The List Abstract Data Type



Computer Science 112
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Abstract Data Types

- An abstract data type (ADT) is a model of a data structure that specifies:
 - the characteristics of the collection of data
 - the operations that can be performed on the collection

Specifying an ADT Using an Interface

In Java, we can use an interface to specify an ADT:

```
public interface List {
    Object getItem(int i);
    boolean addItem(Object item, int i);
    Object removeItem(int i);
    int length();
    boolean isFull();
}
```

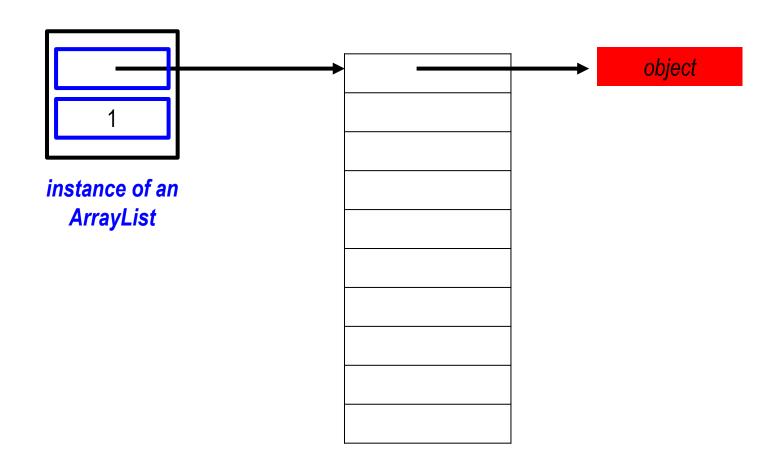
- An interface specifies a set of methods.
 - includes only their headers
 - does not typically include the full method definitions

Implementing an ADT Using a Class

- To implement an ADT, we define a class.
- We specify the corresponding interface in the class header:
 public class ArrayList implements List {
 - tells the compiler that the class will define all of the methods in the interface
 - if the class doesn't define them, it won't compile
- We'll look at two implementations of the List interface:
 - ArrayList uses an array to store the items
 - LLList uses a linked list to store the items

ArrayList Class

Implementing the List interface with an Array



Implementing a List Using an Array

```
public class ArrayList implements List {
        private Object[] items;
        private int length;
        public ArrayList(int maxSize) {
            this.items = new Object[maxSize];
            this.length = 0;
        public int length() {
            return this.length;
        public boolean isFull() {
            return (this.length == this.items.length);
we're showing local variables
outside their stack frame!
                        items
                                                         null
    list
                       length
                                            "if"
    a variable of type
                        an ArrayList object
                                                  "for"
      ArrayList
```

Adding an Item to an ArrayList

Adding at position i (shifting items i, i+1, ... to the right by one):

```
public boolean addItem(Object item, int i) {
       if (i < 0 | | i > length) {
           throw new IndexOutOfBoundsException();
       } else if (isFull()) {
           return false:
       }
       // make room for the new item
       for (int j = length - 1; j >= i; j--) {
           items[j + 1] = items[j];
       }
       items[i] = item;
       length++;
       return true;
   }
example for i = 3:
                            1
                                    3
                         0
                                               6
                                                      8
      items
              6
     length
```

Adding an Item to an ArrayList

Adding at position i (shifting items i, i+1, ... to the right by one): public boolean addItem(Object item, int i) { if (i < 0 | | i > length) { throw new IndexOutOfBoundsException(); } else if (isFull()) { return false: // make room for the new item for (int $j = length - 1; j >= i; j--) {$ items[j + 1] = items[j];items[i] = item; length++; return true; } example for i = 3: items

6

length

Adding an Item to an ArrayList

Adding at position i (shifting items i, i+1, ... to the right by one): public boolean addItem(Object item, int i) { if (i < 0 | | i > length) { throw new IndexOutOfBoundsException(); } else if (isFull()) { return false; } // make room for the new item for (int $j = length - 1; j >= i; j--) {$ items[j + 1] = items[j];} items[i] = item; length++; return true; } example for i = 3: 1 3 5 0 6 8 items

length

Removing item i (shifting items i+1, i+2, ... to the left by one): public Object removeItem(int i) { if (i < 0 || i >= length) { throw new IndexOutOfBoundsException(); Object removed = items[i]; // shift items after items[i] to the left for (int j = i; j < length - 1; j++) { items[length - 1] = null; length--; return removed; example for i = 1: 1 3 8 0 6 items |null|null|null null length "Andy" "Dave" "Kylie" "Jose" "Christine"

