Week 3 - Lecture 17

In each round learner pichs Pe Adversary pichs le

Types of Adversary:

- Oblivious: Before game begins Adversary chooses
 the losses for all rounds
 { l,, l2, --- ln}
- Non-Oblivious: losses can be assigned to each action based on the past observations

Exp3 holds for both kinds of adversary.
Two kinds of E were taken, one for learner 4 one
for adversary, holds for non-oblivious case, which
is the worse of the two.

$$R(n,\pi) = \sum_{t=1}^{n} l_{t} l_{t} - \min_{i} \sum_{t=1}^{n} l_{t} i \xrightarrow{i} \lim_{t \to 1} l_{t} i$$

$$E[R(n,\pi)] = E[\sum_{t=1}^{n} l_{t} l_{t}] - \min_{i} E[\sum_{t=1}^{n} l_{t} i] \xrightarrow{i} (2)$$

The bound on (2) (already derived) will be very

different, may be higher, than That for U

We want to bound $R(n,\pi)$. For a good algorithm we would want to know

$$P_r\left\{R(n,\pi)\leq \mathbb{E}\left[R(n,\pi)\right]\right\} \leftarrow Should be high$$

While finding the bound for (1) we will only consider the Oblivious adversary.

For the bound on a quantity, single value, to be low we would want the estimates to have low variance.

Estimates of Exp3

$$Var(\hat{l}_{ti}) = \frac{1}{P_{ti}}$$

If Pti is small then variance will be very high.

Exp3. P Algorithm

Input:
$$\{n_{\epsilon}\}, K, T \left[Y \in (0, 1)\right], B, \delta$$

for
$$t = 1, 2, ...$$

Hi $P_{ti} = (1-r)e \times p \left(-h_t \hat{L}_{ti}\right)$
 $E \times p \left(-h_t \hat{L}_{tj}\right)$
 $E \times p \left(-h_t \hat{L}_{tj}\right)$

Draw I, ~ P, Yi Îxi = 1xi 1 { Ix= i} + B and Li= L(1-1); + lti

High probability bound for Exp3.P

Also ensures that Pti is not arbitariky close to zero, it is bound by 8/K, thus ensuring low variance

For any SE(0,1) if Exp3P is run with

$$\beta = \sqrt{\frac{\log K/8}{nK}}$$
, $\eta = 0.95 \sqrt{\log K}$, $\gamma = \sqrt{\frac{K \log K}{n}}$

Then,

$$R(n, E \times p3.P) \leq 5.15 \sqrt{n K \log(K/S)}$$

Further if $\beta = \sqrt{\frac{\log K}{n r}}$

 $R(n,E\times \beta 3.P) \leq \sqrt{\frac{n \, K}{\log K}} \log(\frac{1}{8})$ (stup after +5.15 VnKlogK n rounds)

with probability $(1-\delta)$

because we ar considering individual regre

with probability (1-8.

* This situation is when we don't specify & a-priori High Probability bounds are more desired that expected bounds, as it bounds the actual regret we face at every round. But these bounds are difficult to maintain become they are guarenteed by Pr(1-8)

High Probability bounds can be converted to expected bound by considering:

$$\mathbb{E}[X] \leq \int_{\delta}^{1} \frac{1}{s} P_{s} \left\{ x \geq \ln \frac{1}{s} \right\} ds$$

$$\left[R\left(n, E \times \beta 3P\right) - 5.15 \sqrt{n K \log K}\right] \sqrt{\frac{\log K}{n K}} > \log(\frac{1}{N})$$

holds de w.p. 8

Let
$$\left[R(n, E \times p3P) - S.15 \sqrt{n K \log K}\right] \sqrt{\frac{\log K}{n K}} = X$$

Then

$$\mathbb{E}\left[R(n, Exp3P) - 5.15\sqrt{nK\log K}\right]\sqrt{\frac{\log K}{nK}} \leq 1$$

$$\mathbb{E}\left[R(n, E \times b^3 P)\right] \leq \sqrt{\frac{nK}{16g \, \text{K}}} + 5.15 \, \sqrt{nK \log K}$$

Exp3 IX Algorithm

Forces implicit exploration

$$\forall i$$
 $P_{ti} = \frac{e \times p(-\eta_t \hat{L}_{ti})}{\sum_{j} e \times p(-\eta_t \hat{L}_{tj})}$

$$\forall i \quad \hat{l}_{ti} = \underbrace{l_{ti} 1 \{ I_{t} = i \}}_{P_{ti} + Y}$$