

Efficiency; Classifying problems;

Computer Science 111
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Searching & Sorting Algorithms

- We have learned that if n is the size of our list of data:
 - The running time of **sequential search** is proportional to n
 - The running time of **binary search** is proportional to $\log_2 n$
 - The running time of **selection sort** is proportional to n^2
 - The running time of **quick sort** is proportional to $n \log_2 n$
- For example if binary search takes 1 seconds to search for an element in a list with 500,000 elements, then
 - a list with 1,000,000 elements \rightarrow roughly _____
 - a list with 10,000,000 elements \rightarrow roughly _____
 - a list with 100,000,000 elements \rightarrow roughly _____

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 - a list with 1,000,000 elements \rightarrow roughly 2 seconds
 - a list with 10,000,000 elements \rightarrow roughly _____
 - a list with 100,000,000 elements \rightarrow roughly _____

$$1,000,000 = 2 * 500,000$$

$$\log_2 n \sim 1 \text{ sec} \rightarrow \log_2 2n = \log_2 2 + \log_2 n \sim 1 + 1 = 2 \text{ sec}$$

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 - a list with 1,000,000 elements \rightarrow roughly 2 seconds
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 - a list with 100,000,000 elements \rightarrow roughly _____

$$10,000,000 = 20 * 500,000$$

$$\log_2 n \sim 1 \text{ sec} \rightarrow \log_2 20n = \log_2 20 + \log_2 n \sim 4.32 + 1 = 5.32 \text{ sec}$$

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 - a list with 100,000,000 elements \rightarrow roughly 8.64 seconds

$$100,000,000 = 200 * 500,000$$

$$\log_2 n \sim 1 \text{ sec} \rightarrow \log_2 200n = \log_2 200 + \log_2 n \sim 7.64 + 1 = 8.64 \text{ sec}$$

Algorithm Analysis

- Computer scientists characterize an algorithm's efficiency by specifying its *growth function*.
 - the function to which its running time is roughly proportional
- We've seen several different growth functions:

$\log_2 n$	# binary search
n	# sequential/linear search
$n \log_2 n$	# quicksort
n^2	# selection sort
- Others include:

c^n	# exponential growth
$n!$	# factorial growth
- CS 112 develops a mathematical formalism for these functions.

How Does the Actual Running Time Scale?

- How much time is required to solve a problem of size n ?
 - assume the growth function gives the exact # of operations
 - assume that each operation requires 1 μsec (1×10^{-6} sec)

growth function	problem size (n)					
	10	20	30	40	50	60
n						
n^2						
n^5						
2^n						

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n^2	.0001 s					
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2^n	.001 s					



Note for small data sets we **may** get a misleading result of the algorithm's efficiency!

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n^2	.0001 s	.0004 s				
n^5	.1 s	3.2 s				
2^n	.001 s	1.0 s				

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n^5	.1 s	3.2 s	24.3 s			
2^n	.001 s	1.0 s	17.9 min			

But eventually....

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growth function	problem size (n)					
	10	20	30	40	50	60
n	.00001 s	.00002 s	.00003 s	.00004 s		
n^2	.0001 s	.0004 s	.0009 s	.0016 s		
n^5	.1 s	3.2 s	24.3 s	1.7 min		
2^n	.001 s	1.0 s	17.9 min	12.7 days		

The inefficiency of the algorithm becomes apparent!



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n^2	.0001 s	.0004 s	.0009 s	.0016 s	.0025 s	
n^5	.1 s	3.2 s	24.3 s	1.7 min	5.2 min	
2^n	.001 s	1.0 s	17.9 min	12.7 days	35.7 yrs	

..painfully apparent!!



How Does the Actual Running Time Scale?

