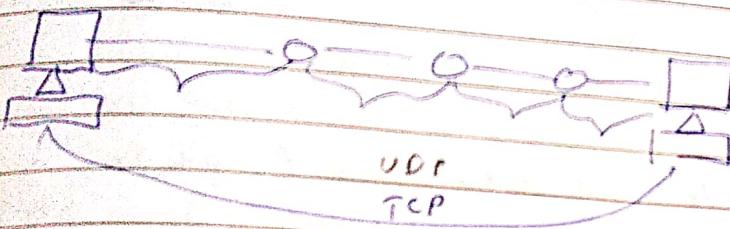
fast ethernet - 100 Mbps

10 G / Ethernet - 10 Gbps

1G

Bus is obsolete, either star or ring is usually used

Properties of Ethernet :-
connectionless, unreliable.



Connection is only possible betw end systems.

④ It creates a frame, & the frame has some format

frame	format	
7	2	6 6 2 46 - 1500 4 bytes
bytes	SFD	DA SA length/type Data CRC
7 bytes	Start	Dest Src 46 bytes
flag	addr addr 2 bytes	46 - 1500 bytes
	Deliminator	
	Created by Physical layer	DLL layer min size 64B
		Preamble, 72B

Initial std given by Robert Metcalf

std taken by DIX (Digital Index Xerox)



modified to IEEE 802.3

Preamble

7 byte info sent from src to dest

Alternating 0's & 1's

SFD

Alternating 0's & 1's

start & end of

indicates start of data

6/3/20

DA

48 bits

MAC address

- divided into 2 parts - manufacturer,
unicode

47: 5B : - - - - - 4 octets.

1 L → R

but within each Sent in rev order 7 before 4.

if all 6 bits 1 → Broadcast

if all octets end w/o 0 → Multicast address.

of odd-numbered
even-numbered

burned on NIC card.

CRC

CRC 32 - fixed part

for error control

Max size of frame = 1580

Length/type - which protocol

length of data

6/3/20

Date:

Yuvan

Cabling w.r.t. Ethernet (Standard)

max length of LAN = 200m
10 Base 5
10 Base 2
10 Base T
10 Base F

- } Physical characteristics -
- length of cable
- length of frame
- encoding

MAC Sublayer Ethernet (CSMA/CD)

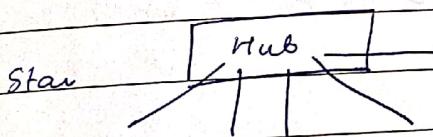
Fast Ethernet - 100 Mbps

Gigabit ethernet - 1000 Mbps - 1 Gbps
10 " " - 10.000 Mbps - 10 Gbps

2017 \Rightarrow IEEE 802.3bs \rightarrow 200 Gbps / 400 Gbps \rightarrow OFC

Bus

Ring 2 most commonly used.
Star



MAC Sublayer Ethernet (CSMA/CD)

