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# Data handling

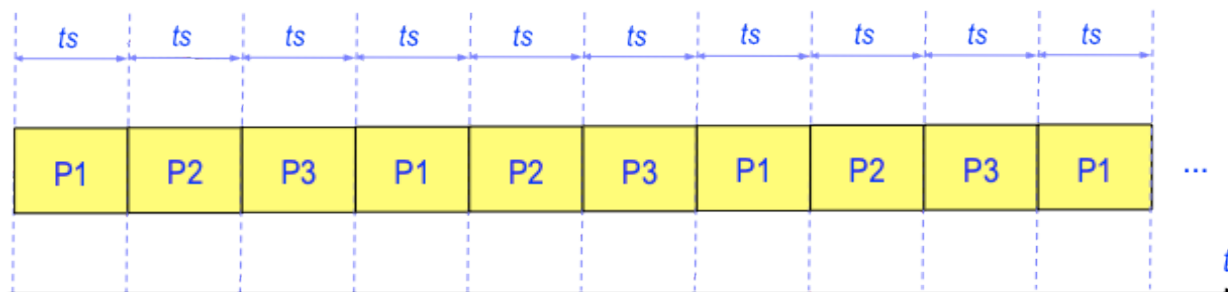
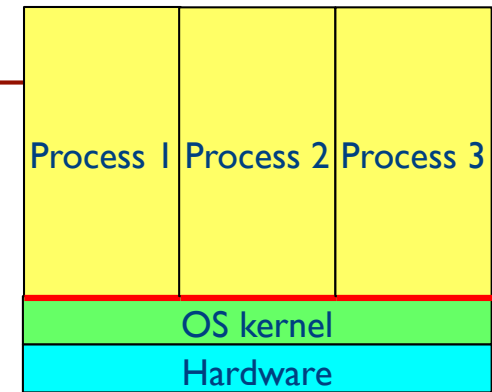
# Concurrency and tasking

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# Processes and Tasks/threads

# Process

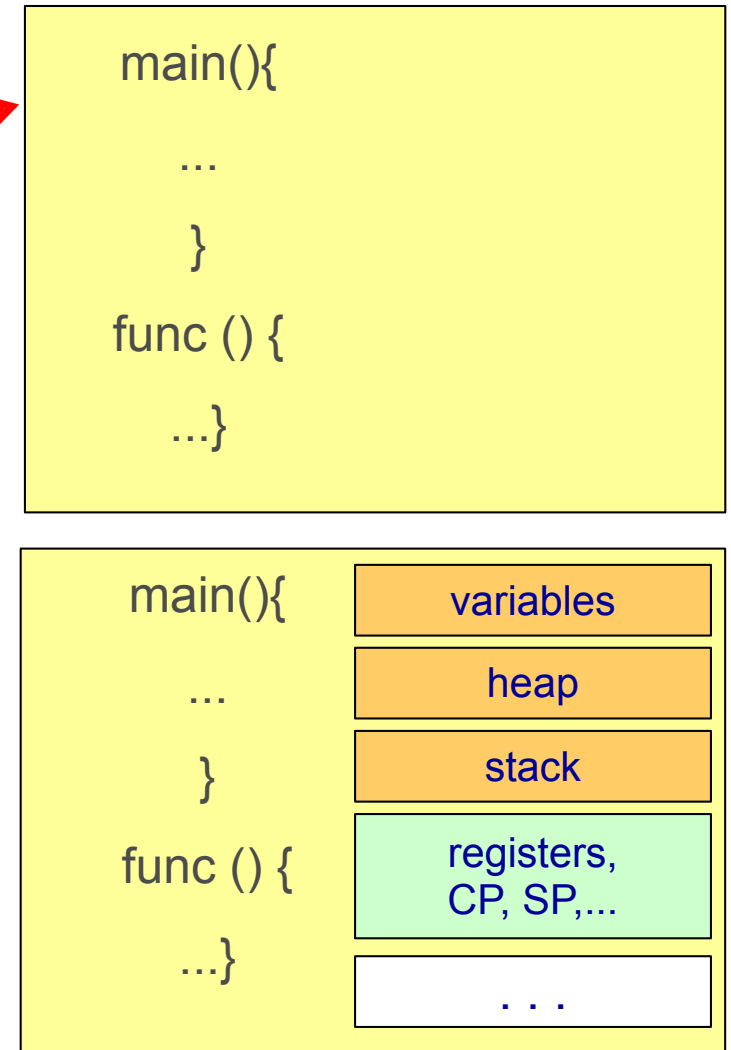
- Most operating systems support **multiprogramming**
- A **process: a program in execution**
  - Needs resources: CPU time, memory, files, E/S devices
- **Processes are isolated:**
  - A process cannot access directly to another
  - Interprocess: invoking OS system calls
- The **OS multiplex process for using the CPU time**
  - Processes creating and dispatching are time consuming



