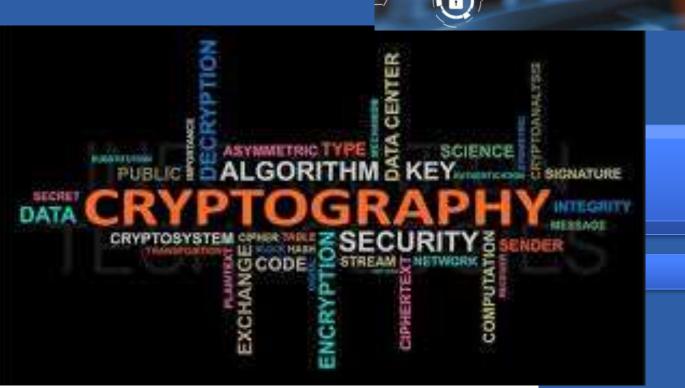
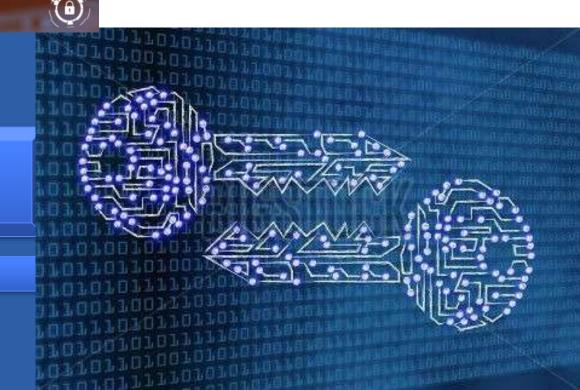




Tema 3 Cifrado Simétrico



Criptografía



Criptografía Clásica



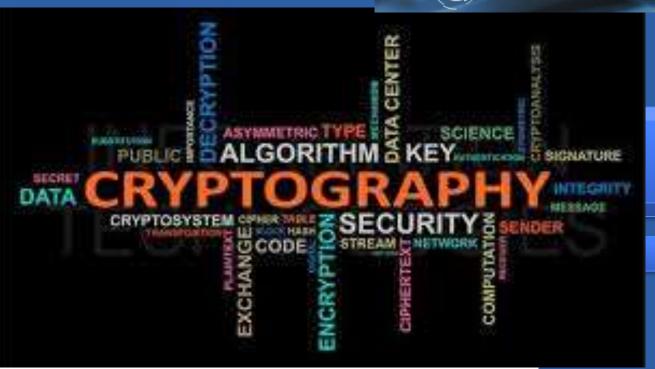
- Introducción
- Algoritmos de cifrado simétrico de flujo
- Algoritmos de cifrado simétrico de bloque
- Modos de operación

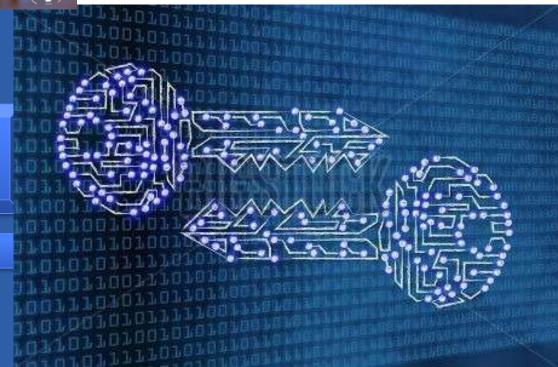






Tema 3.1 Introducción al Cifrado Simétrico





Criptografía Moderna



Podemos clasificar los algoritmos según varias dimensiones:

- Por el tipo de clave usada
 - Simétricos
 - Asimétricos
- Por la cantidad de elementos tomados a la vez:
 - Cifradores de flujo
 - Cifradores de bloque

.

Criptografía Moderna



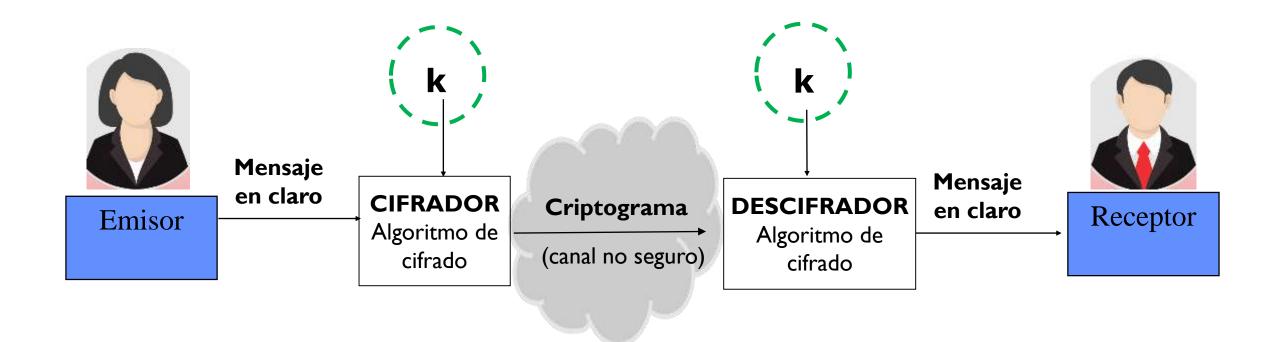
- Por el tipo de clave usada
 - -Simétricos
 - Asimétricos

.

Modelo de Criptosistema



SIMÉTRICOS: Misma clave (k) para cifrar y descifrar

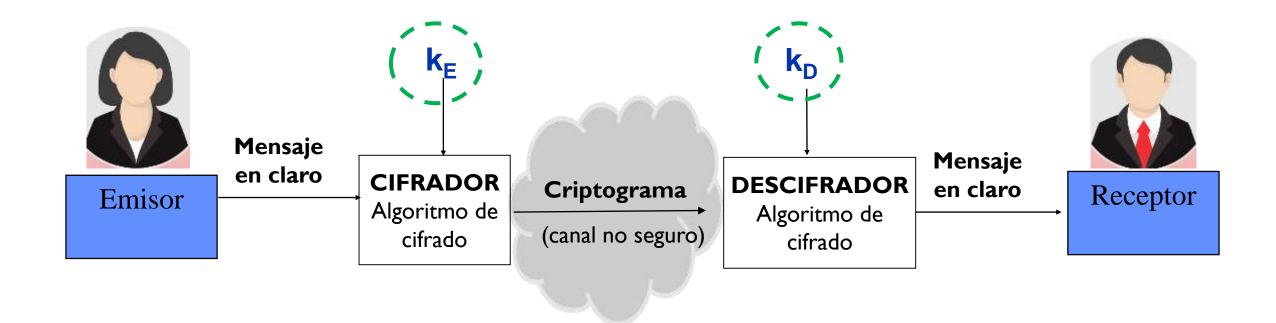


La clave se acuerda y comparte entre las en modo secreto

Modelo de Criptosistema



ASIMÉTRICOS: Una clave (k_E) para cifrar y otra (k_D) para descifrar



Base de los sistemas de clave pública

Criptografía Moderna



Por la cantidad de elementos tomados a la vez:

- Cifradores de flujo

- Cifradores de bloque

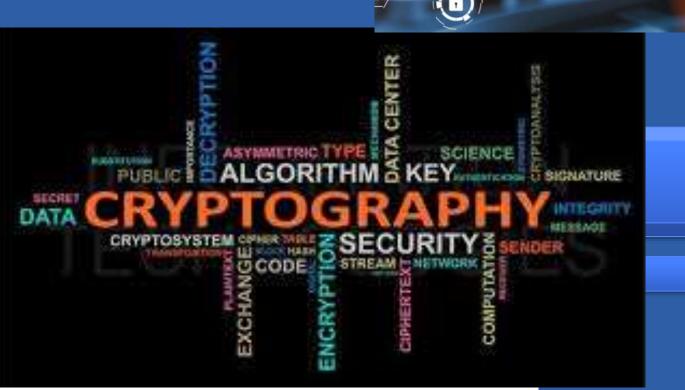
Criptografía

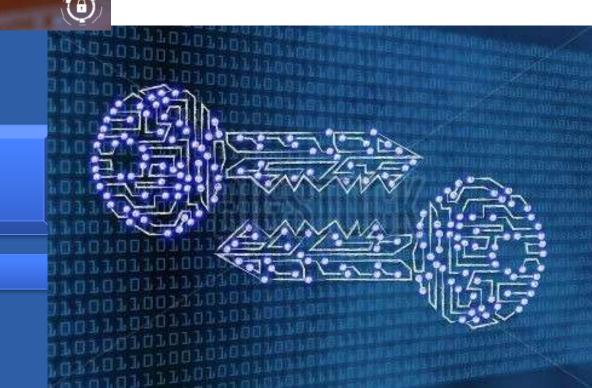




Tema 3.2
Algoritmos de cifrado simétrico de flujo

Criptografía







- Dividen el mensaje en caracteres o símbolos
 - M = m1 m2 m3 m4...
- El flujo de cifrado (clave) también se divide en caracteres o símbolos
 - K = k1 k2 k3 k4...
- Cada carácter se cifra con el carácter correspondiente de la clave:
 - c1 = E(m1, k1)
 - c2 = E(m2, k2)
 - c3 = E(m3, k3)
 - •
- Ejemplos: Vernam, RC4, etc.



