EE360C: Algorithms

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## Quiz #6

## Problem 1: MST [5 points]

Let G be a connected undirected graph with weighted edges, and suppose the weight of every edge of G is distinct. Could there be multiple different minimum spanning trees (i.e., there are some edges in one but not in the other)? Either give an example G with two different minimum spanning trees, or prove that there is exactly one minimum spanning tree in this case.

Hint: Consider two distinct MST  $T_1, T_2$  and use e which is the maximum weight edge in the set  $A = (T_1 \setminus T_2) \cup (T_2 \setminus T_1)$ .

## Solution

## Proof by Contradiction:

Let  $T_1$  and  $T_2$  be two distinct minimum spanning trees. Since they are distinct, there are some edges that appear in exactly one of  $T_1$  or  $T_2$ . Consider the edge of maximum weight among these edges. Without loss of generality, this edge appears only in  $T_1$ , and we can call it  $e_1$ .

Suppose we take  $T_1$  and remove  $e_1$ . This splits the tree into two connected components. Since  $e_1$  is not in  $T_2$ , there must be another edge  $e_2$  in  $T_2$  but not in  $T_1$  such that adding it would connect these two components. Thus taking  $T_1$ , removing  $e_1$ , and adding  $e_2$  forms a new spanning tree  $T'_1$ . Since  $e_2$  is an edge different from  $e_1$  and is contained in exactly one of  $T_1$  or  $T_2$ , it must be that  $w(e_1) < w(e_2)$ . Since  $T'_1$  is the same as  $T_1$  except for removing  $e_1$  and adding  $e_2$ , the total weight of  $T'_1$  is smaller than the total weight of  $T_1$ . But this is a contradiction, since we have supposed that  $T_1$  is a minimum spanning tree.