

# Artificial Intelligence Applied to the Web

## Chapter 3 Part 2 - Web Page Rank Algorithms

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# Outline

**Page Rank**

HITS

SALSA

# PageRank

## Idea

- An Hyperlink that point to me works like a *recommendation* from other page.
- With more recommendation, more **important** is my page.
- The more important are my recommenders, the more **important** is my page.

## Rank

A **numeric value** to represent the importance of a page.

**We aren't inspecting the content.** We are looking for recommendations

# PageRank

Probabilistic point of view.

- We can view the importance of a web page  $i$  as the probability that a random surfer on the Internet opens a browser to any page and starts following **hyperlinks** visits  $i$ .
- Weights on edges in a transition matrix could be assigned in a *probabilistic way*.
- We can model the process as a **random walk** on graphs. Each page has equal probability  $\frac{1}{n}$  to be chosen as a starting point.
- Then, probability that page  $i$  is visited after one step is equal to  $Ax$  and so on.
- The probability that page  $i$  will be visited after  $k$  steps is equal to  $A^k x$ .
- That sequence converges to a **unique probabilistic vector**  $v^*$  called the *stationary distribution*.
- **This will be our PageRank.**

# PageRank

## Mathematical foundation

- $\Pi = (r_i)$  The vector of Rank number, the rank  $r_i$  for page  $p_i$
- **Each page that point to me, add a fraction of its own rank for my total rank**
- That's means it is a **linear combination**:  $r_i = \sum_k T_{ik} r_{ik}$
- $T = (t_{ik})$  The transition matrix; and it is stochastic:  $\sum_k t_{ik} = 1$  (the total rank is conserved).
- **What can we do with this?**

# PageRank

## Mathematical foundation

### Perron-Frobenius Theorem

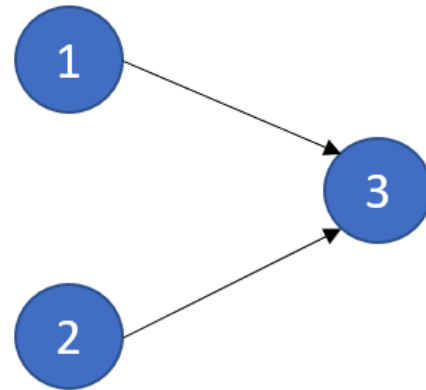
- Is about the **eigenvectors** and **eigenvalues** of *non-negative and irreducible matrices*.
- A matrix is called *non-negative* if all of its entries are  $\geq 0$ .
- A matrix is called *irreducible* if for any of its entries  $(i, j)$  there is  $k$  such that the  $(i, j)$  entry of  $A^k$  is **positive**.
- It says that for  $A$  a non-negative and irreducible matrix, it is an *eigenvalue*  $\lambda_{\max}$  and for all other eigenvalues  $\lambda$  we have  $|\lambda| \leq \lambda_{\max}$ .
- Moreover, if the sum of all entries of each column of  $A$  is 1, then  $\lambda_{\max} = 1$ .

But, there is a little problem: **the matrices we obtain from the Web are not always irreducible**

# PageRank

## Mathematical foundation - Issue 1

There could be: **nodes with no outgoing edges**, also called **dangling nodes**.



$$v_0 = \begin{bmatrix} \frac{1}{3} \\ \frac{1}{3} \\ \frac{1}{3} \end{bmatrix}$$

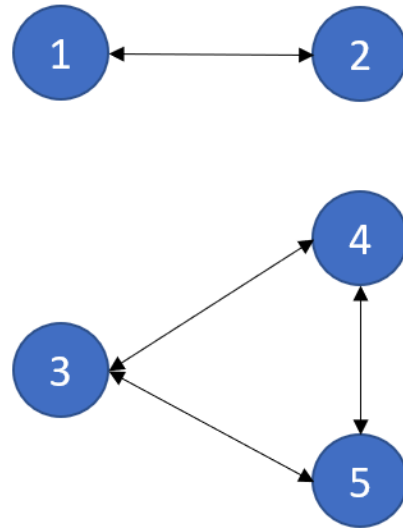
$$v_1 = \begin{bmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 1 & 1 & 0 \end{bmatrix} \cdot \begin{bmatrix} \frac{1}{3} \\ \frac{1}{3} \\ \frac{1}{3} \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ \frac{2}{3} \end{bmatrix}$$

$$v_2 = \begin{bmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 1 & 1 & 0 \end{bmatrix} \cdot \begin{bmatrix} 0 \\ 0 \\ \frac{2}{3} \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}$$

# PageRank

## Mathematical foundation - Issue 2

There could be: **disconnected components**.



