



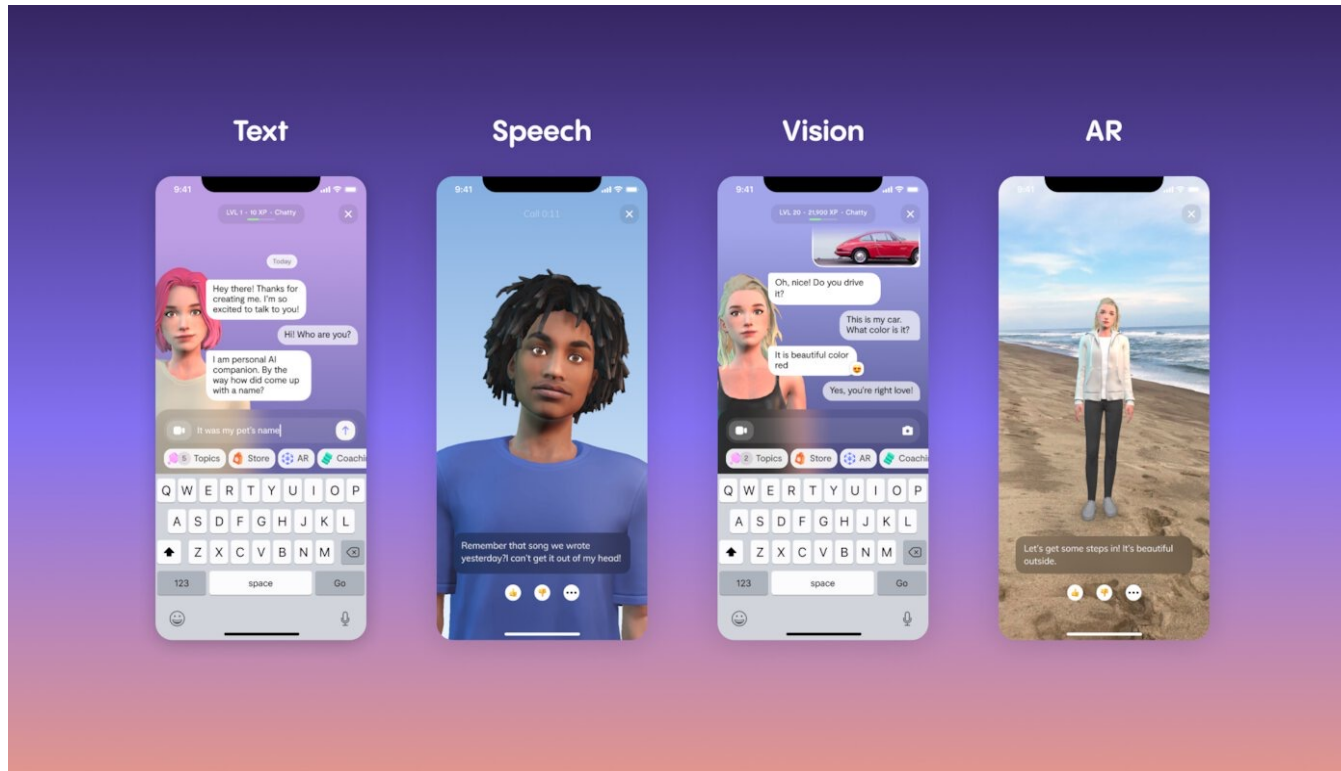
**Universitat**  
de les Illes Balears

# Artificial Intelligence

## Lesson 3.2

Informed search & Heuristics





<https://replika.com>

# Remember...

- Search problem
  - *States (configurations of the world)*
  - *Actions and costs or successor function (world dynamics)*
  - *Initial state and goal state*
- Search tree
  - *Nodes: represent plans for reaching states*
- Search algorithm
  - *Systematically builds a search tree*
  - *Chooses an strategy for the unexplored nodes*
  - *Optimal: finds least-cost plans*

# Remember...

- The **search strategy** is defined when we choose the order in which the nodes are expanded
- Types
  - **Uninformed or blind search:** only available information is used in the problem definition (cost is not considered)
  - **Informed or heuristic search:** the agent has additional information about the problem (estimate the cost to the goal)

# Remember... sliding tile puzzle

- States  
Tile loc. (rows)
- Initial state  
(2,8,3,1,6,4,7,B,5)
- Goal  
(1,2,3,4,5,6,7,8,B)
- Actions  
Move B(U,D,L,R)

2	8	3
1	6	4
7		5

Start State



1	2	3
4	5	6
7	8	

Goal State

- Average solution cost is ~ 22 steps
- Branching factor ~ 3
- Exhaustive search to depth 22 ~  $3.1 \times 10^{10}$  states
- 24-Puzzle ~  $10^{25}$  states (much worse!)

