

These tests cover the three main behaviors of a cache: **Reuse (Locality)**, **Eviction (Capacity)**, and **Hierarchy (L2 Rescue)**.

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### Test Scenario 1: Basic Temporal Locality (The "Happy Path")

**Objective:** Verify that the simulator correctly identifies data that is already present in the cache (Hit) and does not fetch it from memory again.

#### 1. Configuration (In `CacheHierarchy.cpp`):

- **L1 Cache:** Size = 4096 (4KB), Block Size = 64, Associativity = 2
- **L2 Cache:** Size = 65536 (64KB), Block Size = 64, Associativity = 4  
(Standard configuration where L1 is large enough to hold a few items)

#### 2. Input Trace (In `main.cpp`):

C++

```
std::vector<unsigned long> memoryTrace = {  
    0x1000, // Access A  
    0x2000, // Access B  
    0x1000, // Access A (Again)  
    0x2000, // Access B (Again)  
    0x1000 // Access A (Again)  
};
```

#### 3. Expected Behavior:

- **Step 1 & 2:** Cold Misses. `0x1000` and `0x2000` are fetched from memory and stored in L1.
- **Step 3, 4, & 5: L1 Hits.** The simulator should find these addresses residing in L1.
- **L2 Activity:** L2 should have 0 Hits because L1 never missed after the initial load.

#### 4. Expected Output Statistics:

- **L1 Hits:** 3
- **L1 Misses:** 2
- **L2 Hits:** 0
- **L2 Misses:** 2

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### Test Scenario 2: The "Tiny L1" Eviction (Hierarchy Validation)

**Objective:** Verify that when L1 runs out of space, it evicts the oldest item (FIFO) to L2, and subsequent requests for that item are serviced by L2 (L1 Miss -> L2 Hit).

#### 1. Configuration (In `CacheHierarchy.cpp`):

- L1 Cache: Size = 128 (Bytes), Block Size = 64, Associativity = 2  
(CRITICAL: This limits L1 to exactly 2 cache lines)
- **L2 Cache:** Size = 4096 (4KB), Block Size = 64, Associativity = 4

## 2. Input Trace (In `main.cpp`):

```
C++
std::vector<unsigned long> memoryTrace = {
    0x1000, // 1. Fill Slot 1
    0x2000, // 2. Fill Slot 2 (L1 is now FULL)
    0x3000, // 3. New Data. Forces Eviction of 0x1000 (Oldest).
    0x1000 // 4. Retry 0x1000. It is gone from L1, must be found in L2.
};
```

## 3. Expected Behavior:

- **Step 3:** Causes a "Capacity Miss". L1 drops `0x1000` to make room for `0x3000`.
- **Step 4:** The request for `0x1000` fails in L1 (**Miss**) but succeeds in L2 (**Hit**), proving the hierarchy works correctly.

## 4. Expected Output Statistics:

- **L1 Hits:** 0
- **L1 Misses:** 4
- **L2 Hits:** 1 (The rescue of `0x1000`)
- **L2 Misses:** 3 (The initial cold loads)

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## Test Scenario 3: Worst-Case "Thrashing" (Conflict Misses)

**Objective:** Demonstrate the weakness of small associativity. If we alternate between 3 addresses that map to the same set in a 2-way cache, they will constantly kick each other out.

## 1. Configuration (In `CacheHierarchy.cpp`):

- **L1 Cache:** Size = 128 (Bytes), Block Size = 64, Associativity = 2 (Capacity is 2 lines)
- **L2 Cache:** Size = 4096 (4KB), Block Size = 64, Associativity = 4

## 2. Input Trace (In `main.cpp`):

```
C++
std::vector<unsigned long> memoryTrace = {
    0x1000, // Load A
    0x2000, // Load B (L1 Full: [A, B])
    0x3000, // Load C (Evict A. L1: [B, C])
    0x1000, // Load A (Evict B. L1: [C, A])
    0x2000, // Load B (Evict C. L1: [A, B])
    0x3000 // Load C (Evict A. L1: [B, C])
};
```

## 3. Expected Behavior:

- Every single step triggers an eviction.
- The cache is "Thrashing"—spending all its time loading and evicting without ever getting a hit in L1.
- However, L2 (which is larger/4-way) should catch all the re-accesses.

#### 4. Expected Output Statistics:

- **L1 Hits:** 0 (0% Hit Rate)
- **L1 Misses:** 6
- **L2 Hits:** 3 (The repeats of A, B, and C)
- **L2 Misses:** 3 (The first time A, B, and C were seen)