

# MTH251

Data Structures and Algorithms I

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Singapore**

# Course Structure

6 weeks (Jan ~ Mar)/6 seminars/6 labs:

1. Jan 24: Python, Complexity & Big O
2. Feb 07: Array, Stack, Queue, Recursion
3. Feb 14: Linked List
4. Feb 21: Tree
5. Feb 28: Lineary Search, Binary Search, Algorithm Design & Pattern
6. Mar 07: Review, Q&A

Mar 07: Resit Online Session

3 assignments & open book exam:

Assessment	Description	Weight Allocation
Assignment 1	Tutor-Marked Assignment 1	10%
Assignment 2	Tutor-Marked Assignment 2	10%
Assignment 3	Tutor-Marked Assignment 3	10%
Examination	Open book exam	70%
<b>TOTAL</b>		100%

# Learning Objectives

1. python
2. algorithm time & space complexity
3. basic data structure: **array** **stacks** **queues** **list** **tree**
4. recursion algorithm

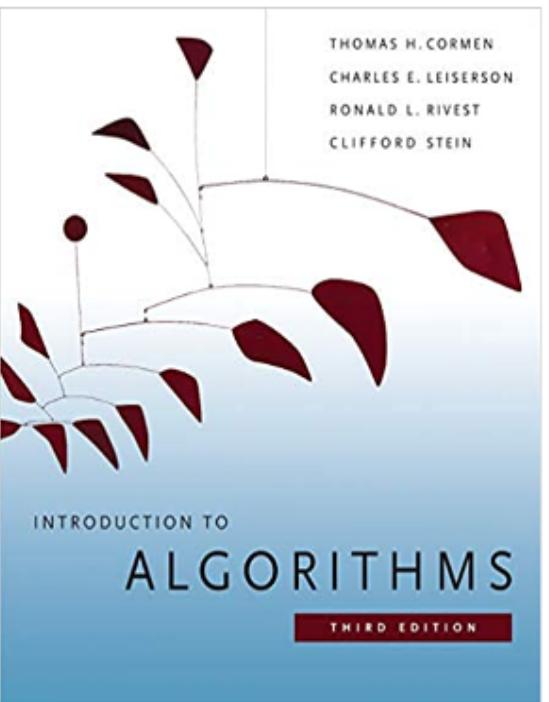
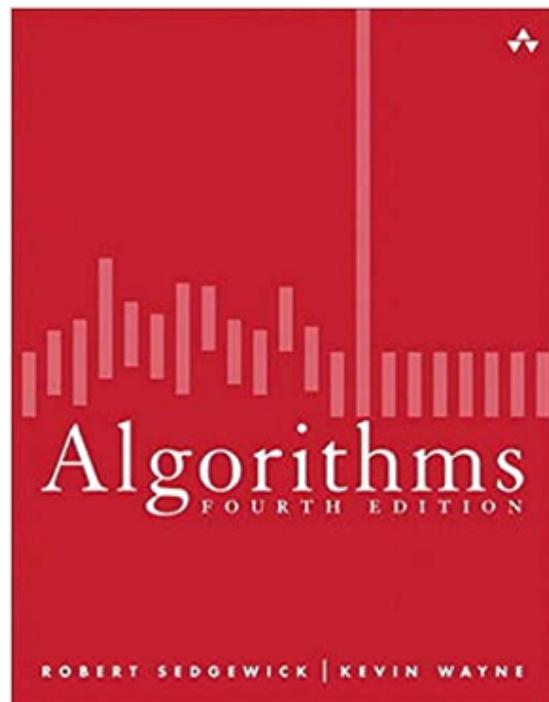
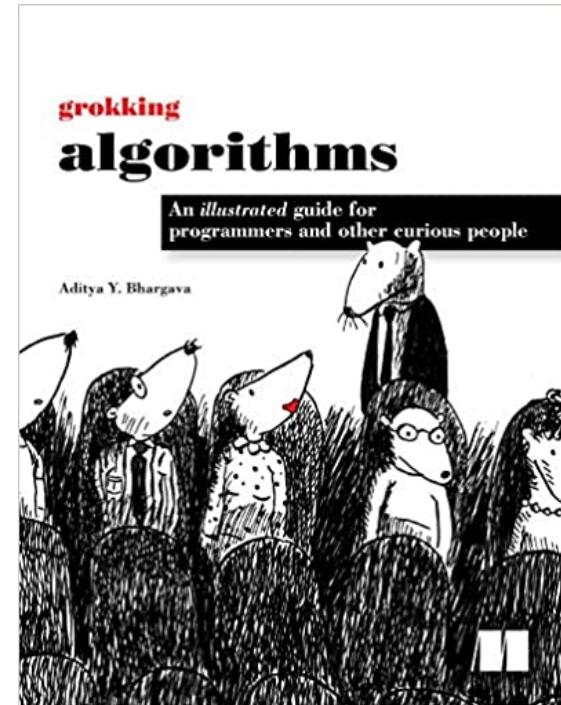
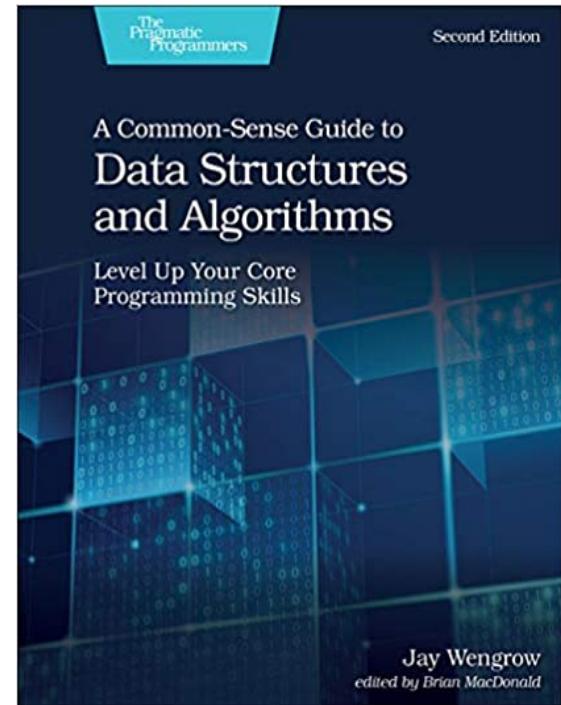
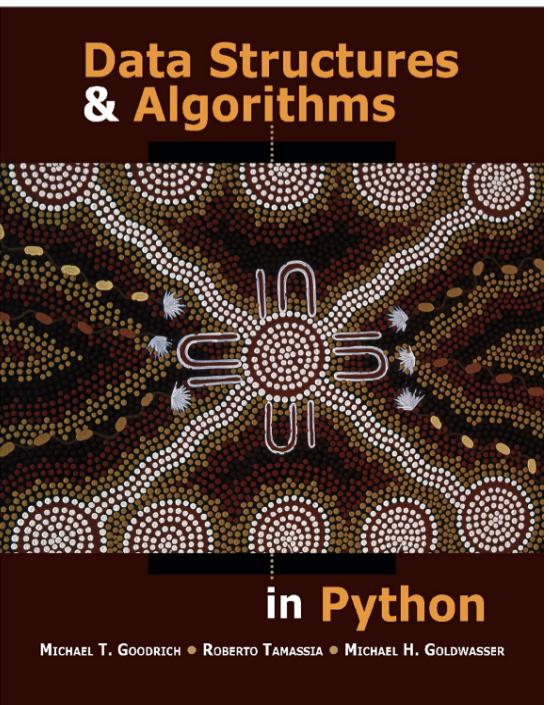
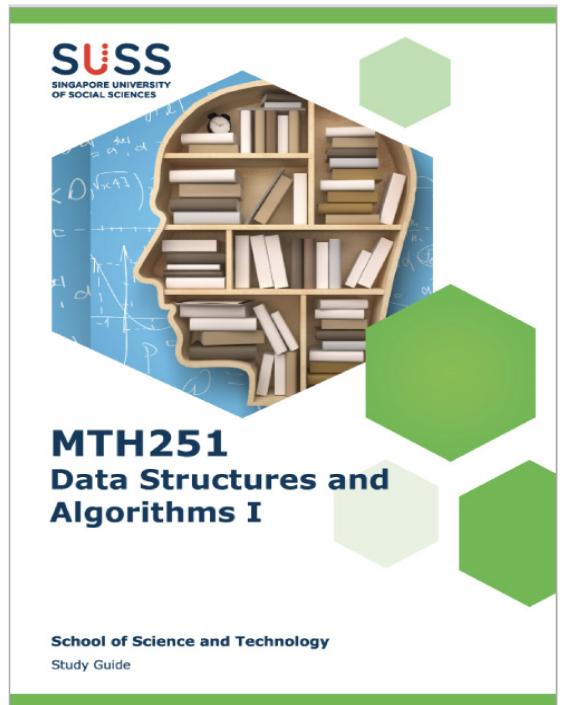
# Slides & Notebooks

slides online: <https://mth251.fastzhong.com/>

 <https://mth251.fastzhong.com/mth251.pdf>

 labs: <https://github.com/fastzhong/mth251/tree/main/public/notebooks>

# Learning Resource



# Learning Resource

A screenshot of the Leetcode website's "Learn" section. The page displays a grid of 12 learning modules, each represented by a colored card with a title, chapter count, item count, and completion percentage. A "Premium" badge is visible at the top right of the grid. Navigation links like "Back to Explore" and "Learn" are at the top left, and account-related icons are at the top right.

Module Title	Chapters	Items	Completion (%)
Detailed Explanation of Graph	6	56	Locked
Introduction to Data Structure Arrays 101	6	31	0%
Introduction to Data Structure Linked List	5	30	0%
Introduction to Data Structure Binary Tree	3	19	0%
Introduction to Algorithms Recursion I	5	21	0%
Introduction to Algorithms Recursion II	4	25	0%
Basic Concepts in ML Machine Learning 101	3	9	0%
Learning about... Binary Search	8	30	0%
Introduction to Data Structure N-ary Tree			
Introduction to Data Structure Binary Search Tree			
Introduction to Data Structure Trie			
Introduction to Data Structure Hash Table			

# Why Python

*The TIOBE Programming Community index is an indicator of the popularity of programming languages.*

Oct 2021	Oct 2020	Change	Programming Language	Ratings	Change
1	3	▲	 Python	11.27%	-0.00%
2	1	▼	 C	11.16%	-5.79%
3	2	▼	 Java	10.46%	-2.11%
4	4		 C++	7.50%	+0.57%
5	5		 C#	5.26%	+1.10%
6	6		 Visual Basic	5.24%	+1.27%
7	7		 JavaScript	2.19%	+0.05%
8	10	▲	 SQL	2.17%	+0.61%

# Why Python

- ✓ Easy To Learn
- ✓ Human Readable
- ✓ Productivity
- ✓ Cross Platform

## The Zen of Python, by Tim Peters

Beautiful is better than ugly.

Explicit is better than implicit.

Simple is better than complex.

Readability counts.

Special cases aren't special enough to break the rules.

There should be one-- and preferably only one --obvious way to do it.

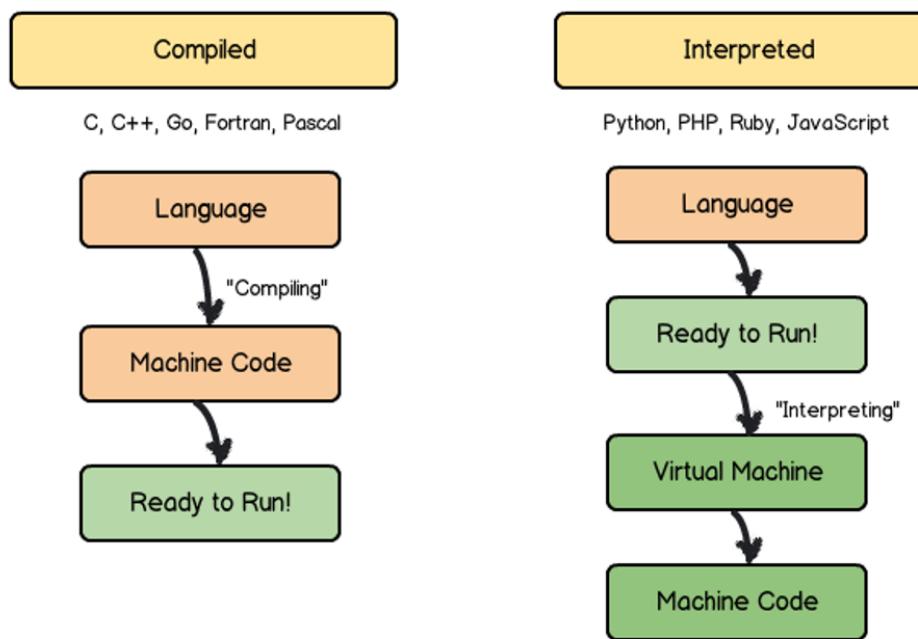
If the implementation is hard to explain, it's a bad idea.

# Python Jobs

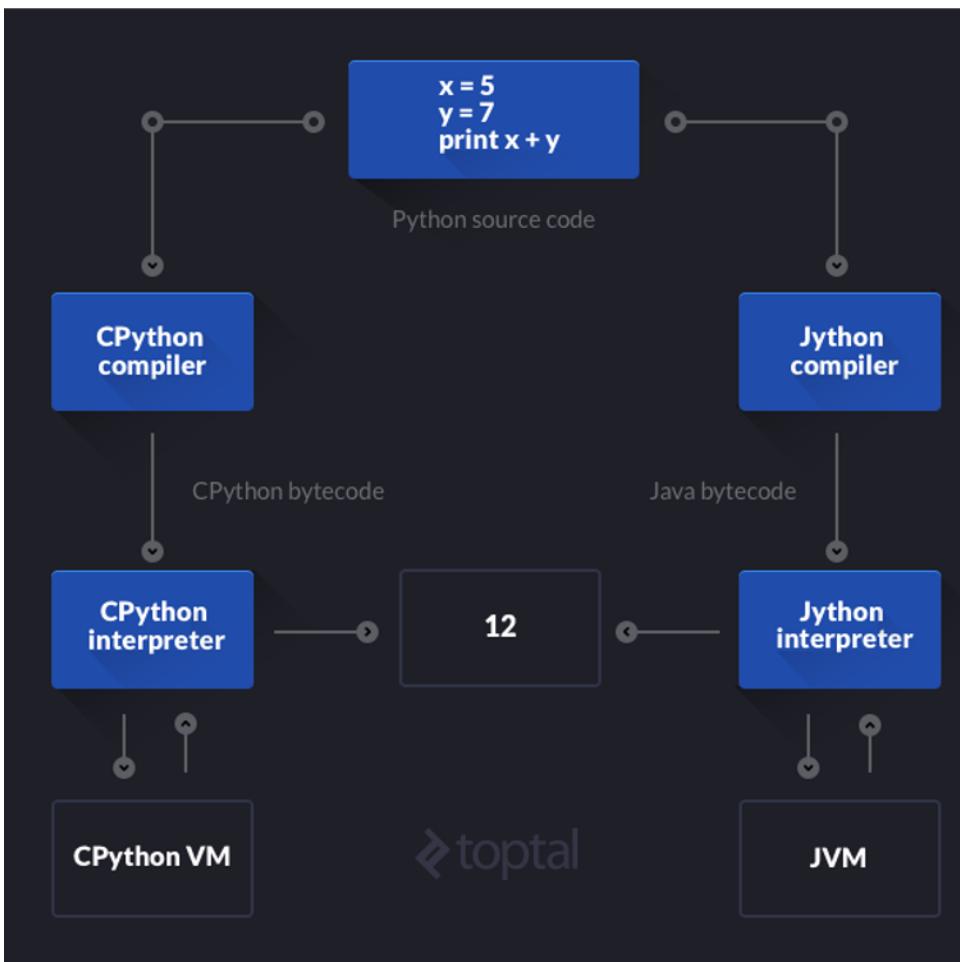
- ✓ backend: **Python** vs. Java, C++, Go, Php
- ✓ devops: **Python** vs. Go, Ruby, Shell
- ✓ test automation: **Python** vs. Groovy, shell
- ✓ data engineering: **Python** vs. Java, C++
- ✓ data analytics & visualization: **Python** vs. R, Java, C++
- ✓ data science & machine learning: **Python** vs. R, Julia, C++

# Python 101

- Compiled vs. Interpreted



- CPython bytecode



- Python implementations

Implementation	Virtual Machine	Ex) Compatible Language
CPython	CPython VM	C
Jython	JVM	Java
IronPython	CLR	C#
Brython	JavaScript engine (e.g. V8)	JavaScript
RubyPython	Ruby VM	Ruby

# Python Data Type & Operators



- numbers: `int` `float` `complex`
  - arithmetic operator: `+` `-` `*` `/` `//` `%` `**`
  - bitwise operator: `&` `|` `^` `>>` `<<` `~`
  - `range()`: a list of integers
- strings: `''` `'''` `''` `"` `\t` `\n` `\r` `\\"` etc.
  - `join()` `split()` `ljust()` `rjust()` `lower()` `upper()` `lstrip()` `rstrip()` `strip()` etc.
- boolean: `True` `False`
  - `True`: non-zero number, non-empty string, non-empty list
  - `False`: `0`, `0.0`, `""`, `[]`, `None`

# Python Data Type & Operators



- boolean: `True` `False`
  - logic operator: `and` `or` `not`
  - comparison operator: `>` `<` `>=` `<=` `==` `!=`
  - identity operator: `is` `is not`
- None
- type conversion/casting: `int()` `float()` `str()` `bool()` `hex()` `ord()`

# Python Collections



- collections: `list` `tuple` `set` `dictionary`
  - membership operator: `in` `not in`
- `list []`: a collection of items, usually the items all have the same type
  - sequence type, sortable, grow and shrink as needed, most widely used
- `tuple ()`: a collection which is ordered and unchangeable
- `set {}`: a collection which is unordered and unindexed
- `dictionary`: a set of `key: value` pairs, unordered, changeable and indexed

# Python Program Structure



- variable
- statement & comments
  - Python uses new lines to complete a command, as opposed to other programming languages often use ; or () ; relies on indentation (whitespace sensitive), to define scope, such as the scope of loops, functions and classes, as opossed to other programming languages often use {}
- control flow
  - if ... elif ... else
  - while for break continue

# Python Program Structure



- function

- `def` `return`

- `main`

- advanced:

