

Starts @ 10:15 am.

ACTIVITY:-

Guess whose logo is it

✓



1

Starbucks



Malibu



Atari



Google



5



CBS



6

Unilever



1

Kodak



BiC



beats



10

You Tube

Constant Operations — includes

- print statements
- if-else condition
- break
- increment/decrement statement
- continue
- return

Their time-complexity is always $O(1)$

for (int i = 1; i <= N; i++)
{
 if (i == 2)
 continue;
}

$$[1, N] = N$$

$$O(N)$$

- Time complexities can be represented by any of the mathematical functions we studied in the last class.

| Sno | Function Name | Function Expression |
|-----|---------------|---------------------|
| 1 | Constant | 1 |
| 2 | Logarithmic | $\log(n)$ |
| 3 | Square root | \sqrt{n} |
| 4 | Linear | n |
| 5 | Linearithmic | $n \cdot \log(n)$ |
| 6 | Quadratic | n^2 |
| 7 | Cubic | n^3 |
| 8 | Exponential | x^n |
| 9 | Factorial | $n!$ |

Linearithmic
↓
 $n \log n$

10 mins break

for i in range(n): [0, n-1]
 for j in range(n/4)
 for k in range(n):
 print("x")

T.C $\Rightarrow ?$ $O(n^3)$ ~~///~~
 $O(n^3)$

| i | j | k |
|---|-----|---|
| 1 | 1 | n |
| | 2 | n |
| | 3 | n |
| | ⋮ | ⋮ |
| | n/4 | n |
| 2 | 1 | n |
| | 2 | n |
| | ⋮ | ⋮ |
| | n/4 | n |
| ⋮ | ⋮ | ⋮ |
| n | 1 | n |
| | 2 | n |
| | ⋮ | ⋮ |
| | n/4 | n |

$$\begin{array}{ccc}
 i & j & k \\
 n & \times & n/4 & \times & n
 \end{array}$$

$$\Rightarrow \frac{n^3}{4}$$

$$(n+1) (\log_2 n)$$

$$= n \log_2 n + \frac{\log_2 n}{2}$$

$$\Rightarrow O(n \log_2 n)$$

```

n=5
c=0
for i in range(1,n+1):
    j=1
    while(j<n):
        c+=1
        print(c)
        j*=2
    
```

```

j=1
while(j<n){
    c+=1
    print(c)
    j*=2
}
    
```

After k iterations, stop.

| iteration | j value after every iteration |
|-----------|---|
| ① | $j=2 \Rightarrow 2 \rightarrow 2^1$ |
| 2 | $j=2 \times 2 \Rightarrow 4 \rightarrow 2^2$ |
| 3 | $j=2 \times 4 \Rightarrow 8 \rightarrow 2^3$ |
| 4 | $j=8 \times 2 \Rightarrow 16 \rightarrow 2^4$ |
| ⋮ | ⋮ |
| k | $\rightarrow 2^k$ |

$$2^k = n \Rightarrow \log_2 2^k = \log_2 n$$

$$k = \log_2 n$$

```

for (int i = 0; i < n; i++) {
    int j = 1
    while (j <= n) {
        j *= 2;
    }
    for (int k = 0; k < n; k++)
        print(=)
}

```

```

n=2
for i in range(n):
    j=1
    while(j<=n):
        j*=2
        for k in range(n):
            print("**")

```

T.C $\rightarrow O(n^2)$

Obs:-

When 2 loops are independent of each other, then take the maximum one.

| i | j | k | max of j, k |
|---|--------------|---|-------------|
| 1 | (\log_2^n) | n | n |
| 2 | (\log_2^n) | n | n |
| ⋮ | | | |
| n | (\log_2^n) | n | n |

$n + n + n \dots n$ times
 $\Rightarrow n + n \Rightarrow n^2$

```
c=0
for i in range(n): →
    for j in range(n):
        for k in range((n/2), n+1, 2):
            c+=1
```

↑ ↑ ↑

for (int k = $n/2$; k < n+1 ; k += 2)

```
for (int i=0; i<n; i++){
```

```
    for (int j=0; j<n; j++){
```

```
        for (int k=0; k<n; k++){
```

```
            print(==)
```

2 times

1

$n \times 2 = 2n$

n

i

j

*

k

n

*

n

*

2

⇒

n^2

⇒

$O(n^2)$

i

0

j

0

1

2

...

n

0

1

...

n

0

1

...

n

k

2

2

2

...