$$A = \left[ \begin{array}{cc|ccc} 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ \hline 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} \\ 0 & 0 & \frac{1}{2} & 0 & \frac{1}{2} \\ 0 & 0 & \frac{1}{2} & \frac{1}{2} & 0 \end{array} \right]$$

$$v = \begin{bmatrix} 1 \\ 1 \\ 0 \\ 0 \\ 0 \end{bmatrix} \quad u = \begin{bmatrix} 0 \\ 0 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

Both eigenvectors with eigenvalue = 1



# PageRank

## Mathematical foundation - The solution of Page and Brin

- They propose a fix positive constant  $d$  between 0 and 1 called the **damplng factor** (with a typical value  $d = 0.85$ ).
- This models the random surfer behaviour!
- Now, the *transition matrix* used to compute PageRank will be a new matrix  $M$  such that:

$$M = dT + \frac{(1 - d)}{n}E$$

with  $E$  as a matrix of ones.

- $M$  is called the **Google matrix**.

# PageRank

## PageRank Algorithm

The assumption is the **importance of a page** is given for the **importance of the pages that pointed it**.

$$r_p^{(k+1)} = (1 - d) \frac{1}{n} r_p^{(k)} + d \sum_{\forall q, p \in P \mid q \rightarrow p} \frac{1}{N_q} r_q^{(k)}$$

# PageRank

## PageRank Algorithm

$$r_p^{(k+1)} = (1 - d) \frac{1}{n} r_p^{(k)} + d \sum_{\forall q, p \in P \mid q \rightarrow p} \frac{1}{N_q} r_q^{(k)}$$

Where:

- $r_p$ : Importance of page  $p$
- $n$ : Number of pages in the web graph
- $d$ : Probability that the surfer follows some out-links of  $q$  when visit that page
- $N_q$ : Number of out-links from page  $q$
- $r_q$ : Importance of page  $q$
- $\frac{1}{N_q}$ : Conditional probability of going to another page

**The PageRank is the fixed point value of this recurrence!**

# PageRank

## PageRank Algorithm

In matrix form:

$$\Pi^{(k+1)} = \left( dT + \left( \frac{1-d}{n} \right) E \right) \Pi^{(k)}$$

- $d$ : The probability of going out the  $n$  node
- $T$ : An transition matrix (stochastic) that is interpreted as the transition probability. But in the Google way they consider **equiprobability**  $\frac{1}{N_q}$
- $E$ : Is the *1 matrix*. A matrix filled with 1.

**Basically, replace the usual transition matrix with the Google matrix and compute the eigenvector with eigenvalue equal to 1**

# PageRank

## PageRank Algorithm

### Intuitive Justification

- Thought as a **model of user behavior**.
- **Random surfer** that keep clicking on links **starting from one random page**, and **some time get bored** and **continue** from other random page.
- **PageRank** value correspond to the **probability that the random surfer visit the page**.
  - The **stationary probability**.
- Then the **d parameter** correspond to the **probability to start again from another page**.

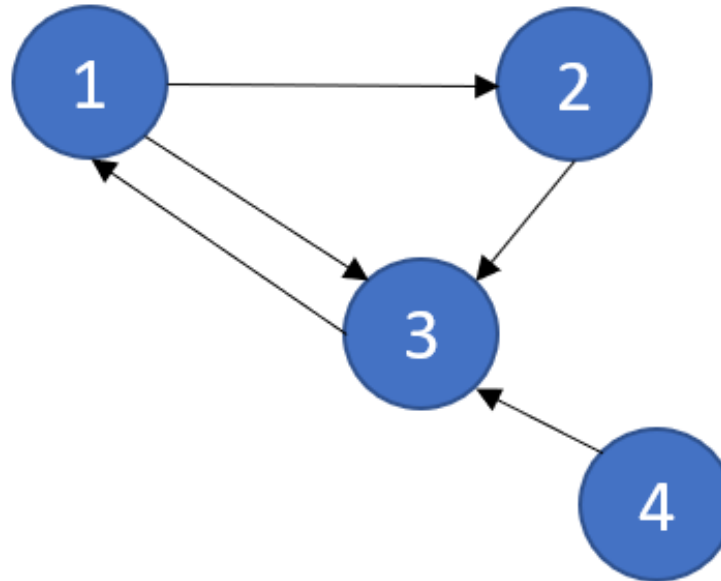
# PageRank

## PageRank Algorithm

1. Initialize:  $\Pi^{(0)} = \left(\frac{1}{n}, \dots, \frac{1}{n}\right)$
2. Set  $d$ , usually  $d = 0.85$
3. Calculate:  $\Pi^{(k+1)} = \left(dT + \frac{(1-d)}{n}E\right)\Pi^{(k)}$
4. If  $\|\Pi^{(k+1)} - \Pi^{(k)}\| < \xi$  stop and return. Else,  $\Pi^{(k)} = \Pi^{(k+1)}$  and go to the point 3.

