$\begin{array}{c} \textbf{Convex Optimization Overview} \\ \textit{Conner DiPaolo} \end{array}$

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$$\mathbf{x} \qquad \theta \mathbf{x} + (1 - \theta) \mathbf{y} \qquad \mathbf{y}$$

Figure 1: Convex Combination. The line segment between \mathbf{x} and \mathbf{y} above represents every possible combination $\{\theta \mathbf{x} + (1 - \theta)\mathbf{y} : 0 \le \theta \le 1\}$.

1 Introduction

Convex optimization in a large way influences the way people think about and phrase machine learning problems. Almost all problems we will see in our studies are developed or can be viewed as optimization problems. Some problems, like the Support Vector Machine you will all see in the coming weeks, are almost entirely based in the heart of convex optimization.

We don't plan on bringing you all up to speed completely on the art of Convex Optimization, but hopefully in two sections you will know enough to be able to think about problems in new, interesting ways that will aid your studies in and out of machine learning.

The only real prerequisite for this material is a strong confidence in linear algebra and familiarity with matrix calculus. Note that much of this material stems from Boyd and Vandenberghe's insanely influential textbook *Convex Optimization*. If you are interested in the topic or need a more in-depth resource, check out the book. It's free online¹.

2 Convex Sets

The first step into examining convexity is defining what a **convex set** is when given, for example, a subset of the real numbers, or the set of matrices.

Definition 2.1 (Convex Combination). In the n = 2 case, the convex combination of points \mathbf{x} and \mathbf{y} in an affine space is

$$\theta \mathbf{x} + (1 - \theta) \mathbf{y}$$

where $0 \le \theta \le 1$. Intuitively, this is the line segment between \mathbf{x} and \mathbf{y} (consider $\theta = 0$ and $\theta = 1$), as seen in Figure 1. More generally, a convex combination of points $\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_n$ in an affine space (vector spaces included) is the combination

$$\boldsymbol{\theta}_1 \mathbf{x}_1 + \boldsymbol{\theta}_2 \mathbf{x}_2 + \dots + \boldsymbol{\theta}_n \mathbf{x}_n = \boldsymbol{\theta}^{\top} \mathbf{x}$$

where $\theta_i \geq 0$ and $\mathbf{1}^{\top}\theta = 1$. That is, $\boldsymbol{\theta}$ lies on the standard probability simplex.

With this definition of a convex combination we can define a convex set:

Definition 2.2 (Convex Set). A set C is convex if, given $\mathbf{x}, \mathbf{y} \in C$, every convex combination of \mathbf{x} and \mathbf{y} is still in C. Mathematically,

$$\theta \mathbf{x} + (1 - \theta)\mathbf{v} \in C. \tag{0 < \theta < 1}$$

See Figure 2 for an illustration. Intuitively, this means that every line segment between any two points in C is contained entirely within C.

2.1 Examples of Convex and Non-Convex Sets

Example 2.1 (All of \mathbb{R}^n). \mathbb{R}^n is convex.

Proof. As a vector space, for any $\theta_1, \theta_2 \in \mathbb{R}$ and \mathbf{x} and \mathbf{y} in \mathbb{R}^n ,

$$\theta_1 \mathbf{x} + \theta_2 \mathbf{y} \in \mathbb{R}^n$$
.

Thus, restricting $\theta_2 = 1 - \theta_1$ and $0 \le \theta_1 \le 1$ does not change this fact, and any convex combination is also in \mathbb{R}^n , making the set convex by definition.

¹http://stanford.edu/~boyd/cvxbook/



Figure 2: **Set convexity.** Intuitively, a set is convex if a line drawn between any two points within the set lies completely in the set. The convex set to the left is a *convex hull* of it's vertices, meaning the set is constructed as all possible convex combinations of its vertices. These sets are subsets of \mathbb{R}^2 .

Example 2.2 (The Non-Negative Orthant \mathbb{R}^n_{\perp}). The set of all vectors

$$\{\mathbf{x}: \mathbf{x} \in \mathbb{R}^n \text{ and } \mathbf{x}_i \geq 0\}$$

is convex.

Proof. Left as an exercise to the reader.