- `lambda`
    - `decorator`
    - `closure`

- error/exception

- handling exception: `try ... except ... else ... finally`

- raise exception: `raise`

# Python OO & Class

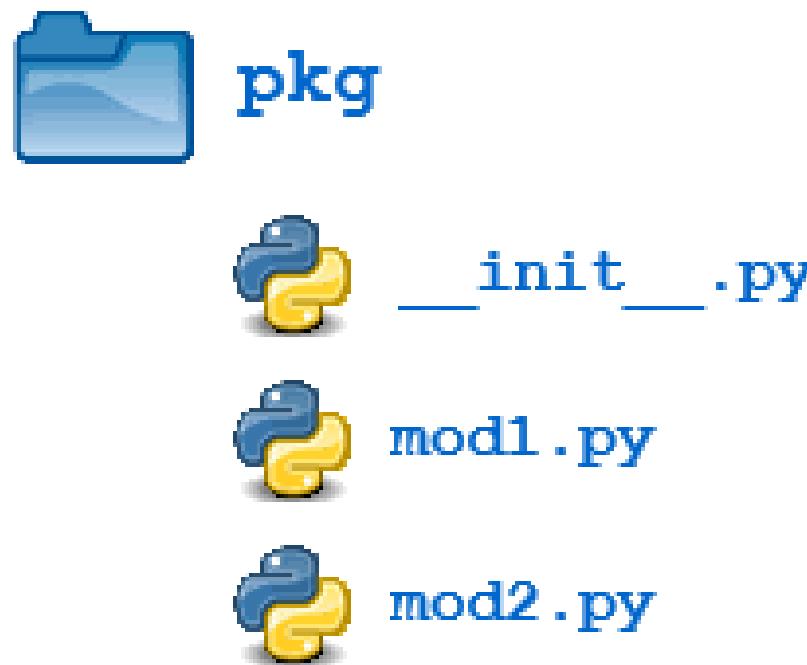


- Procedural vs. OOP vs. FP
- OO Principal
  - Inheritance
  - Encapsulation
  - Polymorphism
- class, instance, attributes, properties, method
- override vs. overload vs. overwrite

# Misc.



- modular programming: function → class → module → package
- Modules: Python module (default main module), C module, Build-in module
- Packages:



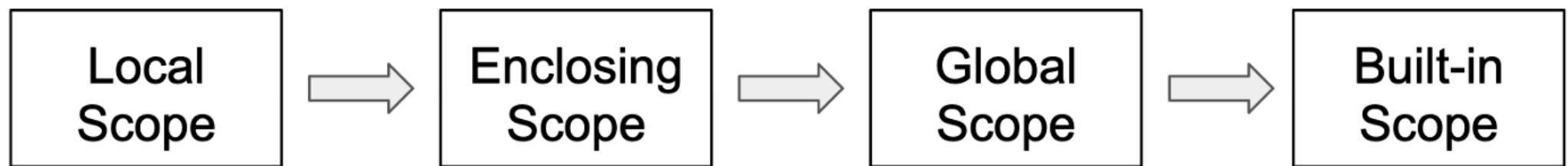
- Standard Lib: math random re os itertools collections

# Misc.



- Namespaces & Scopes: `LEGB` rule

## The LEGB Rule



- memory, copy vs. deepcopy
- help

Cheatsheet:



[Python Crash Course - Cheat Sheets](#)



[Comprehensive Python Cheatsheet](#)

# Python Tutorials

## Programming with Mosh

- ▶ [Python Tutorial - Python for Beginners 2020](#)

## freeCodeCamp

- ▶ [Learn Python - Full Course for Beginners Tutorial](#)
- ▶ [Python for Everybody - Full University Python Course](#)
- ▶ [Intermediate Python Programming Course](#)

## Tech With Tim

- ▶ [Learn Python - Full Course for Beginners Tutorial](#)

## The Hitchhiker's Guide to Python!

# Data Structure & Algorithms (DSA)

A data structure is a way of organizing information so that it can be used effectively by computer

Algorithms provides computer step by step instructions to process the information and solve a problem

**Program = Data Structure + Algorithm**

- Input
- Output
- Definiteness
- Effectiveness
- Finiteness

# Data Structure & Algorithms

Example: keyword searching

**Data Structure:** index

**Algorithm:** looking up the page no.

738

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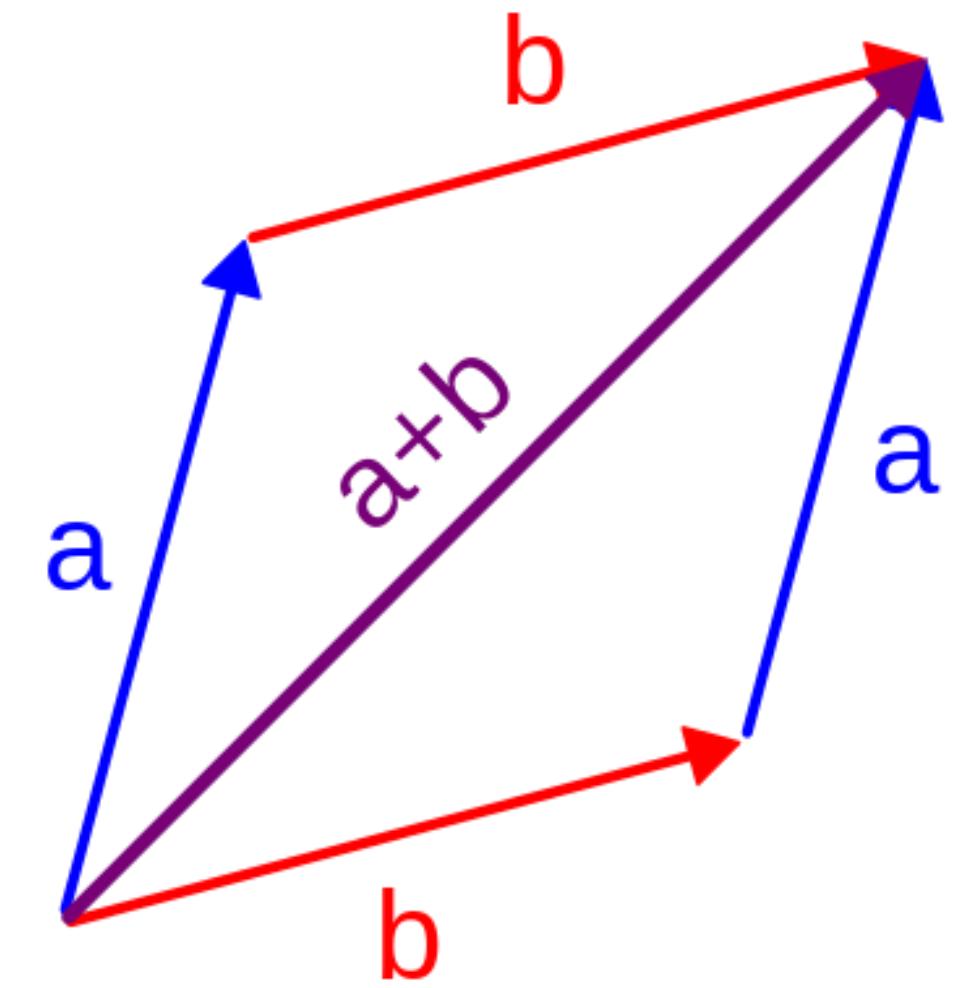
# Data Structure & Algorithms

Example: adding two vectors

**Data Structure:** complex number

**Algorithm:** adding two complex numbers

$$a + b = (x + yi) + (u + vi) = (x + u) + (y + v)i$$



# Data Structure & Algorithms

## Data Structure

- Linear
  - Array, String, Linked List
  - Stack, Queue, Deque, Set, Map/Hash, etc.
- Non-Linear
  - Tree, Graph
  - Binary Search Tree, Red-Black Tree, AVL, Heap, Disjoin Set, Trie, etc.
- Others
  - Bitwise, BloomFilter, LRU Cache

# Algorithms & Data Structure

## Algorithms

- branching: if-else, switch
- iteration: for, while loop
- recursion: divide & conquer, backtrace
- searching: binary search, depth first, breath first, A\*, etc.
- sorting: quick sort, bubble sort, merge sort, etc.
- dynamic programming
- greedy
- ...

# Algorithms & Data Structure

## Why

- ✓ deeper understanding of computer system
- ✓ improve coding skill
- ✓ coding interview
- ✓ building powerful framework and library

## How

 learning by doing, implementing from scratch

 problem solving

# Algorithm Complexity Analysis



How to measure Performance/Efficiency ?

- cpu, memory, io, networking, etc.
- no. of lines
- worst case vs. best case
- code slows as data grows
- ...

# Algorithm Complexity Analysis

**Time Complexity** : by giving the size of the data set as integer N, consider the number of operations that need to be conducted by computer before the algorithm can finish

**Space Complexity** : by giving the size of the data set as integer N, consider the size of extra space that need to be allocated by computer before the algorithm can finish

Good code:

- ✓ readability
- ✓ speed
- ✓ memory

👉 When: Accessing, Searching, Inserting, Deleting, ...

# Big-O



**Lef  $f(n)$  and  $g(n)$  be functions from positive integers to positive integers to positive reals  
 $f = O(g)$  if there is a constant  $c > 0$  such that  $f(n) \leq c \cdot g(n)$  for large  $n$**

⌚  $f(n)$  grows no faster than  $\underline{\underline{g(n)}}$

e.g.

$$f(n) = O(4n^2 + 8n + 16) \rightarrow O(n^2) \rightarrow g(n) = n^2$$

$$O(4n^2 + 8n + 16) = O(n^2)$$

*Big-O describes the trend of algorithm performance when the data size increases*

# Big-O

---

$O(1)$ : constant

---

$O(\log_* n)$ : logarithmic

---

$O(n)$ : linear

---

$O(n \log_* n)$ : linearithmic

---

$O(n^2)$ : polynomial

---

$O(2^n)$ : exponential

---

$O(n!)$ : factorial

---

# Big-O

## 👉 Master theorem (analysis of algorithms)

$$O(f) = f$$

$$O(c \cdot f) = O(f)$$

$$O(f + g) = O(\max(f, g))$$

$$O(f) \cdot O(g) = O(f \cdot g)$$

$$O(f \cdot g) \leq O(f \cdot h) \text{ if & only if } O(g) \leq O(h)$$

$$O(x^a) \leq O(x^b) \text{ if & only if } a \leq b$$

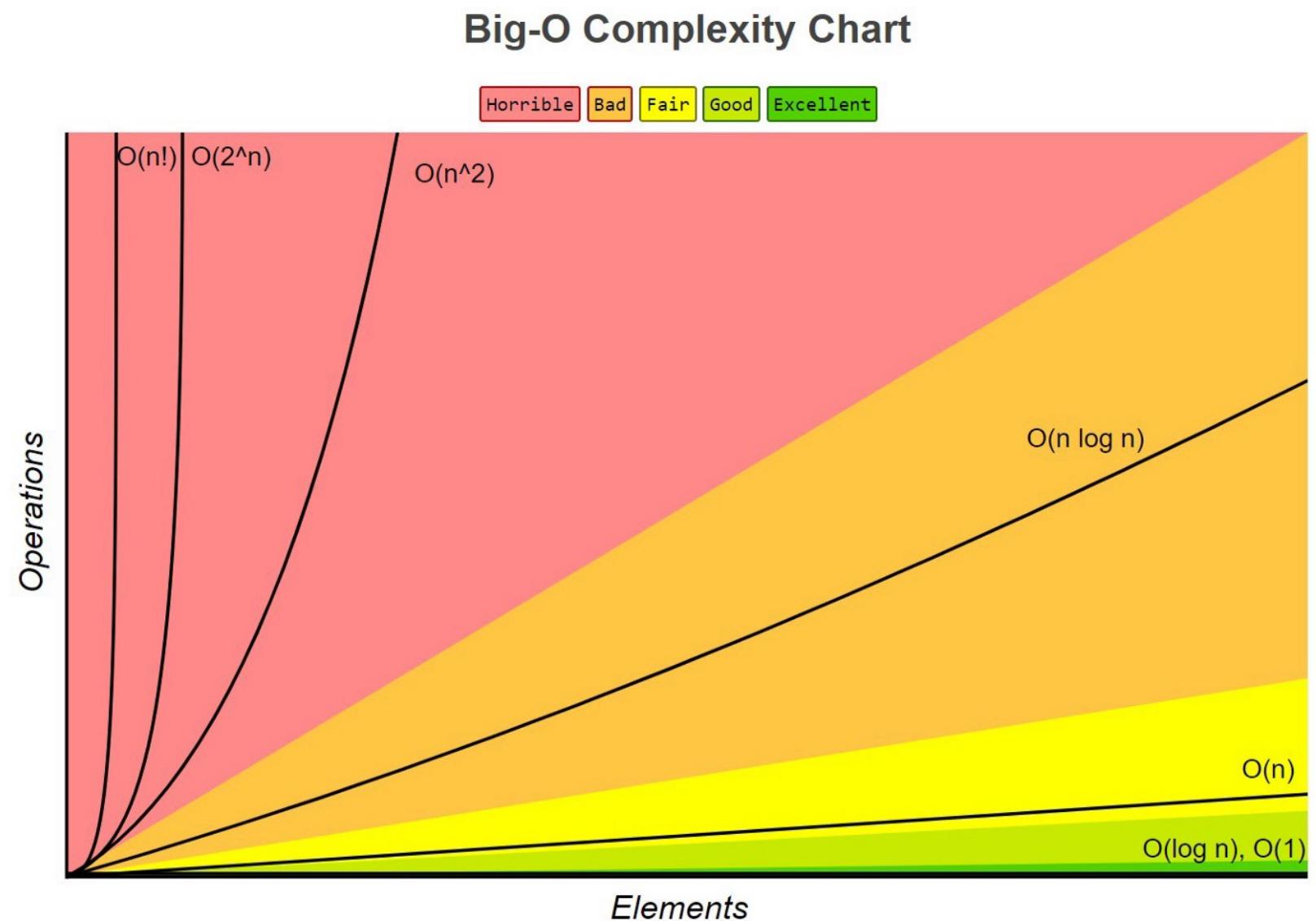
$$O(a^x) < O(b^x) \text{ if & only if } a < b$$

$$O(x^c) < O(d^x) \text{ if & only if } d > 1 \text{ (assuming } c \geq 1 \text{ and } d \geq 1\text{)}$$

$$O(\log_* x) < O(x^c) \text{ if & only if } c > 0$$

# Big-O

😊  $O(1) < O(\log_* n) < O(n) < O(n \log_* n) < O(n^2) < O(2^n) < O(n!)$  😢



↗ <https://www.bigocheatsheet.com/>

# Array

To store a list of similar things, example:

A list of names: ["Alex", "Bob", "Charles", "David"]

A list of numbers: [1, 2, 3, 4]

Each item in the array referred as "**element**"

# Array

- Element Type: same type (array is structured data)
- Element Size: fixed

```
1 # java
2 String[] cars = {"BMW", "Toyota", "Tesla"} // declare & init
3
4 Integer[] scores = new Integer[10] // declare
5 // init
6 scores[0] = 90
7 scores[1] = 80
```

- Element Index: 0, 1, ..., length - 1

# Array 2-D

```
students = [  
    ["Alex", "M", "S1111111A"],  
    ["Bob", "M", "S2222222B"],  
    ["James", "M", "S3333333C"],  
]
```

students[2] → ["James", "M", "S3333333C"]

Students[1][2] → "S2222222B"

Index	0	1	2
0	Alex	M	S1111111A
1	Bob	M	S2222222B
2	James	M	S3333333C

# Array Address

str = "HELLO" = ['H', 'E', 'L', 'L', 'O']

Memory Address	2160	2161	2162	2163	2164	2165	2166	2167	2168	2169	...
Computer's memory	H	E	L	L	L	O	O	O	O	...	
Index Number	0	1	2	3	4	5	6	7	8	9	...

data type: char

data type size: 2 byte (1 byte = 8 bits, 0000 0000 ~ 1111 1111)

**total\_size = array\_size \* data\_type\_size**

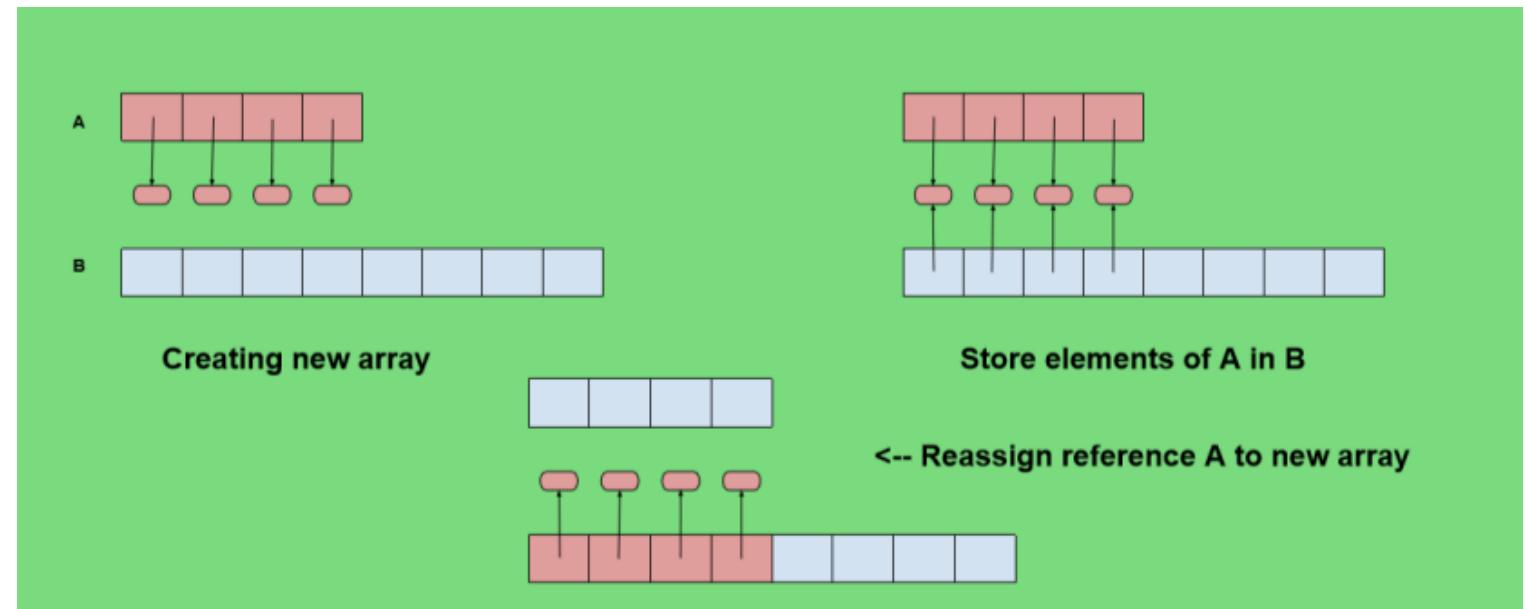
**array[i].address = base\_address + i \* data\_type\_size**

👉 O(1)

# Dynamic Array



- what is the good space/slot size for an array?
- when is the good time to expand/shrink the array?