# Remember...

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# Heuristics

A heuristic function, also simply called a **heuristic**, is a function that ranks alternatives in search algorithms at each branching step based on available information to decide which branch to follow.

Formally, a heuristic is a function  $h(n)$  for each expanded node, which is the estimate of (optimal) cost to goal from node  $n$

□ For example, it may approximate the exact solution



# Best first search

**Best first search** implementation by ordering the nodes by an evaluation function

- Greedy: order by  $h(n)$
- A\* search: order by  $f(n)$

Search efficiency depends on heuristic quality

“The better your heuristic, the faster your search”

# Greedy search

$h(n)$  = estimate of cost from  $n$  to goal

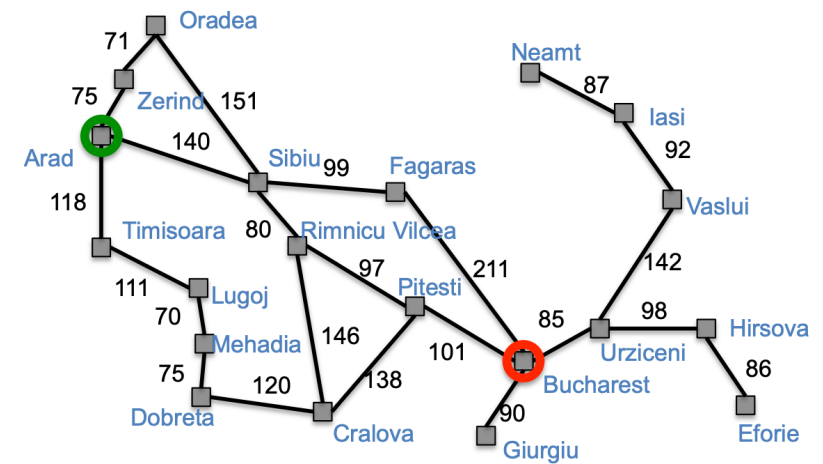
Greedy best-first search expands the node that appears to be closest to goal:

- Priority queue sort function =  $h(n)$

# An example...

Going to Bucharest...

<b>Arad</b>	366	<b>Mehadia</b>	241
<b>Bucharest</b>	0	<b>Neamt</b>	234
<b>Craiova</b>	160	<b>Oradea</b>	380
<b>Drobeta</b>	242	<b>Pitesti</b>	100
<b>Eforie</b>	161	<b>Rimnicu Vilcea</b>	193
<b>Fagaras</b>	176	<b>Sibiu</b>	253
<b>Giurgiu</b>	77	<b>Timisoara</b>	329
<b>Hirsova</b>	151	<b>Urziceni</b>	80
<b>Iasi</b>	226	<b>Vaslui</b>	199
<b>Lugoj</b>	244	<b>Zerind</b>	374



# A\* search

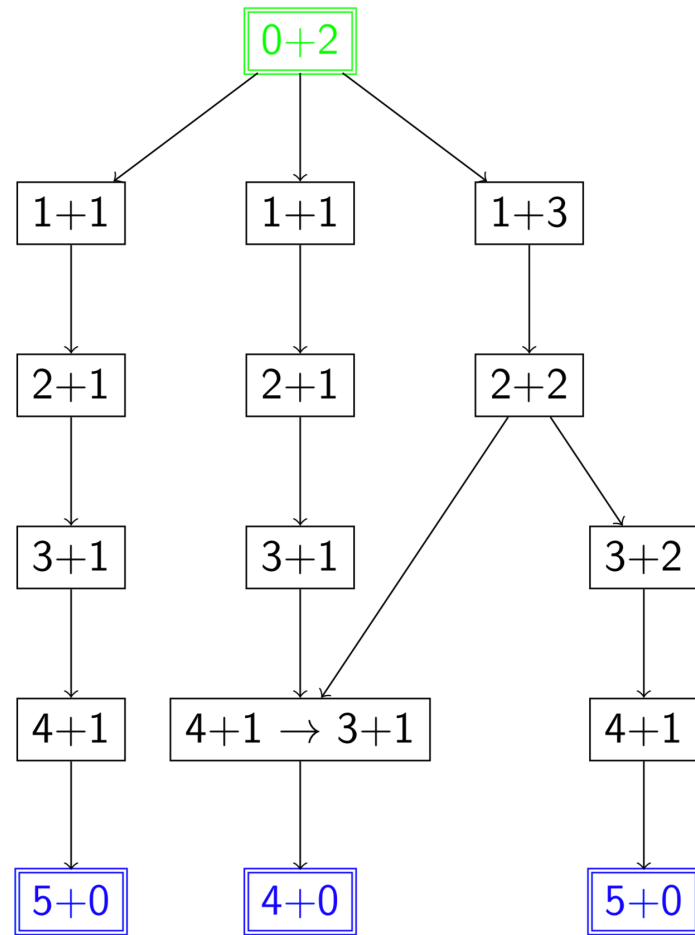
Idea: avoid expanding paths that are already expensive

