

# Baby Tetris First Part Report

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<https://github.com/diarca/baby-tetris/>

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## 1 Question 1: Discounted MDP For Player

The objective of this part is to choose the optimal policy that maximizes the expected cumulative discounted reward over time for the "player".

### States

The state is represented by the current configuration of the grid, and the next piece to be placed.

$S := (\text{current grid}, \text{next piece to place})$

### Actions

The actions correspond to a position and orientation to place the piece on the grid.

$A := (\text{position}, \text{orientation})$

### Transition Function

The transition function is defined as  $P(s'|s, a) := 1/2$

There are two possible pieces (I or L) as the next piece to place, and considering our state representation, and our assumption of uniform distribution between the pieces, each transition has an equal probability of 1/2.

### Reward Function and Discount Factor

The reward function is computed based on the number of lines cleared after placing the piece, and a discount factor  $\lambda$  is used to weigh future rewards.

$R(s, a, s') := \text{coefficient} * \text{number of lines cleared with } \lambda \in [0, 1]$

## 1.1 Finding Optimal Policy

We use the value iteration method to find the optimal policy that maximizes the expected rewards, iteratively updating the value function for each state based on the Bellman equation until it converges to the optimal value (margin  $\epsilon$ ).

$$V(s) = \max_{a \in A} \sum_{s' \in S} P(s'|s, a) \cdot (R(s, a, s') + \lambda V(s'))$$

Value iteration is better suited for maximizing the expected discounted reward since it computes the optimal value then extracts the policy and it's easier to implement.

## 1.2 Implementation

### 1.2.1 Main Architecture

Based on the our design choices from Section 1, we decided to use an object oriented architecture (C++) for our implementation.

For our structure we have declared 7 classes:

1. *Point*: classical point (x,y) class used for grid coordinates
2. *Field*: the class that contain a boolean grid s.t.  $grid[i][j] = 1$  if the corresponding cell is filled in the field
3. *Tromino*: corresponds to the baby Tetris pieces (I Piece or L Piece)
4. *State*: contains both a field grid and an incoming Tromino
5. *Action*: represents a choosable action on a given state (combination for placement and rotation of an incoming Tromino)
6. *Game*: contains the configuration (completed lines scores), the current state and score of a baby Tetris game (the configuration is loaded from a file : *config.txt*)
7. *MDP*: encompasses all the necessary methods in order to compute the value iteration over a game

Now that we presented the overview of the global implementation architecture, we will describe our implementation choices.

### 1.2.2 Implementation Choices

- **Pieces teleported to their final position:** In order to reduce the number of intermediate states when choosing an action, we compute on a state all available actions, then according to the policy, we make a choice over these actions and we teleport the incoming piece in the current state directly to its final position (with the corresponding rotation)
- **Reward Function:** As defined in Section 1, the reward function takes both the current state  $s$  and an available action  $a_s$ . Since we can calculate the available actions, we know in which state  $s'$  we're going to arrive and we can compute directly in  $s'$  the number of completed lines before removing them from the game. Then we abstract the reward function by counting the number of complete lines in the resulting state

- **Game Ending:** According to the subject, the game should end whenever a piece exceeds the height limit of the grid. Since we compute all available actions on a given state, if a piece cannot be placed on the grid, there will be no available action for that state, the game will terminate when there is no available action for the current state.
- **Max Iteration:** Since the value iteration isn't an exact method, if we fix an  $\epsilon$  too small, we might iterate for a very long time. For this reason, we fix a maximum number of iteration for the expected value improvement to have a value iteration method that computes in reasonable time, even if the expected value vector is less qualitative than the  $\epsilon$  margin we fixed

### 1.3 Results

Here, we will present multiple examples of execution of our implementation. In these executions we will use fixed values for:  $\text{width} = 4$ ,  $\epsilon = 0.00000001$ ,  $\text{maxIteration} = 100$ ,  $\text{probAIPiece} = 0.5$  and  $\text{maxGameAction} = 10000$ . Note that we fixed a small  $\epsilon$  in order to approach the optimal policy, keeping in mind that the algorithm is bounded in iteration by the *maxIteration*. And the variant variables will be: the height  $\in \{4, 5\}$  and the discount factor  $\lambda \in [0, 1]$ .

