# **Flash Tensor Product Attention**

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#### **Abstract**

Flash Tensor Product Attention.

## 1 Toward Faster Computation Without Materializing Q, K and V

We now explore whether it is possible to compute the attention scores  $\mathbf{Q} \mathbf{K}^{\top}$  of Tensor Product Attention (TPA) (Zhang et al., 2025) directly from their factorized forms, thereby reducing floating-point operations.

## 1.1 Single-Head Factorization Setup Without Materializing Q and K

Consider a single head i. Each query vector  $\mathbf{Q}_t^{(i)} \in \mathbb{R}^{d_h}$  is factorized (with rank  $R_q$ ):

$$\mathbf{Q}_{t}^{(i)} = \frac{1}{R_{q}} \sum_{r=1}^{R_{q}} a_{q,i}^{(r)}(\mathbf{x}_{t}) \, \mathbf{b}_{q}^{(r)}(\mathbf{x}_{t}),$$

and each key vector  $\mathbf{K}_{\tau}^{(i)} \in \mathbb{R}^{d_h}$  is factorized (with rank  $R_k$ ):

$$\mathbf{K}_{\tau}^{(i)} = \frac{1}{R_k} \sum_{s=1}^{R_k} a_{k,i}^{(s)}(\mathbf{x}_{\tau}) \, \mathbf{b}_k^{(s)}(\mathbf{x}_{\tau}).$$

Their dot-product for tokens  $t, \tau$  is

$$\left[\mathbf{Q}^{(i)}(\mathbf{K}^{(i)})^{\top}\right]_{t,\tau} = \frac{1}{R_q R_k} \sum_{r=1}^{R_q} \sum_{q=1}^{R_k} a_{q,i}^{(r)}(\mathbf{x}_t) \, a_{k,i}^{(s)}(\mathbf{x}_\tau) \, \left\langle \mathbf{b}_q^{(r)}(\mathbf{x}_t), \mathbf{b}_k^{(s)}(\mathbf{x}_\tau) \right\rangle. \tag{1.1}$$

## 1.2 Multi-Head Case

For multi-head attention with h heads, one repeats the factorization across all heads. The  $\mathbf{b}_q^{(r)}, \mathbf{b}_k^{(s)}$  vectors are shared across heads.

## 1.3 Complexity Analysis

We compare the cost of standard multi-head attention versus TPA under two scenarios:

- 1. Naïve: Materialize Q and K from factors, then perform the usual batched GEMM.
- 2. **Specialized:** Attempt to compute  $\mathbf{Q} \mathbf{K}^{\top}$  directly from the rank- $(R_q, R_k)$  factors without explicitly forming  $\mathbf{Q}, \mathbf{K}$ .

**Standard Multi-Head Attention.** For batch size B and sequence length T:

- Projection cost:  $\mathcal{O}(BT d_{\text{model}}^2)$  or  $\mathcal{O}(BT d_{\text{model}} d_h)$ .
- Dot-product:  $\mathbf{Q}(\mathbf{K})^{\top} \in \mathbb{R}^{(B \ h) \times T \times T} \text{ costs } \mathcal{O}(B \ T^2 \ d_{\text{model}}).$

For large T, the  $\mathcal{O}(BT^2d_{\text{model}})$  term dominates.

## TPA: Naïve Implementation.

- Constructing factors:  $\mathcal{O}(BTd_{\text{model}} \times R_a(h+d_h) + R_k(h+d_h) + R_v(h+d_h))$ .
- Materializing  $\mathbf{Q}, \mathbf{K}: \mathcal{O}(BT(R_q h d_h + R_k h d_h)).$
- Dot-product  $\mathbf{Q}(\mathbf{K})^{\top}$ :  $\mathcal{O}(BT^2 d_{\text{model}})$ .

Typically  $R_q, R_k, R_v \ll h$ , so the overhead of constructing factors is small relative to  $\mathcal{O}(T^2 \, d_{\text{model}})$ . Meanwhile, we still gain KV caching benefits.

**TPA: Specialized Implementation.** If we bypass explicitly forming Q, K, each dot product  $Q_t \cdot K_{\tau}$  is a double sum over rank indices. Below we detail its complexity.

