



The Stack ADT

Computer Science 112
Boston University

Christine Papadakis-Kanaris

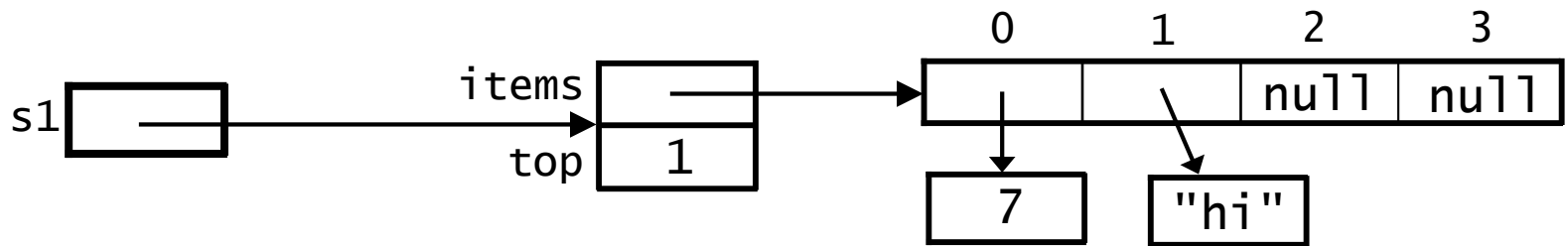
A Stack Interface: First Version

```
public interface Stack {  
    boolean push(Object item);  
    Object pop();  
    Object peek();  
    boolean isEmpty();  
    boolean isFull();  
}
```

- push() returns false if the stack is full, and true otherwise.
- pop() and peek() take no arguments, because we know that we always access the item at the top of the stack.
 - return null if the stack is empty.
- The interface provides no way to access/insert/delete an item at an arbitrary position.
 - encapsulation allows us to ensure that our stacks are manipulated only in ways that are consistent with what it means to be stack

Collection Classes and Data Types

```
public class ArrayStack implements Stack {  
    private Object[] items;  
    private int top;    // index of the top item  
    ...  
}
```



- So far, our collections have allowed us to add objects of any type.

```
ArrayStack s1 = new ArrayStack(4);  
s1.push(7);    // 7 is turned into an Integer object for 7  
s1.push("hi");  
String item = s1.pop();    // won't compile  
String item = (String) s1.pop();    // need a type cast
```

- We'd like to be able to limit a given collection to one type.

```
ArrayStack<String> s2 = new ArrayStack<String>(10);  
s2.push(7);    // won't compile  
s2.push("hello");  
String item = s2.pop();    // no cast needed!
```

Limiting a Stack to Objects of a Given Type

- A *generic* interface and class.

- Here's a generic version of our stack interface:

```
public interface Stack<T> {  
    boolean push(T item);  
    T pop();  
    T peek();  
    boolean isEmpty();  
    boolean isFull();  
}
```

- It includes a *type variable* **T** in its header and body.
 - used as a placeholder for the actual type of the items

A Generic ArrayStack Class

```
public class ArrayStack<T> implements Stack<T> {  
    private T[] items;  
    private int top;    // index of the top item  
    ...  
    public boolean push(T object) {  
        ...  
    }  
    ...  
}
```

- Once again, a type variable **T** is used as a placeholder for the actual type of the items.

Using a Generic Class

```
public class ArrayStack<String> {  
    private String[] items;  
    private int top;  
    ...  
    public boolean push(String item) {  
        ...  
    }  
}
```

```
ArrayStack<String> s1 =  
    new ArrayStack<String>(10);
```

```
public class ArrayStack<T> ... {  
    private T[] items;  
    private int top;  
    ...  
    public boolean push(T item) {  
        ...  
    }  
}
```

```
ArrayStack<Integer> s1 =  
    new ArrayStack<Integer>(25);
```

```
public class ArrayStack<Integer> {  
    private Integer[] items;  
    private int top;  
    ...  
    public boolean push(Integer item) {  
        ...  
    }  
}
```

ArrayStack Constructor

- Java doesn't allow you to create an object or array using a type variable. Thus, we *cannot* do this:

```
public ArrayStack(int maxSize) {  
    items = new T[maxSize];    // not allowed  
    top = -1;  
}
```