## Growth Pattern:

- Python:  $NewAllocated = (size_t)newsize + (newsize >> 3) + (newsize < 9?3 : 6); // 0, 4, 8, 16, 25, 35, 46, 58, 72, 88, \dots$
- Java:  $((size * 3)/2) + 1$
- C#:  $size * 2$

👉 memory management is one of the biggest programming challenges.

# Array Complexity

Operation	Array	Dynamic Array
Accessing	$O(1)$	$O(1)$
Searching	$O(n)$	$O(n)$
Inserting	-	$O(n)$
Deleting	-	$O(n)$

# Python Collection Complexity

List [a, b, c, ...]	Dicts {k:v, ...}	Set {a, b, c,}
mylist.append(val)	O(1)	mydict[key] = val
mylist[i]	O(1)	mydict[key]
val in mylist	O(N)	key in mydict
for val in mylist:	O(N)	for key in mydict:
mylist.sort()	O(NlogN)	

# Trade-offs

```
1  # .. make a list ..
2  if thing in my_list: # O(N)
3
4  # ✅ Good
5  # .. make a set ..
6  if thing in my_set: # O(1)
7
8  # ❌ Bad
9  # .. make a list ..
10 my_set = set(my_list) # O(N)
11 if thing in my_set: # O(1)
12
13 # ✅ Good
14 # .. make a list ..
15 my_set = set(my_list) # O(N)
16 for many_times:
17     if thing in my_set: # O(1)
18
```

# ADT vs. Data Structure

An **abstract data type** (ADT) is an abstraction of a **data structure** which provides only the interface to which a data structure must adhere to. The interface does not give any specific details about how something should be implemented - ADT provides implementation-independent view of a data structure.

Programming language provides different **data types** to implement/represent a specific data structure.

- Array - a linear abstract data type
- Array - a java data type
- Dynamic Array - array with changable size
- List - a python data type, more flexible than a dynamic array
- ArrayList/Vector - java data type, implementation of List

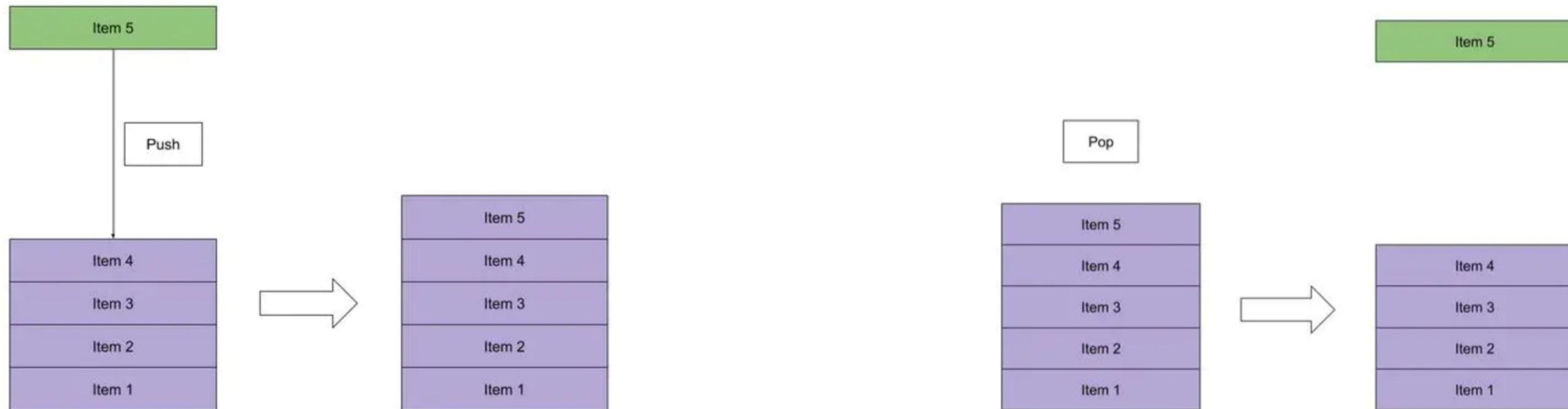
# Stack

- Sequential Access vs Random Access (such as Array)
- **LIFO** (Last In First Out) sequential collection



# Stack: Operations

- **push()** – pushing (storing) an element on the stack
- **pop()** – removing (accessing) an element from the stack
- top()/peek() – get the top data element of the stack, without removing it
- size(), isEmpty(), isFull()



# Stack Complexity

Operation	Stack
Accessing	$O(n)$
Searching	$O(n)$
Inserting	$O(1)$ (push)
Deleting	$O(1)$ (pop)

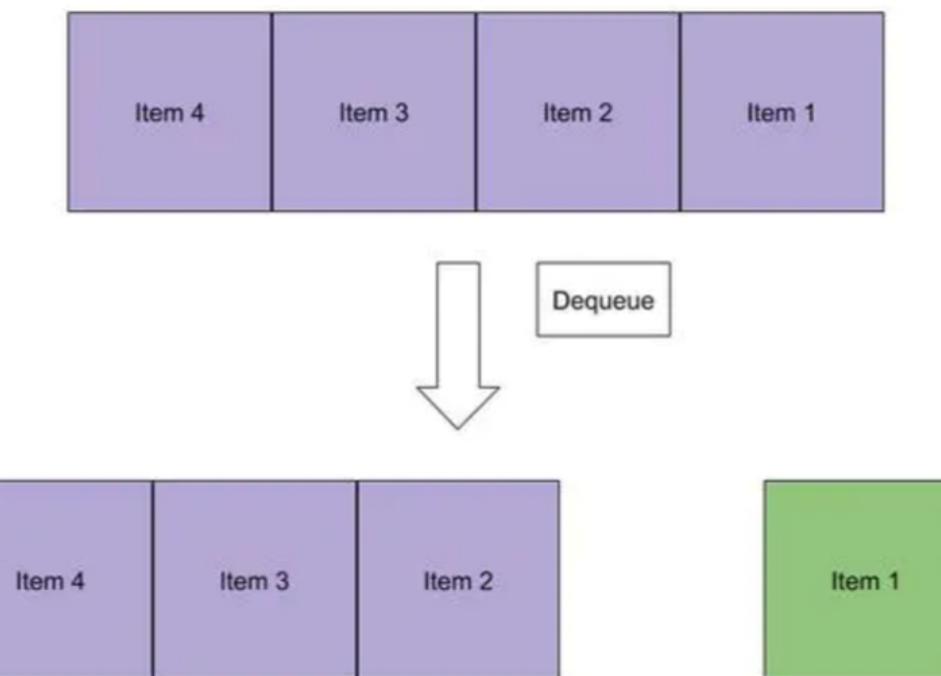
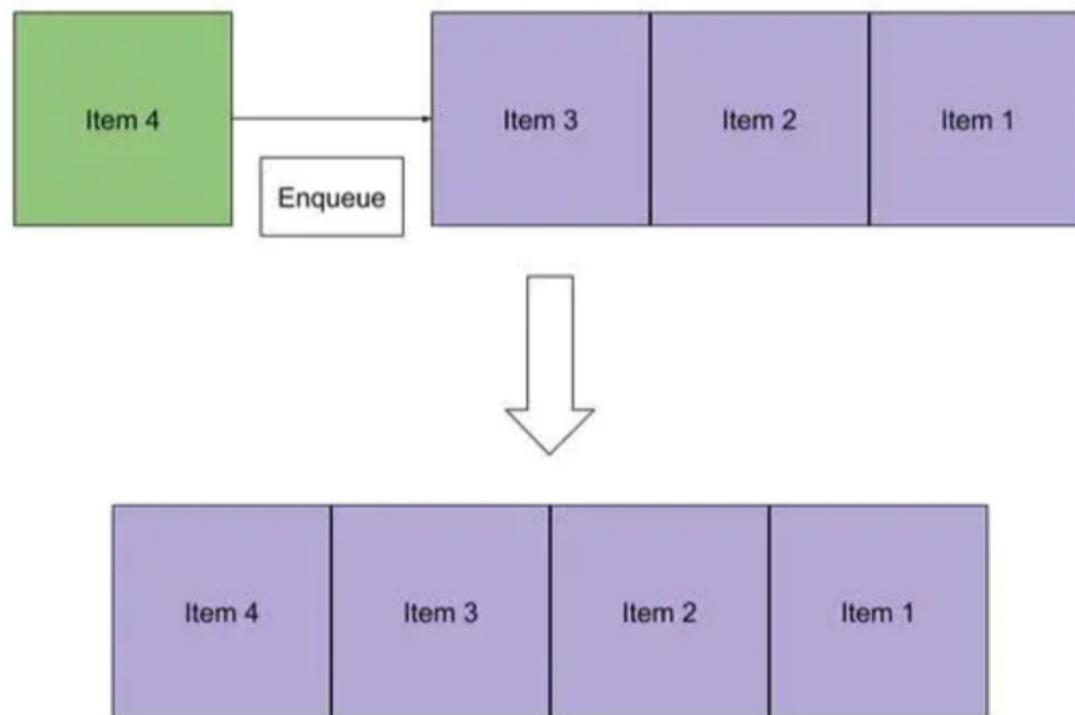
# Queue

- **FIFO** (First In First Out) sequential collection



# Queue: Operations

- **enqueue()** – adding (storing) an element to the queue
- **dequeue()** – removing (accessing) an element from the queue
- `fist()/peek()` – get the first element of the queue, without removing it
- `size()`, `isEmpty()`, `isFull()`



# Queue Complexity

Operation	Queue
Accessing	$O(n)$
Searching	$O(n)$
Inserting	$O(1)$ (enqueue)
Deleting	$O(1)$ (dequeue)

# Linked List

- dynamic linear data structure
- each item contains data & pointer
- data stored in a “Node” class
- each item holds a relative position relative to the other items: 1st, 2nd, ..., last item



# List: Operations (potential)

- add(element)
- append(element)
- insert(position, element)
- delete(element)
- is\_empty()
- \_\_len\_\_()
- set(position, element)
- get(position)
- search(element)
- index(element)
- pop()
- pop(position)

# Linked List Complexity

Operation	Linked List	Dynamic Array
Accessing	$O(n)$	$O(1)$
Searching	$O(n)$	$O(n)$
Inserting	$O(1)$ 🤔	$O(n)$
Deleting	$O(1)$ 🤔	$O(n)$

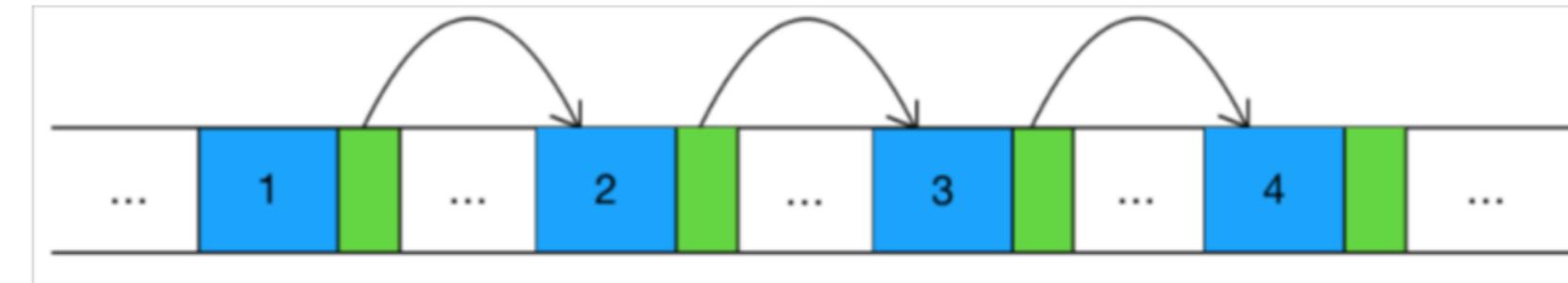
# Linked List vs. Array

- ✓ dynamic, no need to deal with fixed memory size
- ✗ accessing speed

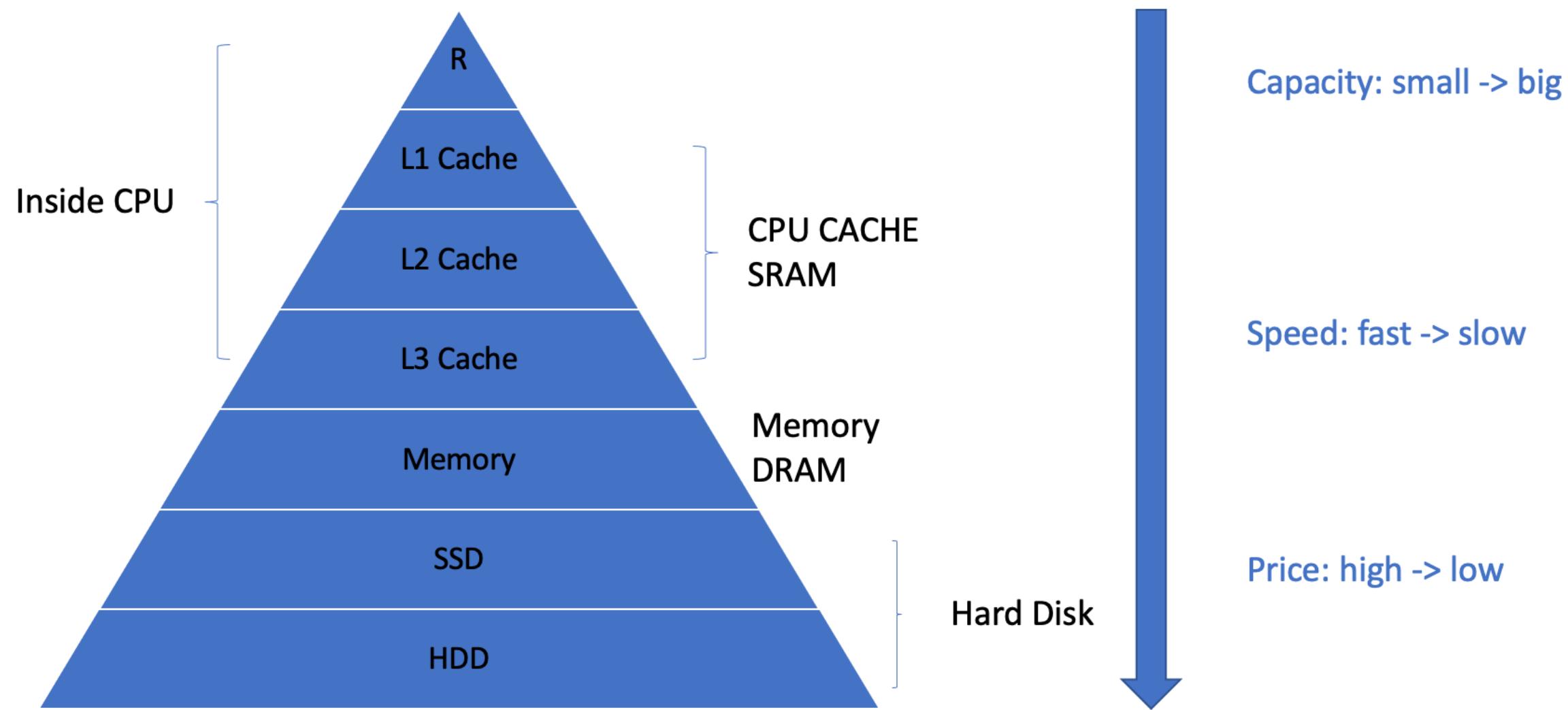
Array:

Memory Address	2160	2161	2162	2163	2164	2165	2166	2167	2168	2169	...
Computer's memory	H		E		L		L		O	...	
Index Number	0		1		2		3		4		

Linked List:

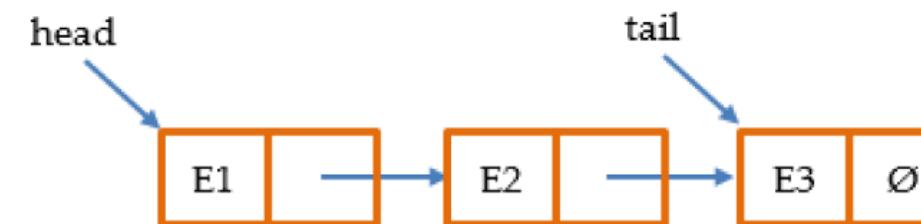


# Linked List vs. Array



# Linked List

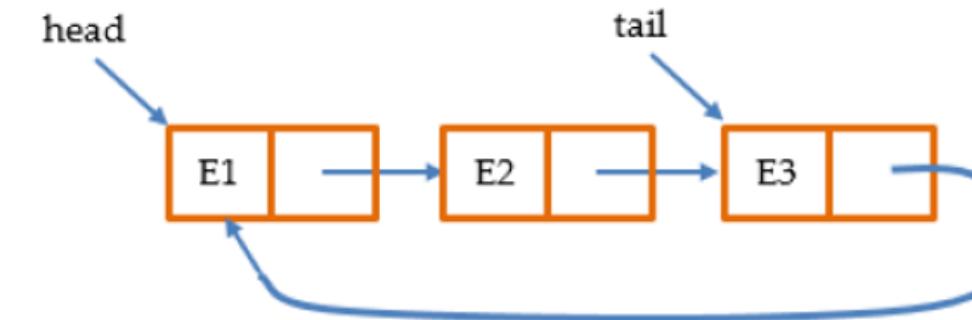
Linked List



Doubly Linked List



Circular Linked List



Positional Linked List



# Circular Linked List

a linked list where all nodes are connected to form a circle

- no null at then end
- can iterate from any node
- e.g. for cpu job list - OS putting running applications in a list and then to cycle through them by giving each of them a slice of time to execute, and then making them wait, when it reaches the end of the list, it can cycle around to the front of the list

circular singly linked list, circular doubly linked list, sorted circular linked list

# Circular Linked List: Operations (potential)

- append(element)
- delete(element)
- search(element)
- is\_empty()
- \_\_len\_\_()

# Positional Linked List

- use Position class instead of index
- remove iteration

# Positional Linked List: Operations (potential)

- first()
- last()
- before(position)
- after(position)
- set(position, element)
- search(element)
- is\_empty()
- \_\_len\_\_()
- add\_first(element)
- add\_last(element)
- add\_before(position, element)
- add\_after(position, element)

# Recursion

Recursion is the process of defining a problem (or the solution to a problem) in terms of (a simpler version of) itself

👉 Recursion by definition is a function that calls itself.