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           items[length - 1] = null;
           length--;
           return removed:
example for i = 1:
                           1
                                2
                                     3
                                                          7
                                                               8
                      0
                                          4
                                                    6
 items
                                             null|null|null|null
                                         null
length
                                              ➤ "Andy"
              "Dave"
                             "Jose"
                                      "Kylie"
                   "Christine"
removed
```

Getting an Item from an ArrayList

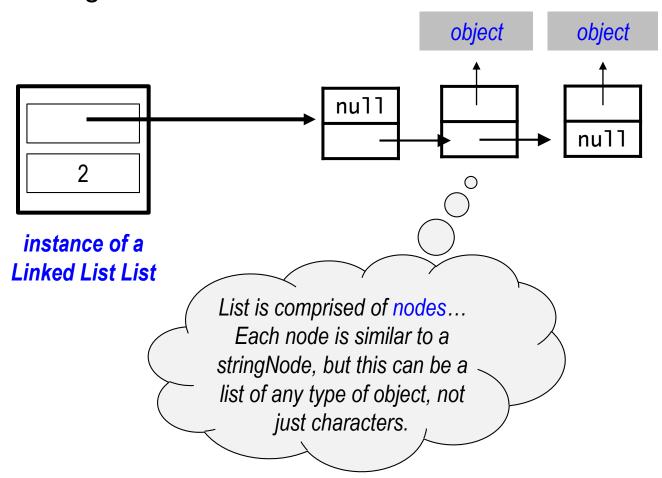
Getting item i (without removing it):
 public Object getItem(int i) {
 if (i < 0 || i >= length) {
 throw new IndexOutOfBoundsException();
 }
 return items[i];
}

Implementing an ADT Using a Class

- To implement an ADT, we define a class.
- We specify the corresponding interface in the class header:
 public class LLList implements List {
 - tells the compiler that the class will define all of the methods in the interface

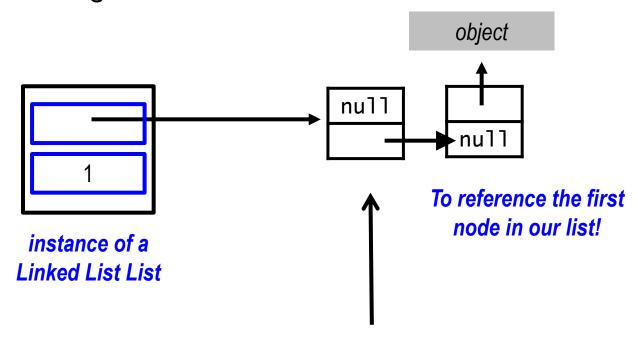
LLList Class

Implementing the List interface with a Linked List



LLList Class

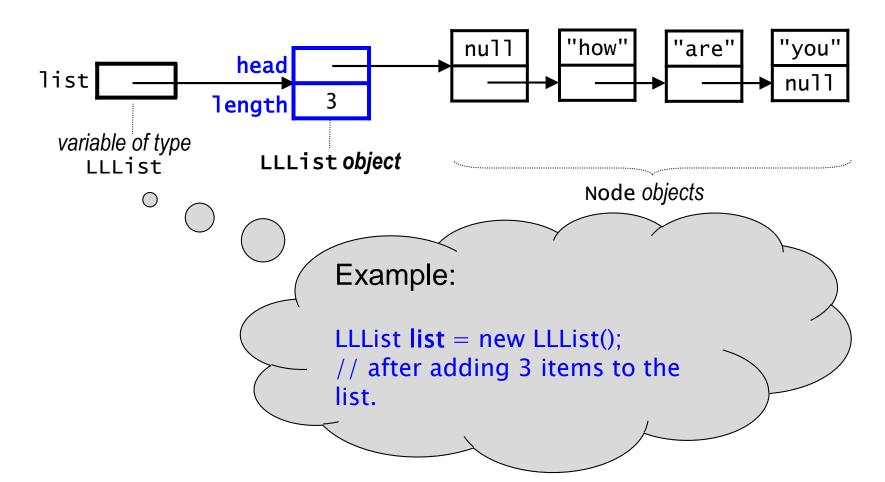
Implementing the List interface with a Linked List



our implementation will use a special first node

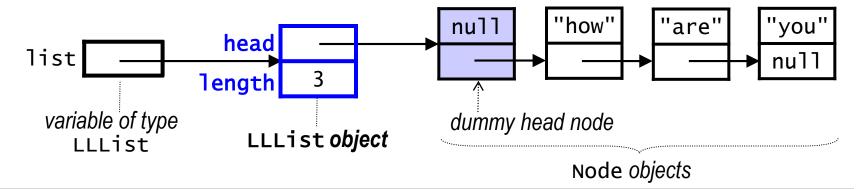
Implementing a List Using a Linked List