#### 1.4 Complexity Analysis for the Specialized Implementation

**Single-Head Complexity.** A single attention head of dimension  $d_h$ . For each query:

$$\mathbf{Q}_{t}^{(i)} = \frac{1}{R_{q}} \sum_{r=1}^{R_{q}} a_{q,i}^{(r)}(\mathbf{x}_{t}) \, \mathbf{b}_{q}^{(r)}(\mathbf{x}_{t}),$$

and for each key:

$$\mathbf{K}_{\tau}^{(i)} = \frac{1}{R_k} \sum_{s=1}^{R_k} a_{k,i}^{(s)}(\mathbf{x}_{\tau}) \, \mathbf{b}_k^{(s)}(\mathbf{x}_{\tau}).$$

Their dot product:

$$\mathbf{Q}_{t}^{(i)} \cdot \mathbf{K}_{\tau}^{(i)} = \frac{1}{R_{q}R_{k}} \sum_{\tau=1}^{R_{q}} \sum_{\tau=1}^{R_{k}} \left[ a_{q,i}^{(r)}(\mathbf{x}_{t}) \, a_{k,i}^{(s)}(\mathbf{x}_{\tau}) \right] \langle \mathbf{b}_{q}^{(r)}(\mathbf{x}_{t}), \mathbf{b}_{k}^{(s)}(\mathbf{x}_{\tau}) \rangle.$$

For each pair (r, s), we pay:

- 1.  $\mathcal{O}(1)$  for multiplying two scalars,
- 2.  $\mathcal{O}(d_h)$  for the dot product  $\mathbf{b}_q^{(r)}(\mathbf{x}_t) \cdot \mathbf{b}_k^{(s)}(\mathbf{x}_\tau)$ .

Since (r, s) runs over  $R_q \times R_k$ , each token-pair  $(t, \tau)$  costs roughly

$$\mathcal{O}(R_q R_k (1 + d_h)) \approx \mathcal{O}(R_q R_k d_h).$$

For T queries and T keys, that is  $\mathcal{O}(T^2 R_q R_k d_h)$  for a single head.

**Multi-Head and Batches (Reusing b-Dot Products).** When extending to h heads, each head i has its own scalar factors  $\mathbf{a}_{q,i}^{(r)}(\mathbf{x}_t)$  and  $\mathbf{a}_{k,i}^{(s)}(\mathbf{x}_\tau)$ , but the **b**-vectors  $\mathbf{b}_q^{(r)}(\mathbf{x}_t)$  and  $\mathbf{b}_k^{(s)}(\mathbf{x}_\tau)$  can still be *shared* across all heads (assuming the same rank-R factors for every head). Hence, one can split the total cost into two stages:

#### 1. b-Dot-Product Stage:

For each token pair  $(t, \tau)$  and each rank pair (r, s), compute the dot product

$$\left\langle \mathbf{b}_{q}^{(r)}(\mathbf{x}_{t}), \, \mathbf{b}_{k}^{(s)}(\mathbf{x}_{\tau}) \right\rangle \in \mathbb{R}.$$

Since each dot product is  $\mathcal{O}(d_h)$  and there are  $R_q R_k$  rank pairs as well as  $T^2$  token pairs, this stage costs:

$$\mathcal{O}(T^2 R_a R_k d_h)$$
.

Crucially, these b-dot products need only be computed *once* and can be cached for reuse by all heads.

#### 2. Per-Head Scalar Multiplications:

After the b-dot products are precomputed (and cached), each head i only needs to multiply each stored dot product by the corresponding scalars  $\mathbf{a}_{q,i}^{(r)}(\mathbf{x}_t) \, \mathbf{a}_{k,i}^{(s)}(\mathbf{x}_\tau)$ . Since this scalar multiplication is  $\mathcal{O}(1)$  per pair, and there are  $T^2$  token pairs and  $R_q R_k$  rank pairs for each of the h heads, this step costs:

$$\mathcal{O}(h\,T^2\,R_qR_k)$$
.

Putting these together, for batch size B, the total cost is

$$\mathcal{O}(BT^2 R_q R_k d_h) + \mathcal{O}(BT^2 h R_q R_k) = \mathcal{O}(BT^2 R_q R_k (d_h + h)).$$

By contrast, the standard multi-head attention dot-product step is  $\mathcal{O}(BT^2hd_h)$ . Hence, for the specialized TPA approach to *reduce* flops,

$$R_q R_k (d_h + h) \leq h d_h.$$

Thus a practical guideline is to ensure  $R_q R_k < h \frac{d_h}{d_h + h}$ . When that holds, bypassing explicit materialization of  $\mathbf{Q}$  and  $\mathbf{K}$  can be beneficial.

## 1.5 Toward Faster Computation Without Materializing Q, K, V

We have explored a two-step procedure for computing  $\mathbf{Q} \mathbf{K}^{\top}$  directly from factorized queries and keys *without* materializing  $\mathbf{Q}$  or  $\mathbf{K}$ . Here, we extend this idea to also avoid explicitly forming  $\mathbf{V}$ . That is, all three activations  $\mathbf{Q}, \mathbf{K}, \mathbf{V}$  remain factorized throughout the attention pipeline. We present a single-head formulation below, and then discuss multi-head and batch extensions.