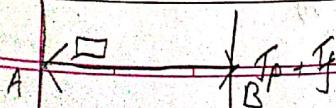
Packet arrived from W/W layer

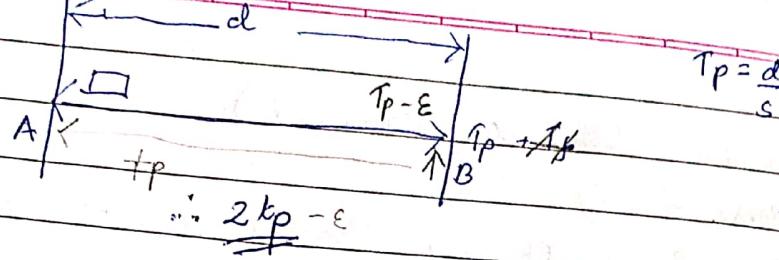
↓ Data encapsulation

Framing Addressing CRC
generate CSMA/CD

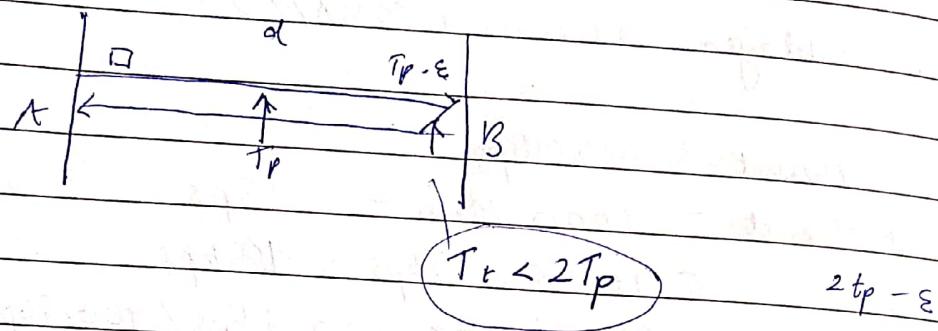
Channel Allocation

Collision Resolution
BEDA





restriction: $|T_t \neq 2T_p|$ to recognize collision



time used to transfer info/
min. frame size
= $2T_p$

If 100 Base, dist is by 10

MAC = 48 bits (address)
(6 bytes)

Last bit of 1st byte = 0 Unicast

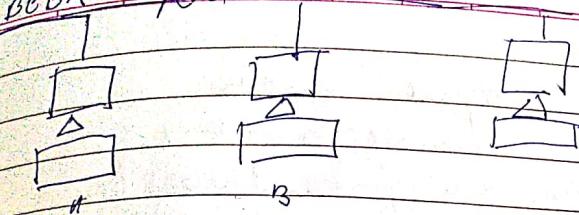
Last bit of 1st byte = 1 Broadcast

FF : FF : FF : FF : FF : FF → indicate broadcast

Ethernet makes use of 1-persistent CSMA.

If collision, it uses

BEBB - Binomial Exponential Back off Algo



$$2^k = \{0, 1, 2, 3\}$$

wait for slot time.

immediately
don't wait transmission after finding channel free

00 } collision
11 }

01 } No collision
10 }

$$\therefore \text{Probability (collision)} = 50\%$$

$$k=0 \dots 7 \quad 00 \\ 11$$

$$\therefore \text{Prob (collision)} = \left(\frac{2}{8}\right)$$

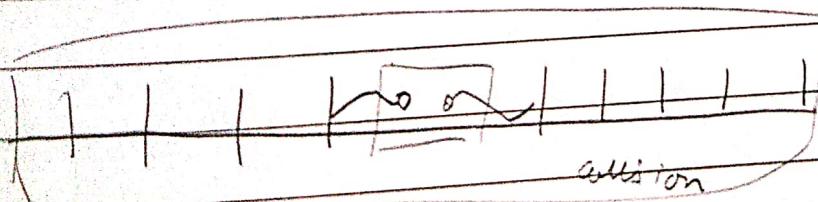
$$k=5$$

$$0 \dots 31$$

Bridged Ethernet & Switched Ethernet

Bridge

2-4 ports max.



collision

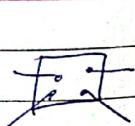
domain reduced

due to bridge

No. of collisions \downarrow

Same no. of broadcast domain

No. of domains \uparrow but belong to same D/W



Switch

multipoint switch
Communication is 1-to-1
For every port, there is a buffer.

Initially works as hub, when it has more info works 1-to-1

Initially no map table.

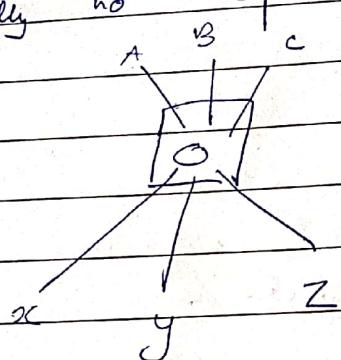
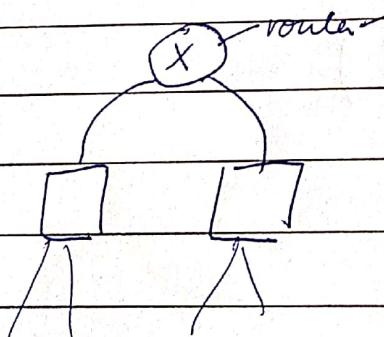


Table is constructed by learning mechanism.



reduces broadcast domain.