EJERCICIO: Cifrador Vernam

- Mensaje = "HOLA"
- Clave = "key"
- Tamaño de símbolo = 1 byte (Codificación ASCII)
- Texto cifrado = ??

•



EJERCICIO: Cifrador Vernam

Mensaje = "HOLA"

$$\rightarrow$$
 "H" = HEX(48) = 0100 1000

$$\succ$$
 "O" = HEX(4F) = 0100 1111

$$\succ$$
 "L" = HEX(4C) = 0100 1100

$$F$$
 "A" = HEX(41) = 0100 0001

Clave = "KEY"

$$\triangleright$$
 "e" = HEX(65) = 0110 0101

$$\rightarrow$$
 "y" = HEX(79) = 0111 1001



EJERCICIO: Cifrador Vernam

Mensaje = "HOLA"

$$\rightarrow$$
 "H" = HEX(48) = 0100 1000

$$\succ$$
 "O" = HEX(4F) = 0100 1111

$$\succ$$
 "L" = HEX(4C) = 0100 1100

$$F$$
 "A" = HEX(41) = 0100 0001

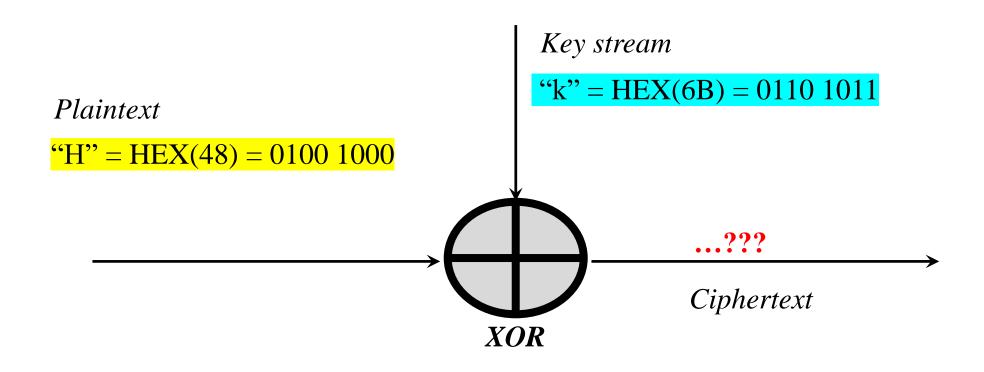
Clave = "KEY"

$$\rightarrow$$
 "k" = HEX(6B) = 0110 1011

$$\triangleright$$
 "e" = HEX(65) = 0110 0101

$$\rightarrow$$
 "y" = HEX(79) = 0111 1001







Generación del keystream

- En la práctica se emplean generadores pseudo-aleatorios
 - > Semilla inicial (clave compartida, usada sólo una vez)
 - > Algoritmo determinista (misma entrada genera misma salida)
 - Útil para que el transmisor y el receptor se entiendan
- Ejemplo: LFSR

$$f(x) = C_n x^n + C_{n.1} x^{n-1} + \dots + C_2 x^2 + C_1 x + 1$$

$$Seed = S_n$$
, $S_{n.1}$, ..., S_2 , S_1

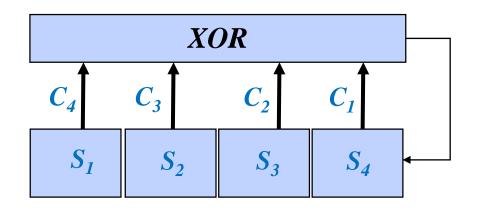


LFSR (Linear Feedback Shift Register)

- Expresión general
 - (para orden de polinomio n=4)

$$f(x) = C_4 x^4 + C_3 x^3 + C_2 x^2 + C_1 x + 1$$

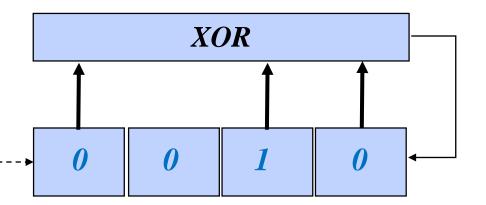
$$Seed = (S1, S2, S3, S4)$$



Ejemplo:

$$f(x) = 1 \cdot x^{4} + 0 \cdot x^{3} + 1 \cdot x^{2} + 1 \cdot x + 1$$

$$Seed = 0 \ 0 \ 1 \ 0$$





PRESTACIONES

VENTAJAS:

- Muy rápidos
- Ideales para prestaciones que requieran tiempo real
 - > Ejemplo: streaming de video / audio
- -Los errores de transmisión no se propagan a otros símbolos

INCONVENIENTES:

- Mala difusión de la información
 - > La información de cada símbolo se traslada íntegramente a su símbolo cifrado
- Seguridad y gestión de la clave
 - > (puramente aleatoria, más larga que el texto a cifrar, sólo usada una vez...)





Home

Profile 1 (SW)

HC-128 Rabbit Salsa20/12 SOSEMANUK

Profile 2 (HW)

Grain v1 MICKEY 2.0 Trivium

eSTREAM: the ECRYPT Stream Cipher Project

Welcome to the home page of eSTREAM, the ECRYPT Stream Cipher Project. The eSTREAM project was a multi-year effort, running from 2004 to 2008, to promote the design of efficient and compact stream ciphers suitable for widespread adoption. As a result of the project, a portfolio of new stream ciphers was announced in April 2008. The eSTREAM portfolio was revised in September 2008, and currently contains seven stream ciphers. This website is dedicated to ciphers in this final portfolio. For information on the eSTREAM project and selection process, including a timetable of the project and further technical background, please visit the original eSTREAM Project website.

The eSTREAM Portfolio

The short report from April 2008 discussing the initial portfolio (with eight stream ciphers) and the end of the eSTREAM project can be found here. The eSTREAM portfolio was revised in September 2008, following the announcement of cryptanalytic results against one of the original algorithms (see here). The portfolio is periodically revisited, as the algorithms mature: the first review of the eSTREAM portfolio was published in October 2009, and is available here; the second review from January 2012 can be found here.

The eSTREAM portfolio ciphers fall into two profiles. Profile 1 contains stream ciphers more suitable for software applications with high throughput requirements. Profile 2 stream ciphers are particularly suitable for hardware applications with restricted resources such as limited storage, gate count, or power consumption.

The eSTREAM portfolio contains the following ciphers:

Profile 1 (SW)	Profile 2 (HW)
HC-128	Grain v1
Rabbit	MICKEY 2.0
Salsa20/12	Trivium
SOSEMANUK	



Algoritmo RC4 (no entra en examen!)

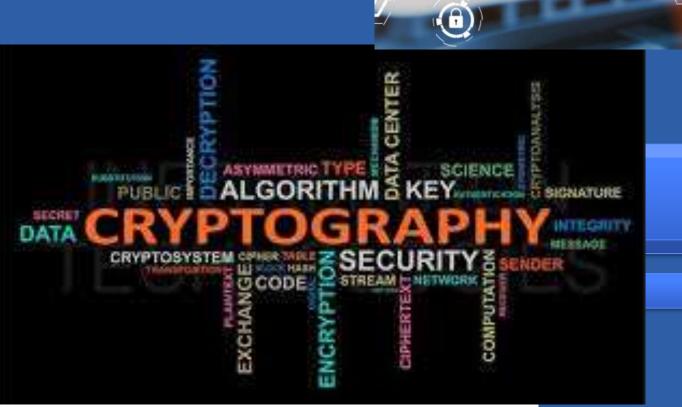


https://www.youtube.com/watch?v=G3HajuqYH2U

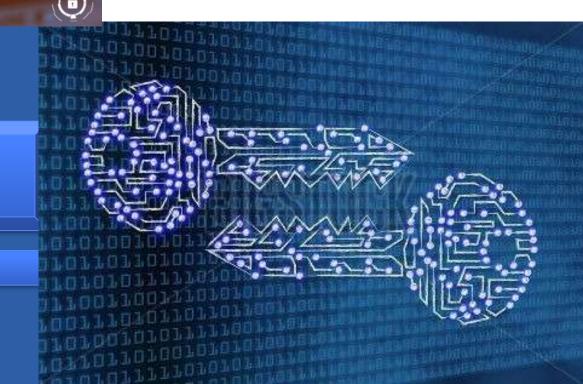




Tema 3.2 Algoritmos de cifrado simétrico de bloque



Criptografía





Dividen el mensaje en bloques de igual tamaño (*)

(*)*Habitualmente:* **64, 128** o **256 bits**

- M = M1 M2 M3 M4...
- Se utiliza la misma clave para el cifrar (y descifrar) todos los bloques
 - C1 = E(M1, K) \rightarrow M1 = D(C1, K)
 - C2 = E(M2, K) \rightarrow M2 = D(C2, K)
 - C3 = E(M3, K) \rightarrow M3 = D(C3, K)
 - •
- Ejemplos: DES, AES, etc.