# PageRank

## PageRank Algorithm - Example



Considering the graph above, get its PageRank

# PageRank

## PageRank Algorithm

### Disadvantages

- If a page is pointed by another one, it means that the page receives a vote for the PageRank calculus.
- If a page is pointed by a lot of pages, it means that the page is important.
- **Only the good pages are pointed by others one**, but:
  - **Reciprocal link**: If the page A links page B, then page B will link page A.
  - **Link Requirements**: Some web pages give electronic gifts, like programs, documents etc., if another page points it.
  - **Near persons community**: For instance, friends and relatives that from their pages point another friend or relative only because of the human relationship between them.



# PageRank

## PageRank Algorithm

- **Not the real Google algorithm.** It is a very carefully hidden secret.
- The original PageRank doesn't consider text content (keywords) on the page. But the real one **does**. The real algorithm also considers **n-grams**
- The real algorithm also considers user behavior. They capture it with:
  - Click on links
  - Google toolbar
  - Google web-accelerator (a Google proxy)
  - Gmail and Youtube

# PageRank

## PageRank sensibility and Link spammers

- Could a group of people artificially reference pages between them in order to increase the rank?
- A **parallel algorithm search for spammer** and lowers its rank.
- The sensibility from a spamming update is:

$$\|\Pi - \tilde{\Pi}\|_1 < \frac{2d}{1-d} \sum_{i \in U} r_i$$

- $U$ : A community of spammers
- $\Pi$ : Change on each  $i \in U$ ,  $r_i$  are the original rank
- The rank is bound nevertheless any action of a group  $U$  of spammers

# Outline

Page Rank

**HITS**

SALSA

# HITS Algorithm

## HITS Algorithm Introduction

- Proposed by John Kleinberg in 1998.
- Stands for Hypertext Induced Topic Search.
- Expand the list of relevant pages returned by a search engine.
- Produce two rankings: Authority ranking and Hub ranking.

### Assumptions

- A credible page will point to credible pages.
- Credible pages are pointed by others.

**The page ranking depends on the user query and the hyperlink structure that follows from paths of the most credible pages**

# HITS Algorithm

## Authority and Hubs Pages

### Definitions

- Authority: A page with many in-links.
- Hub: A page with many out-links.

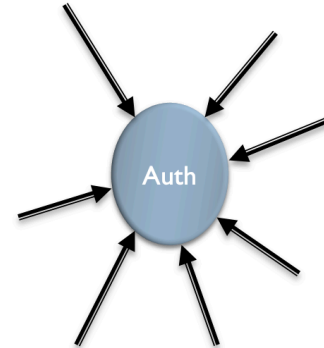
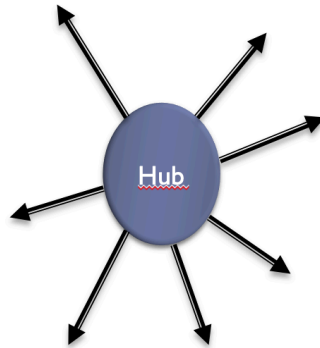
The idea is that:

- A page may have authoritative content on some topic and thus many people trust it and thus link to it.
- When a user comes to a hub page, he/she will find many useful links which take him/her to good content pages on the topic.

# HITS Algorithm

## Authority and Hubs Pages

- If  $q$  it is good hub, then  $q$  point to many good authority pages  $p$ .
- If  $p$  it is a good authority, then  $p$  is pointed by many good hub pages  $q$ .
- Authorities and Hubs have a **mutual reinforcement** relationship.
- We can give a measure of the quality of goodness for authority and hubs. We call it:  
Authority and Hub weight ( $a_p, h_q$ )



# HITS Algorithm

## How it works

### Assumptions

- The authority level (or rank) came from in-edges.
- A **good authority** come from **good hubs** and **a good hub** contains links that point to **good authorities**.

A **simple method** to differentiate the page's relevance is:

- First **assigning non-negative weights**, depending if the page is hub or authoritative. Well, finally the page have both of them.
- Next, the weights are **adjusted by an iterative process** and the **relative page's importance** in the community is calculated.

# HITS Algorithm

## How it works

Given a query  $q$ , HITS collects a set of pages as follows:

1. It send the query  $q$  to a search engine system and collects  $t$  highest ranked pages, which assume to be highly relevant to the search query. This is called the **root** set  $W$ .
2. It then grows  $W$  by including any page pointed to by a page in  $W$  and any page that points to a page in  $W$ . It restrict its size by allowing each page in  $W$  to bring at most  $k$  pages. This set is called the **base set**  $S$ .