2 collision domains

2 broadcast domains

2 x P
Fibre

JT

Fast Ethernet

100 Base Tx - Twisted cable UTP or TCP → Cat.5
100 Base Fx - 150m
100 Base T4 - Twisted pair
4 cables each supports upto 25 mbps \Rightarrow 100 mbps
for K/T

Encoding scheme : 4B/5B & MLT-3

Switch - Full duplex

Base LX
Base SX
Base CX
Base TX

Baseband signal - whatever created by device is sent

encoding scheme : 8B/10T + NRZ-I

IEEE 802.3 65
~~65~~ 200/400

IEEE 802.1

802.2 LLC

802.3 Token Bus

802.4 - Token Bus

802.5 Token Ring

802.6 (MAN)

802.7

802.11 WiFi

802.15 - Bluetooth

802.16 - sensor nodes

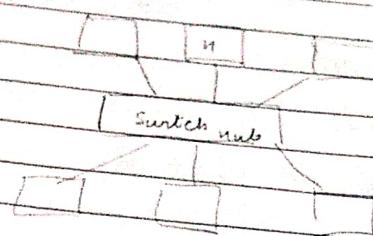
9/3/20

unidirectional, broadcast, wireless
diff stds of ethernet

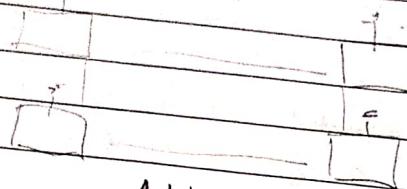
Wireless LAN (802.11)

PACKET TRACER

Page No.:
Date:



— Antenna



Ad hoc - devices

communicate w/o infrastructure.



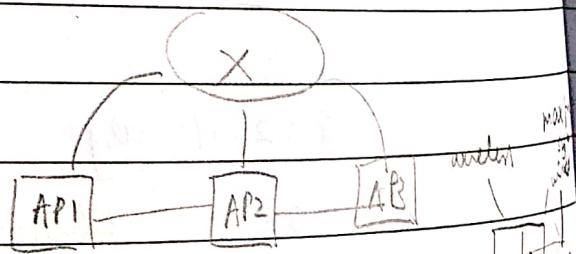
Wired LAN

Infrastructure Based Local N/w

Architecture

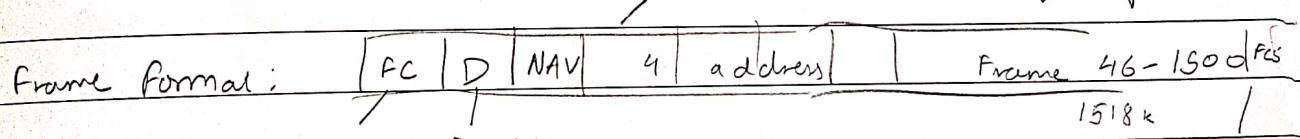
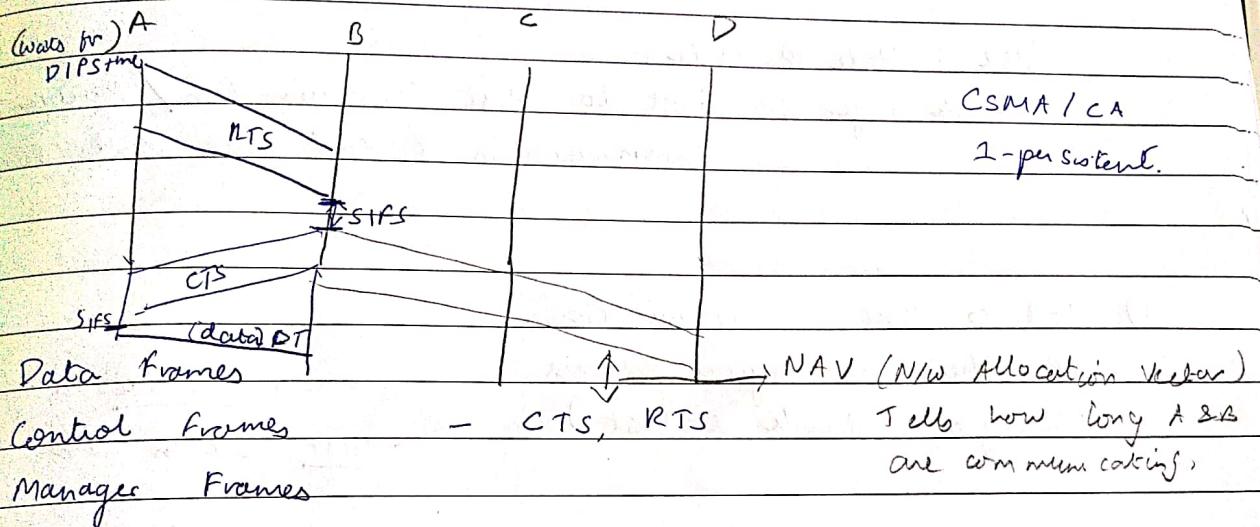
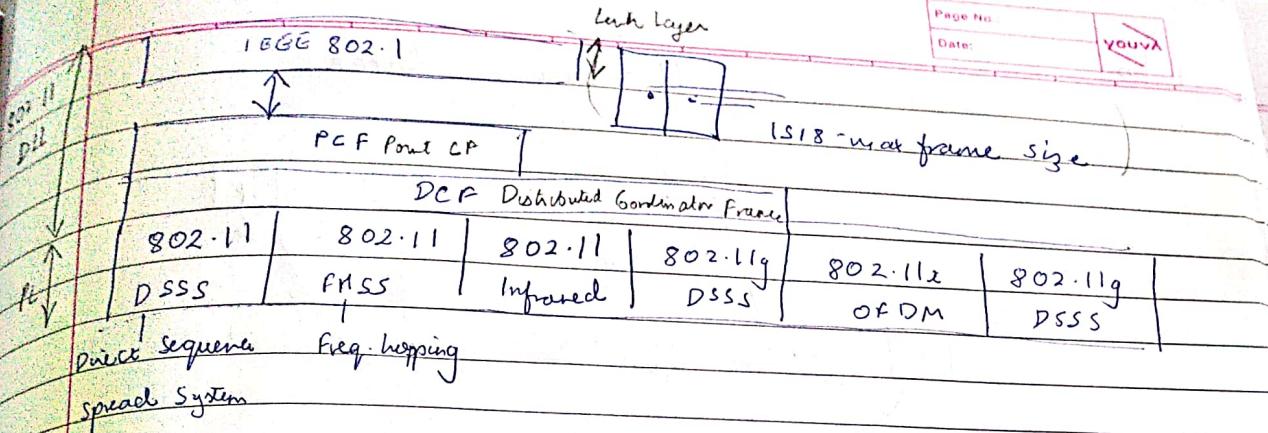
BSS → Basic Service Set