Example 2.3 (Closed Intervals in \mathbb{R}). Let C = [a, b] be a subset of the real numbers where $a \leq b$. Then C is convex.

Proof. Suppose, without loss of generality, that $x_1 \leq x_2$ where $x_1, x_2 \in [a, b]$. Now let $0 \leq \theta \leq 1$. Then

$$\theta x_1 + (1 - \theta)x_2 < \theta x_2 + (1 - \theta)x_2 = x_2$$

because $\theta x_1 \leq \theta x_2$. Similarly,

$$\theta x_1 + (1 - \theta)x_2 \ge \theta x_1 + (1 - \theta)x_1 = x_1$$

because $(1-\theta)x_2 \geq (1-\theta)x_1$. Thus

$$a \le x_1 \le \theta x_1 + (1 - \theta) x_2 \le x_2 \le b$$
,

and hence

$$\theta x_1 + (1 - \theta)x_2 \in [a, b].$$

Therefore, by the definition of convexity, $[a, b] \subseteq \mathbb{R}$ is convex for any $a \leq b$.

Example 2.4 (The Set of All Complex Hermetian Matrices). Let $C = \{A : A \in \mathbb{C}^{n \times n} \text{ and } A^* = A\}$. C is convex.

Proof. Let $A, B \in \mathbb{C}^{n \times n}$ be Hermitian matrices and $0 \le \theta \le 1$. Then

$$(\theta A + (1 - \theta)B)^* = (\theta A)^* + [(1 - \theta)B]^*$$
(1)

$$= \theta A^* + (1 - \theta)B^* \tag{2}$$

$$= \theta A + (1 - \theta)B \qquad \text{(because } A^* = A \text{ and } B^* = B)$$

(3)

Thus every convex combination of Hermitian matrices is Hermitian, and by the definition of convexity the set is convex.

Example 2.5 (The Set of All Real Symmetric Matrices). Let $C = \{A : A \in \mathbb{R}^{n \times n} \text{ and } A^{\top} = A\}$. C is convex.

Proof. Left as an exercise to the reader.

Example 2.6 (The Set of All Linear Matrix Inequalities). Let A_i and B be symmetric $n \times n$ matrices and $\mathbf{x} \in \mathbb{R}^n$. Let $C = \{\mathbf{x} : A(\mathbf{x}) \leq B\}$ where $A(\mathbf{x}) = \mathbf{x}_1 A_1 + \cdots + \mathbf{x}_k A_k$. C is convex.

Proof. Let $0 \le \theta \le 1$ and $\mathbf{x}, \mathbf{y} \in C$. Then

$$A(\theta \mathbf{x} + (1 - \theta)\mathbf{y}) = \sum_{i} [\theta \mathbf{x}_{i} + (1 - \theta)\mathbf{y}_{i}] A_{i}$$

$$= \theta \sum_{i} \mathbf{x}_{i} A_{i} + (1 - \theta) \sum_{i} \mathbf{x}_{i} A_{i}$$

$$= \theta A(\mathbf{x}) + (1 - \theta) A(\mathbf{y})$$

$$\leq \theta B + (1 - \theta) B$$

$$= B.$$

Thus any convex combination of elements of C is contained in C and by definition C is convex.

Example 2.7 (The Space of Probability Distributions). Let \mathcal{P} be the space of continuous probability distributions over \mathbb{R}^n . That is, every element of \mathcal{P} defines a unique probability density function $\mathbb{P}(x) \geq 0$ such that

$$\int_{\mathbb{R}^n} \mathbb{P}(x) dx = 1.$$

 \mathcal{P} is convex.

Proof. Let f and h be valid probability distributions from \mathcal{P} . That is,

$$f, h \in \left\{ \mathbb{P}(x) : \int_{\mathbb{R}^n} \mathbb{P}(x) dx = 1 \text{ and } \mathbb{P}(x) \ge 0 \right\}.$$

Now let $0 < \theta < 1$. Then

$$\theta f(x) + (1 - \theta)h(x) \ge 0$$

as a positive combination of positive functions. Similarly,

$$\int_{\mathbb{R}^n} \left[\theta f(x) + (1 - \theta)h(x) \right] dx = \theta \int_{\mathbb{R}^n} f(x)dx + (1 - \theta) \int_{\mathbb{R}^n} h(x)dx$$
$$= \theta + (1 - \theta) = 1.$$

Thus every convex combination of probability distributions over \mathbb{R}^n is a valid distribution (often called a mixture), and therefore the set of all valid probability distributions over \mathbb{R}^n is convex itself.