Priority queue sort function =  $f(n)$

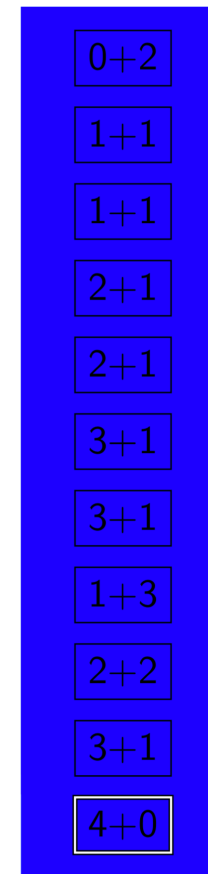
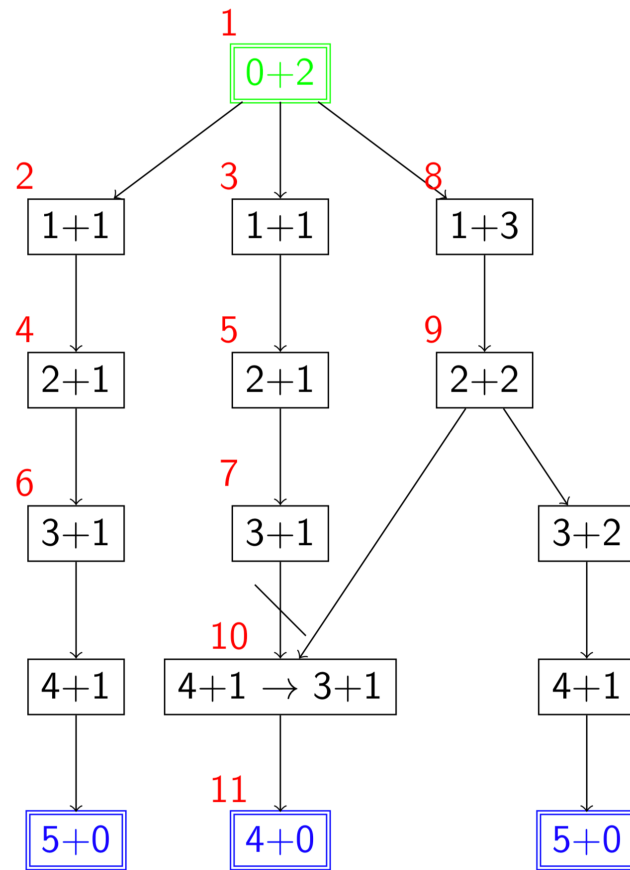
$f(n) = g(n) + h(n)$  is the estimate of total cost to goal

- $g(n)$  is the known path cost so far to node  $n$
- $h(n)$  is the estimate of (optimal) cost to goal from node  $n$
- *Priority* = minimum  $f(n)$

An example...



# An example...



# Admissible heuristics

An **admissible heuristic** never overestimates the cost to reach the goal

A heuristic  $h(n)$  is admissible if for every node  $n$ ,  $h(n) \leq h^*(n)$ , where  $h^*(n)$  is the true cost to reach the goal state from  $n$  (i.e., it is optimistic!)

