Data Structures and Algorithms

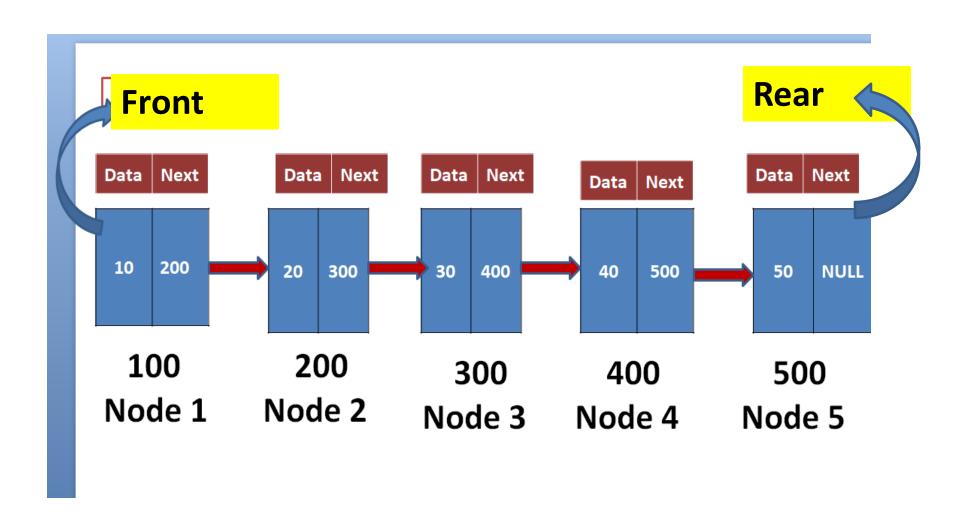
MSc. KhangVQH Faculty Of Information Technology Fall - 2022

Queue Linked List

Terminology: Front, Rear, Node, Data, Next, NULL

Operations: Enqueue(), Dequeue(), Display()

Queue using SLL



Content

- What is Queue Single linked list
- Queue working principle
- What is node?
- Node Structure.
- Operations on Queue SLL.

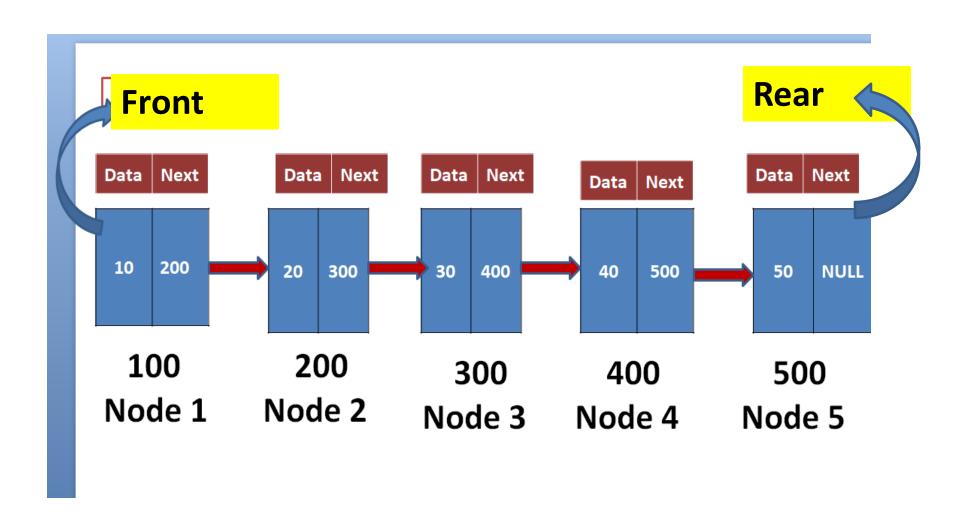
What is Queue Single linked list

- List: collection of number of elements
- SLL: SLL is linear Data Structure.
- It is also a collection of elements(nodes) but every element is linked with next element(node) by address.
- Queue SLL: Implement the queue data structure with single linked list concept.
- The Problem with queue using array is fixed size. But Queue using SLL is unlimited size.

What is Queue Single linked list

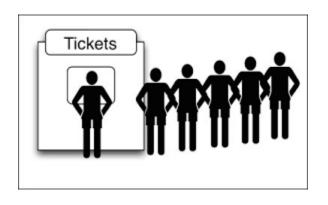
- Queue SLL uses two pointer fields
- 1) Front: it is used for delete the elements from the list
- 2) Rear: it is used to insert the elements into the list

Queue using SLL



Queue working Principle

 In queue data structure, the insertion and deletion operations are performed based on FIFO (First In First Out) principle.





