CSC 226

Algorithms and Data Structures: II

Rich Little
rlittle@uvic.ca
ECS 516

Kruskal's Algorithm

Algorithm Kruskal
Input: a weighted connected graph $G = (V, E)$
Output: an MST T for G
Data structure: Disjoint sets (lists or union-find) *DS*;
    sorted weights priority queue A; and tree T

for each $v \in V$ **do** C(v) ← DS.insert(v) **end** // one cluster per vertex
for each (u,*v*) $\in E$ **do** A.insert((u,v)) **end** // sort edges by weight
$T \leftarrow \emptyset$
while T has fewer than n-1 edges **do**
  (u, v) ← A.deleteMin() // edge with smallest weight
  $C(u)$← DS.findCluster(u);
  $C(v)$← DS.findCluster(v);
  if $C(u) \neq C(v)$ **then**
    add edge (u, v) to T;
    DS.insert(DS.union($C(v)$, $C(u)$)); // merge two clusters
  end
end
return T

Idea:

- Avoid sorting the edge weights by storing the edges in a heap …..

Building up a heap

- m standard insert-operations for a heap result in $O(m \log(m))$ time.
- Can we build up a heap for m given elements faster? Is $O(m)$ possible?

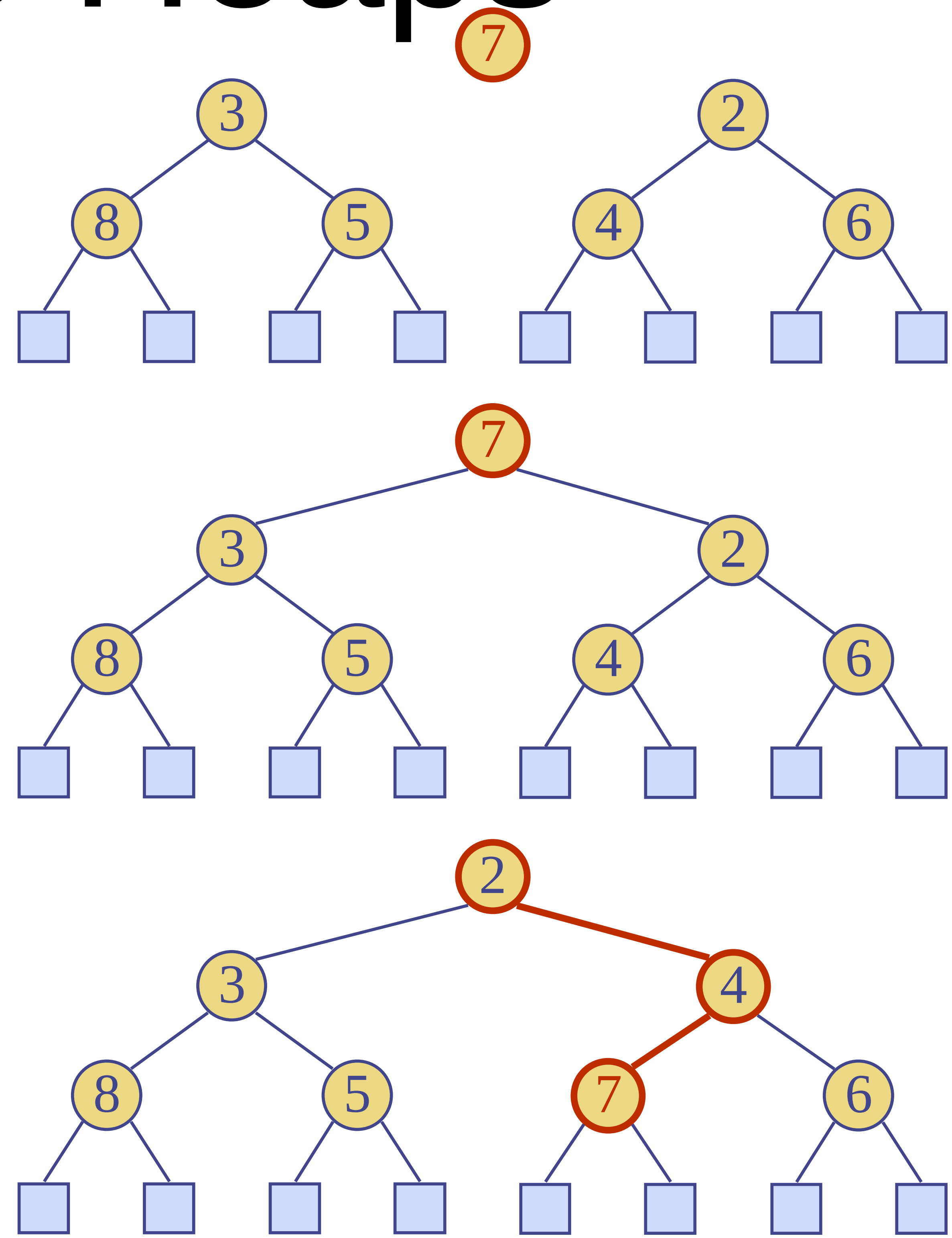
Bottom-Up Heap

```
Algorithm BottomUpHeap(S):
   Input: A list S storing m keys
   Output: A heap T storing the m keys

   if S is empty then
      return external node
   remove the first key, k, from S
   split S in half, lists S1 and S2
   T1 ← BottomUpHeap(S1)
   T2 ← BottomUpHeap(S2)
   T ← merge(k,T1,T2)
   DownHeap(T,root)
   return T
```

Merging Two Heaps

- We are given two heaps and a key *k*
- We create a new heap with the root node storing *k* and with the two heaps as subtrees
- We perform downheap to restore the heap-order property



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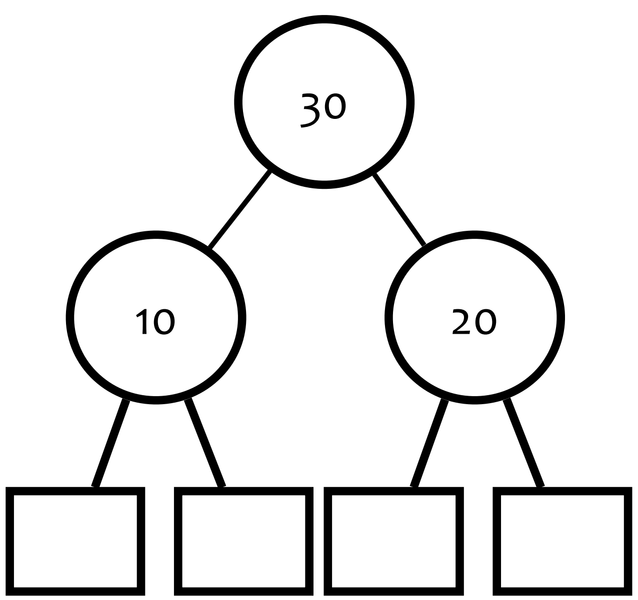

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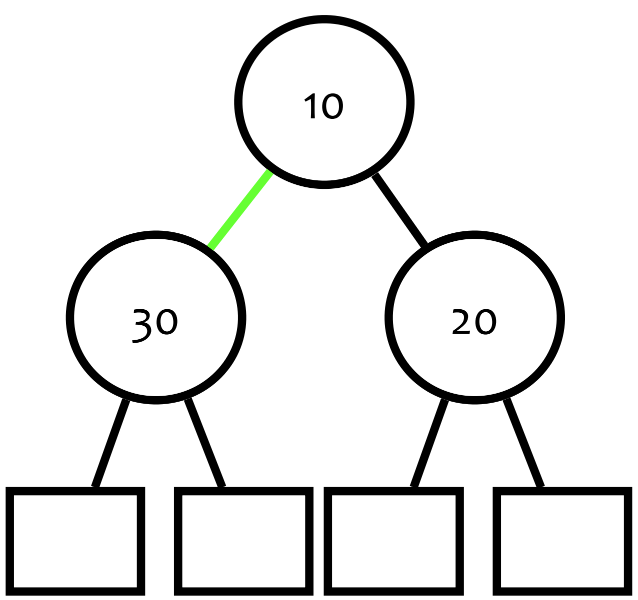

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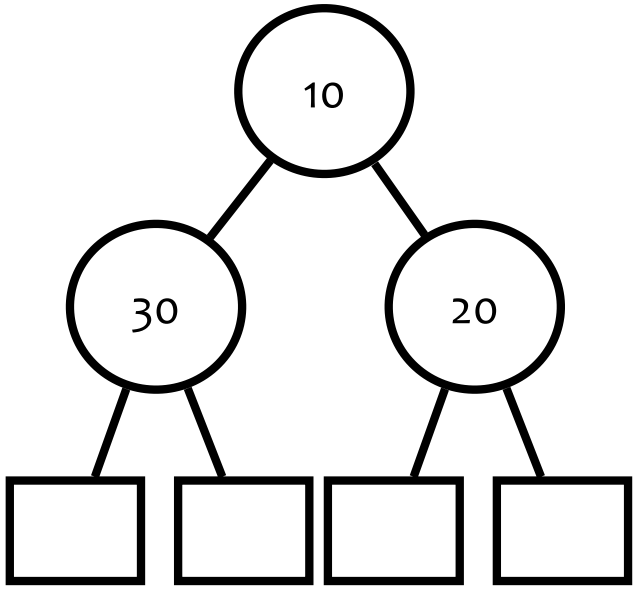

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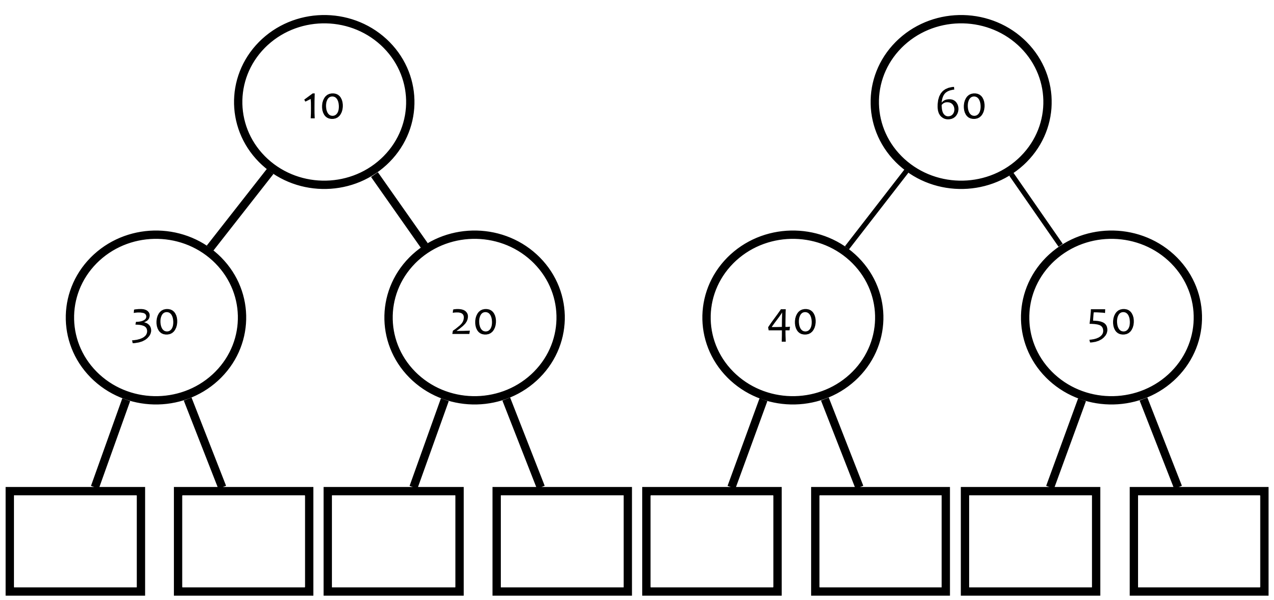

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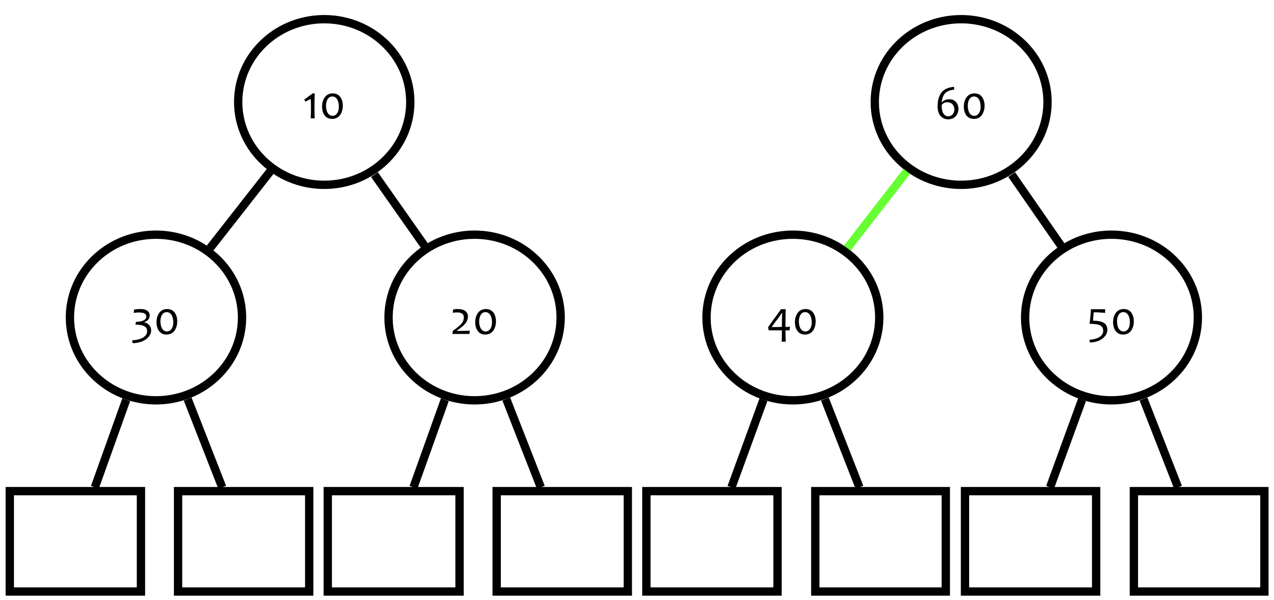

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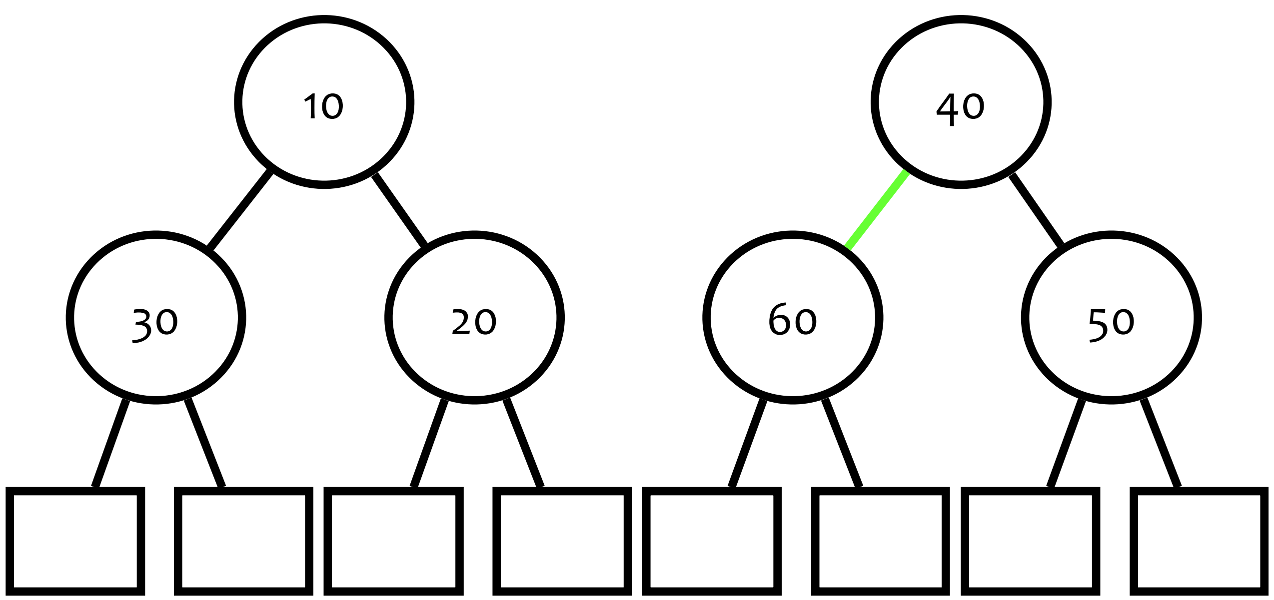