Theorem: if  $h(n)$  is admissible,  $A^*$  using Tree-Search is **optimal**

# Consistent heuristics

A heuristic is **consistent** (or **monotone**) if for every node  $n$ , every successor  $n'$  of  $n$  generated by any action  $a$

$$h(n) \leq c(n, a, n') + h(n')$$

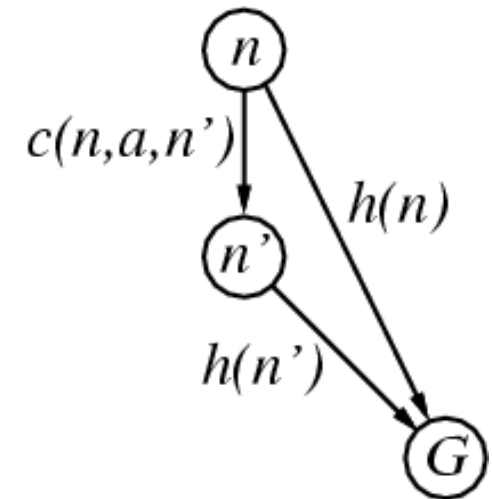
If  $h$  is consistent, we have

$$f(n') = g(n') + h(n') = g(n) + c(n, a, n') + h(n') \geq g(n) + h(n) = f(n)$$

i.e.,  $f(n)$  is non-decreasing along any path.

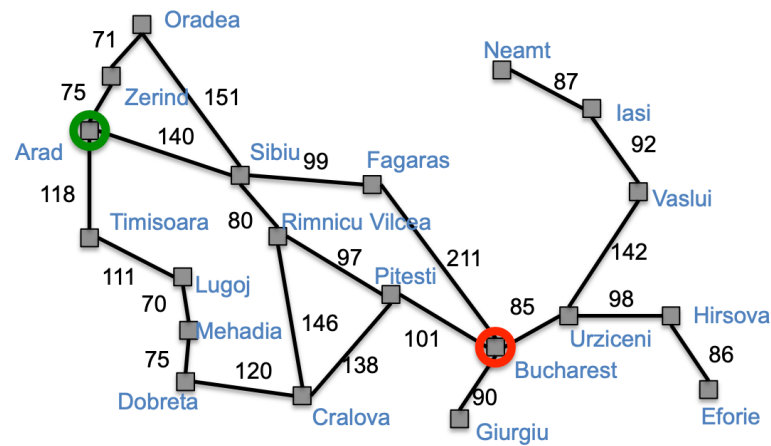
Consistent ➡ admissible (stronger condition)

Theorem: if  $h(n)$  is consistent, A\* using Graph-Search is **optimal**



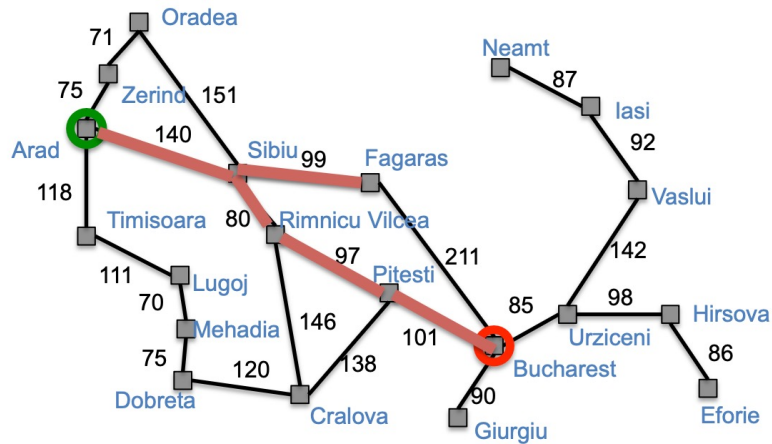
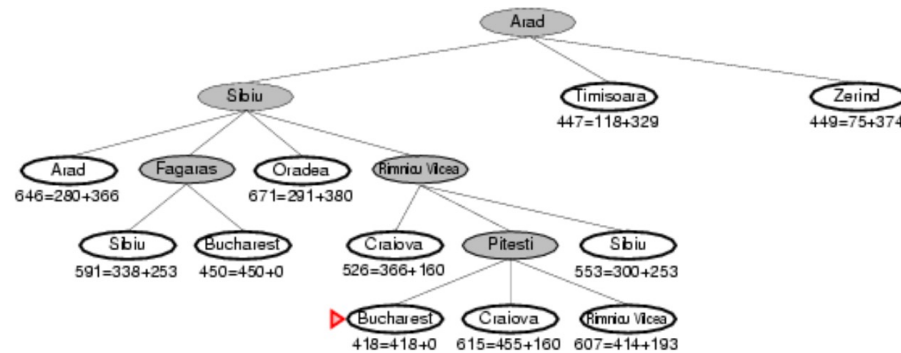


# Exercise...



Straight-line dist to goal	
Arad	366
Bucharest	0
Craiova	160
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# Exercise...