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What is node?

- Every single element in a List is called "Node".
- Node contains two fields
- 1) Data filed-it holds data(element value)
- 2) Next field- it holds address of next node
- Every node has it's own address value in the memory

Data Next

Node 100

Node Structure

Operations on SLL.

- Enqueue(): it is used to insert the node at Rear.
- Dequeue(): it is used to delete the node at Front.
- Display(): it is used to display the nodes from the list.

Algorithm for Enqueue()

- enqueue(value) :It is used inserting an element
 into the Queue
- Step 1 :Create a NewNode with given value and set 'NewNode → next' to NULL.
- Step 2 :Check whether queue
 is Empty (rear == NULL && front==NULL)
- Step 3: If it is Empty then, set front = NewNode and rear = NewNode.
- Step 4: If it is Not Empty then,
- set rear → next = NewNode and rear = NewNode.

Connecting nodes by address

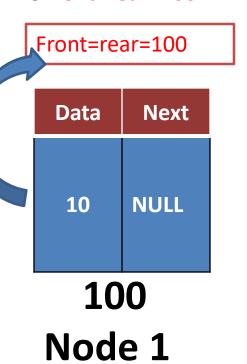
1)Before creating first node :: Assign front=rear=NULL

| Data | Next |
|------|------|
| 10 | NULL |

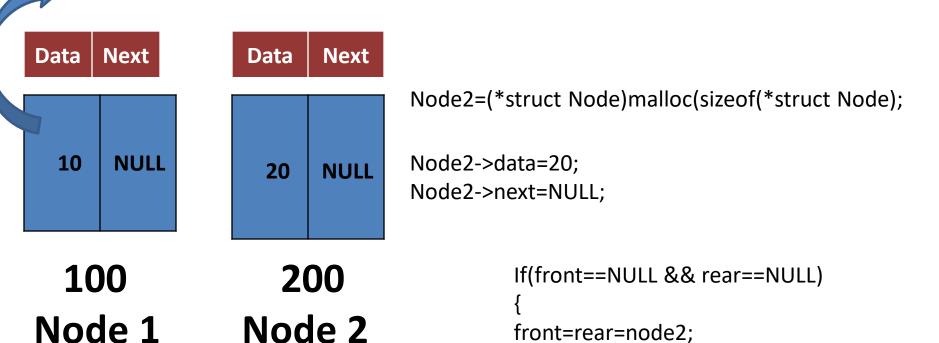
100 Node 1

(first node of list is called "Front" and last node is called "rear" in QSLL.

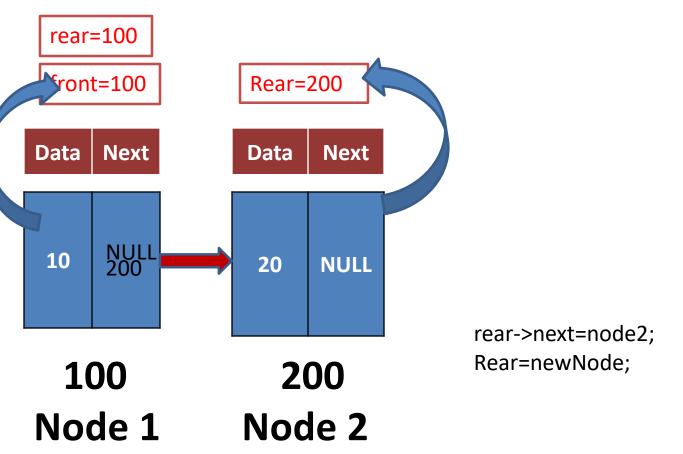
After creating first node Head=first node address i.e front=rear=100