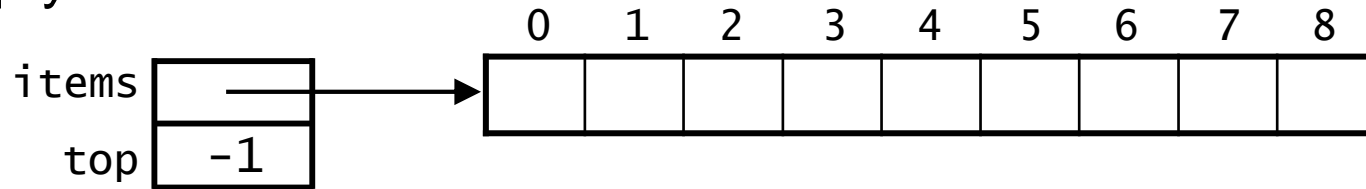
- To get around this limitation, we create an array of type Object and cast it to be an array of type T:

```
public ArrayStack(int maxSize) {  
    items = (T[])new Object[maxSize];  
    top = -1;  
}
```

- The cast generates a compile-time warning, but we'll ignore it.
- Java's built-in ArrayList class takes this same approach.

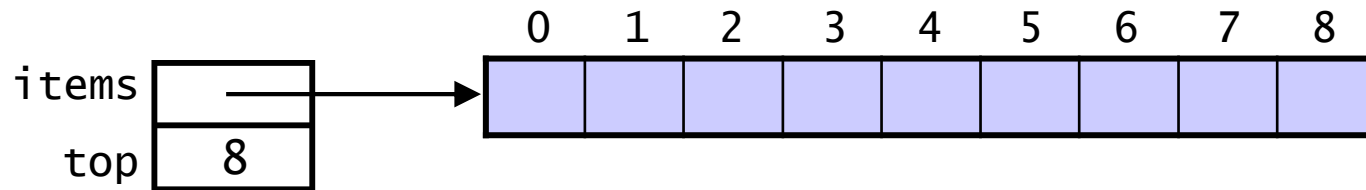
Testing if an ArrayStack is Empty or Full

- Empty stack:



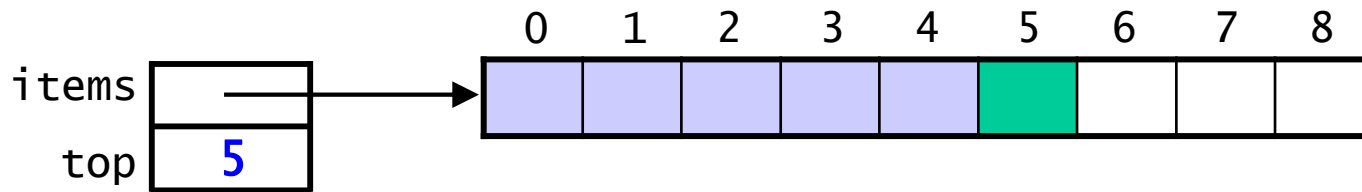
```
public boolean isEmpty() {  
    return (top == -1);  
}
```

- Full stack:



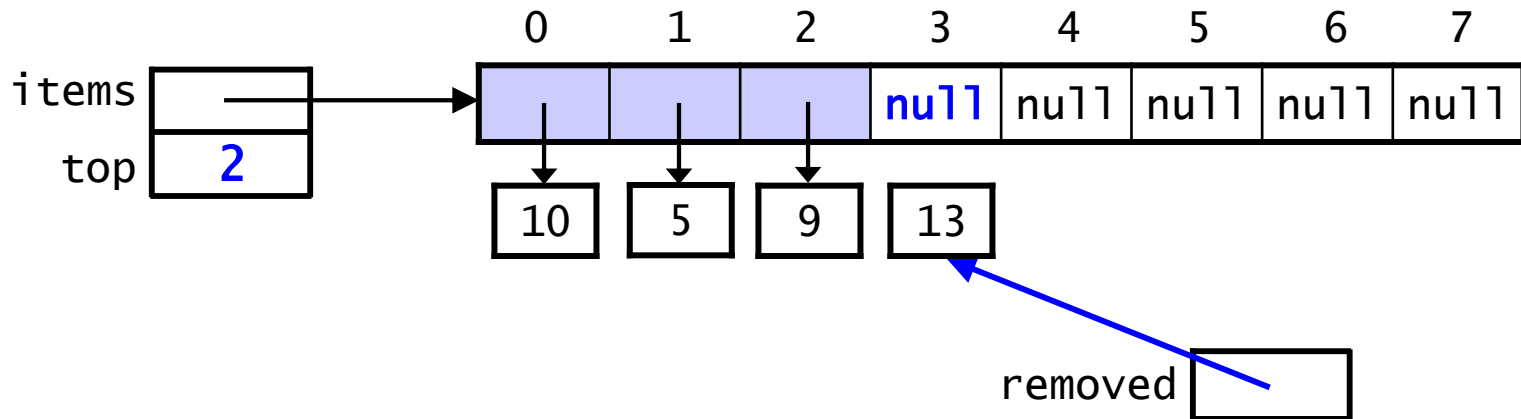
```
public boolean isFull() {  
    return (top == items.length - 1);  
}
```