# HITS Algorithm

## How it works

Assuming we have a set  $S$  of pages connected where:

- $V$  is the set of pages (or nodes).
- $E$  is the set of directed edges (or links).
- Then  $G = (V, E)$
- We use  $L$  to denote the adjacency matrix of the graph, where:

$$L = \begin{cases} 1 & \text{if } (i, j) \in E \\ 0 & \text{otherwise} \end{cases}$$

# HITS Algorithm

## How it works

- Let the authority score of the page  $i$  be  $a(i)$ , and the hub score of the page  $i$  be  $h(i)$ .
- The mutual reinforcing relationship of the two scores is represented as:

$$a(i) = \sum_{(j,i) \in E} h(j)$$

The authority score of the page  $i$  correspond to the sum of all the hubs that are point to it.

$$h(i) = \sum_{(i,j) \in E} a(j)$$

The hub score of the page  $i$  correspond to the sum of all the authorities it points to.

# HITS Algorithm

## How it works

- Then, we use  $a$  to denote the column vector with all the authority scores and  $h$  to denote the column vector with all the hub scores.
- This is equivalent to say:

$$a = L^T h$$

$$h = L a$$

- If we use  $a_k$  and  $h_k$  to denote authority and hub scores at the  $k$ -th iteration, the iterative processes are:

$$a_k = L^T L a_{k-1} \iff a_{k+1} = L^T L a_k$$

$$h_k = L L^T h_{k-1} \iff h_{k+1} = L L^T h_k$$

# HITS Algorithm

## HITS Algorithm Steps

1. Initialize  $a = (1, \dots, 1)$  and  $h = (1, \dots, 1)$

2. Calculate:

$$a_{k+1} = L^T L a_k$$

$$h_{k+1} = L L^T h_k$$

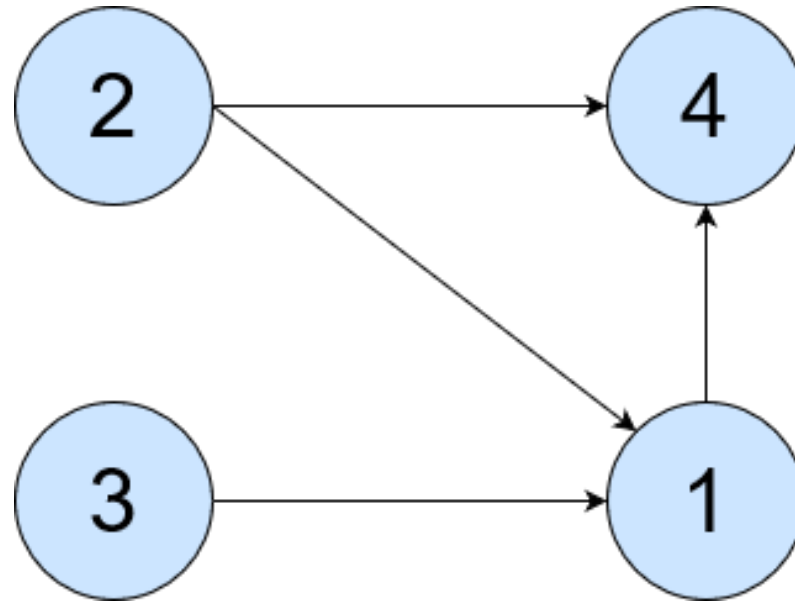
3. Normalize  $a_k$  and  $h_k$

4. Repeat until:  $\|a_{k+1} - a_k\|_1 < \varepsilon_a \wedge \|h_{k+1} - h_k\|_1 < \varepsilon_h$

5. Get the most auth pages and the most hub pages.

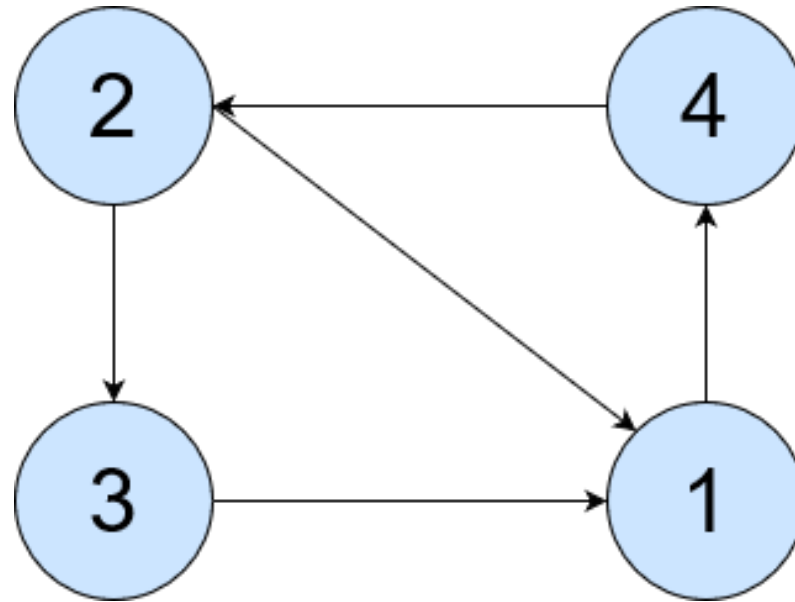
# HITS Algorithm

## Example



# HITS Algorithm

## Example



# HITS Algorithm

## How it works

By construction of HITS:

- After using the query to extract the pages related with the query's terms, the page's text content is ignored in the page's rank task.
- Being the algorithm purely **hyperlink-based computation**.