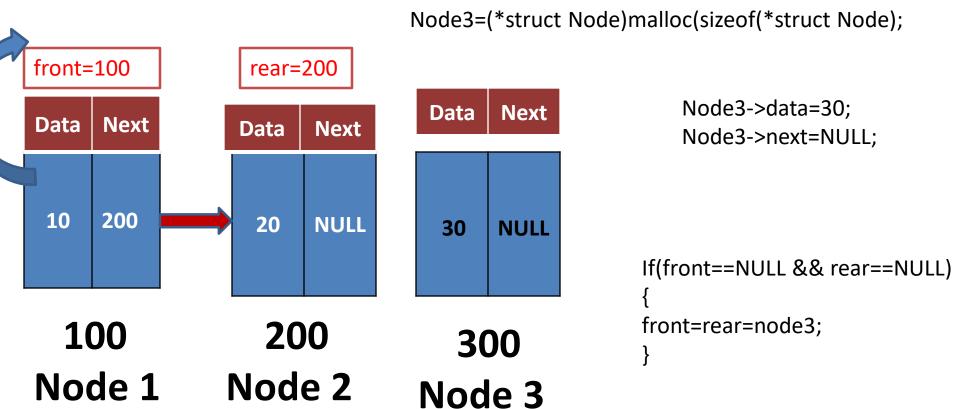


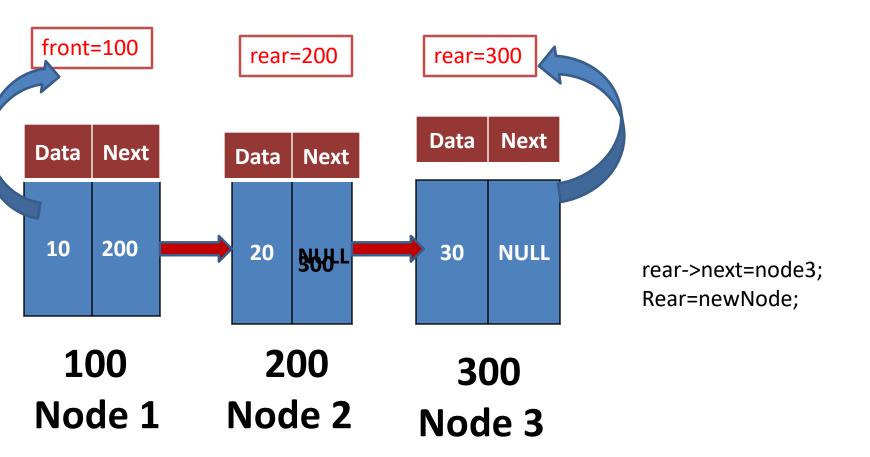
Front=rear=100

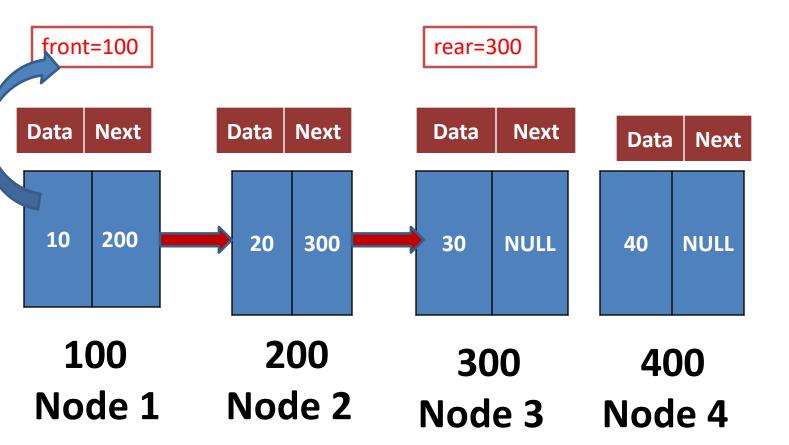


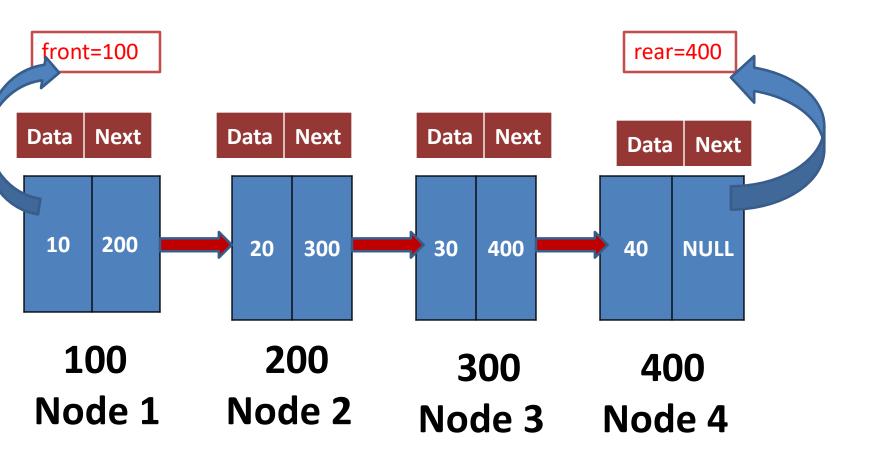
front=rear=node2;