2

2

2

...

2

2

2

...

2

⇒

$T.C \Rightarrow O(n^2)$

```
c=0
for i in range(n):
    for j in range(n):
        for k in range((n/2),n+1,(n/2)):
            c+=1
        print(c)
```

sum n

$$2^k = n$$

$$\log_2 2^k = \log_2 n$$

$$\Rightarrow k = \log_2 n$$

```
n=6  
i=1  
while(i<n):  
    print("**")  
    i=i*2
```

for i in range(1, n):

i=1
while(i<n):
 print("**")
 i=i*2

$$T.C \Rightarrow O(\log_2 n)$$

i

1

2

3

4

k

After every iteration:

$$i \times 2 \Rightarrow 2 \rightarrow 2^1$$

$$i = 4 \rightarrow 2^2$$

$$i = 8 \rightarrow 2^3$$

$$i = 16 \rightarrow 2^4$$

$$i \rightarrow 2^k$$

n=2

```
for i in range(n):  
    for j in range(n**2):  
        print("*")  
    for k in range(n**3):  
        print("*")
```

for(int i=0; i<n; i++){

for(int j=0; j<n*n; j++){

for(int k=0; k<n*n*n; k++){

$\rightarrow O(n^4)$

| i | j | k | max |
|---|-------|-------|-------|
| 1 | n^2 | n^3 | n^3 |
| 2 | n^2 | n^3 | n^3 |
| ⋮ | | | |
| n | n^2 | n^3 | n^3 |

$n^3 + n^3 + n^3 \dots$
 $+ n^3$

$n * (n^3)$

$\rightarrow n^4$

```
n=2
for i in range(n): → n
    j=1
    while(j<=n): → log2 n
        j*=2
    for k in range(n): → n
        print("*")
```

n^2

```
def fun(n):  
    count=0  
    for i in range(int(n/2),n+1):  
        for j in range(1,n+1,2):  
            count+=1  
    print(count)  
  
fun(5)
```

```
n=5
i=n
x=0
while(i>=1):
    for j in range(1,n+1):
        x+=1
    i=int(i/2)
print(x)
```

Time complexity

space complexity

Problem: Find TC and SC

```
def fun(a,b,c):  
    p=a  
    q=b  
    r=c  
    print(p+q+r)  
    return a+b+c  
x=fun(3,4,5)  
print(x)
```

$O(1)$

$O(1)$

$O(1)$

$\rightarrow O(1)$

$\rightarrow O(1)$

$\rightarrow O(1)$

$\rightarrow O(1)$

$T.C. \rightarrow O(1)$

$S.C. \rightarrow O(1)$

```
count = 0;
for (i = 0; i < N; i++){
    count++;
}
j = N;
while(j >= 0){
    count++;
    j--;
}
```

```
count = 0;  
for (i = 0; i < N; i++){  
    count++;  
}  
for (j = 0; j < N; j++){  
    count++;  
}
```

How many times the above code will run?

```
sum = 0; //1
for (i=0;i<N;i++){ //2
    j = 0; //3
    while(j < N){ //4
        sum += i; //5
        j++; //6
    } //7
} //8
```

How many times will the above code run?

A

N

B

N^2

C

$N(\log N)$

D

$\log(N)$


```
count = 0;  
for (i = 0; i < N; i++){  
    for (j = 0; j < N; j++){  
        count++;  
    }  
}
```

How many times will the above code run?

A

N

B

N^2

C

$\text{Log}(N)$

D

None Of These