# Recursion

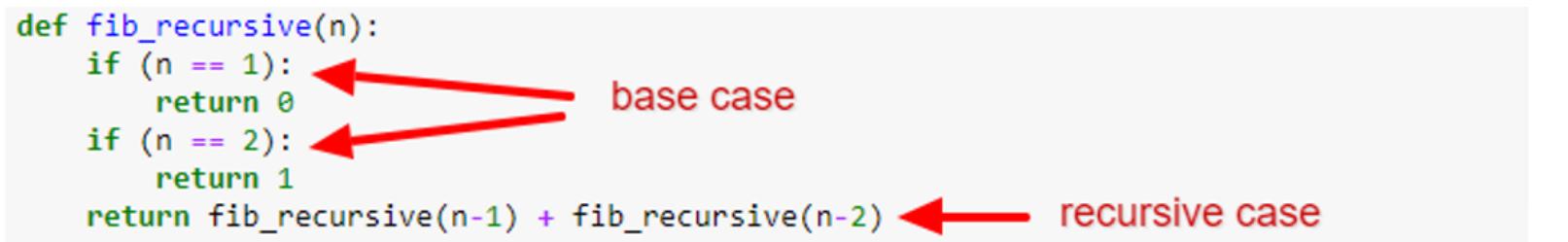
- base case
- recursive case

Example:

Fibonacci sequence 0, 1, 1, 2, 3, 5, 8, ...

- when  $n = 1$ ,  $\text{fib}(1) = 0$
- when  $n = 2$ ,  $\text{fib}(2) = 1$
- when  $n > 2$ ,  $\text{fib}(n) = \text{fib}(n-1) + \text{fib}(n-2)$

```
def fib_recursive(n):
    if (n == 1):
        return 0
    if (n == 2):
        return 1
    return fib_recursive(n-1) + fib_recursive(n-2)
```



The code defines a recursive function for the Fibonacci sequence. It has two base cases: if  $n = 1$ , it returns 0; if  $n = 2$ , it returns 1. For  $n > 2$ , it returns the sum of the previous two Fibonacci numbers, obtained by recursive calls to  $\text{fib\_recursive}(n-1)$  and  $\text{fib\_recursive}(n-2)$ . Red arrows point from the text 'base case' to the first two `if` statements, and from 'recursive case' to the final `return` statement.

# Recursion vs. Iterative

*Anything with a recursion can be done iteratively (loop)*

😊 Intuitive/DRY, code readability

```
def fib_recursive(n):
    if (n == 1):
        return 0
    if (n == 2):
        return 1
    return fib_recursive(n-1) + fib_recursive(n-2)
```

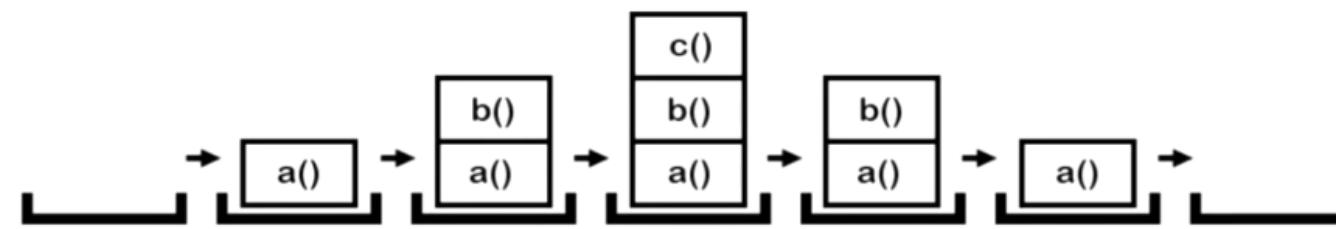
base case  
recursive case

🤔 Optimization, call stack

```
def fib_iterative(n):
    fib = [0,1]
    i = 2
    while (i < n): # index: 2, 3, ..., n-1
        fib.append(fib[i-1] + fib[i-2]) # fib[i]
        i += 1
    return fib[n-1] # the nth number
```

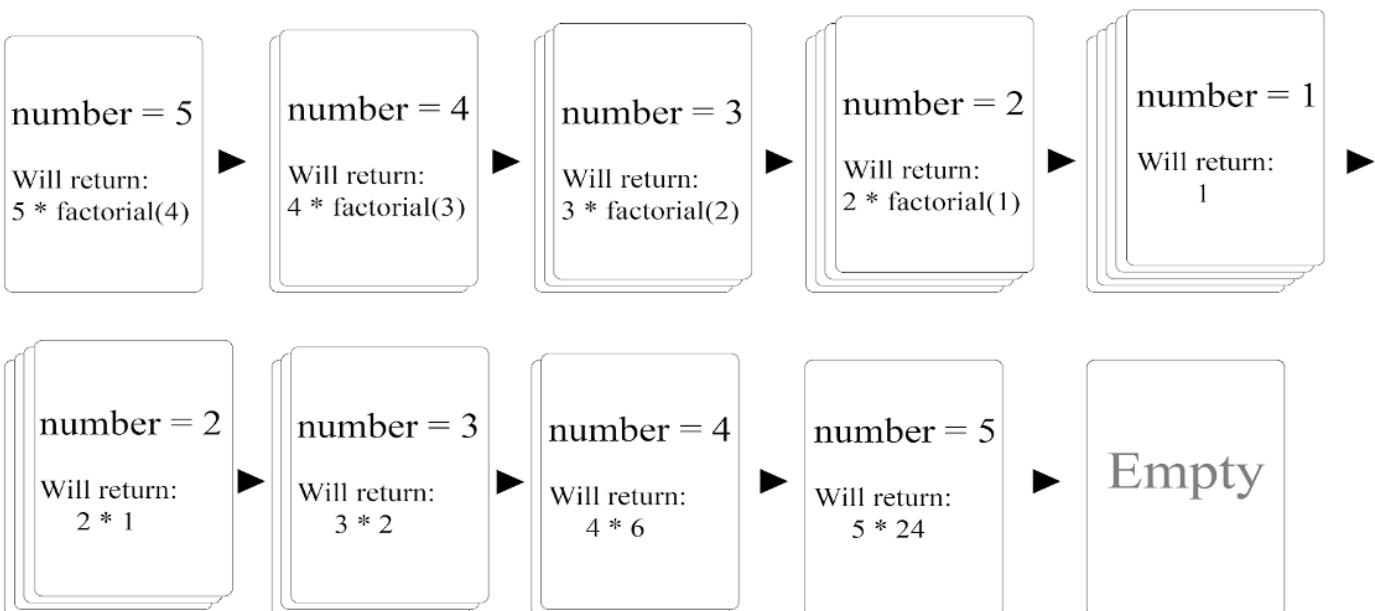
# Recursive Call

Call Stack:



The “call stack” is a stack of “frame objects”.  
(frame object == a function call)

fib\_recursive(5):



# Recursive Call

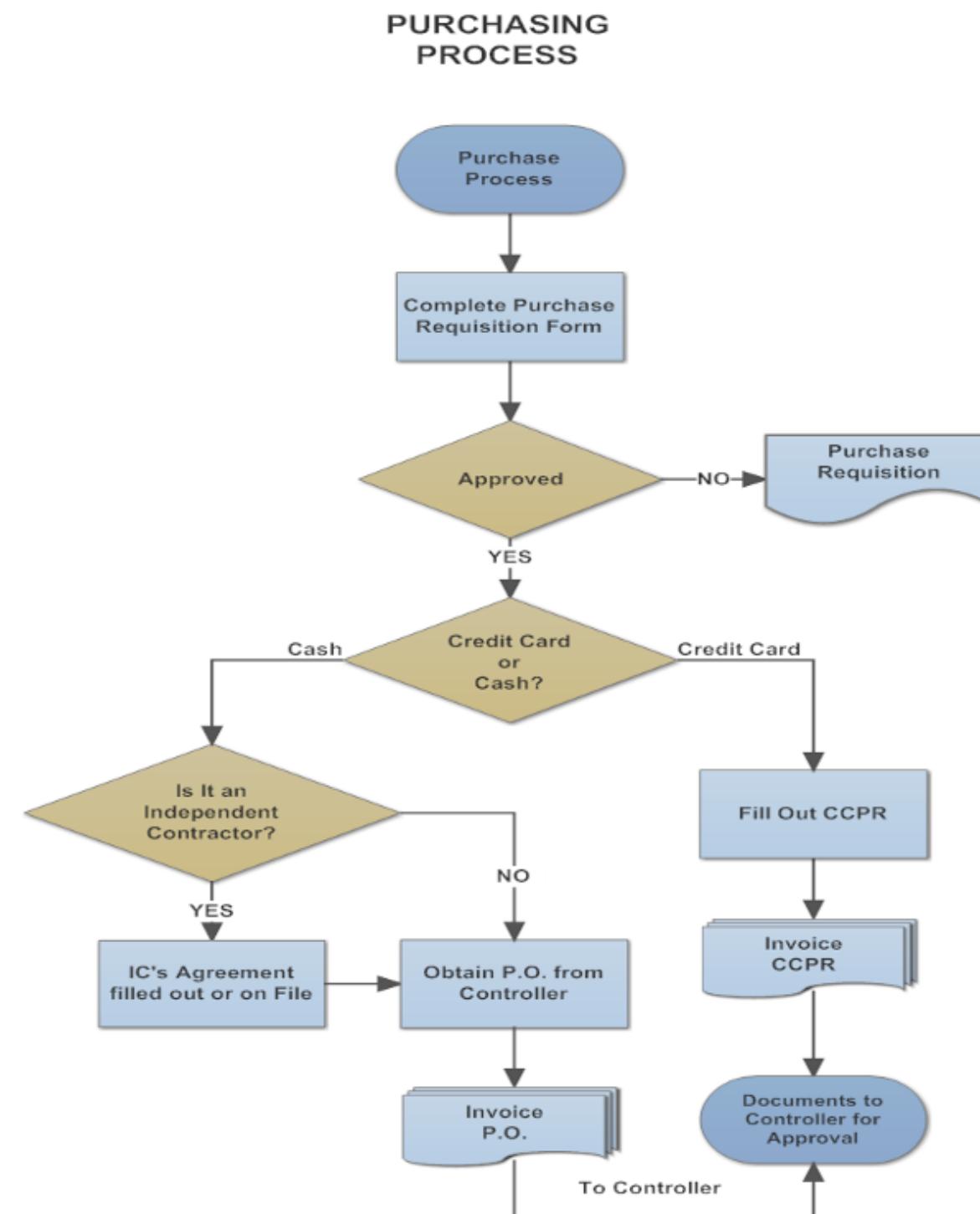
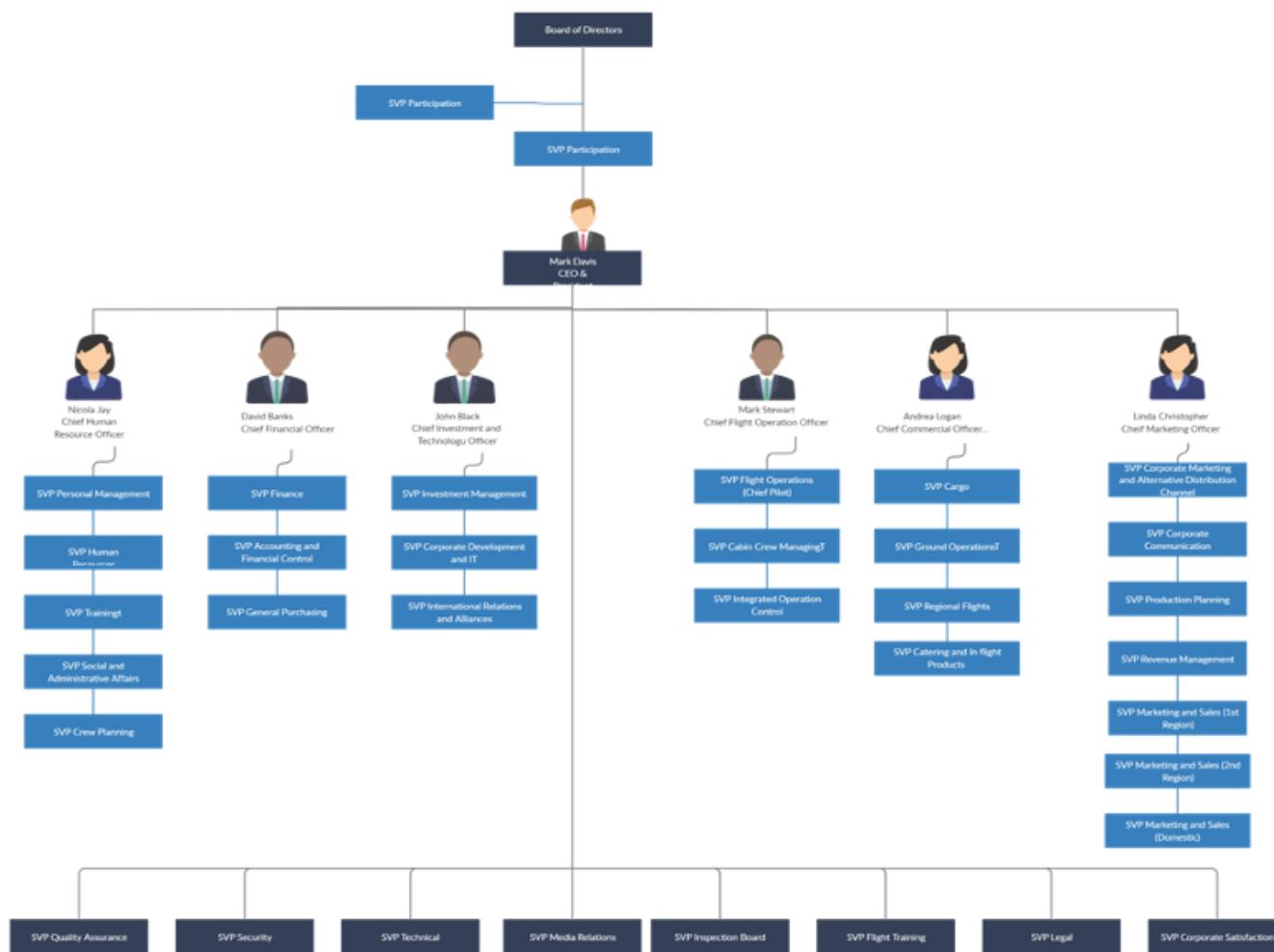
- Max call stack size (stack overflow error)
- Tail Call Optimization
- Memorization

# Recursive Call

Fundamental technique to solve problem:

- Identifying the base case
- Identifying the recursion formula/equation to transform the problem to smaller version
  - Problem requires back-tracking
  - Problem has tree structure