Example 2.8 (Disjoint Intervals in \mathbb{R}). Let $N = [a, b] \cup [c, d]$ where $a \leq b < c \leq d$. N is **not** convex.

Proof. Let $b, c \in N$ be as described above. Then for $0 < \theta < 1$ (not we aren't including inequality),

$$\theta b + (1 - \theta)c \not\in N$$
.

Thus not every convex combination of elements in N is in N, and N is not convex.

Example 2.9 (The Set of All Stochastic Matrices). The set of all matrices A such that for all elements $0 \le A_{ij} \le 1$ and each row sums to 1,

$$M = \{A : A \in \mathbb{R}^{m \times n} \text{ and } 0 \le A_{ij} \le 1 \text{ and } A\mathbf{1} = \mathbf{1}\},$$

is convex.

Note that this set is the set of all matrix representations of every possible Markov Chain.

Proof. Let $A, B \in M$ and $0 \le \theta \le 1$. Then

$$c = [\theta A + (1 - \theta)B]_{ij} = \theta A_{ij} + (1 - \theta)B_{ij}$$

satisfies $0 \le c \le 1$ as $A_{ij}, B_{ij} \in [0, 1]$ and closed intervals on \mathbb{R} are convex as seen in Example 2.3.

Further, because $A\mathbf{1} = \mathbf{1}$ and $B\mathbf{1} = \mathbf{1}$,

$$[\theta A + (1 - \theta)B]\mathbf{1} = \theta A\mathbf{1} + (1 - \theta)B\mathbf{1} = \theta \mathbf{1} + (1 - \theta)\mathbf{1} = \mathbf{1},$$

as desired. Thus every convex combination of elements with M remains in M, and by definition the set M is convex.

Example 2.10 (Norm Balls). For any valid norm $||\cdot||$ and scalar $r \geq 0$, the set

$$C = \{\mathbf{x} : ||\mathbf{x}|| \le r\}$$

is convex.

Proof. Let **x** and **y** be elements from C and $0 \le \theta \le 1$. Then

$$||\theta \mathbf{x} + (1 - \theta)\mathbf{y}|| \le ||\theta \mathbf{x}|| + ||(1 - \theta)\mathbf{y}|| = \theta ||\mathbf{x}|| + (1 - \theta)||\mathbf{y}|| \le \theta r + (1 - \theta)r = r,$$

as desired, where we used the Triangle Inequality for the first step and the second used the homogeneity of norms. Thus every convex combination of elements in C is within C, and therefore C is convex.

3 Convex Functions

You likely have already seen convex functions from your time in Calculus I in high school or something similar. Nevertheless, our treatment will be much more rigorous (though not *too* rigorous) and be very applicable to our later studies.

The most important fact to know about convex functions is that they

- (a) Have a unique global minimizer \mathbf{x}^* .
- (b) Generally have efficient algorithms for finding that minimizer.

3.1 Convex and Concave Functions

Our idea of a convex set will be useful in considering convex functions.

Definition 3.1 (Convex). Given a convex set C, a function $f: C \mapsto \mathbb{R}$ is convex if, for $0 \le \theta \le 1$, given any $\mathbf{x}, \mathbf{y} \in C$,

$$f(\theta \mathbf{x} + (1 - \theta)\mathbf{y}) \le \theta f(\mathbf{x}) + (1 - \theta)f(\mathbf{y}).$$

Intuitively, this means that a convex function always lies under a line between any two points where it is evaluated. This is seen in Figure 3.

