

6/8/2019

# EXPERIMENT - I

## CPU SCHEDULING

(1)

1) Aim:-

Simulate the following non-preemptive CPU scheduling algorithms to find the turn-around time and waiting time.

- a) FCFS
- b) SJF
- c) RR (preemptive)
- d) Priority

2) ALGORITHM:-

a) FCFS:-

Step 1 :- Input the processes along with their burst times (bt)

Step 2 :- Find the waiting time for all processes (wt)

Step 3 :- First process that comes need not wait so waiting time for the first process is zero. ( $wt[0] = 0$ )

Step 4 :- Find the waiting time for all other processes

For process  $i$ ,

$$wt[i] = bt[i-1] + wt[i-1]$$

Step 5 :- Find turnaround time

~~turnaround = waiting + burst~~

## OUTPUT:-

FCFS:-

Enter the number of processes : 5

Enter the burst and arrival time : 50

16 0  
9 0  
4 0  
3 0

SI	Bst	Arr	Turn	Wait
1	5	0	5ms	0ms
2	16	0	21ms	5ms
3	9	0	30ms	21ms
4	4	0	34ms	30ms
5	3	0	37ms	34ms

SJF:-

Enter the number of processes : 4

Enter the burst and arrival time : 40

12 0  
5 0  
2 0

SI	Bst	Arr	Turn	Wait
4	2	0	2ms	0ms
1	4	0	6ms	2ms
3	5	0	11ms	6ms
2	12	0	23ms	11ms

Step 6:- Print the turnaround time, waiting time for each process in the order of arrival.

b) SJF :-

Step 1:- Input the process along with their burst times

Step 2:- Sort the processes in the increasing order of burst time.

Step 3:- Find the waiting time for all the processes

Step 3.1:- First process will have waiting time 0 i.e  $wt[0] = 0$

Step 3.2:- For all other process  $i$ ,

$$wt[i] = bt[i-1] + wt[i-1]$$

Step 4:- Find turnaround time  
turnaround = waiting + burst

Step 5:- Print the turnaround time, waiting time for each process in increasing order of burst time.

c) Round Robin :-

~~Step 1:- Input the processes along with their burst times.~~

RR:-

Enter the number of processes: 4

Enter burst time and arrival time: 4 0

12 0

5 0

2 0

Enter time slice: 2

SI	Bst	Arr	Wait	Turn
1	4	0	6ms	10ms
2	12	0	11ms	23ms
3	5	0	12ms	17ms
4	2	0	6ms	8ms

PRIORITY:-

Enter number of processes: 4

Enter burst and arrival time: 12 0

3 0

5 0

9 0

Enter priority: 1 5 3 4

SI	Bst	Arr	Prio	Turn	Wait
2	12	0	5	12ms	0ms
4	3	0	4	15ms	12ms
3	5	0	3	20ms	15ms
1	9	0	1	29ms	20ms

Step 2 :- Create an array rem-bt[] to track the remaining burst time which is a copy of the values of burst time.

Step 3 :- Create an array wt[] to store the waiting times of processes with initial value of 0.

Step 4 :- Initialize  $t=0$

Step 5 :- Traverse all the processes until they are finished & if  $\text{rem-bt}[i] > \text{quantum}$

$$t = t + \text{quantum}$$

$$\text{bt-rem}[i] = \text{quantum}$$

Else

$$t = t + \text{bt-rem}[i]$$

$$\text{wt}[i] = t - \text{bt}[i]$$

$$\text{bt-rem}[i] = 0$$

d) PRIORITY :-

Step 1 :- Input the process with burst times and priorities

Step 2 :- Sort the process (with burst times) according to the priority.

Step 3 :- Find the waiting times of all processes

~~Step 3 :- First process have waiting time 0, i.e.  $\text{wt}[0]=0$~~

(4)

Step 3.2 :- For all other process, say i  
 $wt[i] = bt[i-1] + wt[i-1]$

Step 4 :- Find the turnaround time

$$\text{turnaround} = \text{waiting} + \text{burst}$$

Step 5 :- Print the waiting time and  
turnaround time of all process according  
to priority.