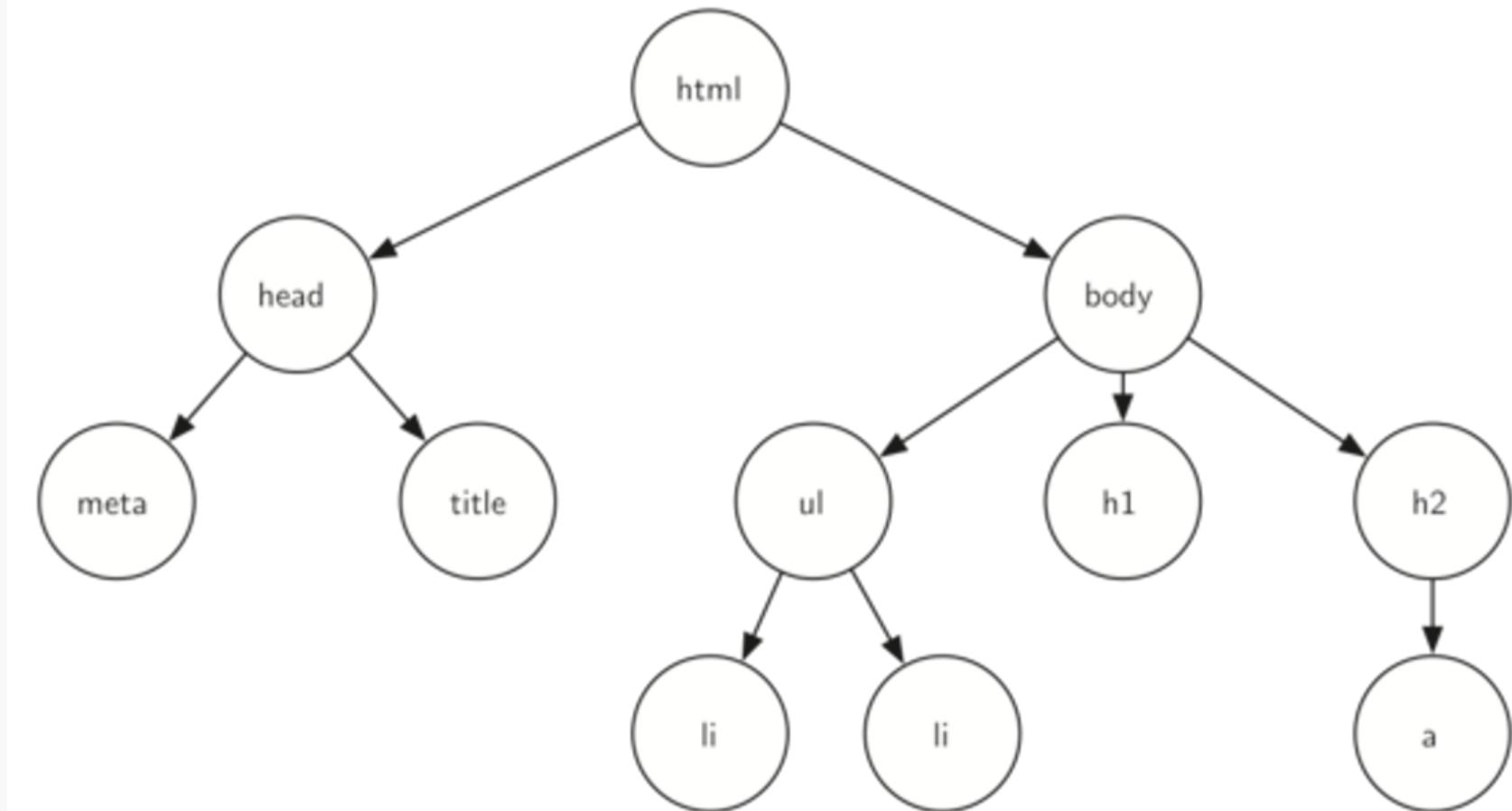
# Tree



# Tree

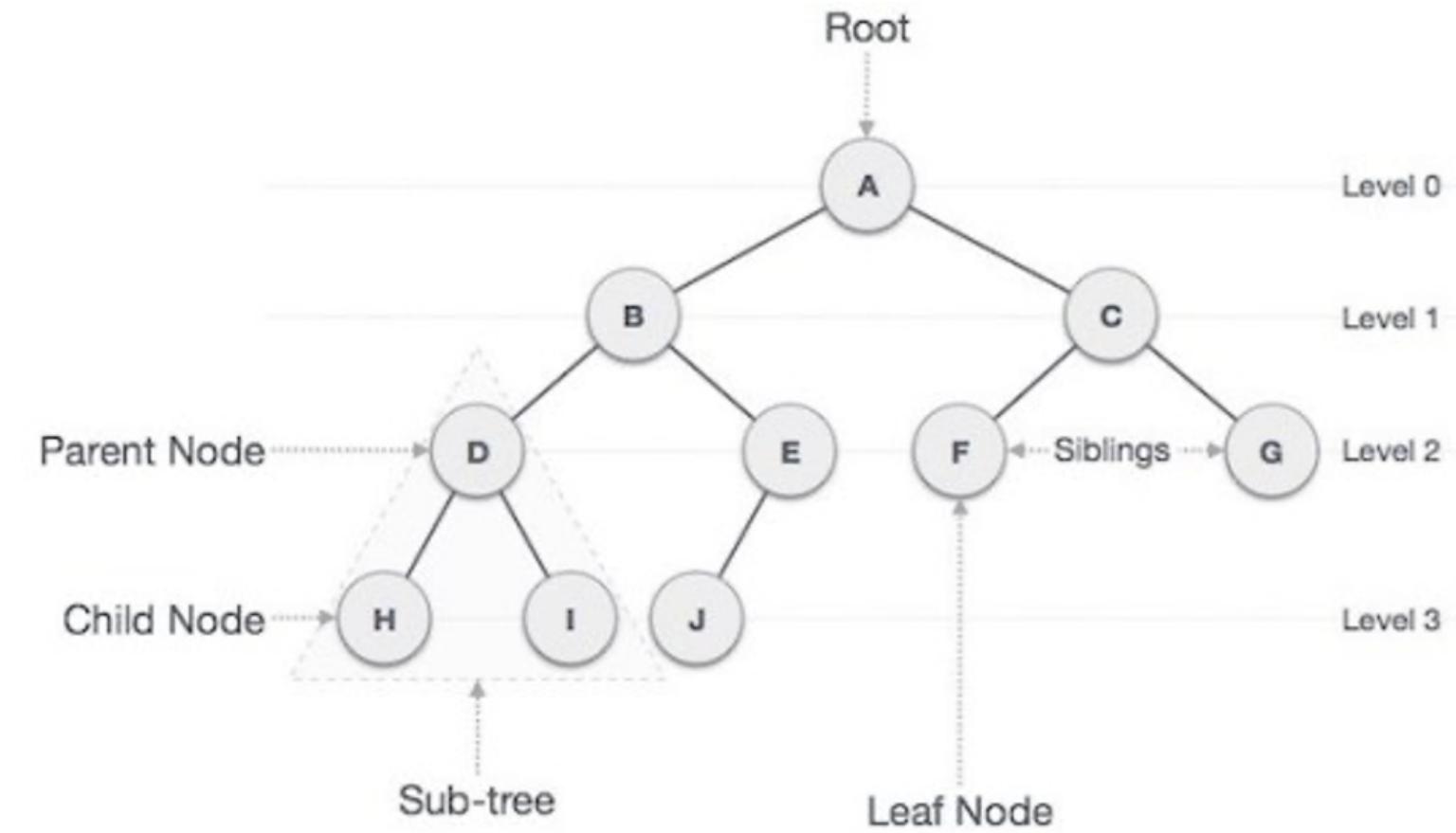
HTML DOM tree

```
1 <html xmlns="http://www.w3.org/1999/xhtml"
2     xml:lang="en" lang="en">
3 <head>
4     <meta http-equiv="Content-Type"
5         content="text/html; charset=utf-8" />
6     <title>simple</title>
7 </head>
8 <body>
9     <h1>A simple web page</h1>
10    <ul>
11        <li>List item one</li>
12        <li>List item two</li>
13    </ul>
14    <h2><a href="http://www.suss.edu.sg">SUSS</a></h2>
15 </body>
16 </html>
```



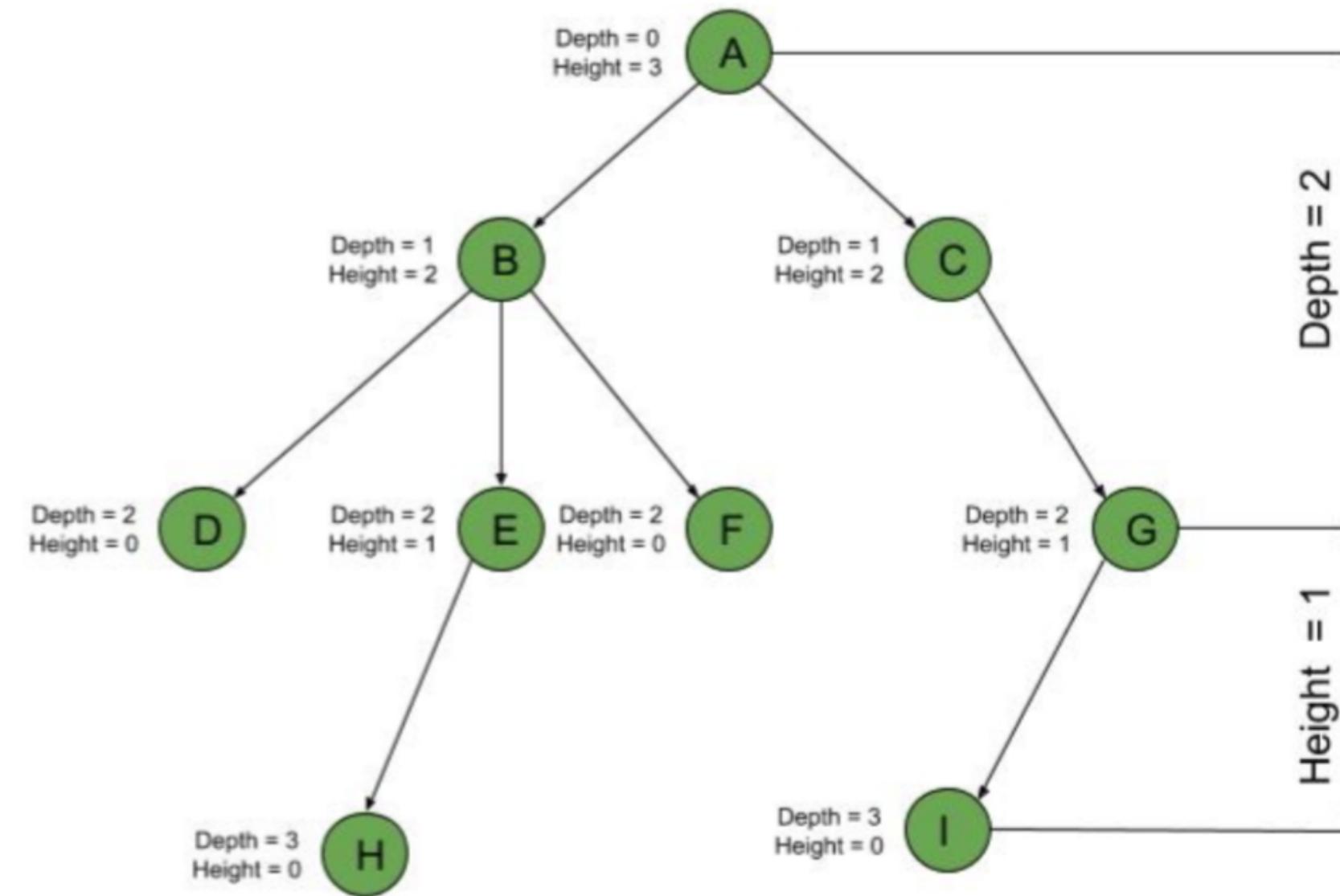
# Tree Terminology

- Node: Root, Leaf, Internal Node
- Node: Parent, Children, Sibling
- Edge
- Degree: no. of outgoing edges or Children
- SubTree
- Path: A → B → D → H
- Level: no. edges in path from root to the node
- Depth: no. edges in path from the node to the root
- Height: no. edges in longest path from the node to the leaf



# Tree Terminology

- **Level**
- **Depth**
- **Height**

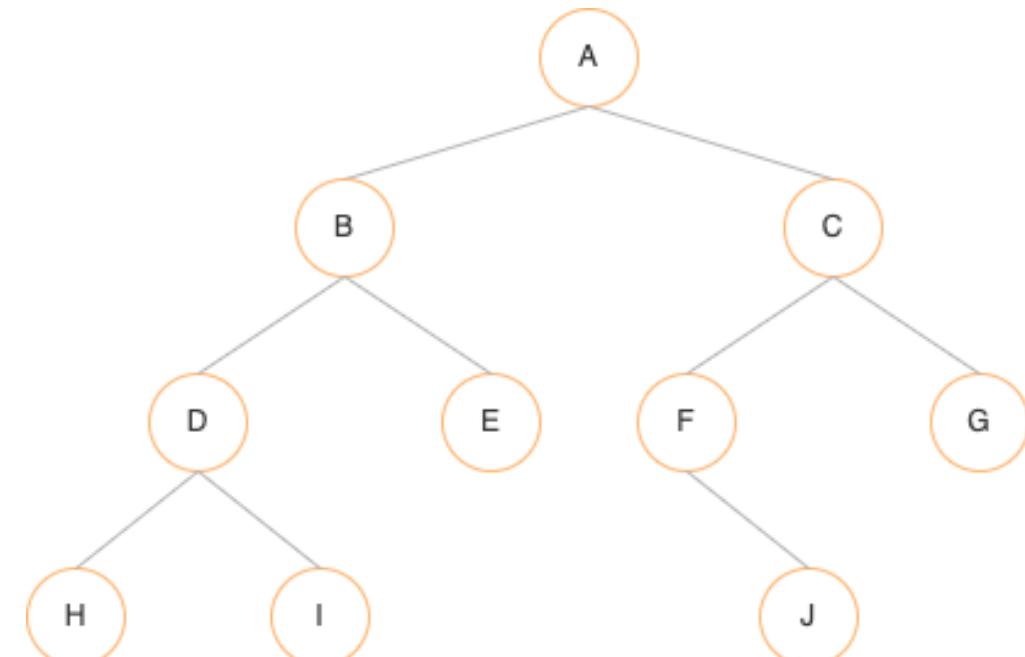


# Binary Tree

- One root
- Max 2 child nodes
- One and only one path from root to each node (path)
- Max nodes on level:  $2^l$
- Max nodes total:  $2^{h+1} - 1$

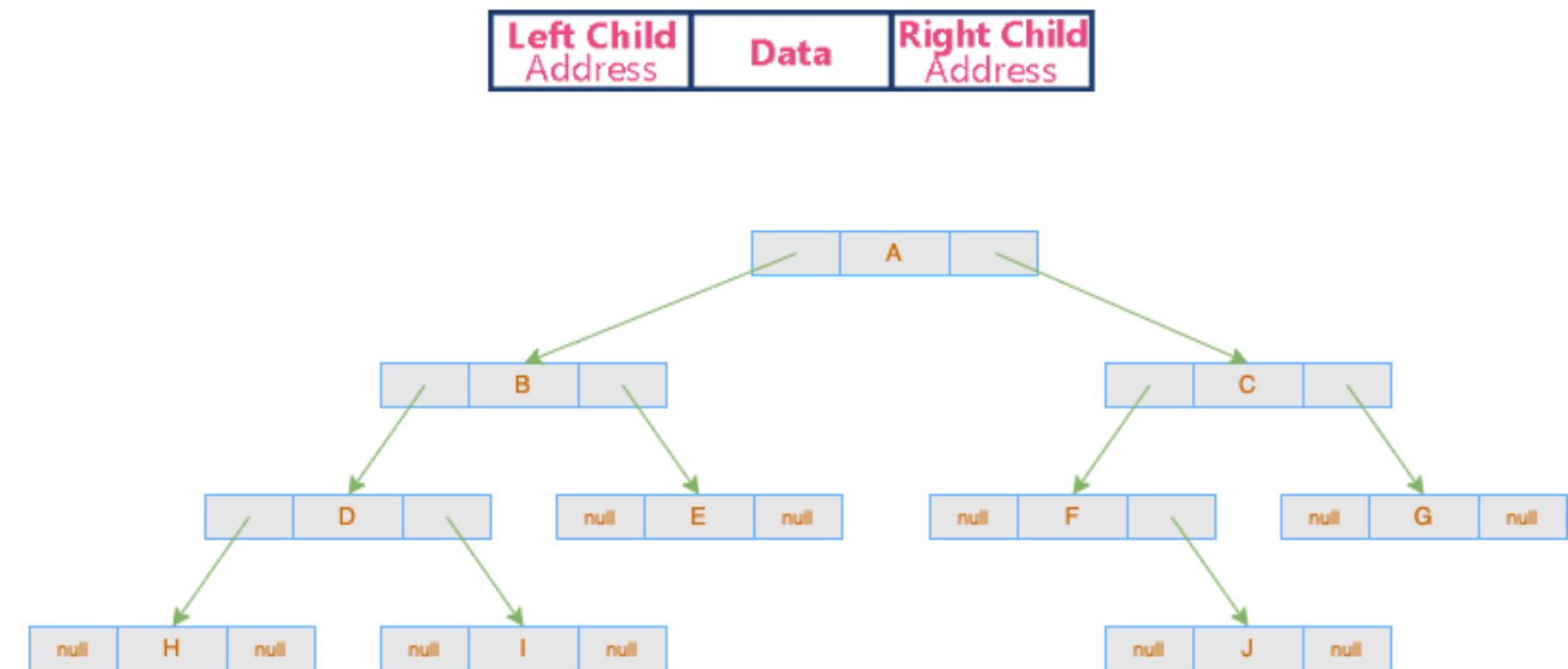
# Binary Tree

Array



['A', 'B', 'C', 'D', 'E', 'F', 'G', 'H', 'I', '#', '#', '#', '#', '#', 'J']

