

CS51 Final Project

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For my final project extension, I decided to implement `eval_l`. It seemed like a natural choice, since OCaml itself is lexically scoped. This would have been much easier to do if I had not chosen to use a helper function with return type expression when writing my `eval_d` method. This choice made sense for `eval_d`, since `eval_d` never deals with Closures. Repeatedly wrapping expressions in `Env.Val` then pattern matching to get rid of the `Env.Val` would have made no sense in `eval_d`, and so I did not retroactively change my `eval_d` method. However, for `eval_l`, a helper function returning an expression would not work, since one has to return a closure for the Function case then use that closure in the function application case. This design difference is why I did not attempt to merge the two methods as was suggested in the writeup.

I started by copy-and-pasting my code from `eval_d`, then modifying it to not require a helper function. The main cases that needed to be changed to implement lexical scoping were

```
| Fun (v, exp1) -> exp
| App (exp1, exp2) -> (match eval_d_h exp1 env with
  | Fun (v, f) -> eval_d_h f
    (Env.extend env v (ref (Env.Val (eval_d_h exp2 env))))
  | _ -> raise (EvalError ("invalid function application")))
```

from `eval_d` becoming

```
| Fun (v, exp1) -> Env.Closure (exp, env)
| App (exp1, exp2) -> (match eval_l exp1 env with
  | Env.Closure (Fun (v, f), env2) -> eval_l f
    (Env.extend env2 v (ref (eval_l exp2 env)))
  | _ -> raise (EvalError ("invalid function application")))
```

The `eval_l` version extends `env2` (the environment passed up from the function) rather than `env` (the current environment) when evaluating function application. This is the critical difference that allows the program to "remember" the scope of the function to make `miniml` lexically scoped rather than dynamically scoped.

For reasons unknown to me and several TFs, who said my code for `let rec` in `eval_d` looked fine for both functions, I also had to change my `let rec` case slightly. The version copied from `eval_d`, which was transcribed into OCaml

from the lecture slides and works perfectly for `eval_d`, throws an `EvalError` in `eval_l` because at some point it tries to directly evaluate `Unassigned`. The code in `eval_l` is more methodical about setting up the recursive pointers and does not throw this error.

I use the test function `test_scope ()` to demonstrate the differences between `eval_l` and `eval_d`. `test_scope ()` runs on two cases involving variable shadowing where `eval_l` and `eval_d` should return different results.