# Process management

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- **Process: program running**
- **Program: code and y static data**
- **Process image in an OS**
  - ▶ Code
  - ▶ Process state:
    - Memory area: values of variables, stack, heap, etc.
    - Registers: program counter, stack pointer
    - . . .



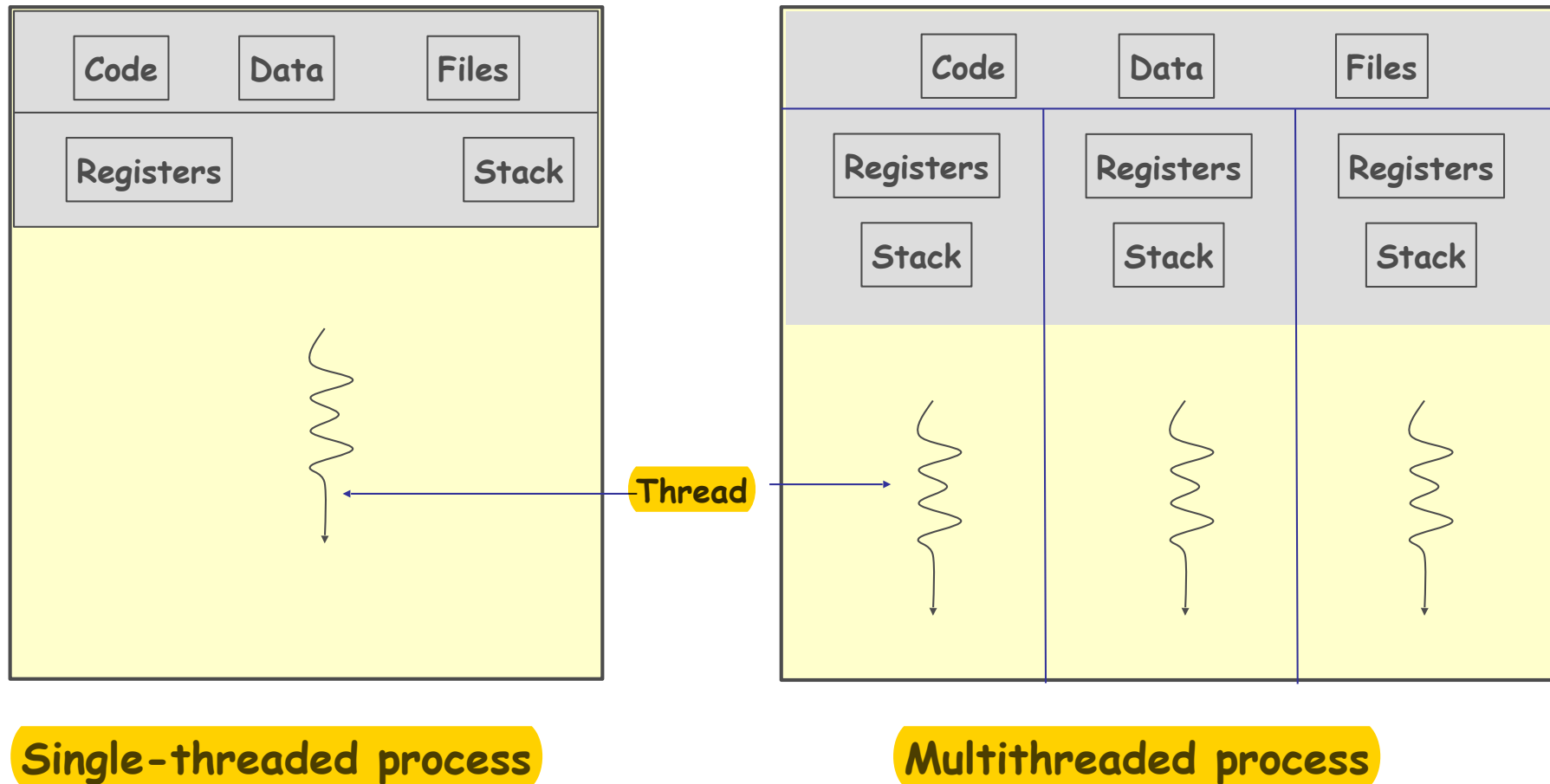
# Thread

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- Most operating systems support multithreading
