

Q3) a) Infix to Postfix.

Expression : $(A + B * (C - D)^E) / (F - G + H * I)$

Symbol	Stack	Postfix
((
A	(A
+	(+	A +
B	(+	A B
*	(+ *	A B *
((+ * (A B *
C	(+ * (A B C
-	(+ * (-	A B C -
D	(+ * (-	A B C D
)	(+ *	A B C D -
^	(+ * ^	A B C D - ^
E	(+ * ^	A B C D - E
)		A B C D - E ^ * +
/	/	A B C D - E ^ * + /

Symbol	Stack	Postfix
(E / (ABCD - E ^ * +
F	E / (ABCD - E ^ * + F
-	/ (-	ABCD - E ^ * + F -
G	/ (-	ABCD - E ^ * + F G
+	/ (+	ABCD - E ^ * + F G -
H	/ (+	ABCD - E ^ * + F G - H
*	/ (+ *	ABCD - E ^ * + F G - H
I	/ (+ *	ABCD - E ^ * + F G - H I
)	/	ABCD - E ^ * + F G - H I * +

So the postfix expression is

ABCD - E ^ * + F G - H I * + /

Question
Nos.Q2) b) Circular Linked List.(1) Insertion at the Start.Step 1: Set $ptr = Start$ Step 2: If $ptr = NULL$ $ptr = \text{New node}$ ~~$\text{New node} = Start$~~ $\text{New node} \rightarrow \text{next} = Start$

exit.

Step 3: If $Start \neq NULL$.

Explanation.

where we point the head to the new node. & then point the new nodes next ptr to the first node.

~~this way~~

Q2) b) Circular Linked List.

(i) Insertion at the start.

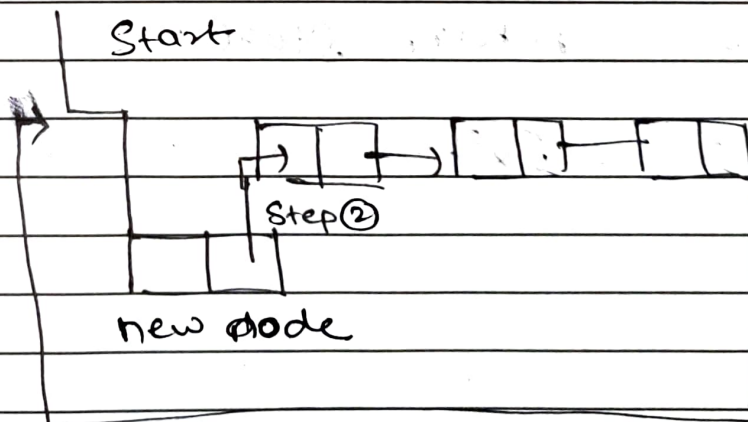
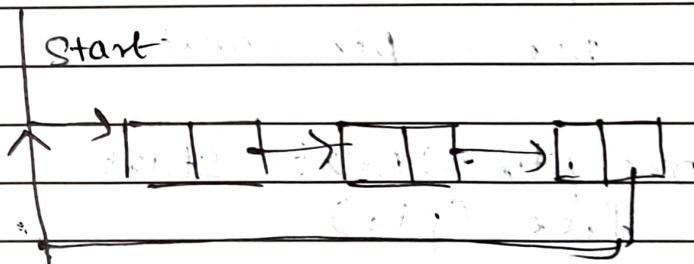
Step 1 : Set $ptr = \text{new Start (head)}$
Input $n = \text{data}$.

Step 2 : Set $\text{newnode} \rightarrow \text{data} = n$.
 $\text{new node} \rightarrow \text{next} = \text{start}$
 $ptr = \text{new node}$.

Step 3 : ~~exit~~ . return start

Step 4 : ~~exit~~

Example.



2.) Deletion of last node.

Step ① : Set $ptr = \text{head Start}$
 $pre\ ptr = ptr$
 Input $n = \text{data}$

Step ② : If $ptr = \text{NULL}$
 print ("List is Empty").
 Exit. Step 6.

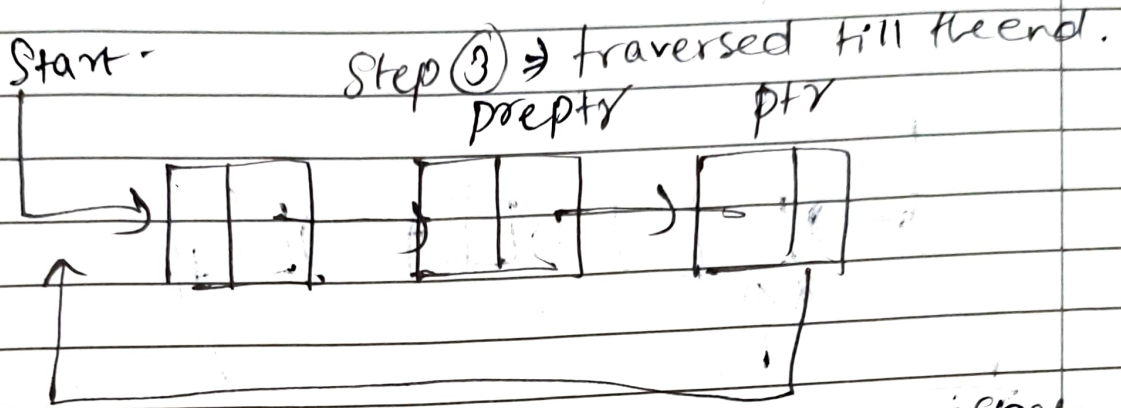
Step ③ : Else ~~if~~ $ptr \neq \text{NULL}$
~~do while~~ ($ptr \neq \text{start}$)
 do while ($ptr \neq \text{start}$)
 $pre\ ptr = ptr$
 $ptr = ptr \rightarrow \text{next}$

Step ④ : $pre\ ptr \rightarrow \text{next} = \text{start}$
 $\text{free}(ptr)$

Step ⑤ : return start .

Step ⑥ : Exit.

Example.



\Downarrow

we used `do while (ptr != null)`

because there are two instances when $ptr = \text{Start}$ once in the Start & once in the end. & what `do while` does is it atleast runs once so we pass the first instance when we got $ptr = \text{Start}$.
next instance when $ptr = \text{Start}$ is at the end of the list. This when the `do while` loop breaks.

we converge the predecessor to Start & free the ptr which is at the end of the node.
removing the last node from the list.

Q1) Asymptotic notations.

→ big O of n.

$O(n)$: ~~worst~~ worst case.

$O(n)$ represents the worst case scenario i.e. the algorithm can't take more ~~the~~ time than this case. Maximum time the algorithm will take to run. Highest time complexity.

→ big ~~Ω~~ Omega of n.

→ $\Omega(n)$: Best case.

→ $\Omega(n)$ represents the best case, i.e. the algorithm / code won't take lesser time in any other case. Minimum time / fastest algorithm can take. Lowest time complexity eg while sorting if you already are given a sorted list.

→ big ~~Θ~~ theta of n.

$\Theta(n)$: Average case.

$\Theta(n)$ represents the average case i.e. the algorithm will take avg time.