

COMPUTER SCIENCE 349A

Handout Number 13

HORNER'S ALGORITHM (NESTED MULTIPLICATION, SYNTHETIC DIVISION)

Given a polynomial $f(x) = \sum_{i=0}^n a_i x^i$ and a value x_0 , this algorithm is used to efficiently evaluate $f(x_0)$ and $f'(x_0)$. To illustrate the basic idea, consider the case $n = 4$:

$$(1) \quad f(x) = a_0 + a_1 x + a_2 x^2 + a_3 x^3 + a_4 x^4$$

can be rewritten in the form

$$(2) \quad f(x) = a_0 + x \times (a_1 + x \times (a_2 + x \times (a_3 + x \times a_4))) .$$

Evaluation of (1) at x_0 requires 7 multiplications and 4 additions, whereas (2) requires only 4 multiplications and 4 additions.

The general case (for a polynomial of order n):

form (1) requires $2n - 1$ multiplications and n additions,

form (2) requires n multiplications and n additions.

An algorithm to evaluate $f(x_0)$, assuming that $f(x) = \sum_{i=0}^n a_i x^i$ is written in the “**nested**” form, as in (2):

Let $b_n = a_n$. Then compute

$$\left. \begin{array}{l} b_{n-1} = a_{n-1} + b_n x_0 \\ b_{n-2} = a_{n-2} + b_{n-1} x_0 \\ \vdots \\ b_0 = a_0 + b_1 x_0 \end{array} \right\} \quad \begin{array}{l} b_k = a_k + b_{k+1} x_0 \\ \text{for } k = n-1, n-2, \dots, 1, 0 \end{array}$$

Then $b_0 = f(x_0)$.

NOTE that execution of this algorithm requires exactly n multiplications and n additions.

Development of an algorithm to evaluate $f'(x_0)$:

Let b_n, b_{n-1}, \dots, b_0 be defined as above, and now define

$$Q(x) = b_1 + b_2x + b_3x^2 + \dots + b_nx^{n-1}.$$

Then

$$\begin{aligned} & (x - x_0)Q(x) + b_0 \\ &= (x - x_0)(b_1 + b_2x + \dots + b_nx^{n-1}) + b_0 \\ &= (b_1x + b_2x^2 + \dots + b_nx^n) - (b_1x_0 + b_2x_0x + \dots + b_nx_0x^{n-1}) + b_0 \\ &= (b_0 - b_1x_0) + (b_1 - b_2x_0)x + \dots + (b_{n-1} - b_nx_0)x^{n-1} + b_nx^n \\ &= a_0 + a_1x + \dots + a_{n-1}x^{n-1} + a_nx^n \\ & \quad \text{using the definition of the } b_k \text{'s} \\ &= f(x). \end{aligned}$$

Differentiating this with respect to x gives

$$f'(x) = Q(x) + (x - x_0)Q'(x),$$

which implies that

$$f'(x_0) = Q(x_0).$$

Thus, to evaluate $f'(x_0)$, one first needs to evaluate $f(x_0)$ as above, which gives the coefficients b_n, b_{n-1}, \dots, b_0 , and then evaluate $Q(x_0)$. The most efficient way to evaluate $Q(x_0)$, of course, is to use the nested form for the polynomial $Q(x)$. The following algorithm evaluates both $f(x_0)$ and $f'(x_0) = Q(x_0)$, using nested multiplication to evaluate both of the polynomials.

HORNER'S ALGORITHM

Given values a_0, a_1, \dots, a_n and x_0 , compute:

$$\begin{array}{ll} b_n = a_n & c_n = b_n \\ b_{n-1} = a_{n-1} + b_nx_0 & c_{n-1} = b_{n-1} + c_nx_0 \\ b_{n-2} = a_{n-2} + b_{n-1}x_0 & c_{n-2} = b_{n-2} + c_{n-1}x_0 \\ \vdots & \vdots \\ b_1 = a_1 + b_2x_0 & c_1 = b_1 + c_2x_0 \\ b_0 = a_0 + b_1x_0 & \end{array}$$

Then

$$b_0 = f(x_0) \text{ and } c_1 = f'(x_0).$$

EXAMPLE

Let $n = 4$ and

$$f(x) = x^4 - 2x^3 + 2x^2 - 3x + 4.$$

Using Horner's algorithm to evaluate $f(1)$ and $f'(1)$:

$$\begin{array}{ll} b_4 = 1 & c_4 = 1 \\ b_3 = -2 + (1)(1) = -1 & c_3 = -1 + (1)(1) = 0 \\ b_2 = 2 + (-1)(1) = 1 & c_2 = 1 + (0)(1) = 1 \\ b_1 = -3 + (1)(1) = -2 & c_1 = -2 + (1)(1) = -1 \\ b_0 = 4 + (-2)(1) = 2 & \end{array}$$

giving $f(1) = b_0 = 2$ and $f'(1) = c_1 = -1$.

Note that the explicit form of $f'(x)$, namely

$$f'(x) = 4x^3 - 6x^2 + 4x - 3,$$

is not obtained; only the value of $f'(1)$ is computed. Since $Q(x)$ depends on the value of x_0 , which is equal to 1 above, all computations must be re-done in order to evaluate $f'(x)$ at a different value of x .