```
public class LLList implements List {
    private Node head;
    private int length;
    ...
}
```



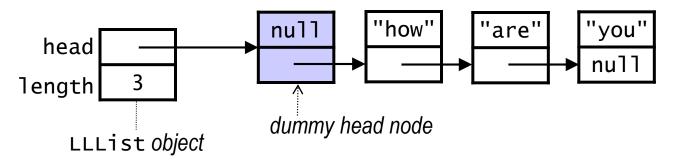
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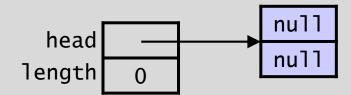


- Differences from the linked lists we used for strings:
 - we "embed" the linked list inside another class
 - users of our LLList class won't actually touch the nodes
 - we use non-static methods instead of static ones myList.length() instead of length(myList)
 - we use a special dummy head node as the first node

Using a Dummy Head Node



- The dummy head node is always at the front of the linked list.
 - like the other nodes in the linked list, it's of type Node
 - it does not store an item
 - it does not count towards the length of the list
- Using it allows us to avoid special cases when adding and removing nodes from the linked list.
- An empty LLList still has a dummy head node:



Implementing a List Using a Linked List

```
public class LLList implements List {
    private Node head;
    private int length;
    ...
}

Implies we need a
    class Node!

variable of type
    LLList

LLList object

Node objects
```

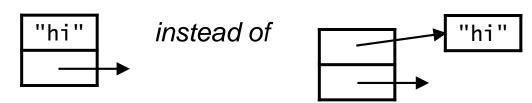
- Differences from the linked lists we used for strings:
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An Inner Class for the Nodes

```
public class LLList implements List {
    private class Node {
        private Object item;
        private Node next;

    private Node(Object i, Node n) {
        item = i;
        next = n;
    }
}
```

- We make Node an inner class, defining it within LLList.
 - allows the LLList methods to directly access Node's private fields, while restricting access from outside LLList
 - the compiler creates this class file: LLList\$Node.class
- For simplicity, our diagrams may show the items inside the nodes.



Other Details of Our LLList Class

```
public class LLList implements List {
   private class Node {
       // see previous slide
   private Node head;
   private int length;
   public LLList() {
        head = new Node(null, null);
        length = 0;
   public boolean isFull() {
        return false:
```

 Unlike ArrayList, there's no need to preallocate space for the items. The constructor simply creates the dummy head node.

Other Details of Our LLList Class

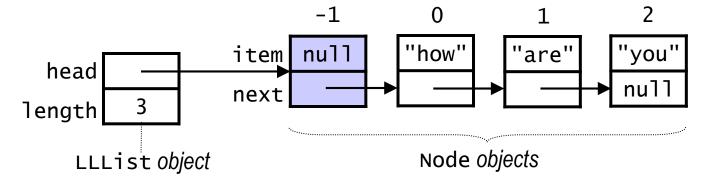
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   private Node head;
   private int length;
   public LLList() {
        head = new Node(null, null);
        length = 0;
    public boolean isFull() {
        return false:
```

- Unlike ArrayList, there's no need to preallocate space for the items. The constructor simply creates the dummy head node.
- The linked list can grow indefinitely, so the list is never full!

