PCFGs: Parsing & Evaluation

LING 571 — Deep Processing Techniques for NLP
October 12, 2022
Shane Steinert-Threlkeld

Announcements

- HW2 due Friday at 11:59pm
 - readme.{txtlpdf}
 - Separate upload to Canvas
 - NOT in hw2.tar.gz
 - Run check_hw2.sh before submitting!
- Start symbol: either "%start S" or first nonterminal
 - NB: needs to be readable by nltk's grammar loading methods
- nltk.data.load: "file:path/to/grammar.cfg"
- Condor job: runs conversion, but not parsing before/after; that can be done separately

Roadmap

- CKY + back-pointers
- PCFGs
- PCFG Parsing (PCKY)
- Inducing a PCFG
- Evaluation
- [Earley parsing]
- HW3 + collaboration

CKY Parsing: Backpointers

Backpointers

- Instead of list of possible nonterminals for that node, each cell should have:
 - Nonterminal for the node
 - Pointer to left and right children cells
 - Either direct pointer to cell, or indices

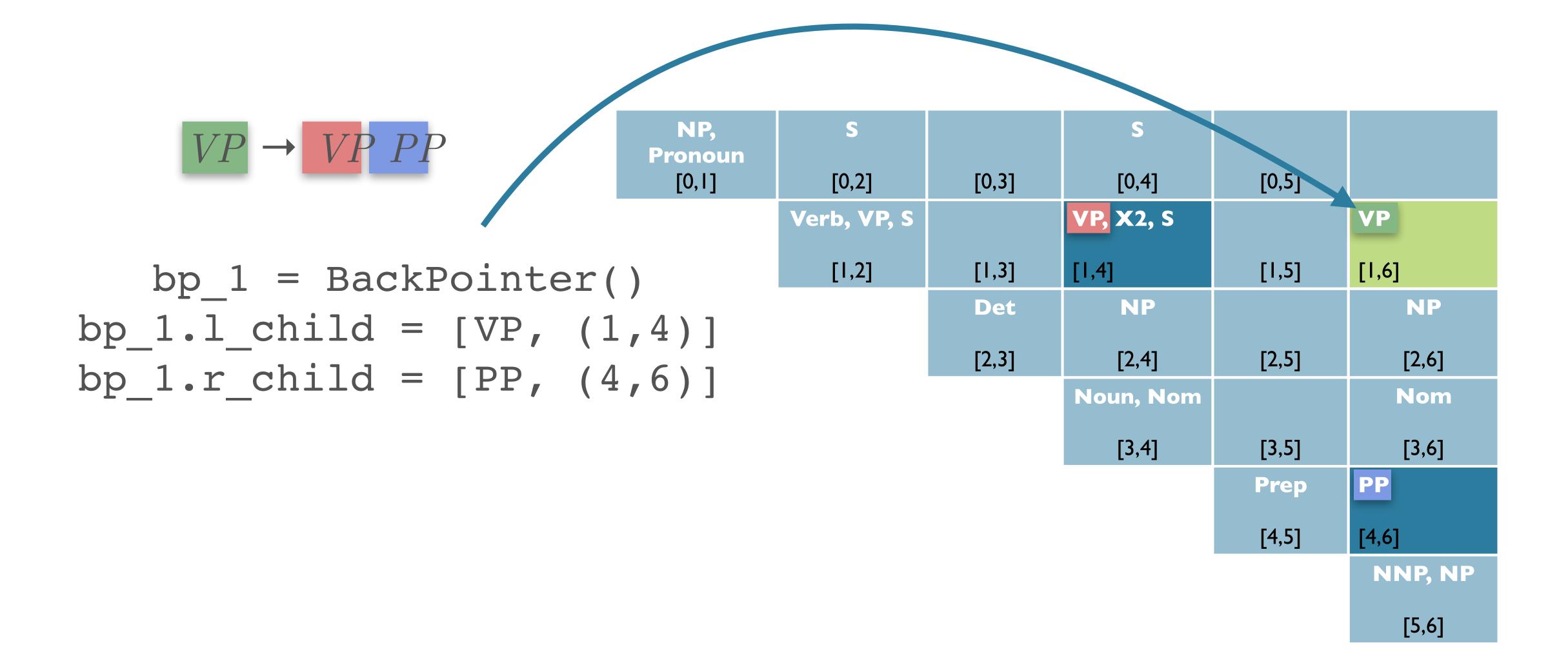
For example:

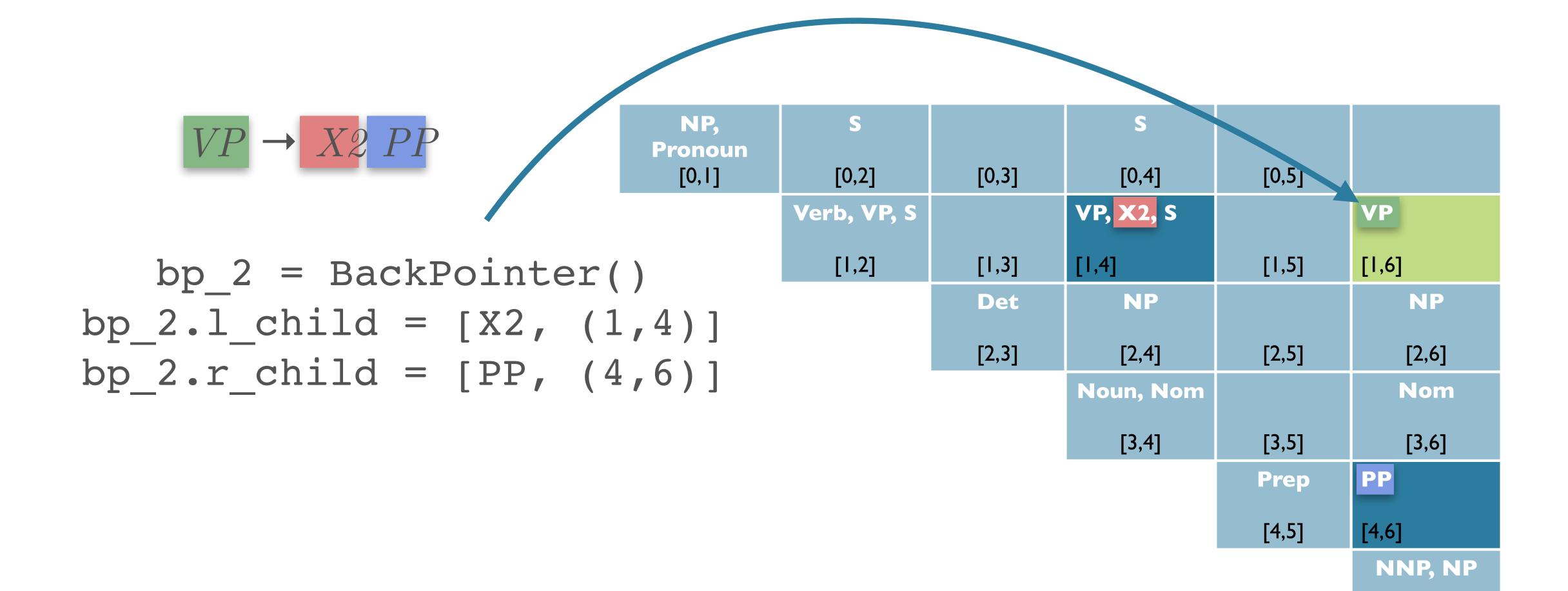
```
bp_2 = BackPointer()
bp_2.l_child = [X2, (1,4)]
bp 2.r child = [PP, (4,6)]
```

CKY Parser

- Pair each nonterminal with back-pointer to cells from which it was derived
- Last step:
 - construct trees from back-pointers in [0, n]

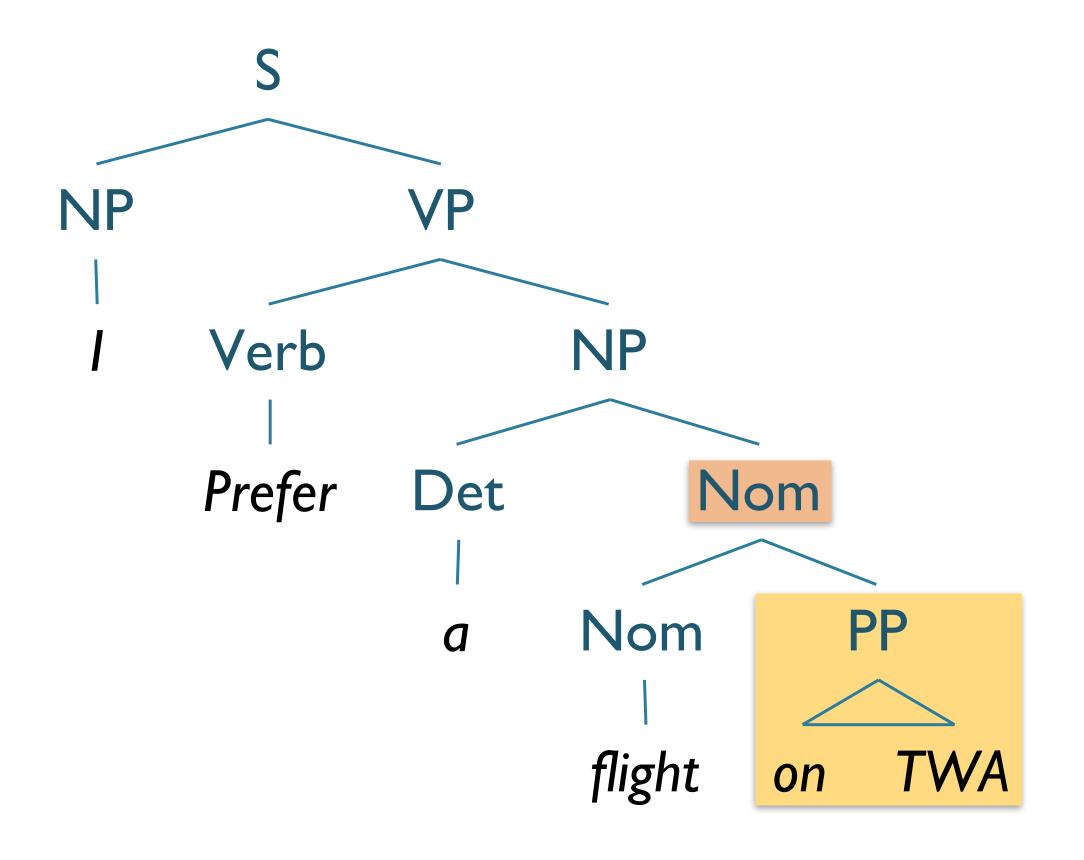
| NP, Pronoun | S | | S | | |
|----------------|-------------|-------|-----------|-------|---------|
| [0,1] | [0,2] | [0,3] | [0,4] | [0,5] | |
| | Verb, VP, S | | VP, X2, S | | VP |
| | [1,2] | [1,3] | [1,4] | [1,5] | [1,6] |
| | | Det | NP | | NP |
| | | [2,3] | [2,4] | [2,5] | [2,6] |
| | | | Noun, Nom | | Nom |
| | | | [3,4] | [3,5] | [3,6] |
| | | | | Prep | PP |
| | | | | [4,5] | [4,6] |
| | | | | | NNP, NP |
| | | | | | [5,6] |

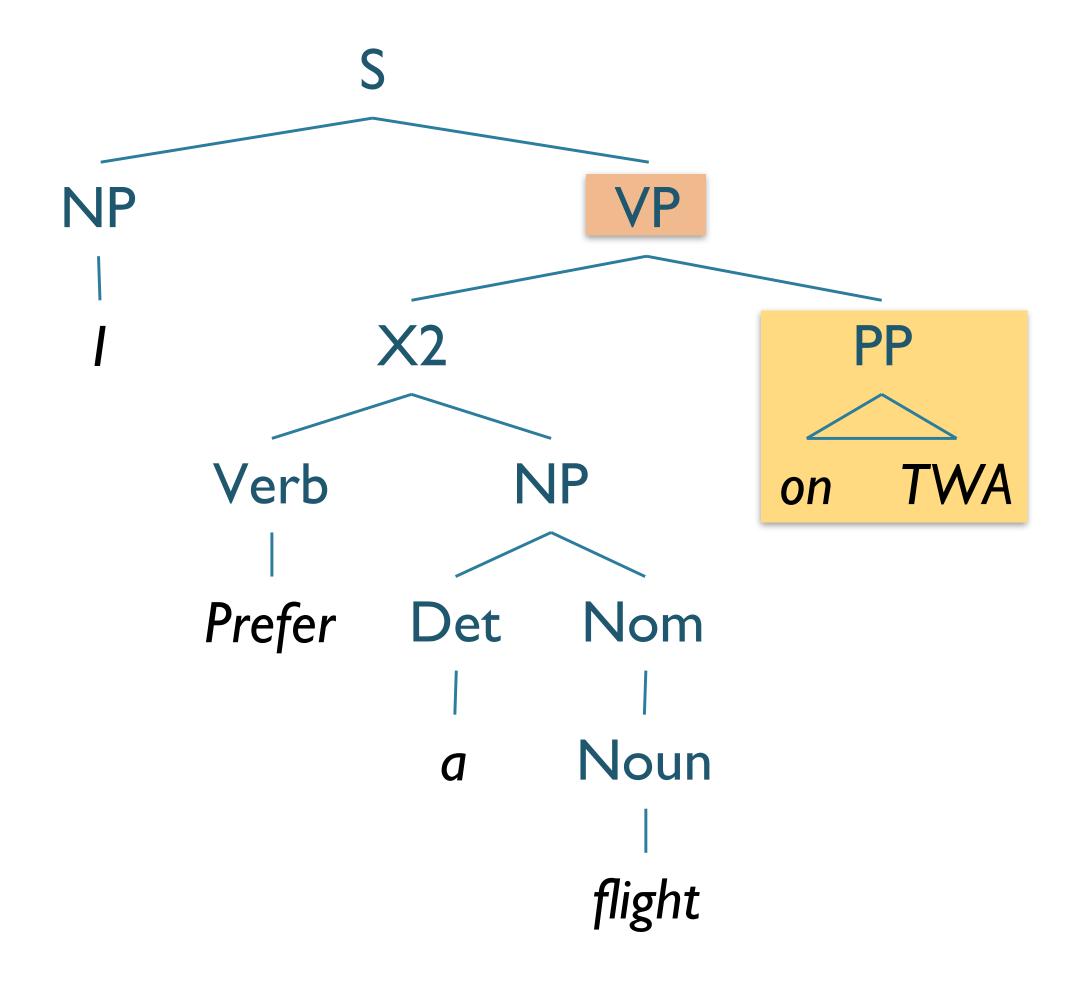




[5,6]