Pushing an Item onto an ArrayStack



```
public boolean push(T item) {  
    if (isFull()) {  
        return false;  
    }  
    top++;  
    items[top] = item;  
    return true;  
}
```

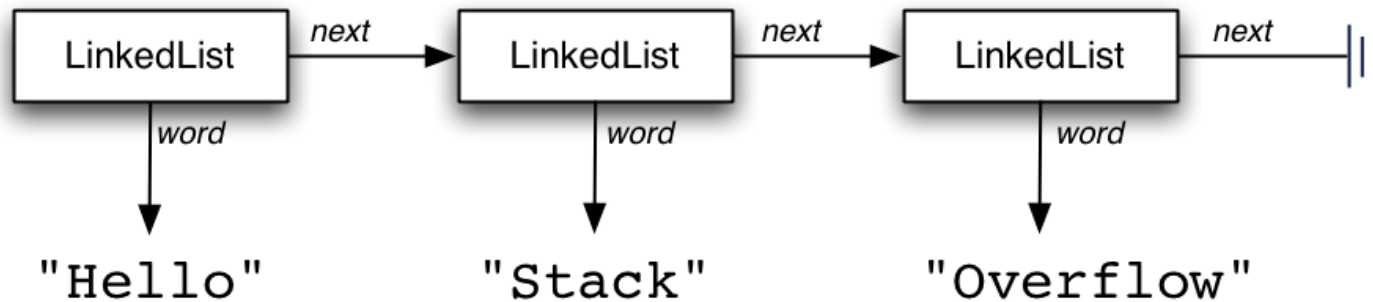
ArrayStack pop() and peek()



```
public T pop() {  
    if (isEmpty()) {  
        return null;  
    }  
    T removed = items[top];  
    items[top] = null;  
    top--;  
    return removed;  
}
```

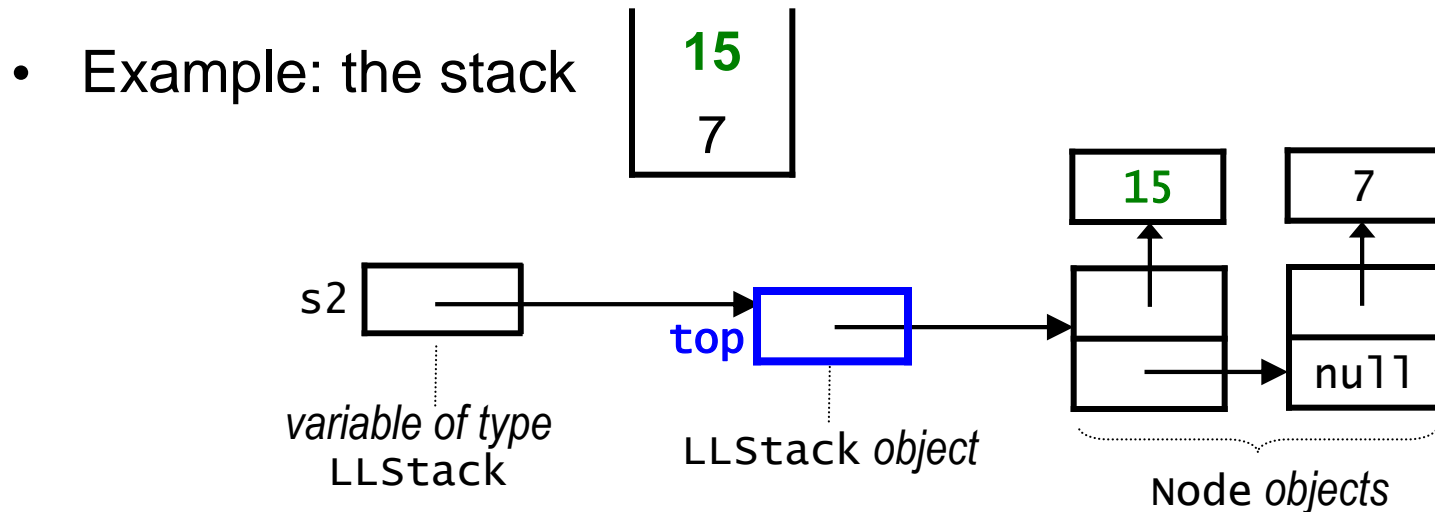
- peek just returns `items[top]` without decrementing `top`.