- A process can include a set of running threads
  - ▶ Threads are running in concurrency
  - ▶ Thread state: registers, CP, SP
  - ▶ Threads in a process share the additional resources
    - Are less safe than processes.
- Threads creating and multiplexing are efficient
- Embedded systems: usual to have only threads

# Processes/Threads

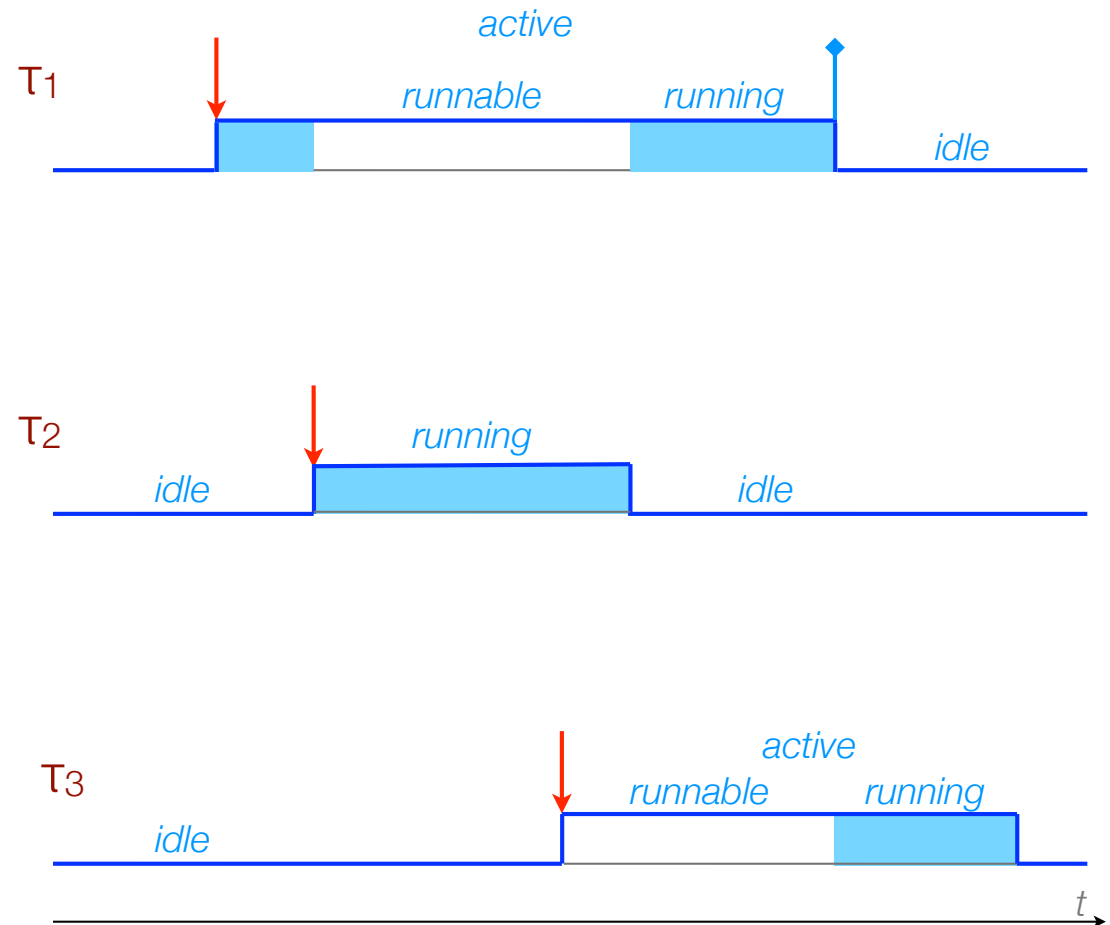
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# Concurrent tasks

