

Lexical and Syntax Analysis

Chapter 2 Part 2

Traditional Compilation Process

- ▶ Lexical analysis – Used to find the names and numeric and operator literals (tokens)
 - Use a scanner
- ▶ Syntax analysis –
 - Determine if the program is syntactically correct
 - Find the parse tree for the program.
 - Use a parser
- ▶ Usually separate processes

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Lexical Analysis

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Identify the tokens in these programs

C

```
sum = 0;
prod = 1;
for( int j = 1; j<= 10 ; j++)
{
    sum += j;
    prod *= j;
}
```

Pascal

```
sum := 0;
prod := 1;
for j := 1 to 10 do
begin
    sum  := sum + j;
    prod := prod * j
end
```

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Calculator Language Extended BNF (Handout)

```

<progr> → <stmt-lst> .
<stmt-lst> → <stmt> { ; <stmt> }*
<stmt> → <read-stmt> | <write-stmt> | <assign-stmt>
<read-stmt> → read <id>
<write-stmt> → write <id>
<assign-stmt> → <id> = <expr>
<expr> → <expr> { <op> <expr> }* | ( <expr> ) |
<number> | <id>
<op> → + | - | * | /
<id> → <letter> { <letter> | <digit> }*
<number> → <digit> { <digit> }* [ . { <digit> }* ]

```

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Calculator Language Programs

```

height = 67.0;
width = 3.4;
area = height * width;
write area .

```

```

-----
read x;
y = 5 * ( x + 10 );
write y .

```

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Tokens for Calculator Language

<u>Token</u>	<u>Pattern</u>
read	read
write	write
id	letter { letter digit }*
number	digit {digit}* [. {digit}]*
assign	=
op	+ - * /

more tokens on
next page

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Tokens for Calculator Language –2

<u>Token</u>	<u>Pattern</u>
lparen	(
rparen)
semicolon	;
period	.

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Scanner to recognize tokens

- ▶ Scan text of program, looking for tokens
- ▶ Skip over white space (blanks, newlines,...)
- ▶ As soon as you find a token match, report token type found and continue scanning
- ▶ Match the longest substring possible
 - “3.1415 “ is the number 3.1415 , not a 3 and . and 1415
 - “reader” is the identifier *reader* , not *read* and *er*

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Pseudocode for scanner

- ▶ Distribute copies in class

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Pseudocode Scanner for Calculator Language

```

set nextchar to first char in program
while not at end of program
{
    skip over white space
    if nextchar is in set { ( , ) , + , - , * , / , ; , = , . }
        return appropriate single char token
    else if nextchar is a digit
        read additional digits
        if nextchar is a .
            read any additional digits
            return number token with number
        else
            return number token with number
    else if nextchar is a letter
        read any additional letters and digits
        check to see if resulting string is read or write
        if so, return corresponding token
        else return id token with id
    else
        return error
}

```

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Programming Project #1 Part A

- ▶ Write a Java or C++ program to implement the scanner for the Calculator Language.
- ▶ Testing: Test on lexically “correct” and “incorrect” sample programs.

See posted assignment for more details.

Due: Sept 26, 2013

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Syntax Analysis

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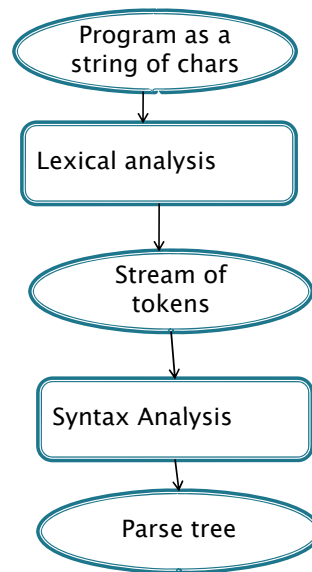
Parsing

- ▶ Given a BNF grammar description of the language
 - Determine if a program is a legal program
 - Create an (implicit) parse tree for the program
 - If program is not a legal program, return a helpful error message
- ▶ There are software tools to create a syntax analyzer given a BNF description (eg. yacc)

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Lexical + Syntax Analysis Steps



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Two Basics Kinds of Parsers

- ▶ Top down parser : Create a parse tree for a string from the top of the tree down
 - Easier to code
 - Serious Grammar Limitations
- ▶ Bottom-up parsers: Create a parse tree for a string from the bottom of the tree up
 - Complex to code but can be more efficient
 - Can handle more grammar types than top down parsing

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Illustrate top down parsing vs.
bottom up parsing with abstract
grammar:

$S \rightarrow A B \mid B S B$
 $A \rightarrow a$
 $B \rightarrow b$

Here the capital letters are the
nonterminals and lowercase letters
are the terminals.