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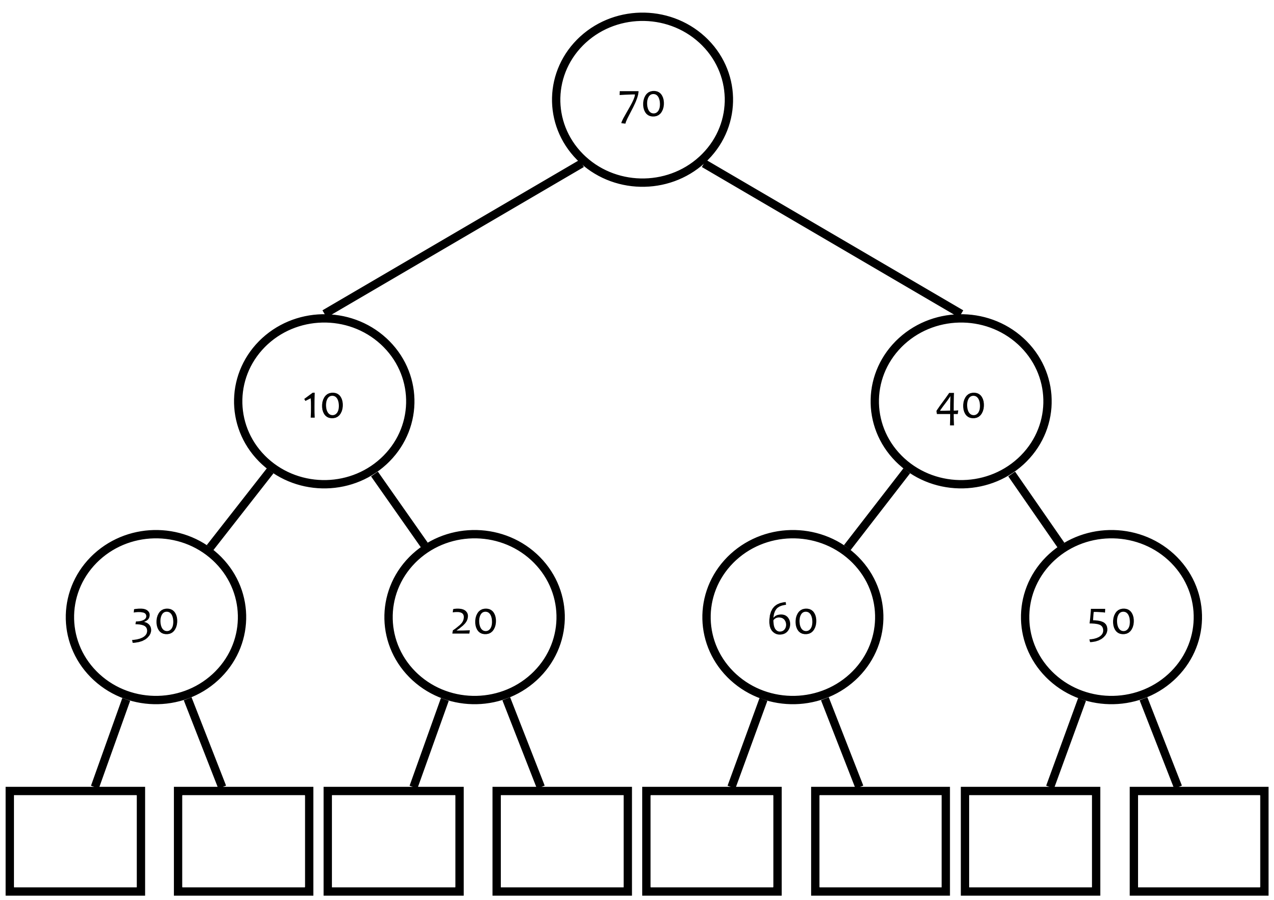

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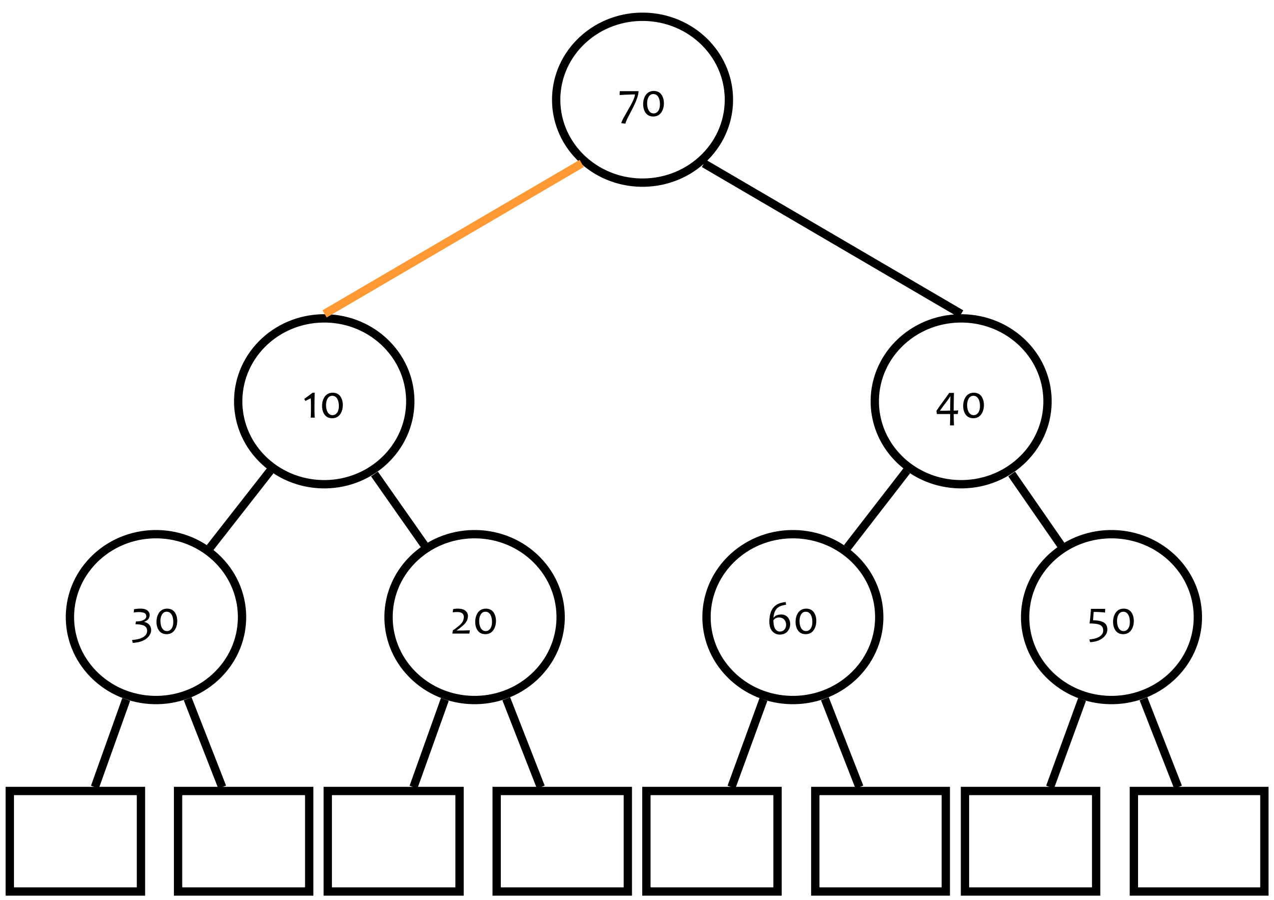

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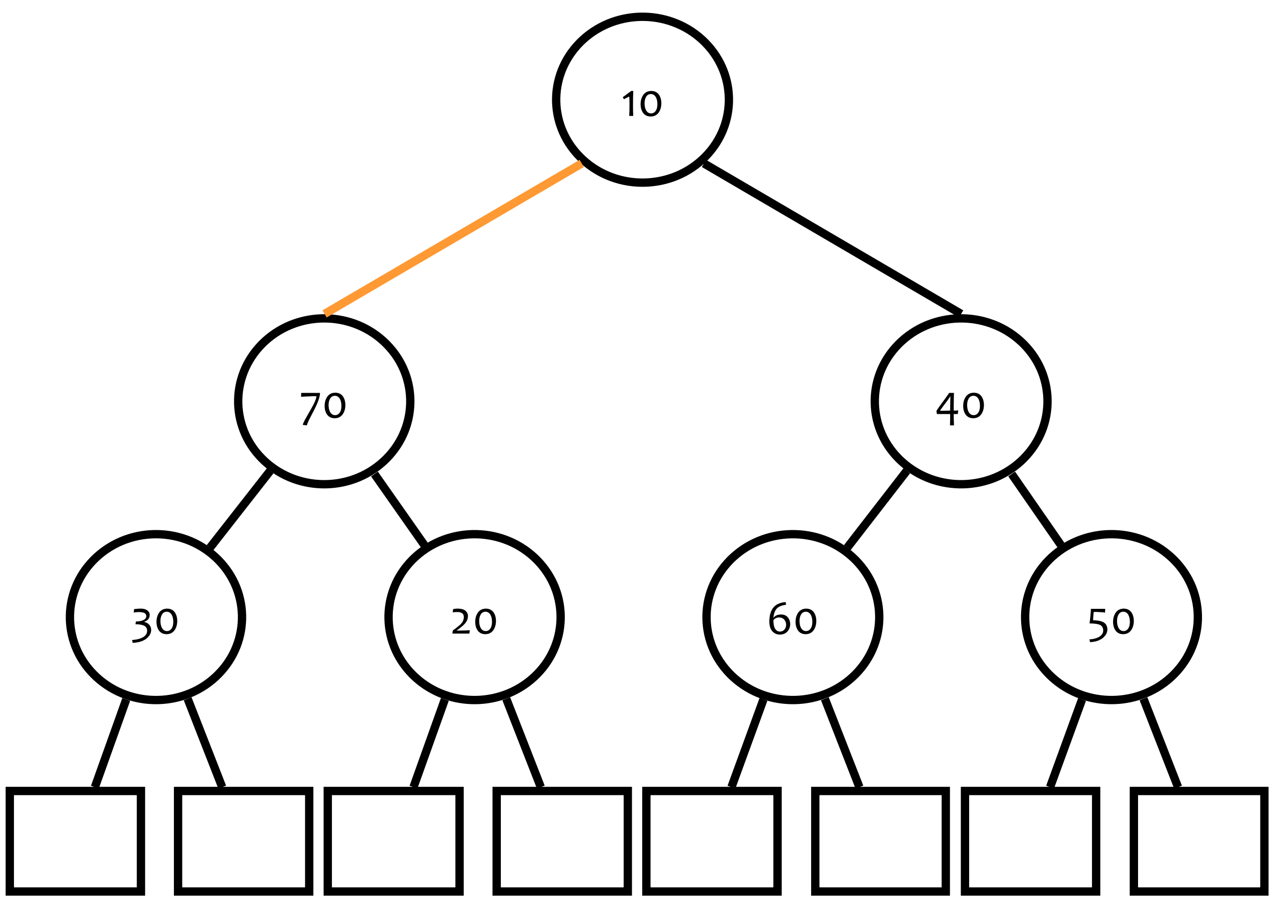

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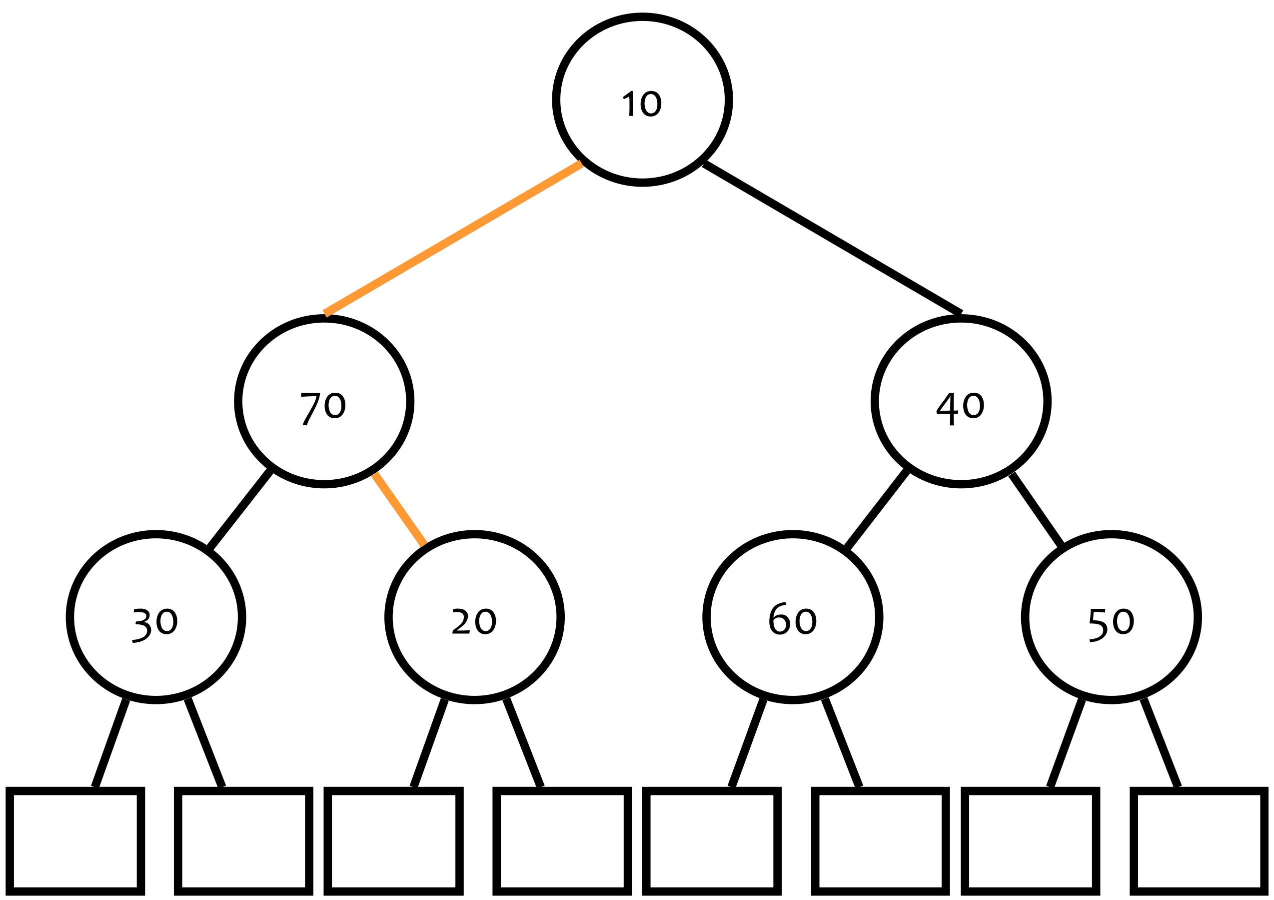

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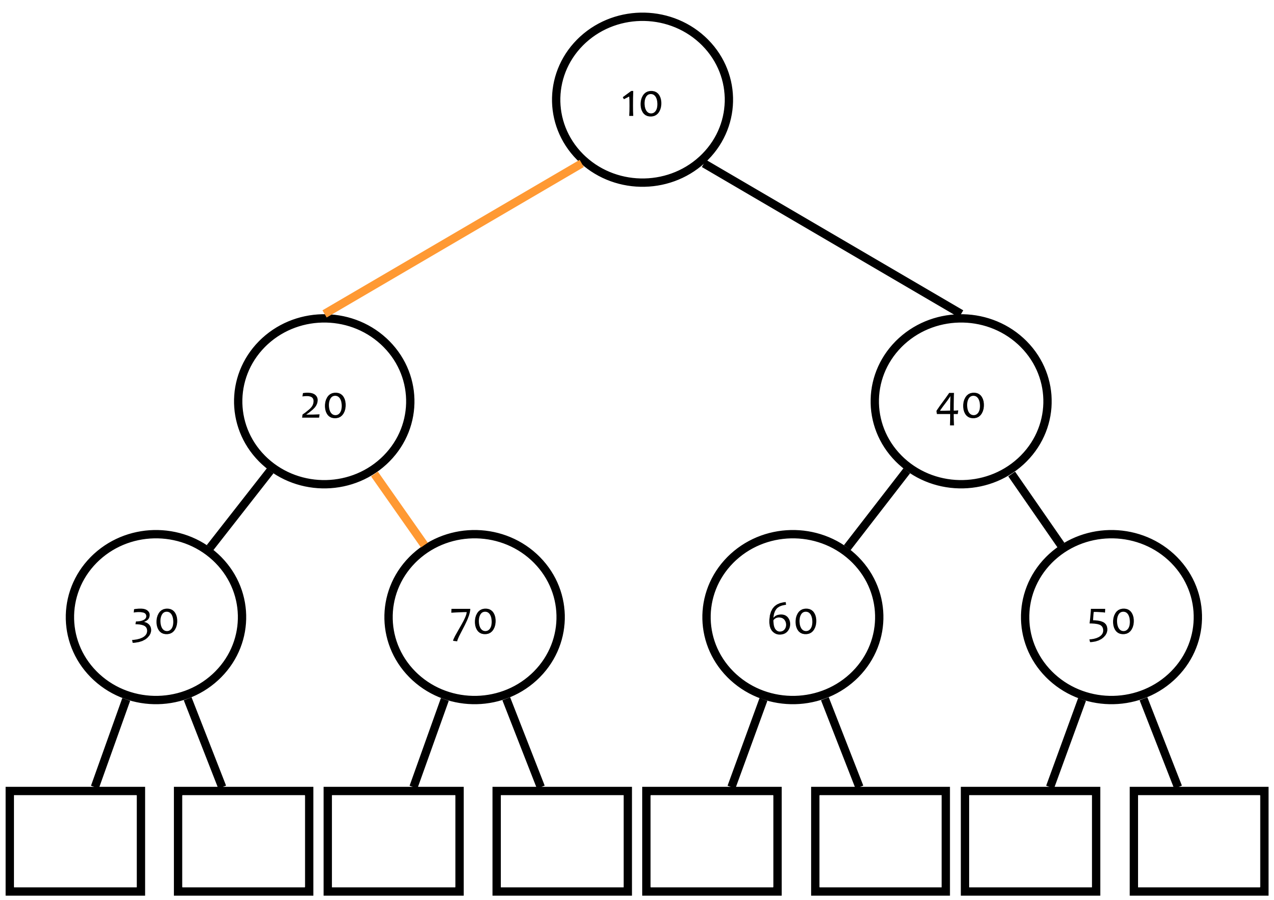