Linked List Implementation of a Stack



Implementing a Generic Stack Using a Linked List

```
public class LLStack<T> implements Stack<T> {  
    private Node top;    // top of the stack  
    ...  
}
```



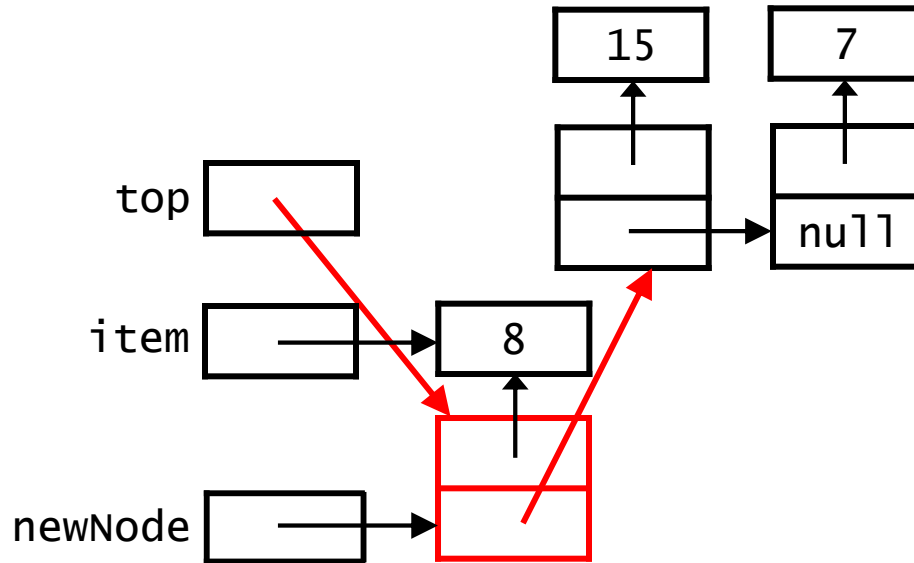
- Things worth noting:
 - our `LLStack` class needs only a single instance variable—a reference to the first node, which holds the top item
 - top item = leftmost item (vs. rightmost item in `ArrayStack`)
 - we don't need a dummy node
 - only one case: always insert/delete at the front of the list!

Other Details of Our LLStack Class

```
public class LLStack<T> implements Stack<T> {
    private class Node {
        private T item;
        private Node next;
        ...
    }
    private Node top;
    public LLStack() {
        top = null;
    }
    public boolean isEmpty() {
        return (top == null);
    }
    public boolean isFull() {
        return false;
    }
}
```

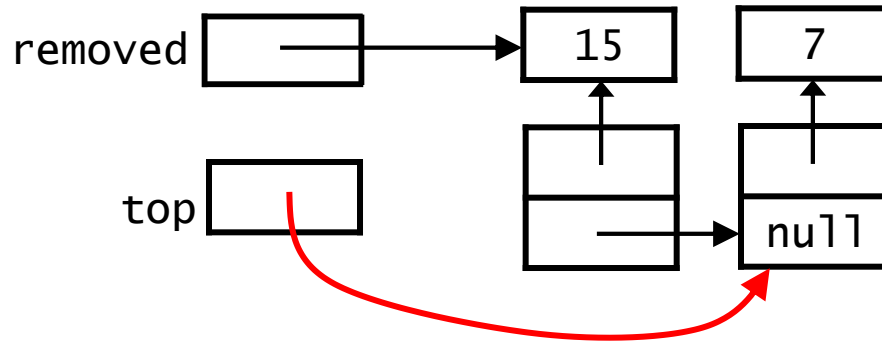
- The inner Node class uses the type parameter T for the item.
- We don't need to preallocate any memory for the items.
- The stack is never full!

