

Problem Set 5, Part I

Problem 1: Sorting practice

1-1){3, 4, 18, 24, 33, 40, 8, 10, 12}

1-2){4, 10, 18, 24, 33, 40, 8, 3, 12}

1-3){4, 10, 18, 8, 3, 12, 24, 33, 40}

1-4){10, 18, 4, 24, 12, 3, 8, 40, 33}

1-5){10, 18, 4, 3, 12, 24, 8, 40, 33}

1-6){4, 10, 18, 24, 33, 40, 8, 3, 12}

Problem 2: Practice with big-O

2-1)

function	big-O expression
$a(n) = 5n + 1$	$a(n) = O(n)$
$b(n) = 5 - 10n - n^2$	$b(n) = O(n^2)$
$c(n) = 4n + 2\log(n)$	$c(n) = O(n)$
$d(n) = 6n\log(n) + n^2$	$d(n) = O(n^2)$
$e(n) = 2n^2 + 3n^3 - 7n$	$e(n) = O(n^3)$

2-2) $O(n^2)$. The runtime of the outermost loop is 3 and the middle loop is n and the inner loop is n since it counts 0 to n until n times. So run-time is $3 * n * n = 3n^2$. So the run time is n^2

2-3) $O(n\log n)$ For 0 to n is n times and n to 0, divided by integer division 2 is $\log n$ with base 2. So the $O(n)$ run time is $n * \log_2 n = n\log(n)$.

Problem 3: Comparing two algorithms

worst-case time efficiency of algorithm A: $(n-1) * (n/2) = n^2/2 - n/2 = O(n^2)$

Explanation: Basically run the whole algorithm from start with $n-1$ when $i = 0$ until 0 when i adds up to array's length - 1. So there are total n number of i and j is decreasing by 1 for each j start from $n-1$ when $i = 0$. So Addition of $n-1, n-2, n-3, \dots, 2, 1, 0$ is the multiple of $(n-1)$ and $(n/2)$.

worst-case time efficiency of algorithm B: $O(n\log n)$

Explanation: This algorithm contains Mergesort and its' own for loop. The worst case of mergesort and the best case of runtime is $O(n\log n)$ no matter what. And it takes n times

to search duplicate through from the beginning until the end. So it would takes $O(n \log n + n) = O(n \log n)$

Problem 4: Practice with references

4-1)

Expression	Address	Value
n	0x100	0x712
n.ch	0x712	'n'
n.prev	$0x712 + 6 = 0x718$	0x064
n.prev.prev	$0x064 + 6 = 0x070$	0x360
n.prev.next.next	$0x712 + 2 = 0x714$	null
n.prev.prev.next	$0x360 + 2 = 0x362$	0x064

4-2)

```
n.prev.next = x;  
x.next = n;  
x.prev = n.prev;  
n.prev = x;
```

4-3)

```
public static void initPrevs(DNode first){  
    DNode trav = first;  
    DNode trail = null;  
  
    while(trav != null){  
        trail = trav.next;  
        trail.prev = trav;  
        trav = trav.next;  
    }  
}
```