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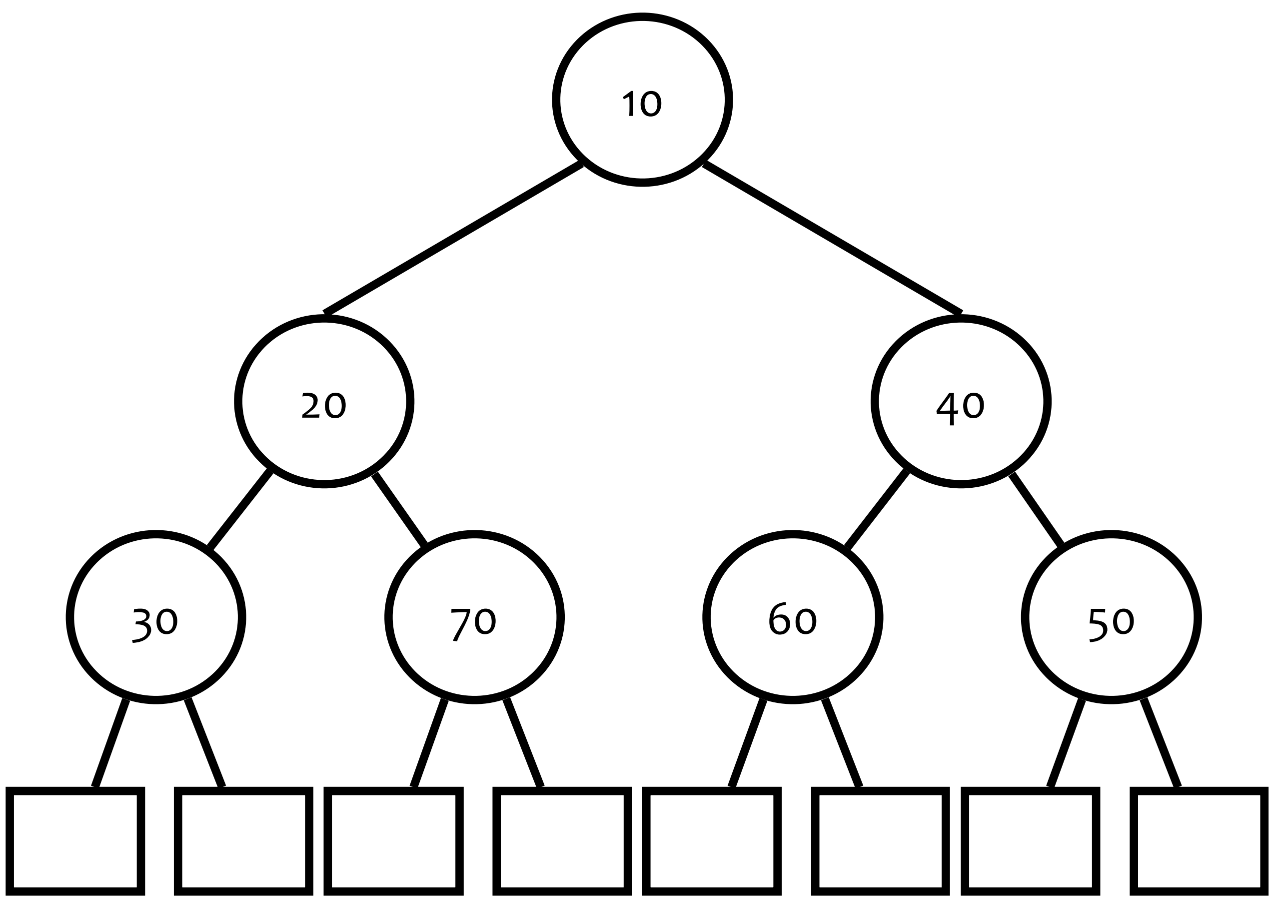

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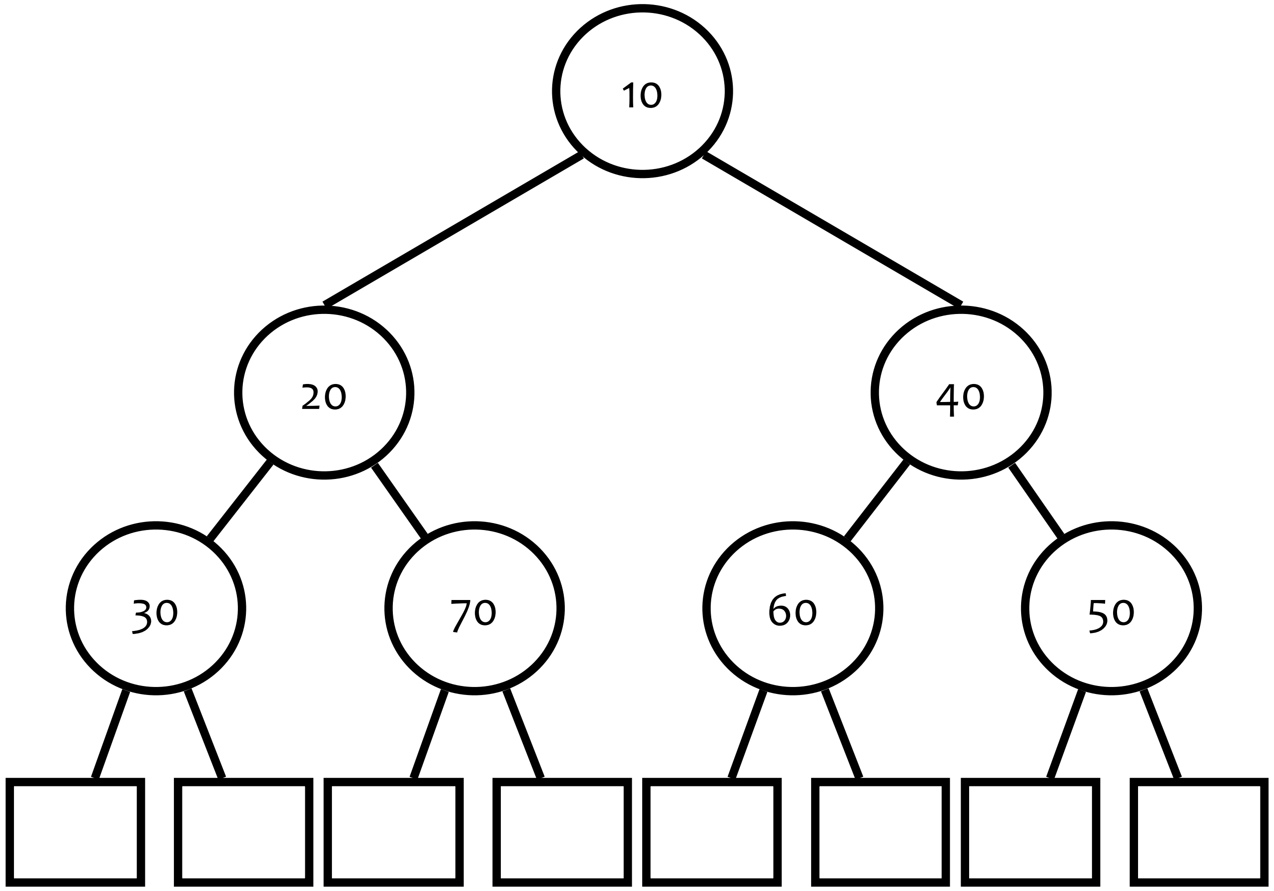

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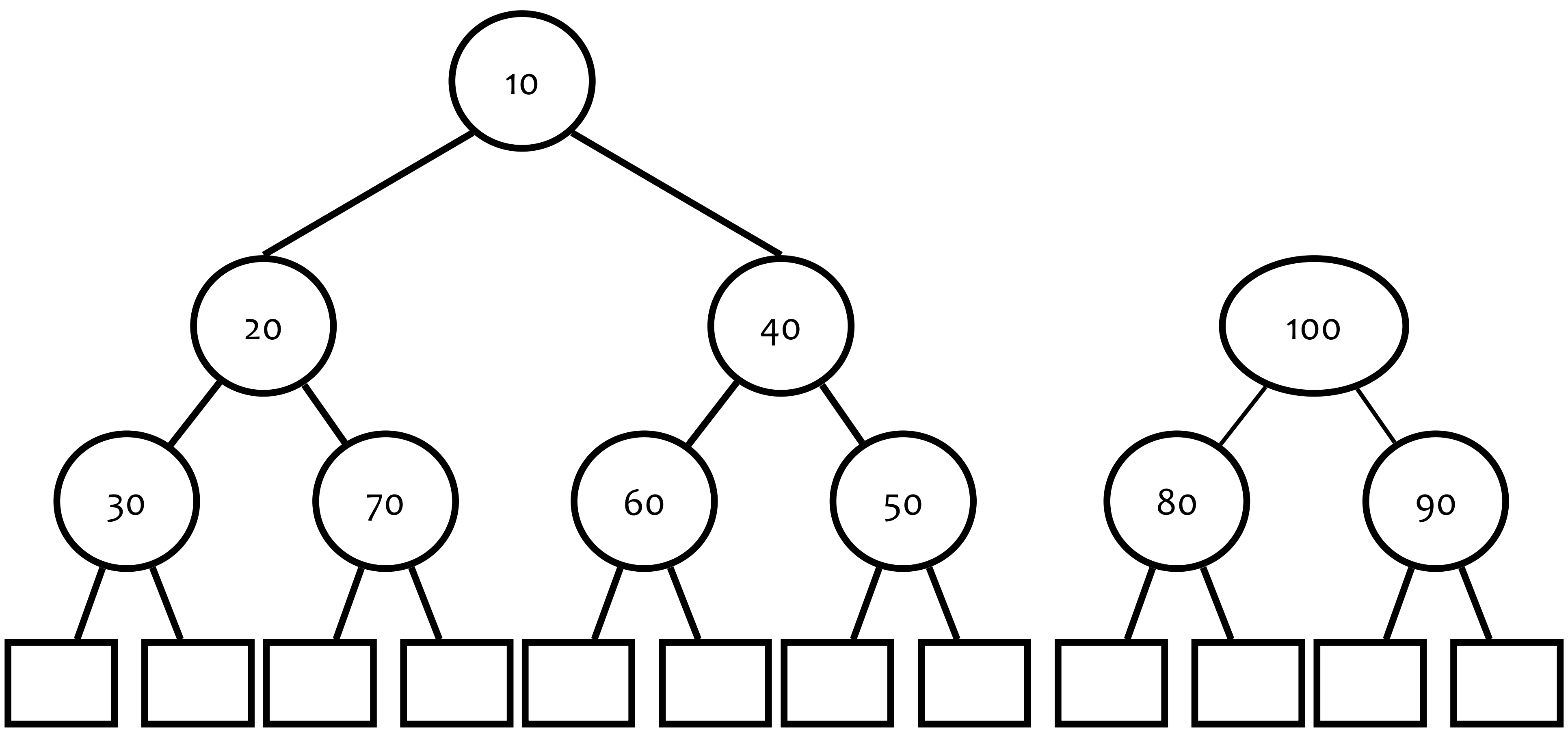

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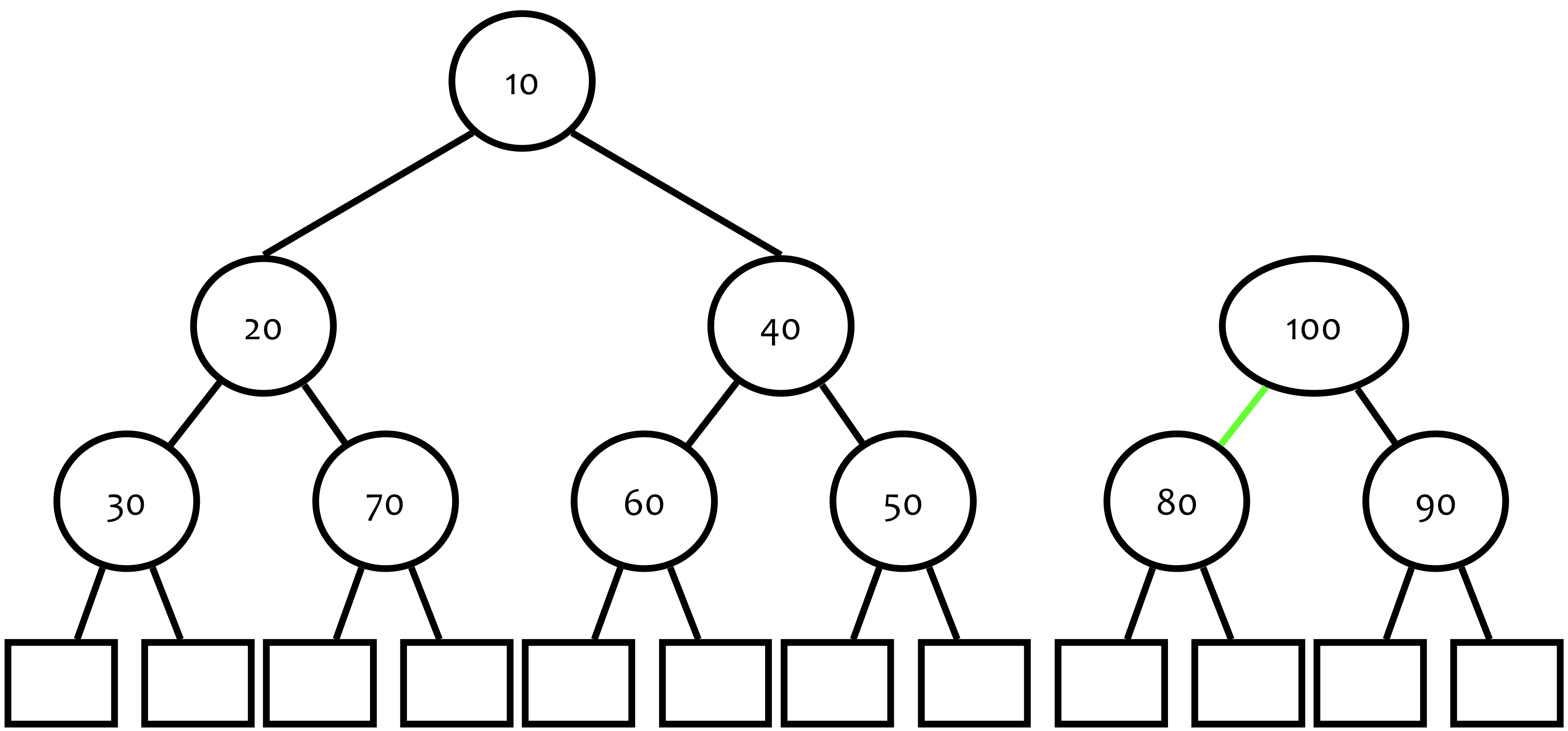

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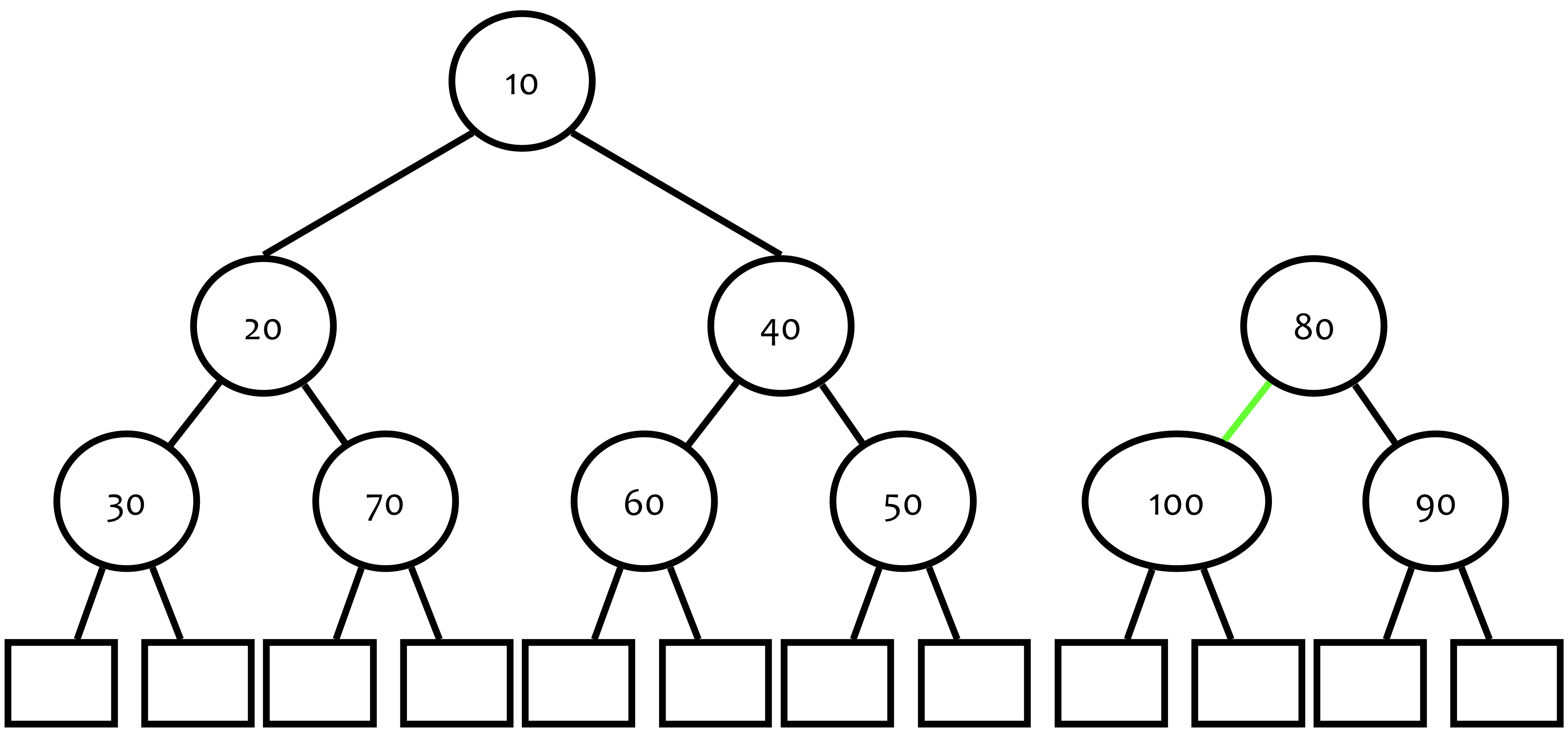