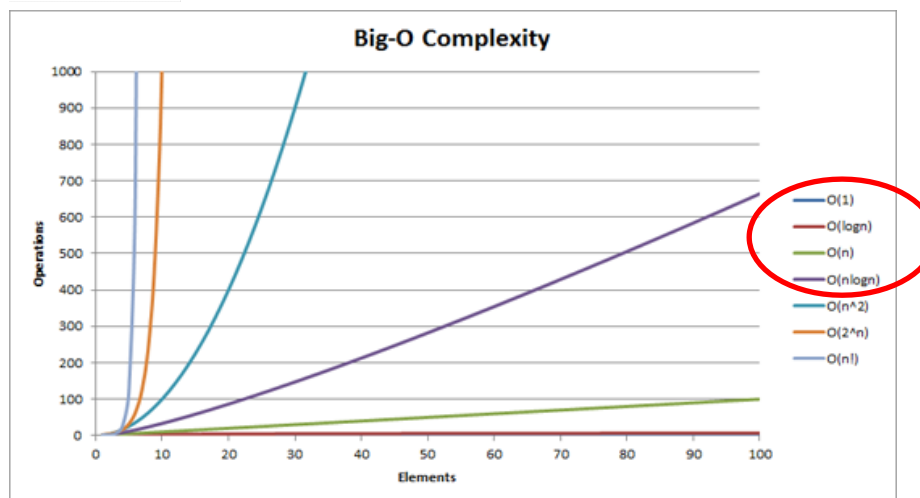
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n^2	.0001 s	.0004 s	.0009 s	.0016 s	.0025 s	.0036 s
n^5	.1 s	3.2 s	24.3 s	1.7 min	5.2 min	13.0 min
2^n	.001 s	1.0 s	17.9 min	12.7 days	35.7 yrs	36,600 yrs

..excruciatingly apparent!!



Algorithm Analysis



We use selection sort to sort a list of length 10,000, and it takes 2 seconds to complete the task.

If we now use selection sort to sort a list of length 80,000, roughly how long should it take?

- A. 2 seconds
- B. 128 seconds
- C. 16 seconds
- D. 256 second
- E. it is impossible to predict
- F. none of these

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 - F. none of these
- 10,000 \rightarrow 80,000
 - the list is 8x longer.
 - The growth function of selection sort is proportional to n^2
 - thus, running time is ~64x longer.

$$80,000 = 8 * 10,000$$

$$(n)^2 \sim 2 \text{ sec} \rightarrow (8n)^2 = 64n^2 \sim 64*2 = 128 \text{ sec}$$

We use **the best algorithms** to search and find an element in the following list of length 5,000,000, and it takes 1 second to complete the task.

lst = [2, 10000, -759, 450, 3, 10, ...]

If we now use **the best algorithms** to search and find another element in the following list of length 20,000,000, roughly how long should it take?

lst = [20, 1010, 759, -401, 9, -10, ...]

- A. 4 seconds
- B. 3 seconds
- C. 128 seconds
- D. 5 seconds
- E. 64 seconds
- D. it is impossible to predict

Available Algorithms:

- sequential search $\sim n$
- binary search $\sim \log_2 n$
- selection sort $\sim n^2$
- quick sort $\sim n \log_2 n$

$$\log_2 4n = \log_2 4 + \log_2 n = 2 + \log_2 n$$

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As list is not sorted, we cannot use binary search directly. So, we have two options:

1. Using **sequential search**.
2. Using **quick sort** to sort it and then **binary search**

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1. Using sequential search

$$n \sim 1 \text{ sec}$$

$$20,000,000 = 4 * 5,000,000$$

$$n \sim 1 \text{ sec} \rightarrow 4n \sim 4 \text{ sec}$$

It takes 4 seconds to find the element using sequential search.

