

Lecture 18

Week 4

Adversarial Bandits for Exp 3, Exp 3P, Exp 3I

$$\bar{R}(n, \pi) \leq O(\sqrt{nk \log k}) \quad \leftarrow \text{Covered in Week}$$

$$\bar{R}(n, \pi) \geq \Omega(\sqrt{nk})$$

Online Convex Optimization - Assume that:

Input: A convex set $S \subset \mathbb{R}^d$

For $t = 1, 2, \dots$

predict vector $w_t \in S$

receive convex loss function $f_t: S \rightarrow \mathbb{R}$

Suffer loss $f_t(w_t)$

Adversary is choosing convex function in each round

$$R(n, \pi) = \sum_{t=1}^n f_t(w_t) - \min_{u \in S} \sum_{t=1}^n f_t(u) \quad \text{--- (1)}$$

For the full information setting, with weighted majority algorithm, S can be considered as

\dots These are the

$\mathcal{S} = \{ w \mid w_t \geq 0 \forall t \text{ \& } \sum w_i = 1 \}$ [same weights that are assigned to the experts]

In the weighted majority,

after choosing w_t , one would incur V_t from the environment (adversary). Hence loss was:

$$\text{loss} = \langle w_t, v_t \rangle = f_t(w_t) \leftarrow \text{can be written like this for convex setting, it's a linear function}$$

Now let the regret be:

$$R(n, \pi) = \sum_{t=1}^n f_t(w_t) - \min_{u \in U} \sum_{t=1}^n f_t(u) \quad \text{--- (1a)}$$

s.t. $U \subset S$ containing all unit vectors

$$\text{thus } U = \{ (1, 0, \dots, 0), (0, 1, 0, \dots), \dots, (0, \dots, 1, \dots, 0), \dots, (0, \dots, 0, 1) \}$$

$$\text{then } f_t(u) = f_t(e_i) = \sum_{t=1}^n v_{ti}$$

Thus eq. 1a becomes

$$R(n, \pi) = \sum_{t=1}^n \langle v_t, w_t \rangle - \min_{t=1} \sum_{t=1}^n v_{ti} \quad \text{--- (1b)}$$

The aforementioned case was for a linear function, but convex optimization can be applied for any function, as long as it's convex.

Hence reconsider eq ①

$$R(n, \pi) = \sum_{t=1}^n f_t(w_t) - \min_{u \in S} \sum_{t=1}^n f_t(u) \quad \text{--- ①}$$

For any convex function, at round t ,

once w_t is chosen, f_t will be revealed by the em

The goal now is to find out an algorithm to choose w

Follow the Leader (FTL)

Till round t , we know

$$\{J_1, J_2, J_3, \dots, J_{t-1}\}$$

Hence we can choose w_t as :

$$w_t = \underset{w \in S}{\operatorname{argmin}} \sum_{i=1}^{t-1} J_i(w)$$