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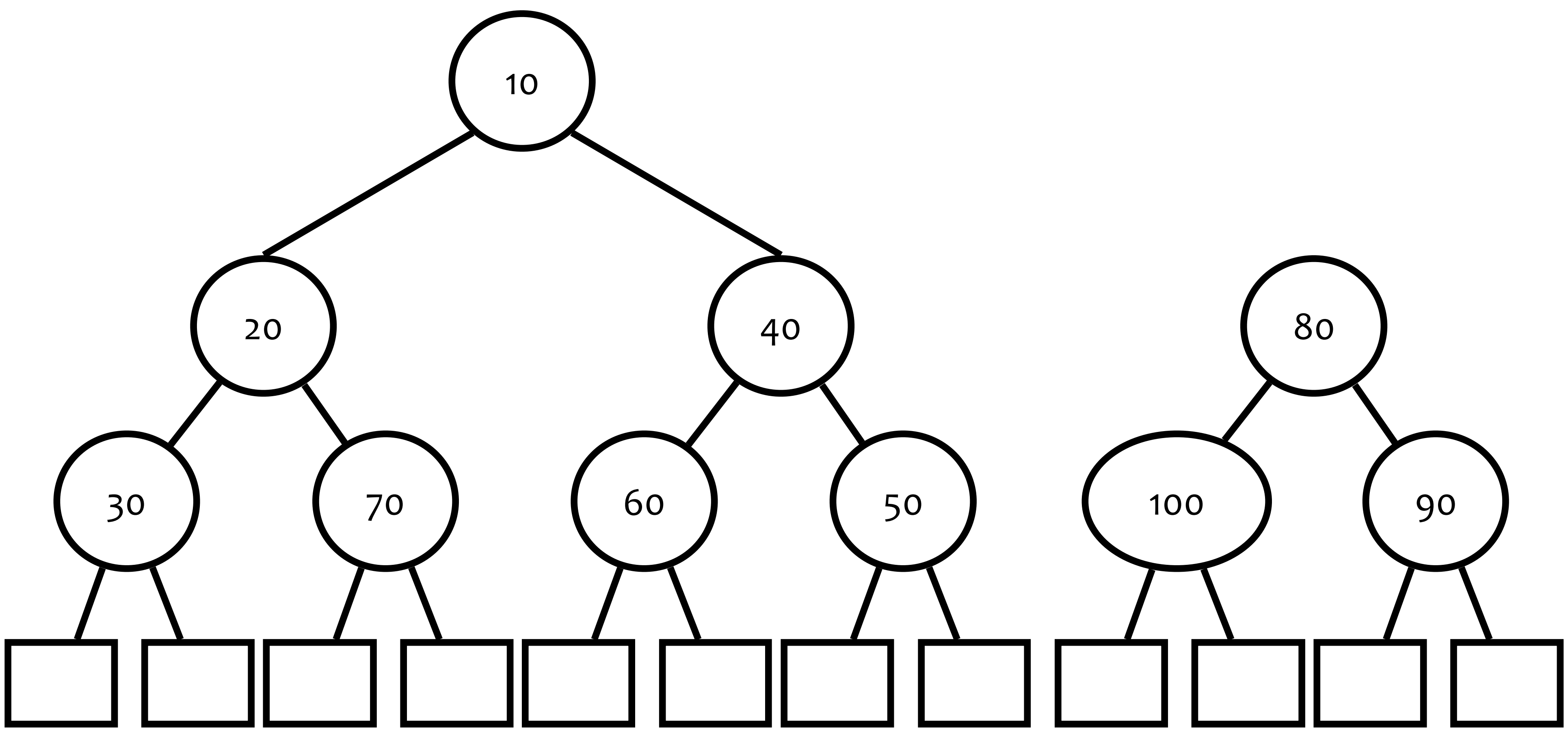

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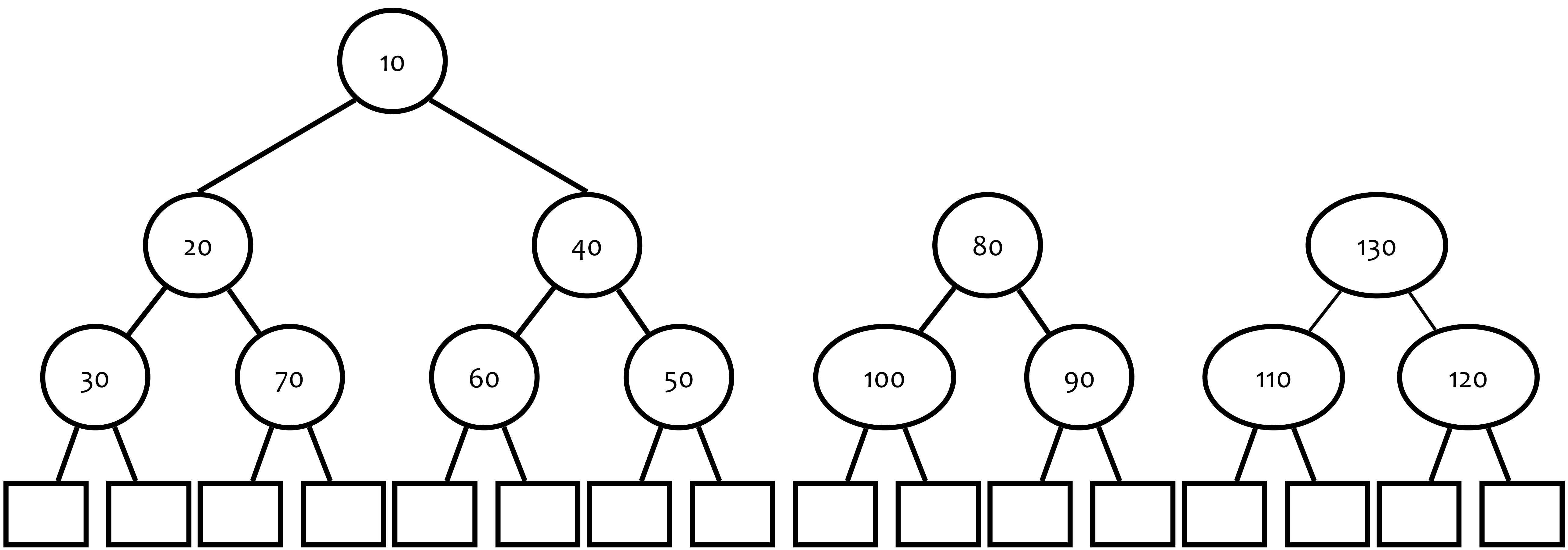

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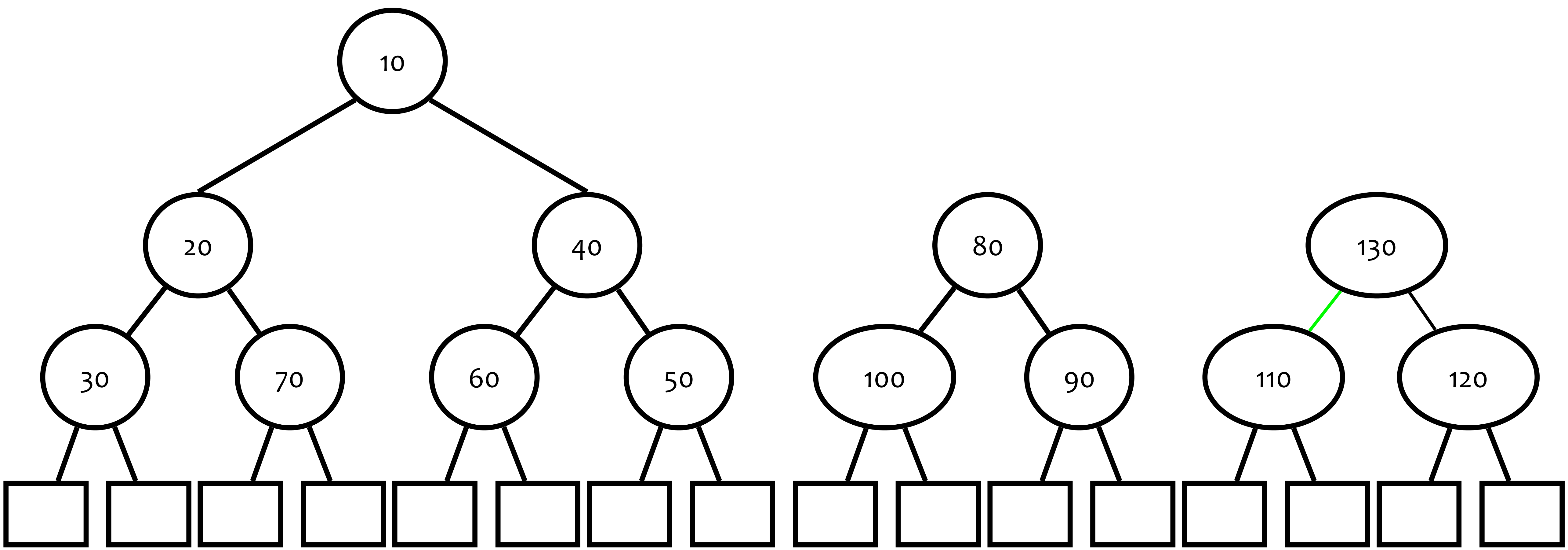

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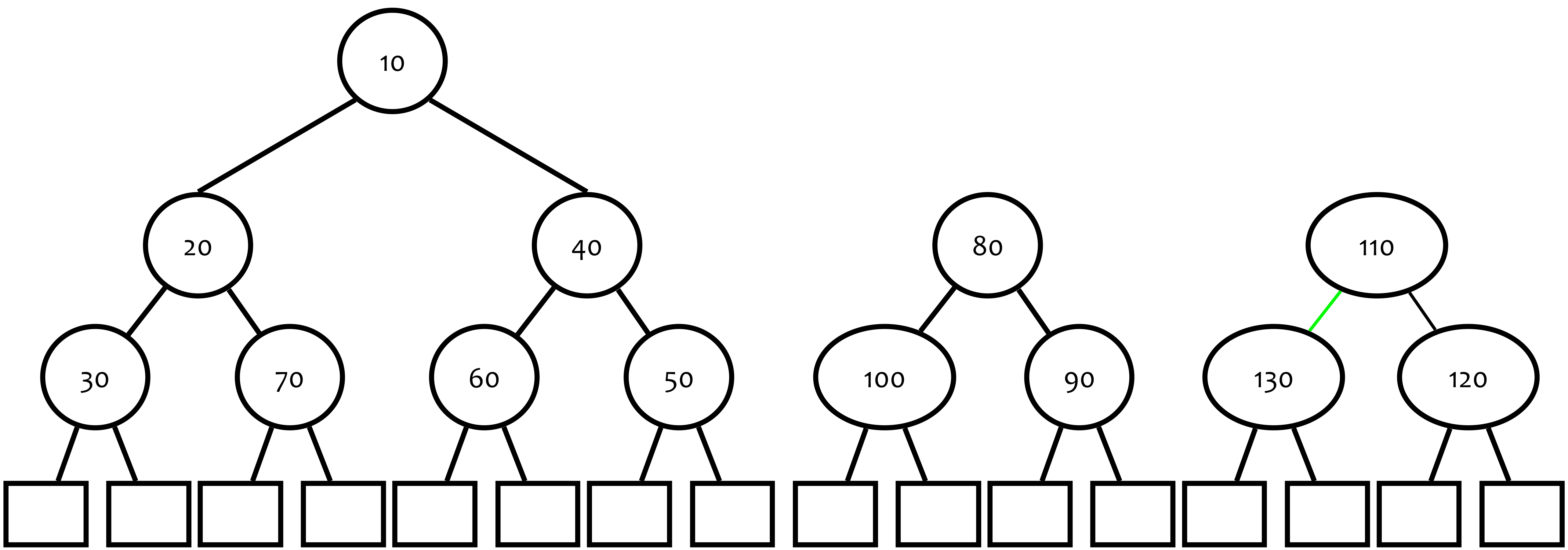

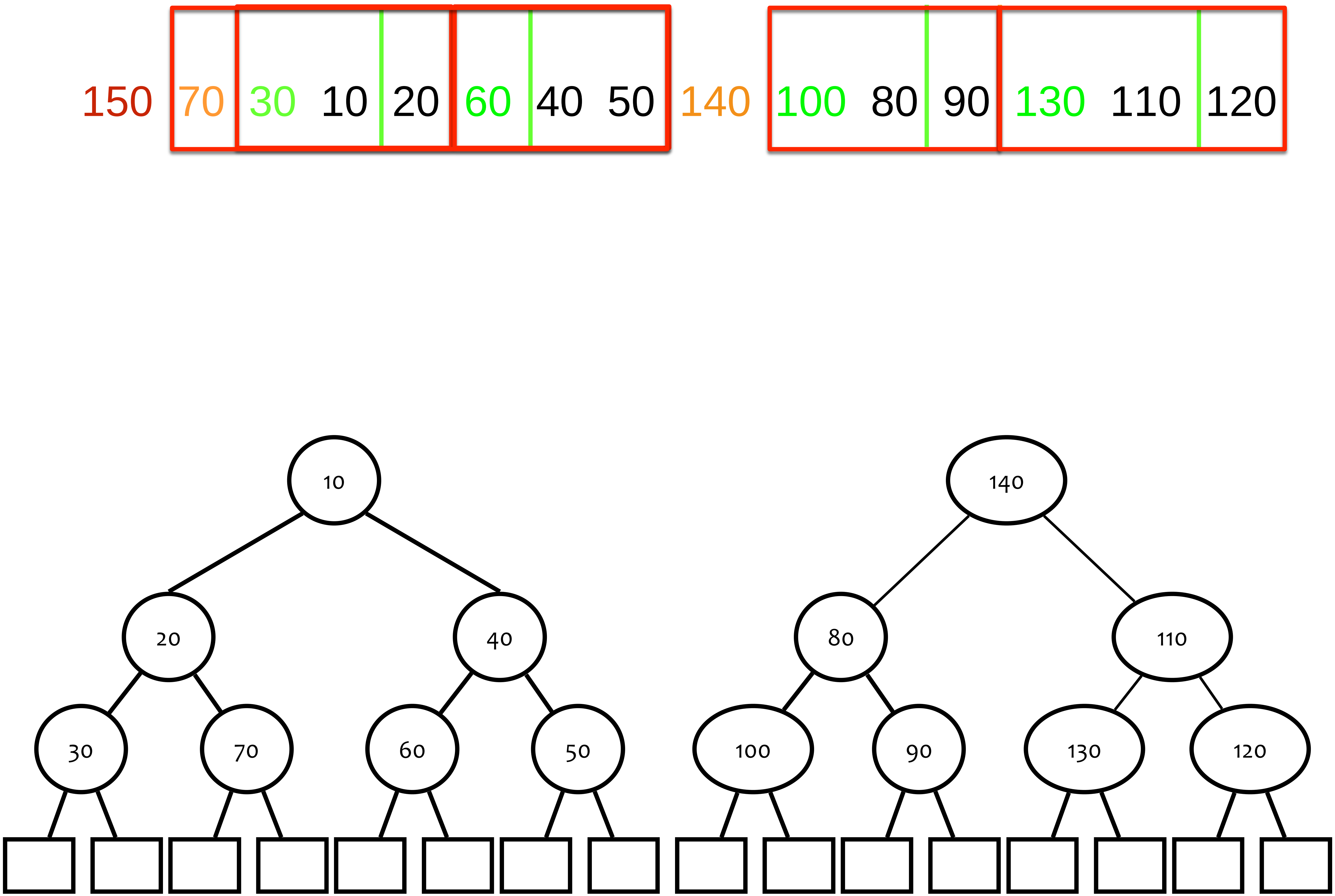
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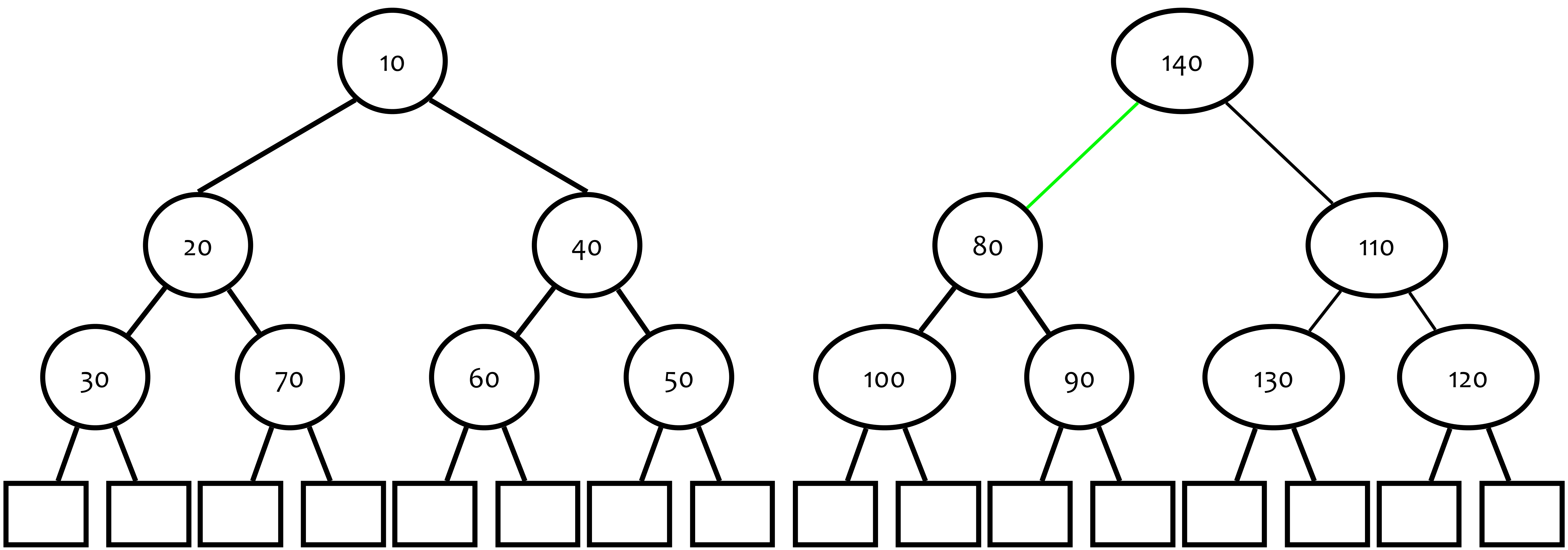