Extending the Two-Step Approach to Avoid V Materialization. After we obtain  $\mathbf{Q}\mathbf{K}^{\top}$ , we apply  $\alpha_{t,\tau} = \operatorname{softmax}\left(\frac{1}{\sqrt{d_h}}\left(QK^{\top}\right)_{t,\tau}\right)$ . The final attention output at token t (single head) is

head(t) = 
$$\sum_{\tau=1}^{T} \alpha_{t,\tau} \mathbf{V}_{\tau}$$
.

Using the factorization  $\mathbf{V}_{\tau} = \frac{1}{R_v} \sum_{u=1}^{R_v} a_v^{(u)}(\mathbf{x}_{\tau}) \mathbf{b}_v^{(u)}(\mathbf{x}_{\tau})$ , we write:

head(t) = 
$$\frac{1}{R_v} \sum_{\tau=1}^{T} \alpha_{t,\tau} \sum_{u=1}^{R_v} a_v^{(u)}(\mathbf{x}_{\tau}) \mathbf{b}_v^{(u)}(\mathbf{x}_{\tau}).$$

Rearrange sums:

head(t) = 
$$\frac{1}{R_v} \sum_{v=1}^{R_v} \left[ \sum_{\tau=1}^{T} \left( \alpha_{t,\tau} a_v^{(u)}(\mathbf{x}_\tau) \right) \mathbf{b}_v^{(u)}(\mathbf{x}_\tau) \right].$$

We *still* do not explicitly form  $\mathbf{V}_{\tau}$ . Instead:

Stage 1: Calculating  $\mathbf{b}_v^{(u)}(\mathbf{x}_\tau)$  for all tokens. We simply observe that each output  $\operatorname{head}(t)$  can be computed by summing vectors  $\mathbf{b}_v^{(u)}(\mathbf{x}_\tau) \in \mathbb{R}^{d_h}$  weighted by  $\alpha_{t,\tau} \, a_v^{(u)}(\mathbf{x}_\tau)$ . The complexity for constructing  $\mathbf{b}_v^{(u)}(\mathbf{x}_\tau) \, \forall u, \tau$  is  $\mathcal{O}(T \, R_v \, d_h)$ .

Stage 2: Weighted Summation by  $\alpha_{t,\tau} a_v^{(u)}(\mathbf{x}_{\tau})$ . For each token t, the final attention head output is

$$\frac{1}{R_v} \sum_{\tau=1}^T \alpha_{t,\tau} \sum_{u=1}^{R_v} a_v^{(u)}(\mathbf{x}_\tau) \, \mathbf{b}_v^{(u)}(\mathbf{x}_\tau) \, = \, \frac{1}{R_v} \sum_{u=1}^{R_v} \left[ \sum_{\tau=1}^T \left( \alpha_{t,\tau} \, a_v^{(u)}(\mathbf{x}_\tau) \right) \mathbf{b}_v^{(u)}(\mathbf{x}_\tau) \right].$$

We still never explicitly materialize V. Instead, for each pair (t,u), we must accumulate the sum of T vectors  $\mathbf{b}_v^{(u)}(\mathbf{x}_\tau) \in \mathbb{R}^{d_h}$ , each scaled by the scalar  $\alpha_{t,\tau} \, a_v^{(u)}(\mathbf{x}_\tau)$ . Because each vector is  $d_h$ -dimensional, each (t,u) summation costs  $\mathcal{O}(T\,d_h)$ . Summed over  $t=1\ldots T$  and  $u=1\ldots R_v$ , the total work is  $\mathcal{O}(T^2\,R_v\,d_h)$  for the entire sequence.

In practice, one precomputes all  $\mathbf{b}_v^{(u)}(\mathbf{x}_\tau)$  for  $\tau=1\dots N$ , so each accumulation can be implemented as a simple "scalar-times-vector add" in a tight loop. This cost is usually smaller than the  $\mathbf{Q}\mathbf{K}^\top$  factorized cost if  $R_v\ll h$ .