-



Objetivos de seguridad (para dificultar criptoanálisis):

Difusión

- Que la estructura del mensaje quede disipada en la estructura del texto cifrado
 - Dificultar análisis de frecuencias, etc.
 - Cada modificación de 1 bit del mensaje tiene que afectar a muchos bits del texto cifrado, y viceversa.
- Se consigue mediante funciones de permutación

Confusión

- Que la relación entre la clave y el texto cifrado sea tan compleja como sea posible
 - > Y no se infiera información de la clave a partir del texto cifrado.
- Se consigue mediante algoritmos de sustitución

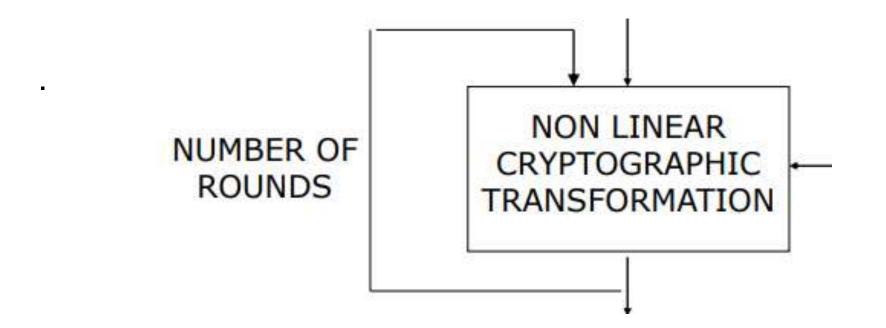


Esquema general

INPUT INITIAL TRANSFORMATION NON LINEAR SUBKEY NUMBER OF CRYPTOGRAPHIC **KEY** GENERATION ROUNDS TRANSFORMATION ALGORITHM FINAL TRANSFORMATION OUTPUT



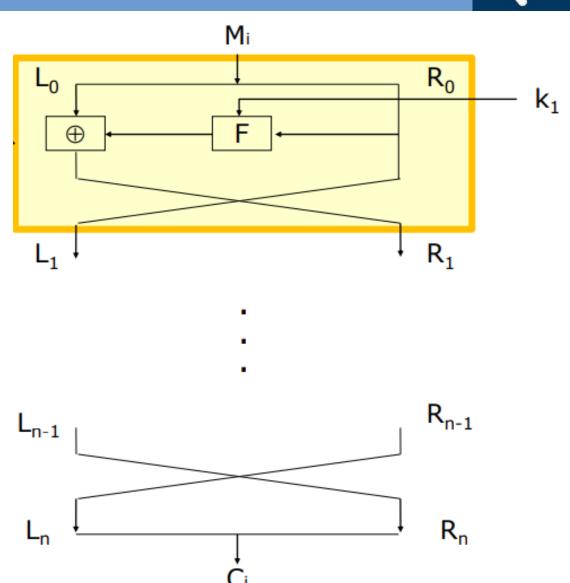
- Principales esquemas de transformación criptográfica no lineal:
 - Feistel (ejemplo: DES)
 - Esquema de Sustitución Permutación (ejemplo: AES)



Esquema de Feistel



- Procedimiento del esquema de Feistel:
- Divide el bloque M en dos mitades, L₀ y
 R₀
- Repite los siguientes pasos durante nondas. En cada ronda "i":
 - Aplica una función F sobre la mitad derecha (R_i) y la subclave de ronda k_i)
 - Realiza un XOR entre la salida de F y la mitad izquierda L_i
 - Intercambia la mitad derecha y la izquierda



Esquema de Feistel

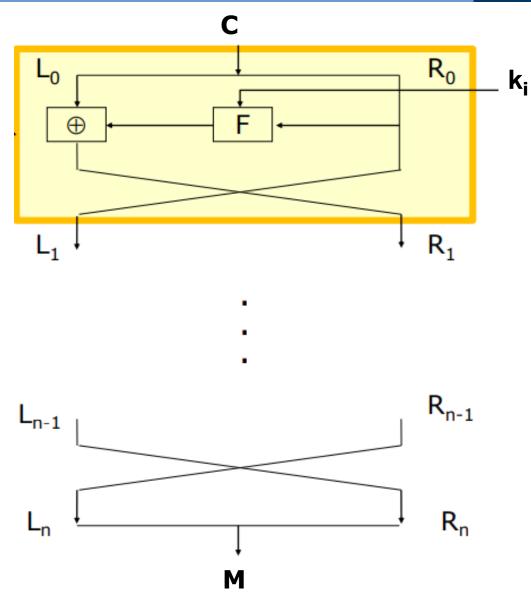


Para descifrar:

- 1. Se utiliza el mismo circuito que para cifrar
- 2. Pero las subclaves de ronda se utilizan en el orden inverso

$$\checkmark k_n \rightarrow k_{n-1} \rightarrow \dots k_1$$

 Ojo con el orden de la permutación inicial y final



Esquema de Feistel



Seguridad en esquema de Feistel basada en:

- Diseño de una buena función F
- 2. Diseño de una buen algoritmo de generación de subclaves

Parámetros típicos del esquema de Feistel:

- 1. Tamaño de bloque: cuanto más grande, mayor seguridad pero más lento
 - √ 64 bits o más.
- 2. Tamaño de clave: cuanto más grande, mayor seguridad pero más lento
 - √ 64 bits o más.
- 3. Número de rondas: cuanto más grande, mayor seguridad pero más lento
 - ✓ Valor típico: 16 rondas.





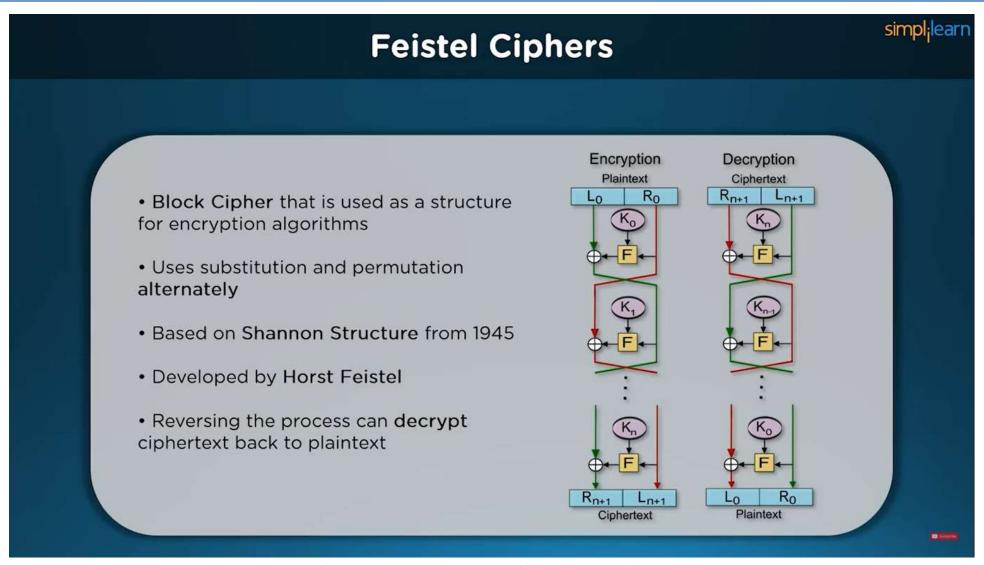
Tema 3.2.1 Algoritmos de cifrado simétrico de bloque: DES