### a) SINGLE-LEVEL

Enter name of directory - D1

- 1. Create File
- 2. Delete File
- 3. Search File
- 4. Display File
- 5. Exit

Enter choice - 1

Enter name - F1

Enter choice - 1

Enter name - F2

Enter choice - 4

The Files are - F1 F2

Enter choice - 2

Enter name - F1

File F1 is deleted

Enter choice - 4

The Files are - F2

Enter choice - 3

Enter name - F1

File F1 not Found

Enter choice - 5

(5)

## EXPERIMENT - 2

### FILE ORGANIZATION

1) AIM:-

Simulate the following file organization techniques

- a) Single-level directory
- b) Two level directory
- c) Hierarchical directory

2) ALGORITHM:-

a) Single-level Directory

It is the simplest of all directory structures. There is only one directory containing all the files.

Algorithm :-

Step 1 :- Input the number of files, say N

Step 2 :- While ( $N \neq 0$ )

    Step 2.1 :- Input filename.

    Step 2.2 :- Decrement N

Step 3 :- Display the file structure.

b) Two-level Directory

~~It contains a root directory, under which each user has a private directory.~~

- 6) 1) Create Directory 2) Create File 3) Delete File  
4) Search File 5) Display 6) Exit

Enter choice - 1

Enter name of Directory - D1

Directory created

Enter choice - 2

Enter name of Directory - D2

Directory created.

Enter choice - 5

Directory Files

D1

D2

Enter choice - 2

Enter directory name - D1

Enter name of file - F1

File created

Enter choice - 5

Directory Files

D1

F1

D2

Enter choice - 4

Enter name of directory - D1

Enter name of file - F1

File F1 is found.

Step 1:- Initialize the root directory

Step 2:- Input the number of users, and the number of files they have.

Step 3:- Initialize the user directory and add it to the tree structure

Step 4:- Add files of each user to each user directory.

c) Hierarchical.

It allows subdirectories to be created within user directories and itself.

Step 1:- Initialize the root directory

Step 2:- Input the number of users, the number of subdirectories/files they have.

Step 3:- Initialize each user directory and their subdirectories and files

Step 4:- Display as a tree structure.

Enter choice - 3

Enter name of directory - D1

Enter name of file - F1

File F1 is deleted

Enter choice - 5

Directory Files

D1

D2

Enter choice - 6

c) HIERARCHICAL

You are on root

- 1) list everything 2) change directory 3) go to parent directory 4) add new file 5) delete file 6) create new directory 7) delete directory 8) exit

Enter choice : 6

Enter name : D1

Enter choice : 4

Enter name : F1

Enter choice : 2

Enter name : D1

Directory changed

Enter choice : 3

Directory changed

Enter choice : 1

\* D1 \* F1

Enter choice : 5

Enter name of directory or file to delete : d1  
Successfully deleted

Enter choice : 1

Empty directory

Enter choice : 3

Enter choice : 7

Enter name of file or directory to delete : d1  
Successfully deleted

Enter choice : 1

F1

Enter choice : 8

Exit

(6)

## EXPERIMENT - 3

### BANKER'S ALGORITHM

**Aim:-**

Implement the bankers algorithm for deadlock.

**ALGORITHM:-**

a) Safety Algorithm:-

Step 1:- Let work and finish be arrays of length  $m$  and  $n$ , respectively  
initialize:

work = Available

Finish[i] = false for  $i = 0, 1, \dots, n-1$

Step 2:- Find  $i$  such that both

Finish[i] = false

Need[i]  $\leq$  work

If no such  $i$  exists, go to step 4

Step 3:- work = work + Allocation[i]

Finish[i] = true

Go to step 2

Step 4:- If Finish[i] = true for all  $i$ , then  
the system is in safe state

OUTPUT:-

Enter the number of processes and resources : 5

Enter the resource allocation table :

0 1 0  
2 0 0  
3 0 2  
2 1 1  
0 0 2

Enter the number of resources available of  
Type 1 : 10

Enter the number of resources available of  
Type 2 : 5

Enter the number of resources available of  
Type 3 : 7

Enter the maximum resource table

7 5 3  
3 2 2  
9 0 2  
2 2 2  
4 3 3

System is in safe state

Order is : P2 - P4 - P5 - P1 - P3

## b) Resource Request Algorithm

Step 1:- If  $\text{Request}[i] \leq \text{Need}[i]$ , go to Step 2  
Else raise an error

Step 2:- If  $\text{Request}[i] \leq \text{Available}$ , go to Step 3  
Otherwise  $P[i]$  must wait

Step 3 :-  $\text{Available}[i] = \text{Available} - \text{Request}[i]$   
 $\text{Allocation}[i] = \text{Allocation}[i] + \text{Request}[i]$   
 $\text{Need}[i] = \text{Need}[i] - \text{Request}[i]$



OUTPUT:-

Enter buffer size : 1

Enter your choice :- 1

1. Produce    2. Consume    3. Exit

Your choice : 1

Produced item ..

