CSC 225 SPRING 2014 ALGORITHMS AND DATA STRUCTURES I MIDTERM EXAMINATION UNIVERSITY OF VICTORIA

| L. | Student ID: | |
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- 3. DATE: 21 FEBRUARY 2014 DURATION: 50 MINUTES INSTRUCTOR: V. SRINIVASAN
- 4. THIS QUESTION PAPER HAS FIVE PAGES (INCLUDING THE COVER PAGE).
- 5. THIS QUESTION PAPER HAS FOUR QUESTIONS.
- 6. ALL ANSWERS TO BE WRITTEN ON THIS EXAMINATION PAPER.
- 7. THIS IS A CLOSED BOOK EXAM. NO CALCULATORS OR ANY OTHER ITEMS ARE ALLOWED.
- 8. KEEP YOUR ANSWERS SHORT AND PRECISE.

| Q1(4) | |
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| Q2(4) | |
| Q3 (4) | |
| Q4(3) | |
| TOTAL(15) = | |

- 2
- 1. The two parts of Question 1 test the basics of **asymptotic analysis**. All logarithms in this question are to the base 2.
 - (a) [2 Marks] Order the following functions by increasing growth rates. Justify whenever needed.

$$n^{3/2}$$
, $2n(\log n)^2$, $\sqrt{\log n}$, $n^{0.5}$, $2^{\sqrt{n}}$, 100.

(b) [2 Marks] Compute the running time of the following algorithm in terms of n using O-notation. Show your calculations.

Algorithm Loop(n)

$$s \leftarrow 0$$

for $i \leftarrow 1$ to $2n$ do
for $j \leftarrow 1$ to i do
 $s \leftarrow s + i$

2. Question 2 checks if we can solve **recurrence equations**.

[4 Marks] Solve the following recurrence equation to get a closed-formula for T(n). You can assume the n is a power of two.

$$T(n) = 1 \text{ if } n = 1$$

= $2T\left(\frac{n}{2}\right) + \log n \text{ if } n \ge 2$

- 3. Question 3 is about **Sorting Algorithms**. We assume that the output should be sorted in ascending order from smallest to largest.
 - (a) [2 Marks] Suppose that your Quick-Sort algorithm uses a pivot rule that picks the element in the "middle". That is, for an array $A[0,1,\ldots,n-1]$ of size n, it uses the element in A[n/2] as pivot if n is even and the element in A[(n-1)/2] as pivot if n is odd. Illustrate how this algorithm works using a quick-sort tree on the input:

$$[7 \ 6 \ 5 \ 4 \ 3 \ 2 \ 1]$$

(b) [2 Marks] What is the running time of this version of quick-sort on sequences of size n that are already sorted from largest to smallest? Explain why by writing down the recurrence equation for the running time of your algorithm on such sequences.

4. Question 4 is based on **Proof by Induction**.

We say that a tree is binary if every internal node has exactly two children. Using induction on the size of the tree, prove the following claim: In any binary tree, the number of external nodes (also called as leaves) is one more than the number of internal nodes. (Hints: For the base case, use a tree of size 1. For the induction step, use the fact that the left and right subtrees of the root node are smaller in size.) [3 Marks]