# Concurrent tasks

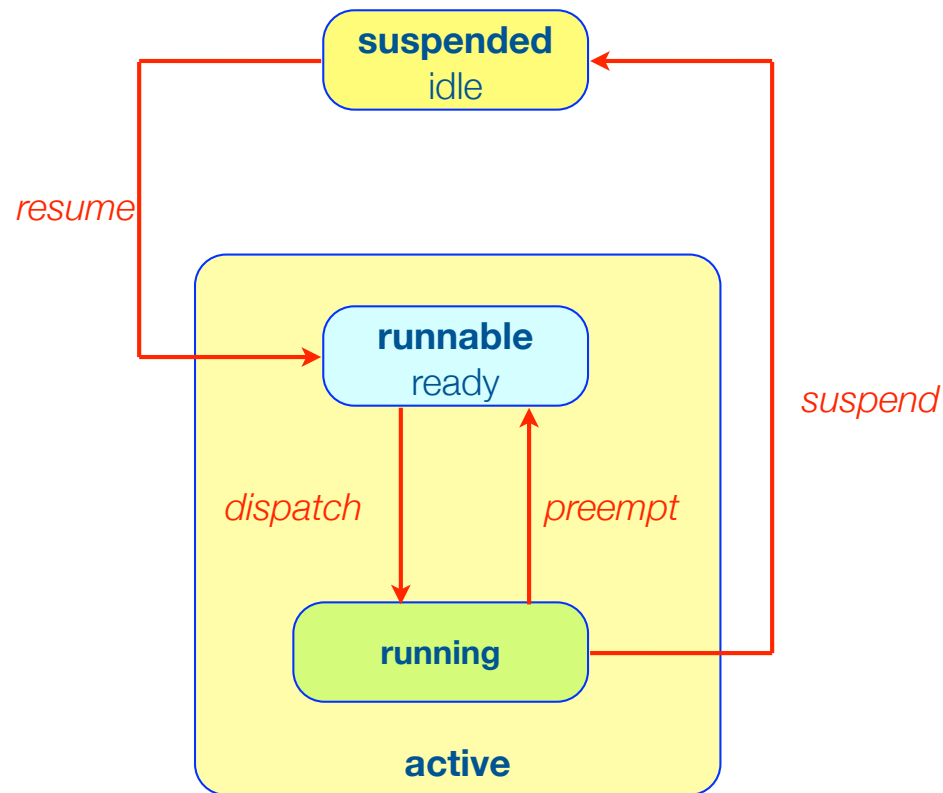
- Tasks are most often implemented using *threads*
  - ▶ concurrent flows of execution that progress asynchronously
  - ▶ processor time is multiplexed among active threads
  - ▶ requires operating system or run-time system support
    - e.g. POSIX threads, Ada tasks





# Task states

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- A task may be in different states
- Runnable tasks are dispatched for execution according to a scheduling method
- Scheduling is implemented by the OS kernel or the RTS

# Tasks in Ada

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- Ada tasks are program units
  - but not compilation units
  - must be compiled within packages
- Composed of specification and body
  - specification contains the visible interface
    - often empty
  - body contains the executable statements
- Tasks belong to a task type
  - explicit or anonymous

# Example

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```
package Temperature_Control is
    task type Temperature_Controller; -- specification
end Temperature_Control;

package body Temperature_Control is
    task body Temperature_Controller is -- body
        ...                             -- local declarations
    begin
        ...                             -- sequence of statements
    end Temperature_Controller;

    task Driver;                        -- single task object
    task body Driver is                -- body
        ...
    begin
        ...
    end Driver;
end Temperature_Control;
```

# Example (continued)

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```
with Temperature_Control;  
procedure Control is          -- main procedure  
    use Temperature_Control;  
    TC1, TC2 : Temperature_Controller;  
begin                          -- TC1, TC2, PC, and Driver run in parallel  
    null;                      -- the main procedure does nothing  
end Control;                  -- and does not terminate until TC1, TC2,  
                              -- and Driver complete
```