### 1.6 Overall Complexity for Single-Head

Combining the four bullet-point stages from above (ignoring smaller overheads like the softmax) yields:

- (i) **QK** b-Dot Product Stage:  $\mathcal{O}(T^2 R_q R_k d_h)$ .
- (ii) QK Scalar-Multiply Stage:  $\mathcal{O}(T^2 R_q R_k)$ .
- (iii) Computing  $\mathbf{b}_v^{(u)}(\mathbf{x}_\tau)$  for all tokens:  $\mathcal{O}(T R_v d_h)$ .
- (iv) Weighted Summation by  $\alpha_{t,\tau} a_v^{(u)}(\mathbf{x}_{\tau}) : \mathcal{O}(T^2 R_v d_h)$ .

Hence, for a single head, the total cost is:

$$\mathcal{O}\Big(T^2 R_q R_k d_h + T^2 R_q R_k + T R_v d_h + T^2 R_v d_h\Big).$$

In many cases (especially for large T), the  $\mathcal{O}(T^2)$  terms dominate, so one often focuses on

$$\mathcal{O}\Big(T^2 R_q R_k d_h + T^2 R_q R_k + T^2 R_v d_h\Big).$$

## 1.7 Multi-Head and Batch Extensions (Reuse of b-Dot Products)

When extending to h heads and batch size B, all sequence-length-dependent terms are multiplied by  $\sim B\,h$ . However, crucial b-dot products can be shared across heads:

**QK** b-Dot Products. Since each head has distinct scalar factors  $a_{q,i}$ ,  $a_{k,i}$  but the same  $\mathbf{b}_q^{(r)}$ ,  $\mathbf{b}_k^{(s)}$  across heads, each pairwise dot product

$$\langle \mathbf{b}_q^{(r)}(\mathbf{x}_t), \, \mathbf{b}_k^{(s)}(\mathbf{x}_\tau) \rangle$$

is computed just once per batch. That cost remains

$$\mathcal{O}(BT^2R_qR_kd_h),$$

not multiplied by h. After caching these dot products, each of the h heads pays  $\mathcal{O}(BT^2hR_qR_k)$  total for the head-specific scalar multiplications (the " $\alpha_{t,\tau}$ "-like factors).

**V** b-Evaluations. Likewise, the  $\mathbf{b}_v^{(u)}$  factors are shared across heads (i.e. one set of  $\mathbf{b}_v$ -vectors for all heads). Hence, computing all  $\mathbf{b}_v^{(u)}(\mathbf{x}_\tau)$  for  $\tau=1\dots T$  (across the batch) is a one-time cost:

$$\mathcal{O}(BTR_v d_h)$$
.

Then each head i has its own scalar factors  $a_{v,i}^{(u)}(\mathbf{x}_{\tau})$ , so the final accumulation  $\sum_{\tau=1}^{T} \alpha_{t,\tau} \, a_{v,i}^{(u)}(\mathbf{x}_{\tau}) \, \mathbf{b}_{v}^{(u)}(\mathbf{x}_{\tau}) \cos \mathcal{O} \left( B \, T^2 \, h \, R_v \, d_h \right)$  in total (for all t,u).

Putting it all together, the total flops for multi-head attention with batch size B are:

$$\underbrace{\mathcal{O}\big(B\,T^2\,R_q\,R_k\,d_h\big)}_{\mbox{QK b-dot products (shared across heads)}} + \underbrace{\mathcal{O}\big(B\,T^2\,h\,R_q\,R_k\big)}_{\mbox{per-head QK scalar mult.}} + \underbrace{\mathcal{O}\big(B\,T\,R_v\,d_h\big)}_{\mbox{Compute }\mathbf{b}_v\mbox{ for all tokens (shared across heads)}} + \underbrace{\mathcal{O}\big(B\,T^2\,h\,R_v\,d_h\big)}_{\mbox{final accumulations (per head)}}.$$

**Discussion.** By contrast, standard multi-head attention typically requires  $\mathcal{O}\big(B\,T^2\,h\,d_h\big)$  flops for the  $\mathbf{Q}\mathbf{K}^{\top}$  dot product (plus a similar  $\mathcal{O}\big(B\,T^2\,h\,d_h\big)$  for multiplying by  $\mathbf{V}$ ). The factorization can yield savings provided  $R_qR_k\ll h$  (for  $\mathbf{Q}\mathbf{K}$ ) and  $R_v\ll h$  (for  $\mathbf{V}$ ), though actual speedups depend on how well these multi-stage kernels are implemented and on hardware efficiency. By retaining  $\mathbf{Q},\mathbf{K}$ , and  $\mathbf{V}$  in factorized form, one can forgo the usual steps:

$$\mathbf{x}_t \mapsto \mathbf{Q}_t, \, \mathbf{K}_\tau \mapsto (\mathbf{Q} \, \mathbf{K}^\top) \mapsto \operatorname{softmax}(\mathbf{Q} \, \mathbf{K}^\top) \, \mathbf{V} \mapsto \operatorname{final output}.$$

Instead, the large  $\mathbf{Q}, \mathbf{K}, \mathbf{V}$  tensors (of size  $T \times d_h$ ) are never materialized. The cost is replaced by rank-based b-dot-product computations plus per-head scalar multiplications. The main challenge is to keep the factor ranks  $(R_q, R_k, R_v)$  sufficiently small relative to  $d_h$  and to implement the necessary multi-stage kernels efficiently. When  $R_q, R_k, R_v \ll h$ , fully factorized QKV attention can yield substantial gains in both computation and memory footprint.