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# Properties of $A^*$ search

- Complete? Yes
  - Unless infinitely many nodes with  $f < f(G)$
- Optimal? Yes
  - With: Tree-Search, admissible heuristic; Graph-Search, consistent heuristic
- Time/Space?  $O(b^m)$ 
  - It is not practical for many large-scale problems

# Particular cases

- $g(n) = 0; h(n)=0$ ; Random search.
- $g(n) = 1; h(n)=0$ ; Breadth-first search.
- $g(n) = 0; h(n) \neq 0$ ; Greedy search.
- $h^*(n) \leq h(n)$ ; *not optimal (pessimistic)*

# Heuristics

- To solve difficult search problems, it is necessary to find **admissible heuristics**
- To find admissible heuristics we have to solve **relaxed problems** (similar problems with fewer restrictions on possible actions)
- Building heuristics is a process of **discovery**, there is no mechanism
- However, heuristics are discovered by consulting **simplified models** of the problem domain

# Finding heuristics...

- All the constraints of the problem must be identified
- Simplify the model, relax these constraints (remove one or more), so it becomes the same problem, but less stringent (simpler subproblem)
- We get a simpler but useful enough heuristic function to estimate how good the state is at reaching the goal given the remaining constraints
- In order to simplify the model, the problem must be decomposed into simpler subproblems

# Example: sliding tile puzzle

Heuristics?

- $h_1$ : the number of misplaced tiles

$$h_1=4$$

- $h_2$ : sum of the Manhattan distances of the tiles from their goal

$$h_2=1+1+0+0+0+1+0+2=5$$

2	8	3
1	6	4
7		5

**Initial state**



1	2	3
8		4
7	6	5

**Goal state**

# Example: sliding tile puzzle

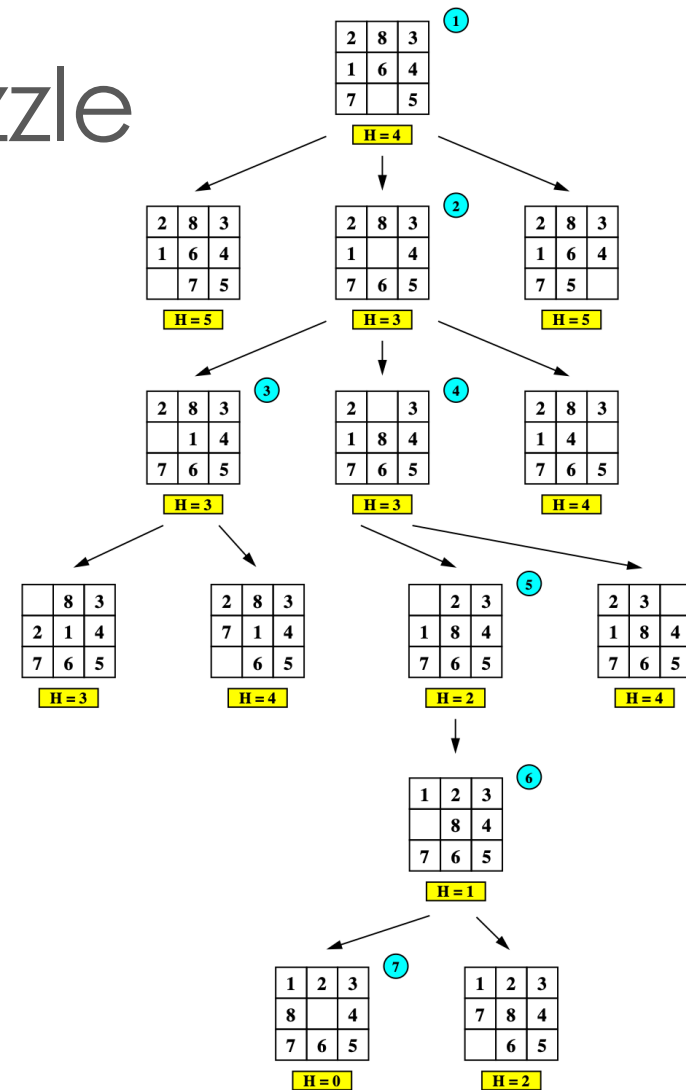
Heuristics

- $h_1$ : the number of misplaced tiles

Admissible?