Resulting Parses

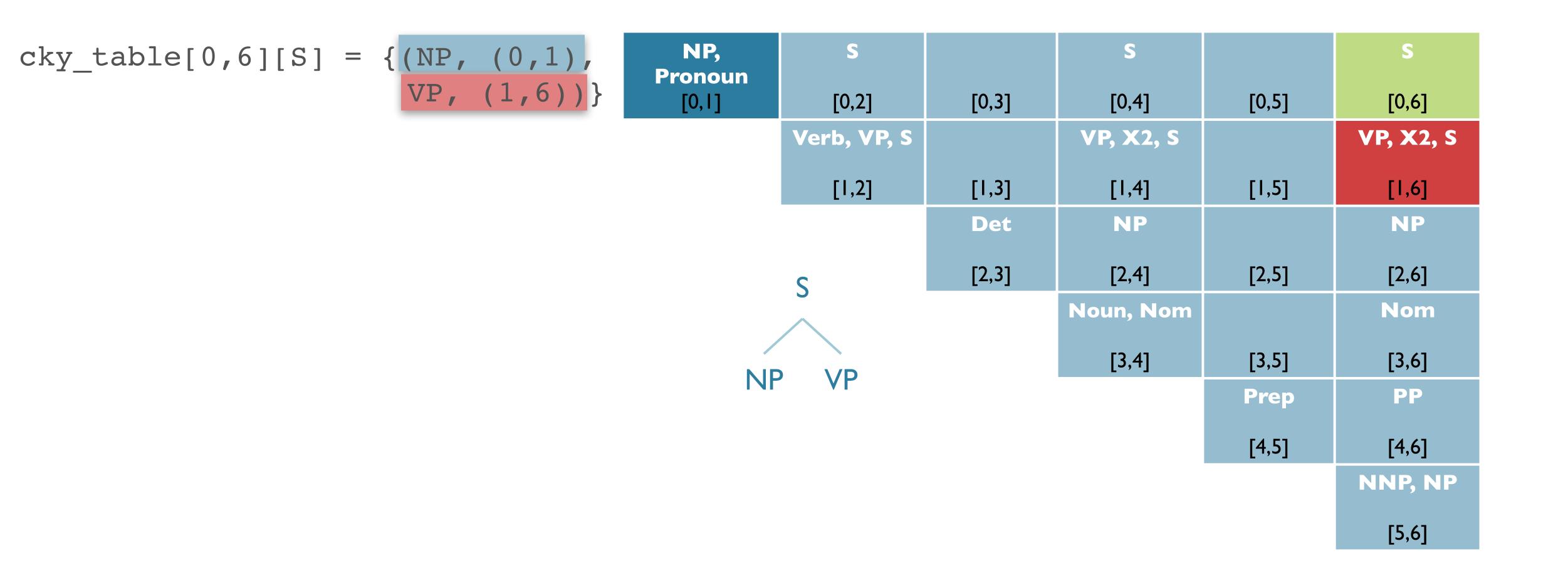


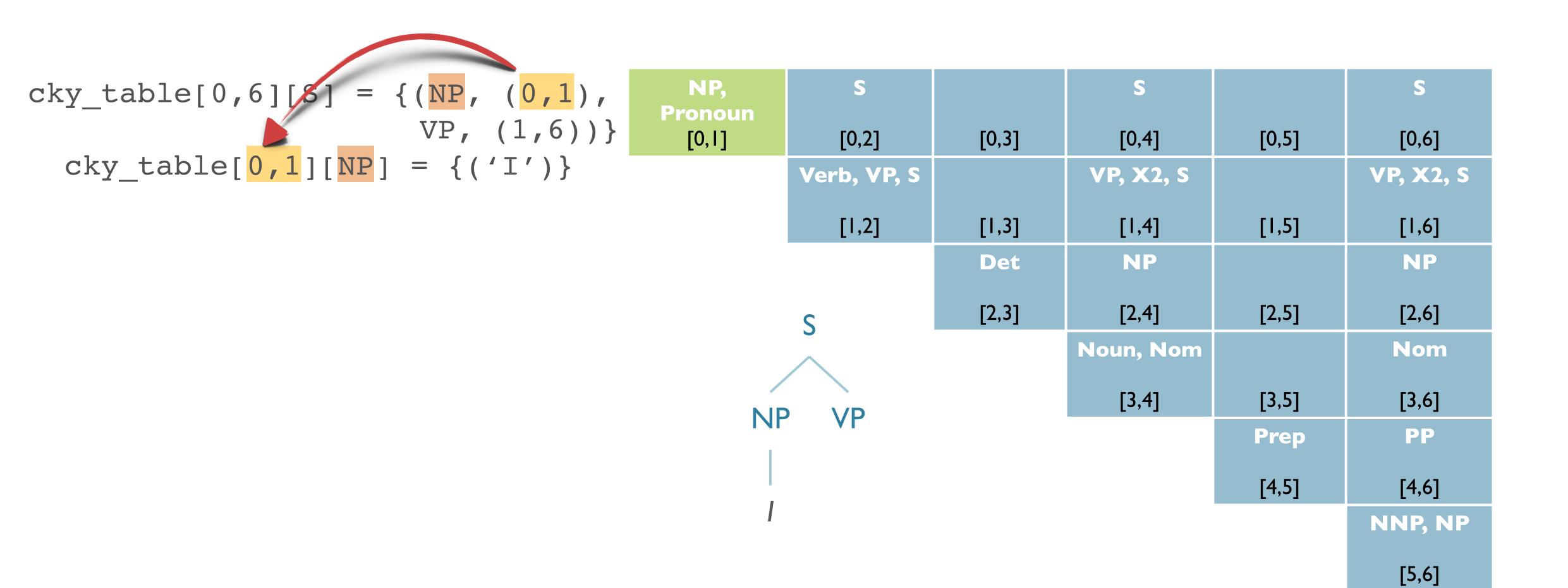


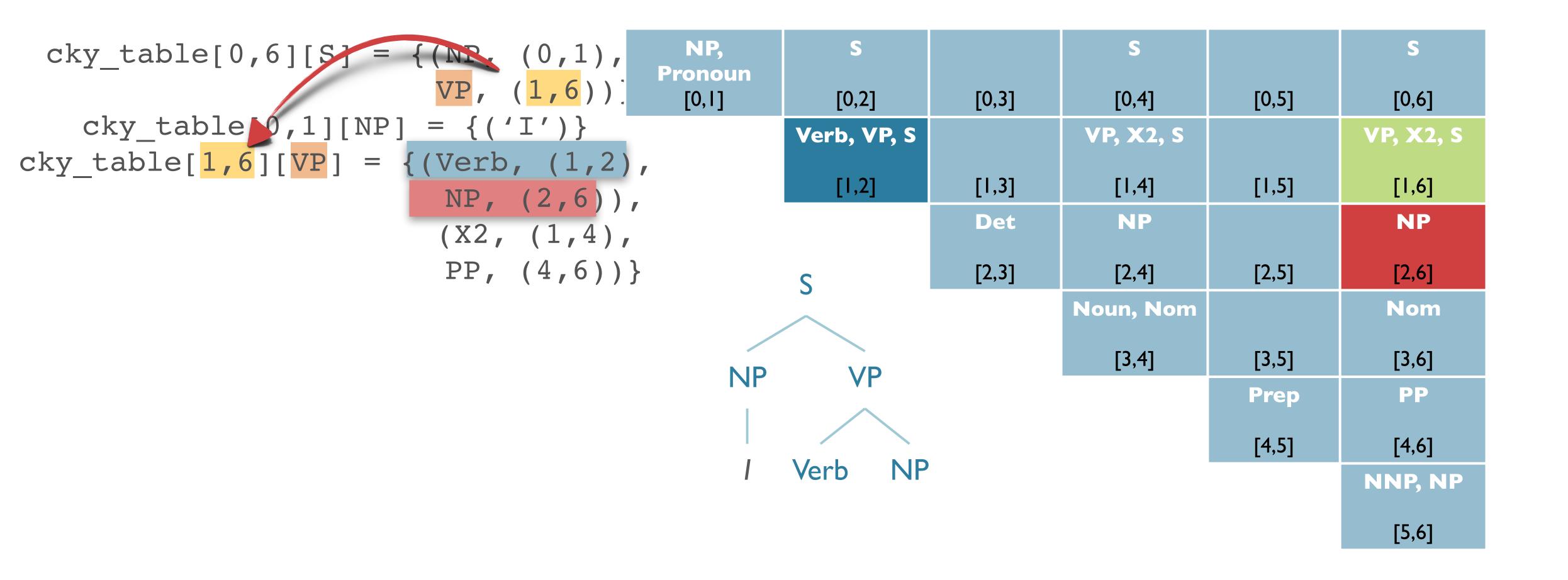
CKY Discussion

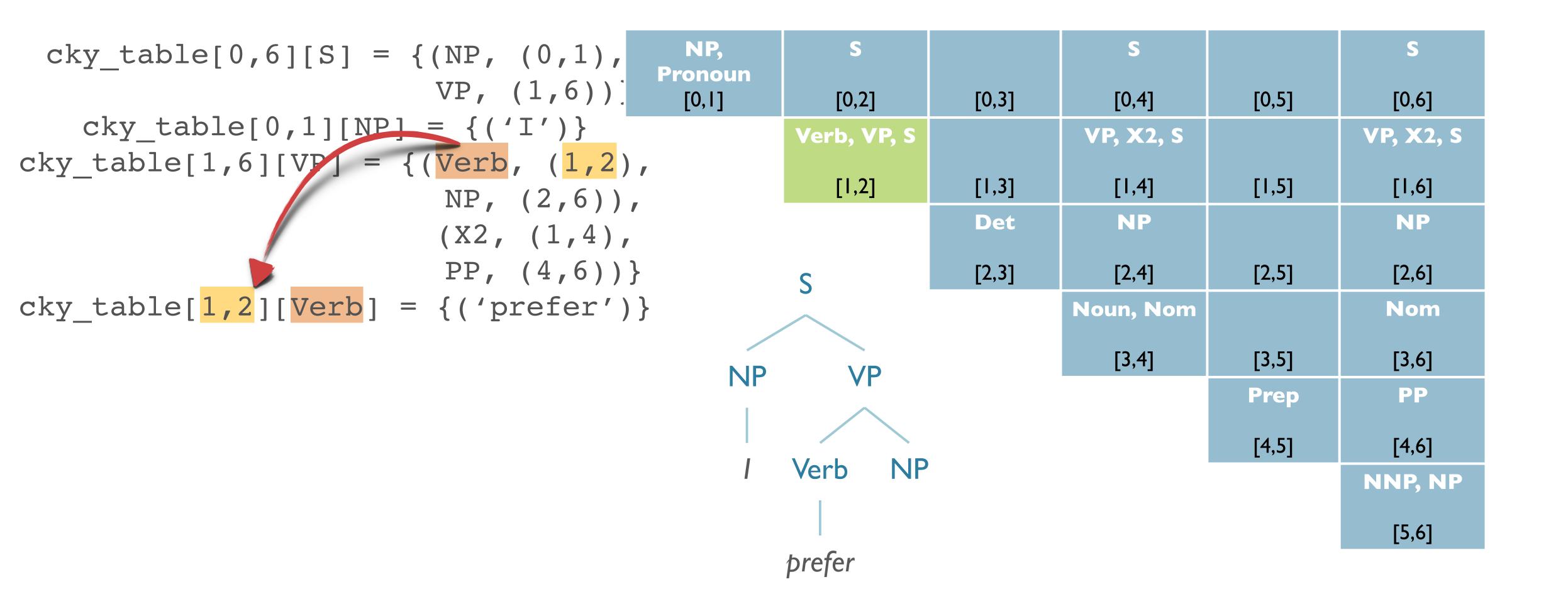
- Running time:
 - ullet $O(n^3)$ where n is the length of the input string
 - Inner loop grows as square of # of non-terminals
- Expressiveness:
 - As implemented, requires CNF
 - Weak equivalence to original grammar
 - Doesn't capture full original structure
 - Back-conversion?
 - Can do binarization, terminal conversion
 - Unit productions requires change in CKY