Definition 3.2 (Concave Functions). Given a convex set C, function $f: C \mapsto \mathbb{R}$ is concave if -f is convex. That is, for $0 \le \theta \le 1$, given any $\mathbf{x}, \mathbf{y} \in C$,

$$f(\theta \mathbf{x} + (1 - \theta)\mathbf{y}) \ge \theta f(\mathbf{x}) + (1 - \theta)f(\mathbf{y}).$$

Theorem 3.1 (Restriction To a Line). A function $f: \mathbb{R}^n \to \mathbb{R}$ is convex if and only if f is convex when restricted to any line that intersects its domain. Mathematically, f is convex if and only if for all $x \in C$ and any $\mathbf{v} \in \mathbb{R}^n$ and $t \in \mathbb{R}$,

$$g(t) = f(\mathbf{x} + t\mathbf{v})$$

is convex where $g: \{t: \mathbf{x} + t\mathbf{v} \in C\} \mapsto \mathbb{R}$

Proof. Omitted. See *Convex Optimization* by Boyd and Vandenberghe.

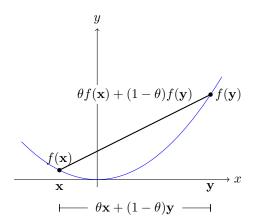


Figure 3: Convex Functions. This function $f:[a,b] \mapsto \mathbb{R}$ is convex because it's domain is convex and every line between two points on the function lies above the function.

3.2 First Order Conditions: $f(x) \ge f(y) + \nabla f(y)^{\top}(x-y)$

Generally speaking, determining convexity from the definition is hard. In this section we will develop more tractable conditions that will help us determine if an arbitrary function is convex. The following necessary and sufficient condition for convexity is called the First Order Condition because it relies only on the first derivative.

Theorem 3.2 (First Order Conditions). Let f be a function mapping from some convex set C to \mathbb{R} . Then f is convex if and only if, given \mathbf{x} and \mathbf{y} from C,

$$f(\mathbf{x}) \ge f(\mathbf{y}) + \nabla f(\mathbf{y})^{\top} (\mathbf{x} - \mathbf{y}).$$

Proof. Omitted. See *Convex Optimization* by Boyd and Vandenberghe.

3.2.1 Examples On Determining Convexity

Theorem 3.3 (Linear Functions are Both Concave and Convex). Given a function $f: C \to \mathbb{R}$ where C is a convex set and for any $a, b \in \mathbb{R}$ and $\mathbf{x}, \mathbf{y} \in \mathbb{R}$

$$f(a\mathbf{x} + b\mathbf{v}) = af(\mathbf{x}) + bf(\mathbf{v}),$$

f is both convex and concave. Note that f is the definition of a linear function.

Proof. Let \mathbf{x}, \mathbf{y} be any elements of C and $0 \le \theta \le 1$. Then by linearity

$$f(\theta \mathbf{x} + (1 - \theta)\mathbf{y}) = \theta f(\mathbf{x}) + (1 - \theta)f(\mathbf{y}).$$

Thus by definition f is both convex and concave.

Example 3.1. $f(x) = x^2$ is convex.

Proof. We will use the much more convenient second order condition for this problem later. Consider $x, y \in \mathbb{R}$ and $0 \le \theta \le 1$. Then

$$f(\theta x + (1 - \theta)y) = (\theta x + (1 - \theta)y)^{2}$$

$$= (\theta x + (1 - \theta)y)(\theta x + (1 - \theta)y)$$

$$= \theta^{2}x^{2} + 2\theta(1 - \theta)xy + (1 - \theta)^{2}y^{2}.$$

f will be convex if and only if

$$\theta f(x) + (1 - \theta)f(y) - f(\theta x + (1 - \theta)y) \ge 0$$

for all x, y. We have

$$g = \theta f(x) + (1 - \theta)f(y) - f(\theta x + (1 - \theta)y) = \theta x^{2} + (1 - \theta)y^{2} - \theta^{2}x^{2} - 2\theta(1 - \theta)xy - (1 - \theta)^{2}y^{2}$$
$$= \theta(1 - \theta)x^{2} - 2\theta(1 - \theta)xy + \theta(1 - \theta)y^{2}$$
$$= \theta(1 - \theta)(x - y)^{2}$$

We know $(x-y)^2 \ge 0$ for all x and y. Similarly, when $\theta \in [0,1]$, both θ and $1-\theta$ are positive, so this expression is positive. Thus f must be convex, as desired.

3.3 Second Order Conditions: $\nabla^2 f \succ 0$

Here we discuss the most useful condition for determining convexity.