Left/Right Linked List



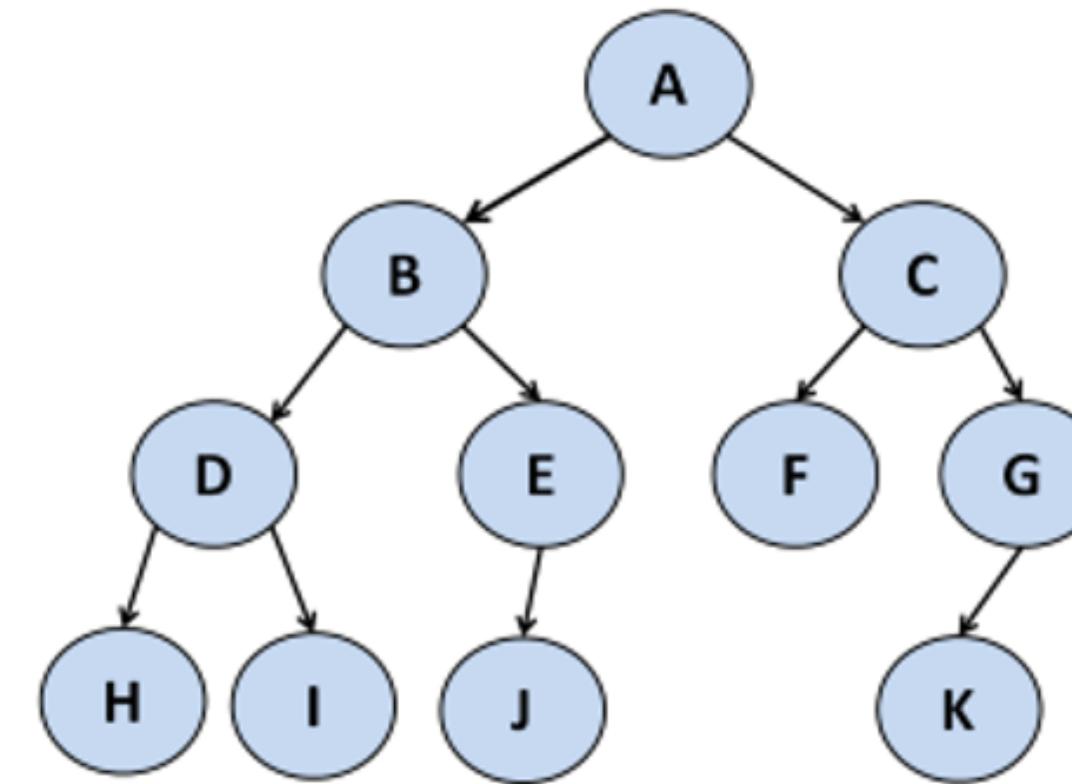
# Binary Tree Traverse (DFS): pre-order



**ROOT → Left → Right**

1. Visit the root
2. Traverse the left subtree
3. Traverse the right subtree

**A B D H I E J C F G K**



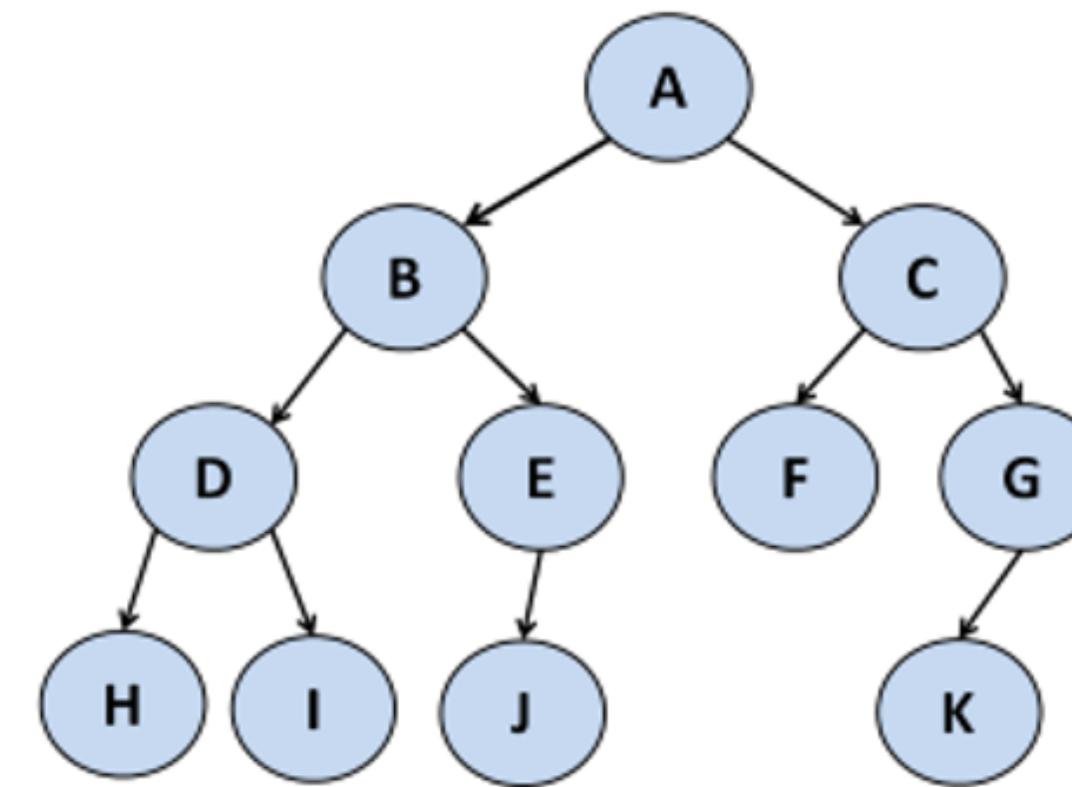
# Binary Tree Traverse (DFS): in-order



**Left → Root → Right**

1. Traverse the left subtree
2. Visit the root
3. Traverse the right subtree

**H D I B J E A F C K G**



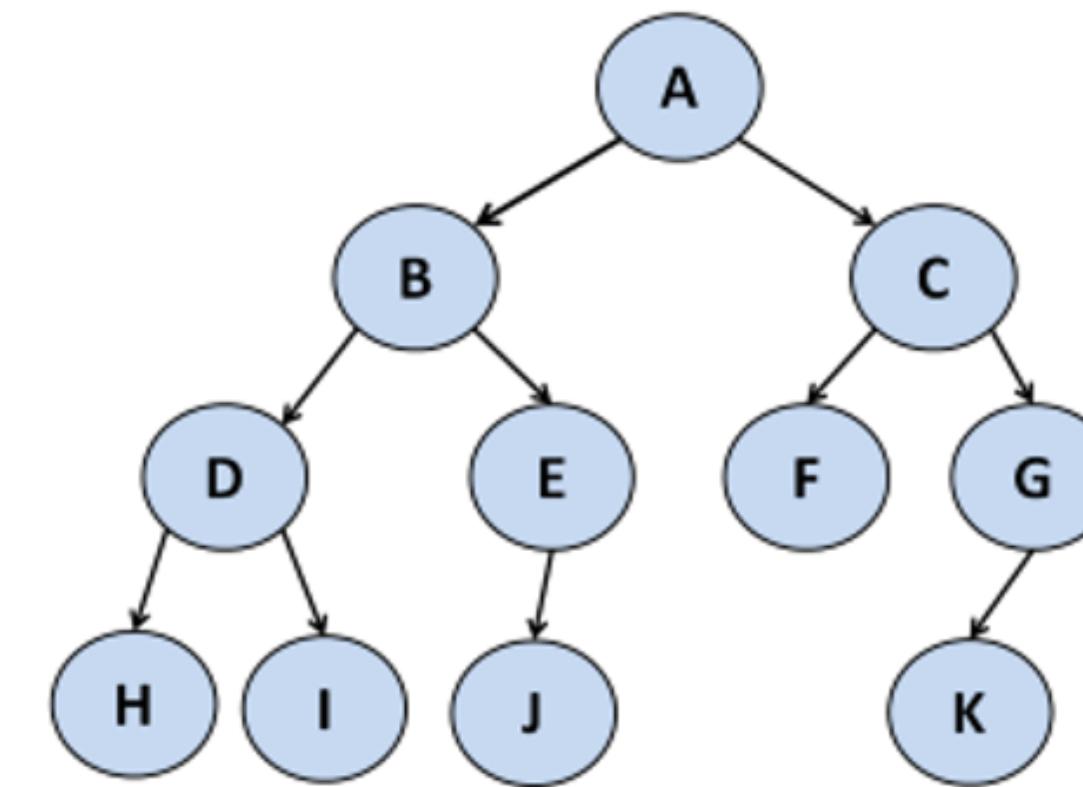
# Binary Tree Traverse (DFS): post-order



**Left → Right → Root**

1. Traverse the left subtree
2. Traverse the right subtree
3. Visit the root

**H I D J E B F K G C A**

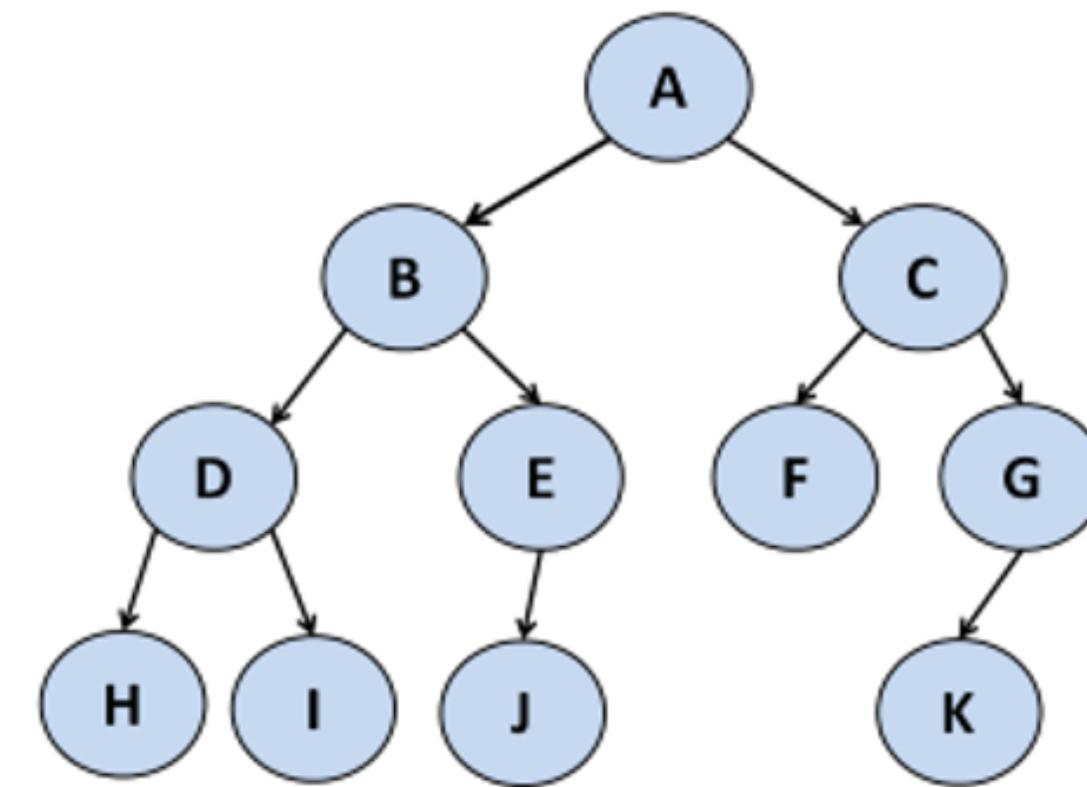


# Binary Tree Traverse (BFS): level order



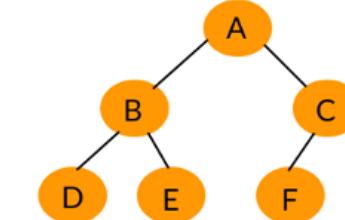
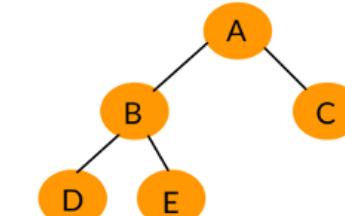
1. Visit the root
2. Visit the left node
3. Visit the right node
4. Go to next level

**A B C D E F G H I J K**

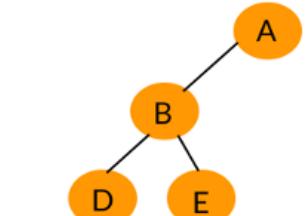


# Binary Tree

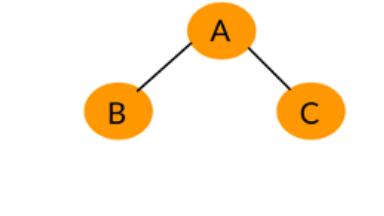
- **Complete Tree** : every level is completely filled except the last (leaf) and all nodes are as far left as possible
- **Full Binary Tree** : every node has two child nodes except leaf
- **Perfect Binary Tree** : every node has two child nodes except leaf and all leaves on same level



Complete Binary Tree

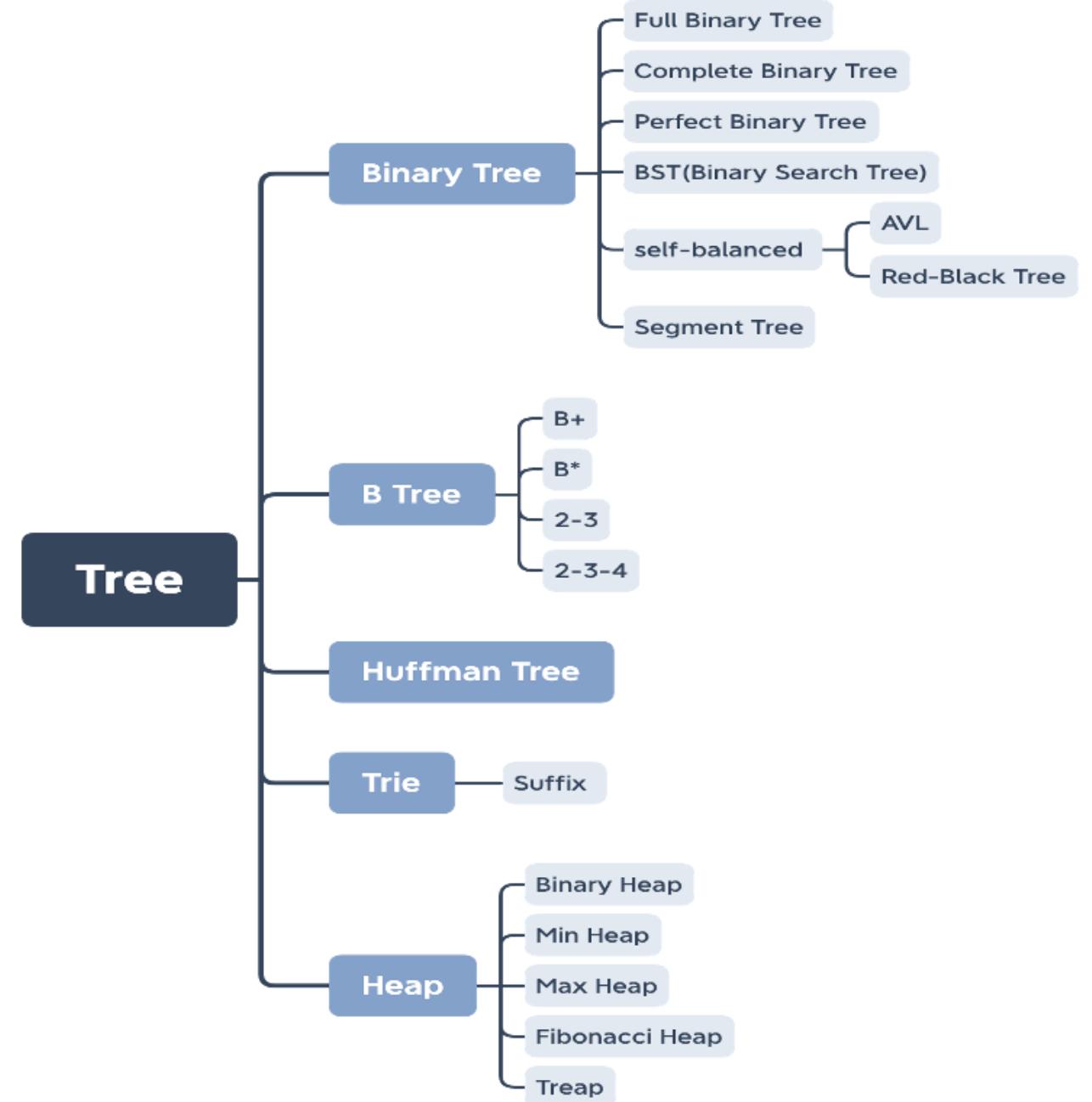


Full Binary Tree



Perfect Binary Tree

# Trees



# Summary

# Array

Concept	<ul style="list-style-type: none"><li>consecutive memory space: <math>\text{arr}[i].\text{address} = \text{base\_address} + i * \text{data\_type\_size}</math></li><li>same data type → same size for each element</li><li>fixed length</li></ul>
Complexity	<b>Accessing O(1)</b> , Searching O(n)
Notes	<ul style="list-style-type: none"><li>not memory friendly</li><li>cpu cacheable</li><li>index from 0</li><li>fundamental data structure to implement others such as stack, queue, heap</li><li>data type (programming language) vs. data structure</li></ul>
Hands-on	dynamic array, stack/queue, binary search, etc.

# Stack

Concept	<b>LIFO/FILO</b>
Complexity	Accessing O(n), Searching O(n), Inserting/push O(1), Deleting/pop O(1)
Notes	Stack implementation by dynamic array or linked list
Hands-on	function call stack, expression matching, etc.

# Queue

Concept	<b>FIFO/LIFO</b>
Complexity	Accessing O(n), Searching O(n), Inserting/enqueue O(1), Deleting/dequeue O(1)
Notes	Queue implementation by dynamic array or linked list
Hands-on	priority queue, circular queue, job queue, resource pool, etc.

# Linked List

Concept	<ul style="list-style-type: none"><li>nonconsecutive memory space</li><li>node: data + pointer</li><li>Single Linked List, Doubly Linked List, Circular Linked List, Positional Linked List</li></ul>
Complexity	Accessing O(n), Searching O(n), <b>Inserting O(1), Deleting O(1)</b>
Notes	<ul style="list-style-type: none"><li>accessing slower than array</li><li>with/without head/tail node (which don't store any data)</li><li>fundamental data structure to implement others such as skip list, hash table, etc.</li></ul>
Hands-on	stack, queue, traverse/reverse/update/merge, etc.

# Binary Tree

Concept	<ul style="list-style-type: none"><li>■ one root</li><li>■ max 2 child nodes</li><li>■ height &amp; depth</li><li>■ 4 traversal (DFS/BFS): in-order(left-root-right), pre-order(root-left-right), post-order(left-right-root), level-order</li><li>■ proper, perfect, full, complete binary tree</li></ul>
Complexity	<ul style="list-style-type: none"><li>■ DFS: time <math>O(n)</math>, space <math>O(h)</math></li><li>■ BFS: time <math>O(n)</math>, space <math>O(n)</math></li></ul>
Notes	stored in array or linked nodes
Hands-on	4 traversal

# Algorithm Design & Patterns

# Binary Search

- Input: array, target element
- Output: position (-1 if not existing)

As we know, whenever we are given a sorted Array or LinkedList or Matrix, and we are asked to find a certain element, the best algorithm we can use is the

- Shrink the search space every iteration (recursion)
- Cannot exclude potential answers during each shrinking

Although the basic idea of binary search is comparatively straightforward, the details can be surprisingly tricky

— Donald Knuth

# Binary Search

- Contains
- First occurrence of a key
- Last occurrence of a key
- Least element greater than
- Greatest element less than
- Closest element

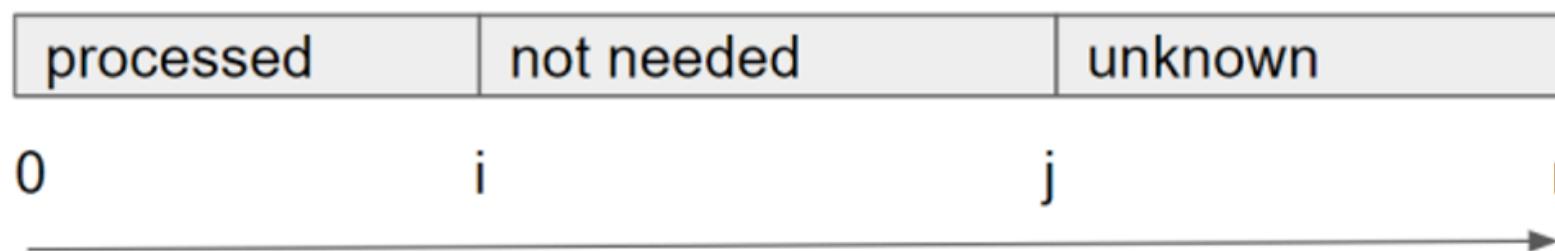
Further reading: Generalized Binary Search with predicates and main theorem:

<https://www.topcoder.com/community/competitive-programming/tutorials/binary-search>

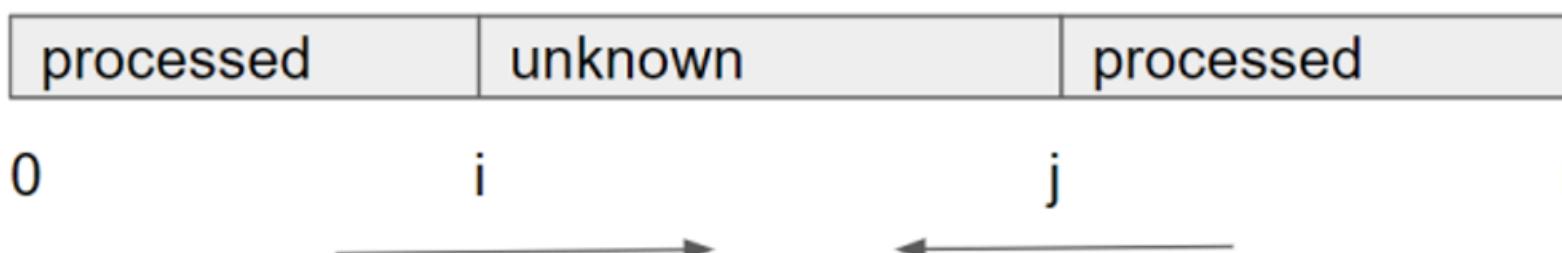
# Two Pointers

In problems where we deal with sorted arrays (or LinkedLists) and need to find a set of elements that fulfill certain constraints, the **Two Pointers** approach becomes quite useful. The set of elements could be a pair, a triplet or even a subarray.

- same direction

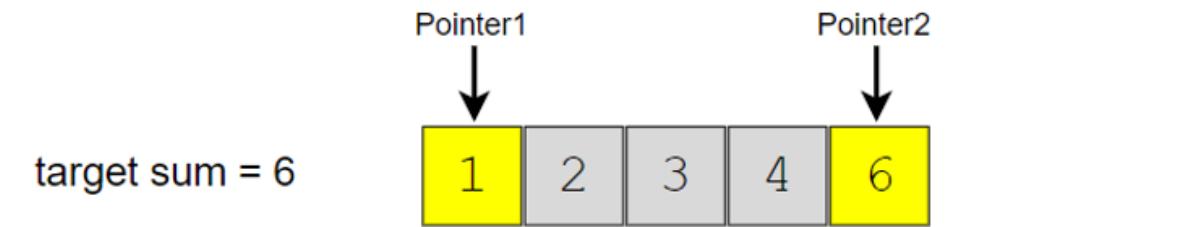


- reverse direction

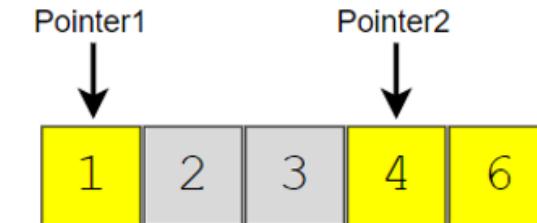


# two sum

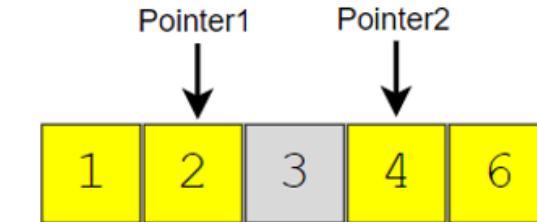
Given an array of sorted numbers and a target sum, find a pair in the array whose sum is equal to the given target.