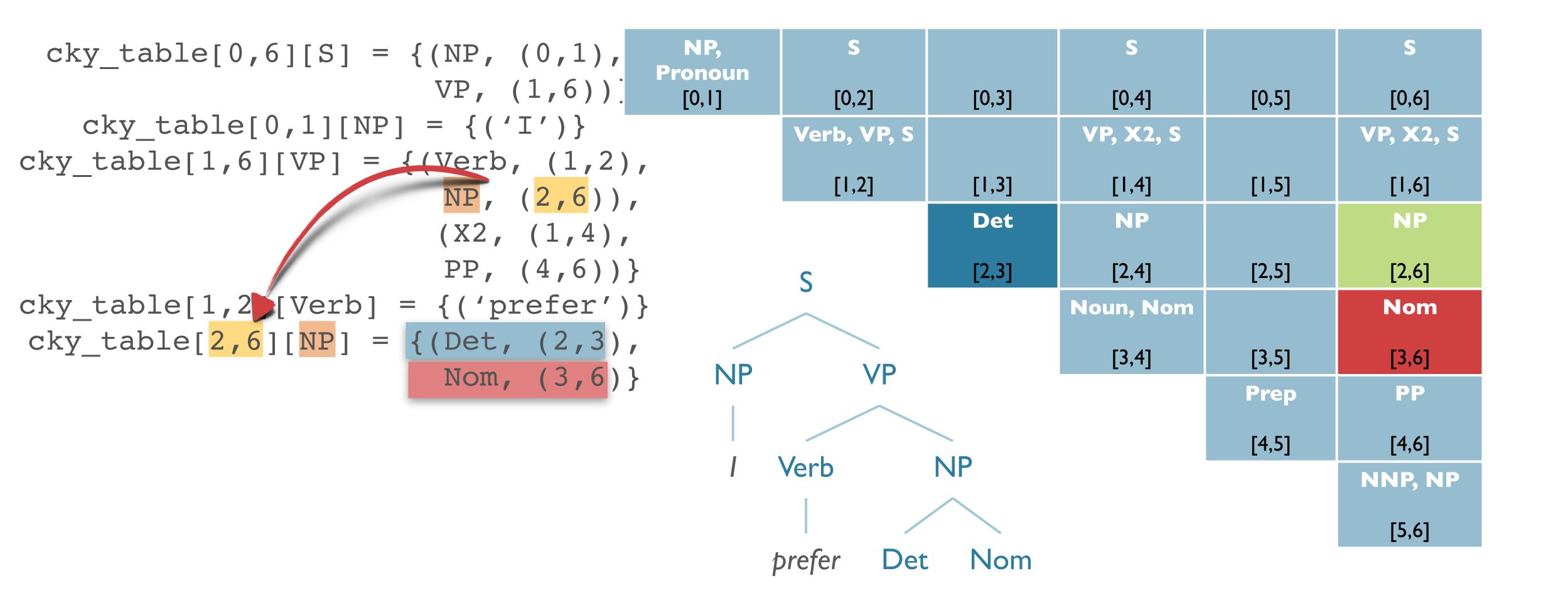
CKY + Back-pointers Example

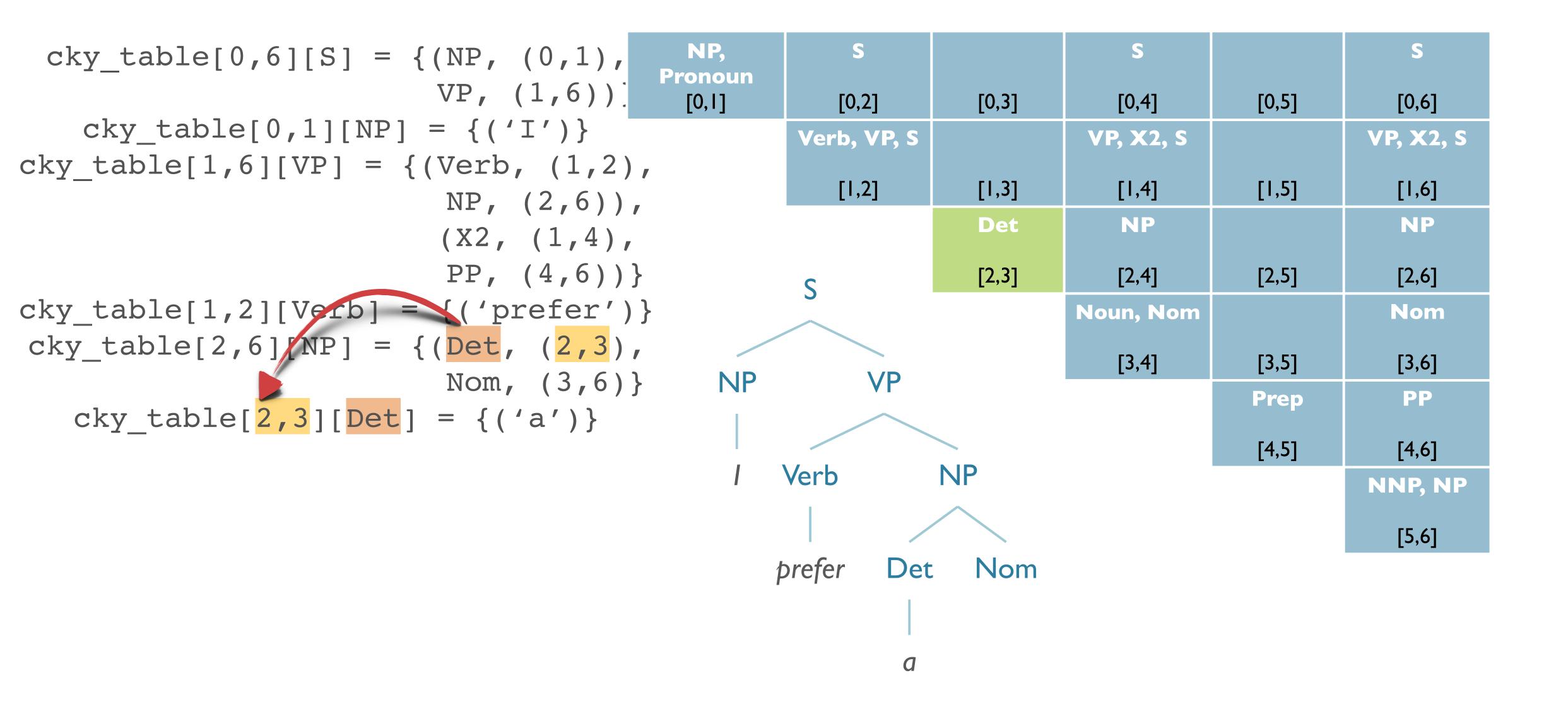




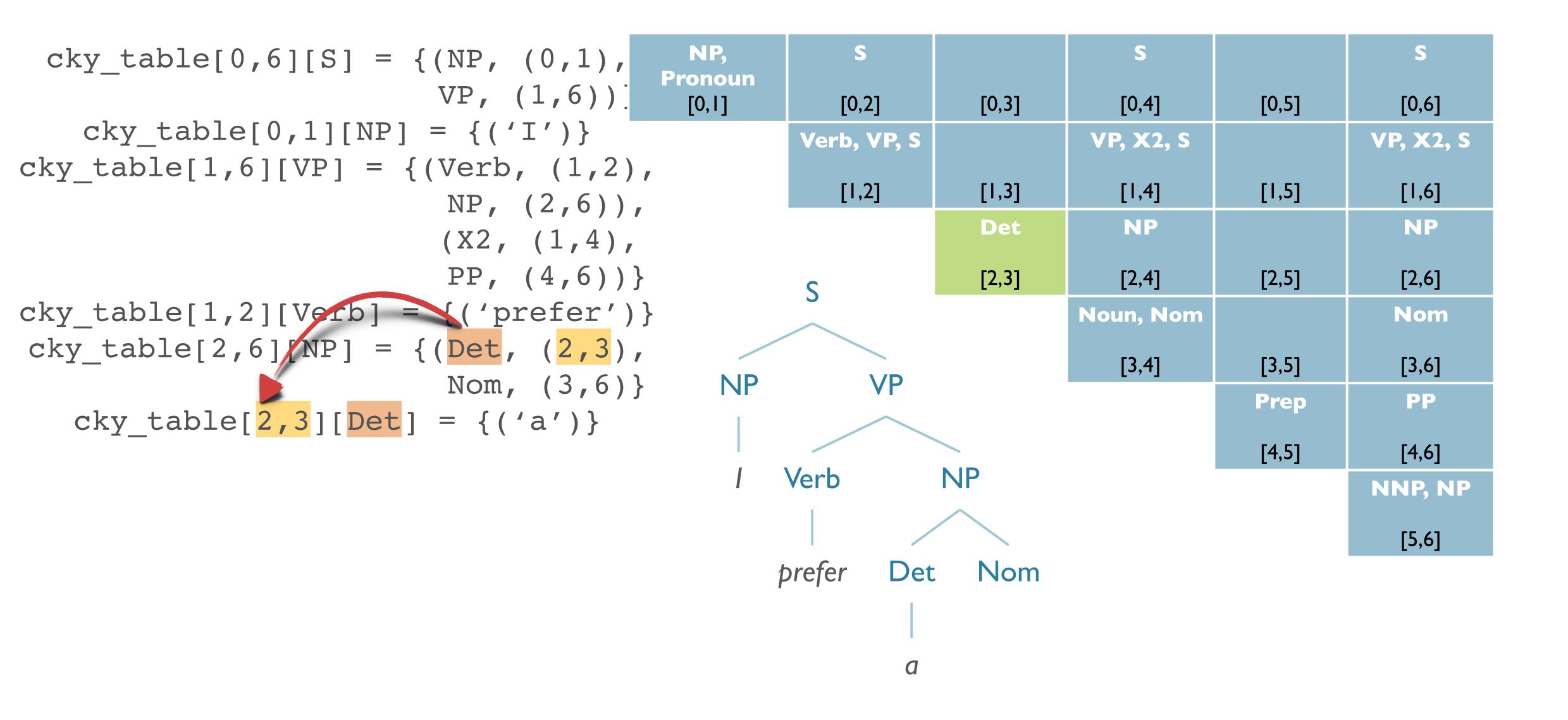


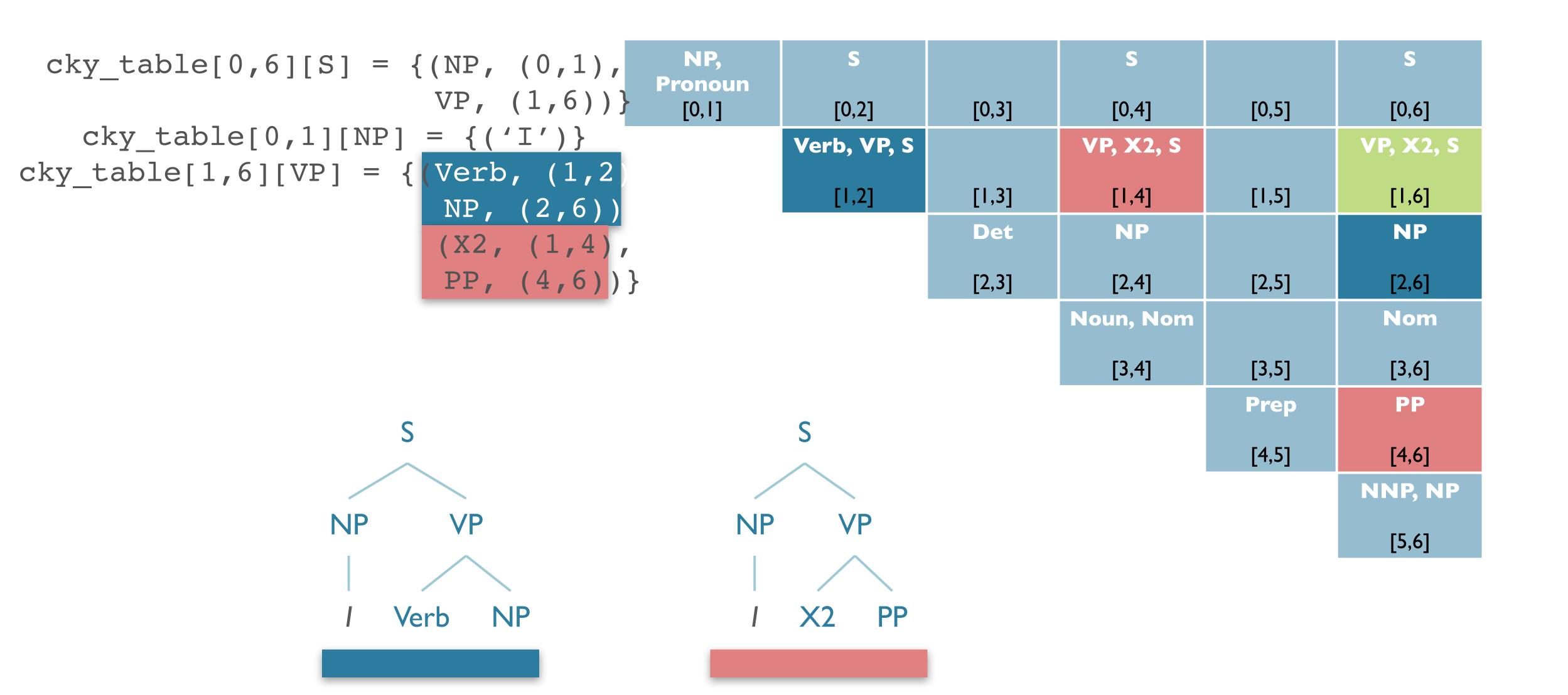






TWA





Probabilistic Context-Free Grammars

Probabilistic Context-free Grammars: Roadmap

Motivation: Ambiguity

Approach:

Definition

Disambiguation

Parsing

Evaluation

Enhancements

Motivation

What about ambiguity?

Current algorithm can *represent* it

...can't resolve it.

Probabilistic Parsing

- Provides strategy for solving disambiguation problem
 - Compute the probability of all analyses
 - Select the most probable

- Employed in language modeling for speech recognition
 - N-gram grammars predict words, constrain search
 - Also, constrain generation, translation

N

a set of non-terminal symbols (or variables)

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|--------|---|--|
| \sum | a set of terminal symbols (disjoint from N) | |

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| R | a set of rules of productions, each of the form $A	o eta[p]$, where A is a non-terminal where A is a non-terminal, eta is a string of symbols from the infinite set of strings $(\Sigma\cup N)_*$ and p | | |
| | is a number between 0 and 1 expressing $P(oldsymbol{eta} oldsymbol{A})$ | | |

| N | a set of non-terminal symbols (or variables) |
|----------------|---|
| \sum | a set of terminal symbols (disjoint from N) |
| R | a set of rules of productions, each of the form $A 	o \beta[p]$, where A is a non-terminal where A is a non-terminal, β is a string of symbols from the infinite set of strings $(\Sigma \cup N)*$ and p is a number between 0 and 1 expressing $P(\beta A)$ |
| \overline{S} | a designated start symbol |

PCFGs

- Augment each production with probability that LHS will be expanded as RHS
 - $\bullet P(A \rightarrow \beta)$
 - $\bullet P(A \rightarrow \beta | A)$
 - $\bullet P(\beta|A)$
 - \bullet $P(RHS \mid LHS)$
- NB: the first is often used; but the latter are what's really meant.

PCFGs

Sum over all possible expansions is 1

$$\sum_{\beta} P(A \to \beta) = 1$$

- A PCFG is <u>consistent</u> if sum of probabilities of all sentences in language is 1
 - Recursive rules often yield inconsistent grammars (Booth & Thompson, 1973)

Example PCFG: Augmented \mathcal{L}_1

| | Lexicon |
|-------|--|
| [.80] | $Det \rightarrow that [.10] \mid a [.30] \mid the [.60]$ |
| [.15] | $Noun \to book [.10] \mid flight [.30] \mid meal [.15] \mid money [0.95]$ |
| [.05] | $\mid flights \ [0.30] \mid dinner \ [.10]$ |
| [.35] | $Verb \rightarrow book [.30] \mid include [.30] \mid prefer [.40]$ |
| [.30] | $Pronoun \rightarrow I[.40] \mid she[.05] \mid me[.15] \mid you[.40]$ |
| [.20] | $Proper-Noun \rightarrow Houston [.60] \mid NWA [.40]$ |
| [.15] | $Aux \rightarrow does [.60] \mid can [.40]$ |
| [.75] | $Preposition \rightarrow from [.30] \mid to [.30] \mid on [.20] \mid near [.15]$ |
| [.20] | $\mid through \; [.05]$ |
| [.05] | |
| [.35] | |
| [.20] | |
| [.10] | |
| [.15] | |
| [.05] | |
| [.15] | |
| [1.0] | |
| | [.15] [.05] [.35] [.30] [.20] [.15] [.75] [.20] [.35] [.20] [.10] [.15] [.05] [.15] |

Example PCFG: Augmented \mathcal{L}_1

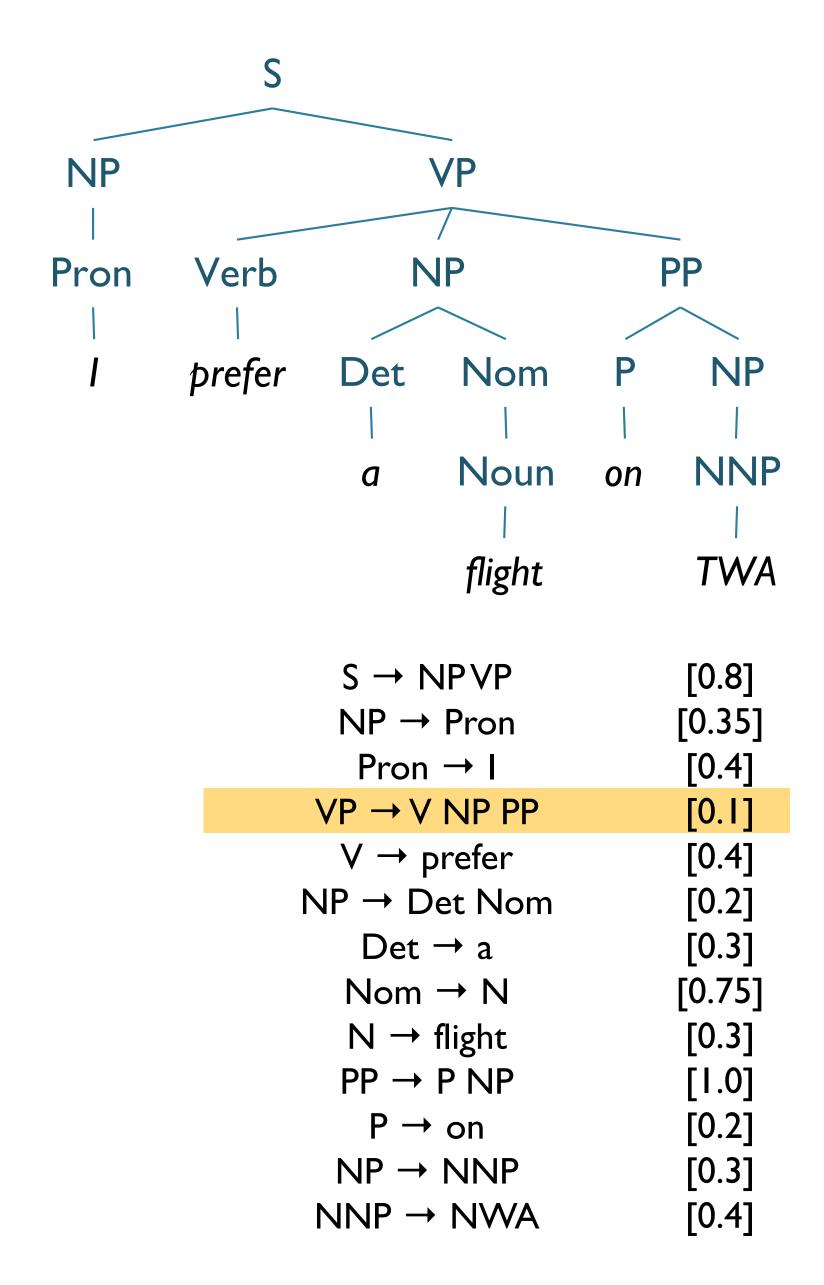
| Grammar | | | Lexicon |
|---------|--------------------------------------|-------|--|
| | $S \rightarrow NP VP$ | [.80] | $Det \rightarrow that [.10] \mid a [.30] \mid the [.60]$ |
| | $S \rightarrow Aux NP VP$ | [.15] | $Noun \rightarrow book [.10] \mid flight [.30] \mid meal [.15] \mid money [0.5]$ |
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| | $Nominal \rightarrow Nominal \ Noun$ | [.20] | $\mid through \; [.05]$ |
| | $Nominal \rightarrow Nominal PP$ | [.05] | |
| | $VP \rightarrow Verb$ | [.35] | |
| | $VP \rightarrow Verb NP$ | [.20] | |
| | $VP \rightarrow Verb NP PP$ | [.10] | |
| | $VP \rightarrow Verb PP$ | [.15] | |
| | $VP \rightarrow Verb NP NP$ | [.05] | |
| | $VP \rightarrow VP PP$ | [.15] | |
| | $PP \rightarrow Preposition NP$ | [1.0] | |
| | | | |

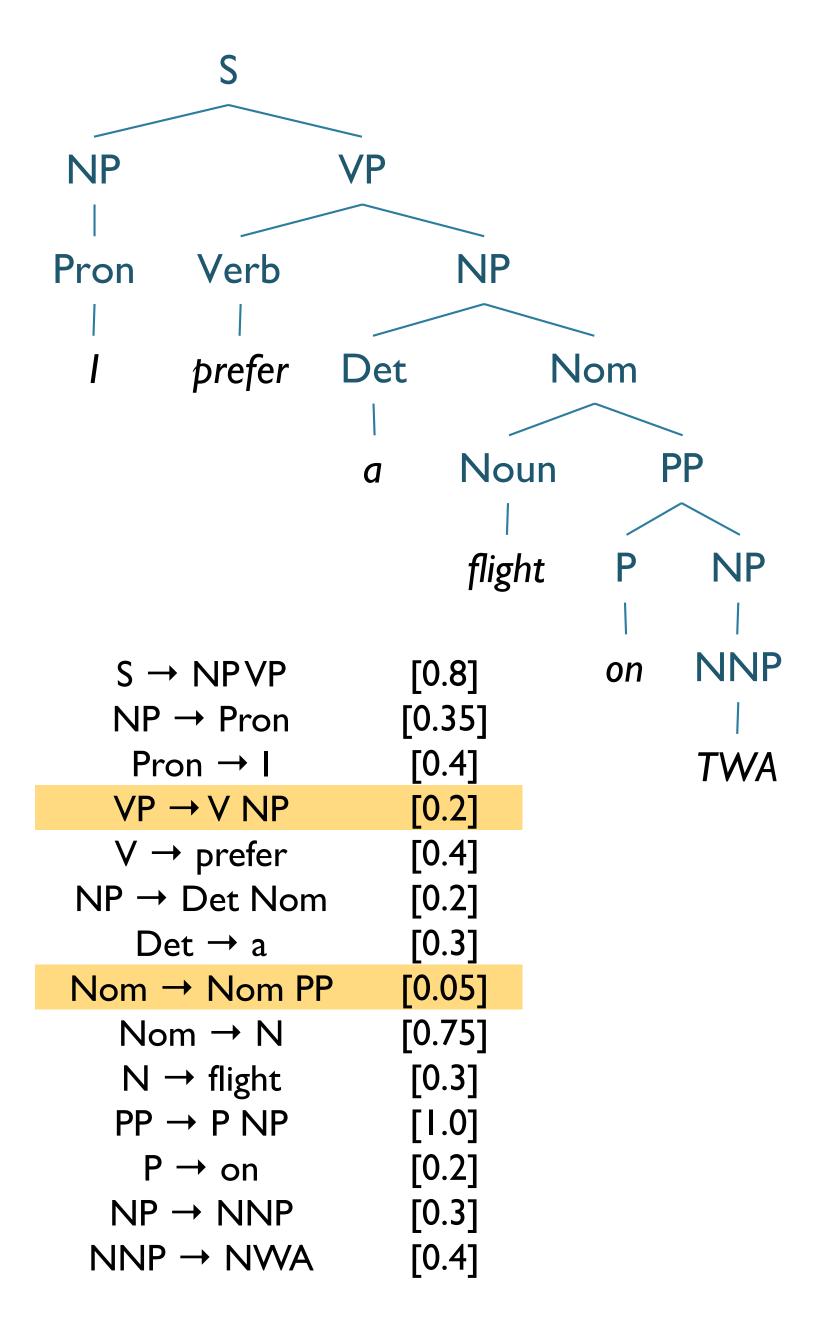
Disambiguation

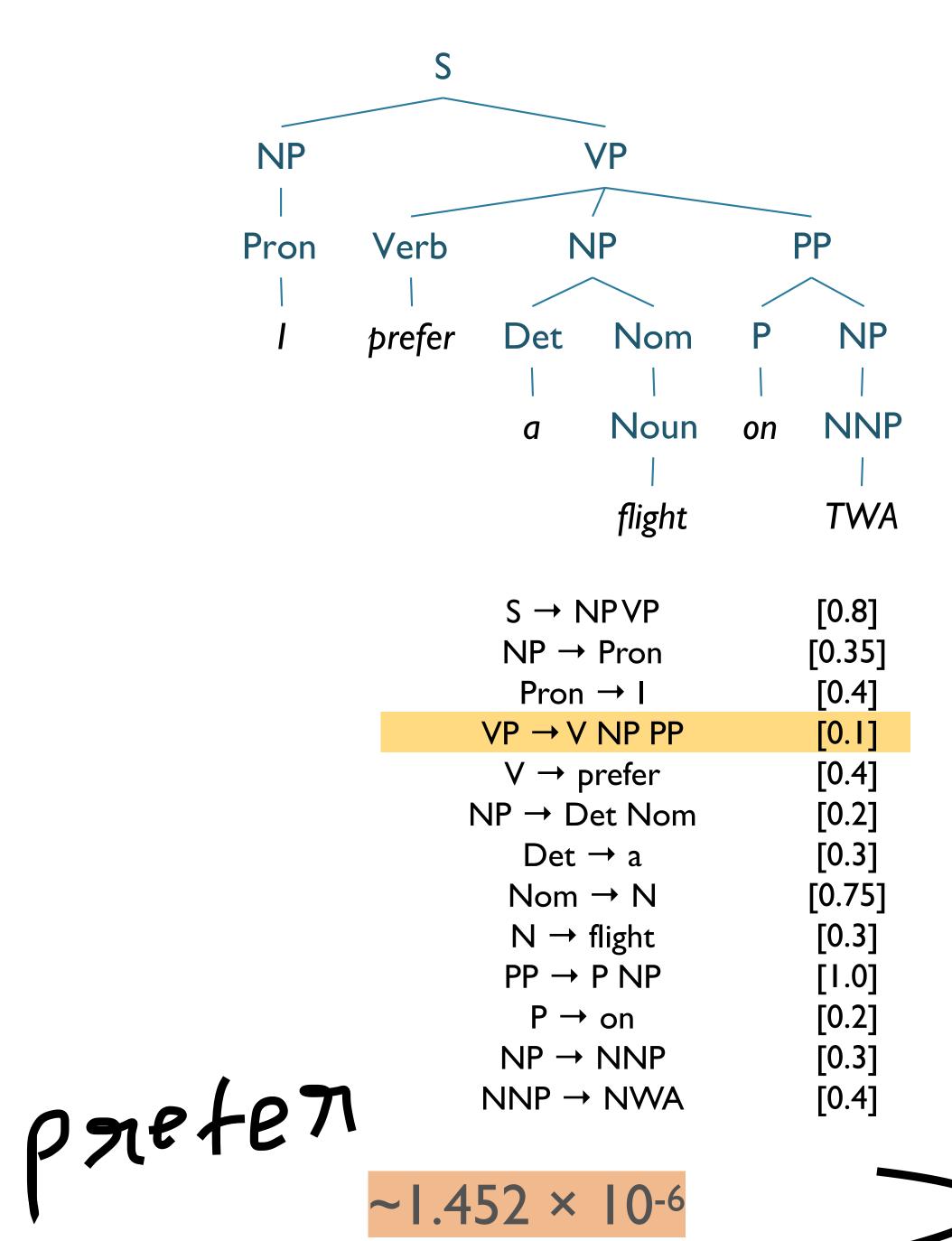
- ullet A PCFG assigns probability to each parse tree T for input S
- ullet Probability of T: product of all rules used to derive T