Produced: 1   Consumed: 0

Your choice : 1

Buffer is full!!!

Produced: 1   Consumed: 0

Your choice : 2

Consumed item

Produced: 1   Consumed: 1

Your choice : 2

Buffer empty!!

Produced: 1   Consumed: 1

Your choice : 3

(6)

## EXPERIMENT - 4

### PRODUCER - CONSUMER PROBLEM

1) Aim:-

To implement the producer-consumer problem using semaphores.

2) ALGORITHM:-

The algorithm uses three semaphores, two counting semaphores full and empty and a binary semaphore mutex.

Step 1:- Enter the size of buffer say  $n$ .

Step 2:- Initialize full=0 and empty=n, mutex=1

Step 3:- Input operation to be performed  
i.e produce or consume or exit

Step 4:- If operation is produce

Step 4.1:- If mutex=1 and buffer not full

begin Step 4.1.1:- mutex = mutex - 1

Step 4.1.2:- full = full + 1

Step 4.1.3:- empty = empty - 1

Step 4.1.4:- print "Item produced"

end Step 4.1.5:- mutex = mutex + 1

Step 4.2:- else print "Buffer Full"

Step 5:- If operation is consume.

Step 5.1:- If mutex=1 and buffer not empty

begin

Step 5.1.1:- mutex = mutex - 1

Step 5.1.2:- signal = signal + 1

Step 5.1.2. :-  $full = full - 1$

Step 5.1.3. :-  $empty = empty + 1$

Step 5.1.4. :- print "Consumed Item"

Step 5.1.5. :-  $mutex = mutex + 1$

end

Step 5.2. :- else print "Buffer Empty"

Step 6. :- Repeat steps 3,4,5 until exit

~~Step 6. :- Repeat steps 3,4,5 until exit~~

OUTPUT:-

Enter philosopher number (1-5) : 1

Enter operation

1. start Eating 2. Stop Eating 3. Exit

1

Philosopher 1 is eating

Enter philosopher number (1-5) : 2

Enter operation

1

Philosopher 2 cannot eat

Enter philosopher number (1-5) : 4

Enter operation

1

Philosopher 4 is eating

Enter philosopher number (1-5) : 1

Enter operation

2

Philosopher 1 is thinking

Enter philosopher number (1-5) : 2

Enter operation

1

Philosopher 2 is eating

(13)

## EXPERIMENT - 5

### DINING PHILOSOPHER'S PROBLEM

1) Aim:-

To simulate the working of dining philosopher's problem.

2) ALGORITHM:-

Step 1:- Declare set of states THINKING,  
HUNGRY, EATING for philosophers  
and flag values and initialize each philosopher  
to think.

Step 2:- Input philosopher's id/number.

Step 3:- Enter the operation to be performed  
ii Eating , Stop Eating , Exit

Step 4:- If operation is Eating

Step 4.1:- Set state as HUNGRY.

Step 4.2:- If the philosophers to left and  
to right are not EATING

Step 4.2.1:- Set state as EATING

Step 4.2.2:- Set flag as 1.

Step 4.3:- If flag is set to 1

Step 4.3.1:- Output philosopher is eating

Step 4.3.2:- else output cannot eat

Step 5:- If operation is Stop eating

Step 5.1:- Set philosopher state as  
THINKING

Step 5-2:- Test for states of philosophers  
to the left and right

Step 5-3:- print philosopher is THINKING

Step 6:- Repeat 3,4,5 until exit

OUTPUT:

Enter the range: 200

Enter the number of disk requests : 3

Enter the head position (<200) : 60

Enter the requests : 50 70 80

Enter your choice

1. FCFS 2. SCAN 3. C-SCAN 4. EXIT

1

No Disk Req

0	60
1	50
2	70
3	80

Seek time : 40ms

Enter your choice

1. FCFS 2. SCAN 3. C-SCAN 4. EXIT

2

No Disk Req

0	60
1	50
2	0
3	70
4	80

Seek Time : 140ms

Enter your choice

1. FCFS 2. SCAN 3. C-SCAN 4. EXIT

# EXPERIMENT - 6

## DISK SCHEDULING

1) AIM:-

To simulate disk scheduling algorithms

- i) FCFS
- ii) SCAN
- iii) C-SCAN

2) ALGORITHM:-

i) FCFS :-

1) Start

2) input no. of requests, say  $n$

3) for  $i = 2$  to  $n+1$

    3.1) input each request, say  $t[i]$

4) for  $i = 1$  to  $n+1$

    4.1)  $\text{totim}[i] = t[i+1] - t[i]$

    4.2) if ( $\text{totim} < 0$ )