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2. Using **quick sort** to sort it and then **binary search**

$$(n \log_2 n + \log_2 n) \sim 1 \text{ sec}$$

$$20,000,000 = 4 * 5,000,000$$

$$(4n \log_2 4n + \log_2 4n) = [4n (2 + \log_2 n) + 2 + \log_2 n]$$

$$= 2 + 8n + (n \log_2 n + \log_2 n) + 3n \log_2 n$$

$$= 2 + 8n + 1 + 3n \log_2 n = 3 + 8n + 3n \log_2 n$$

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2. Using **quick sort** to sort it and then **binary search**

$3 + 8n + 3n \log_2 n \sim ???$ Seconds

- It is not possible to compute!!!
- The only available relation is $(n \log_2 n + \log_2 n) \sim 1 \text{ sec}$

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- A. 4 seconds
 - Sequential $\sim n \sim 4 \text{ seconds}$
- B. 3 seconds
 - Quicksort + binary $\sim n \log_2 n + \log_2 n \sim ??? \text{ seconds}$
 - Which one is the answer A or F?
- C. 128 seconds
 - Comparing Big O complexity of **n** & **$n \log_2 n$** shows that $O(n) < O(n \log_2 n)$
- D. 5 seconds
 - Thus, $O(\mathbf{n}) < O(\mathbf{n \log_2 n} + \log_2 n)$
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- D. 5 seconds
 - Thus, $O(n) < O(n \log_2 n + \log_2 n)$
 - Hence, using sequential search is the best algorithm.
- E. 64 seconds
 - And, the answer is choice A, i.e. 4 seconds.
- F. it is impossible to predict

Classifying Problems

- **"Easy" problems:** can be solved using an algorithm with a growth function that is a *polynomial* of the problem size, n .
 - $\log_2 n$
 - n
 - $n \log_2 n$
 - n^2
 - n^3
 - etc.
- we can solve large problem instances in a reasonable amount of time

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But even an easy problem can
be solved inefficiently!

- we can solve large problem instances in a *reasonable* amount of time

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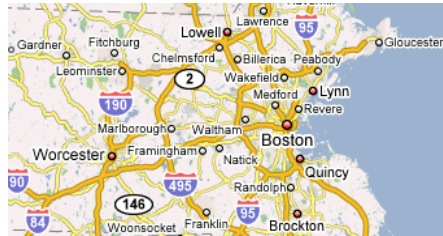
n^3

etc.

- we can solve large problem instances in a reasonable amount of time
- **"Hard" problems:** their only known solution algorithm has an *exponential* or *factorial* growth function.
 - c^n
 - $n!$
 - they can only be solved exactly for small values of n

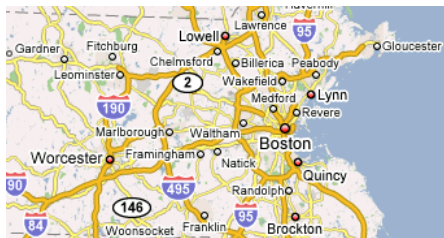
Example of a "Hard" Problem: Map Labeling

- **Given:** the coordinates of a set of *point features* on a map
 - cities, towns, landmarks, etc.
- **Task:** determine positions for the point features' *labels*



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Valid Labeling Optimal Labeling

- Because the point features tend to be closely packed, we may get overlapping labels.
- **Goal:** find the labeling with the fewest overlaps

Map Labeling (cont.)

- One possible solution algorithm: **brute force!**
 - try all possible labeling
- How long would this take?

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- thus, running time will "grow proportionally" to 4^n

exponential time!

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- example: 30 points $\rightarrow 4^{30}$ possible labelings
 - if it took 1 μ sec to consider each labeling,
it would take over 36,000 years to consider them all!

Can Optimal Map Labeling Be Done Efficiently?

- In theory, a problem like map labeling could have a yet-to-be discovered efficient solution algorithm.
- How likely is this?

Can Optimal Map Labeling Be Done Efficiently?

- In theory, a problem like map labeling could have a yet-to-be discovered efficient solution algorithm.
- How likely is this?
- **Not very!**
- If you could solve map labeling efficiently, you could also solve many other hard problems!
 - the *NP-hard* problems
 - another example: the traveling salesperson problem in the reading

Dealing With "Hard" Problems

- When faced with a hard problem, we resort to approaches that quickly find solutions that are "good enough".
- Such approaches are referred to as *heuristic* approaches.
 - heuristic = rule of thumb
 - no guarantee of getting the optimal solution
 - typically get a good solution