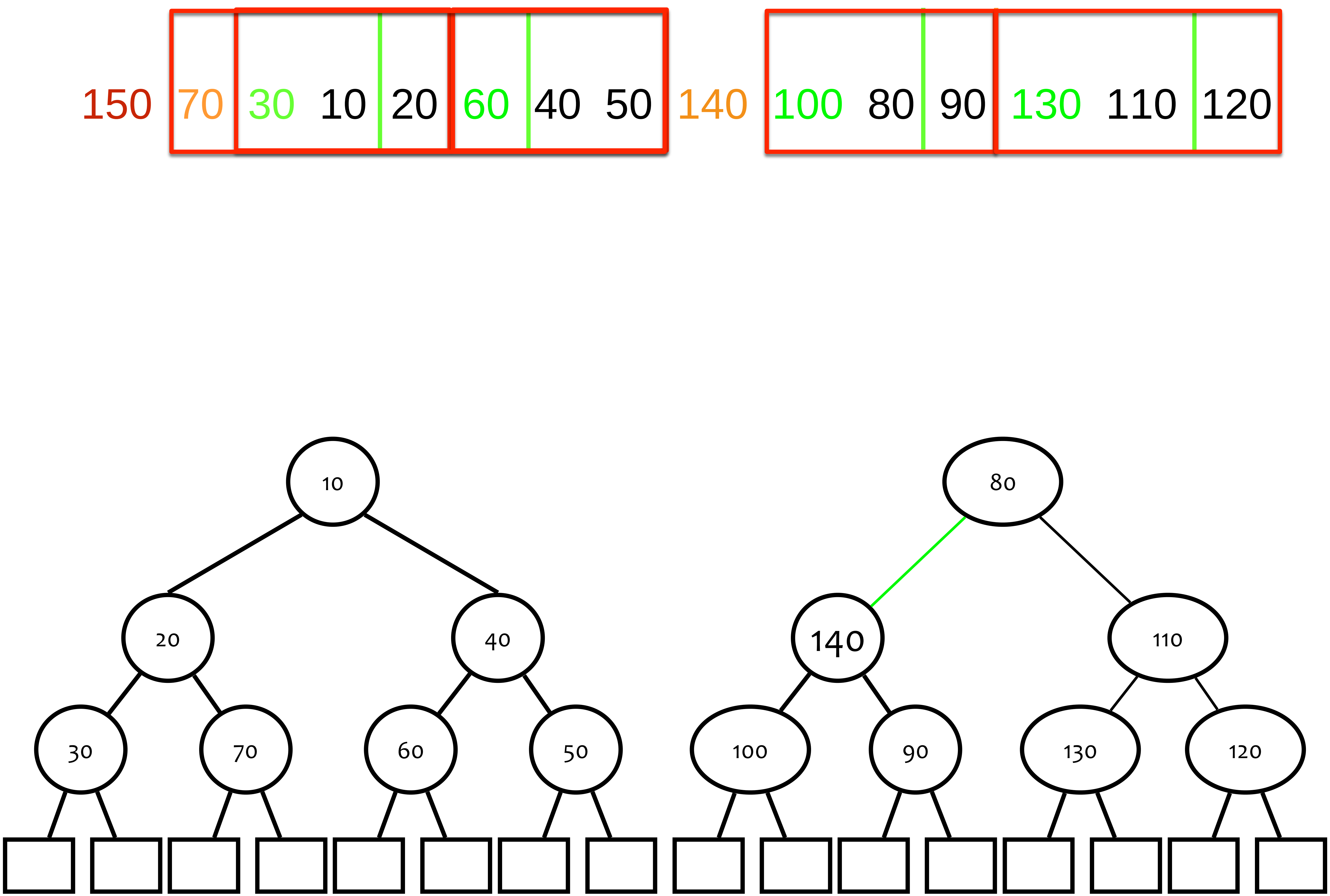
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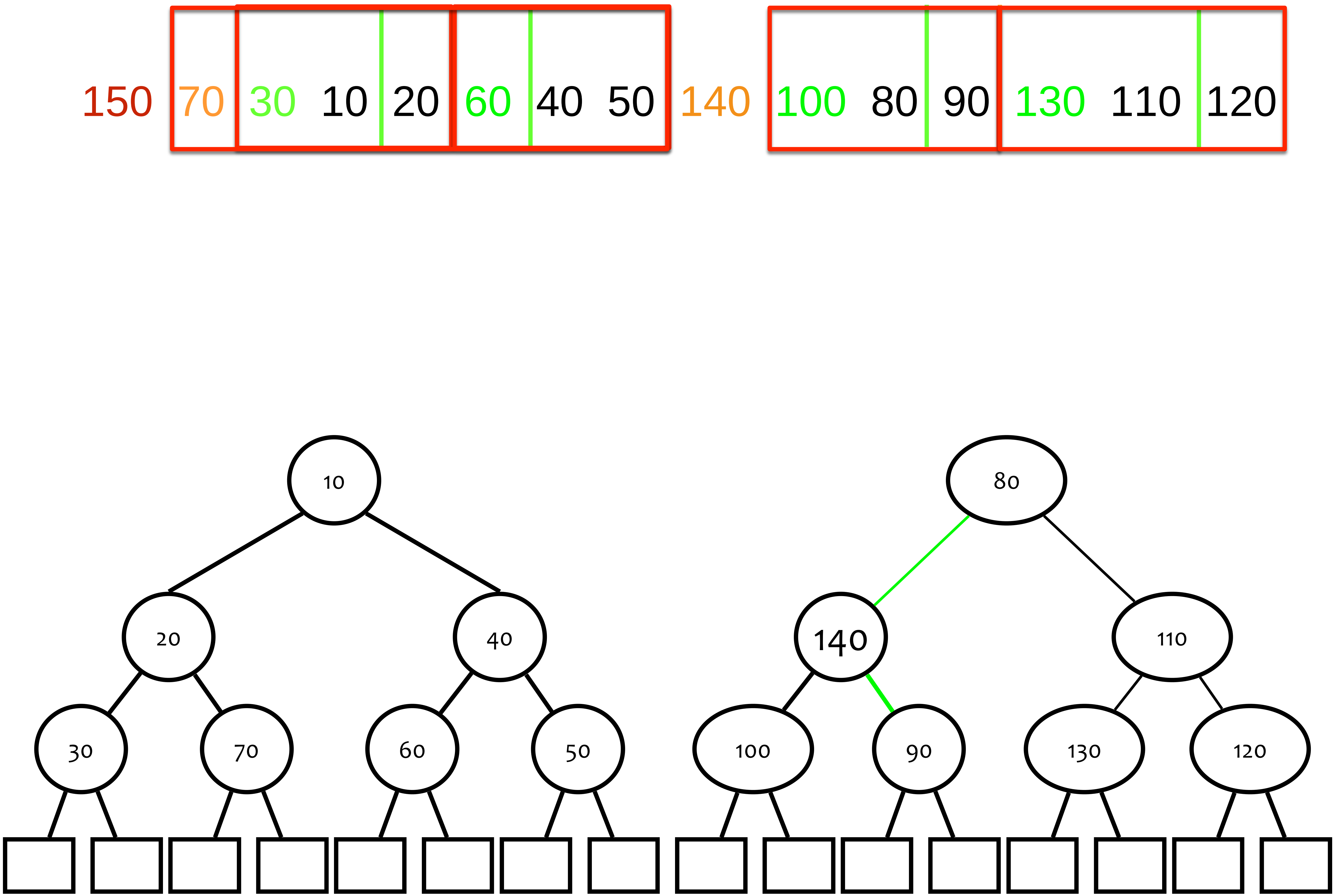
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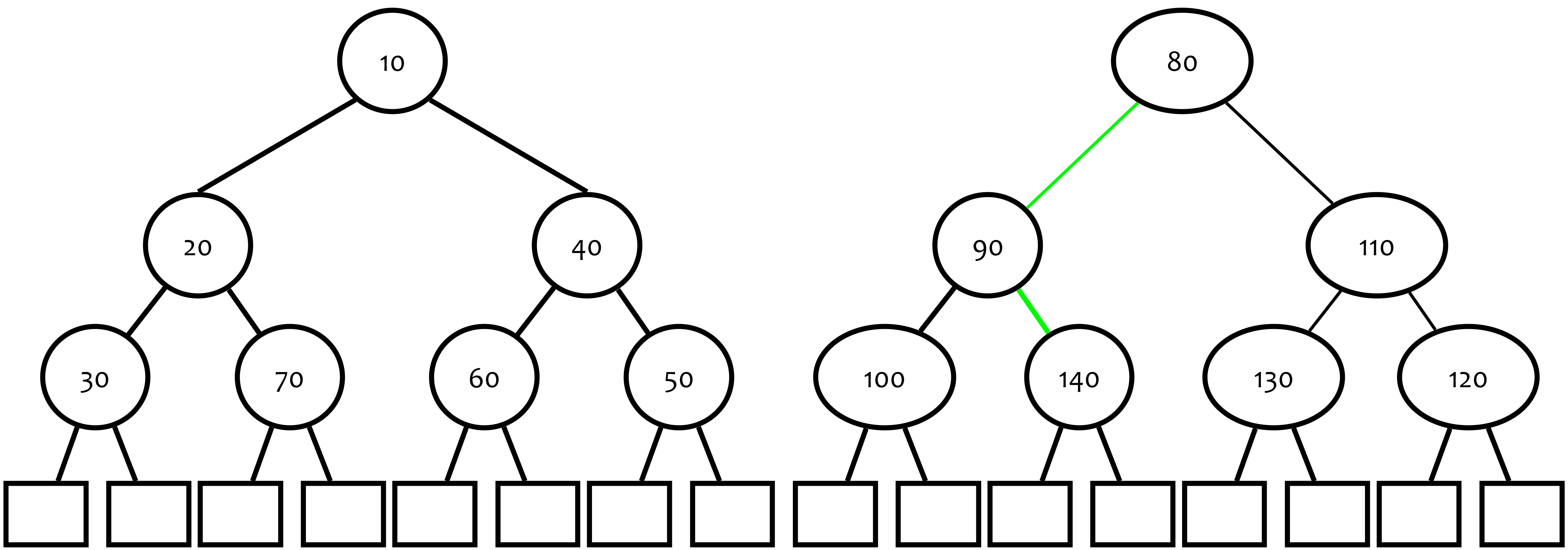