        4.2.1)  $\text{totim}[i] = \text{totim}[j] + i$

5) for  $i = 1$  to  $n+1$

    5.1)  ~~$\text{tot} + \text{totim}[i]$~~

6)  ~~$\text{avtim} = (\text{float})[\text{tot}/n]$~~

7) for  $i = 0$  to  $n$

    7.1) print  $t[i], \text{totim}[i], \text{avtim}$

8) Stop

3

No	Disk Req
0	60
1	50
2	0
3	199
4	70

Seek time : 270ms

Enter choice

- 1. FCFS    2. SCAN    3. C-SCAN    4. EXIT
- 4

~~Others~~

SCAN :-

- 1) Start
- 2) Enter no of inputs,  $n$
- 3) Enter head position,  $h$
- 4) Initialize  $t[0]=0$  &  $t[n]=L$
- 5) for  $i=2$  to  $n+2$ 
  - 5.1) Input each request,  $t[i]$
- 6) for  $i=0$  to  $n+2$ 
  - 6.1) for  $j=0$  to  $(n+2)-i-1$ 
    - 6.1.1) if ( $t[j] > t[j+1]$ )
    - 6.1.1.1) Swap values,  $t[j], t[j+1]$
- 7) for  $i=0$  to  $n+2$ 
  - 7.1) if  $t[i] == h$   
 $j=1, k=i$
- 8) for  $p=k+1$  to 2
  - 8.1) ~~atr[p] = t[k+1]~~
- 9) for  $j=0$  to  $n$ 
  - 9.1) if ( $atr[j] > atr[j+1]$ )
  - 9.2) ~~dr[j] = atr[j] - atr[j+1]~~  
else
  - 9.3) ~~dr[j] = atr[j+1] - atr[j]~~
- 10) print average movement (float) sum/n.

C-SCAN:

- 1) Start
- 2) No. of requests,  $n$
- 3) Head position say  $h$
- 4) Head tracks, tot
- 5) for  $i = 3$  to  $n+2$ 
  - 5.1) output disk req.,  $t[i]$
  - 5.2) for  $j = 0$  to  $n+2-i$ 
    - 5.2.1) if ( $t[i] > t[j+1]$ )
      - 5.2.1.1) swap values,  $t[j], t[j+1]$
- 6) for  $i = 0$  to  $n+2$ 
  - 6.1) if ( $t[i] = h$ )
    - 6.1.1)  $j = 1$ , break.
- 7)  $p = 0$
- 8) while ( $t[j] \neq p_{t-1}$ )
  - 8.1)  $atr[p] = t[j]$
  - 8.2)  $j++$ ,  $p++$
- 9) while ( $p_1 = (x+s) \& t[i] \neq t[h]$ )
  - 9.1)  ~~$atr[p] = t[i]$~~
- 10) for  $j = 0$  to  $n+2$ 
  - 10.1) if ( $atr[j] > atr[j+1]$ )
    - 10.1.1)  $d[j] = atr[j] - atr[j+1]$
    - else
      - 10.1.2)  $d[j] = atr[j+1] - atr[j]$

(0.1.3) ~~sum<sub>i</sub> = d[i]~~

11) print sum, sum/n.

## INPUT:

COPY	START	1000
-	LDA	ALPHA
-	ADD	ONE
-	SUB	TWO
-	STA	BETA
ALPHA	BYTE	C'KLNCE
ONE	RESB	2
TWO	WORD	5
BETA	RESW	1
-	END	-

## OUTPUT:

	COPY	START	1000
1000	-	LDA	ALPHA
1003	-	ADD	ONE
1006	-	SUB	TWO
1009	-	STA	BETA
1012	ALPHA	BYT <del>E</del> A	C'KLNCE
1017	ONE	RESB	2
1019	TWO	WORD	5
1022	BETA	RESW	1
1025	-	END	-

Program Length = 25

## EXPERIMENT - 7

(19)

## TWO PASS ASSEMBLER

1) Aim:-

To implement pass 1 and pass 2 of a 2-pass assembler by creating a symbol table with suitable hashing.