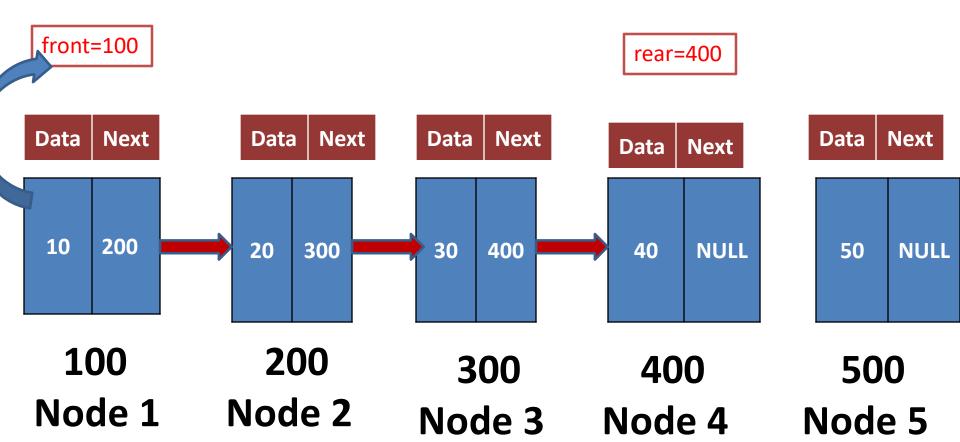


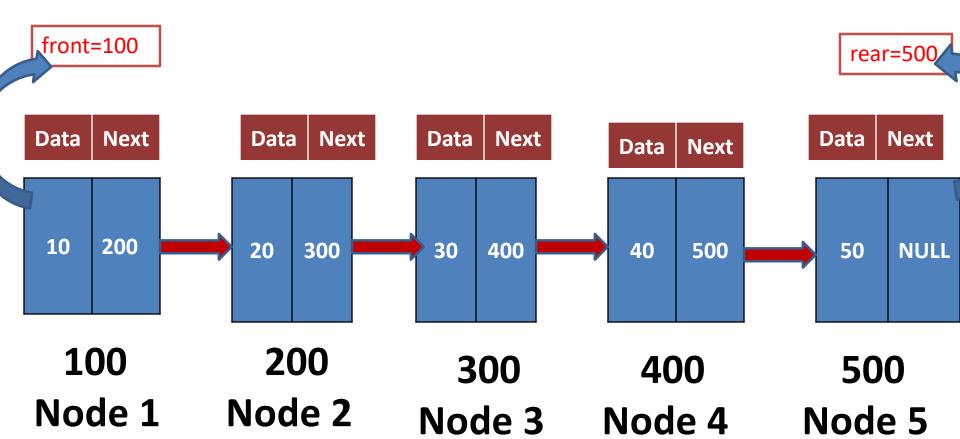






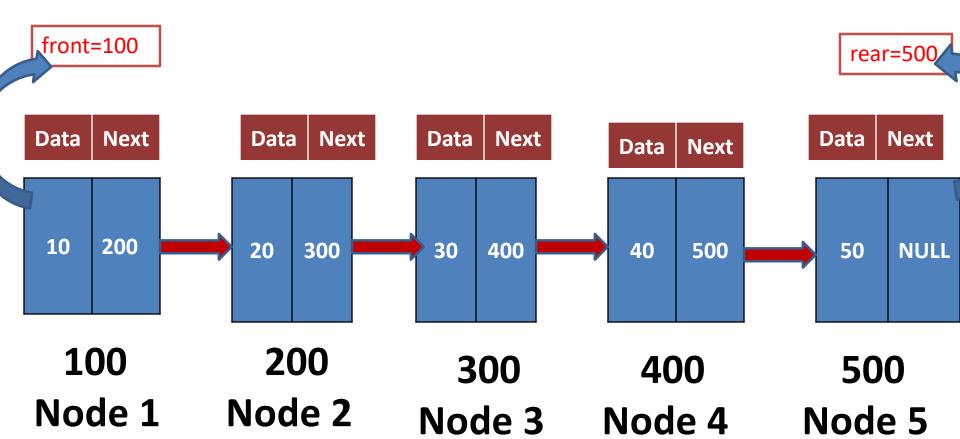


