Week 7 - Lecture 33

We lenau:

$$\overline{\mathbb{R}}_{n} = \underbrace{\mathbb{E}}_{i=1}^{K} \mathbb{E}\left[T_{i}(n)\right] \triangle_{i} = \underbrace{\mathbb{E}}_{(i:\Delta_{i}>0)} \mathbb{E}\left[T_{i}(n)\right] \triangle_{i}$$

Theorem: UCB pull each sub-optimal arm K
in expectation at most $\Delta_i = \max_j M_i - M_i$

$$\mathbb{E}\left[T_{K}(n)\right] \leq \frac{6\log n}{\Delta_{K}^{2}} + \pi^{2}/_{3} + 1 \quad \text{times}.$$

And Psuedo-regret is bounded as:

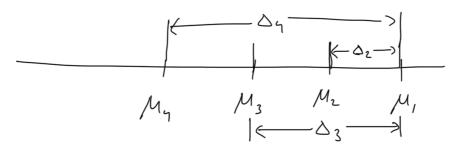
$$\overline{R}_{n} \leq 6 \leq 6 \log K + \leq (\pi^{2}/3 + 1) \cdot \Delta_{i}$$

$$= \frac{1}{1 \cdot \Delta_{i} > 0} \Delta_{K} + \frac{1}{1 \cdot \Delta_{i} > 0}$$

$$= \frac{1}{1}$$

Assuming arm 1 as optimal then $\triangle_1 = 0$, $\triangle_2 > 0$, $\triangle_3 > 0 - - ...$

$$\frac{\text{Defn}:}{\sum_{i>1}} = \min_{i>1} \Delta_i$$



 \triangle is smallest between \triangle_2 , \triangle_3 $A \triangle_4$: the gap between highest mean and next highest mean

 $\therefore \triangle \angle \triangle_i \forall i > 1$

Plugging in defn of a in ea eq 1

$$\overline{R}_{n} \leq 6\left(K-1\right) \frac{\log n}{\Delta} + \frac{1}{2} \left(\frac{\pi^{2}}{3} + 1\right) \Delta_{i}$$

$$= O\left(\frac{\log n}{\Delta}\right) \longrightarrow Sub-linear in n$$

- -> If \triangle is smaller then identifying the best arm is more challenging than if \triangle is larger.

 If the sub-oftimal Ms are closer to the optimal Men figuring and the one optimal arm becomes harder.
- —) The \bar{R}_n bound given in the theorem captures this as $\bar{R}_n = O\left(K \frac{\log n}{\Delta}\right)$: as Δ becomes

$$\widehat{\mathcal{M}}_{i}(t) = \frac{1}{T_{i}(t-1)} \underbrace{\leq}_{S=1}^{T_{i}(t-1)} X_{si}$$

$$\underbrace{\int_{i}^{T_{i}(t-1)} X_{si}}_{S=1} X_{si}$$
dirst $T_{i}(t-1)$
observed namples

-> S does not run like t, it is to be taken forly for the rounds where i was played

$$UCB_{i}(t-1) = \hat{\mu}_{i}(t-1) + \sqrt{2 \frac{\log t}{T_{i}(t-1)}}$$
Index of arm i

-> Suppose in round t sub-oftimal arm i is played, this could have happened because for that round t, index of arm i is greater than the index of all other arms and in particular is greater than the index of arm 1 (assumed to be optimal)

At round t,
$$I_t = i$$
 (i is sub-optimal)

i]: $\hat{\mathcal{M}}_i(t-1) + \sqrt{\frac{2 \log t}{T_i(t-1)}} \geq \hat{\mathcal{M}}_j(t-1) + \sqrt{\frac{2 \log t}{T_j(t-1)}}$
 $\forall j$

$$\hat{\mathcal{U}}_{i}(t-1) + \sqrt{\frac{2 \log t}{T_{i}(t-1)}} > \hat{\mathcal{U}}_{i}(t-1) + \sqrt{\frac{2 \log t}{T_{i}(t-1)}}$$

-) Thus the index of i-th arm was overestimated and index of optimal arm was underestimated, because i-th arm was not played enough no. of times

That the estimates were no happened that the estimates were not estimated properly and
$$\mu$$
, lies above and $\hat{\mu}$, lies below the respective confidence bounds.