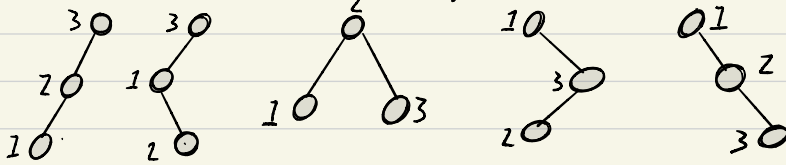


A) Binary Search Trees

- 5 distinct binary search trees on 3 nodes



• Rooted Binary tree

- Either empty or consists of a node called the root, together with 2 rooted binary trees called the left and right subtrees
- The order among subtrees matter so we can distinguish b/t the left & right nodes

• Binary Search Tree

- labels each node in a binary tree with a single key such that for any node labeled x , all nodes in the left subtree of x have keys $< x$ while all nodes in the right subtree of x have keys $> x$

B) Implementing Binary search Trees

- Binary tree nodes have left & right pointer fields, an optional parent pointer, & a data field

- Basic Operations

i) Searching

- Start at the the root
- Unless it contains the query key x , proceed left or right depending on whether x occurs before or after the root key
- $O(h)$ where h = height of tree

ii) Finding minimum & maximum

□ Minimum

- By definition the smallest key must reside in the left subtree of the root since all keys in the left subtree have values less than that of the root

- Therefore, the min element must be the leftmost descendent of the root

□ Maximum

- Maximum element must be the rightmost descendent of the root by the definition of minimum.

iii) Traversal in a Tree

- Each item is processed once
- Runs in $O(n)$ where $n = \# \text{ of nodes}$ in the tree

iv) Insertion

- Only 1 place to insert an item x into a binary search tree
- Must replace the NULL pointer found in T after an unsuccessful query for the key k
- The implementation uses recursion to combine the search and node insertion stages of key insertion
- After the search is implemented, allocating the node & linking it to the tree is a constant-time operation

c) Balanced Search Tree

- Red-black trees and splay trees