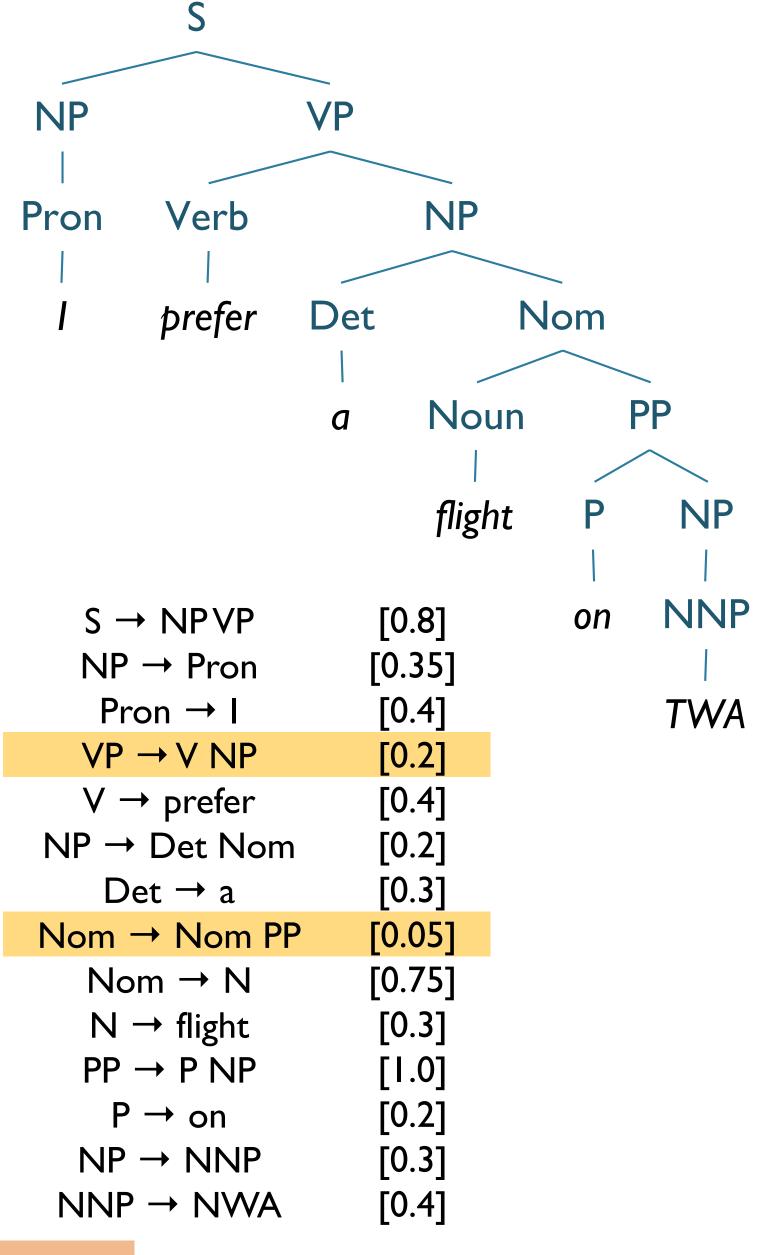
$$P(T,S) = \prod_{i=1}^{n} P(RHS_i | LHS_i)$$

$$P(T, S) = P(T)P(S | T) = P(T)$$









 $\sim 1.452 \times 10^{-7}$

Parsing Problem for PCFGs

• Select T such that (s.t.)

$$\hat{T}(S) = \underset{T \text{ s.t. } S = yield(T)}{\operatorname{argmax}} P(T)$$

- String of words S is yield of parse tree
- Select the tree \hat{T} that maximizes the probability of the parse

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- PCFGs are able to give probability of entire string without as bad sparsity
- Model probability of syntactically valid sentences
 - Not just probability of sequence of words

PCFGs: Parsing

Probabilistic CKY (PCKY)

- Like regular CKY
 - Assumes grammar in Chomsky Normal Form (CNF)
 - \bullet $A \rightarrow B C$
 - \bullet $A \rightarrow w$
 - Represent input with indices b/t words:
 - Book 1 that 2 flight 3 through 4 Houston 5

Probabilistic CKY (PCKY)

- ullet For input string length n and non-terminals V
 - Cell [i, j, A] in (n+1) × (n+1) × V matrix
 - Contains probability that A spans [i, j]

```
function Probabilistic-CKY-PARSE(words, grammar) returns most probable parse and its probability
for j ← from 1 to LENGTH(words) do
for all \{A \mid A \rightarrow words[j] \in grammar\}
       table[j-1, j, A] \leftarrow P(A \rightarrow words[j])
 for i \leftarrow \text{from } j-2 \text{ downto } 0 \text{ do}
  for k \leftarrow i + 1 to j-1 do
  for all \{A \mid A \rightarrow B \ C \in grammar,
         and table[i, k, B] > 0 and table[k, j, C] > 0
  if [table[i, j, A]] < P(A \rightarrow BC) \times table[i, k, B] \times table[k, j, C]) then b \in \mathcal{S} \subset \mathcal{S} \subset \mathcal{S}
      table[i, j, A] \leftarrow P(A \rightarrow BC) \times table[i, k, B] \times table[k, j, C]
                                                                                choose
     back[i, j, A] \leftarrow \{k, B, C\}
  return Build_Tree(back[1, Length(words), S]), table[1, Length(words), S]
```

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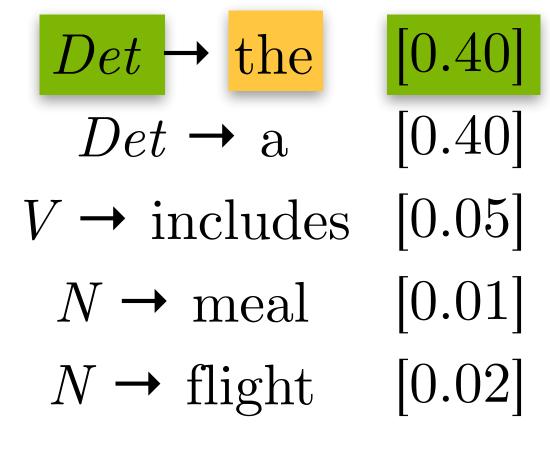
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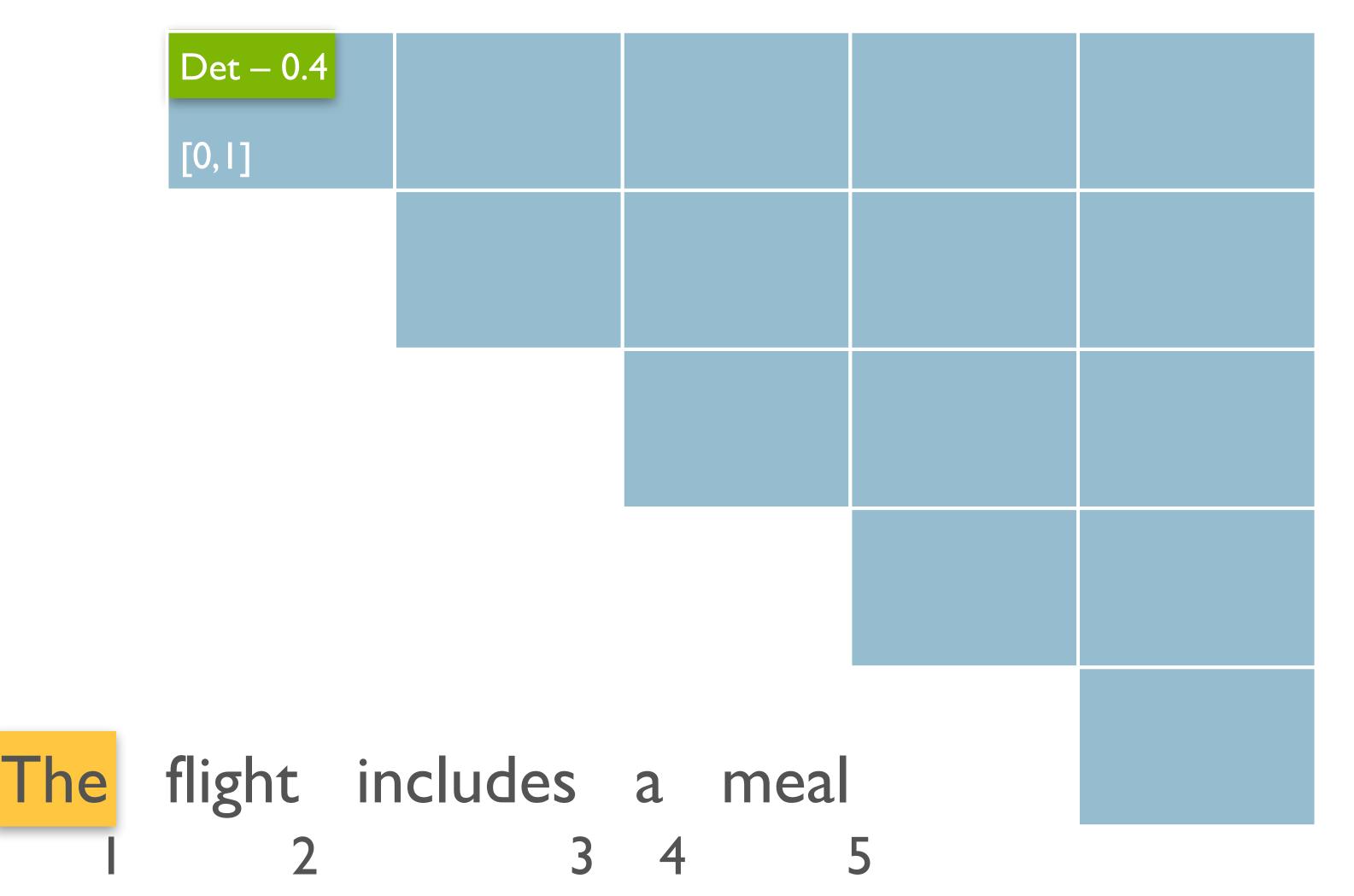
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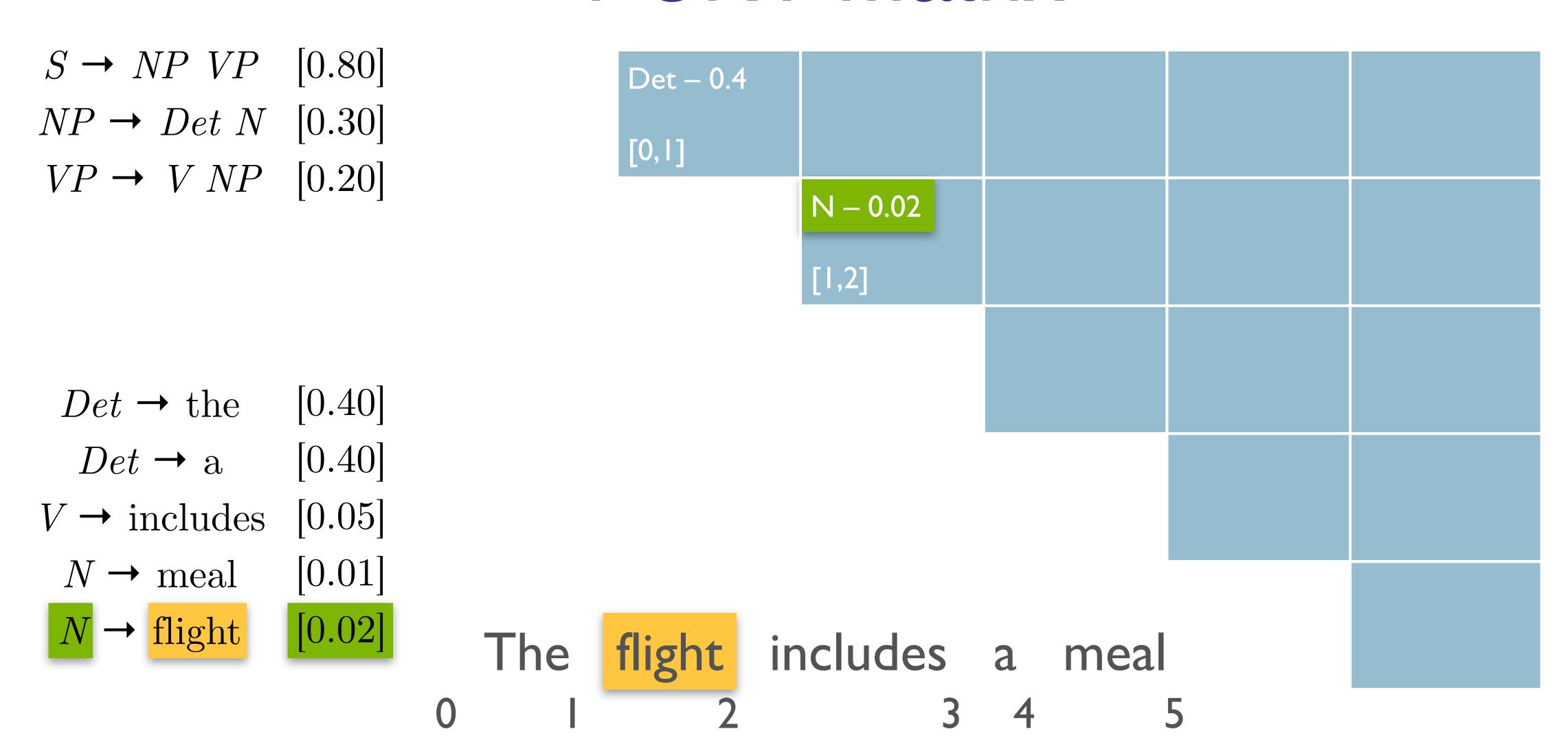
PCKY Grammar Segment

```
S \rightarrow NP \ VP \quad [0.80] Det \rightarrow the \quad [0.40] NP \rightarrow Det \ N \quad [0.30] Det \rightarrow a \quad [0.40] VP \rightarrow V \ NP \quad [0.20] V \rightarrow includes \quad [0.05] N \rightarrow meal \quad [0.01] N \rightarrow flight \quad [0.02]
```

$$S \rightarrow NP \ VP \quad [0.80]$$
 $NP \rightarrow Det \ N \quad [0.30]$
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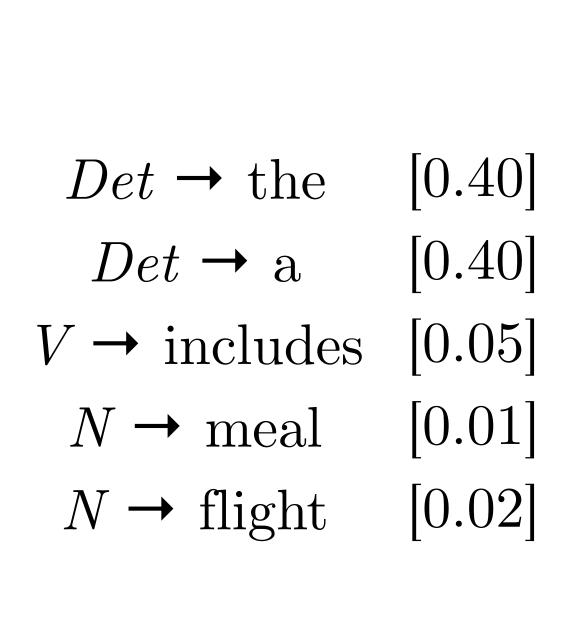


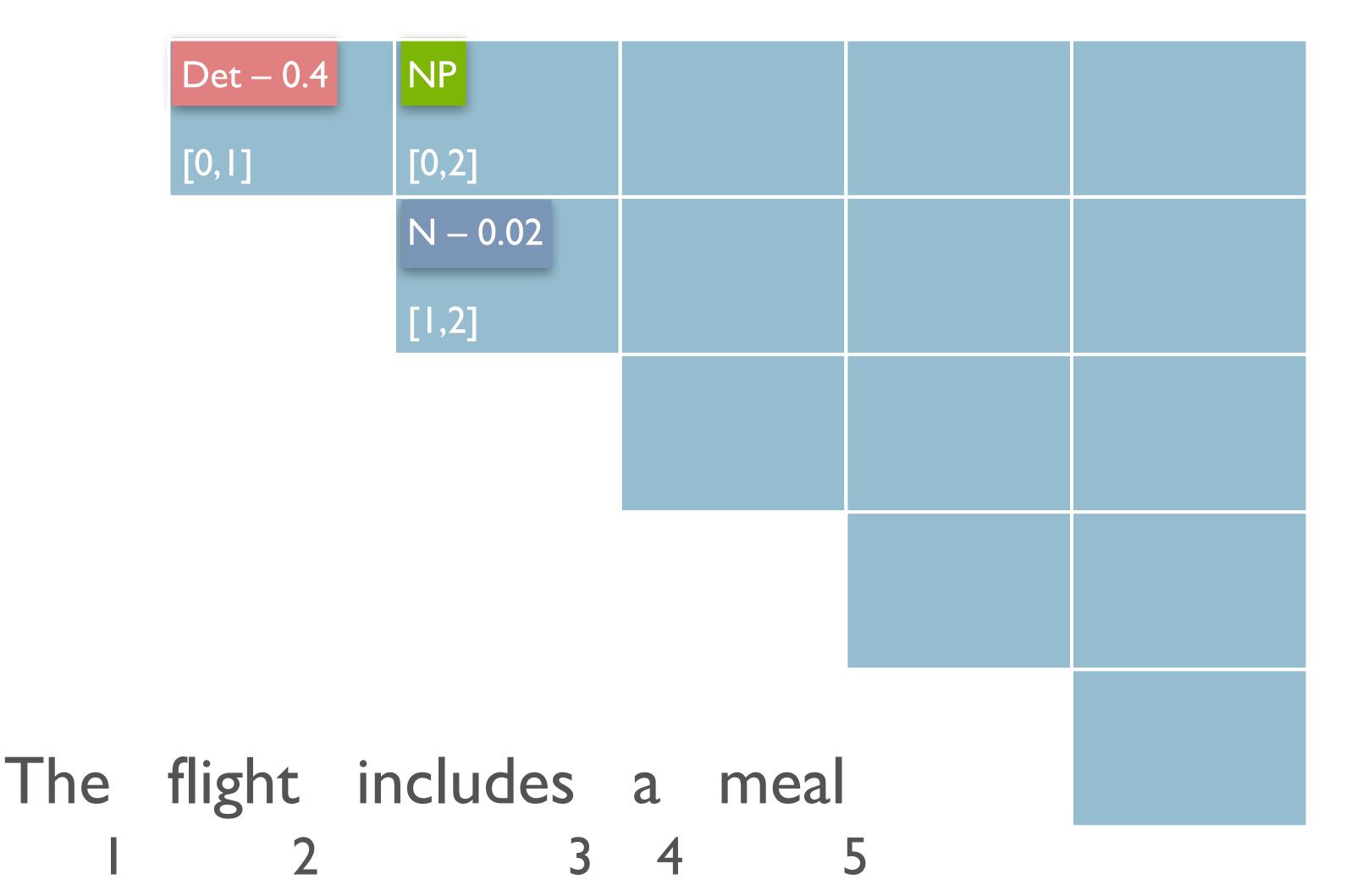


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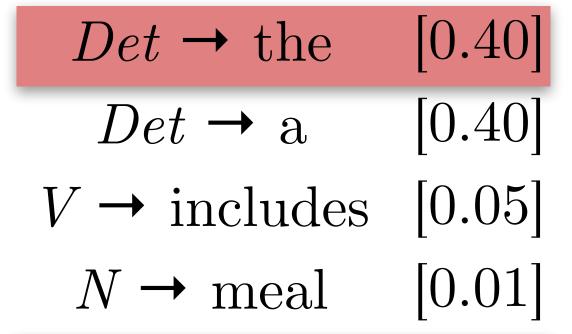
```
Det \rightarrow the \quad [0.40]
Det \rightarrow a \quad [0.40]
V \rightarrow includes \quad [0.05]
N \rightarrow meal \quad [0.01]
N \rightarrow flight \quad [0.02]
```

```
Det – 0.4
                 NP
       [0,1]
                 [0,2]
                 N - 0.02
                 [1,2]
The flight includes a meal
```

$$S \rightarrow NP \ VP \quad [0.80]$$

$$NP \rightarrow Det \ N \quad [0.30]$$

$$VP \rightarrow V \ NP \quad [0.20]$$



[0.02]