Getting a Node

- Private helper method for getting node i
 - to get the dummy head node, use i = -1

```
private Node getNode(int i) {
    // private method, so we assume i is valid!
    Node trav = head;
    int travIndex = -1;
    while (travIndex < i) {</pre>
        travIndex++;
        trav = trav.next;
    }
    return trav;
}
```



Getting a Node

Private helper method for getting node i

 to get the dummy head node, use i private Node getNode(int i) Note travIndex is initialized // private method, so w/ to -1 because: a. head is referencing the Node trav = head; dummy node and that is its int travIndex = -1; position in the list. while (travIndex < i</pre> b. you may need to return a travIndex++; reference to the dummy node trav = trav.next; when adding/removing node at position 0. return trav; "how" "you" null "are" item head nul next length Node objects LLList object

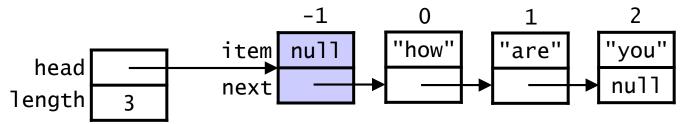
Getting an Item

```
public Object getItem(int i) {
    if (i < 0 | | i >= length) {
         throw new IndexOutOfBoundsException();
    }
    Node n = getNode(i);
    return n.item;
 example for i = 1:
                          -1
                                                     2
                                   0
                                            1
                   item
                                  "how"
                                                   "you"
                         null
                                          "are"
    head
                                                   null
                   next
           3
  length
                                   Node objects
       LLList object
```

Adding an Item to an LLList

```
public boolean addItem(Object item, int i) {
    if (i < 0 || i > length) {
        throw new IndexOutOfBoundsException();
    }
    Node newNode = new Node(item, null);
    Node prevNode = getNode(i - 1);
    newNode.next = prevNode.next;
    prevNode.next = newNode;
    length++;
    return true;
}
```

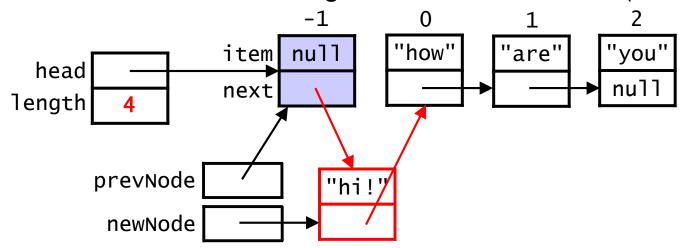
This works even when adding at the front of the list (i = 0):



Adding an Item to an LLList

```
public boolean addItem(Object item, int i) {
    if (i < 0 || i > length) {
        throw new IndexOutOfBoundsException();
    }
    Node newNode = new Node(item, null);
    Node prevNode = getNode(i - 1);
    newNode.next = prevNode.next;
    prevNode.next = newNode;
    length++;
    return true;
}
```

This works even when adding at the front of the list (i = 0):



addItem() Without a Dummy Head Node

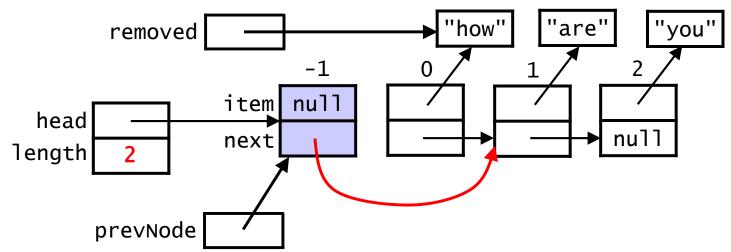
```
public boolean addItem(Object item, int i) {
    if (i < 0 || i > length) {
        throw new IndexOutOfBoundsException();
    Node newNode = new Node(item, null);
                                 // case 1: add to front
   if (i == 0) {
        newNode.next = head;
        head = newNode;
                                 // case 2: i > 0
    } else {
        Node prevNode = getNode(i - 1);
        newNode.next = prevNode.next;
        prevNode.next = newNode;
    length++;
    return true;
}
```

(the gray code shows what we would need to add if we didn't have a dummy head node)

Removing an Item from an LLList

```
public Object removeItem(int i) {
    if (i < 0 || i >= length) {
        throw new IndexOutOfBoundsException();
    }
    Node prevNode = getNode(i - 1);
    Object removed = prevNode.next.item;
    prevNode.next = prevNode.next.next;
    length--;
    return removed;
}
```