LLStack push()



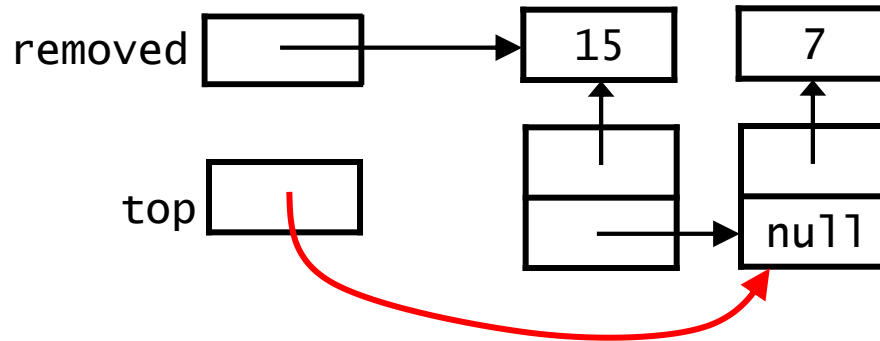
```
public boolean push(T item) {  
    Node newNode = new Node(item, top);  
    top = newNode;  
    return true;  
}
```

LLStack pop() and peek()



```
public T pop() {  
    if (isEmpty()) {  
        return null;  
    }  
  
    T removed = top.item;  
    top = top.next;  
    return removed;  
}
```

LLStack pop() and peek()



```
public T pop() {  
    if (isEmpty()) {  
        return null;  
    }  
  
    T removed = top.item;  
    top = top.next;  
    return removed;  
}
```

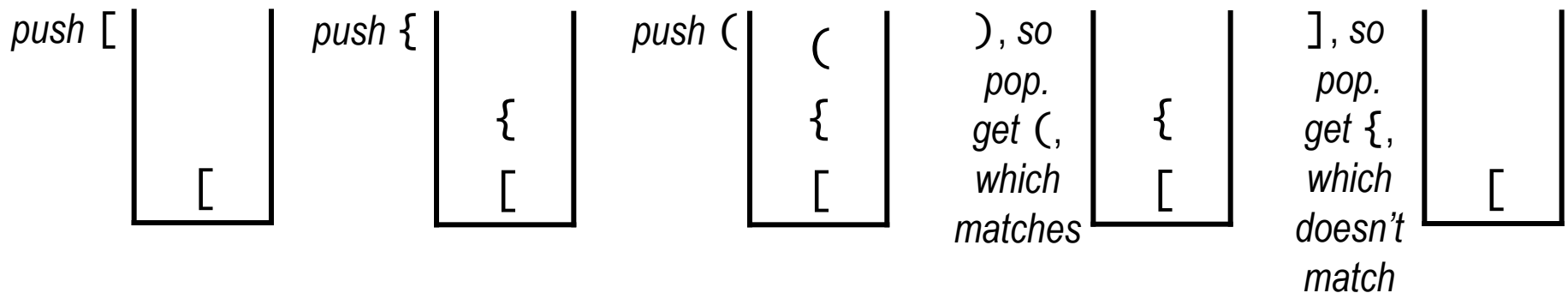
```
public T peek() {  
    if (isEmpty()) {  
        return null;  
    }  
    return top.item;  
}
```


Efficiency of the Stack Implementations

	ArrayStack	LLStack
push()	$O(1)$	$O(1)$
pop()	$O(1)$	$O(1)$
peek()	$O(1)$	$O(1)$
space efficiency	$O(m)$ where m is the <i>anticipated</i> maximum number of items	$O(n)$ where n is the number of items currently on the stack

Applications of Stacks

- The runtime stack in memory
- Converting a recursive algorithm to an iterative one by using a stack to emulate the runtime stack
- Making sure that delimiters (parens, brackets, etc.) are balanced:
 - push open (i.e., left) delimiters onto a stack
 - when you encounter a close (i.e., right) delimiter, pop an item off the stack and see if it matches
 - example: $5 * [3 + \{ (5 + 16 - 2) \}]$



- Evaluating arithmetic expressions