# Shared data and synchronisation

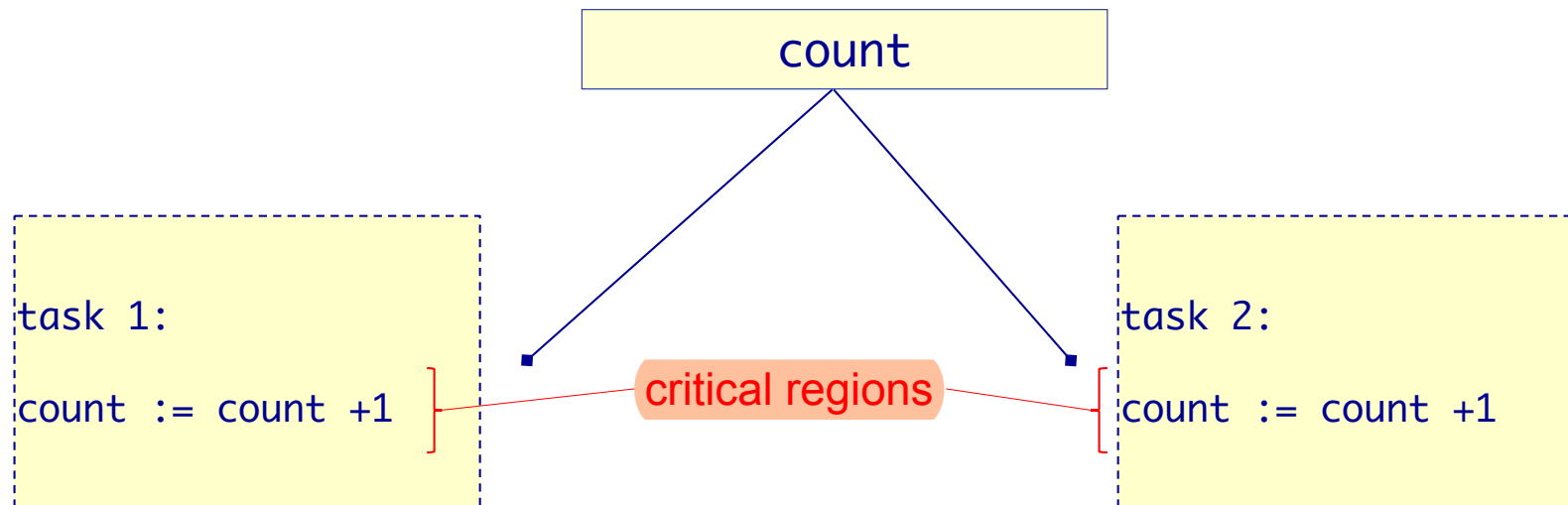
# Shared data

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- Concurrent access to shared data may lead to race conditions
  - ▶ the value of the data may depend on the relative speeds of concurrent threads
- In order to prevent race conditions, access to shared data must be carried out in mutual exclusion

# Example: shared counter

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# Shared counter (1)