Criptografía

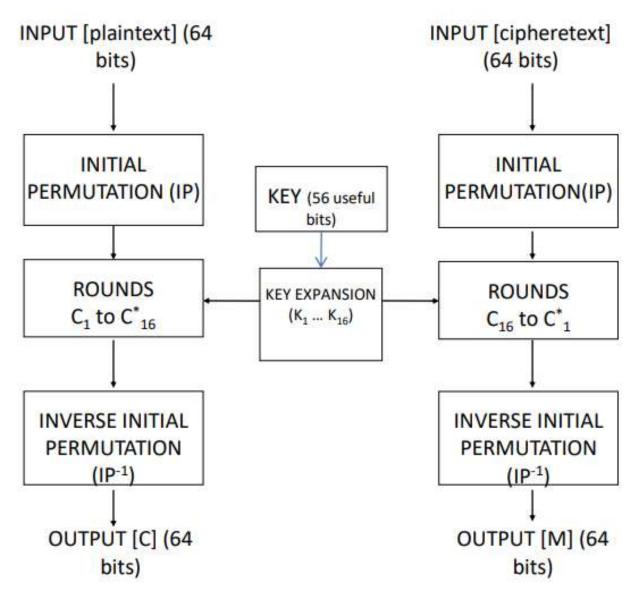




https://www.youtube.com/watch?v=S918rR4VdqQ



- Basado en esquema de Feistel:
 - 1. Tamaño de bloque: 64 bits
 - Tamaño de clave: 64 bits
 - 8 son de paridad
 - Clave inicial de 56 (64-8=56)
 - 3. Número de rondas: 16
 - Algoritmo de expansión de claves
 - > 16 Subclaves de 48 bits
 - Complejidad basada en:
 - 1. Sustituciones lineales
 - 2. Sustituciones no lineales
 - 3. Permutaciones





Especificación técnica

Federal Information Processing Standards Publication 46-3



1999 October 25

Announcing the

DATA ENCRYPTION STANDARD

Federal Information Processing Standards Publications (FIPS PUBS) are issued by the National Institute of Standards and Technology after approval by the Secretary of Commerce pursuant to Section 5131 of the Information Technology Management Reform Act of 1996 (Public Law 104-106), and the Computer Security Act of 1987 (Public Law 100-235).

- Name of Standard. Data Encryption Standard (DES).
- Category of Standard. Computer Security, Cryptography.
- 3. Explanation. The Data Encryption Standard (DES) specifies two FIPS approved cryptographic algorithms as required by FIPS 140-1. When used in conjunction with American National Standards Institute (ANSI) X9.52 standard, this publication provides a complete description of the mathematical algorithms for encrypting (enciphering) and decrypting (deciphering) binary coded information. Encrypting data converts it to an unintelligible form called cipher. Decrypting cipher converts the data back to its original form called plaintext. The algorithms described in this standard specifies both enciphering and deciphering operations which are based on a binary number called a key.

FIPS PUB 46-3

FEDERAL INFORMATION PROCESSING STANDARDS PUBLICATION

> Reaffirmed 1999 October 25

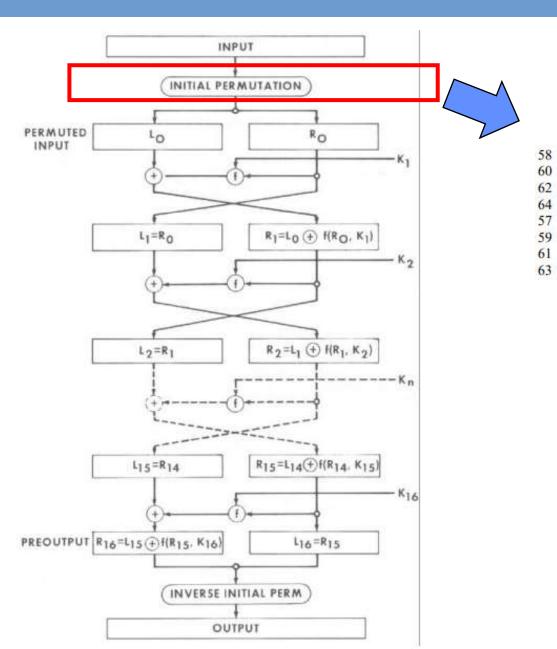
U.S. DEPARTMENT OF COMMERCE/National Institute of Standards and Technology

DATA ENCRYPTION STANDARD (DES)

CATEGORY: COMPUTER SECURITY SUBCATEGORY: CRYPTOGRAPHY



Permutación Inicial



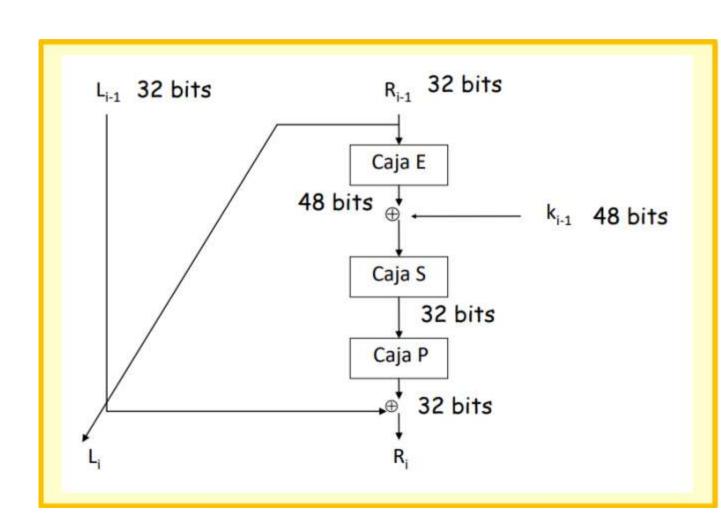
IP

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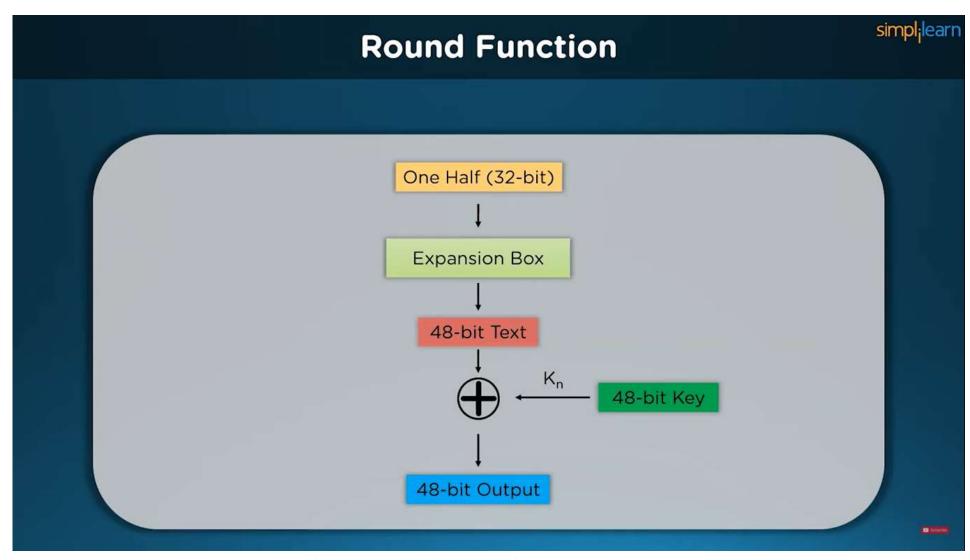
52



- Basado en esquema de Feistel:
 - 1. Tamaño de bloque: 64 bits
 - 2. Tamaño de clave: 64 bits
 - 8 son de paridad
 - Clave inicial de 56 (64-8=56)
 - 3. Número de rondas: 16
 - Algoritmo de expansión de claves
 - > 16 Subclaves de 48 bits