 $N \rightarrow \text{flight}$

```
Det – 0.4
                 NP
       [0,1]
                 [0,2]
                 N - 0.02
                 [1,2]
               = 0.00024
The flight includes a meal
```

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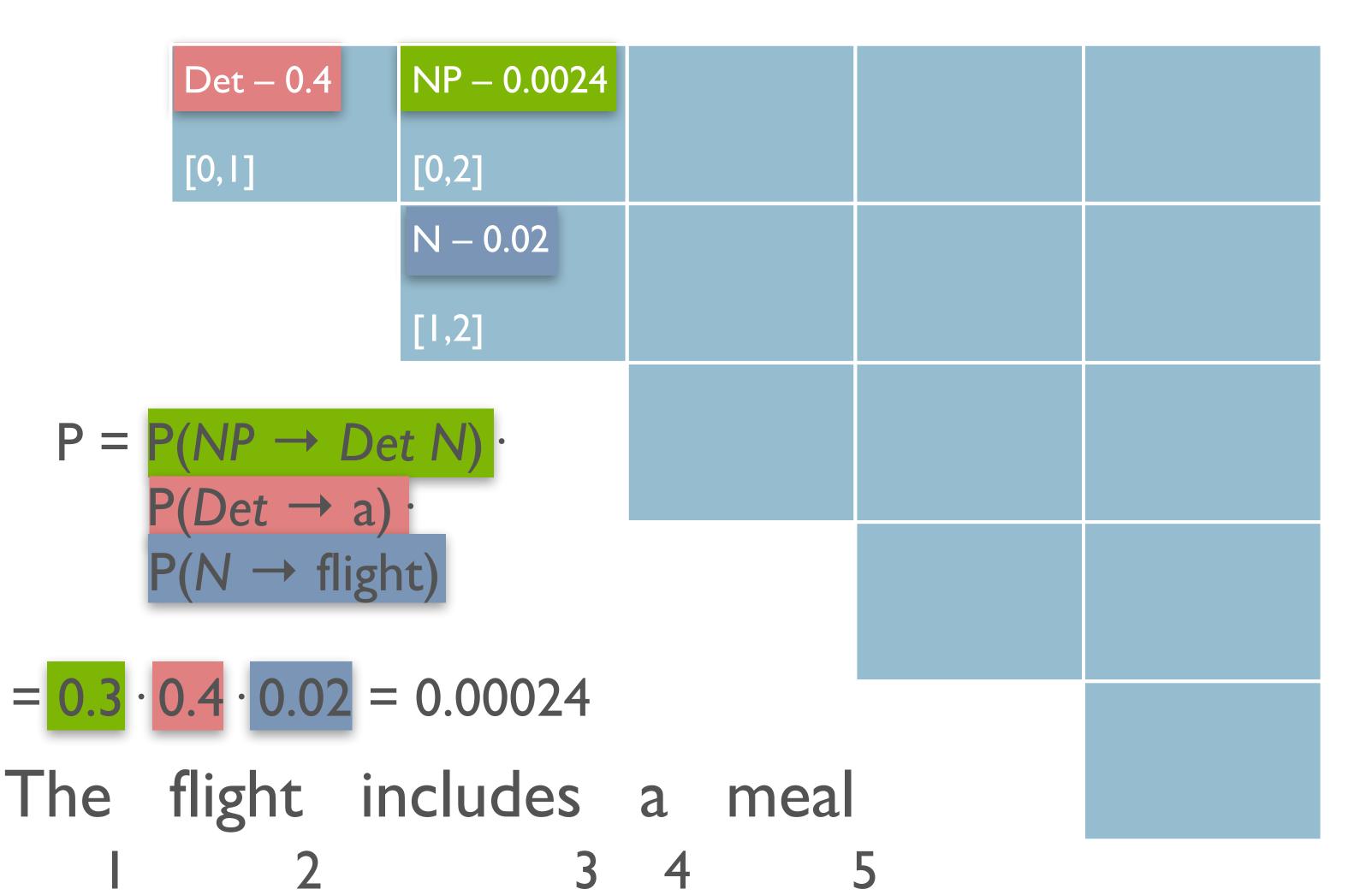
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$$N \rightarrow \text{meal}$$
 [0.01]

$$N \rightarrow \text{flight} \quad [0.02]$$



| $S \rightarrow NP VP$ | [0.80] |
|------------------------|--------|
| $NP \rightarrow Det N$ | [0.30] |
| $VP \rightarrow V NP$ | [0.20] |

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|---------------------------------|--------|
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| $N \rightarrow \text{flight}$ | [0.02] |

The

| Det - 0.4 | NP - 0.0024 | | | S - 2.304×10-8 |
|------------------------|-------------|----------|-----------|----------------|
| [0,1] | [0,2] | [0,3] | [0,4] | [0,5] |
| | N - 0.02 | | | |
| | [1,2] | [1,3] | [1,4] | [1,5] |
| | | V — 0.05 | | VP - 1.2×10-5 |
| | | [2,3] | [2,4] | [2,5] |
| | | | Det – 0.4 | NP - 0.0012 |
| | | | [3,4] | [3,5] |
| | | | N - 0.01 | |
| flight includes a meal | | | [4,5] | |

Inducing a PCFG

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 - Number of times a nonterminal is expanded by a given rule:

$$P(\alpha \to \beta \mid \alpha) = \frac{Count(\alpha \to \beta)}{\sum_{\gamma} Count(\alpha \to \gamma)} = \frac{Count(\alpha \to \beta)}{Count(\alpha)}$$

$$\Sigma_{\gamma} Count(\alpha \rightarrow \gamma)$$

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- Alternative: Learn probabilities by re-estimating
 - (Later)

Probabilistic Parser Development Paradigm

| | Train | Dev | Test |
|-------|---------------------------------|---|---------------------------------|
| | Large | Small | Small/Med |
| Size | (eg.WSJ 2-21, 39,830 sentences) | (e.g.WSJ 22) | (e.g. WSJ, 23, 2,416 sentences) |
| Usage | Estimate rule probabilities | Tuning/Verification, Check for Overfit | Held Out, Final Evaluation |

Assume a 'gold standard' set of parses for test set

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 - Partial credit:
 - Constituents in output match those in reference
 - Same start point, end point, non-terminal symbol

Parseval

- How can we compute parse score from constituents?
- Multiple Measures:

```
Labeled Recall (LR) = 
# of correct constituents in hypothetical parse
# of total constituents in reference parse

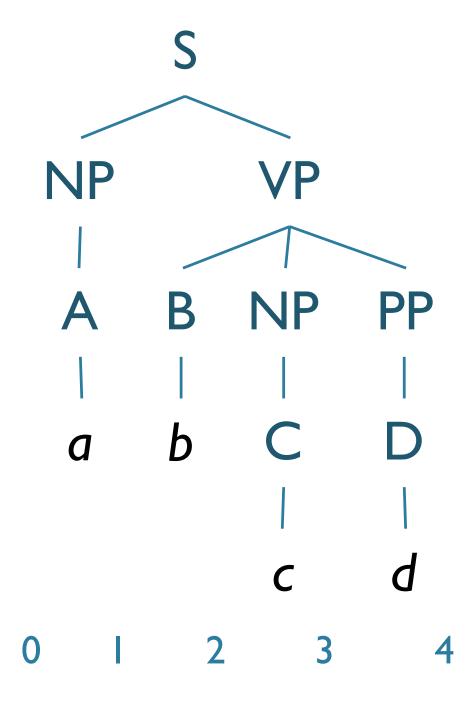
Labeled Precision (LP) = 
# of correct constituents in hypothetical parse
# of total consituents in hypothetical parse
```

Parseval

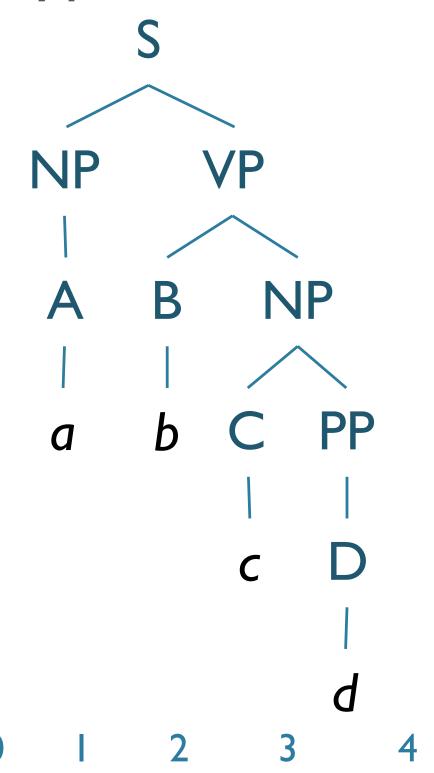
• F-measure:

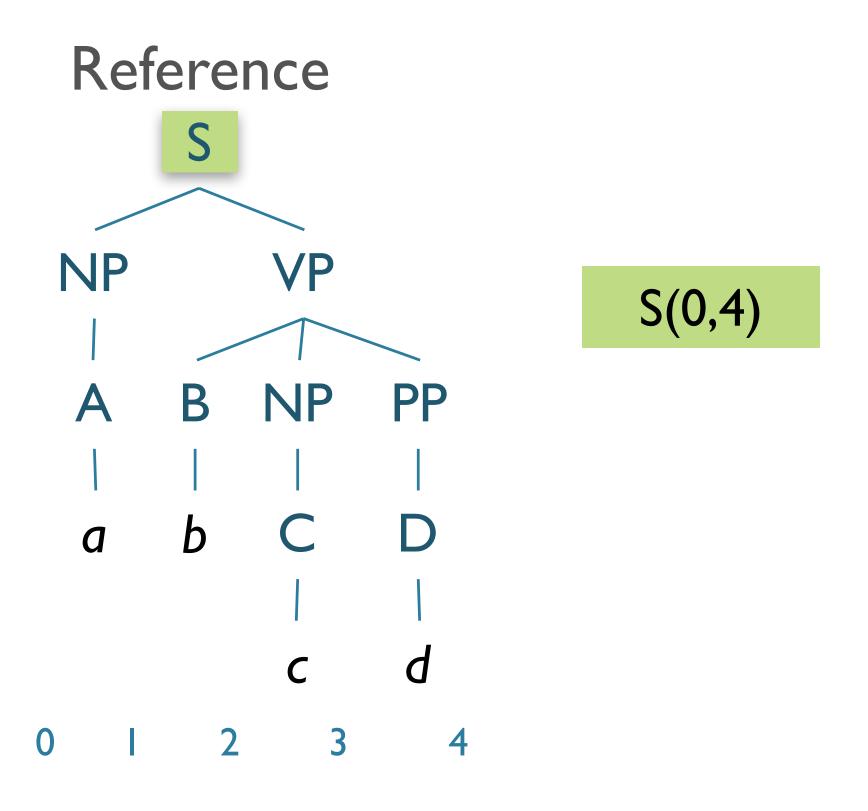
- Combines precision and recall
- Let $\beta \in \mathbb{R}$, $\beta > 0$ that adjusts P vs. R s.t. $\beta \propto \frac{R}{P}$
- F_{β} -measure is then: $F_{\beta} = (1 + \beta^2) \cdot \frac{P \cdot R}{\beta^2 \cdot P + R}$
- With F1-measure as $F_1 = \frac{2PR}{P+R}$