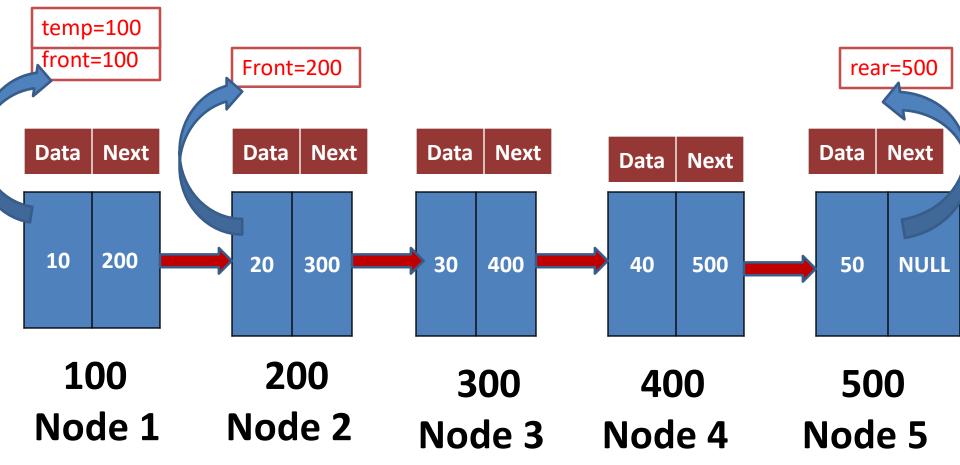


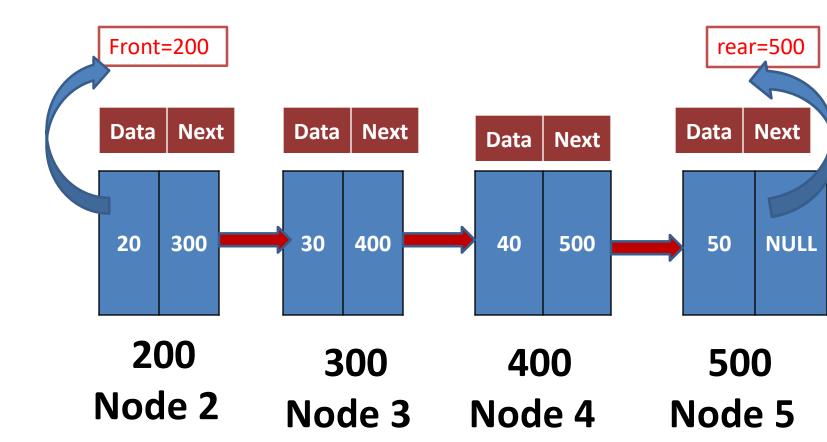


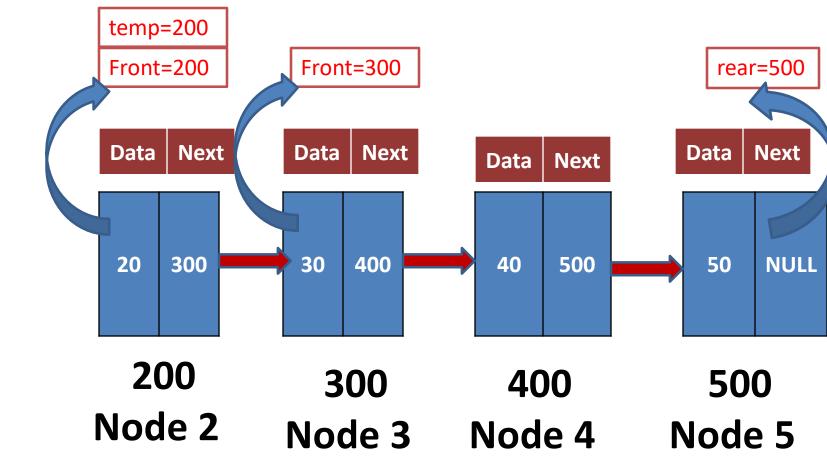
Algorithm for Dequeue()