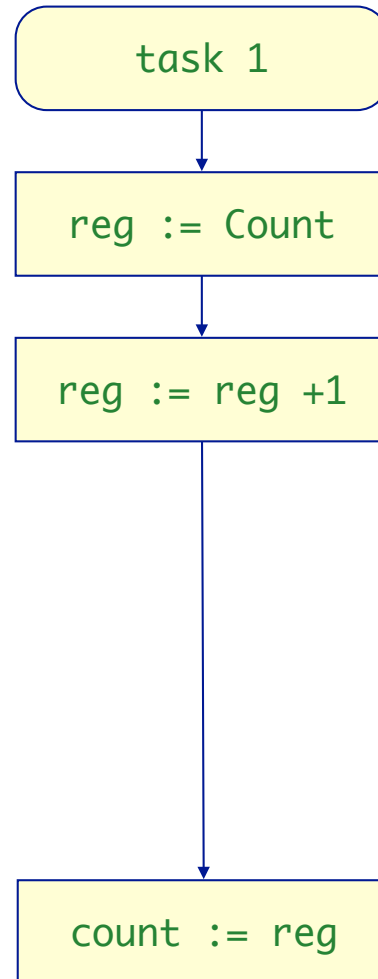
count == 0

reg ← 0

reg ← 1

count ← 1

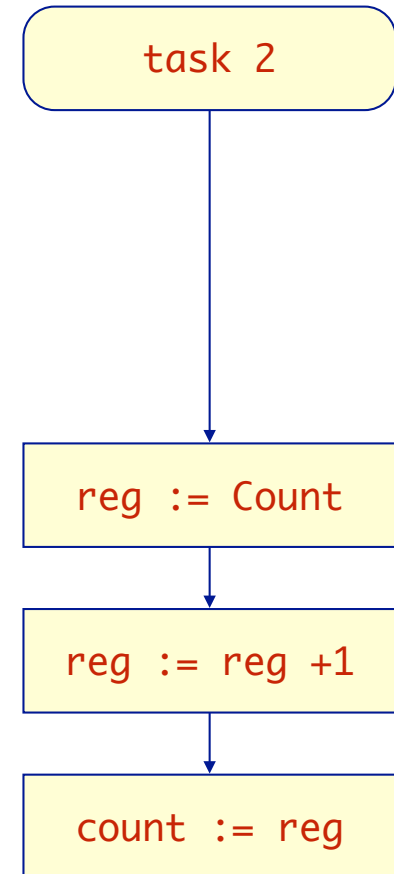
count == 1



reg ← 0

reg ← 1

count ← 1





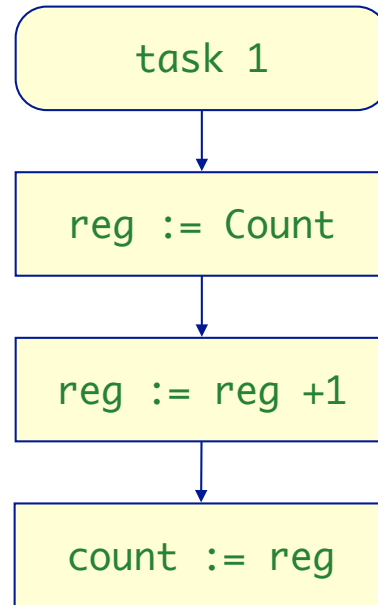
# Shared counter (2)

count == 0

reg ← 0

reg ← 1

count ← 1

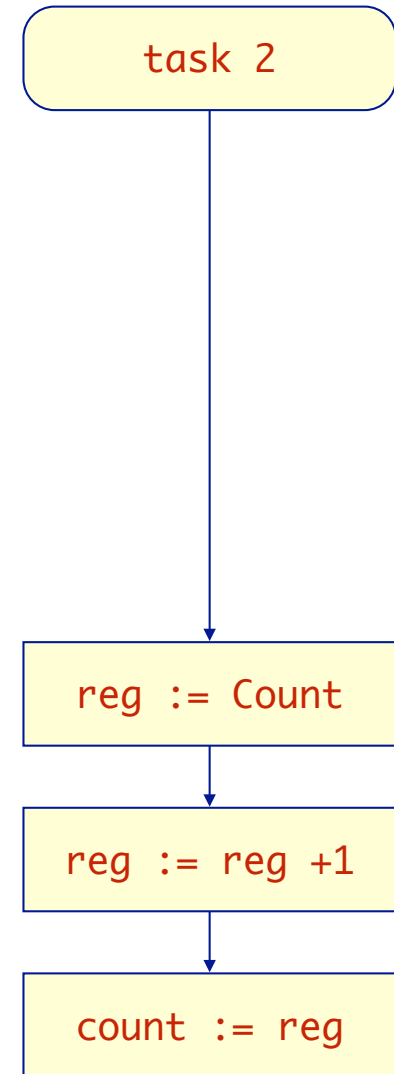


count == 2

reg ← 1

reg ← 2

count ← 2



# Race condition

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- The result may be 1 or 2
  - ▶ depending on the respective speeds of task execution
    - "who runs faster"
- Such a program is incorrect
  - ▶ each execution is different
    - behaviour is not deterministic
  - ▶ the outcome of the program cannot be verified
- Solution: mutual exclusion
  - ▶ while one task is using the shared data, no other task can do
  - ▶ critical regions are mutually exclusive