$1 + 6 >$  target sum, therefore let's decrement Pointer2



$1 + 4 <$  target sum, therefore let's increment Pointer1

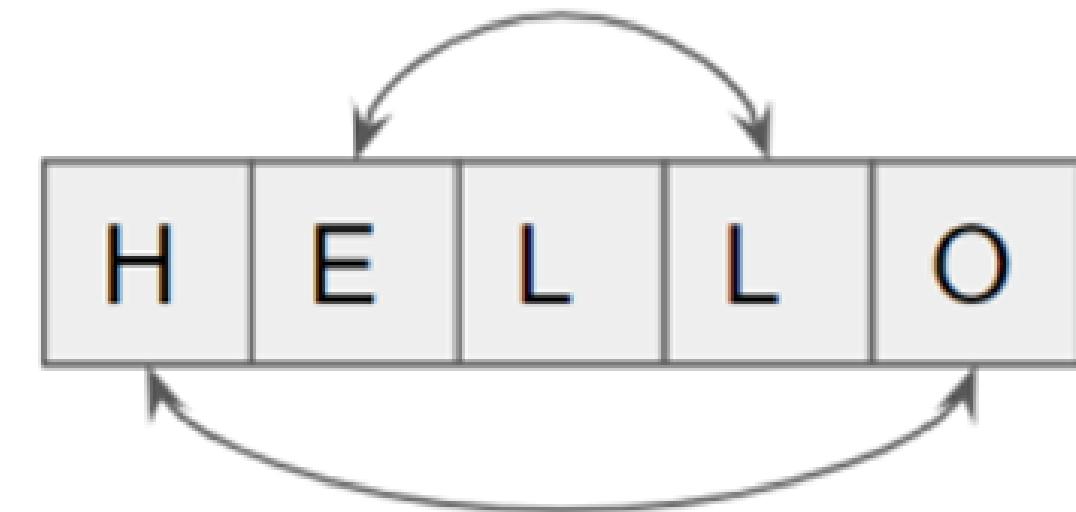


$2 + 4 ==$  target sum, we have found our pair!

# reverse a string

Write a function that reverses a string. The input string is given as an array of characters `char[]`. Do not allocate extra space for another array, you must do this by modifying the input array in-place with  $O(1)$  extra memory.

You may assume all the characters consist of printable ascii characters.



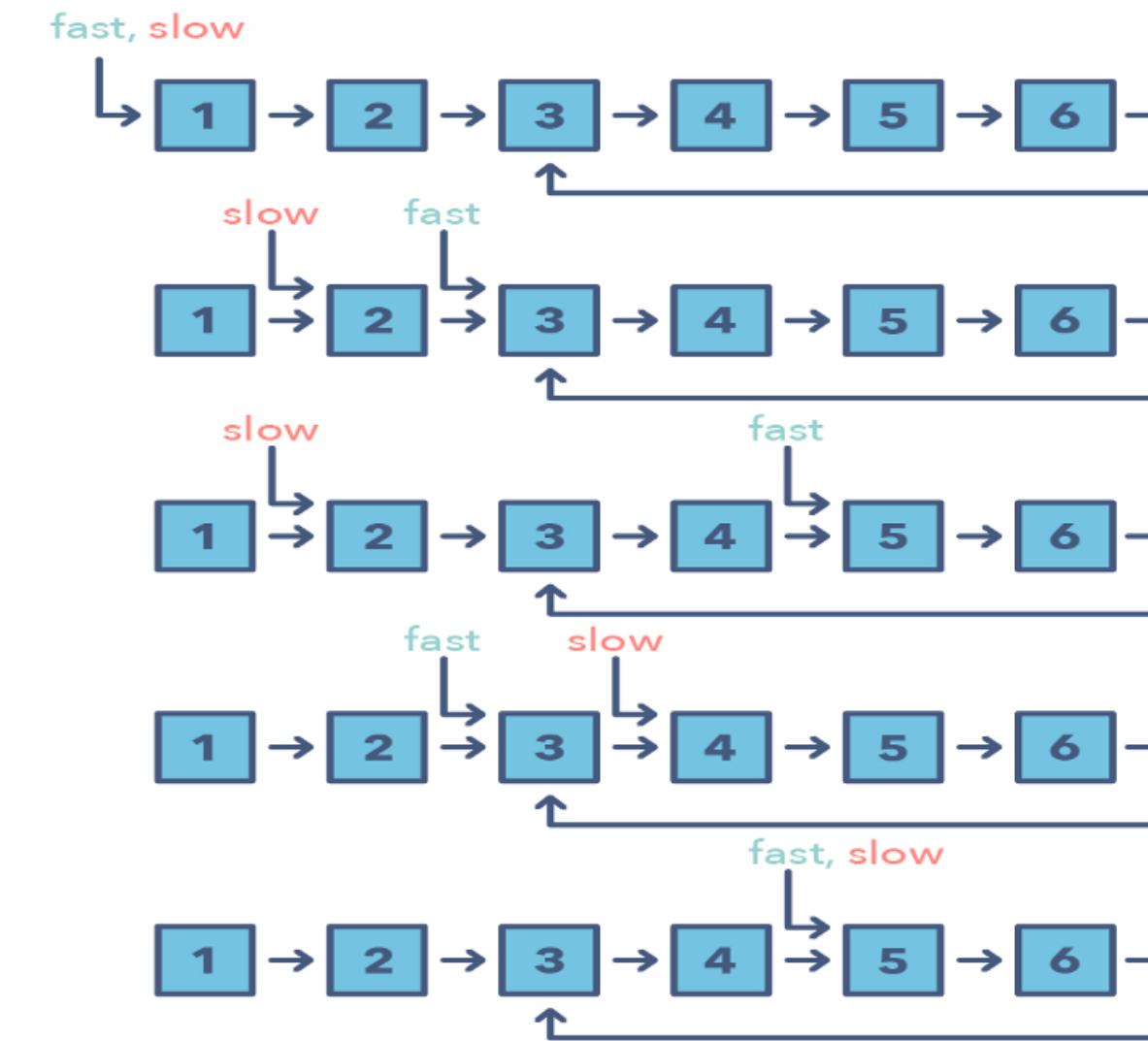
# Fast & Slow Pointers

The **Fast & Slow pointer** approach, also known as the Hare & Tortoise algorithm, is a pointer algorithm that uses two pointers which move through the array (or sequence/LinkedList) at different speeds.

- Fast & Slow Pointers, distance btw them
- Speed

# reverse a string

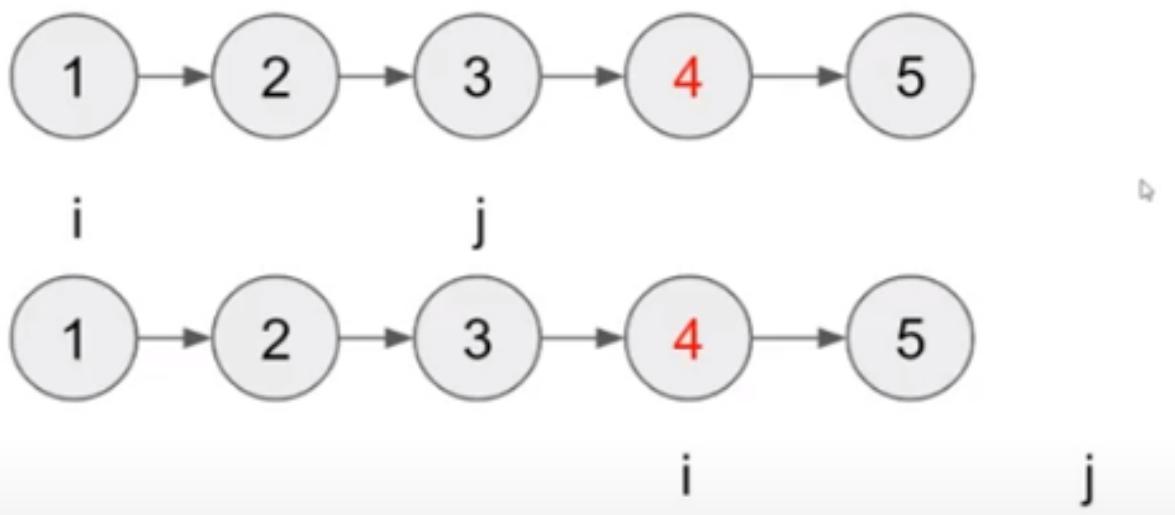
Given the head of a Singly LinkedList, write a function to determine if the LinkedList has a cycle in it or not.



# $k$ 'th node from the end of linked list

Given the head of a Singly LinkedList, Find  $k$ 'th node from the end of a linked list.

Example:  $k = 2$



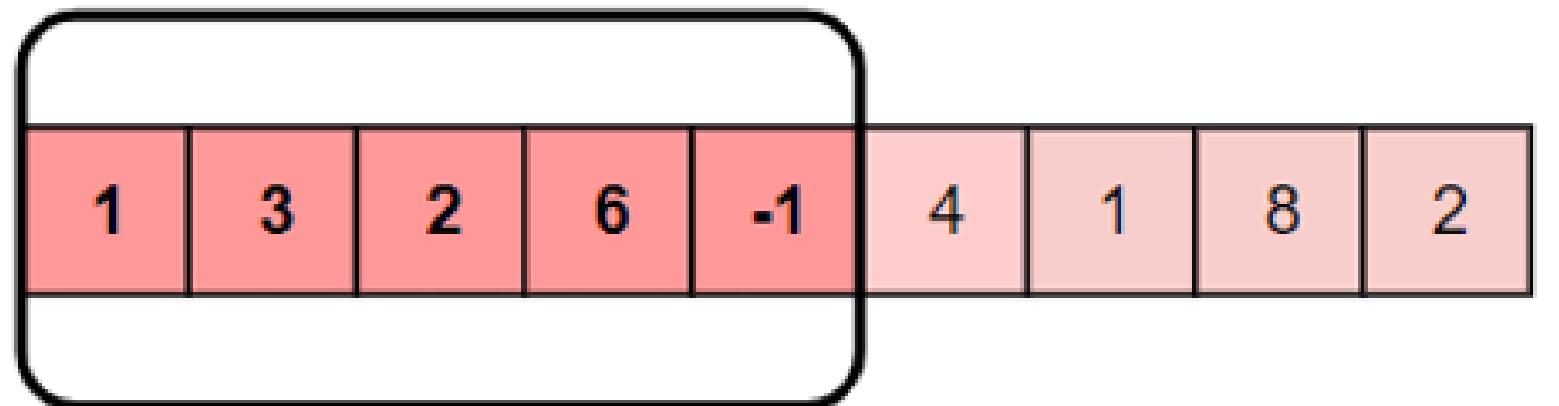
# Sliding Window

In many problems dealing with an array (or a LinkedList), we are asked to find or calculate something among all the contiguous subarrays (or sublists) of a given size.

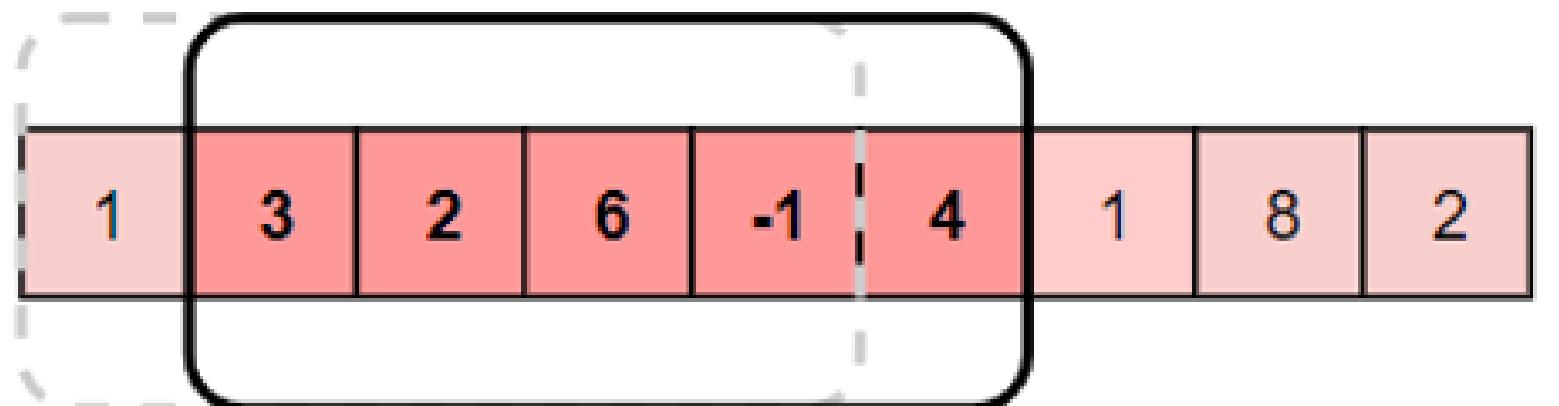
# longest substring

Given an array, find the average of all contiguous subarrays of size 'K' in it.

Sliding window -->



Slide one element forward



# longest substring

Given a string, find the length of the longest substring which has no repeating characters.

Example1 Input: String="aabccbb" Output: 3

Explanation: The longest substring without any repeating characters is "abc".

Example2 Input: String="abbbb" Output: 2

Explanation: The longest substring without any repeating characters is "ab".

Example3 Input: String="abccde" Output: 3

Explanation: Longest substrings without any repeating characters are "abc" & "cde".

# SUSS

SINGAPORE UNIVERSITY  
OF SOCIAL SCIENCES

# LAB



# Lab 1



- download and install Anaconda
- create and Activate your Anaconda Python env
- (Optional) install and setup VS Code
- familiar yourself with Python and do exercise [lab1.ipynb](#)

# Lab 1

## Download and install Anaconda

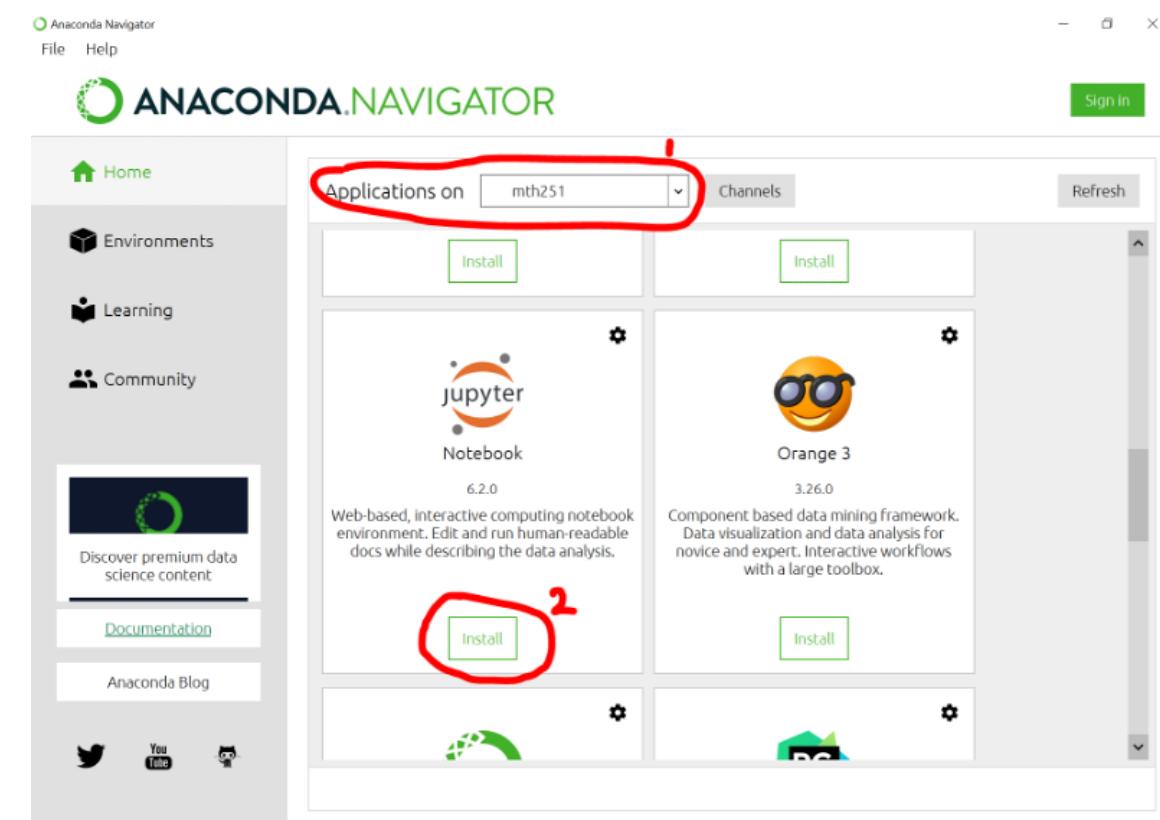
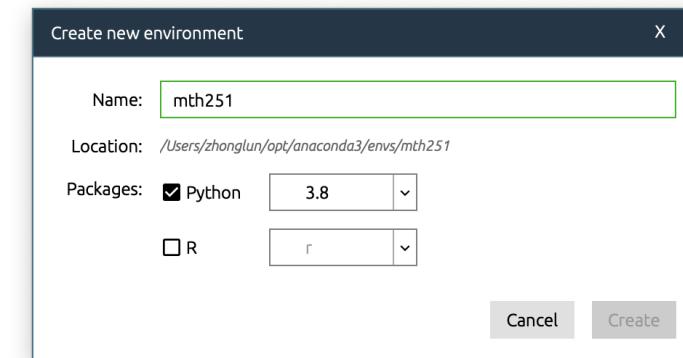
<https://www.anaconda.com/products/individual>

The screenshot shows the Anaconda Individual Edition landing page. At the top, there's a navigation bar with links for Products, Pricing, Solutions, Resources, Blog, and Company. A prominent 'Get Started' button is located in the top right. Below the navigation, the 'Individual Edition' section features a large green 'Q' icon and the text 'Your data science toolkit'. A descriptive paragraph explains that the Individual Edition is the easiest way to perform Python/R data science and machine learning on a single machine, developed for solo practitioners. It highlights over 20 million users worldwide and thousands of open-source packages and libraries. A 'Download' button is visible. At the bottom, there are three links: 'Open Source' (with a Python logo), 'Conda Packages' (with a package icon), and 'Manage Environments' (with a server icon).

