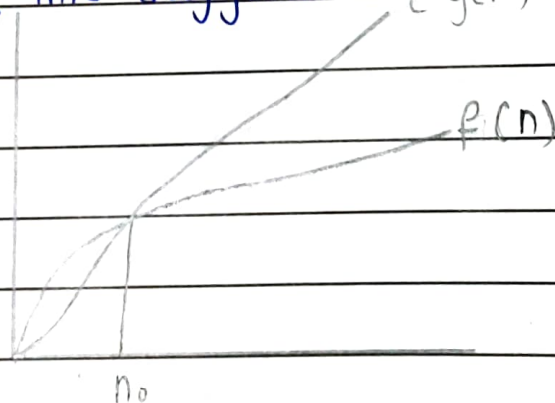


Q. 11

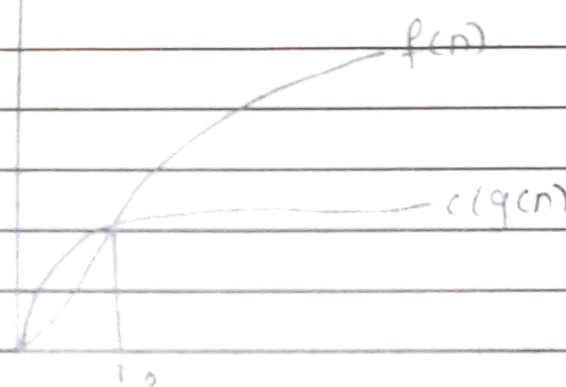
1] Asymptotic notions are used to check the time complexity of algorithms and to find out which algorithm is the most time efficient when input size is large enough.

2] Big O  $\rightarrow$  This is used to represent the worst case of the time complexity. This suggests the <sup>max</sup> time that an algorithm can take for completion.



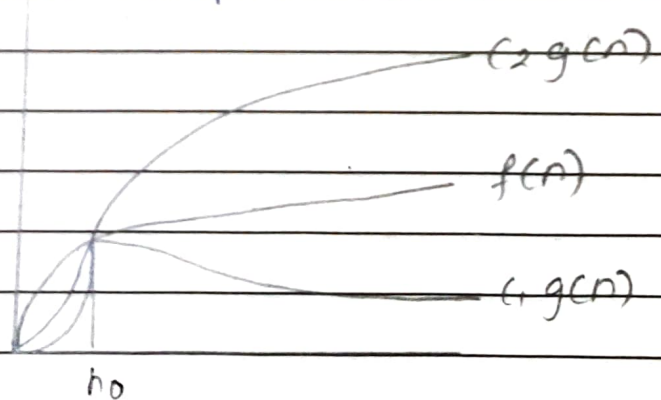
In the above graph  $c_g(n)$  and  $f(n)$  are functions and we can say that  $c_g(n) = O(f(n))$  if and only if  $c_g(n) \geq f(n)$  and  $n > n_0$ .

2] Big.  $\omega$   $\rightarrow$  This is used to represent the best of time complexity of an algorithm. This suggests the minimum time algorithm can take for its completion.



From the above graph we can say that  $f(n) = \omega(g(n))$  if and only if  $f(n) > c(g(n))$  and  $n > n_0$ .

3] Big  $\Theta$   $\rightarrow$  This is used to represent the average of time complexity of an algorithm. This suggests the average time an algorithm can take for its completion.



From above graph we can say that

$$c(g(n)) = f(n) = \Theta(c(g(n))) \text{ and } f(n) = \Theta(f(n))$$

$$c(g(n)) = \Theta(f(n)) \text{ if and only if}$$

Question  
Nos.Marks  
Awarded

$C(g(n)) \leq f(n) \leq C(g(n))$  and

$n > n_0$  for some  $n_0$  and for all  $n$  such that

$n > n_0$  and  $n$  is a power of 2, we have

$f(n) \leq C(g(n))$  and  $f(n) \geq C(g(n))$ .

Therefore,

Let  $n$  be a power of 2. Then, we have

$f(n) \leq C(g(n))$  and  $f(n) \geq C(g(n))$ .

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