CLEVER Project (Chakrabarti S.)

- Addresses the problem by considering query's terms in the calculus of the above equations.
- A non-negative weight, whose initial value is basis on the text that surround the hyperlink expression (a tag in HTML)

# HITS Algorithm

## Advantages and Disadvantages

### Advantages

- Double ranking (by Authority and by Hub).
- Rank pages according to a query topic.

### Disadvantages

- Doesn't have the anti-spam capability of PageRank.
- Topic drift.
- Query-dependence.
- Query time evaluation.



# Outline

Page Rank

HITS

**SALSA**

# SALSA

## What is it?

### Source

“Google’s PageRank and Beyond: The Science of Search Engine Rankings”. Amy N. Langville and Carl D. Meyer

- Proposed by R. Lempel and S. Moran in 2001.
- SALSA stands for Stochastic Approach to Link Structure Analysis.
- Like HITS, SALSA create both Authority and Hub scores for webpages.
- SALSA creates a neighborhood graph showing the closeness between Authority pages and Hub pages.

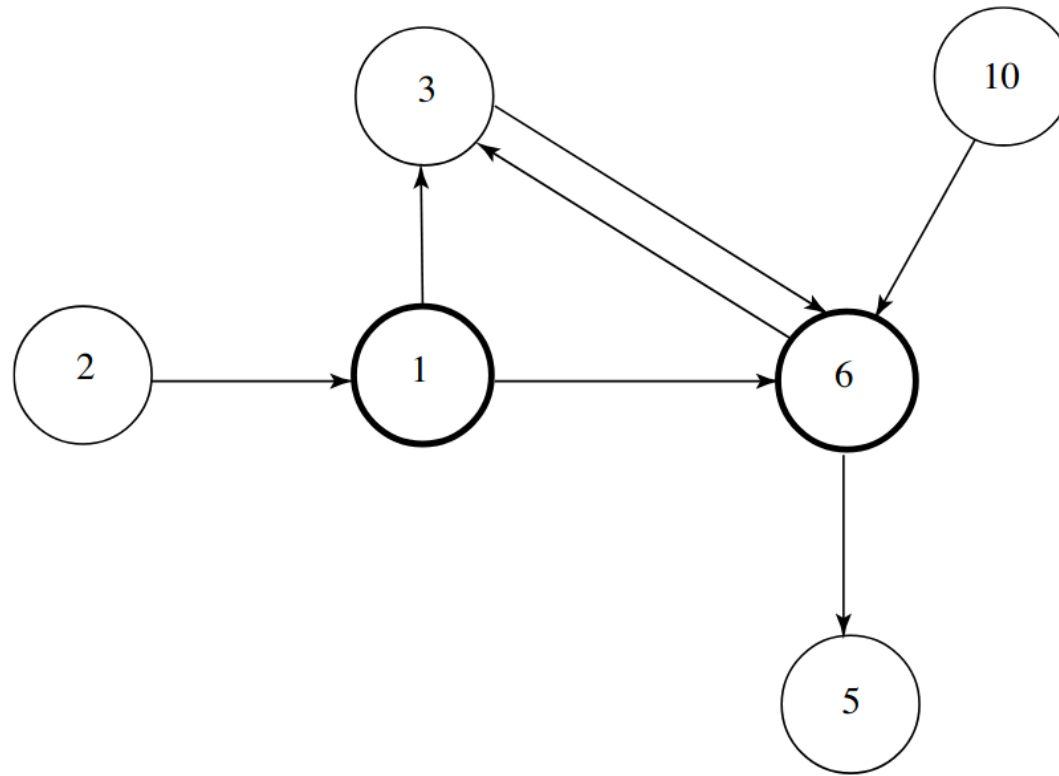
# Other Algorithms

## SALSA

- Rather than forming an adjacency matrix  $L$  for the neighborhood graph  $N$ , a bipartite undirected graph, denoted  $G$ , is built.
- $G$  is defined by the sets:
  - $V_h$ : Set of Hub nodes (all nodes in  $N$  with outdegree  $> 0$ ).
  - $V_a$ : Set of Authority nodes (all nodes in  $N$  with indegree  $> 0$ ).
  - $E$ : Set of directed edges in  $N$ .
- Note that a node in  $N$  may be in both  $V_h$  and  $V_a$ .