Theorem 3.4 (Second Order Conditions). A twice-differentiable function $f: C \mapsto \mathbb{R}$ is convex if and only if and only if the Hessian

$$\nabla^2 f \succeq 0$$

Proof. Omitted.

3.3.1 More Examples On Determining Convexity

Example 3.2 (Quadratic Functions). The function $f: \mathbb{R}^n \to \mathbb{R}$ defined by

$$f(\mathbf{x}) = \mathbf{x}^{\top} A \mathbf{x} + \mathbf{b}^{\top} \mathbf{x} + c$$

is convex if $A \succeq 0$.

Proof. We know

$$\nabla^2 f = 2A.$$

Similarly we know that $2A \succeq 0$ if and only if $A \succeq 0$. Thus by the second order conditions f is convex if and only if $A \succeq 0$.

We can also show that f is concave if $A \leq 0$.

Example 3.3. $f(x) = a e^x$ is convex if $a \ge 0$.

Proof. The Hessian

$$\nabla^2 f = \frac{d^2 f}{dx^2} = a \ e^x \ge 0$$

as long as $a \ge 0$. Thus by the second order conditions f is convex if $a \ge 0$.

Example 3.4. The function $f: \mathbb{R}^n \to \mathbb{R}$ defined by

$$f(\mathbf{x}) = e^{\mathbf{x}^{\top}\mathbf{x}} = e^{\mathbf{x}_1^2 + \dots + \mathbf{x}_n^2}$$

is convex.

Proof. Each element of the Hessian

$$\nabla^2 f_{ij} = \frac{\partial^2 f}{\partial x_i \partial x_j} = 4 \mathbf{x}_i \mathbf{x}_j e^{\mathbf{x}^\top \mathbf{x}}.$$

Therefore

$$\nabla^2 f = 4\mathbf{x}\mathbf{x}^{\top} e^{\mathbf{x}^{\top} \mathbf{x}}.$$

Is this positive semidefinite? Consider $\mathbf{z} \in \mathbb{R}^n$. Then

$$\mathbf{z}^{\top} \nabla^2 f(x) \mathbf{z} = \mathbf{z}^{\top} 4 \mathbf{x} \mathbf{x}^{\top} e^{\mathbf{x}^{\top} \mathbf{x}} \mathbf{z} = 4 (\mathbf{z}^{\top} \mathbf{x})^2 e^{\mathbf{x}^{\top} \mathbf{x}} \ge 0$$

because $\exp(\cdot)$ and $(\mathbf{z}^{\top}\mathbf{x})^2$ are both non-negative. Thus $\nabla^2 f \succeq 0$ and by the second order conditions f is convex.

3.4 Operations Preserving Convexity

In this section we examine some handy tools for constructing convex functions from other convex functions. There are *many* more properties than those shown here. See *Convex Optimization* by Boyd and Vandenberghe for many others.

Theorem 3.5 (Non-Negative Weighted Sum). The function $F: C \to \mathbb{R}$ where C is a convex set, $f_i: C \to \mathbb{R}$ is a convex function, $w_i \geq 0$ and

$$F(\mathbf{x}) = \sum_{i} w_i f_i(\mathbf{x})$$

is convex.

Proof. Let $\mathbf{x}, \mathbf{y} \in C$ and $0 \le \theta \le 1$. Then

$$F(\theta \mathbf{x} + (1 - \theta)\mathbf{y}) = \sum_{i} w_{i} f_{i}(\theta \mathbf{x} + (1 - \theta)\mathbf{y})$$

$$\leq \sum_{i} w_{i} \left[\theta f_{i}(\mathbf{x}) + (1 - \theta)f_{i}(\mathbf{y})\right]$$

$$= \theta F(\mathbf{x}) + (1 - \theta)F(\mathbf{y}),$$

as desired.

Theorem 3.6 (Pointwise Maximum). Given convex functions $f_i: C \to \mathbb{R}$, and $F: C \to \mathbb{R}$ defined as

$$F(\mathbf{x}) = \max_{i} f_i(\mathbf{x}),$$

F is convex.