ESS → Extended Service Set.



Switch, Access Point - DLL

Access Point - Does conversion as wireless LAN uses diff protocol
from wired LAN.



2 bits: To DS From DS
of FC

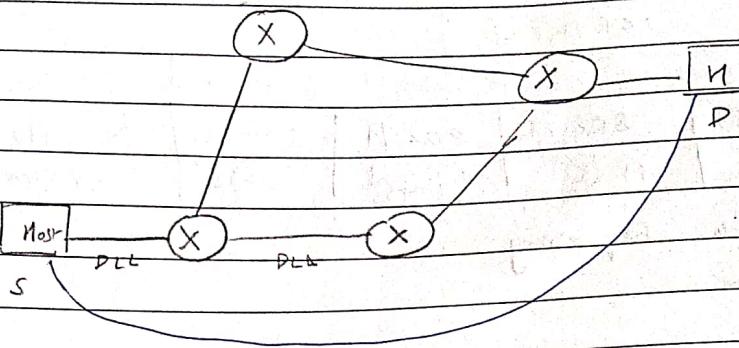
0	0	0
0	1	1
1	0	2

3 Src, dest, AP
4 Src, dest, AP, etc

Indicate how many addresses
should be there in the
frame.

* DLL

NETWORK LAYER

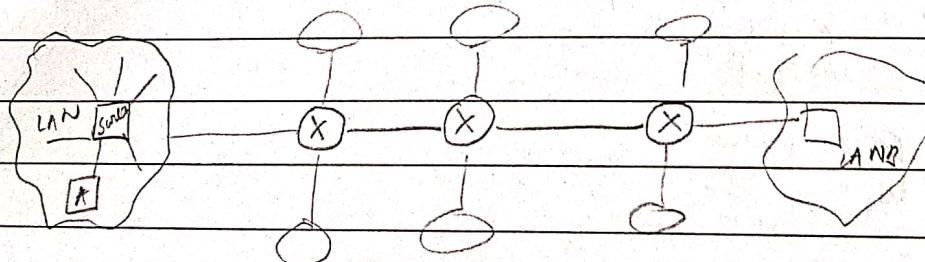


DLL : Node to Node

Network layer : Host to Host Communication / Machine to Machine Communication / end to end communication

(1) Host to Host communication -

(2) IP address or Logical Address.
(identifies network & host) ie which network to go to
Host within network
∴ requires Network id & Host id.