2) ALGORITHM:-

Pass 1:-

begin

read first input line

if OPCODE = 'START' then

begin

save # [OPERAND] as starting address

initialize LOCCTR to starting address

write to intermediate file

read next input line

end (if START)

else

initialize LOCCTR to 0

while OPCODE ≠ 'END' do

begin

if this is not a comment line then

begin

if there is a symbol in LABEL field then

begin

search SYMTAB for LABEL

if found then

get error flag (duplicate symbol)

else

insert {LABEL, LOCCTR} into SYMTAB

end (if symbol)

Pass 2:

H^ COPY^ 1000^ 25

T^ 001000^ 001012^ 001017^ 051019^ 2310221

7576786769^ 00005^

search OPTAB for OPCODE  
if found then  
    add 3 to LOCCTR  
else if OPCODE = 'RESW' then  
    add 3 + #{\$OPERAND} to LOCCTR  
else if OPCODE = 'WORD' then  
    add 3 to LOCCTR  
else if OPCODE = 'RESB' then  
    add #{\$OPERAND} to LOCCTR  
else if OPCODE = 'BYTE' then  
begin  
    find length of constant in bytes  
    add length to LOCCTR  
end (if BYTE)  
else  
    set error flag (invalid opcode)  
end (if not a comment)  
write line to intermediate file  
read next input line  
end (while not END)  
~~write last line to intermediate file~~  
say {LOCCTR - starting address} as program length  
end (pass1)

✓  
✓  
✓

Pass 2:

begin

read first output line (from intermediate file)

if OPCODE = 'START' then

begin

write listing line

read next output line

end (if START)

write Header record to object program

initialize first Text Record.

while OPCODE ≠ 'END' do

begin

if this is not a comment line then

begin

search OPTAB for OPCODE

if found then

begin

if there is a symbol

in operand field

then

begin

search SYMTAB for OPERAND

if found then

store symbol value  
as operand address

else

begin

store 0 as operand  
address

set error flag

end.

end (if symbol)

else  
store 0 as operand address  
assemble the object code instruction  
end(if opcode found)  
else if OPCODE = 'BYTE' or 'WORD' then  
convert constant to object code  
if object code will not fit on current  
TEXT Record then  
begin  
write TEXT record to the object program  
initialize new Text Record  
end  
add object code to Text Record  
end(if not comment)  
write listing line  
read next output line  
end(while not END)  
write last Text Record to object program  
write END record to object program  
write last listing line  
end(Pass 2)

~~Pass 2~~  
~~4/11/19~~

### INPUT:

H COPY 001000 00107A  
T 001000 1E 141033 482039 001036 281030  
301015 482061 3C1003 00102A 0C1039 001030  
T 00101E 15 0C1036 482061 081033 4C0000  
454F46 000003 000000  
:  
E 001000

### OUTPUT

MEMORY ADDRESS	CONTENTS			
1000	14103348	20390010	36281030	301015
1010	20613C10	0300102A	0C103900	10200
:				
1040	XXXXXX	XXXXXX04	10300010	30E0
:				
1090	20792C10	36		

# EXPERIMENT-8

(23)

## ABSOLUTE LOADER

1) Aim:-

To implement an absolute loader

2) ALGORITHM:-

begin

read Header Record

verify program name and length

read first Text Record.

while record type ≠ 'E' do

begin

{if object code is in character form, convert  
into internal representation}

move object code to specified location in  
memory

read next object program record

end

jump to address specified in End Record

end.

~~Surf~~  
~~Surf~~  
4/11/17

**INPUT:-**

COPY START 1000  
- LDA ALPHA  
- STA BETA  
ALPHA RESW 1  
BETA RESW 1  
- END -

**OUTPUT**

H ^ COPY ^ 1000 ^ 0c  
T ^ 001000 ^ 0c ^ 000000 ^ 23000  
T ^ 1001102 ^ 1006  
T ^ 1004102 ^ 1009  
E ^ 001000

✓

EXPERIMENT-9  
SINGLE PASS ASSEMBLER

(25)

1) Aim:-

To implement a single pass assembler

2) ALGORITHM:-

```
begin
    if starting address is given then
        LOCCTR = starting address
    else
        LOCCTR = 0
    while OPCODE != END do
        begin
            read a line from the code
            if there is a label then
                begin
                    if this label is in SYMTAB then
                        error
                    else
                        insert (label, LOCCTR) onto SYMTAB
                end
            search OPTAB for the OPCODE
            if found then
                LOCCTR += N
            else if this is an assembly directive then
                update LOCCTR as directed
            else
                error
            write line to intermediate file
        end
    program size = LOCCTR - starting address
end.
```