Proof. By the definition of convexity, given $\mathbf{x}, \mathbf{y} \in C$ and $0 \le \theta \le 1$ we have

$$F(\theta \mathbf{x} + (1 - \theta)\mathbf{y}) = \max_{i} f_{i}(\theta \mathbf{x} + (1 - \theta)\mathbf{y})$$

$$\leq \max_{i} \{\theta f_{i}(\mathbf{x}) + (1 - \theta)f_{i}(\mathbf{y})\} \qquad \text{(because } f_{i} \text{ is convex)}$$

$$\leq \theta \max_{i} f_{i}(\mathbf{x}) + (1 - \theta) \max_{i} f_{i}(\mathbf{y})$$

$$= \theta F(\mathbf{x}) + (1 - \theta)F(\mathbf{y}),$$

as desired.

4 Optimization Problems

An optimization problem is a problem of the form

minimize:
$$f_0(\mathbf{x})$$

subj. to: $f_i(\mathbf{x}) \leq 0$, $i = 1, ..., m$
 $h_i(\mathbf{x}) = 0$, $i = 1, ..., p$

where $f_0: \mathbb{R}^n \to \mathbb{R}$ is called the **objective function**, $f_i: \mathbb{R}^n \to \mathbb{R}$ is an **inequality constraint**, and $h_i: \mathbb{R}^n \to \mathbb{R}$ is an **equality constraint**. The goal of an optimization problem, as you might be able to guess, is to find \mathbf{x} in the problem domain domain

$$\mathcal{D} = \bigcap_{i=0}^{m} \mathbf{dom} \ f_i \ \cap \ \bigcap_{i=1}^{p} \mathbf{dom} \ h_i$$

that minimizes $f_0(\mathbf{x})$ subject to the given conditions. The problem is said to be *feasible* if such an \mathbf{x} exists, and *infeasible* otherwise.

The optimal value p^* of the problem above is defined to be

$$p^* = \inf\{f_0(\mathbf{x}) : f_i(\mathbf{x}) \le 0, \ i = 1, \dots, m, \ h_i(\mathbf{x}) = 0, \ i = 1, \dots, p\}.$$

If there are feasible points \mathbf{x}_k with $f_0(\mathbf{x}_k) \to -\infty$ as $k \to \infty$, then we say the problem is unbounded below.

Note that we could have a similarly defined *maximization* problem. Everything is the same except for the definition of *unbounded below*, obviously.

4.1 Optimal Points

We say \mathbf{x}^* is an *optimal point* (or that it solves the above problem) if \mathbf{x}^* is feasible and $f_0(\mathbf{x}^*) = p^*$. If an optimal point exists, we say the problem is *solvable*.

We say a feasible point **x** is *locally optimal* if there exists some R > 0 such that **x** solves

minimize:
$$f_0(\mathbf{z})$$

subj. to: $f_i(\mathbf{z}) \leq 0$, $i = 1, ..., m$
 $h_i(\mathbf{z}) = 0$, $i = 1, ..., p$
 $||\mathbf{z} - \mathbf{x}||_2 \leq R$

for variable \mathbf{z} . This intuitively means that \mathbf{x} minimizes f_0 over nearby points in the feasible set.

4.2 Equivalent Problems

Given an optimization problem of the form

minimize:
$$f_0(\mathbf{x})$$

subj. to: $f_i(\mathbf{x}) \leq 0$, $i = 1, ..., m$
 $h_i(\mathbf{x}) = 0$, $i = 1, ..., p$

we may want to express our problem in a cleaner or easier to solve form. Under certain conditions changing our problem's form will not actually change the solution. Note that, as usual, there are a few more possible transformations that are kosher. Check out Boyd and Vandenberghe for these, but the ones shown here are likely the only ones you will need.

4.2.1 Scaling

The problem

minimize:
$$\alpha_0 f_0(\mathbf{x})$$

subj. to: $\alpha_i f_i(\mathbf{x}) \leq 0$, $i = 1, ..., m$
 $\beta_i h_i(\mathbf{x}) = 0$, $i = 1, ..., p$

for $\alpha > 0$ and $\beta \neq 0$ is equivalent to the original problem. This should be intuitive since, for example x^2 is minimized at 0, changing the scale of your axes maintains that minimum. Similarly, if equality holds (ie. 4x + 2 = 0) multiplying by any non-zero number on both sides maintains that equality.