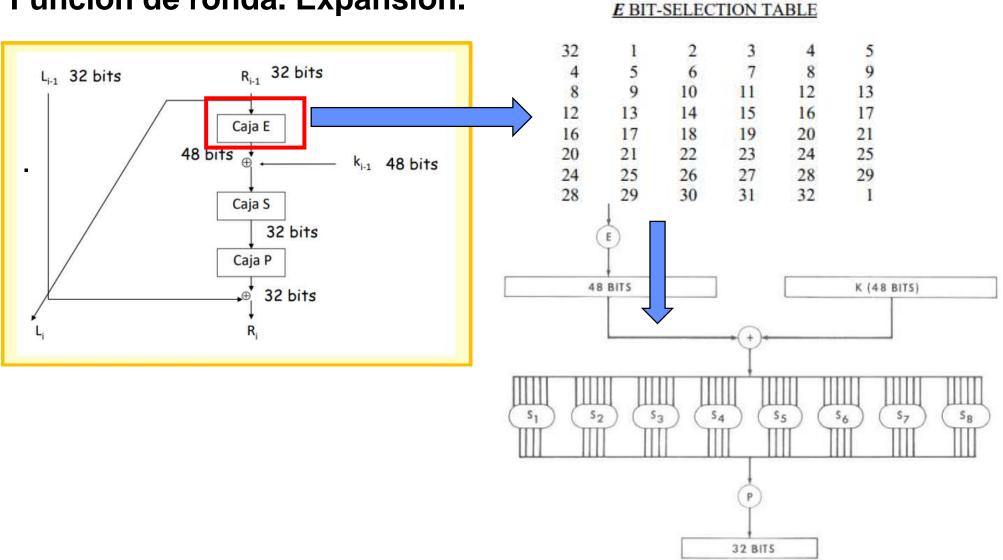




https://www.youtube.com/watch?v=S918rR4VdqQ

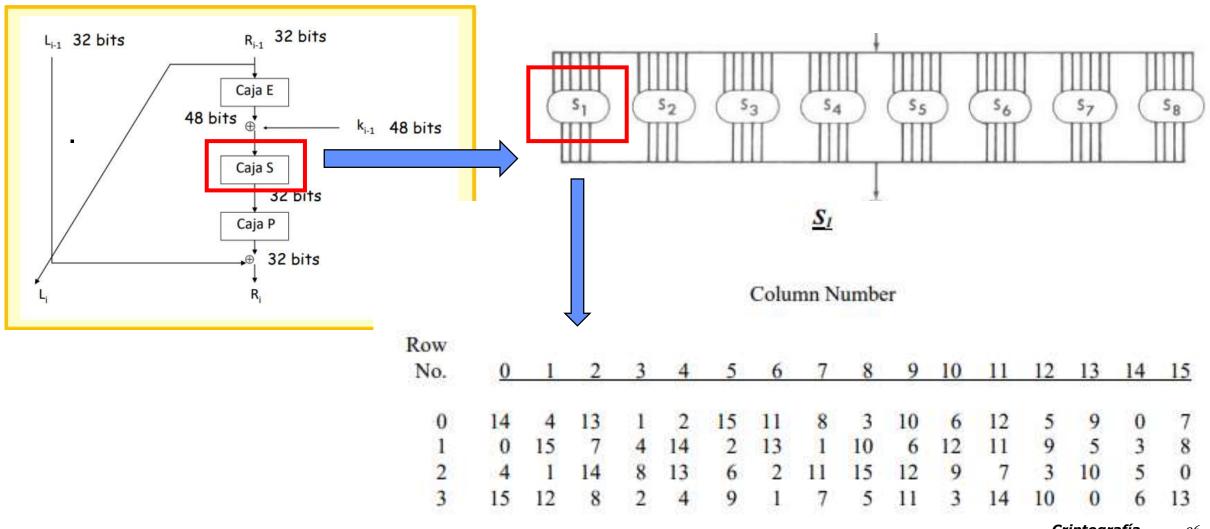


Función de ronda. Expansión.

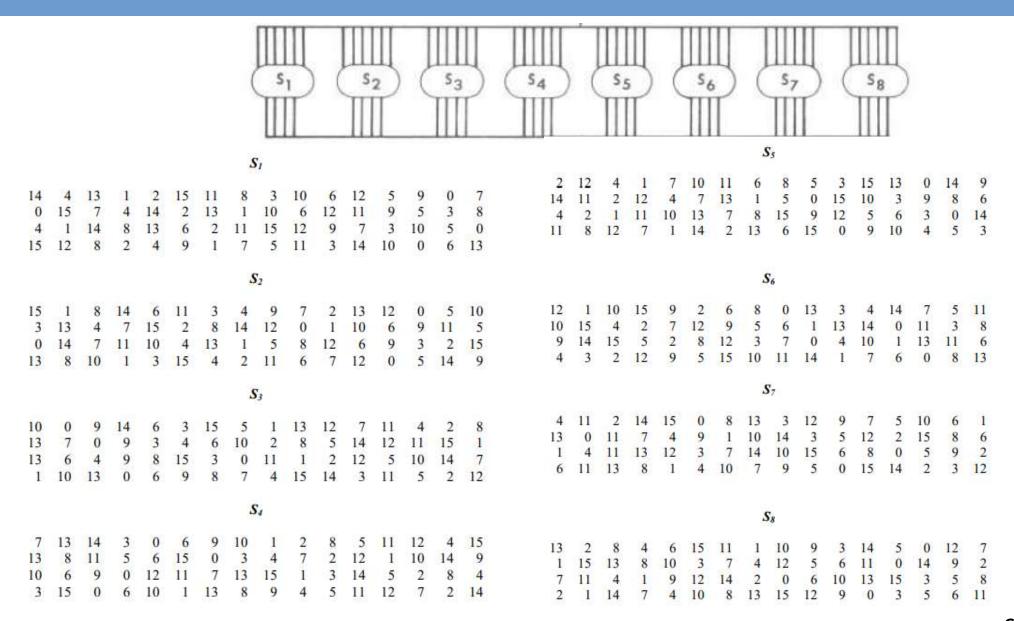




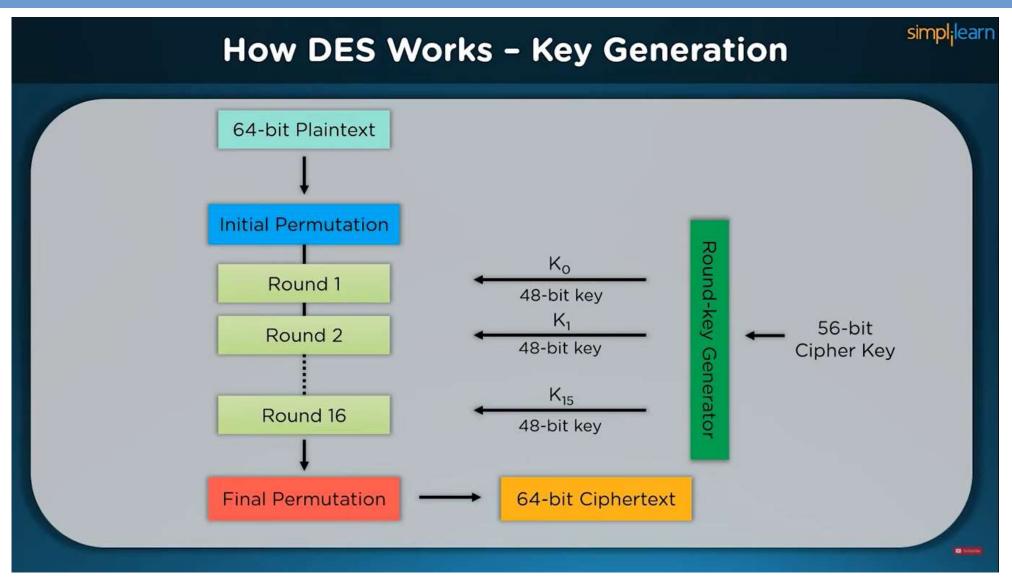
Función de ronda. Sustitución.







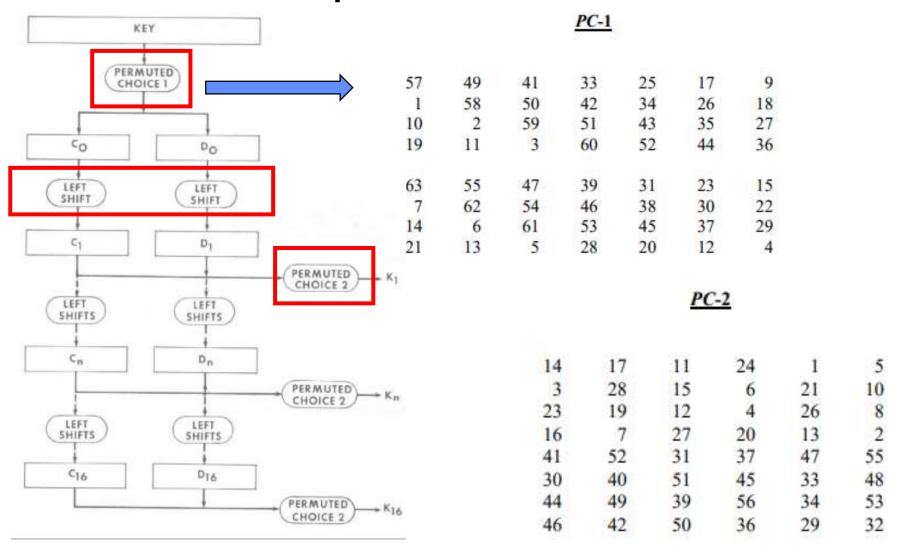




https://www.youtube.com/watch?v=S918rR4VdqQ

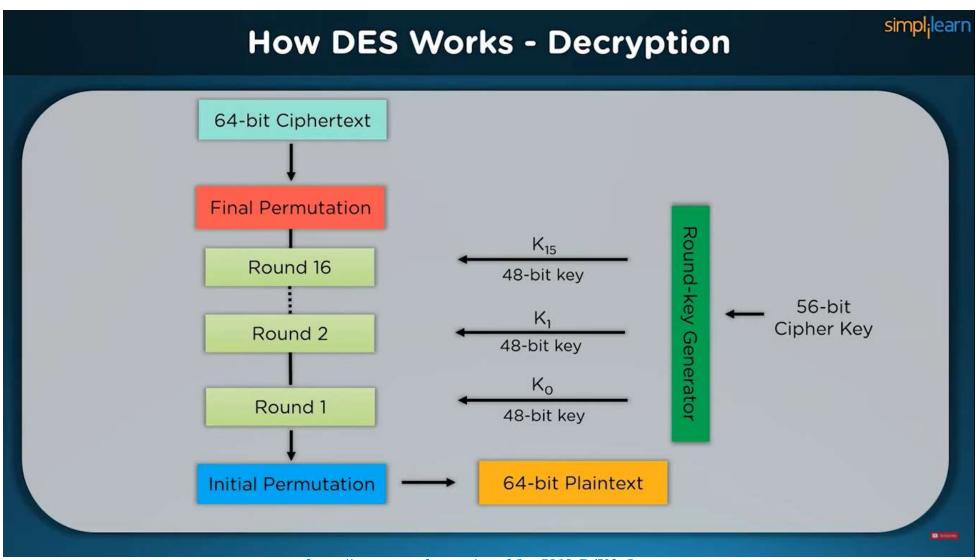


Función de ronda. Expansión de Claves.