# Protected objects in Ada

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- Protected objects encapsulate shared data and operations that can be carried out on them
- Mutual exclusion is automatically provided
- POs are instances of (possibly anonymous) protected types
- Protected types are declared in two separate parts
  - ▶ specification: visible interface of the protected objects
    - name of the type and, optionally, parameters
    - structure of protected data
    - specifications of protected operations
  - ▶ body: implementation of the protected operations

# Example

```
type Operating_Mode is (Off, Safe, Nominal, Latency);
protected type Mode_Manager (Initial_Mode : Operating_Mode) is
    function Current_Mode return Operating_Mode;
    procedure Set_Mode (New_Mode : Operating_Mode);
private
    Mode : Operating_Mode := Initial_Mode;
end Mode_Manager;

protected body Mode_Manager is
    function Current_Mode return Operating_Mode is
    begin
        return Mode;
    end Current_Mode;
    procedure Set_Mode (New_Mode : Operating_Mode) is
    begin
        Mode := New_Mode;
    end Set_Mode;
end Mode_Manager;
```

# Example (continued)

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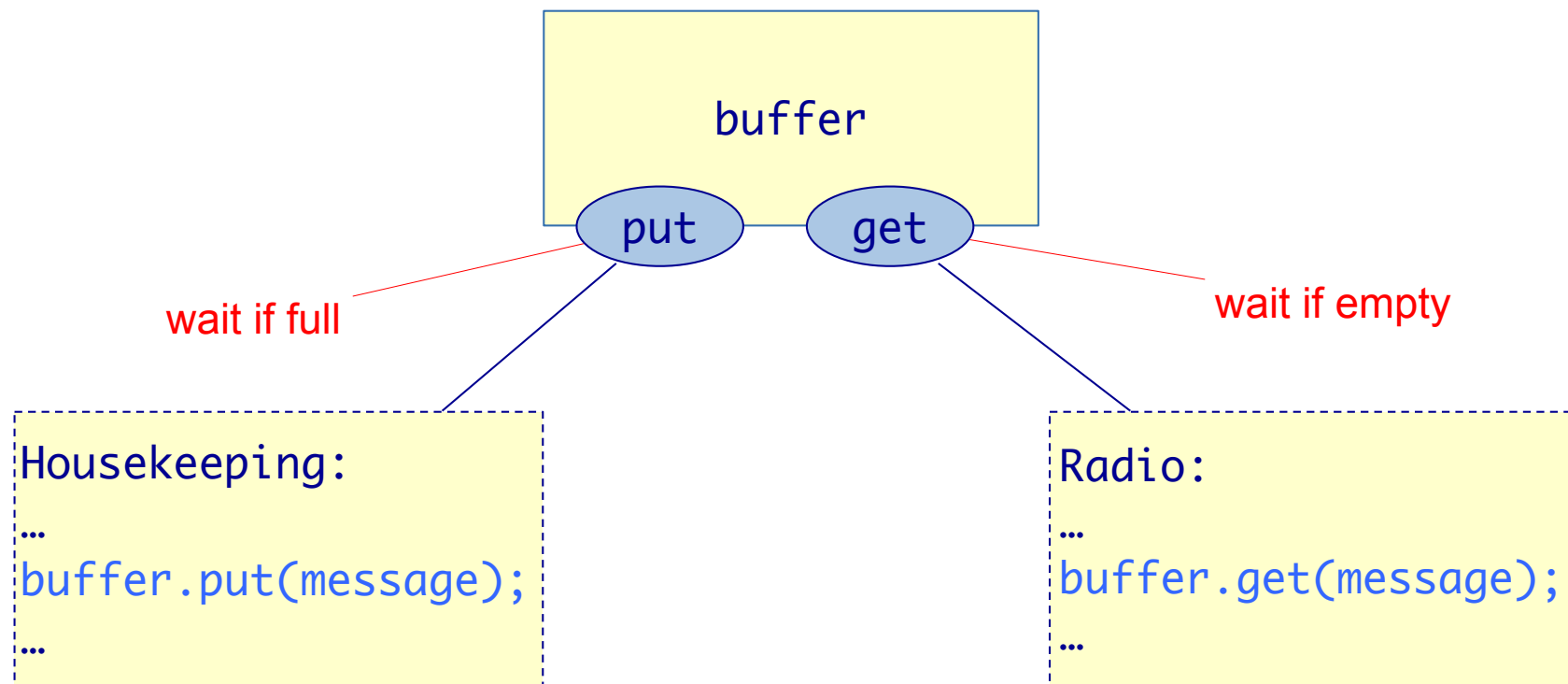
```
System_Mode : Mode_Manager (Off);

case System_Mode.Current_Mode is
  when Off => null;
  when Safe => ...;
  when Nominal => ...;
  when Latency => ...;
end case;

if (Battery_Low) then
  System_Mode.Set_Mode(Safe);
end if;
```