  
W.M.U.R

## INPUT:

EXI MACRO &A,&B

- LDA &A
- STA &B
- MEND -

SAMPLE START 1000

- EXI N1,N2
- N1 RESW 1
- N2 RESW 1
- END -

## OUTPUT

SAMPLE START 1000

- EXI N1,N2
- LDA N1
- STA N2
- N1 RESW 1
- N2 RESW 1
- END -

# EXPERIMENT - 10

## MACRO PROCESSOR

- 1) Aim:-  
To implement a two pass macro processor

- 2) ALGORITHM:-

Pass 1:

begin

read first output line

while OPCODE != 'MEND' do

begin

if OPCODE = 'MACRO' then

begin

insert label {macro name} onto

NAMTAB

insert label, operand onto DEFTAB

end

else

insert opcode, operand onto DEFTAB

read next output line

end

insert opcode onto DEFTAB

end

Pass 2:

begin

read first output line

while OPCODE != 'END' do

begin

if OPCODE = 'MACRO' then

begin

read next output line

while OPCODE != 'MEND' do

begin

read next input line  
end  
end  
else  
begin  
    Search NAMTAB for OPCODE  
    if found then  
        begin  
            insert operand into ARGTAB  
        read next row of DEFTAB  
        while OPCODE != 'MEND' do  
            begin  
                if OPERAND[0] = '&'amp; then  
                    Search ARGTAB for argument  
                end  
            read next row of DEFTAB  
            end  
        end  
    end  
    read next input line  
end  
end.

### Symbol Table menu

1. Create a Symbol Table.
2. Search in the Symbol Table.
3. Exit

Enter your choice : 1

Enter address : 567

Enter label : abc

Do you want to continue ? (Y/N) ? Y

Enter address : 678

Enter label : ghi

Do you want to continue ? (Y/N) ? Y

Enter address : 123

Enter label : hij

Do you want to continue ? (Y/N) ? N

### Symbol Table

Hash	values	label
0	0	
1	0	
2	123	hij
3	0	
4	0	
5	0	
6	567	abc
7	678	ghi
8	0	
9	0	
10	0	

# EXPERIMENT - II

## SYMBOL TABLE WITH HASHING

1) Aim:-

To implement a symbol table with suitable hashing.

ALGORITHM:-

Set MAX = 11

Create structure symb with add and label[]  
declare structure variable sy[]

void search()

{ open file symbol.txt onto fp1 couth mode  
read la.

for (i=0; i<MAX; i++)

{ Read j and sy[i].add from fp1

if (sy[i].add != 0)

Read sy[i].label from fp1

}

for (i=0; i<MAX; i++)

{ if (sy[i].add == 0)

{ if (strcmp(sy[i].label, la) == 0)

{ set = 1;

s = sy[i].add;

}

}

}

if (set == 1)

Display label is present

else

Display label is not present

3

void display (int a[MAX])

Enter your choice : 2

Enter label : hij

The label is not present in the symbol table.

{ open file symbol.txt into fp with w mode  
Display the symbol table.

for(i=0; i<MAX; i++)

    write i, sy[i].add, sy[i].label onto  
    fp

}

ont create (ont num)

{ ont key;

    key = num % 11;

return key;

}

void display(ont a[MAX], ont key, ont num)

{ Declare flag=0, count=0

if (a[key] == 0)

{ a[key] = num;

    sy[key].add = num;

    strcpy(sy[key].label, l);

}

else

{ i=0;

    while (i < MAX)

{ if (a[i] != 0)

        count++;

        i++;

}

if (count == MAX)

{ Display Hash table is full

    display(a)

}

exit(1);

```

for (i=key+1; i<MAX; i++) {
    if (a[i] == 0)
        {
            a[i] = num;
            flag = 1;
            sy[key].add = num;
            strcpy(sy[key].label, l);
            break;
        }
}

for (i=0; i<key && flag == 0; i++) {
    if (a[i] == 0)
        {
            a[i] = num;
            flag = 1;
            sy[key].add = num;
            strcpy(sy[key].label, l);
            break;
        }
}
}

void main()
{
    Declare a[MAX], num, key, ch, ans = 'y';
    for (i=0; i<MAX; i++)
        a[i] = 0;
    do
    {
        Read ch
        switch(ch)
        {
            case 1: while (ans == 'y')
                {
                    Read num
                    key = create(num);
                    Read l
                    insert(a, key, num);
                }
        }
    }
}

```

Read one

3

display(a)

silak:

case 2: Search

break

case 3: exit()

3

3 while (ch<=3);