Iteration	Number of
Number	Left Shifts
. 1	1
2	1
3	1 2 2 2 2 2 2 2
4	2
5	2
6	2
7	2
8	2
9	1
10	2 2 2
11	2
12	2
13	2 2 2
14	2
15	2
16	1
Criptogra	fía 30





https://www.youtube.com/watch?v=S918rR4VdqQ

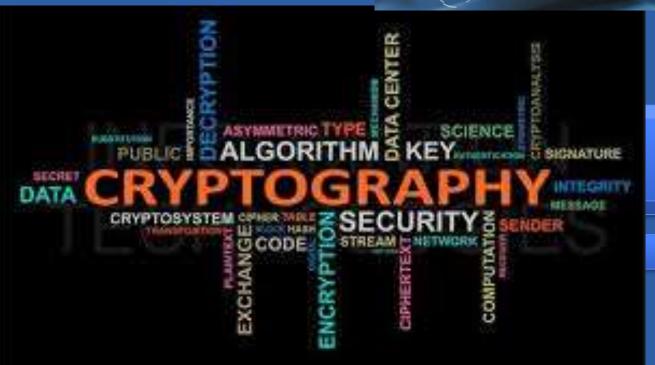


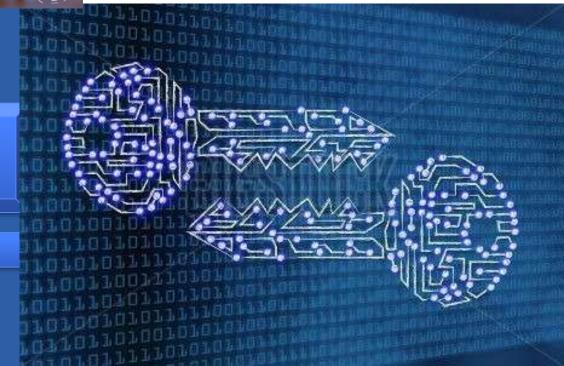
Criptografía





Tema 3.2.2
Algoritmos de cifrado simétrico de bloque: AES







Especificación técnica

Federal Information

Processing Standards Publication 197

November 26, 2001

Announcing the

ADVANCED ENCRYPTION STANDARD (AES)

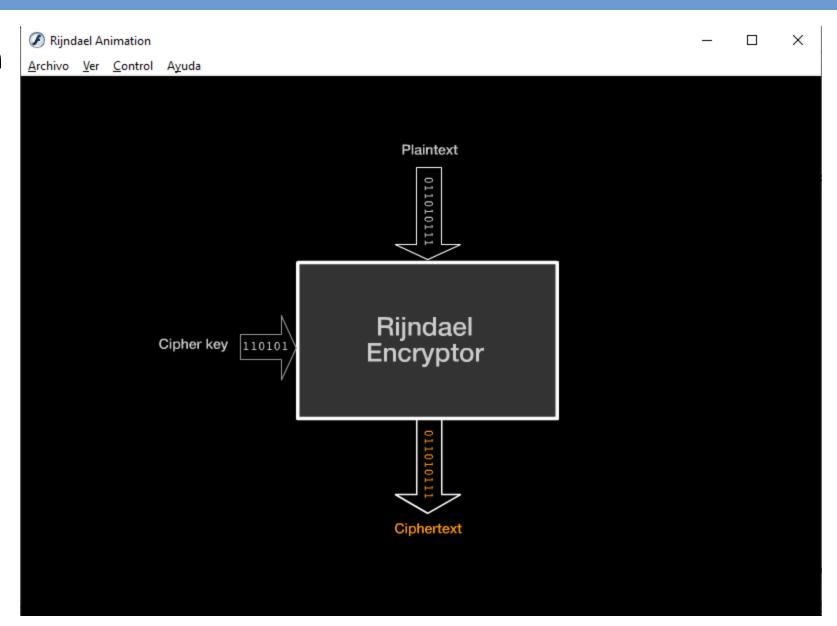
Federal Information Processing Standards Publications (FIPS PUBS) are issued by the National Institute of Standards and Technology (NIST) after approval by the Secretary of Commerce pursuant to Section 5131 of the Information Technology Management Reform Act of 1996 (Public Law 104-106) and the Computer Security Act of 1987 (Public Law 100-235).

- Name of Standard. Advanced Encryption Standard (AES) (FIPS PUB 197).
- Category of Standard. Computer Security Standard, Cryptography.
- 3. Explanation. The Advanced Encryption Standard (AES) specifies a FIPS-approved cryptographic algorithm that can be used to protect electronic data. The AES algorithm is a symmetric block cipher that can encrypt (encipher) and decrypt (decipher) information. Encryption converts data to an unintelligible form called ciphertext; decrypting the ciphertext converts the data back into its original form, called plaintext.

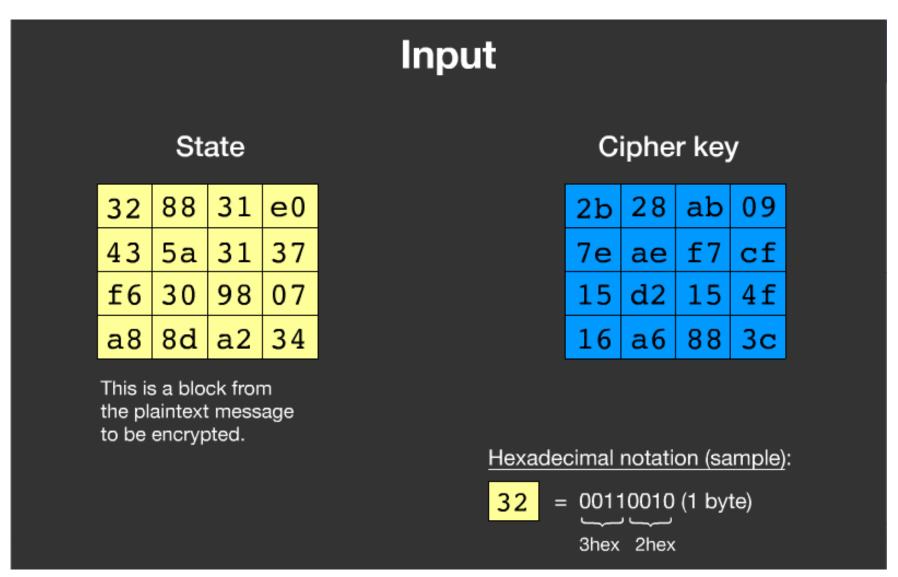
The AES algorithm is capable of using cryptographic keys of 128, 192, and 256 bits to encrypt and decrypt data in blocks of 128 bits.

- 4. Approving Authority. Secretary of Commerce.
- Maintenance Agency. Department of Commerce, National Institute of Standards and Technology, Information Technology Laboratory (ITL).