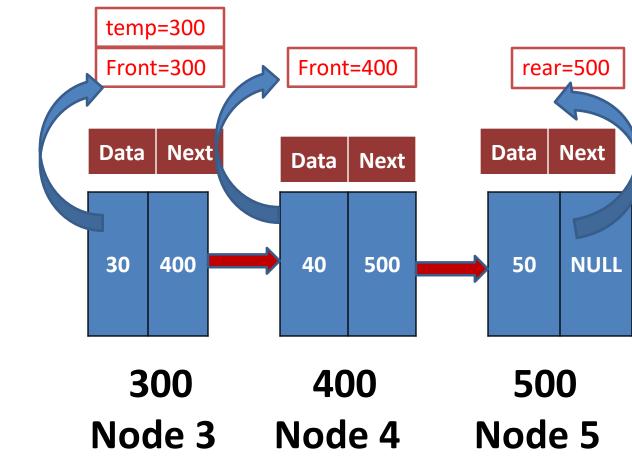
- dequeue(): Deleting an Element from Queue
- Step 1: Check whether queue is Empty
 - (front == NULL && rear==NULL).
- Step 2: If it is Empty, then display "Queue is Empty, Deletion is not possible" and end from the function
- Step 3: If it is Not Empty then, define a Node pointer 'temp' and set it to 'front'.
- Step 4 : Then set 'front = front → next' and delete
 'temp' (free(temp)).

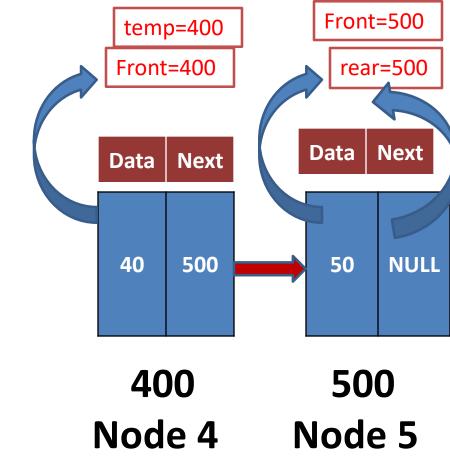






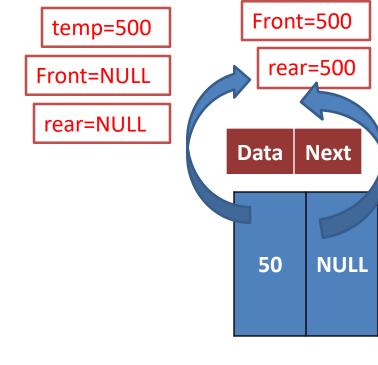






front == NULL && rear==NULL).

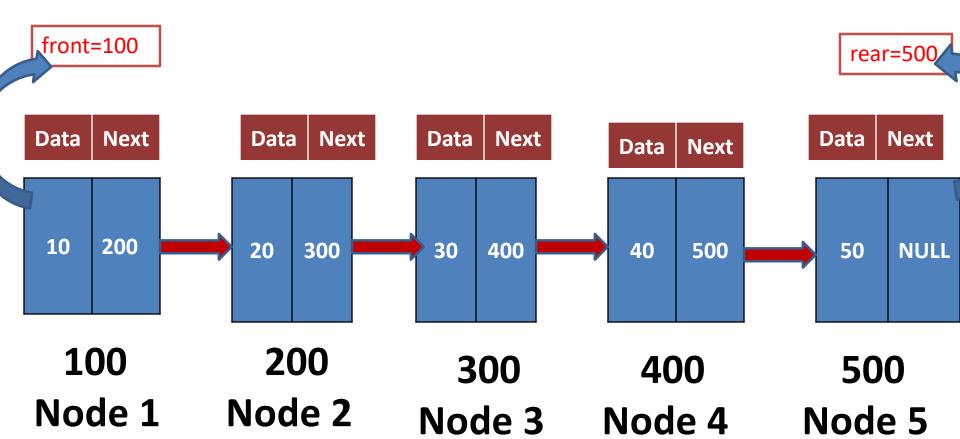
Queue is Empty, Deletion is not possible



500 Node 5

Algorithm for Display()

- display(): Display an Element from Queue
 Step 1: Check whether queue is Empty
 (front == NULL && rear==NULL).
- Step 2: If it is Empty, then display "Queue is Empty, Display is not possible" and end from the function
- **Step 3:** If it is **Not Empty** then, define a Node pointer 'temp' and initialize with **front**.
- Step 4: Display 'temp → data' and move it to the next node. Repeat the same until 'temp' reaches to 'rear' (temp → next != NULL).
- Step 5 : Finally! Display 'temp → data ---> NULL'.



Thank You