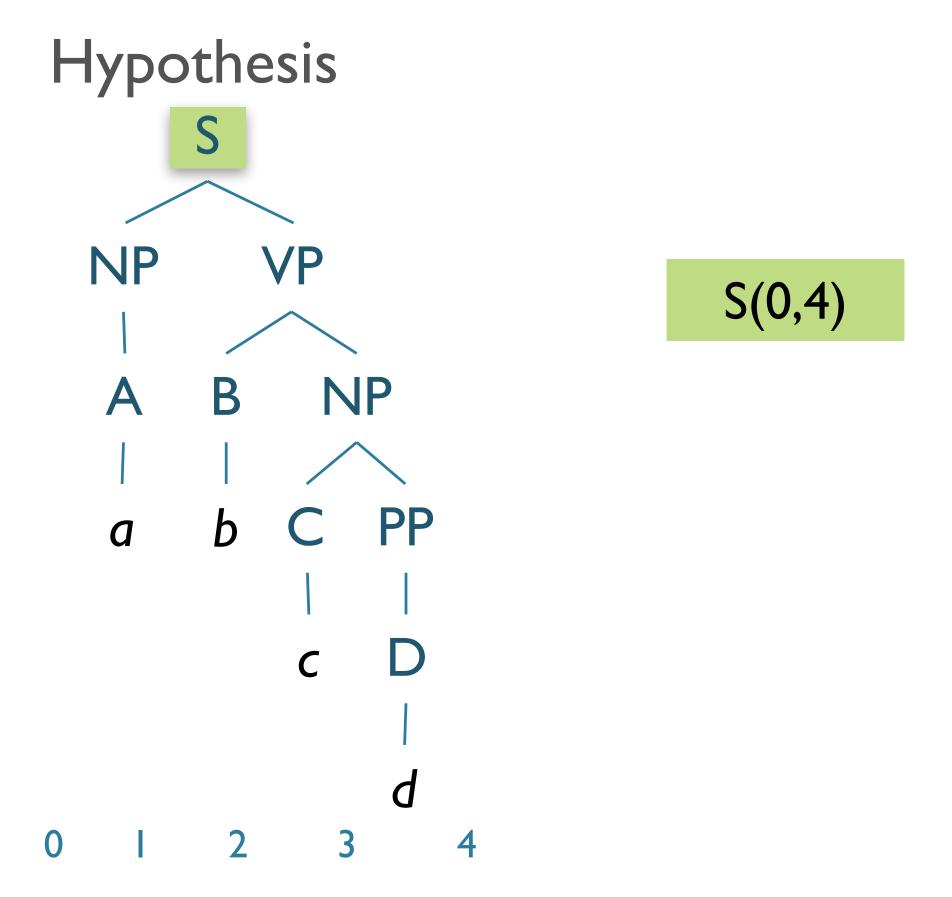
Reference

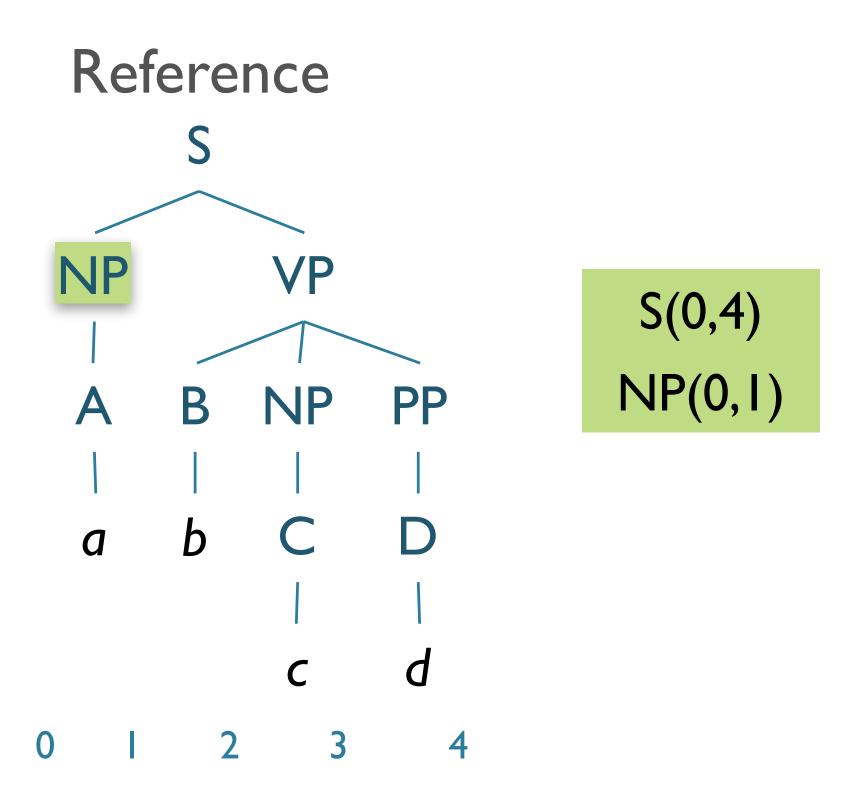


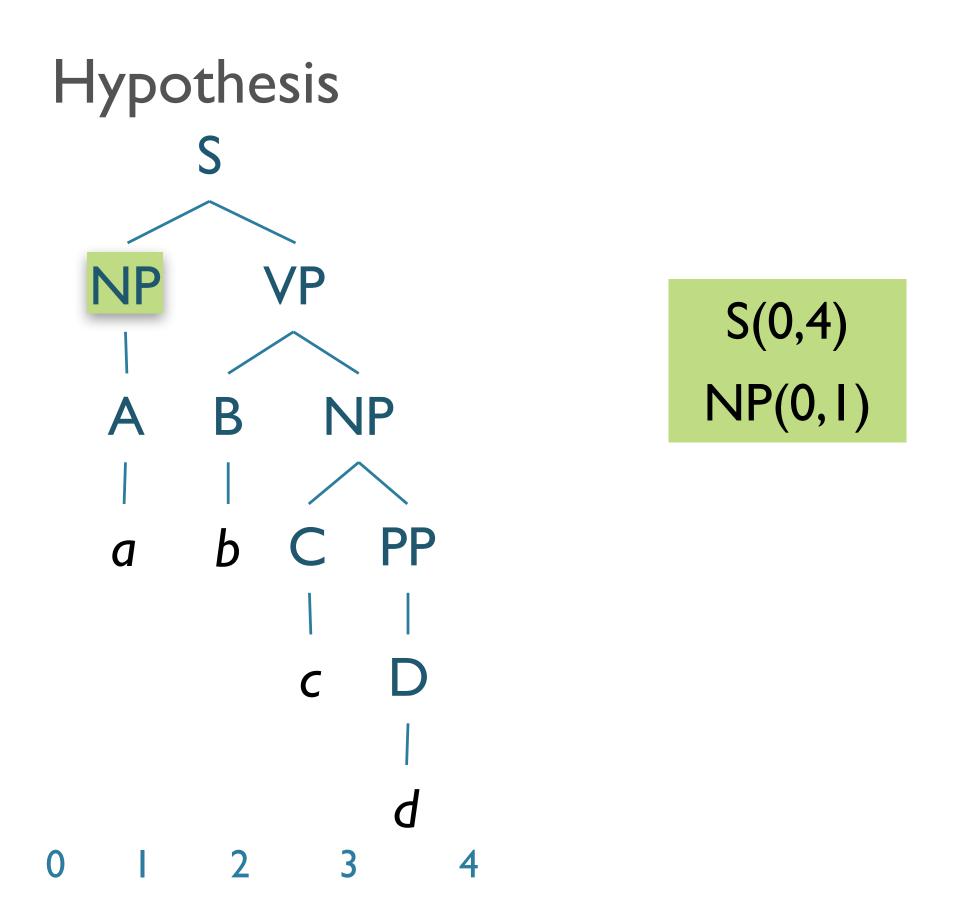
Hypothesis

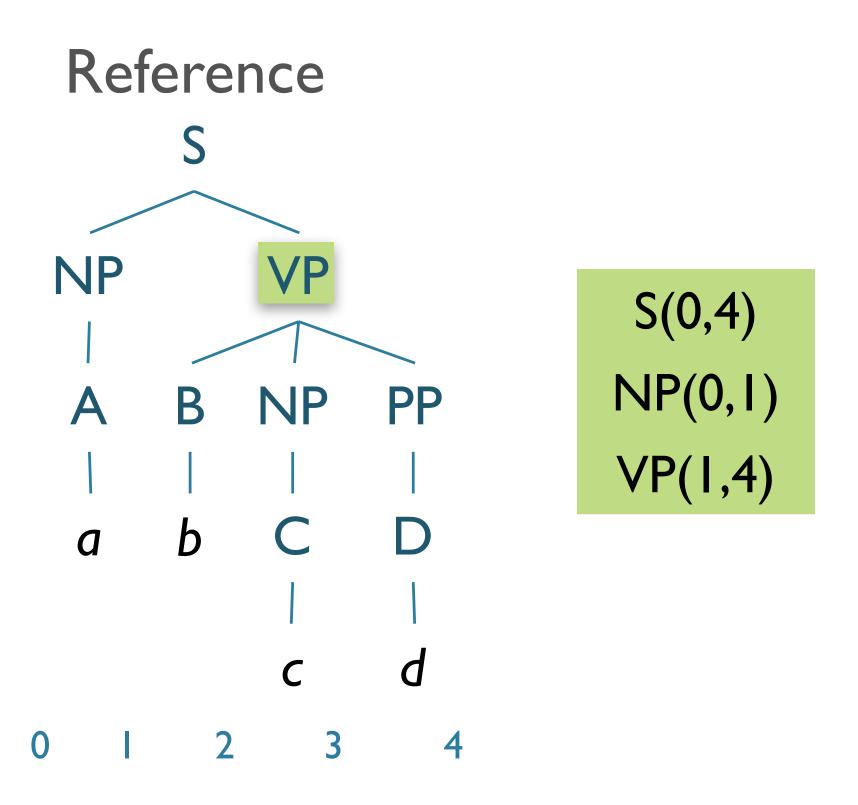


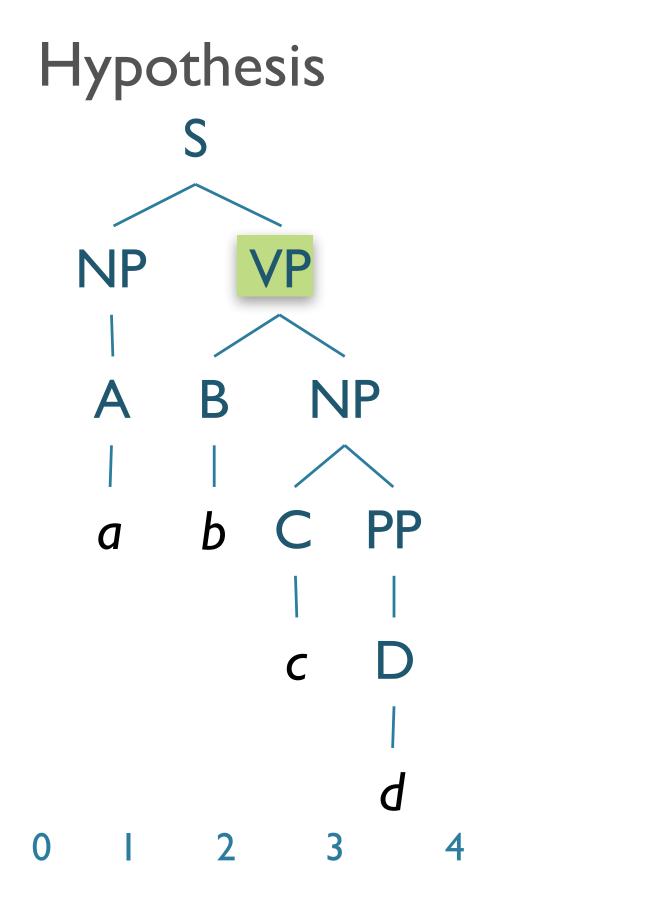


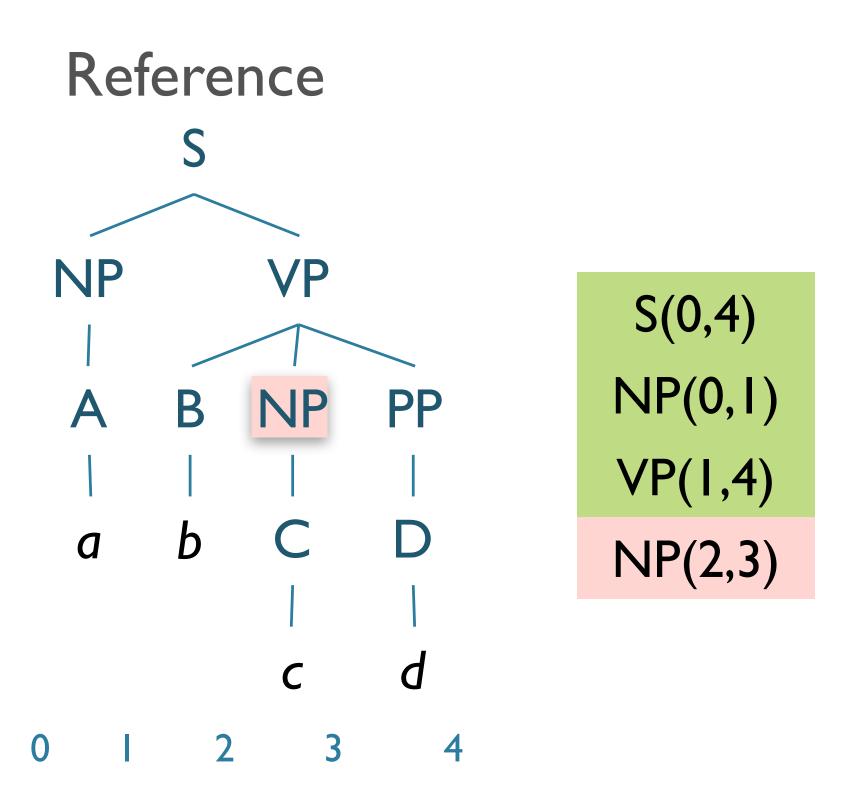


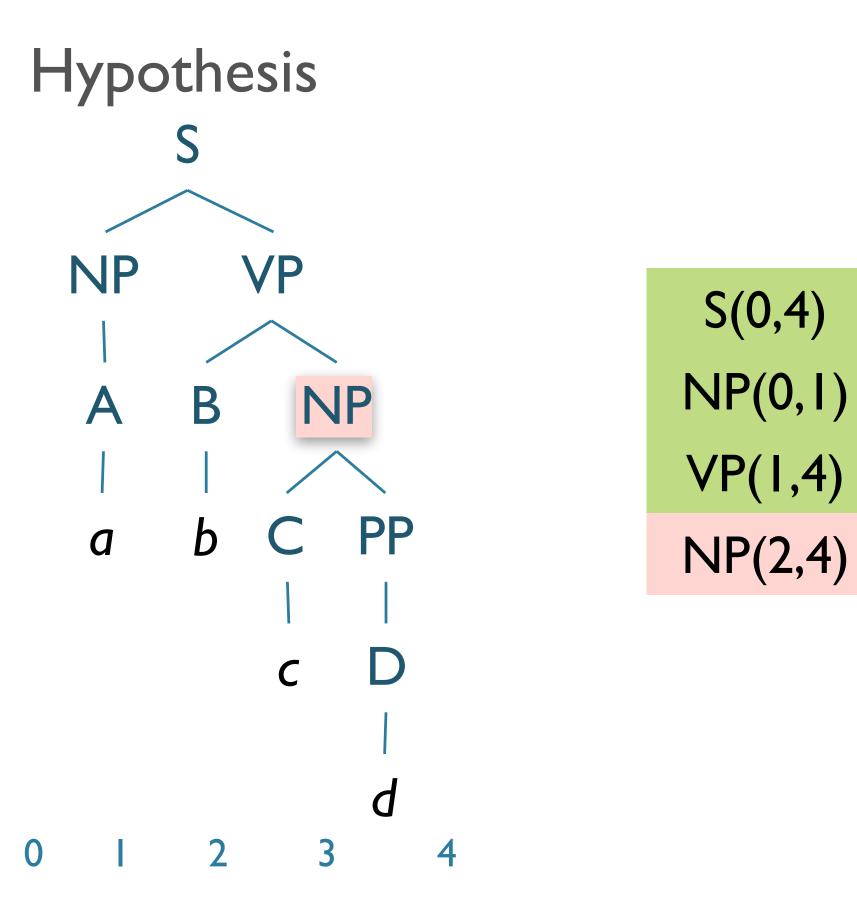


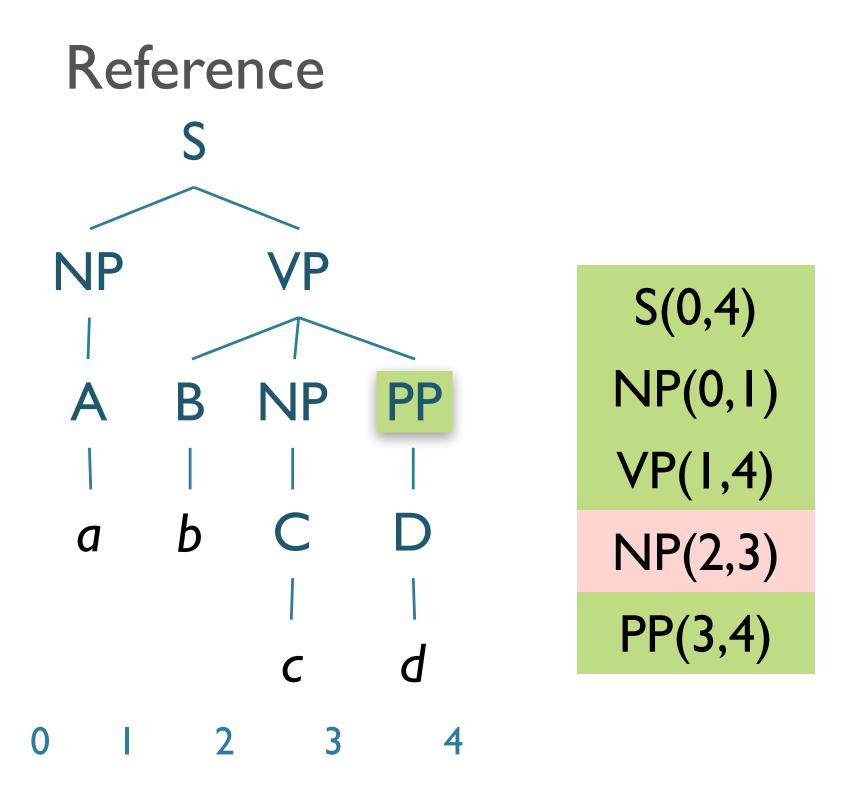


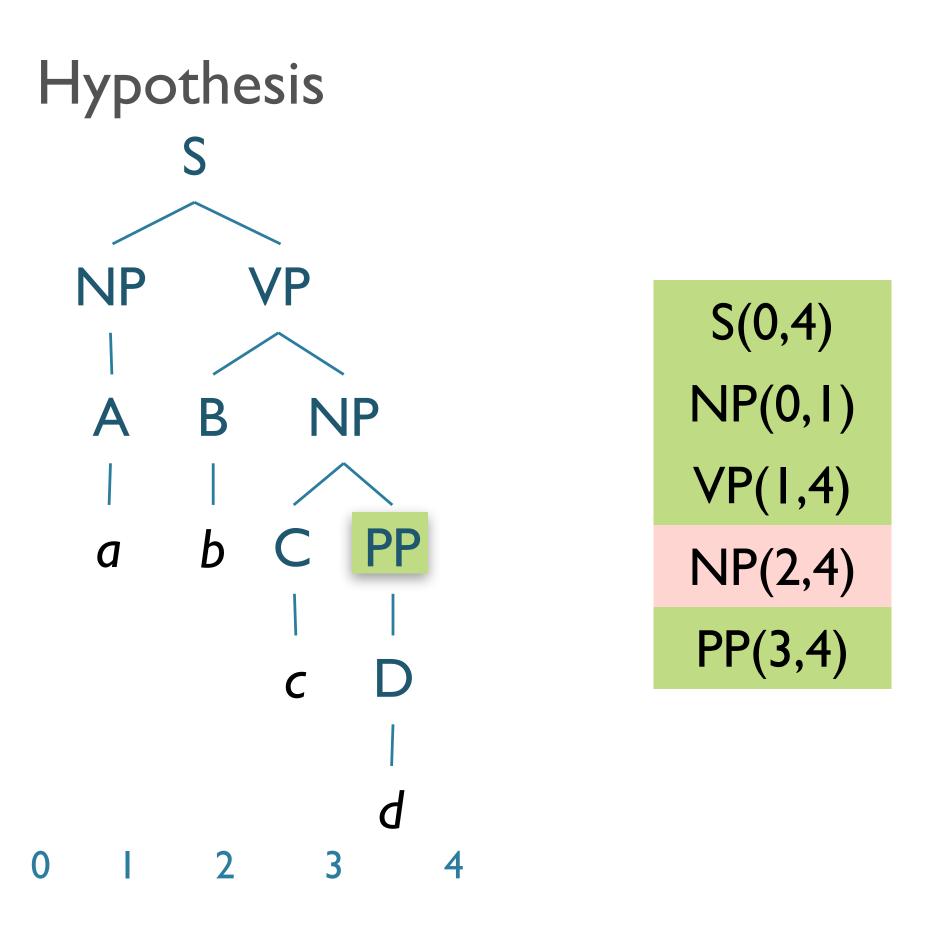




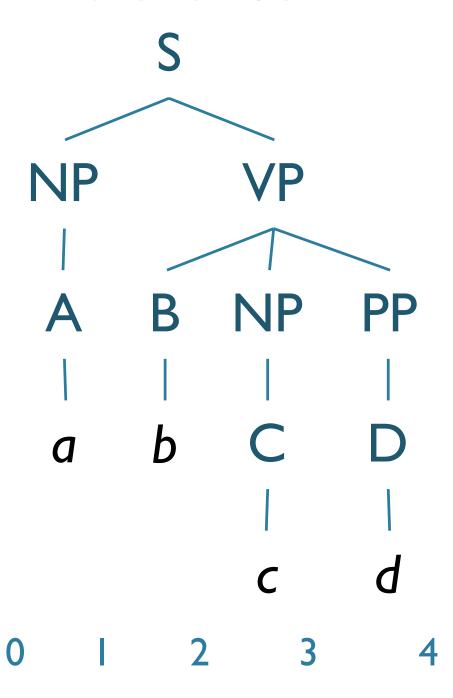








Reference



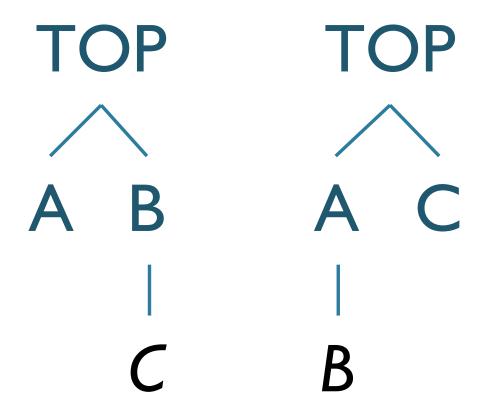
S(0,4) NP(0,1) VP(1,4) NP(2,3) PP(3,4)

> LP: 4/5 LR: 4/5

Hypothesis **VP** NP NP a

S(0,4) NP(0,1) VP(1,4) NP(2,4) PP(3,4)

- Crossing Brackets:
 - # of constituents where produced parse has bracketings that overlap for the siblings:
 - ((A B) C) { (0,2), (2,3) }
 and hyp. has
 (A (B C)) { (0,1), (1, 3) }



```
/* crossing is counted based on the brackets */
/* in test rather than gold file (by Mike) */
for(j=0;j<bn2;j++){
  for(i=0;i<bn1;i++){
    if(bracket1[i].result != 5 &&
        bracket2[j].result != 5 &&
        ((bracket1[i].start < bracket2[j].start &&
        bracket1[i].end > bracket2[j].start &&
        bracket1[i].end < bracket2[j].end) ||
        (bracket1[i].start > bracket2[j].start &&
        bracket1[i].start < bracket2[j].end &&
        bracket1[i].end > bracket2[j].end))){
```

from evalb.c

State-of-the-Art Parsing

- Parsers trained/tested on Wall Street Journal PTB
 - LR: 90%+;
 - LP: 90%+;
 - Crossing brackets: 1%

- Standard implementation of Parseval:
 - evalb

Evaluation Issues

- Only evaluating constituency
- There are other grammar formalisms:
 - LFG (Constraint-based)
 - Dependency Structure
- Extrinsic evaluation
 - How well does getting the correct parse match the semantics, etc?