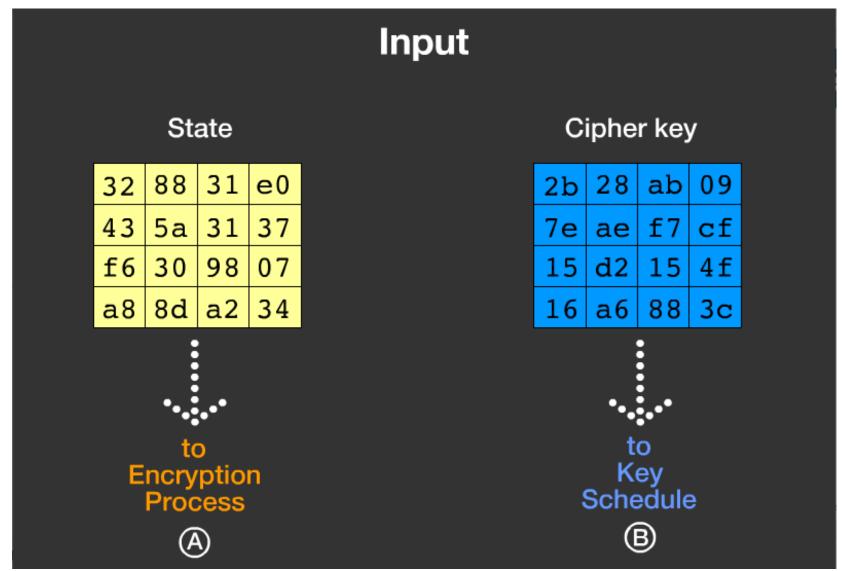








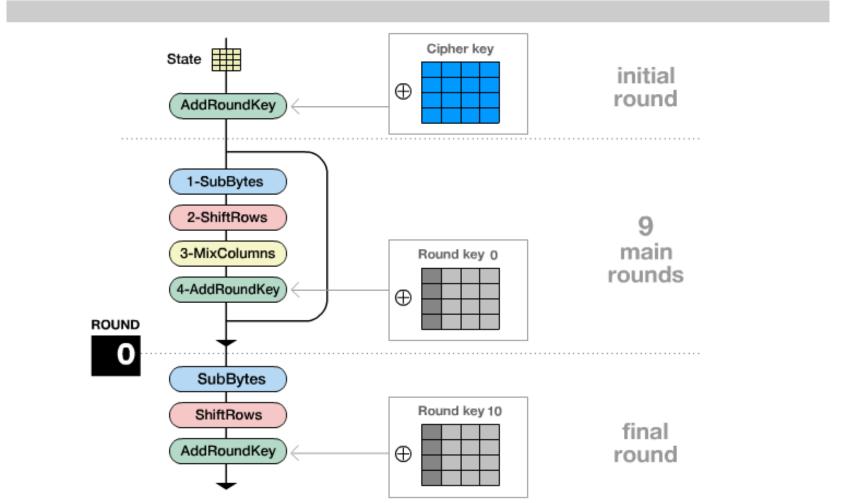




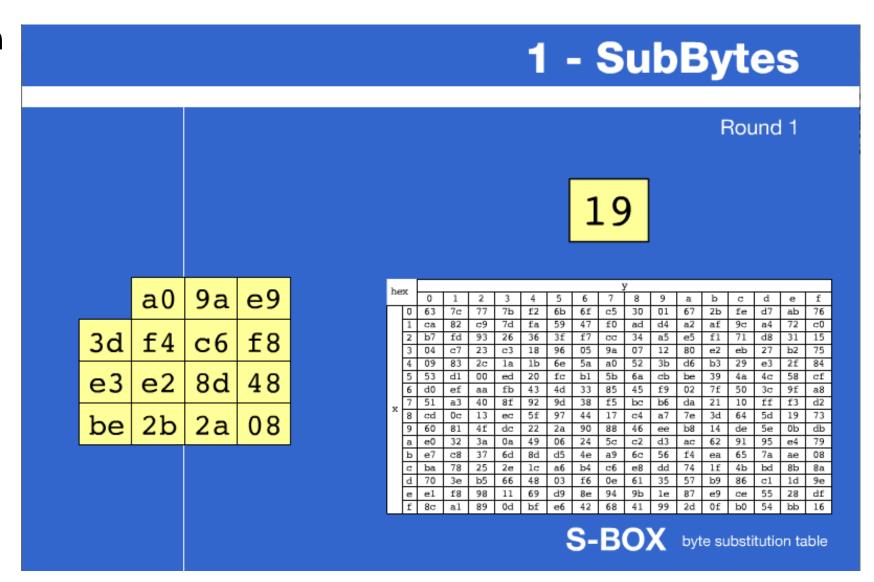


Animation

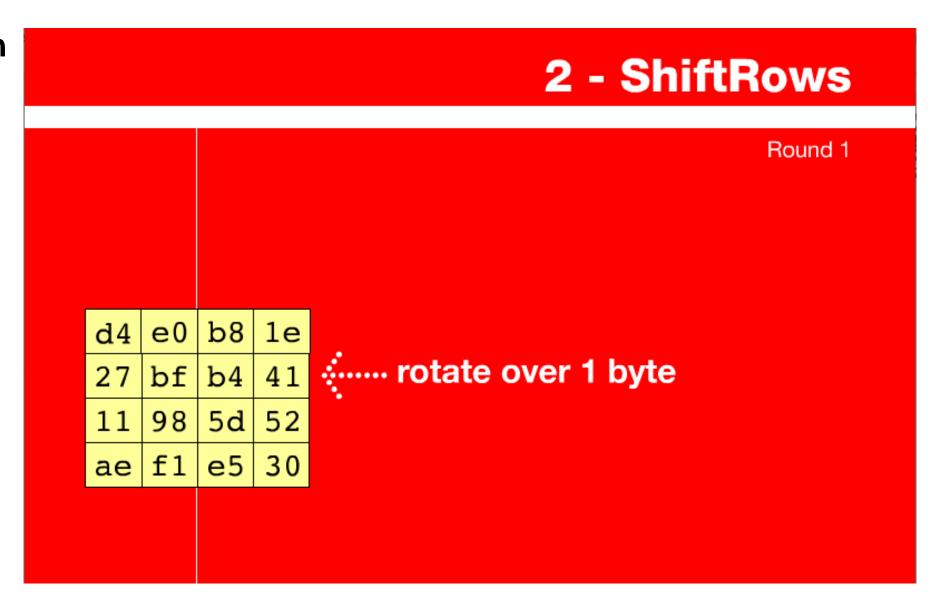
Encryption Process











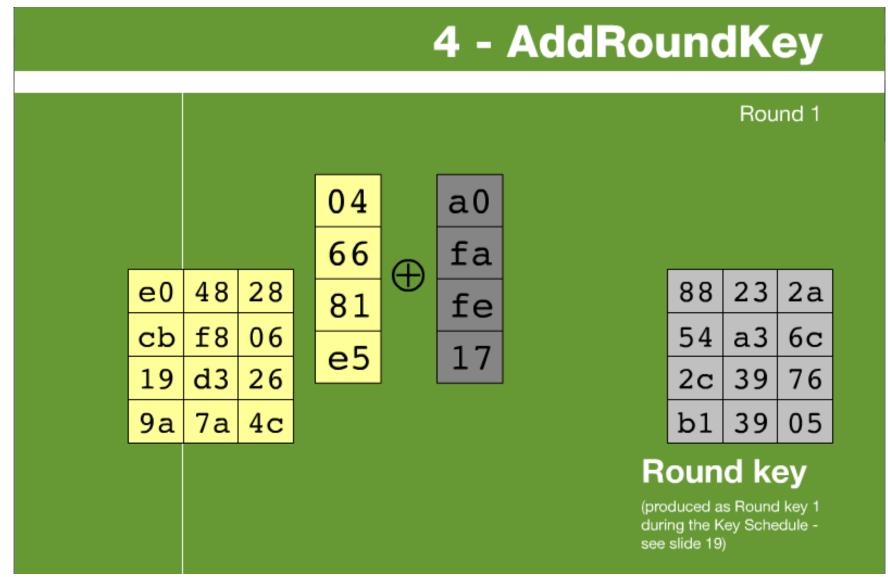


Animation

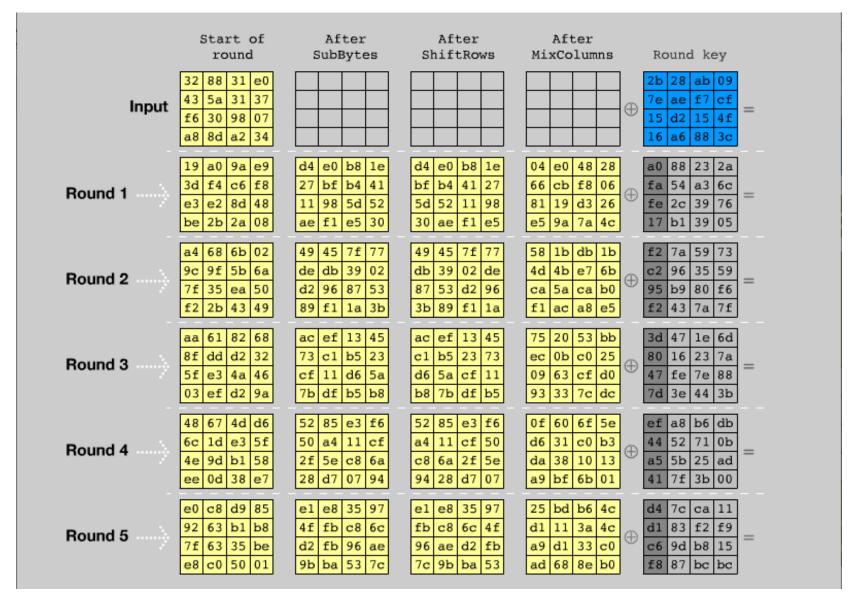
3 - MixColumns Round 1 04 02 03 01 01 bf 66 02 03 01 b8 e0 1e 5d 81 41 b4 03 01 01 02 30 e5 52 11 98 The four numbers of one column f1 e5 ae are modulo multiplied in Rijndael's

Galois Field by a given matrix.