This works even when removing the first item (i = 0):





Efficiency of the List Implementations



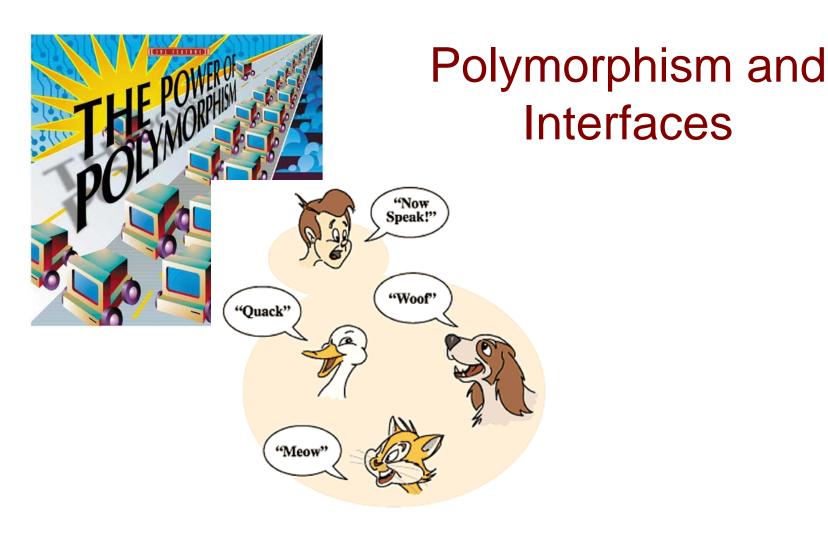
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Efficiency of the List Implementations

summary

| | ArrayList | LLList |
|----------------------|---|---|
| <pre>getItem()</pre> | O(1) because arrays provide random access | O(n) in the average and worst cases because you need to traverse the linked list |
| addItem() | O(n) in the average and worst cases because you need to shift items right O(1) if adding to the back of the list | O(n) in the average and worst cases because you need to traverse the linked list O(1) if adding to the front could make adding to the back O(1) by maintaining a reference to the last node |
| removeItem() | O(n) in the average and worst cases because you need to shift items left O(1) if removing the last item | O(n) in the average and worst cases O(1) if removing the first item could we improve the efficiency of removing the last item? not unless we use a doubly-linked list |
| space efficiency | O(m) where m is the anticipated maximum length | <i>O</i> (n) |



Recall: Polymorphism

 A reference variable of type T can hold a reference to an object from a subclass of T:

```
Rectangle r1 = new Square(50, "cm");
```

- this works because Square is a subclass of Rectangle
- a square is a rectangle!

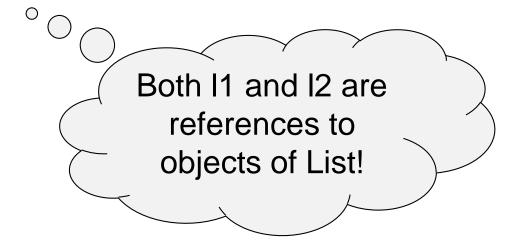
Another Example of Polymorphism

An interface can be used as the type of a variable:

```
List myList;
```

 We can then assign an object of any class that implements the interface:

```
List | 11 = new ArrayList(20);
List | 12 = new LLList();
```



Another Example of Polymorphism

An interface can be used as the type of a variable:

```
List myList;
```

 We can then assign an object of any class that implements the interface:

```
List | 11 = \text{new ArrayList}(20);
List | 12 = new LLList();
              More specifically, both
              11 and 12 are references
               to objects of classes
                that implement the
                   List Interface!
```

Another Example of Polymorphism

An interface can be used as the type of a variable:

```
List myList;
```

 We can then assign an object of any class that implements the interface:

```
List 11 = new ArrayList(20);
List 12 = new LLList();
```

This allows us to write code that works with any implementation of an ADT:

```
public static void processList(List vals) {
   for (int i = 0; i < vals.length(); i++) {
     ...
}</pre>
```

- vals can be an object of any class that implements List
- regardless of which class vals is from,
 we know it has all of the methods in the List interface