(3) Routing & Forwarding

Determine the route. (best path :- router is intelligent device)

Forwarding info to route.

- uses routing tables.

Path	Port
A	1
A	2
B	3
C	4
D	5

OS TF - Open Shortest Path Follow.
BGBGP

Router works in Layer 3 It contains s/w

(4) Providing services to transport layer.

Connectionless service \rightarrow IP

Connection-oriented \rightarrow ATM

Transport layer does sequencing.
Virtual circuit switching

(5) Fragmentation.

\because Buffer capacity might not be there.

\therefore Handshaking is reqd to find buffer size.

(6) Congestion Control.

20 Services:

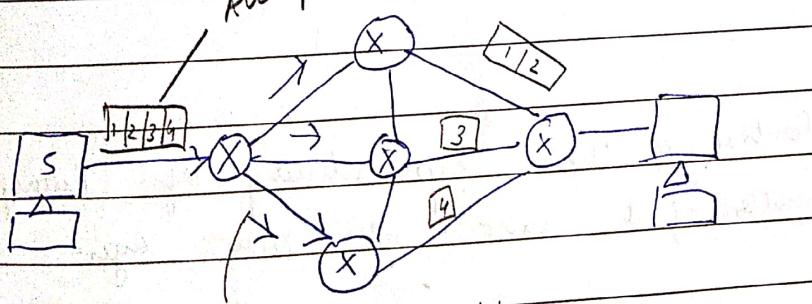
data gram service or Packet switching (connectionless)

Virtual circuit (VC) - connection oriented.

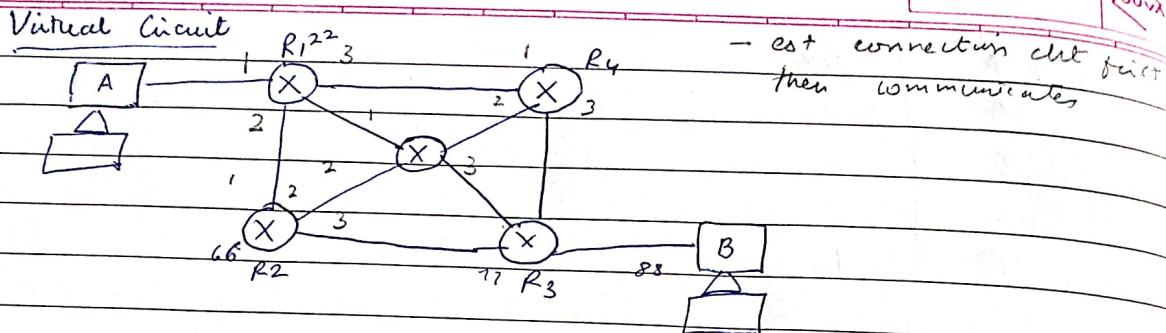
All packets do not follow same path

Read!

Dijkstra's Algo
Bellman Ford Algo



determining the path
is responsibility of
Network Layer.



- est connection at first
then communicates

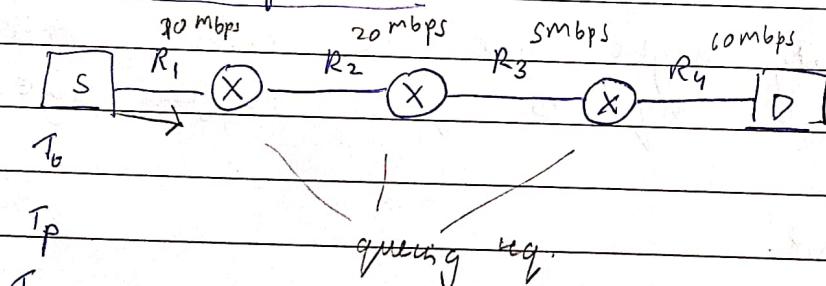
Input Port	Output Port
Label	Port
e.g. 22	1