# Other Algorithms

## SALSA Example



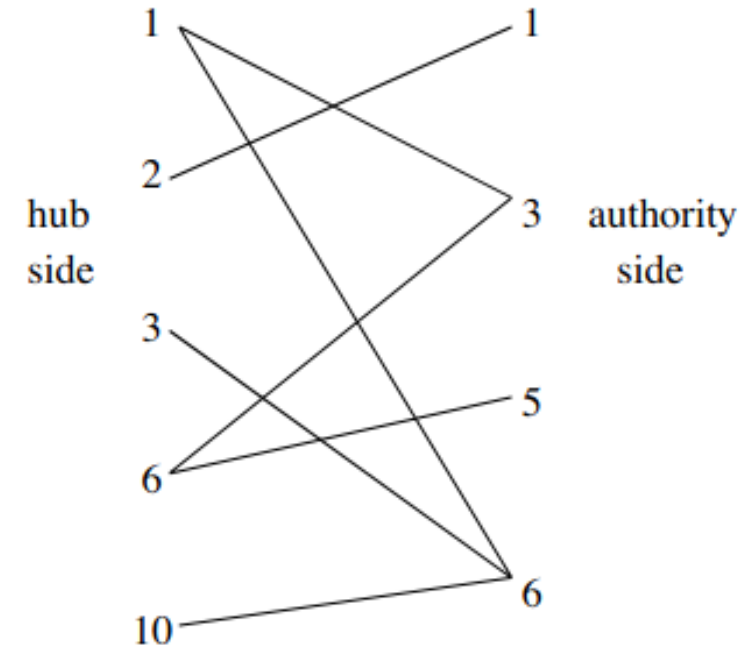
# Other Algorithms

## SALSA Example

- From the previous graph, we have that:

$$V_h = \{1, 2, 3, 6, 10\} \quad V_a = \{1, 3, 5, 6\}$$

- Then, we can create the bipartite undirected graph  $G$ , as shown.
- Every directed edge in  $E$  is represented by an undirected edge in  $G$ .



# Other Algorithms

Where does SALSA fit in?

- Matrices  $H$  and  $A$  can be derived from the adjacency matrix  $L$  used in the HITS and PageRank methods.
- HITS used unweighted matrix  $L$ .
- PageRank uses a row weighted version of matrix  $L$ .
- SALSA uses both row and column weighting.

# Other Algorithms

How are A and H computed in SALSA?

To get A and H matrices, we have to calculate normalize versions of  $L$ :

- Let  $L_r$  be  $L$  with each nonzero row divided by its row sum.
- Let  $L_c$  be  $L$  with each nonzero column divided by its column sum.

In the previous example, it will be:

$$\begin{array}{c}
 \begin{array}{c} 1 \\ 2 \\ 3 \\ 5 \\ 6 \\ 10 \end{array}
 \begin{pmatrix}
 & \begin{array}{cccccc} 1 & 2 & 3 & 5 & 6 & 10 \end{array} \\
 \begin{array}{c} 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \end{array} & \begin{array}{c} 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{array} & \begin{array}{c} 1 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \end{array} & \begin{array}{c} 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{array} & \begin{array}{c} 1 \\ 0 \\ 1 \\ 0 \\ 0 \\ 1 \end{array} & \begin{array}{c} 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{array}
 \end{pmatrix}
 \end{array}
 L_r =
 \begin{array}{c}
 \begin{array}{c} 1 \\ 2 \\ 3 \\ 5 \\ 6 \\ 10 \end{array}
 \begin{pmatrix}
 & \begin{array}{cccccc} 1 & 2 & 3 & 5 & 6 & 10 \end{array} \\
 \begin{array}{c} 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \end{array} & \begin{array}{c} 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{array} & \begin{array}{c} \frac{1}{2} \\ 0 \\ 0 \\ 0 \\ \frac{1}{2} \\ 0 \end{array} & \begin{array}{c} 0 \\ 0 \\ 0 \\ 0 \\ \frac{1}{2} \\ 0 \end{array} & \begin{array}{c} \frac{1}{2} \\ 0 \\ 1 \\ 0 \\ 0 \\ 1 \end{array} & \begin{array}{c} 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{array}
 \end{pmatrix}
 \end{array}
 L_c =
 \begin{array}{c}
 \begin{array}{c} 1 \\ 2 \\ 3 \\ 5 \\ 6 \\ 10 \end{array}
 \begin{pmatrix}
 & \begin{array}{cccccc} 1 & 2 & 3 & 5 & 6 & 10 \end{array} \\
 \begin{array}{c} 0 \\ 1 \\ 0 \\ 0 \\ 0 \\ 0 \end{array} & \begin{array}{c} 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{array} & \begin{array}{c} \frac{1}{2} \\ 0 \\ 0 \\ 0 \\ \frac{1}{2} \\ 0 \end{array} & \begin{array}{c} 0 \\ 0 \\ 0 \\ 0 \\ 1 \\ 0 \end{array} & \begin{array}{c} \frac{1}{3} \\ 0 \\ \frac{1}{3} \\ 0 \\ 0 \\ \frac{1}{3} \end{array} & \begin{array}{c} 0 \\ 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{array}
 \end{pmatrix}
 \end{array}
 \end{array}$$