For each execution we will measure:

1. the number of iterations for the value iteration algorithm
2. the average expected gain of the optimal policy
3. the final score either at the end of the game or when we reach the maximum number of actions
4. the execution time of the said execution

### 1.3.1 Executions

```
All constants:
width = 4, height = 4, probaIPiece
= 0.5, maxGameAction = 10000
epsilon = 1e-08, maxIteration =
100, lambda = 0.1

i = 0 and delta = 6
i = 1 and delta = 0.6
i = 2 and delta = 0.06
i = 3 and delta = 0.00525
i = 4 and delta = 0.00045
i = 5 and delta = 0

average over final V 0.412926
...
...
**..
**..

Game Over! Global score: 7849 in
10000 actions
make run 3.02s user 0.08s system
92% cpu 3.340 total
```

height: 4 and  $\lambda$ : 0.1

```
All constants:
width = 4, height = 4, probaIPiece
= 0.5, maxGameAction = 10000
epsilon = 1e-08, maxIteration =
100, lambda = 0.4

i = 0 and delta = 6
i = 1 and delta = 2.4
i = 2 and delta = 0.96
i = 3 and delta = 0.336
i = 4 and delta = 0.1152
i = 5 and delta = 0

average over final V 0.473608
***.
***.
***.
***.

Game Over! Global score: 225 in
285 actions
make run 2.93s user 0.08s system
93% cpu 3.237 total
```

height: 4 and  $\lambda$ : 0.4

```
All constants:
width = 4, height = 5, probaIPiece
= 0.5, maxGameAction = 10000
epsilon = 1e-08, maxIteration =
100, lambda = 0.1

i = 0 and delta = 6
i = 1 and delta = 0.6
i = 2 and delta = 0.06
i = 3 and delta = 0.006
i = 4 and delta = 0.00058125
i = 5 and delta = 3.1875e-05
i = 6 and delta = 0

average over final V 0.462044
...
...
...
...
...*
.***

Game Over! Global score: 7864 in
10000 actions
make run 58.26s user 0.32s system
99% cpu 58.837 total
```

height: 5 and  $\lambda$ : 0.1

```
All constants:
width = 4, height = 5, probaIPiece
= 0.5, maxGameAction = 10000
epsilon = 1e-08, maxIteration =
100, lambda = 0.9

i = 0 and delta = 6
i = 1 and delta = 5.4
i = 2 and delta = 4.86
i = 3 and delta = 4.374
i = 4 and delta = 3.81358
i = 5 and delta = 1.88219
i = 6 and delta = 0

average over final V 0.676855
...
...
...
...
...*
.***

Game Over! Global score: 8513 in
10000 actions
make run 58.20s user 0.34s system
99% cpu 58.809 total
```

height: 5 and  $\lambda$ : 0.9

### 1.3.2 Observations

Over our executions, we remark that the  $\lambda$  has a huge impact on the optimal policy computed. On the smaller grid ( $4 \times 4$ ), it is really difficult to reach the *maxGameAction* with  $\lambda = 0.4$ , since the model will prioritize short term decisions in order to increase its score faster. In comparison, with  $\lambda = 0.1$ , the model will not gain that much per action, prioritizing safer actions that allow a higher score on the long term.

That result was as expected, since in a  $4 \times 4$  grid, the space for the risky moves is limited. The study of the game with the  $4 \times 5$  grid shows that bigger grid provides a larger space to take risks in order to increase the global score. The game with  $height = 5$  and  $\lambda = 0.9$  illustrates this perfectly because it scores higher than the game with  $height = 5$  and  $\lambda = 0.1$ , since the risky actions can be better exploited on a bigger grid.

In both cases, the total computational time is pretty acceptable (max 1 minute) even if it can be largely optimized (see Section 1.4). About the exact same improvement, the average expected gain  $V$  cannot be really discussed here since there is a huge part of states that aren't reachable in practice.

Finally, we notice that the value iteration algorithm choice described in Section 1 was pretty relevant because we can observe the quick convergence to the optimal policy (at the last iteration,  $\Delta$  is always equal to 0).

### 1.3.3 Usage

Different constants can be adjusted by modifying the corresponding `#define` statements in *src/Tetris.cpp*, *hdr/MDP.h* and *hdr/State.h*.

In order to recompile and run the whole project, execute the following command:

```
make run
```

Unlike the textual output shown in Section 1.3.1, in the submitted version of the code, there is an exhaustive description of the played game (evolution of the state at each action) as follows:

```
....          .*..          .*..  
...*.          ***.          ***.  
***. --- LPiece --> ***. -- Completion -> ***.  
***.          ***.          ***.  
Current score: 734  
  
.*.          .***          ....  
***.          ****          .***  
***. --- LPiece --> ***. -- Completion -> ***.  
***.          ***.          ***.  
Current score: 735  
  
.*.          ***.          ***.  
.***          ***.          .***  
***. --- IPiece --> ***. -- Completion -> ***.  
***.          ***.          ***.  
Current score: 735
```

## 1.4 Incoming Improvement

For now we compute naively the value iteration algorithm on all state combinations, which represents

$$\begin{aligned} &= 2^{\text{gridHeight} \times \text{gridWidth}} * \text{nbOfTrominoTypes} \\ &= 2^{4 \times 4} * 2 \\ &= 2^{17} \\ &= 131072 \end{aligned}$$

states in the smallest ( $4 \times 4$ ) version of the game. In the near future we want to reduce this number by computing the value iteration for only the reachable states (from a given starting state  $s_0$ ). For this purpose, we will make use of a Breadth First Search algorithm which will explore iteratively all reachable states from the current one.