↑ unknown during set up phase \therefore 66

VCI - Virtual Circ. Identifiers

Only 1st communication req IP address, all others req VCI

Network Performance



\because transmission from S, R1, R2, R3

$$(n+1)(T_p + T_b + T_{proc}) + n(T_{que})$$

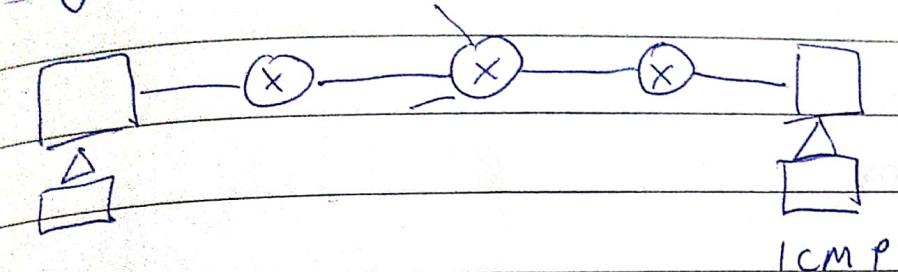
/ bandwidth

$$\text{Throughput} = \min(R_1, R_2, R_3, R_4)$$

Congestion Control is the responsibility of Transport Layer.

Some functionalities of it exist at N/w layer.

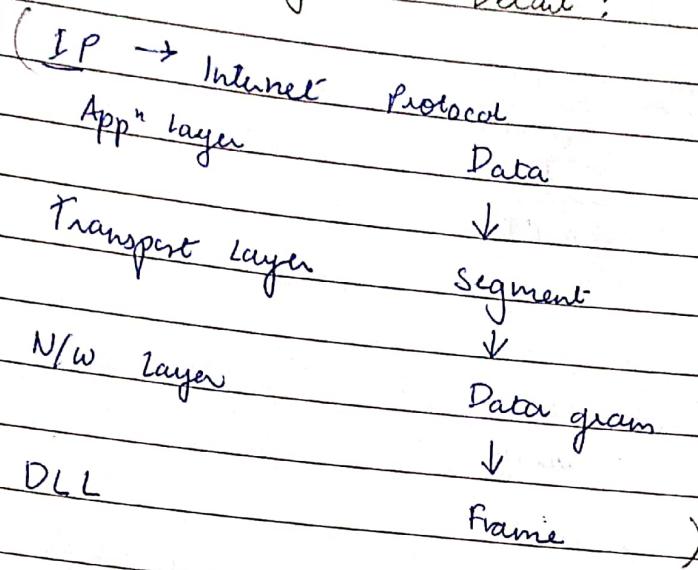
Congestion Control



ICMP

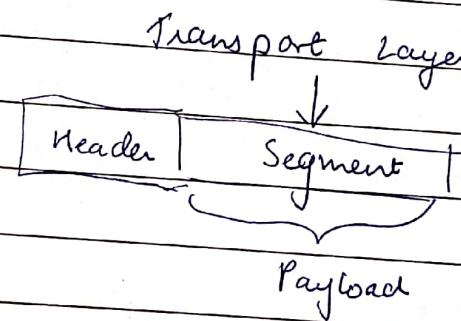
Chalk packets →
Back ←

Network Layer in Detail:

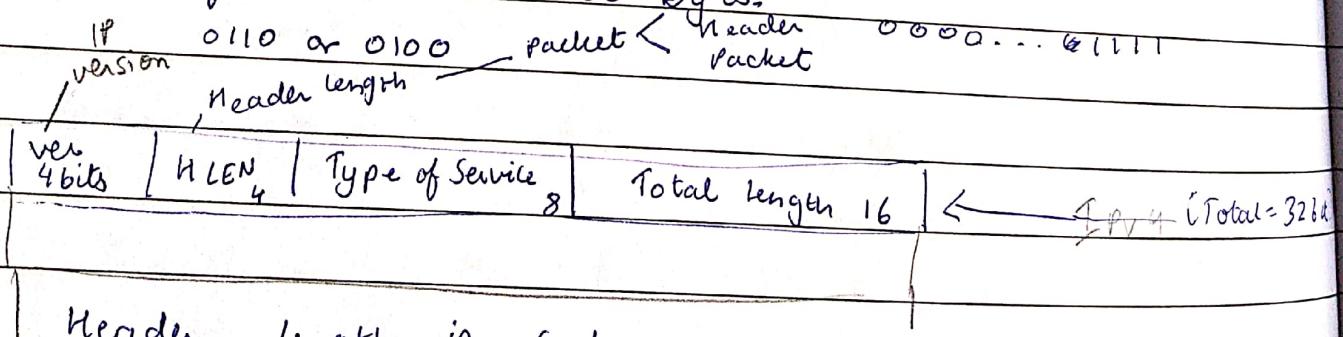


Packetisation

Packet



Size of header: 20 - 60 Bytes.