# Other Algorithms

How are A and H computed in SALSA?

- $H$ : SALSA's hub matrix, consists of the nonzero rows and columns of  $L_r L_c^T$
- $A$ : SALSA's authority matrix, consists of the nonzero rows and columns of  $L_c^T L_r$

$$L_r L_c^T = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 5 & 6 & 10 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 5 \\ 6 \\ 10 \end{matrix} & \begin{pmatrix} \frac{5}{12} & 0 & \frac{1}{6} & \frac{1}{2} & \frac{1}{6} & \frac{1}{6} \\ 0 & 1 & 0 & 0 & 0 & 0 \\ \frac{1}{3} & 0 & \frac{1}{3} & 0 & 0 & \frac{1}{3} \\ 0 & 0 & 0 & 0 & 0 & 0 \\ \frac{1}{4} & 0 & 0 & \frac{1}{2} & \frac{1}{4} & 0 \\ \frac{1}{3} & 0 & \frac{1}{3} & 0 & 0 & \frac{1}{3} \end{pmatrix} \end{matrix}$$

$$L_c^T L_r = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 5 & 6 & 10 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 5 \\ 6 \\ 10 \end{matrix} & \begin{pmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & \frac{1}{2} & \frac{1}{4} & \frac{1}{4} & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & \frac{1}{6} & 0 & \frac{5}{6} & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{pmatrix} \end{matrix}$$



# Other Algorithms

How are A and H computed in SALSA?

$$H = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 6 & 10 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 6 \\ 10 \end{matrix} & \begin{pmatrix} \frac{5}{12} & 0 & \frac{1}{6} & \frac{1}{6} & \frac{2}{12} \\ 0 & 1 & 0 & 0 & 0 \\ \frac{1}{3} & 0 & \frac{1}{3} & 0 & \frac{1}{3} \\ \frac{1}{4} & 0 & 0 & \frac{3}{4} & 0 \\ \frac{1}{3} & \frac{1}{3} & 0 & 0 & \frac{1}{3} \end{pmatrix} \end{matrix}$$

$$A = \begin{matrix} & \begin{matrix} 1 & 3 & 5 & 6 \end{matrix} \\ \begin{matrix} 1 \\ 3 \\ 5 \\ 6 \end{matrix} & \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & \frac{1}{2} & \frac{1}{4} & \frac{1}{4} \\ 0 & \frac{1}{2} & \frac{1}{2} & 0 \\ 0 & \frac{1}{6} & 0 & \frac{5}{6} \end{pmatrix} \end{matrix}$$

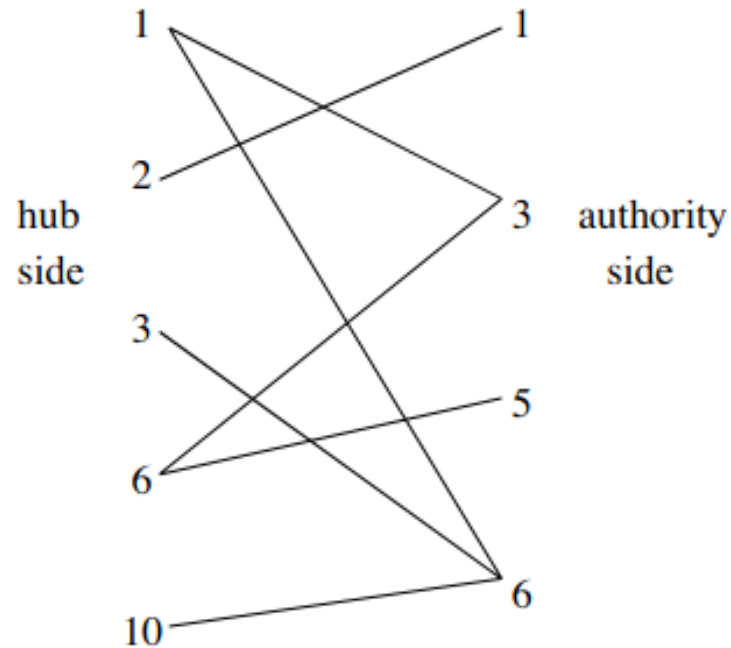
# Other Algorithms

## The Power Method

- If the bipartite graph  $G$  is *connected*, then  $A$  and  $H$  are both irreducible Markov chains and  $\pi_h^T$ , the stationary vector of  $H$ , gives the hub scores for the query with neighborhood graph  $N$ , and  $\pi_a^T$  gives the authority scores.
- If  $G$  is not *connected*, then  $A$  and  $H$  contain multiple irreducible components. In this case, the global authority and hub scores must be pasted together from the stationary vectors for each individual irreducible component.