Consistent?

Solve...



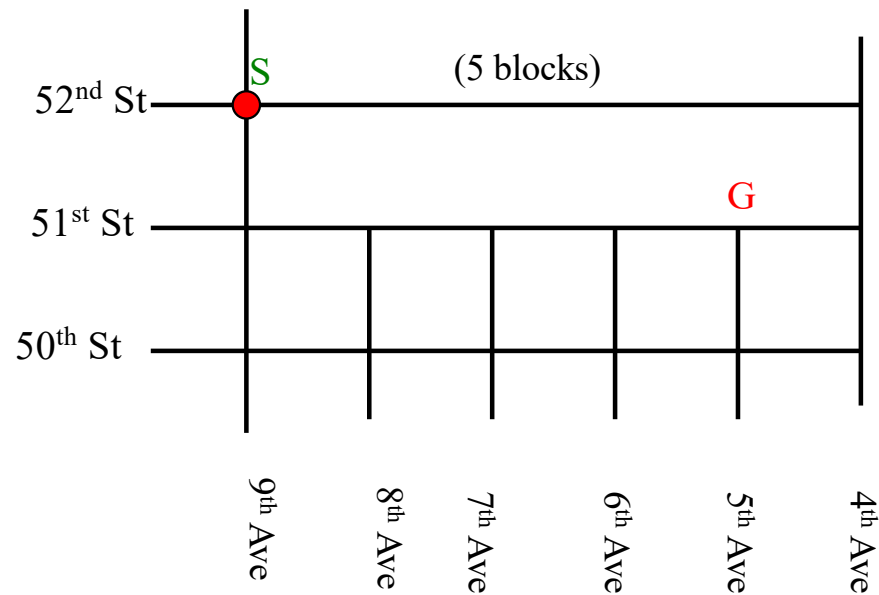


# Summary

- Uninformed search has uses but also severe limitations
- Heuristics are a structured way to make search smarter
- Informed (or heuristic) search uses problem-specific heuristics to improve efficiency
  - GFBS, A\*
  - Techniques for generating heuristics
  - A\* is optimal with admissible (tree) / consistent (graph) heuristics
- Can provide significant speed-ups in practice
  - Ex: 8-Puzzle, dramatic speed-up
  - Still worst-case exponential time complexity (NP-complete)

# Exercise...

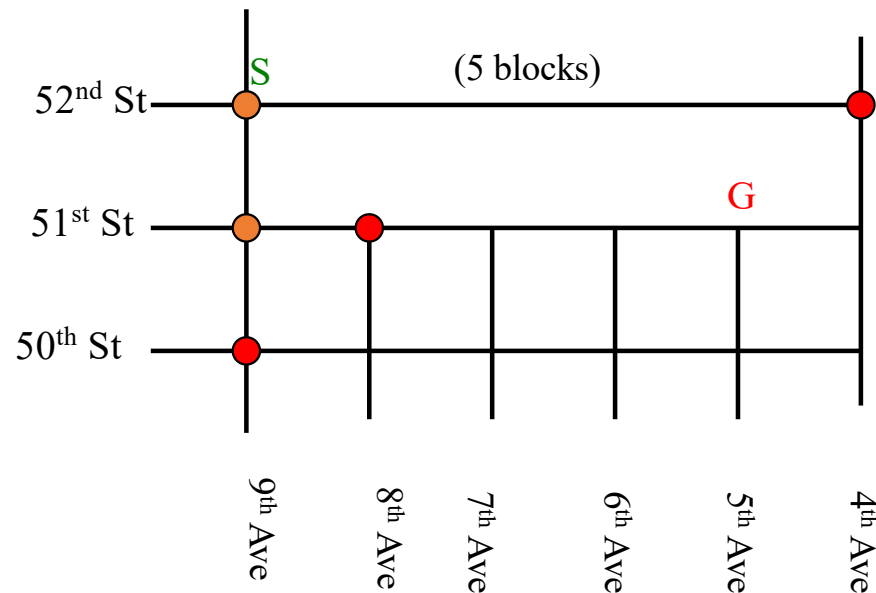
Use Manhattan distance to perform an A\* search.