# Conditional synchronisation

- Some times a task must wait for another task to fulfil some condition
  - e.g. message buffer



# Conditions and entry barriers

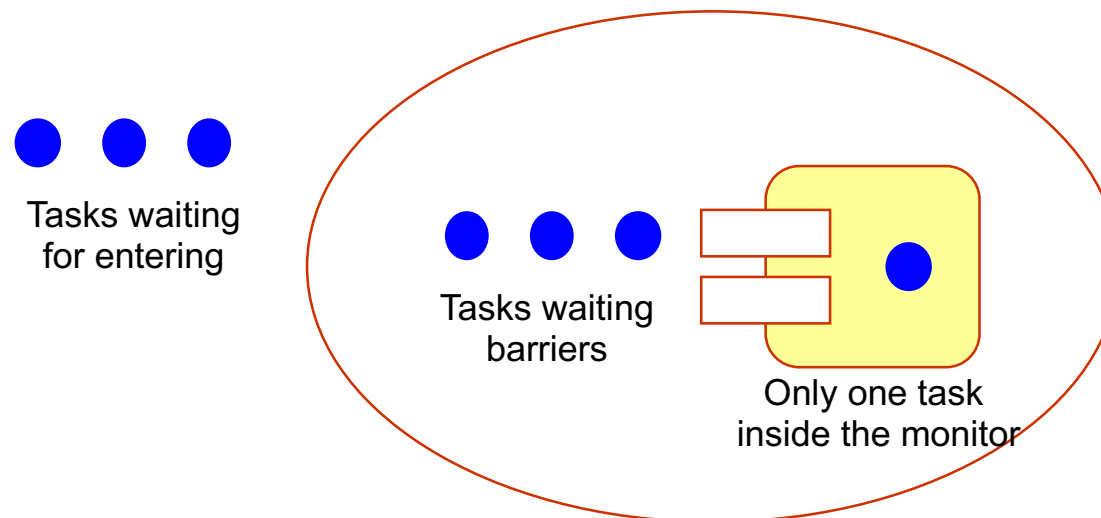
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- Conditional synchronisation is achieved by using special operations called **entries**
- Each entry has a **barrier** (a Boolean expression)
  - ▶ a task calling an entry with a closed (false) barrier is suspended
  - ▶ whenever another operation on the same object ends the barriers are re-evaluated and opened if true
  - ▶ one task waiting on a previously closed barrier is resumed
  - ▶ the resumed task executes the entry body

# Mutual exclusion and barriers

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- Tasks waiting on a barrier have priority for accessing the monitor
  - ▶ This approach avoids race conditions
- This model is called eggshell





# Example: message buffer

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```
type Message is ...;
Size : constant Positive := 32;
type Index is mod Size;
type Message_Queue is array (Index) of Message;

protected type Message_Buffer is    -- specification
    entry Put(M : in Message);
    entry Get(M : out Message);
private
    First    : Index := Index'First;
    Last     : Index := Index'First;
    Number   : Natural := 0;
    Store    : Message_Queue;
end Message_Buffer;
```

# Example: message buffer (continued)

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protected body Message\_Buffer is

entry Put(M : in Message) when Number < Size is  
begin

    Store(Last) := M;

    Last := Last + 1;

    Number := Number + 1;

end Put;

entry Get(M : out Message) when Number > 0 is  
begin

    M := Store (First);

    First := First + 1;

    Number := Number - 1;

end Get;

end Message\_Buffer;

