

Type Checking (Part-2)

Course : 2CS701/IT794 – Compiler
Construction

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Ref : Ch.6 Compilers Principles, Techniques, and Tools by Alfred Aho, Ravi Sethi, and Jeffrey Ullman

Glimpse

- Quick Recap of Type Checking Part-1
- Type System (contd....)
 - Checking of Identifier Re-declaration and Un-declaration
 - Checking of Function Declaration and Invocation
 - Type Conversion and Coercion
- Implementation in YACC
 - Constructing Type Graph in Yacc
 - Type Checking in Yacc
 - Type Coercion in Yacc
 - Checking L-Values and R-Values in Yacc

Checking of Identifier Re-declaration and Un-declaration

$E \rightarrow \mathbf{true} \mid \mathbf{false}$	{ $E.type = \mathit{boolean}$ }
$E \rightarrow \mathbf{num}$	{ $E.type = \mathit{integer}$ }
...	
$E \rightarrow \mathbf{id}$	{ $lookup(\mathbf{id}.entry)$ if ($\mathbf{id}.type = \mathbf{NULL}$) then print ‘ Identifier un-declared ’ $E.type = \mathit{type_error}$ else $E.type = lookup(\mathbf{id}.entry)$ }
$D \rightarrow \mathbf{id} : T$	{ $lookup(\mathbf{id}.entry)$ if ($\mathbf{id}.type \neq \mathbf{NULL}$) then print ‘ <i>Identifier re-declaration</i> ’ $D.type = \mathit{type_error}$ else $addtype(\mathbf{id}.entry, T.type)$ }
$T \rightarrow \mathbf{boolean}$	{ $T.type := \mathit{boolean}$ }
$T \rightarrow \mathbf{char}$	{ $T.type := \mathit{char}$ }
...	

A Simple Language Example: Functions

$$T \rightarrow T \rightarrow T$$

Function type declaration

$$E \rightarrow E (E)$$

Function call

Example:

```
v : integer;  
odd : integer -> boolean;  
if odd(3) then  
    v := 1;
```


Simple Language Example: Function Declarations

$$T \rightarrow T_1 \rightarrow T_2 \{ T.\text{type} := \text{function}(T_1.\text{type}, T_2.\text{type}) \}$$


Parametric type:
type constructor

Simple Language Example: Checking Function Invocations

$$E \rightarrow E_1 (E_2) \quad \{ E.type := \textbf{if } E_1.type = \textit{function}(s, t) \textbf{ and } E_2.type = s \\ \textbf{then } t \textbf{ else } type_error \}$$

Type Conversion and Coercion

- *Type conversion* is explicit, for example using type casts
- *Type coercion* is implicitly performed by the compiler to generate code that converts types of values at runtime
- Both require a *type system* to check and infer types from (sub)expressions

Example of Type Coercion



1) $E \rightarrow E_1 \text{ op } E_2 \quad \{ E.\text{type} :=$

if $((E_1.\text{type} = \text{integer} \text{ and } E_2.\text{type} = \text{float}) \text{ or } (E_2.\text{type} = \text{integer} \text{ and } E_1.\text{type} = \text{float}))$ **then** float

else if $(E_2.\text{type} = \text{float} \text{ and } E_1.\text{type} = \text{float})$ **then** float

else if $(E_2.\text{type} = \text{integer} \text{ and } E_1.\text{type} = \text{integer})$ **then** integer

else if ... else type_error }

2) $E \rightarrow E_1 + E_2 \quad \{ E.\text{type} :=$ **if** $((E_1.\text{type} = \text{type_error} \text{ or } E_2.\text{type} = \text{type_error})$ **then** type_error

else $\max(E_2.\text{type}, E_1.\text{type})$ }

Example of Type Conversion

$$E \rightarrow (T) E_1 \quad \{ E.type := \text{if } ((E_1.type = type_error) \\ \text{then } type_error \\ \text{else } T.type) \}$$

Constructing Type Graphs in Yacc

Type *mkint() construct int node if not already
 constructed

Type *mkarr(Type*, int) construct array-of-type node
 if not already constructed

Type *mkptr(Type*) construct pointer-of-type node
 if not already constructed

Constructing Type Graphs in Yacc

```
%union
{ Symbol *sym;
  int num;
  Type *typ;
}
%token INT
%token <sym> ID
%token <int> NUM
%type <typ> type
%%
decl : type ID          { addtype($2, $1); }
     | type ID '[' NUM ']' { addtype($2, mkarr($1, $4)); }
     ;
type : INT              { $$ = mkint(); }
     | type '*'          { $$ = mkptr($1); }
     | /* empty */       { $$ = mkint(); }
     ;
```


Type Checking in Yacc

```
%{  
enum Types {Tint, Tfloat, Tpointer, Tarray, ... };  
typedef struct Type  
{ enum Types type;  
  struct Type *child; // at most one type parameter  
} Type;  
%}
```

```
%union  
{ Type *typ;  
}
```

```
%type <typ> expr
```

```
%%
```

```
...
```


Type Checking in Yacc (contd...)

...

%%

```
expr : expr '+' expr { if ($1->type != Tint
                        || $3->type != Tint)
                        semerror("non-int operands in +");
                        $$ = mkint();
                        emit(iadd);
                        }
```


Type Coercion in Yacc

...

%%

```
expr : expr '+' expr
    { if ($1->type == Tint && $3->type == Tint)
      { $$ = mkint(); emit(iadd);
      }
      else if ($1->type == Tfloat && $3->type == Tfloat)
      { $$ = mkfloat(); emit(fadd);
      }
      else if ($1->type == Tfloat && $3->type == Tint)
      { $$ = mkfloat(); emit(i2f); emit(fadd);
      }
      else if ($1->type == Tint && $3->type == Tfloat)
      { $$ = mkfloat(); emit(swap); emit(i2f); emit(fadd);
      }
      else semerror("type error in +");
        $$ = mkint();
    }
}
```


Checking L-Values and R-Values in Yacc

```
%{  
typedef struct Node  
{ Type *typ; // type structure  
  int islval; // 1 if L-value  
} Node;  
%}  
  
%union  
{ Node *rec;  
}  
  
%type <rec> expr  
  
%%  
...
```


Checking L-Values and R-Values in Yacc (contd...)

```
expr : expr '+' expr
    { if ($1->typ->type != Tint || $3->typ->type != Tint)
      semerror("non-int operands in +");
      $$->typ = mkint();
      $$->islval = FALSE;
      emit(...);
    }
| expr '=' expr
    { if (!$1->islval || $1->typ != $3->typ)
      semerror("invalid assignment");
      $$->typ = $1->typ;
      $$->islval = FALSE;
      emit(...);
    }
| ID
    { $$->typ = lookup($1);
      $$->islval = TRUE;
      emit(...);
    }
}
```


Q and A