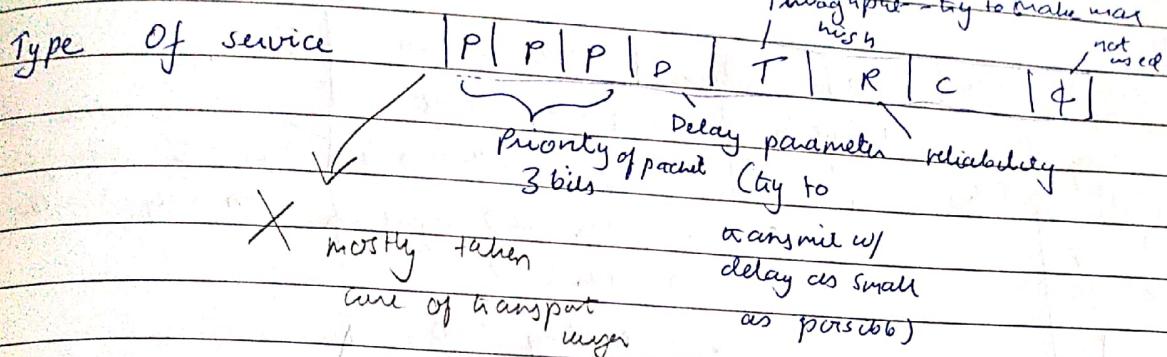


Header length is scaled to get 20 - 60B
Scale factor of 4

e.g. If HLEN = 5

$$= 5 \times 4 = 20$$

In N/W layer, when packets are sent, based on type of service priority can be given.



Not reliable - ∵ IP is connectionless protocol.

Total length of packet: $2^{16} B = 65535$

Max payload possible - 620,60

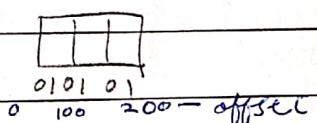
Max of DLL : 46 - 1500 (Payload)

Field	Via	MLEN	Type of Service	Total Length	To indicate no. of bytes preceding the packet.
header fields			Identification 16	Flag 3	Fragment offset (13) if $10 = 10 \times 8 = 80$ bytes given as offset.
IP header			TTL		

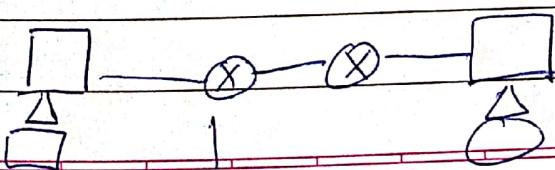
DLL supports : 1500 B

Transport layer gives 3000 B segment

Then N/W layer does fragmentation based on max size of payload.



∴ Each fragment has same fragment id, on receiving, N/W layer puts them together based on fragment id.



fragmentation may happen at each station dep on max size

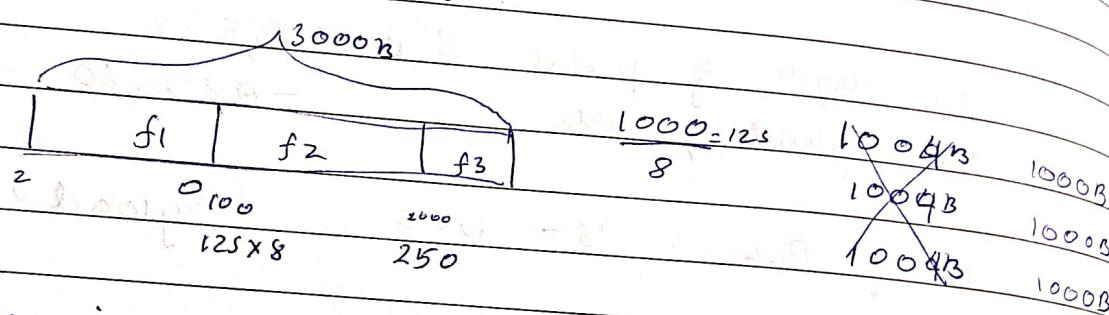
IP - Internet Protocol

Page No.
Date

As it moves from one host to another it drops 1 hop.		32		connectless / connection-oriented.	
ver / HLEN	Type of service	Total length			
Identification	Flags	frag offset			
TTL	Protocol	Header checksum 4			
8	8	IP Addrs (32) Logical addrs			
		DSI IP Addrs (32)			
		Options & Padding 40B			

It is calculated for every IP as TTL changes on every router.