#### 1.8 Decoding Speed during Inference Time of MHA, MQA, GQA, MLA, and TPA

Suppose we are in an autoregressive setting, decoding the current token  $\mathbf{x}_T$  given cached keys and values (KV) from all previous tokens  $\mathbf{x}_1, \dots, \mathbf{x}_{T-1}$ . For each attention head  $i \in \{1, \dots, h\}$ , we store  $\mathbf{K}_i \in \mathbb{R}^{T \times d_h}$ ,  $\mathbf{V}_i \in \mathbb{R}^{T \times d_h}$ . Below, we compare the flops needed by MHA, MQA, GQA, MLA, and TPA to compute the next-token logits during inference.

**MHA, MQA, and GQA.** Despite sharing or grouping keys/values in MQA and GQA, the *decoding* cost for MHA, MQA, and GQA remains of the same order. Specifically, for each head *i*, we compute:

$$\mathbf{Q}_i(\mathbf{x}_T) \in \mathbb{R}^{d_h}, \quad \mathbf{K}_i \in \mathbb{R}^{T imes d_h}, \quad \mathbf{Q}_i(\mathbf{x}_T) \mathbf{K}_i^ op \in \mathbb{R}^{1 imes T},$$
 and  $\operatorname{Softmax}(\mathbf{Q}_i(\mathbf{x}_T) \mathbf{K}_i^ op) \mathbf{V}_i \in \mathbb{R}^{d_h}.$ 

Hence, the flops scale linearly in h,  $d_h$ , and T. For example, forming  $\mathbf{Q}_i(\mathbf{x}_T)\mathbf{K}_i^{\top}$  for each head i costs roughly  $\mathcal{O}(h\,d_h\,T)$ .

**MLA.** During inference, MLA can be seen as MQA but uses a larger head dimension to accommodate both RoPE and compressed representations (e.g.,  $d_h' = d_{\rm rope} + d_c$ ). In typical configurations,  $d_{\rm rope} + d_c$  can be significantly larger (e.g.,  $d_h' = 576$  rather than  $d_h = 64$  or 128), thus inflating the dot-product cost by roughly  $4.5 \times$  to  $9 \times$  compared to MHA/MQA/GQA.

**TPA.** Recall that TPA factorizes  $\mathbf{Q}$  and  $\mathbf{K}$  into rank- $(R_q, R_k)$  terms (see Section 1), potentially avoiding large  $\mathbf{Q}$ ,  $\mathbf{K}$  materializations. At inference, TPA's dot-product cost can be broken into two parts:

$$\underbrace{R_q\,R_k\,d_h\,T}_{\text{QK b-dot products (shared across all heads)}} + 2 \underbrace{R_q\,R_k\,h\,T}_{\text{per-head scalar multiplications}},$$

where T is the current sequence length. For concrete values  $d_h = 128$ , h = 64,  $R_q = 8$ , and  $R_k = 2$  (or  $R_q = 16$ ,  $R_k = 1$ ), we obtain:

MHA, MQA, GQA: 
$$128 \times 64 \times T = 8192\,T$$
, MLA:  $576 \times 64 \times T = 36,384\,T$ , TPA:  $\left(8 \times 2 \times 128 \times T\right) + \left(2 \times 8 \times 2 \times 64 \times T\right) = 4096\,T$ .

Thus, in this setup, TPA can significantly reduce the flops needed for computing the  $\mathbf{Q}(\mathbf{x}_T)\mathbf{K}^{\top}$  operation at each decoding step. The actual end-to-end wall-clock speedup also depends on kernel fusion, caching strategies, and hardware implementation details, but the factorized formulation offers a pathway to more efficient decoding than standard attention.

#### References

Yifan Zhang, Yifeng Liu, Huizhuo Yuan, Zhen Qin, Yang Yuan, Quanquan Gu, and Andrew Chi-Chih Yao. Tensor product attention is all you need. *arXiv preprint arXiv:2501.06425*, 2025.