See board work.

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Recursive Descent Parsing

- ▶ Top-down parser
- ▶ Collection of recursive subprograms
 - One for each nonterminal
- ▶ Grammar should not have any left recursive rules
 - $\langle \text{st_list} \rangle \rightarrow \langle \text{st_list} \rangle ; \langle \text{stmt} \rangle \mid \langle \text{stmt} \rangle$ BAD
 - $\langle \text{st_list} \rangle \rightarrow \langle \text{stmt} \rangle ; \langle \text{st_list} \rangle \mid \langle \text{stmt} \rangle$ OK

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Write a recursive descent parser for the following grammar.

Assignment Language

```
<prog> → <stmtList>
<stmtList> → <stmt> | <stmt> ;<stmtList>
<stmt> → <id> = <num>
<id> → x | y | z
<num> → 0|1|2|3|4|5|6|7|8|9
```

Token List: { x,y,z,0,1,2,3,4,5,6,7,8,9, ; , = }

Parser should report “successful parse” or “syntax error”.
Project does not require you to create a syntax tree.

Input: string of characters representing list of tokens.
Output: “successful parse” or “syntax error” as appropriate

Examples:

tokens

Parser returns

“x=6;y=5\$”

“successful parse”

“x=78\$”

“syntax error”

“y = 0; z = 8\$”

“syntax error” (white space)

“y=z=4\$”

“syntax error”

“x=5;y=2”

“syntax error”

“ x =8;\$”

“syntax error”

Distribute Recursive Descent Parser for
Assignment Language

Go over it in class

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Pseudocode for recursive descent parser

```
//global variables  
char nextchar; String tokens
```

```
void program() { }
```

```
void stmtList() { }
```

```
void stmt() { }
```

```
void id() {}
```

```
void num() { }
```

```
void getChar() { //gets next char from tokenList }
```

Note: tokens string
will contain only chars
from token list and
will end with a "\$" to
indicate end of input.

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```
void main() {  
    //read in tokens  
    tokens = "x=5;y=8$"  
  
    getChar();  
  
    program();  
  
    print( " Successful parse");  
}
```

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```
void program() {  
  
    stmtList();  
  
    if ( nextchar == '$')  
        print "success"  
    else  
        error();  
}
```

Grammar Rule <prog> → <stmtList>

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```
void stmtList() {  
    stmt();  
    if (nextchar == ';') {  
        getChar();  
        stmtList();  
    }  
    else  
        return;  
}
```

Grammar Rule
 $\langle \text{stmtList} \rangle \rightarrow \langle \text{stmt} \rangle$
 $| \langle \text{stmt} \rangle ; \langle \text{stmtList} \rangle$

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```
void stmt() {  
    id();  
  
    if (nextchar == '=' )  
    {  
        getchar();  
        num();  
    }  
    else  
        error();  
}
```

Grammar Rule
 $\langle \text{stmt} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{num} \rangle$

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```
void id()
{
    if (nextchar == 'x' || nextchar == 'y'
        || nextchar == 'z')
    {
        getchar();
        return;
    }
    else
        error();
}
```

Grammar Rule
<id> → x | y | z

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```
void num()
{
    if (isdigit(nextchar) )
    {
        getchar();
        return;
    }
    else
        error();
}
```

Grammar Rule
<num> → 0|1|2|3|4|5|6|7|8|9

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Programming Project #1 Part B

- ▶ Modify the recursive descent parser for the Assignment Language so that it prints helpful error messages. Hint: Modify the error method so that it accepts a string as a parameter. This will be the error message.
- ▶ Thoroughly test your program to show all of the errors it can detect and report.
- ▶ Download parser from course Moodle site. See Moodle for complete description of Project 1 Part B

Due Sept 26, 2013