# Lab 1

## Create and Activate your Anaconda Python env

1. Launch Conda Navigator
2. Environments → +Create to create a new env (**mth251**)
3. Switch to your Python env **mth251** from “Application on”
4. Install Jupyter Notebook
5. After that, Launch Jupyter Notebook



# Lab 1

## Create and Activate your Anaconda Python env

💬 Also possible to perform via command line:

```
1 > # create the env
2 > conda create -n mth251 python=3.8
3 > # activate the env
4 > conda activate mth251
5 > # install jupyter
6 > conda install -c conda-forge notebook
7 > # multipledispatch for lab1
8 > conda install -c anaconda multipledispatch
9 > # start jupyter notebook
10 > jupyter notebook
```

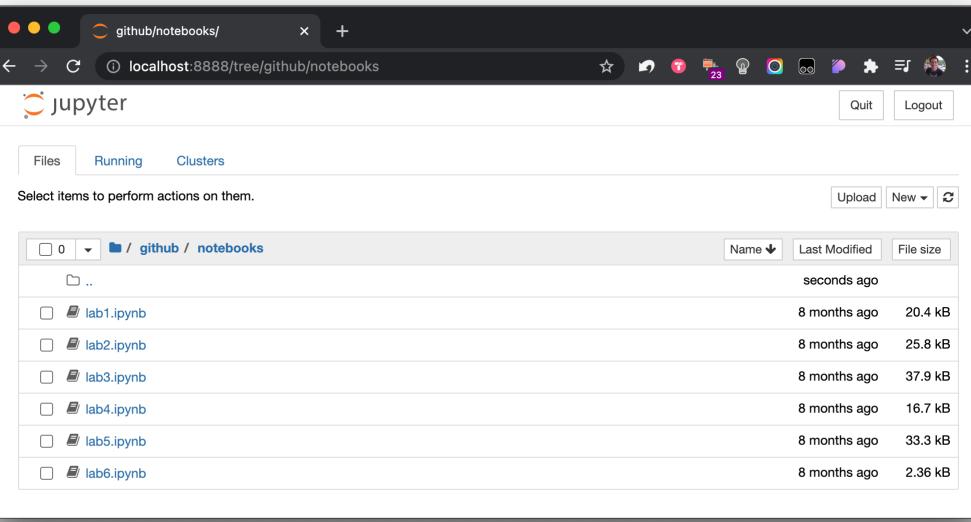
- copy & paste the Jupyter link in the prompt to your browser
- save notebooks to your local
- Control-C to stop Jupyter from the command line

# Lab 1

## Create and Activate your Anaconda Python env

6. Now you are ready to create, edit and run

Jupyter notebooks (lab1.ipynb):



# Lab 1



VS Code now fully integrated with Jupyter notebook, refer to this link:

[Jupyter Notebooks in VS Code](#)



Google provides online Jupyter env:

<https://colab.research.google.com/>

notebooks: <https://github.com/fastzhong/mth251/tree/main/public/notebooks>

colab: <https://colab.research.google.com/github/fastzhong/mth251/blob/main/public/notebooks/lab1.ipynb>

# Lab 1

## **lab1.ipynb**

- Python
  - 1. Data Type & Operators
  - 2. Collections
  - 3. Program Structure
  - 4. OO & Class
- Big O

# Lab 2



- review Array, Stack, Queue, Recursion
- exercise [lab2.ipynb](#)
- priority queue
- circular queue

# Lab 2

## Exercise: two sum

Given an array of integers `nums` and an integer `target`, return indices of the two numbers such that they add up to `target`.

You may assume that each input would have exactly one solution, and you may not use the same element twice.

### Example 1

Input: `nums = [2,7,11,15]`, `target = 9`

Output: `[0,1]`

Because `nums[0] + nums[1] == 9`, we return `[0, 1]`

### Example 2

Input: `nums = [3,2,4]`, `target = 6`

Output: `[1,2]`

### Example 3

Input: `nums = [3,3]`, `target = 6`

Output: `[0,1]`

# Lab 2

## Exercise: valid parentheses

Given a string  $s$  containing just the characters  $'($ ,  $)$ ,  $'\{$ ,  $\}$ ,  $'[$  and  $']$ , determine if the input string is valid.

An input string is valid if:

- Open brackets must be closed by the same type of brackets.
- Open brackets must be closed in the correct order.

Example 1

Input:  $s = "()\[\{\}]"$  Output: true

Example 2

Input:  $s = "(\[)"$  Output: false

Example 3

Input:  $s = "(\[\])"$  Output: false

Example 4

Input:  $s = "\{\[\]\}"$  Output: true

# Lab 2

**Priority Queue** is similar to queue but the element with higher priority can be moved forward to the front.  
Use existing Queue class to implement a priority queue (element with lower value has higher priority).

- Priority Queue can be used in Printer Jobs or Schedule Tasks.

# Lab 2

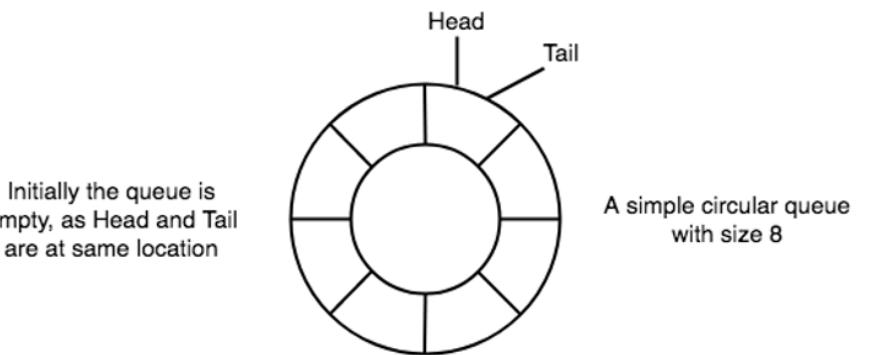
**Circular Queue** is a linear data structure in which the operations are performed based on FIFO (First In First Out) principle and the last position is connected back to the first position to make a circle. It is also called "Ring Buffer".

Design your implementation of circular queue.

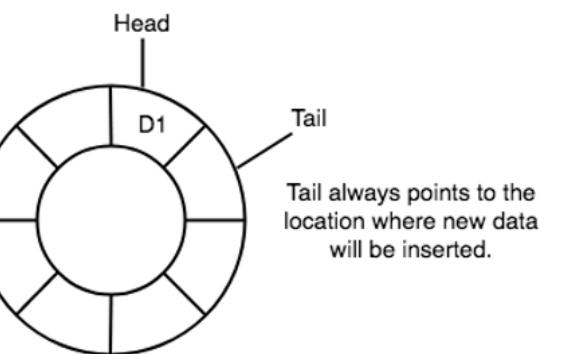
# Lab 2

## Circular Queue

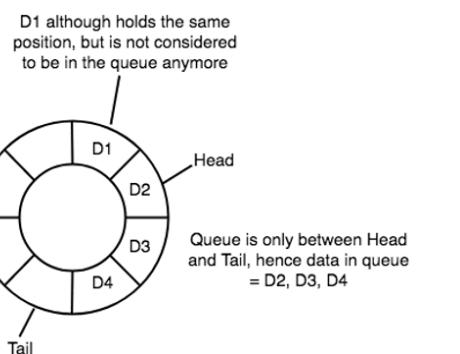
1. init



2. enqueue D1



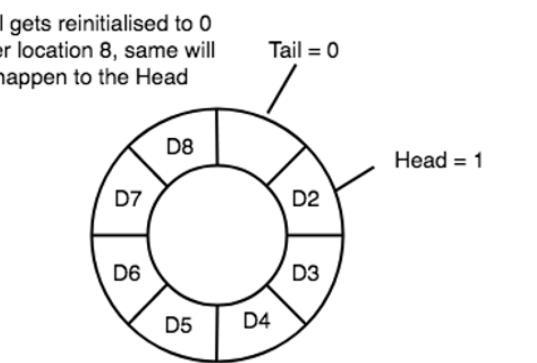
3. enqueue D2, D3, D4 and dequeue D1



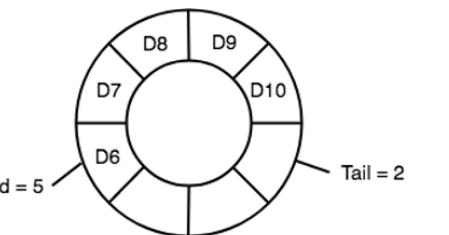
# Lab 2

## Circular Queue

4. enqueue D5, D6, D7, D8



5. dequeue D2, D3, D4, D5 and enqueue D9, D10



In such a situation the value of the Head pointer will be greater than the Tail pointer

# Lab 2

## Circular Queue

- ❑ One of the benefits of the circular queue is that we can make use of the spaces in front of the queue. In a normal queue, once the queue becomes full, we cannot insert the next element even if there is a space in front of the queue (and it does not prevent the program accidentally creates a large queue or stack and use up the memory).

### Implementation of CircularQueue class:

- **enqueue()**: insert the element
- **dequeue()**: delete the element
- **front()**: return the first element in the queue, if queue is empty, return None
- **rear()**: return the last element in the queue, if queue is empty, return None
- **is\_empty()**: return true if queue is empty
- **is\_full()**: return true if queue is full

# Lab 3



- Python (lab1)
  - 5. misc.
  - 6. PEP8
- review Linked List, Doubly Linked List
- exercise [lab3.ipynb](#)

# Lab 3

- implement Stack by linked list
- implement Queue by linked list
- reverse a linked list
  - recursive implementation
  - iterative implementation

# Lab 4



- review Binary Tree and 4 traverse methods
- exercise [lab4.ipynb](#)

# Lab 4

Exercise: convert binary tree from linked list to array

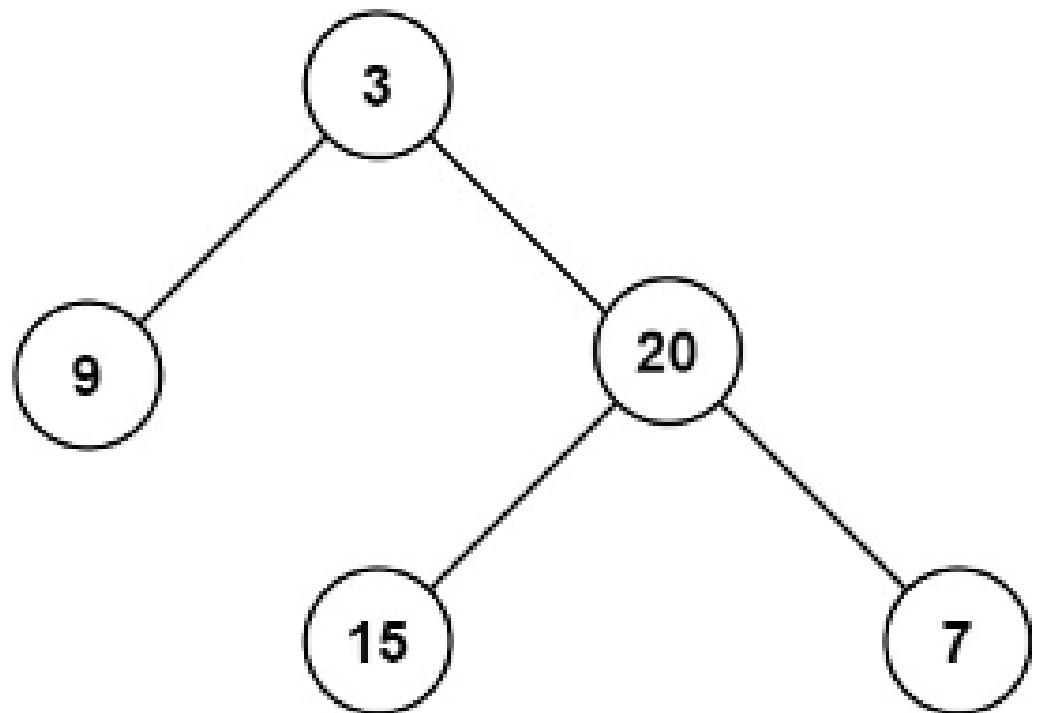
Exercise: convert binary tree from array to linked list

# Lab 4

## Exercise: get maximum depth of binary tree

A binary tree's maximum depth is the number of nodes along the longest path from the root node down to the farthest leaf node.

*For example: the maximum depth is 3*

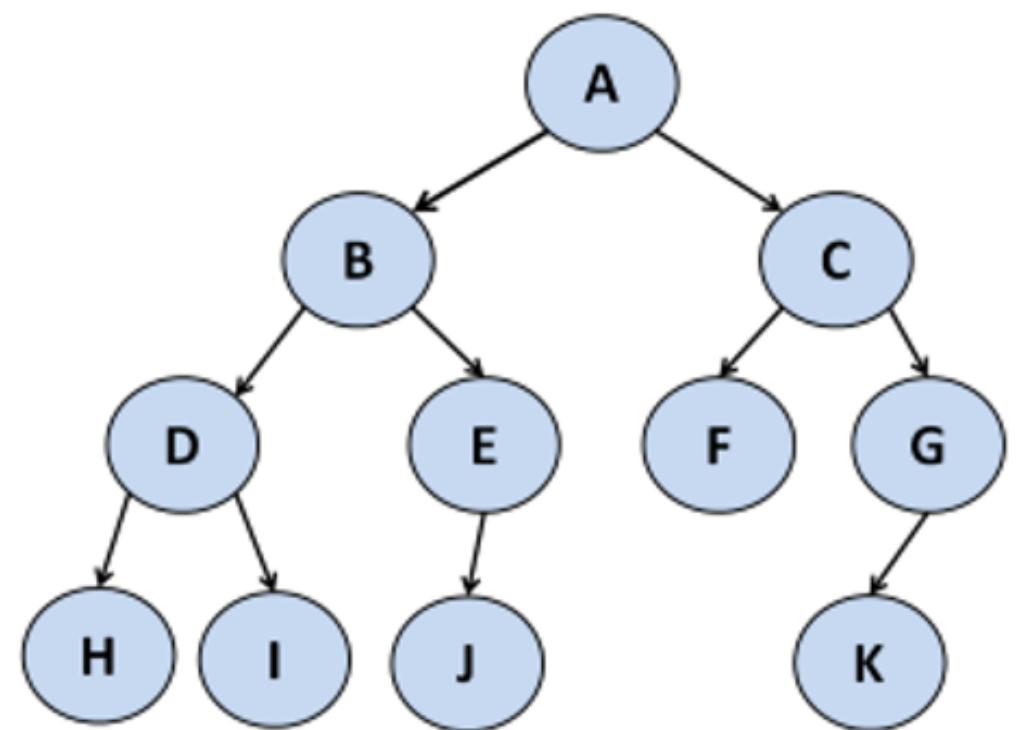


# Lab 4

## Exercise: get minimum depth of binary tree

A binary tree's minimum depth is the number of nodes along the shortest path from the root node down to the nearest leaf node.

*For example: the minimum depth is 3*

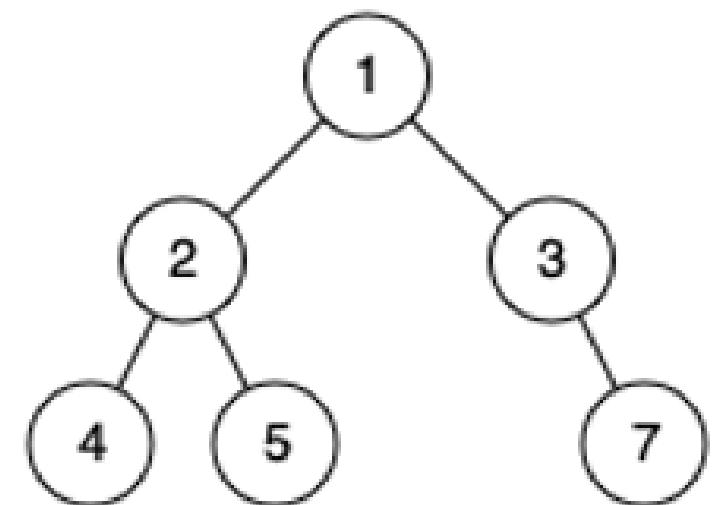


# Lab 4

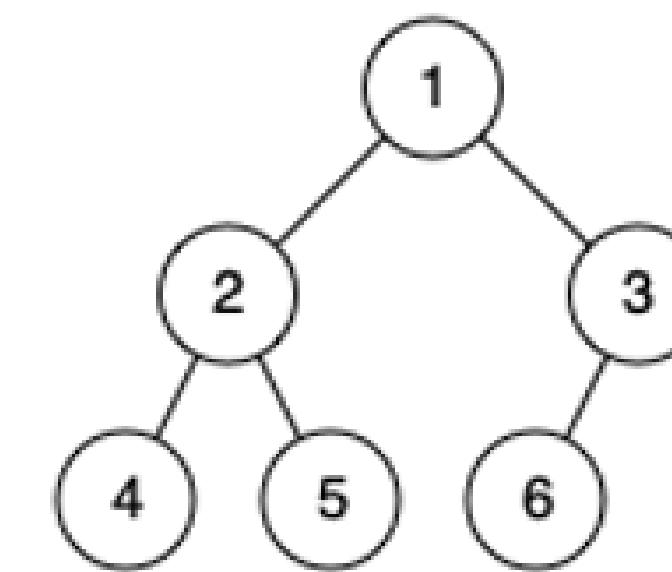
## Exercise: check a complete binary tree

In a complete binary tree, every level, except possibly the last, is completely filled, and all nodes in the last level are as far left as possible. It can have between 1 and  $2^h$  nodes inclusive at the last level  $h$ .

*complete = true:*



*complete = false:*



# Lab 5



- review linear search, binary search
- let us do some leetcode exercises  [lab5.ipynb](#)

# Lab 5

## lineary search & binary search

1. Go to <https://www.cs.usfca.edu/~galles/visualization/Search.html> to understand how Linear Search & Binary Search is working

2. Implement Linear Search & Binary Search in Python by yourself:

- familiar with Python coding style
- understand the input, output, steps and ending condition
- learn and compare different approaches (time & space complexity)
- test code reliability with different cases

# Lab 5

## Exercise: palindrome

Implement a Python function to determine if a string is a palindrome, for example, 'racecar' and 'level' are palindromes

# Lab 5

## Exercise: remove duplicate numbers

Given a sorted array `nums`, remove the duplicates in-place such that each element appears only once and returns the new length.

Do not allocate extra space for another array, you must do this by modifying the input array in-place with O(1) extra memory.

### Example 1

Input: `nums = [1,1,2]`

Output: `2, nums = [1,2]`

Explanation: Your function should return length = 2, with the first two elements of `nums` being 1 and 2 respectively. It doesn't matter what you leave beyond the returned length.

### Example 2

Input: `nums = [0,0,1,1,1,2,2,3,3,4]`

Output: `5, nums = [0,1,2,3,4]`

Explanation: Your function should return length = 5, with the first five elements of `nums` being modified to 0, 1, 2, 3, and 4 respectively. It doesn't matter what values are set beyond the returned length.

# Lab 6



-  lab6.ipynb
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