

Explain why sorting all the buckets in bucket sort using an $O(n^2)$ algorithm like insertion sort results in an $O(n)$ sorting algorithm.

If nb is number of elements per bucket, and n is the number of buckets, then the runtime for any individual bucket utilizing insertion sort would be:

$$O(nb^2)$$

The probability of any one element falling into a bucket is $1/n$. Extrapolating an expected value from the statistical concept of a binomial distribution we get $2-1/n$ for nb^2 .

The variance is $1-1/n$. The mean of samples is 1.

$$\text{expected value} = \text{variance} + (\text{mean})^2$$

$$\text{so } nb^2 = 2-1/n$$

So now our runtime becomes:

$$n \times O(2n-1)$$

which the largest factor in there is n , so it simplifies to $O(n)$ regardless of using an sorting algorithm of $O(n^2)$ runtime.

In parlance that I did not have to thoroughly pore over to understand, the buckets are relatively small and the values are evenly distributed (as best as we can get with `rand()`). Any call to insertion sort does not have to take that many steps to complete its task.