4.2.2 Change of Variables

Given an one-to-one mapping $\phi : \mathbb{R}^n \to \mathbb{R}^n$, where the image (range) of ϕ covers the domain of your problem \mathcal{D} , the problem

minimize:
$$f_0(\phi(\mathbf{x}))$$

subj. to: $f_i(\phi(\mathbf{x})) \leq 0$, $i = 1, ..., m$
 $h_i(\phi(\mathbf{x})) = 0$, $i = 1, ..., p$

is equivalent to the original. This should be clear. If x solves the original problem, then $z = \phi^{-1}(x)$ solves the transformed problem. The converse also holds.

4.2.3 Transformation of Objective and Constrain Functions

Suppose $\psi_0 : \mathbb{R} \to \mathbb{R}$ is monotonically increasing, $\phi_1, \dots, \psi_m : \mathbb{R} \to \mathbb{R}$ satisfy $\phi_i(u) \leq 0$ if and only if $u \leq 0$, and $\psi_{m+1}, \dots, \phi_{m+p} : \mathbb{R} \to \mathbb{R}$ satisfy $\phi_i(u) = 0$ if and only if u = 0. The problem

minimize:
$$\psi_0(f_0(\mathbf{x}))$$

subj. to: $\psi_i(f_i(\mathbf{x})) \leq 0$, $i = 1, ..., m$
 $\psi_{m+i}(h_i(\mathbf{x})) = 0$, $i = 1, ..., p$

is equivalent to the original. This should be evident from the conditions we placed on each ψ_i .

4.2.4 Slack Variables

This will come in very handy. A prudent observation is that $f_i(x) \leq 0$ if and only if for some $s_i \geq 0$, $f_i(x) + s_i = 0$. Thus the problem

minimize:
$$f_0(\mathbf{x})$$

subj. to: $f_i(\mathbf{x}) + s_i = 0, \quad i = 1, \dots, m$
 $h_i(\mathbf{x}) = 0, \quad i = 1, \dots, p$
 $s_i \ge 0, \quad i = 1, \dots, m$

is equivalent to the original.

4.2.5 Epigraph Form

The *epigraph form* of the original problem is

minimize:
$$t$$

subj. to: $f_0(\mathbf{x}) - t \le 0$,
 $f_i(\mathbf{x}) \le 0$, $i = 1, \dots, m$
 $h_i(\mathbf{x}) = 0$, $i = 1, \dots, p$

where $t \in \mathbb{R}$. It should be clear that this is the same as the original problem as the first constraint can be viewed as $f_0(\mathbf{x}) \leq t$. Thus the objective must always lie under t and minimizing t will minimize the highest possible value of f_0 .

4.3 Convex Optimization Problems

We will now study the branch of optimization problems we will see *extremely* often in our studies. We say a problem of the form

minimize:
$$f_0(\mathbf{x})$$

subj. to: $f_i(\mathbf{x}) \leq 0$, $i = 1, ..., m$
 $h_i(\mathbf{x}) = 0$, $i = 1, ..., p$

is a **convex optimization problem** if the objective function and inequalities f_i are convex, and the equality constraints h_i are linear. Convex problems are, as a whole, extremely nice in the sense that we have efficient (read 'polynomial time') algorithms for finding global optima.

- 4.3.1 Linear Programs
- 4.3.2 Quadratic Programs
- 4.3.3 Global Optimality of a Convex Optimum
- 4.4 Examples of Convex Optimization Problems
- 4.4.1 Optimal Flow As a Convex Optimization Problem

5 Algorithms for Solving Convex Problems

- 5.1 Gradient Descent
- 5.2 Newton's Method
- 6 Duality
- 6.1 Duality Gap
- 6.2 The Lagrange Dual Function
- 6.3 Finding the Dual
- 6.3.1 The Dual of The Linear Program
- 6.4 Solving Problems Using Duality
- 6.4.1 Minimize a Quadratic Under Quadratic Constraints
- 6.5 Complementary Slackness of Duality
- 6.5.1 Solving The Dual Linear Program via a Primal Solution
- 6.6 The KKT (Karush-Kuhn-Tucker) Conditions