Minimum Header size = 20
Max " " = 60

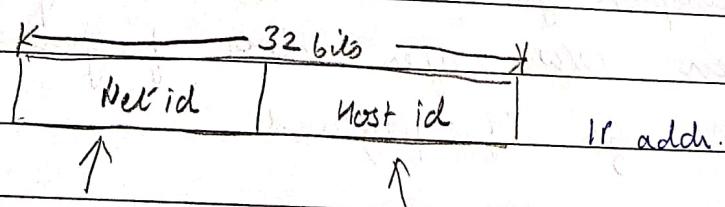


Offset is add of 1st fragment
offset of next is size of prev + offset.

Multiples of 8 should be fragmentation no.

IPv4 = 32 bits

IPv6 = 128 bits



To identify

n/w
no of bits
used depends
on type
of add

In that network,

host is identified w/

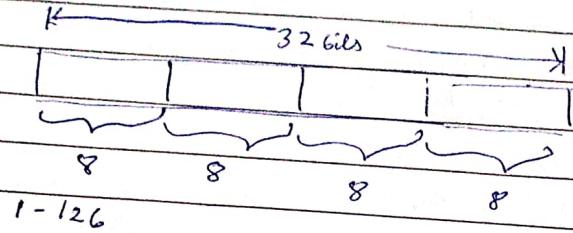
$$8 \times 4 = 32$$

1

Octet

0 - 255

A-class



If first bit = 0, it belongs to class A.

all 0s & 1s - not possible

$\begin{array}{ccccccc} 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ \hline 0 & 1 & 1 & 1 & 1 & 1 & 1 \end{array}$

N/w id indicates no of n/w possible.

No of n/w possible = 126

'A': No. of IP addresses not used $2^{31} - 1$ is fixed

126 n/w ids are possible.

2^{24} bits IP addresses computer are possible

$\begin{array}{ccccccc} 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ \hline 0 & 0 & 1 & 1 & 1 & 1 & 1 \\ \hline 0 & 0 & 2 & 2 & 2 & 2 & 1 \\ \vdots & & & & & & \\ 255 & 255 & 255 & 255 & 255 & 255 & 1 \end{array}$

$2^{24} - 2$ - ea IP addr is possible.

1. 0 0 0

? - n/w id

:
1. 255. 255. 255

{ These 2 are
not possible

{ checked

i.e. broadcast add. add.
i.e. ~~pro~~

56. 0 0 0 - n/w id

56. 0 0 0

:

56. 255. 255. 255 - broadcast id.

- Dividing big network into many small subnetwork
→ subnets.

(Subnet /) Default masks

255. 0. 0. 0

A mask used to identify the n/w idea of
a given given network

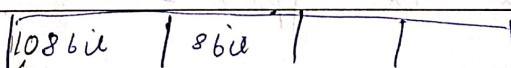
56. 26. 36. 72

255. 0. 0. 0

56. 0 0 0

B Class

128 - 191



no fd

no fd

first 2 bits are always 1, 0

or all 0's or all 1's

14 bits are variable

- 2^{14} possible networks

Subnets / $2^{16} - 2$ hosts

132. 166. 0. 0

132. 0. 0. 1

132. 0. 255. 255

Default mask 255. 255. 0. 0,
 " network position
 255 - broadcast address.

IP address

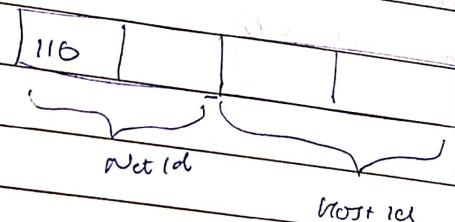
168. 38. 22. 48

168. 38. 6. 0

168. 38. 255. 254,

168. 38. 255. 255

C class



prefix : 110

2^{21}

possible networks available

each h/w supports $2^{8-2} = 256$ sys. per system.

192. 223

D class

Prefix : 1110

2^{20}

Class E: prefix : 1111

used / private or public
within organization

can be duplicated in an orgn eg. 10. 0. 0. 0 10. 0. 0. 0

Page No.:	
Date:	
YOU	

10.2555 . 255 . 255 - A class private add

172.16 0 0 - B "



172.31.00

192.168.0.0 - C "