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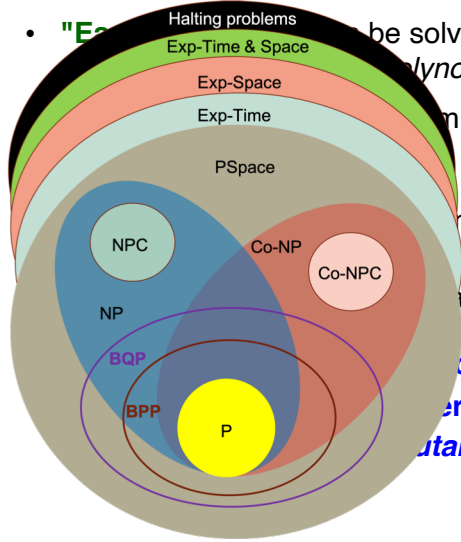
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- **A third class: *Impossible* problems!**
 - **can't be solved, no matter how long you wait!**

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 - **referred to as *uncomputable* problems**

Classifying Problems

-
- "Easy" problems can be solved using an algorithm with a polynomial of the problem size, n .
- Some instances in a problem can be solved in a polynomial of the problem size, n .
- Known solution algorithm with polynomial growth function.
- Not known solution algorithm for all instances of n .
- NP-complete problems!
- How long you wait!
- NP-complete problems



Classifying Pro

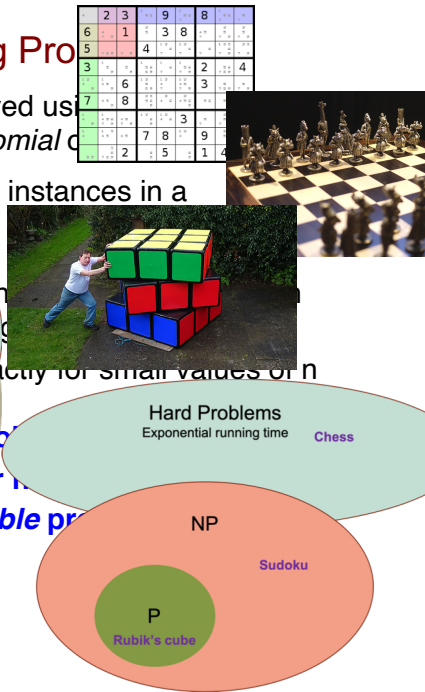
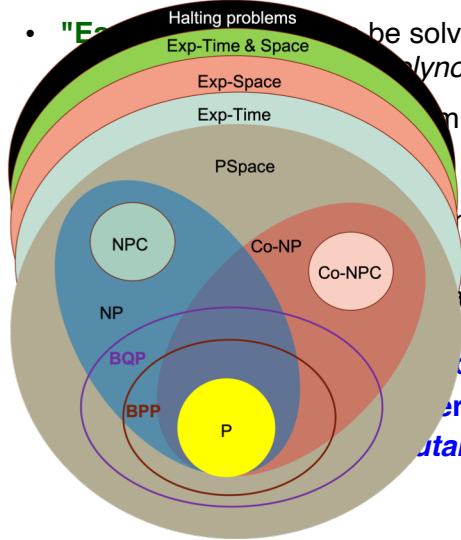
- Complexity classes and their relationships:

 - Halting problems** (top layer)
 - Exp-Time & Space** (second layer)
 - Exp-Space** (third layer)
 - Exp-Time** (fourth layer)
 - PSpace** (fifth layer)
 - NP** (Sixth layer)
 - NPC** (Seventh layer)
 - Co-NP** (Eighth layer)
 - Co-NPC** (Ninth layer)
 - BQP** (Tenth layer)
 - BPP** (Eleventh layer)
 - P** (Twelfth layer)

Other elements:

 - 4x4 grid of numbers (top right):

7	6	7	3
8	6	6	3
7	8	3	9
2	5	1	4
 - Chessboard (top right)
 - Person pushing a large stack of colorful blocks (middle right)
 - Hard Problems (bottom right)
 - Exponential running time (bottom right)
 - NP (bottom right)



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