Earley Parsing

Earley vs. CKY

- CKY doesn't capture full original structure
 - Can back-convert binarization, terminal conversion
 - Unit non-terminals require change in CKY

Earley vs. CKY

- CKY doesn't capture full original structure
 - Can back-convert binarization, terminal conversion
 - Unit non-terminals require change in CKY
- Earley algorithm
 - Supports parsing efficiently with arbitrary grammars
 - Top-down search
 - Dynamic programming
 - Tabulated partial solutions
 - Some bottom-up constraints

Earley Algorithm

- Another dynamic programming solution
 - Partial parses stored in "chart"
 - Compactly encodes ambiguity
 - O(N³)
- Chart entries contain:
 - Subtree for a single grammar rule
 - Progress in completing subtree
 - Position of subtree w.r.t. input

Earley Algorithm

- First, left-to-right pass fills out a chart with *N+1* states
 - Chart entries sit between words in the input string
 - Keep track of states of the parse at those positions
 - For each word position, chart contains set of states representing all partial parse trees generated so far
 - e.g. chart[0] contains all partial parse trees generated at the beginning of sentence

Chart Entries

- Three types of constituents:
 - Predicted constituents
 - In-progress constituents
 - Completed constituents

Parse Progress

- Represented by Dotted Rules
 - Position of indicates type of constituent
- 0 Book 1 that 2 flight 3
 - $S \rightarrow VP$ [0,0] (predicted)
 - $NP \rightarrow Det \cdot Nom$ [1,2] (in progress)
 - $VP \rightarrow VNP$ [0,3] (completed)
- [x,y] tells us what portion of the input is spanned so far by rule
- Each state *s_i*: <*dotted rule*>, [<*back pointer*>, <*current position*>]

o Book 1 that 2 flight 3

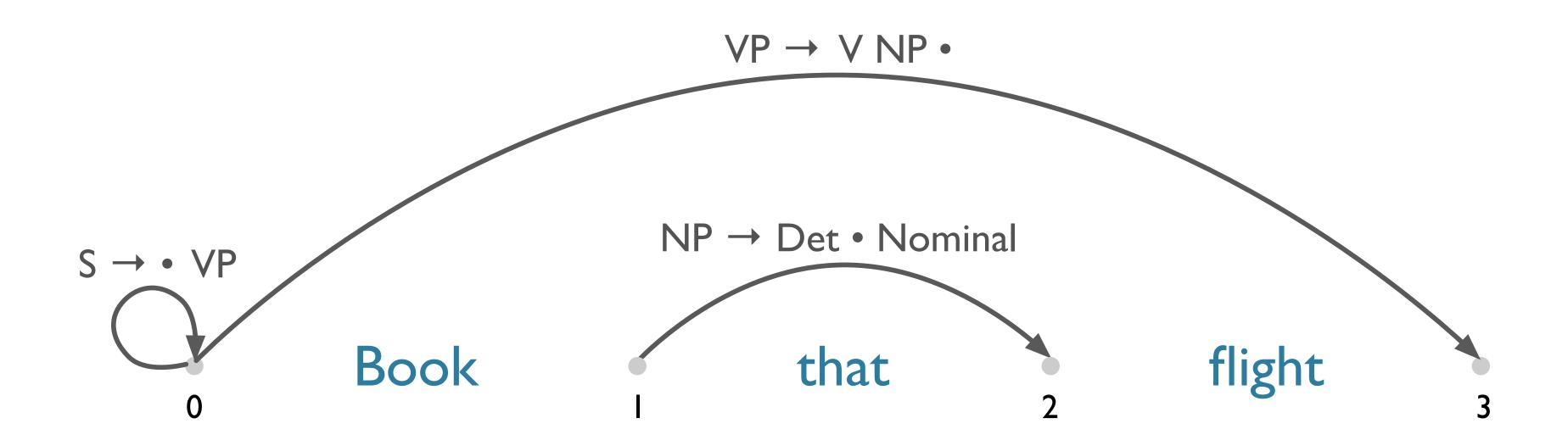
- $S \to VP, [0,0]$
 - First 0 means S constituent begins at the start of input
 - Second 0 means the dot is here too
 - So, this is a top-down prediction

o Book 1 that 2 flight 3

- $S \rightarrow VP$, [0,0]
 - First 0 means S constituent begins at the start of input
 - Second 0 means the dot is here too
 - So, this is a top-down prediction
- *NP* → *Det Nom*, [1,2]
 - the NP begins at position 1
 - the dot is at position 2
 - so, Det has been successfully parsed
 - Nom predicted next

o Book 1 that 2 flight 3 (continued)

- $V \rightarrow V NP \cdot [0,3]$
 - Successful VP parse of entire input



Successful Parse

- Final answer found by looking at last entry in chart
- If entry resembles $S \rightarrow \alpha \cdot [0,N]$ then input parsed successfully
- Chart will also contain record of all possible parses of input string, given the grammar

Parsing Procedure for the Earley Algorithm

- Move through each set of states in order, applying one of three operations:
 - predictor: add predictions to the chart
 - scanner: read input and add corresponding state to chart
 - completer: move dot to right when new constituent found
- Results (new states) added to current or next set of states in chart
- No backtracking and no states removed: keep complete history of parse

Earley Algorithm

```
function Earley-Parse(words, grammar) returns chart
 ENQUEUE((\gamma \rightarrow \bullet S, [0,0]), chart[\theta])
 for i \leftarrow \text{from } 0 \text{ to LENGTH}(words) \text{ do}
   for each state in chart[i] do
      if INCOMPLETE?(state) and
          Next-Cat(state) is not a part of speech then
        \mathbf{PREDICTOR}(state)
      elseif Incomplete: (state) and
          NEXT-CAT(state) is a part of speech then
        SCANNER(state)
      else
        Complete Energy (state)
      end
    end
 return(chart)
```

Earley Algorithm

```
procedure PREDICTOR((A \rightarrow \alpha \bullet B \beta, (i,j)))
  for each (B \rightarrow \gamma) in Grammar-Rules-For(B,grammar) do
    ENQUEUE((B \rightarrow \bullet \gamma, [j,j]), chart[j])
  end
procedure SCANNER((A \rightarrow a \bullet B \beta, [i,j]))
  if B c Parts-of-Speech(word/j/) then
    ENQUEUE((B \rightarrow word[j] \bullet, [j,j+1]), chart[j+1])
procedure Completer (B \rightarrow \gamma \bullet, [j,k])
 for each (A \rightarrow \alpha \bullet B \beta, [i,j]) in chart[j] do
    ENQUEUE((A \rightarrow a B \bullet \beta, [i,k]), chart[k])
  end
```

3 Main Subroutines of Earley

- Predictor
 - Adds predictions into the chart
- Scanner
 - Reads the input words and enters states representing those words into the chart
- Completer
 - Moves the dot to the right when new constituents are found

Predictor

- Intuition:
 - Create new state for top-down prediction of new phrase
- Applied when non part-of-speech non-terminals are to the right of a dot:
 - $S \rightarrow VP[0,0]$
- Adds new states to current chart
 - One new state for each expansion of the non-terminal in the grammar

$$VP \rightarrow \cdot V$$
 [0,0]
 $VP \rightarrow \cdot V NP$ [0,0]

Chart[0]

| S0 | $\gamma \rightarrow \cdot S$ | [0,0] | Dummy start state |
|------------------------------|--|----------------------------------|---|
| S1 | $S \rightarrow \cdot NP VP$ | [0.0] | Predictor |
| S2 | $S \rightarrow \cdot Aux NP VP$ | [0,0] | Predictor |
| S3 | $S \rightarrow \cdot VP$ | [0,0] | Predictor |
| S4 | NP → • Pronoun NP → • Proper-Noun NP → • Det Nominal | [0,0] | Predictor |
| S5 | | [0,0] | Predictor |
| S6 | | [0,0] | Predictor |
| S7 S8 S9 S10 S11 | $VP \rightarrow \cdot Verb$ $VP \rightarrow \cdot Verb NP$ $VP \rightarrow \cdot Verb NP PP$ $VP \rightarrow \cdot Verb PP$ $VP \rightarrow \cdot VP PP$ | [0,0] [0,0] [0,0] [0,0] | Predictor Predictor Predictor Predictor Predictor |

Chart[1]

| S12 | Verb → book • | [0,1] | Scanner |
|--------------------------|---|-------------------------|--|
| S13 S14 S15 S16 | $VP \rightarrow Verb \cdot NP$ $VP \rightarrow Verb \cdot NP$ $VP \rightarrow Verb \cdot NP PP$ $VP \rightarrow Verb \cdot PP$ | [0,1] [0,1] [0,1] | Completer Completer Completer Completer |
| S17 | $S \rightarrow VP$. | [0,1] | Completer |
| S18 | $VP \rightarrow VP \cdot PP$ | [0,1] | Completer |
| S19 S20 S21 S22 | $NP \rightarrow \cdot Pronoun$ $NP \rightarrow \cdot Proper-Noun$ $NP \rightarrow \cdot Det Nominal$ $PP \rightarrow \cdot Prep NP$ | [1,1] [1,1] [1,1] | Predictor Predictor Predictor Predictor |

Book that flight

S0: $\gamma \to S[0,0]$



Book that flight

S0:
$$\gamma \rightarrow \cdot S[0,0]$$

S3:
$$S \rightarrow VP[0,0]$$

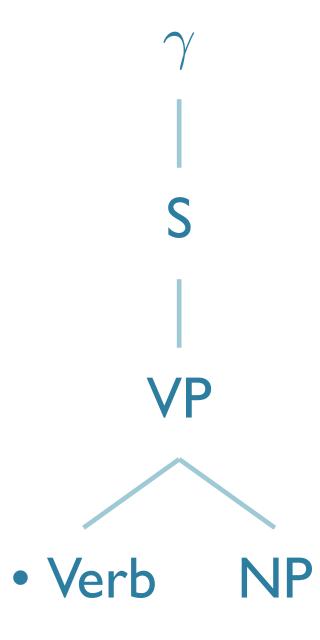


Book that flight

S0: $\gamma \to S[0,0]$

S3: $S \rightarrow VP[0,0]$

S8: *VP* → • *Verb NP* [0,0]

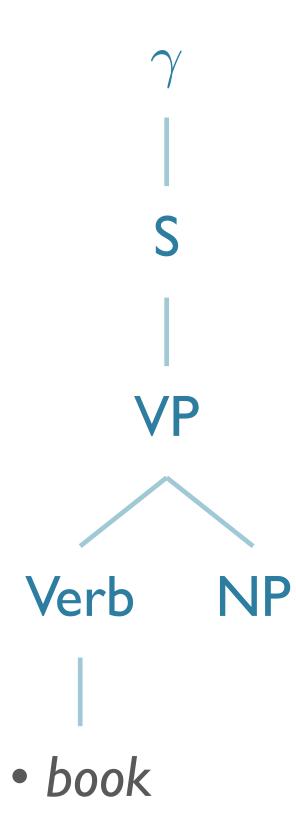


```
S0: \gamma \to S[0,0]
```

S3: $S \rightarrow VP[0,0]$

S8: *VP* → • *Verb NP* [0,0]

S12: *Verb* → • *book* [0,0]

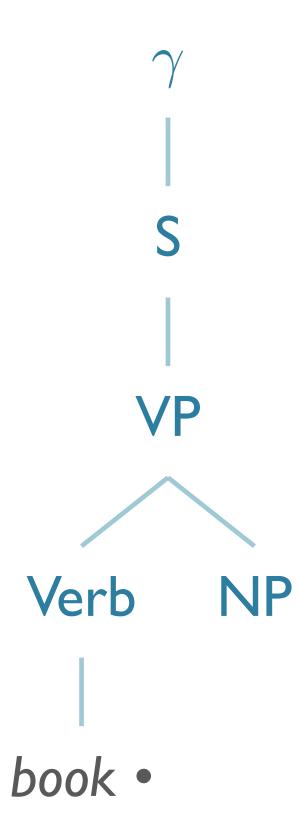


```
S0: \gamma \to S[0,0]
```

S3: $S \rightarrow VP[0,0]$

S8: *VP* → • *Verb NP* [0,0]

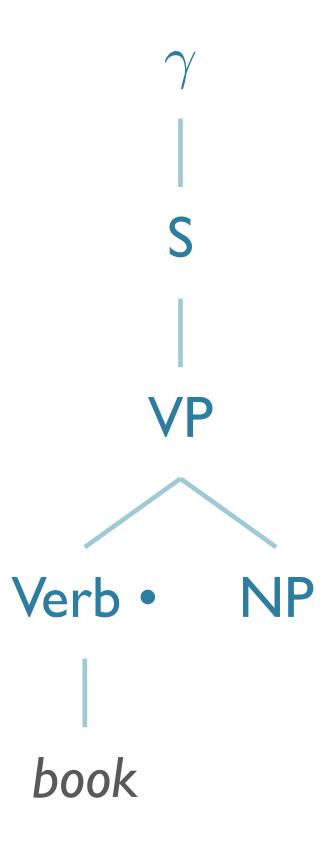
S12: *Verb* → *book* • [0,1]



```
S0: \gamma \to S[0,0]
```

S3: $S \rightarrow VP[0,0]$

S8: *VP* → *Verb* • *NP* [0,1]



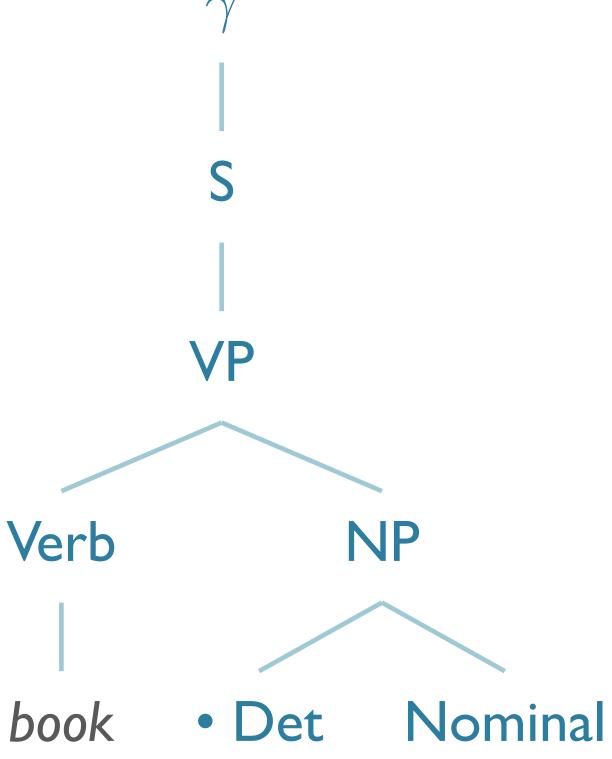
```
S0: \gamma \to S[0,0]
```

S3: $S \rightarrow VP \cdot [0,1]$

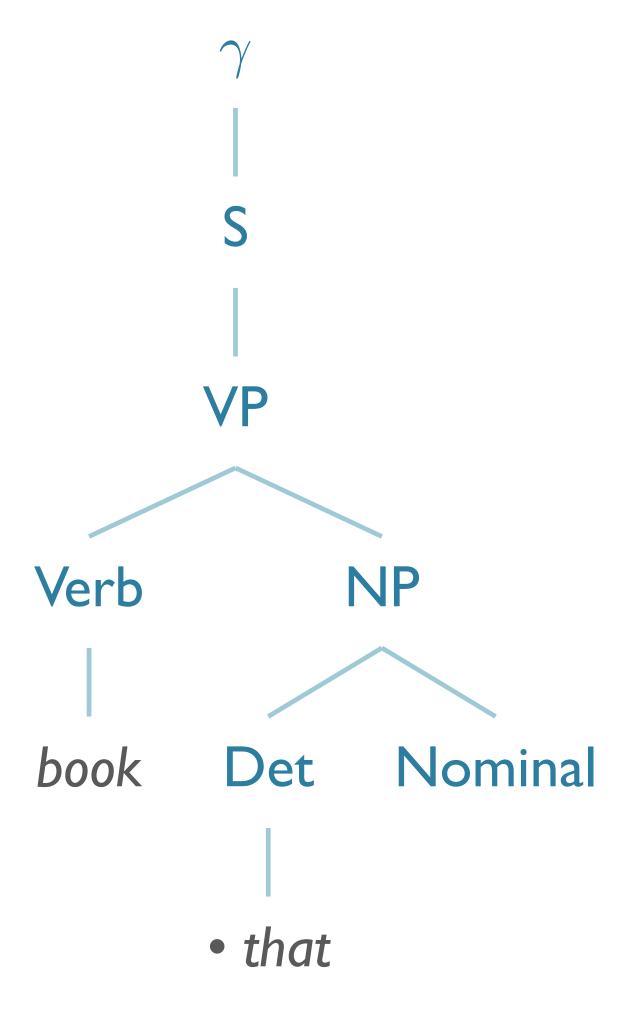
S8: *VP* → *Verb* • *NP* [0,1]



```
S0: \gamma \rightarrow \cdot S[0,0]
S3: S \rightarrow VP \cdot [0,1]
S8: VP \rightarrow Verb \cdot NP[0,1]
S21: NP \rightarrow \cdot Det Nominal[1,1]
```