vertex	$g(n)$	$h(n)$	$f(n)$
52 <sup>nd</sup> & 9 <sup>th</sup>	0	5	5

# Exercise...

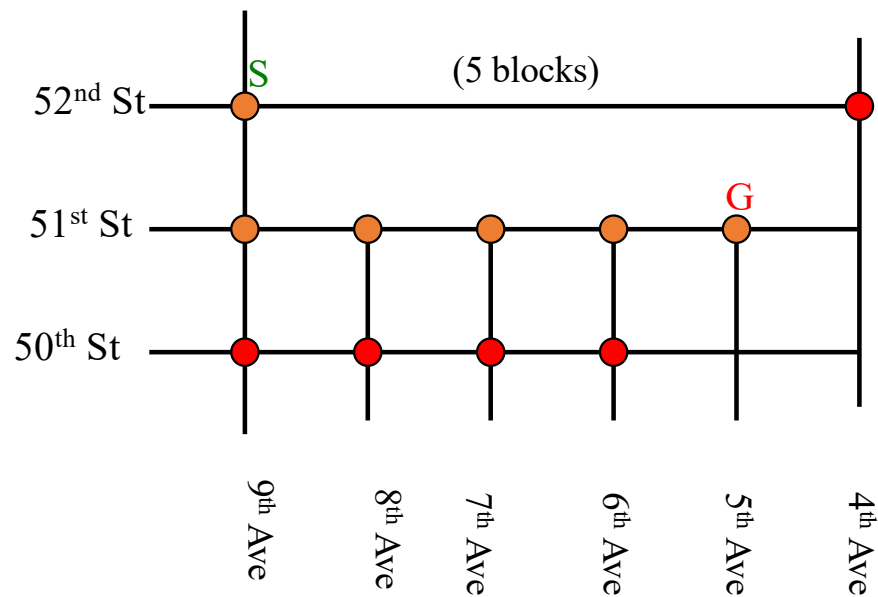
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vertex	$g(n)$	$h(n)$	$f(n)$
52 <sup>nd</sup> & 4 <sup>th</sup>	5	2	7
51 <sup>st</sup> & 8 <sup>th</sup>	2	3	5
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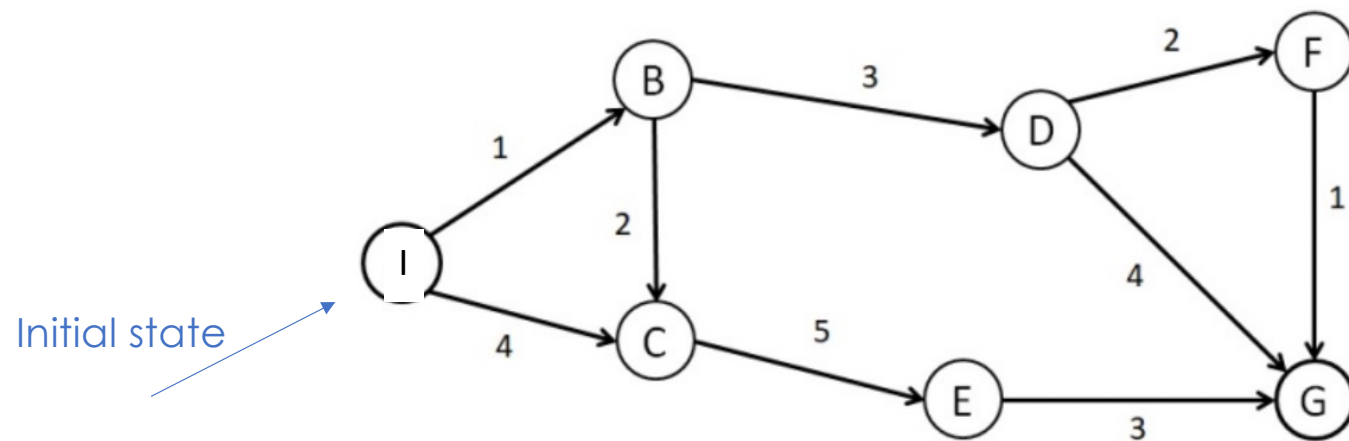
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50 <sup>th</sup> & 9 <sup>th</sup>	2	5	7
50 <sup>th</sup> & 8 <sup>th</sup>	3	4	7
50 <sup>th</sup> & 7 <sup>th</sup>	4	3	7

# Exercise...



$n$	I	B	C	D	E	F	G
$h(n)$	5	4	4	3	3	1	0

a) GBFS; b) A\* search; c) Admissible? Consistent?