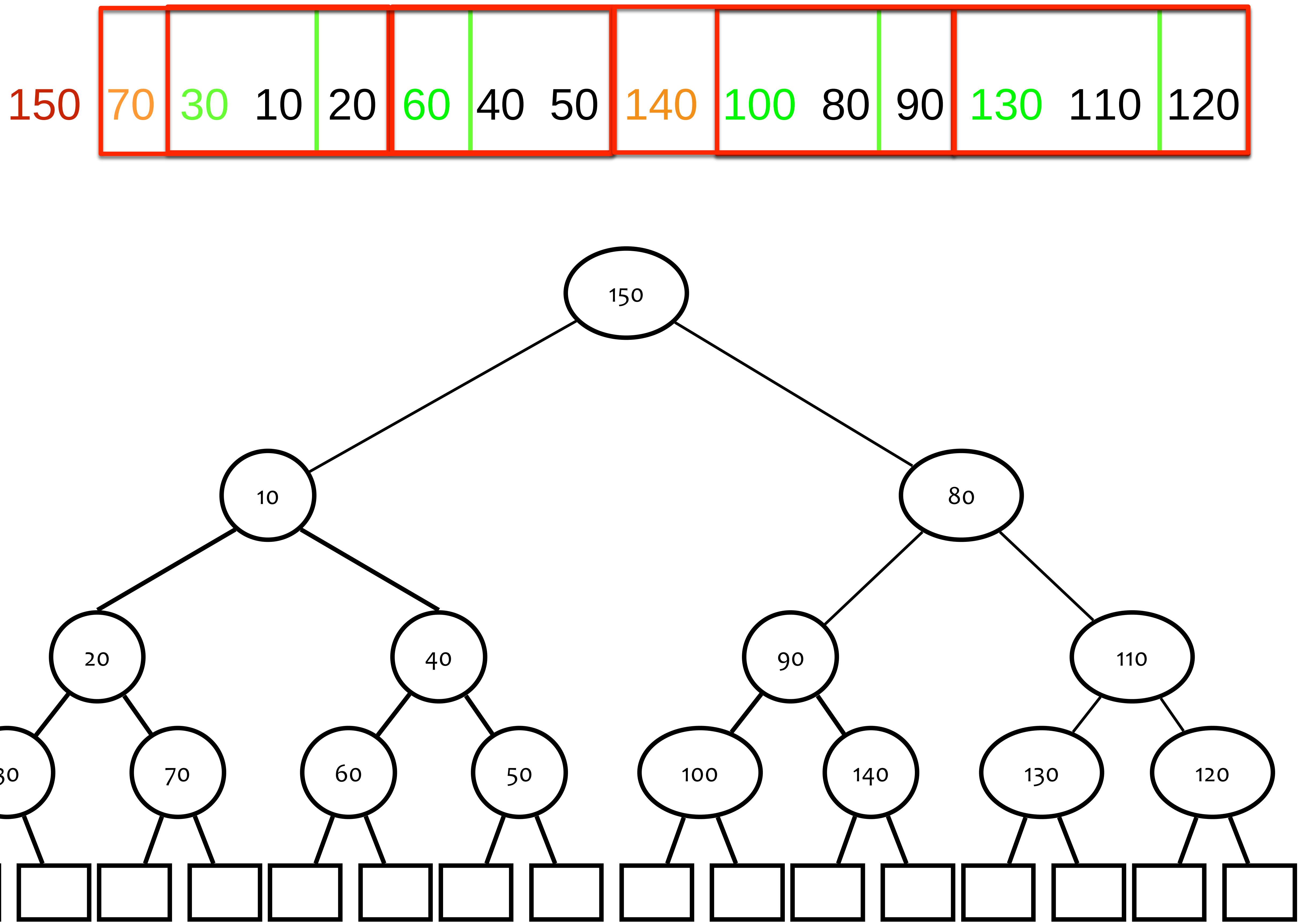

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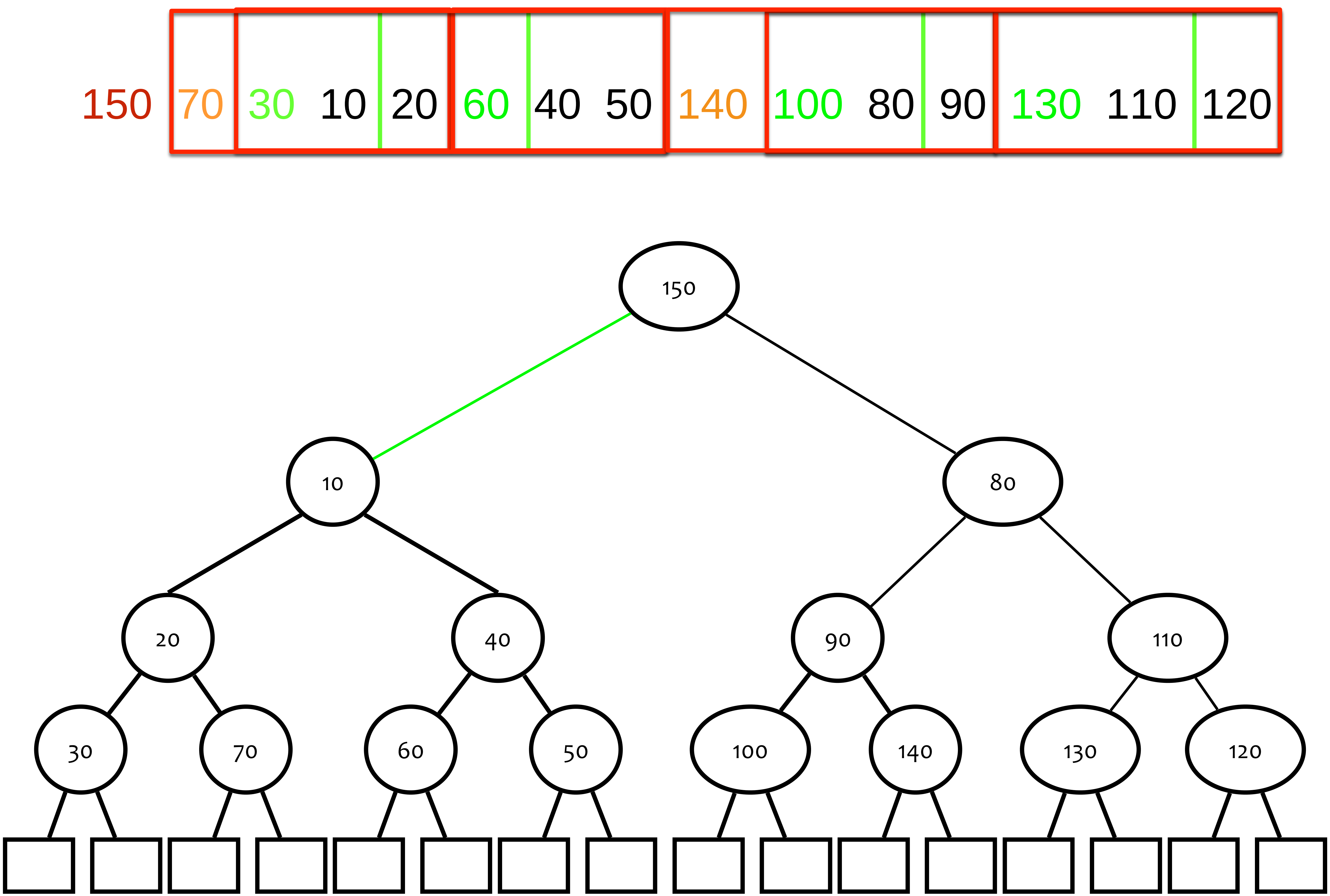
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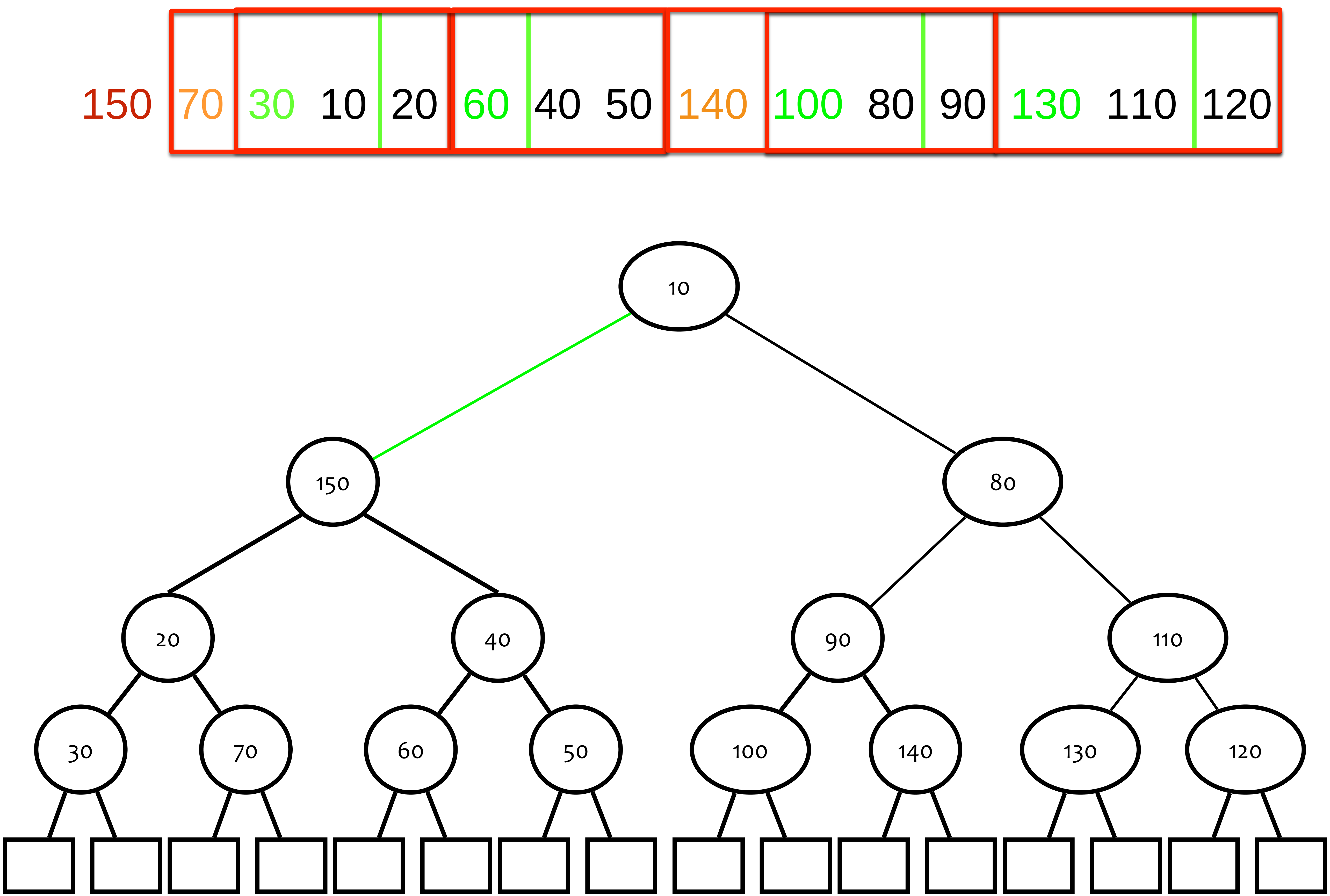

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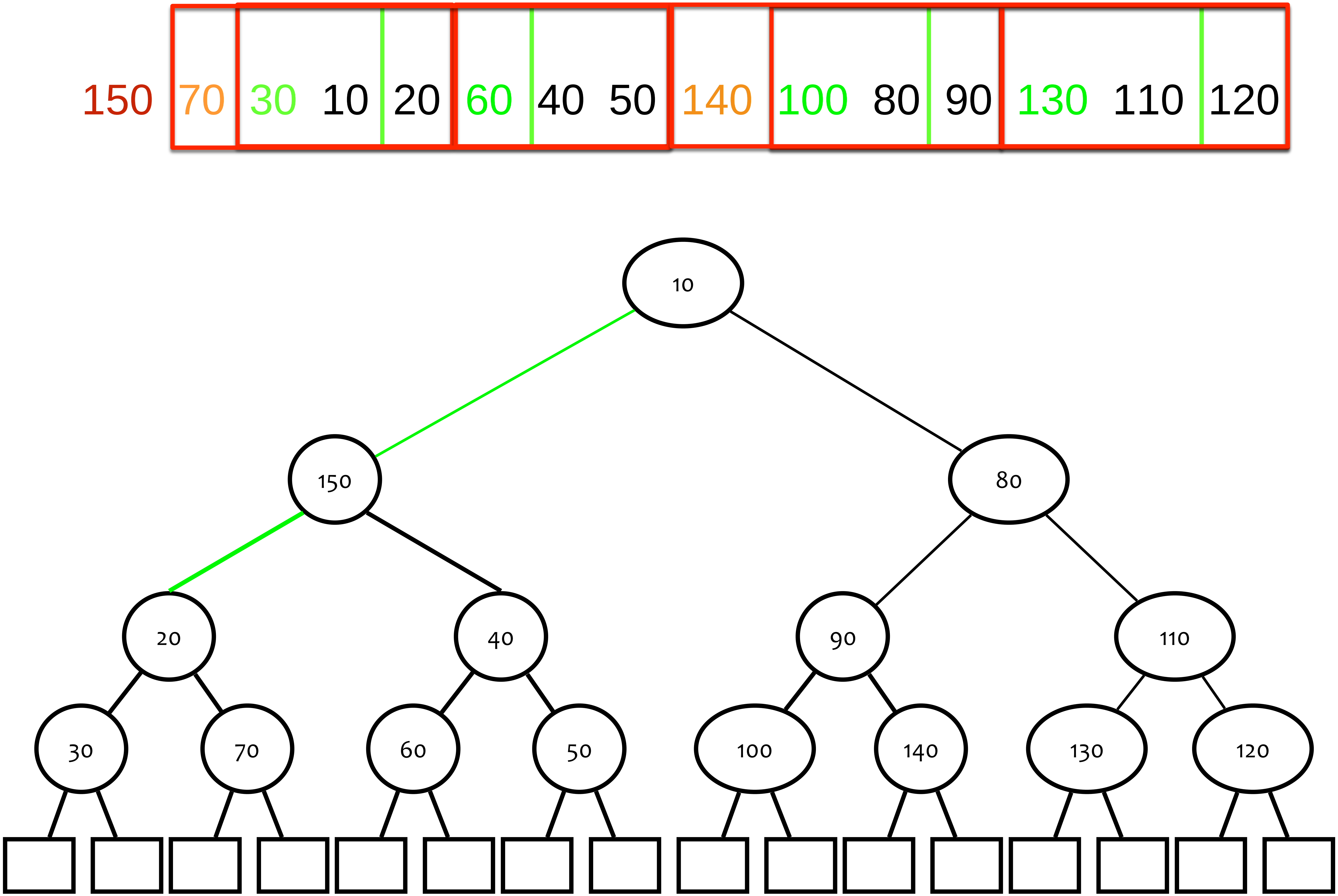
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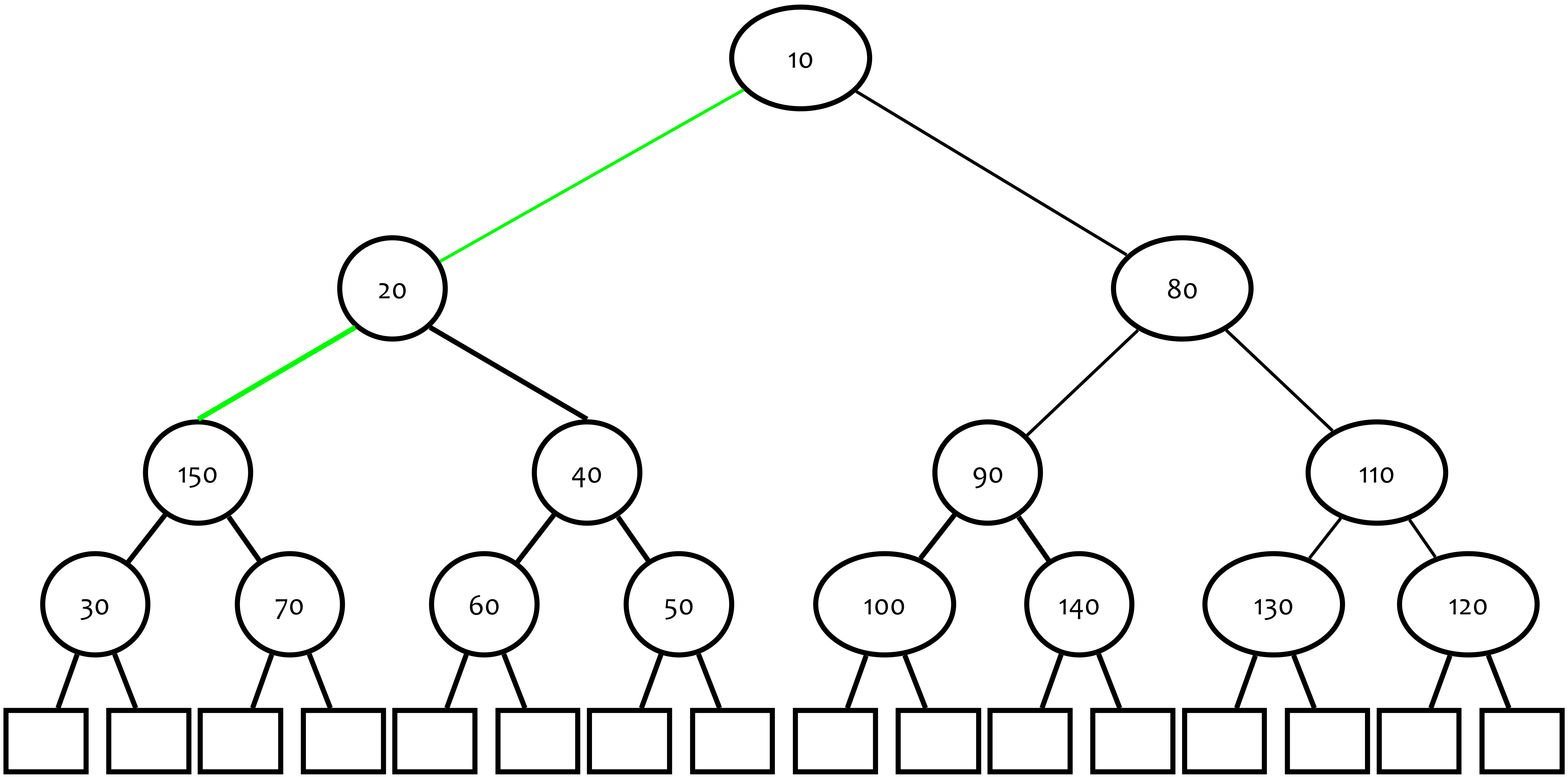