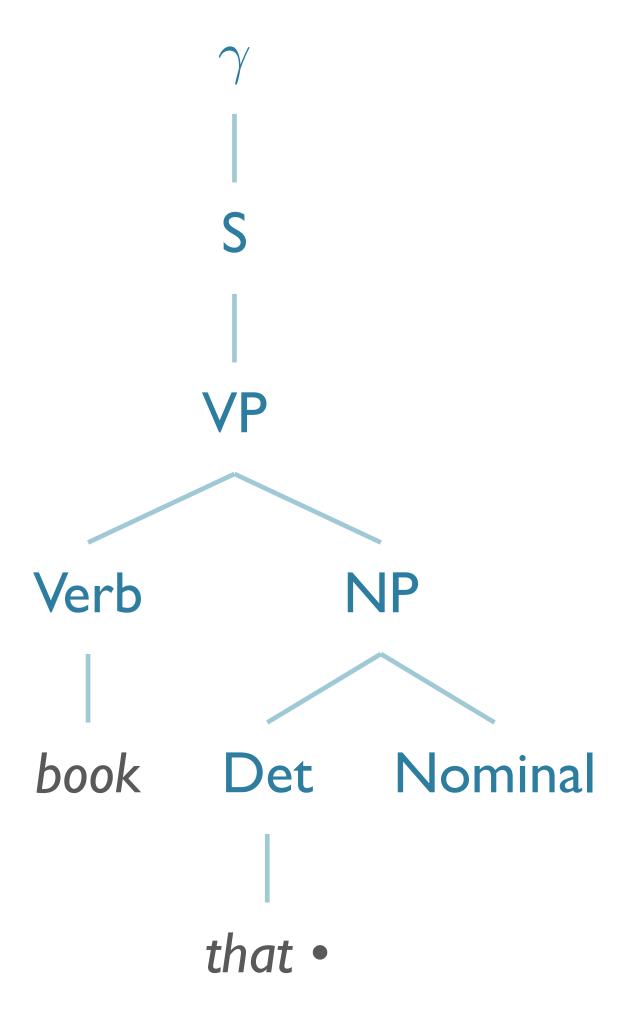


```
S0: \gamma \to \cdot S[0,0]
S3: S \to VP \cdot [0,1]
S8: VP \to Verb \cdot NP[0,1]
S21: NP \to \cdot Det Nominal[1,1]
S23: Det \to \cdot \text{"that"}[1,1]
```



```
S0: \gamma \rightarrow \cdot S[0,0]
S3: S \rightarrow VP \cdot [0,1]
S8: VP \rightarrow Verb \cdot NP[0,1]
S21: NP \rightarrow \cdot Det Nominal[1,1]
```

S23: *Det* → "that" • [1,2]



```
S0: \gamma \to S[0,0]
S3: S \rightarrow VP \cdot [0,1]
S8: VP → Verb • NP [0,1]
S21: NP → Det • Nominal [1,2]
                                              VP
                                      Verb
                                                    NP
                                                      Nominal
                                       book
                                              Det •
                                               that
```

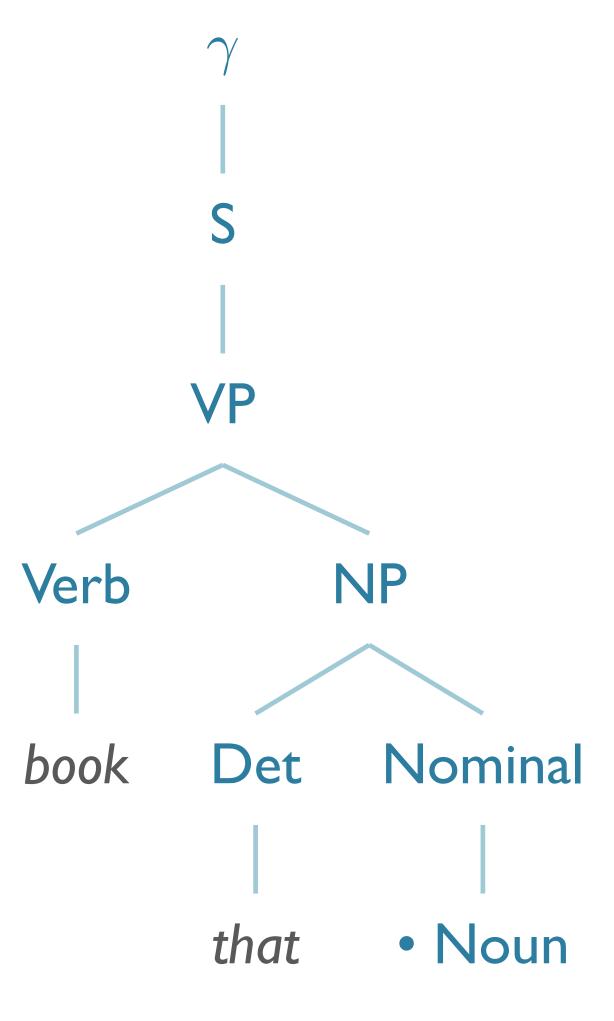
```
S0: \gamma \to S[0,0]
```

S3: $S \rightarrow VP \cdot [0,1]$

S8: *VP* → *Verb* • *NP* [0,1]

S21: *NP* → *Det* • *Nominal* [1,2]

S25: *Nominal* → • *Noun* [2,2]



```
S0: \gamma \to S[0,0]
```

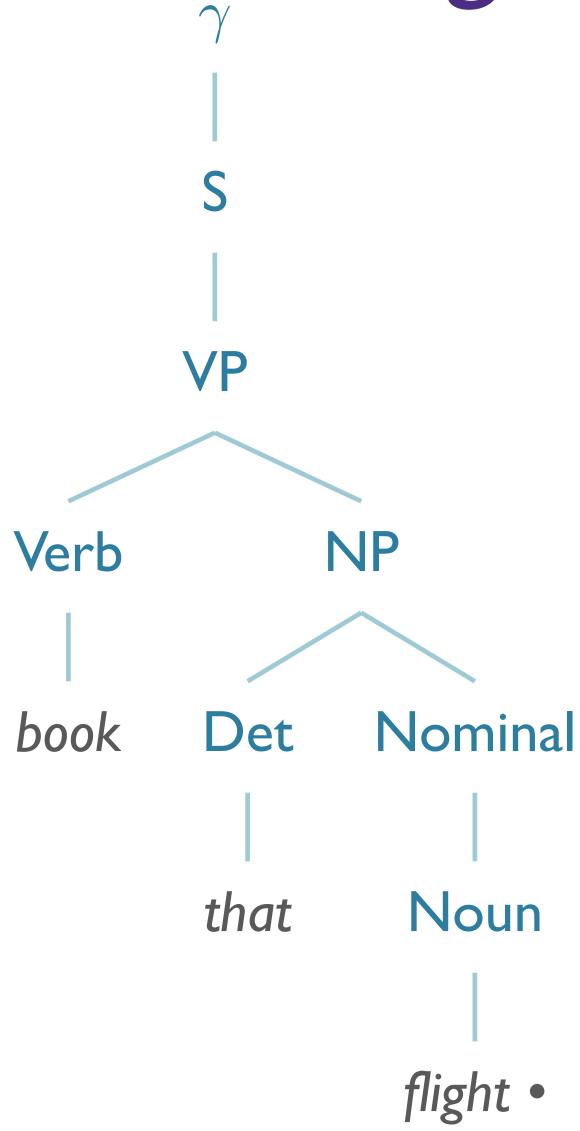
S3:
$$S \rightarrow VP \cdot [0,1]$$

S8: *VP* → *Verb* • *NP* [0,1]

S21: *NP* → *Det* • *Nominal* [1,2]

S25: *Nominal* → • *Noun* [2,2]

S28: *Noun* → "*flight*" • [2,3]



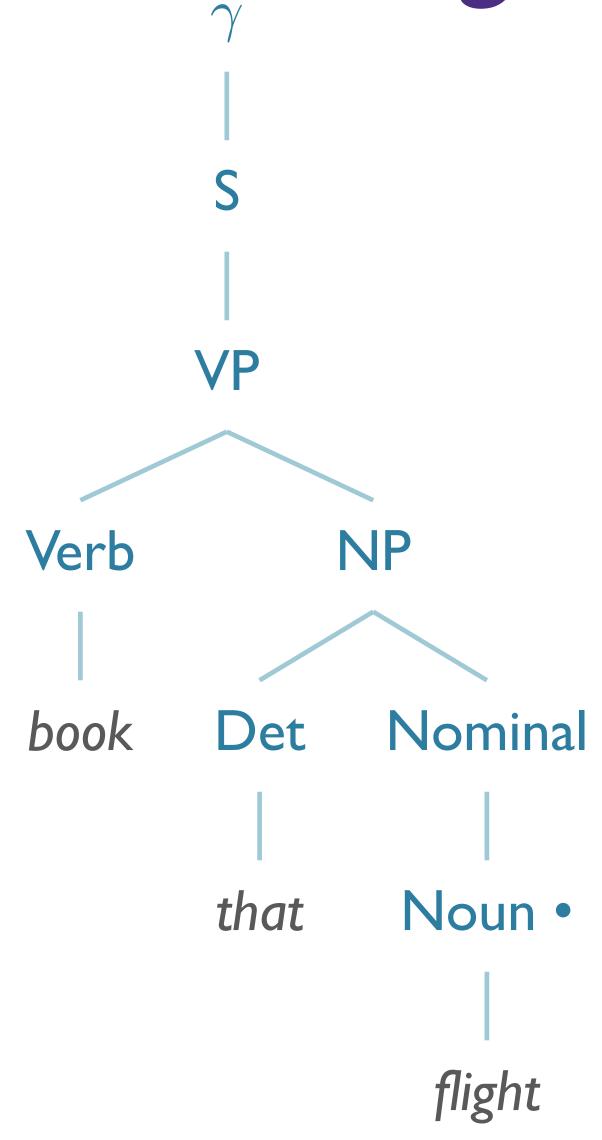
```
S0: \gamma \rightarrow \cdot S[0,0]
```

S3: $S \rightarrow VP \cdot [0,1]$

S8: *VP* → *Verb* • *NP* [0,1]

S21: *NP* → *Det* • *Nominal* [1,2]

S25: *Nominal* → *Noun* • [2,3]

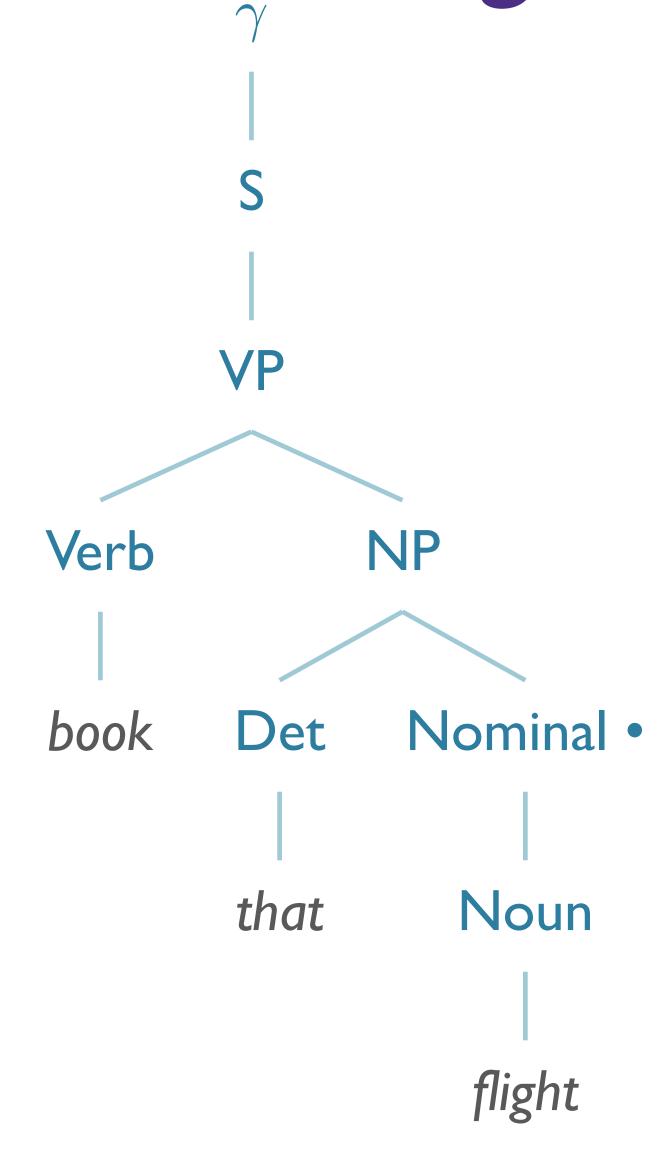


S0: $\gamma \rightarrow \cdot S[0,0]$

S3: $S \rightarrow VP \cdot [0,1]$

S8: *VP* → *Verb* • *NP* [0,1]

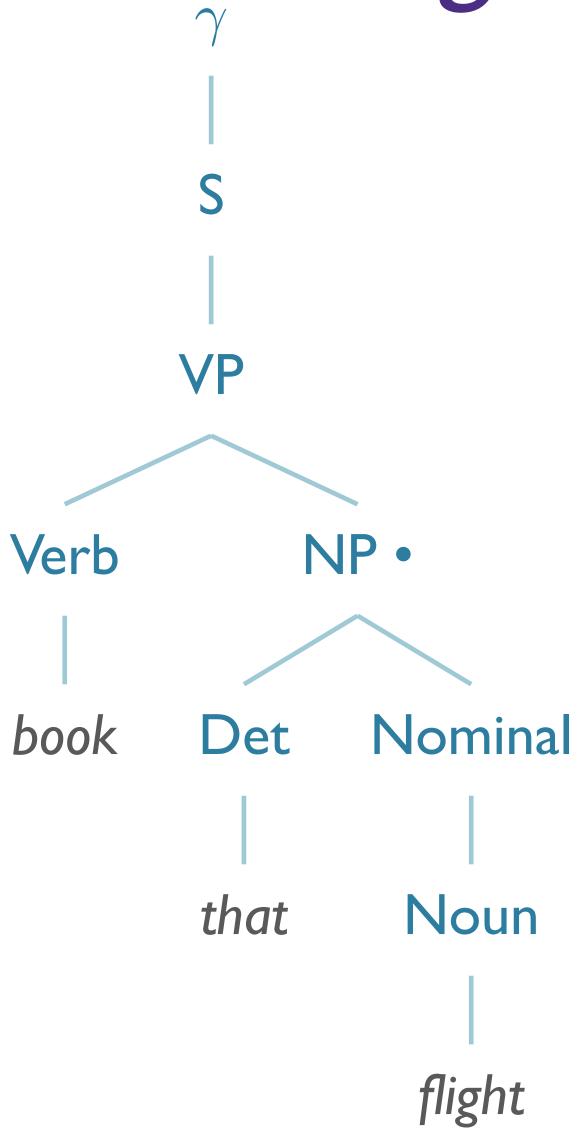
S21: *NP* → *Det Nominal* • [1,3]



S0: $\gamma \to S[0,0]$

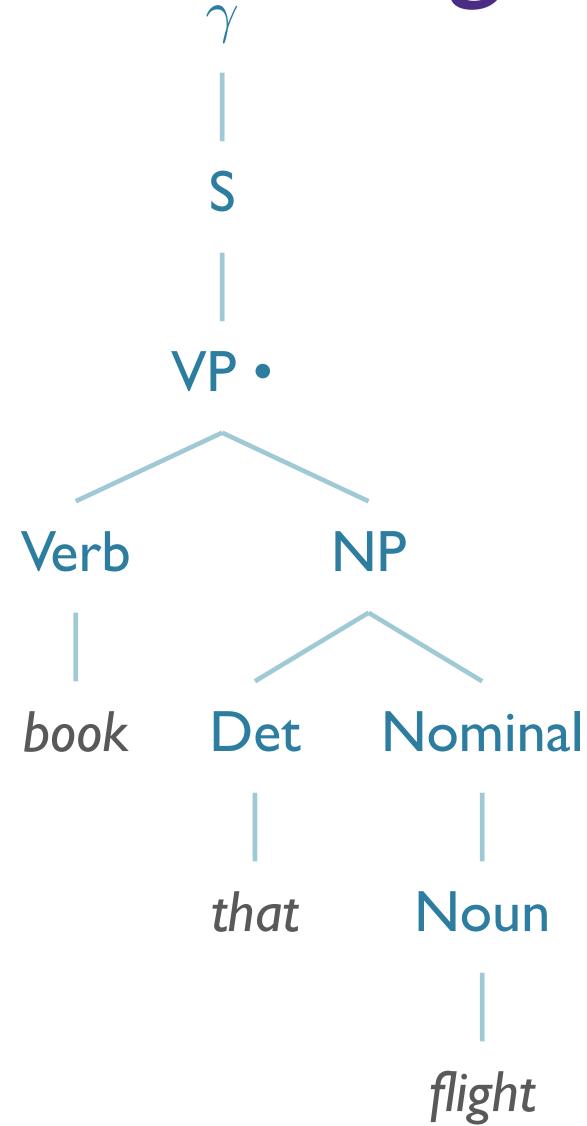
S3: $S \rightarrow VP \cdot [0,1]$

S8: *VP* → *Verb NP* • [0,3]



S0: $\gamma \to S[0,0]$

S3: $S \rightarrow VP \cdot [0,3]$



What About Dead Ends?

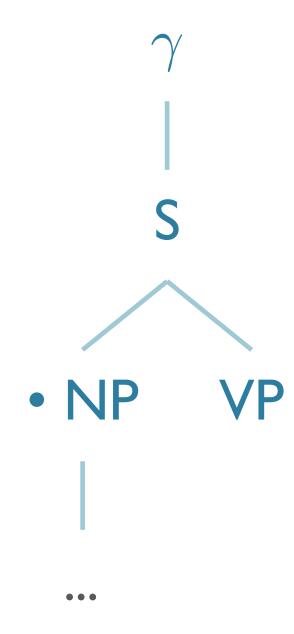
S0: $\gamma \to S[0,0]$

S1: $S \rightarrow \cdot NP VP$ [0,0]

NP → • Pronoun

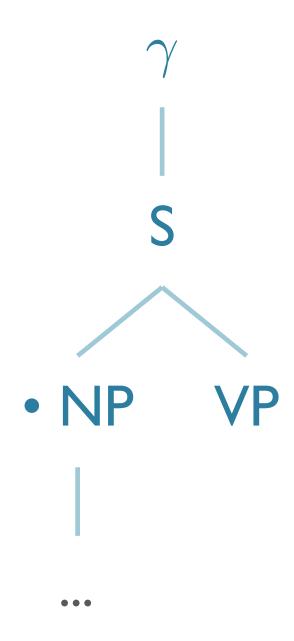
NP → • Proper-Noun

NP → • Det Nominal



book

```
S0: \gamma \to S[0,0]
S1: S \to NP VP[0,0]
```



book

 We now have a top-down parser in hand. Does it enter infinite loops on rules like S -> S 'and' S?

- We now have a top-down parser in hand. Does it enter infinite loops on rules like S -> S 'and' S?
- No!

```
procedure ENQUEUE(state, chart-entry)
  if state is not already in chart-entry then
    Push(state, chart-entry)
  end
```

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- No!

```
procedure ENQUEUE(state, chart-entry)
if state is not already in chart-entry then
    Push(state, chart-entry)
end
```

Exercise: parse 'table and chair' using the very simple grammar Nom -> Nom 'and' Nom | 'table' | 'chair'