# Other Algorithms

## The Power Method Example



- Because  $G$  is **not connected**,  $A$  and  $H$  contain multiple connected components.

- $H$  contains two connected components:

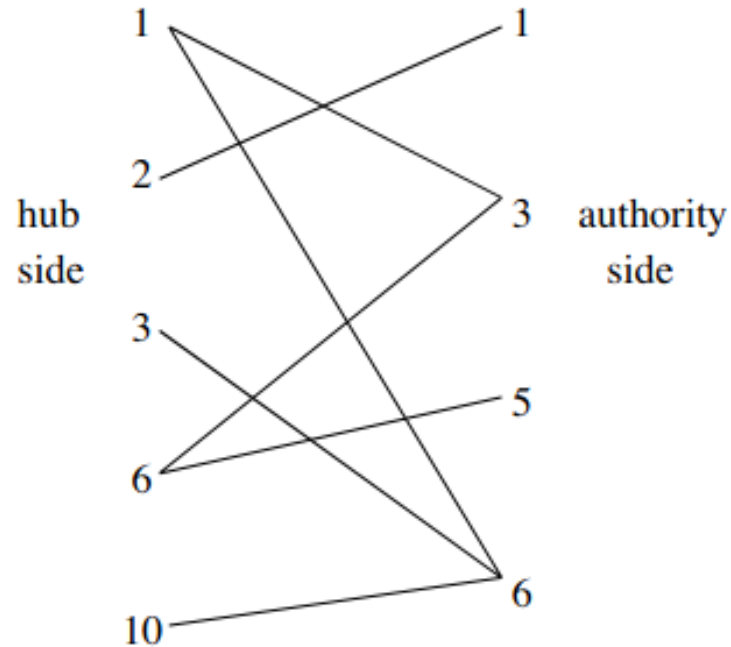
$$C = \{2\} \quad D = \{1, 3, 6, 10\}$$

- $A$  contains two connected components:

$$E = \{1\} \quad F = \{3, 5, 6\}$$

# Other Algorithms

## The Power Method Example



- The stationary vectors for the two irreducible components of  $H$  are:

$$\pi_h^{T(C)} = (1) \quad \pi_h^{T(D)} = \left(\frac{1}{3} \quad \frac{1}{6} \quad \frac{1}{3} \quad \frac{1}{6}\right)$$

- The stationary vectors for the two irreducible components of  $H$  are:

$$\pi_a^{T(E)} = (1) \quad \pi_a^{T(F)} = \left(\frac{1}{3} \quad \frac{1}{6} \quad \frac{1}{2}\right)$$

# Other Algorithms

## The Power Method Example

- Now, we can join the two components together for each matrix.
- We must multiply each entry in the vector by its appropriate weight.

# Other Algorithms

## The Power Method Example

- Thus the global Hub vector is:

$$\pi_h^T = \left( \frac{4}{5} \times \frac{1}{3} \quad \frac{1}{5} \times 1 \quad \frac{4}{5} \times \frac{1}{6} \quad \frac{4}{5} \times \frac{1}{3} \quad \frac{4}{5} \times \frac{1}{6} \right)$$

$$\pi_h^T = (0.2667 \quad 0.2 \quad 0.1333 \quad 0.2667 \quad 0.1333)$$

# Other Algorithms

## The Power Method Example

- And the global Authority vector is:

$$\pi_a^T = \left( \frac{1}{4} \times 1 \quad \frac{3}{4} \times \frac{1}{3} \quad \frac{3}{4} \times \frac{1}{6} \quad \frac{3}{4} \times \frac{1}{2} \right)$$

$$\pi_a^T = (0.25 \quad 0.25 \quad 0.125 \quad 0.375)$$

# Other Algorithms

## Advantages and Disadvantages of SALSA

### Advantages

- Not affected as much my topic drift like HITS.
- Less affected susceptible to spamming.
- Dual rank (Authority and Hubs).

### Disadvantages

- Query-dependence.
- Query time evaluation.



# Artificial Intelligence Applied to the Web

## Chapter 3 Part 2 - Web Page Rank Algorithms

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