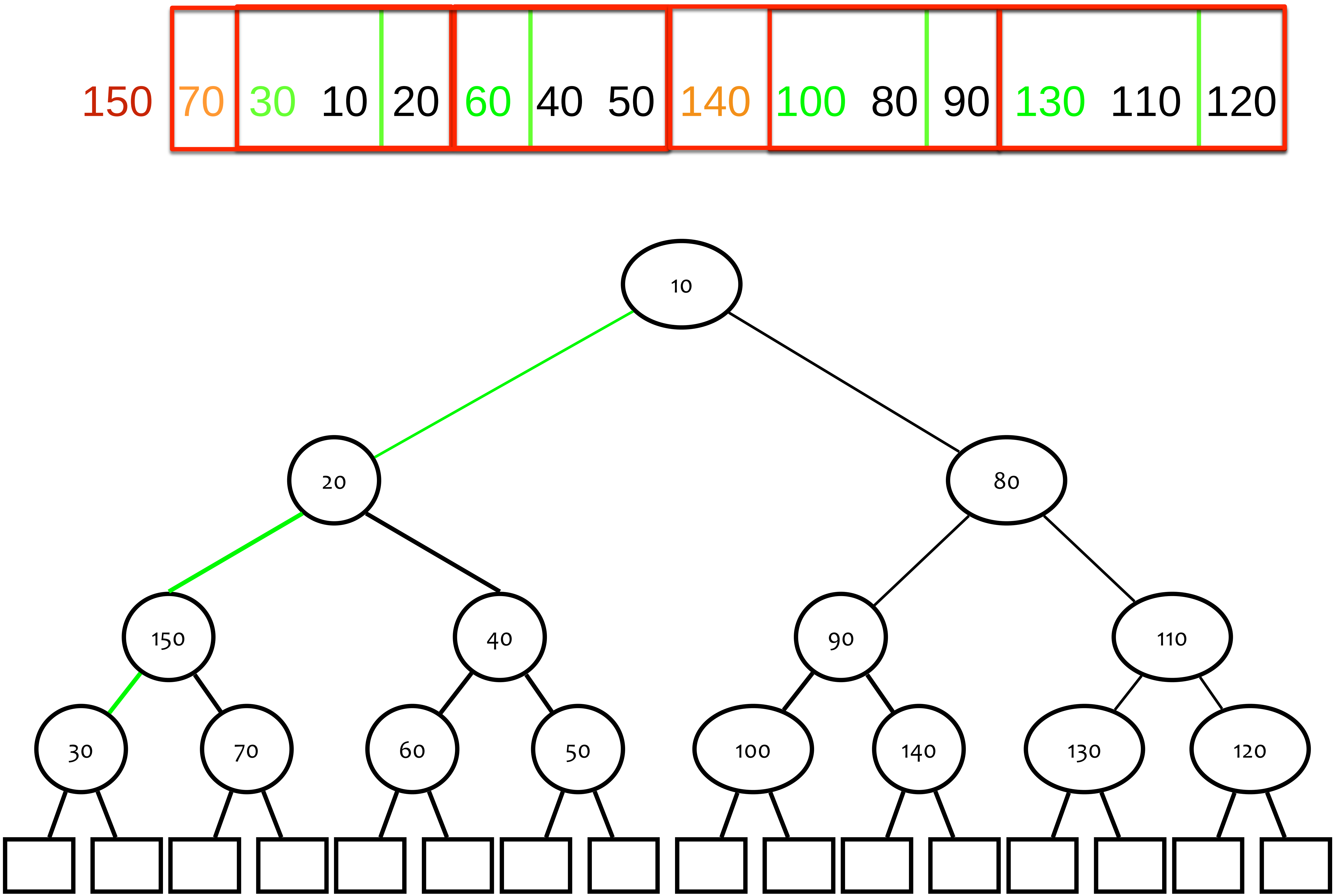
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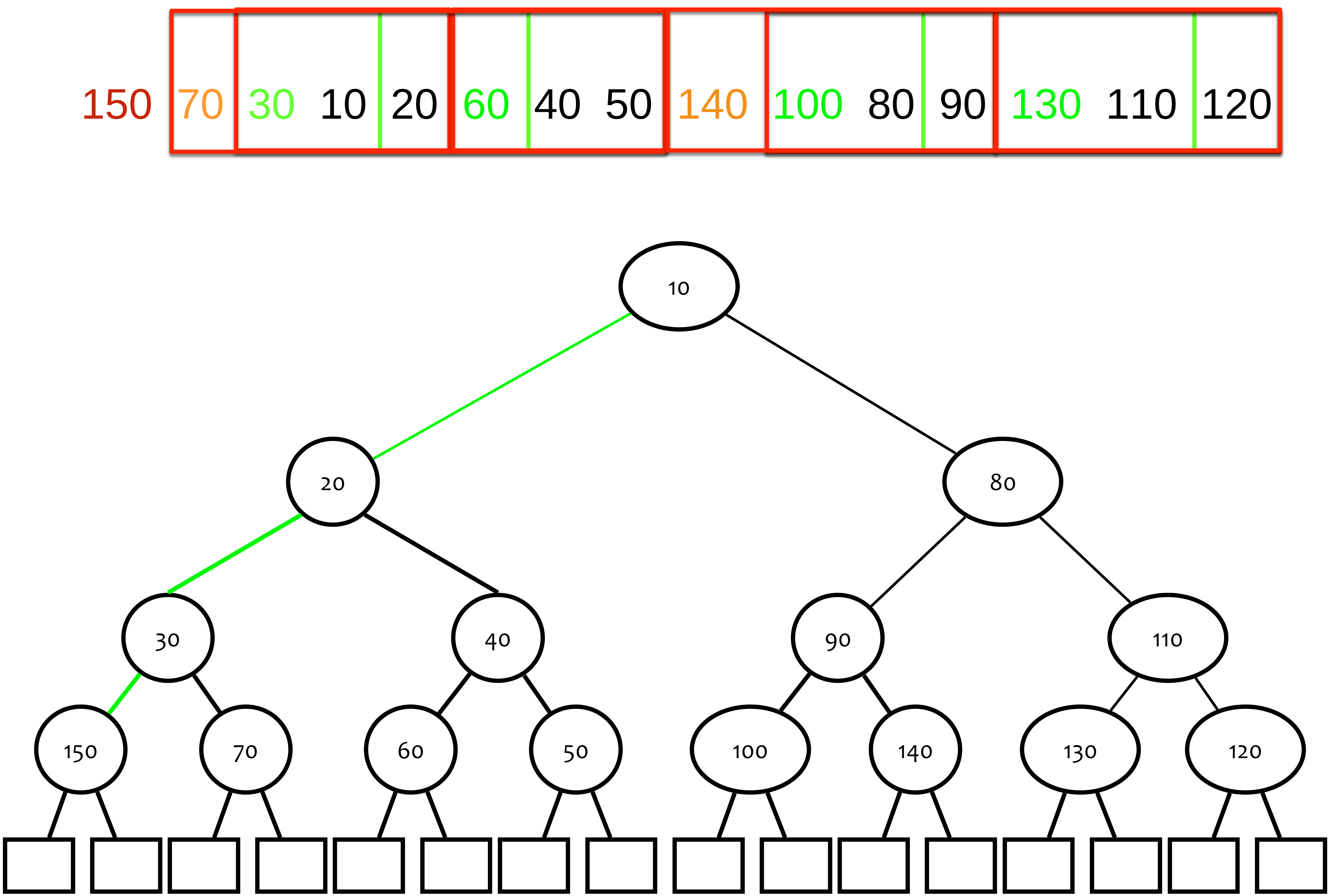
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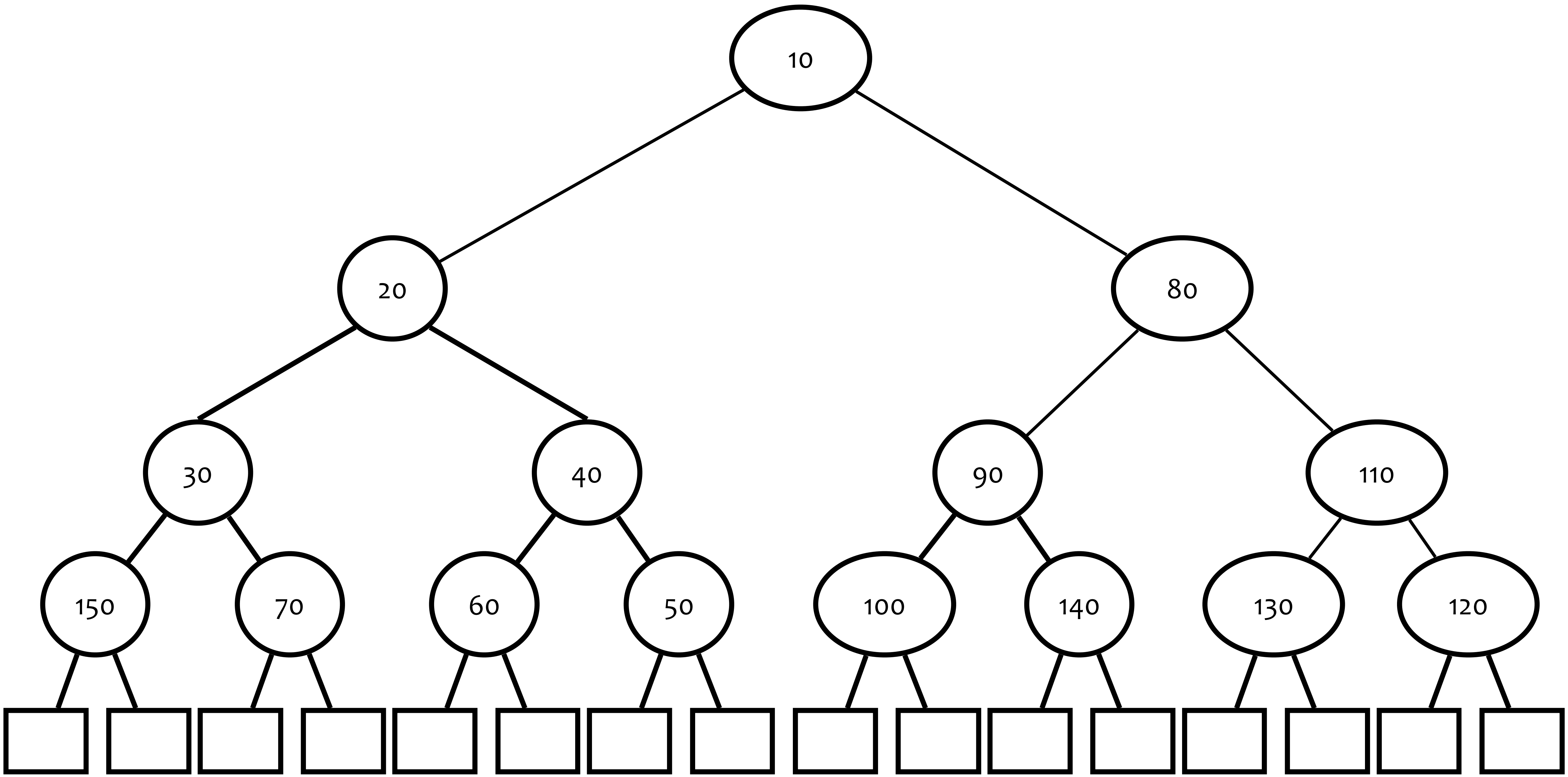
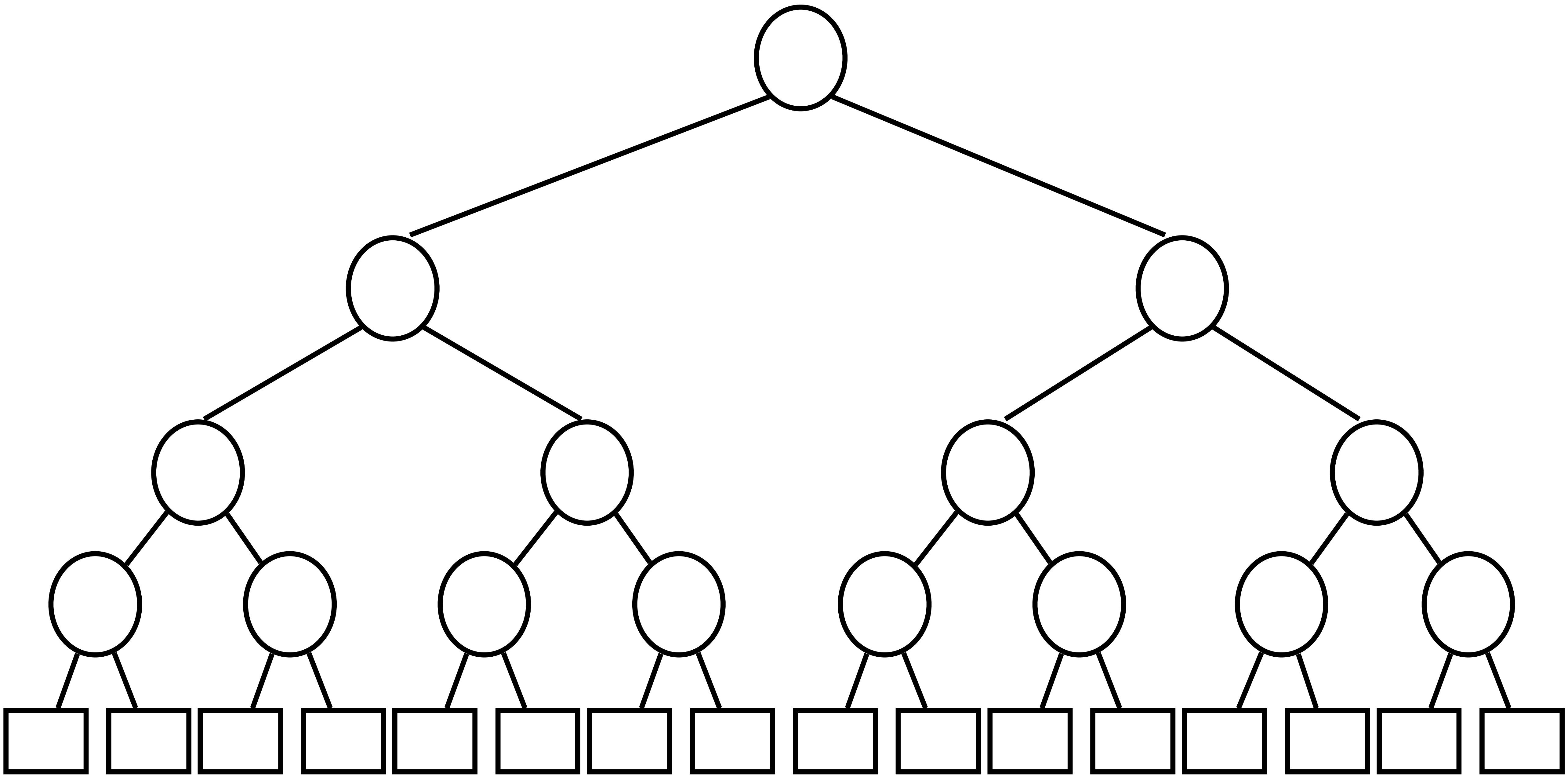
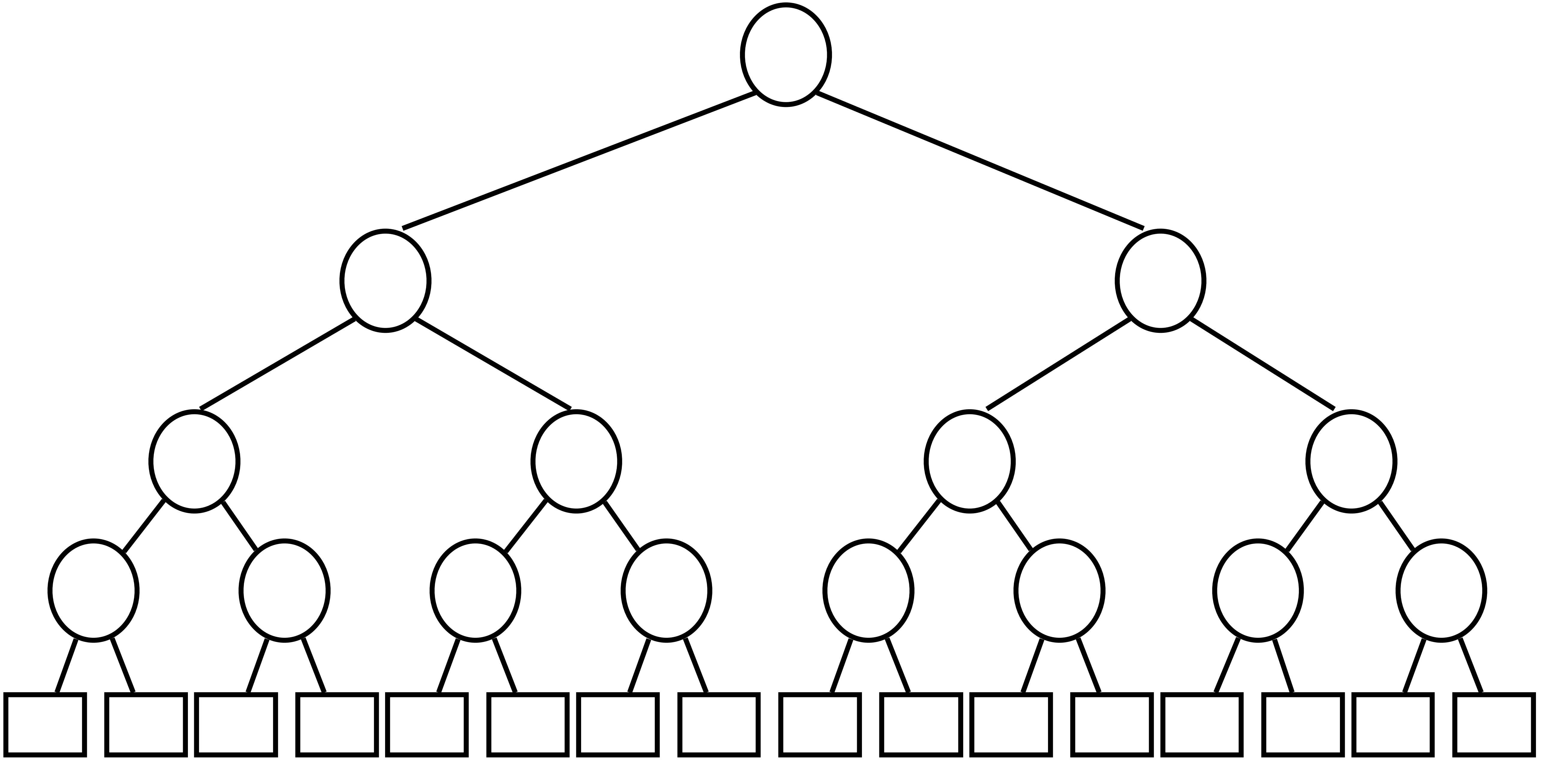


Did we really insert all m elements in $O(m)$ time??



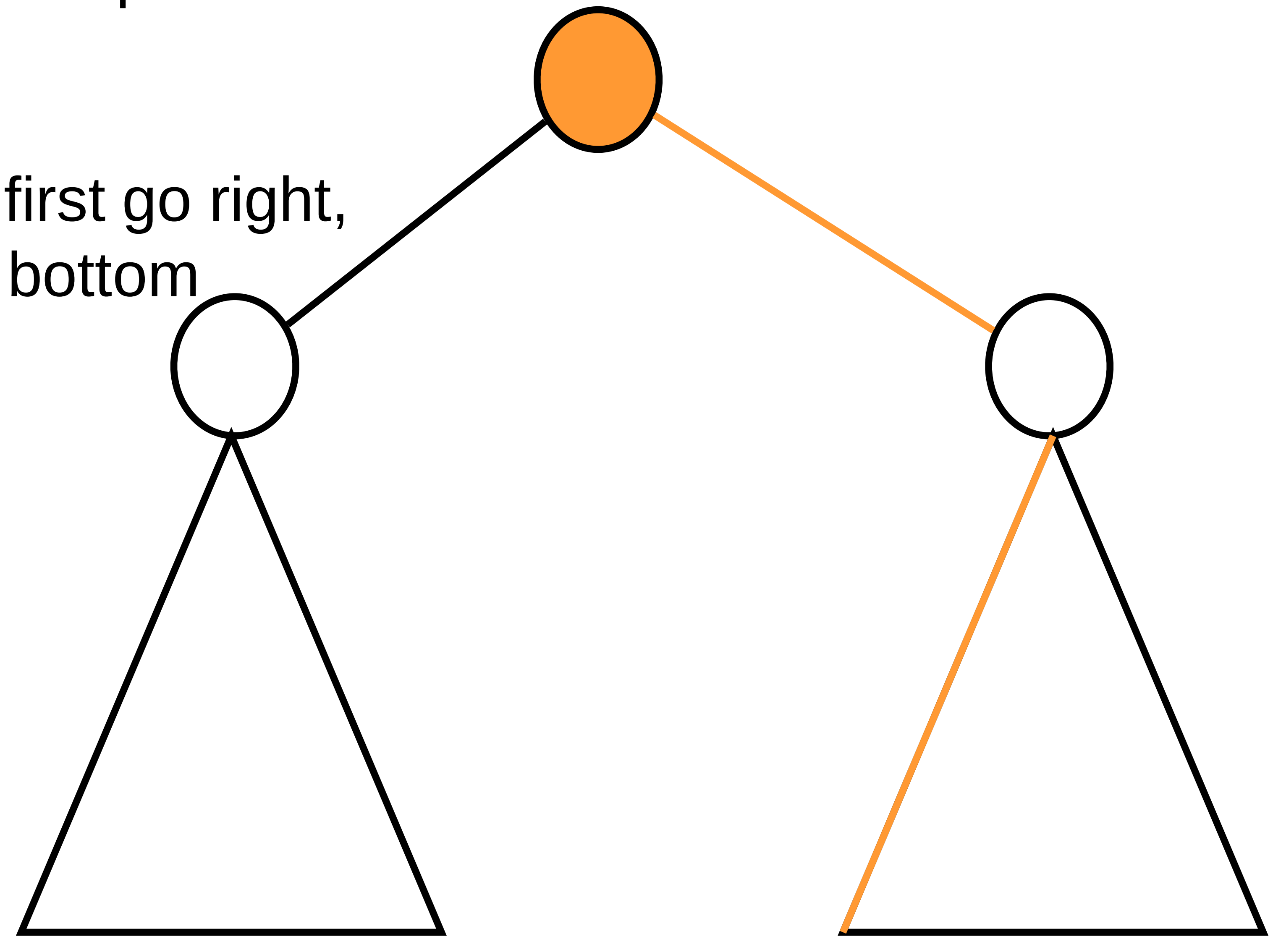
We show: max# bubble down ops < # heap edges