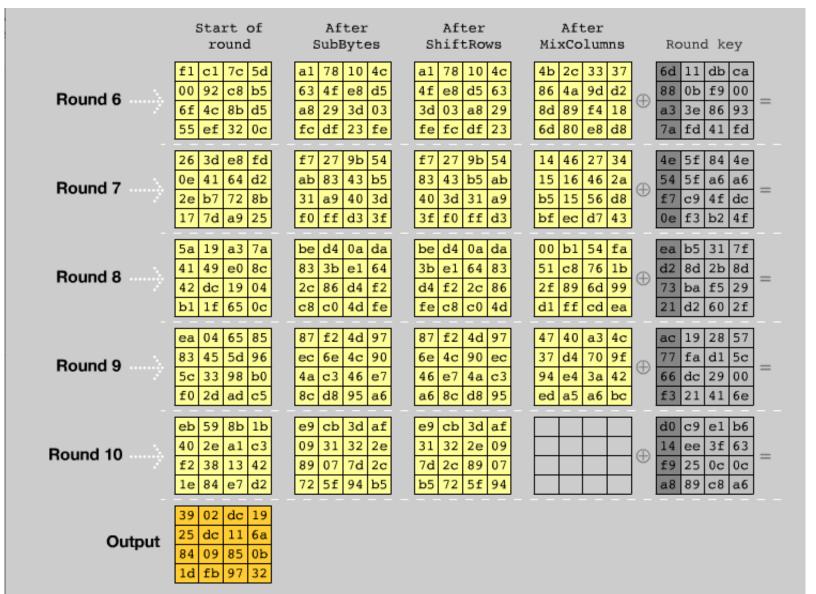




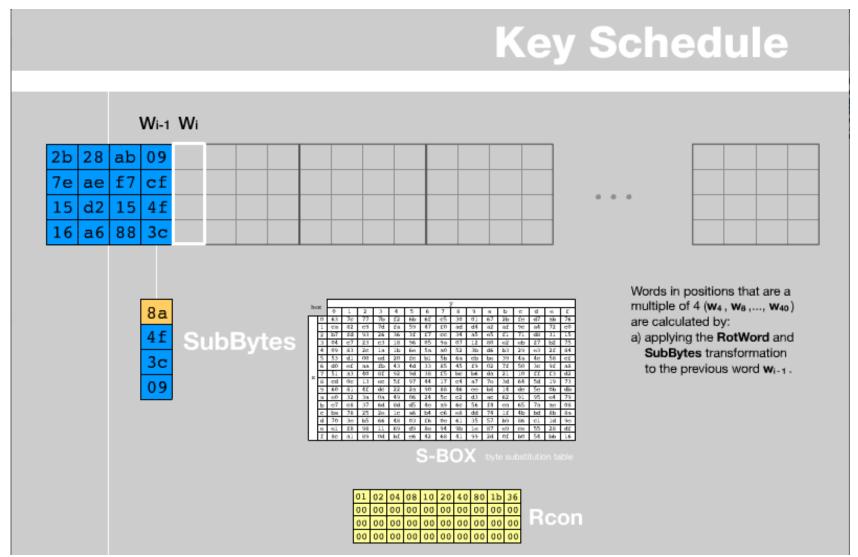




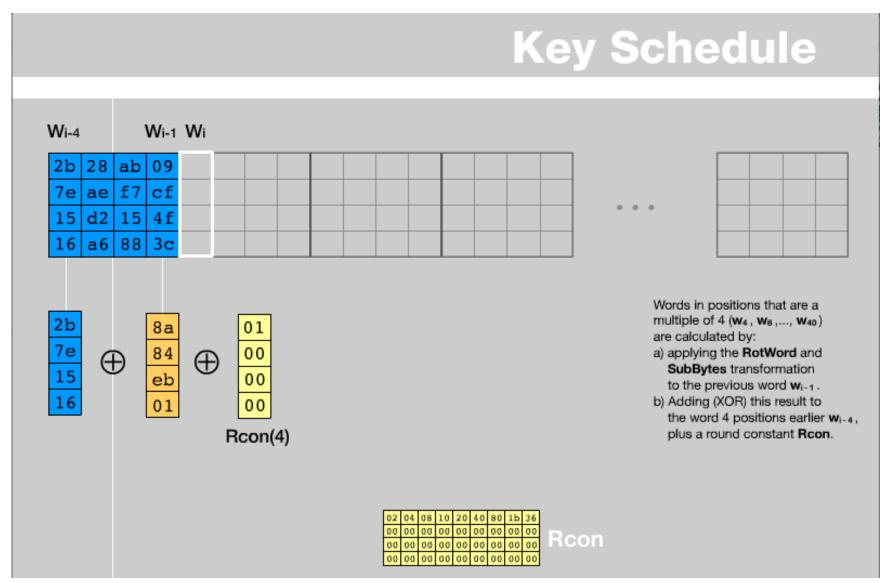




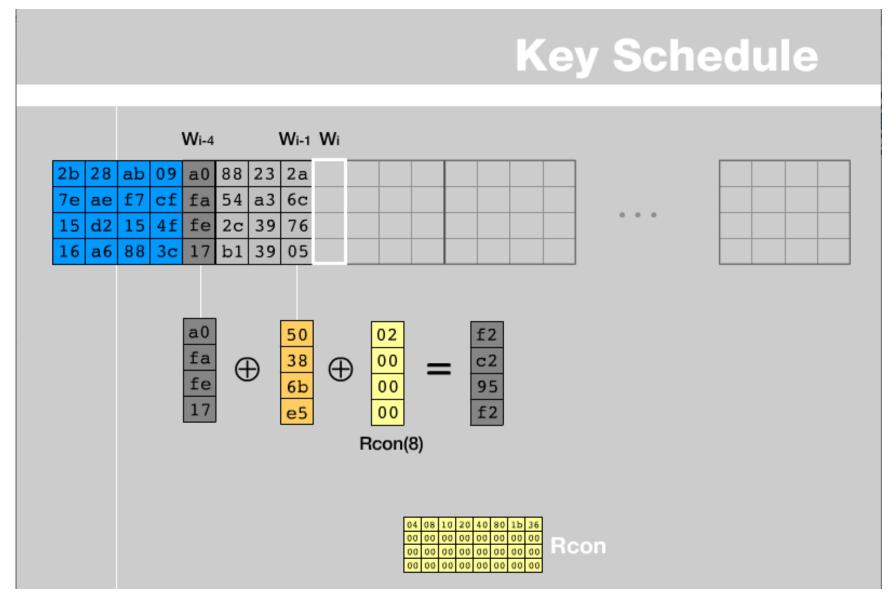






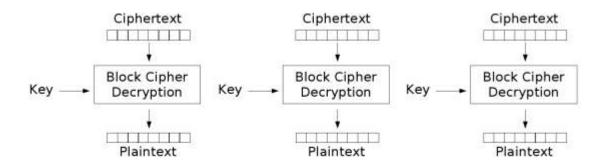




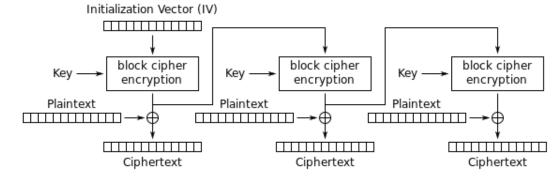


Modos de Operación

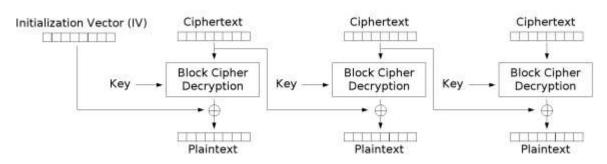




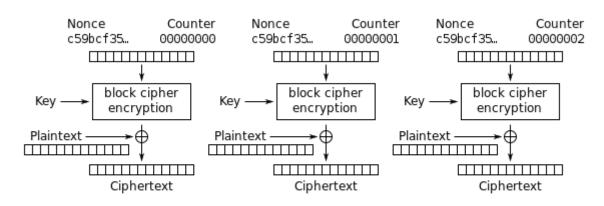
Electronic Codebook (ECB) mode decryption



Output Feedback (OFB) mode encryption



Cipher Block Chaining (CBC) mode decryption



Counter (CTR) mode encryption







Tema 3
Cifrado
Simétrico