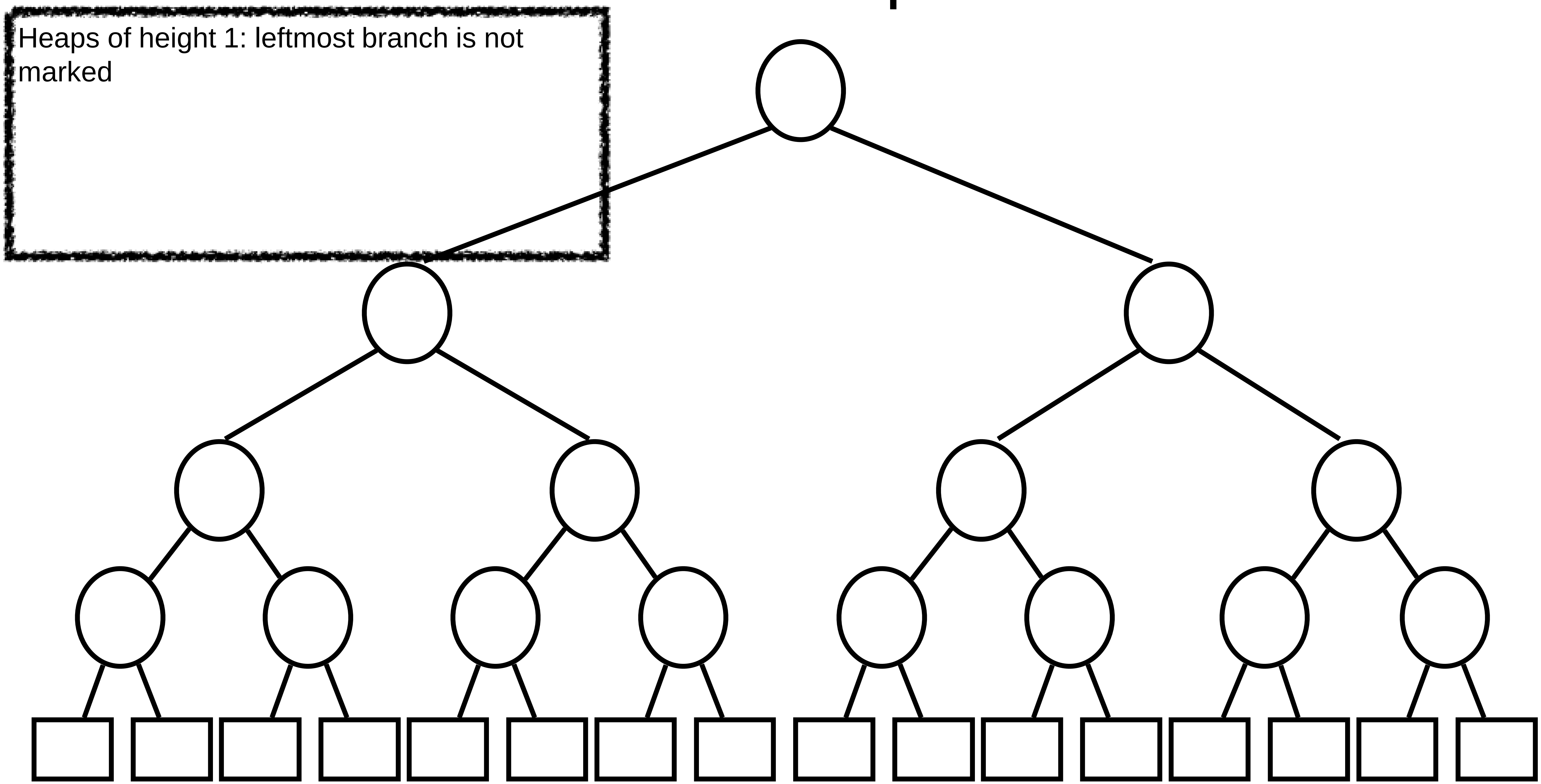
Proof idea

For each new node joining two heaps: mark path of maximum number of bubble-down operations

For marking: first go right, then left until bottom



For each new node joining two heaps: mark path with maximum number of bubble-down operations



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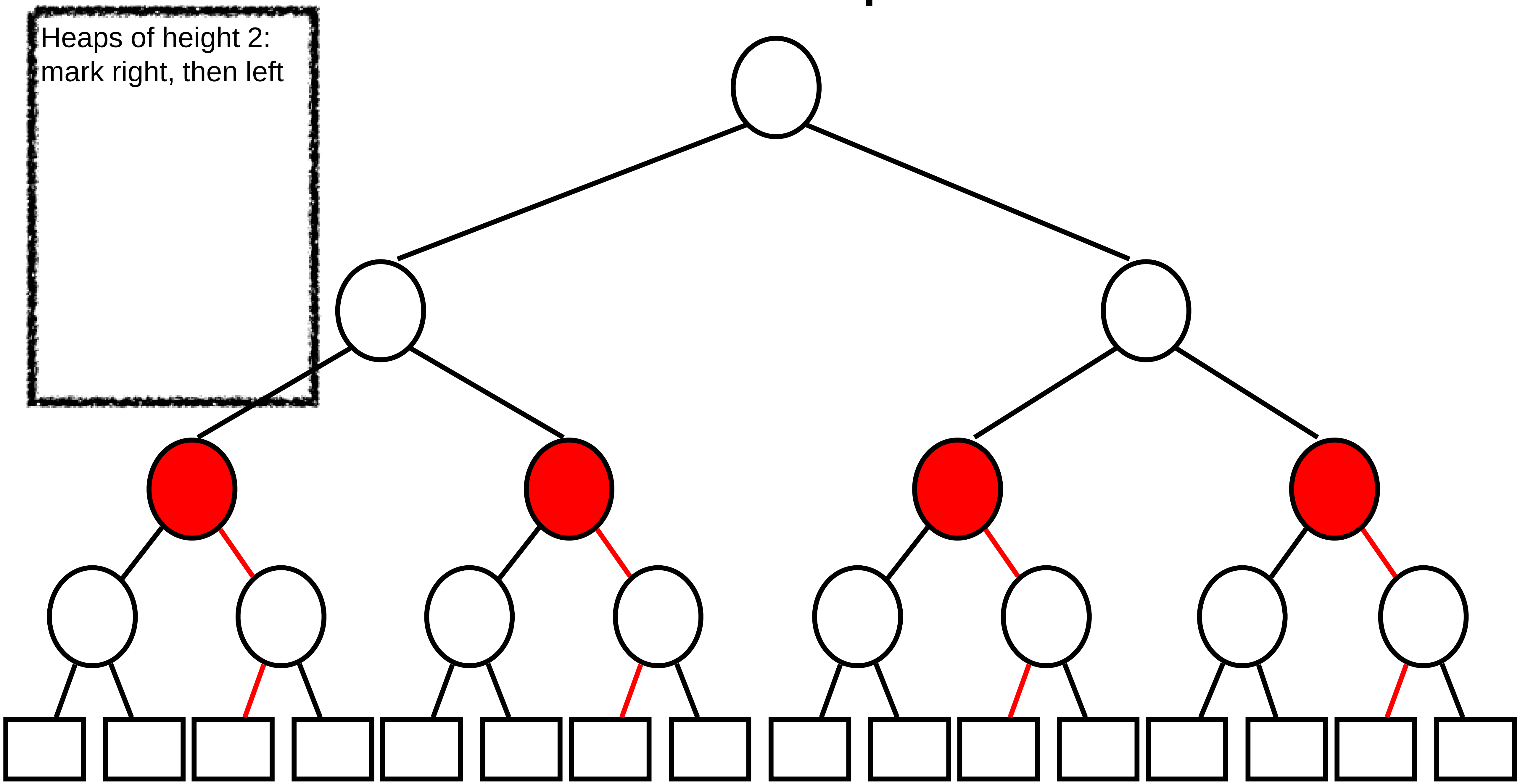
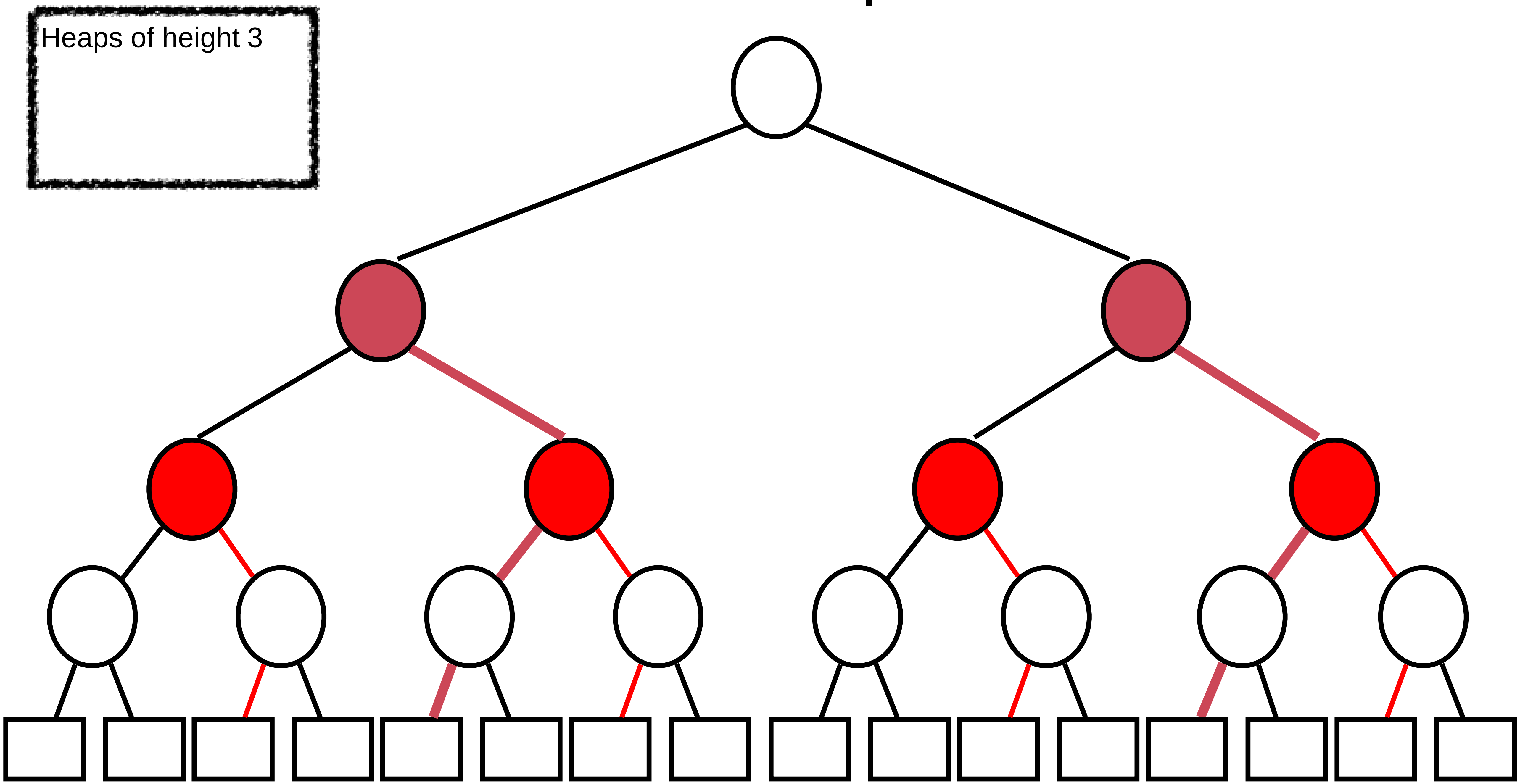


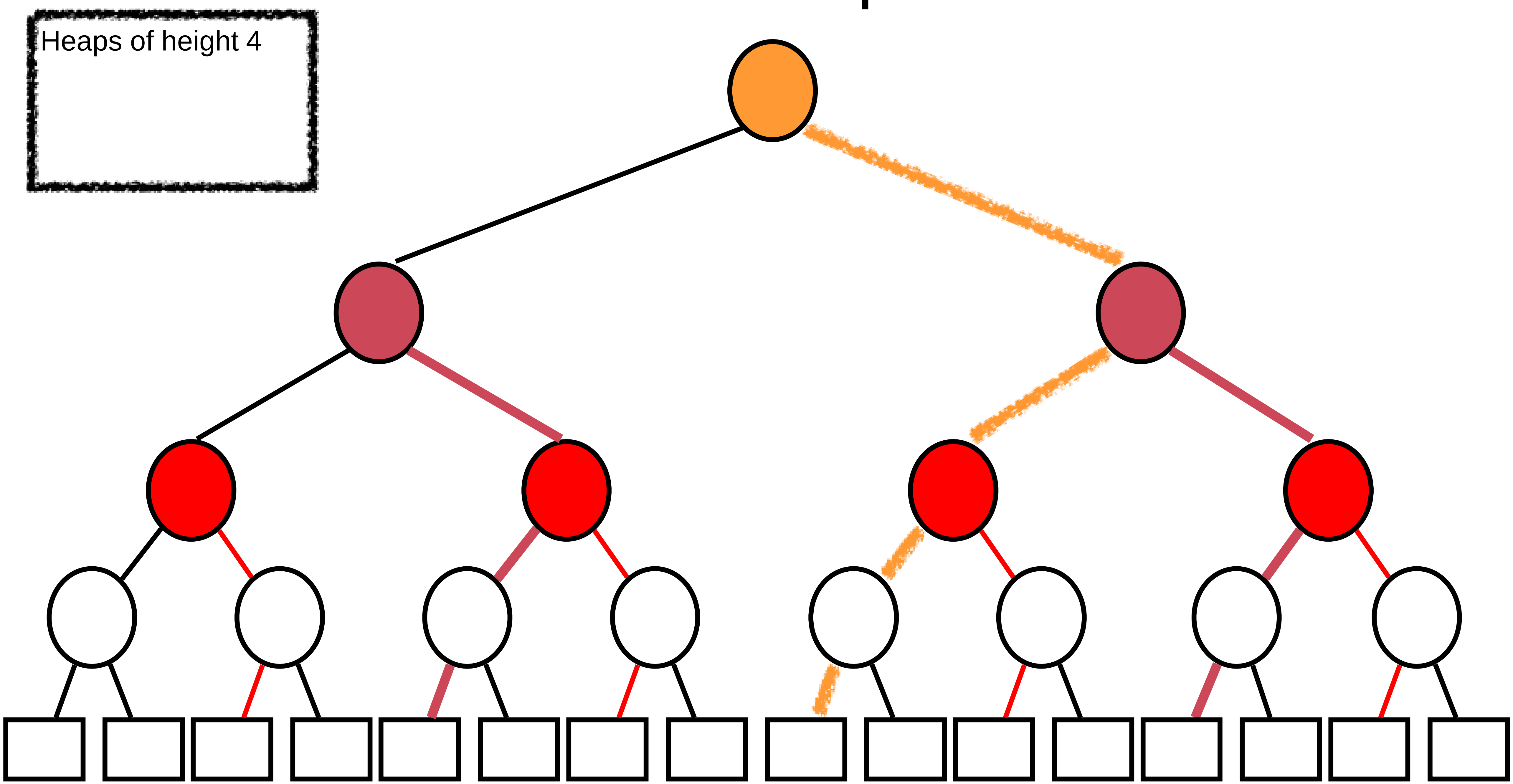
For each height-2 heap, leftmost branch not marked

For each new node joining two heaps: mark path with maximum number of bubble-down operations



For each height-3 heap, leftmost branch not marked

For each new node joining two heaps: mark path with maximum number of bubble-down operations



For height-4 heap, leftmost branch not marked

Inductive argument: marking procedure will never mark all edges in heap, since the leftmost branch is never marked

- Note: leftmost branch in height-h heap: not marked
- When joining 2 heaps of height h to heap of height $h + 1$: new edges to be marked are
  - edge joining new node and right heap of height h, and
  - edges on left path in the right heap of height h
- We conclude: leftmost branch in height $(h+1)$ heap is not marked

Build Heap In-place

Algorithm buildHeap(A,n):
    for $i \leftarrow \lfloor n/2 \rfloor$ **to** 1 **do**
        downHeap(A,i)

Algorithm downHeap(A,i):
    $l \leftarrow 2i$
    $r \leftarrow 2i + 1$
    if $l \leq n \wedge A[l] < A[i]$ **then**
        $min \leftarrow l$
    else
        $min \leftarrow i$
    if $r \leq n \wedge A[r] < A[min]$ **then**
        $min \leftarrow r$
    if $i \neq min$ **then**
        swap(i,min)
        downHeap(A,min)