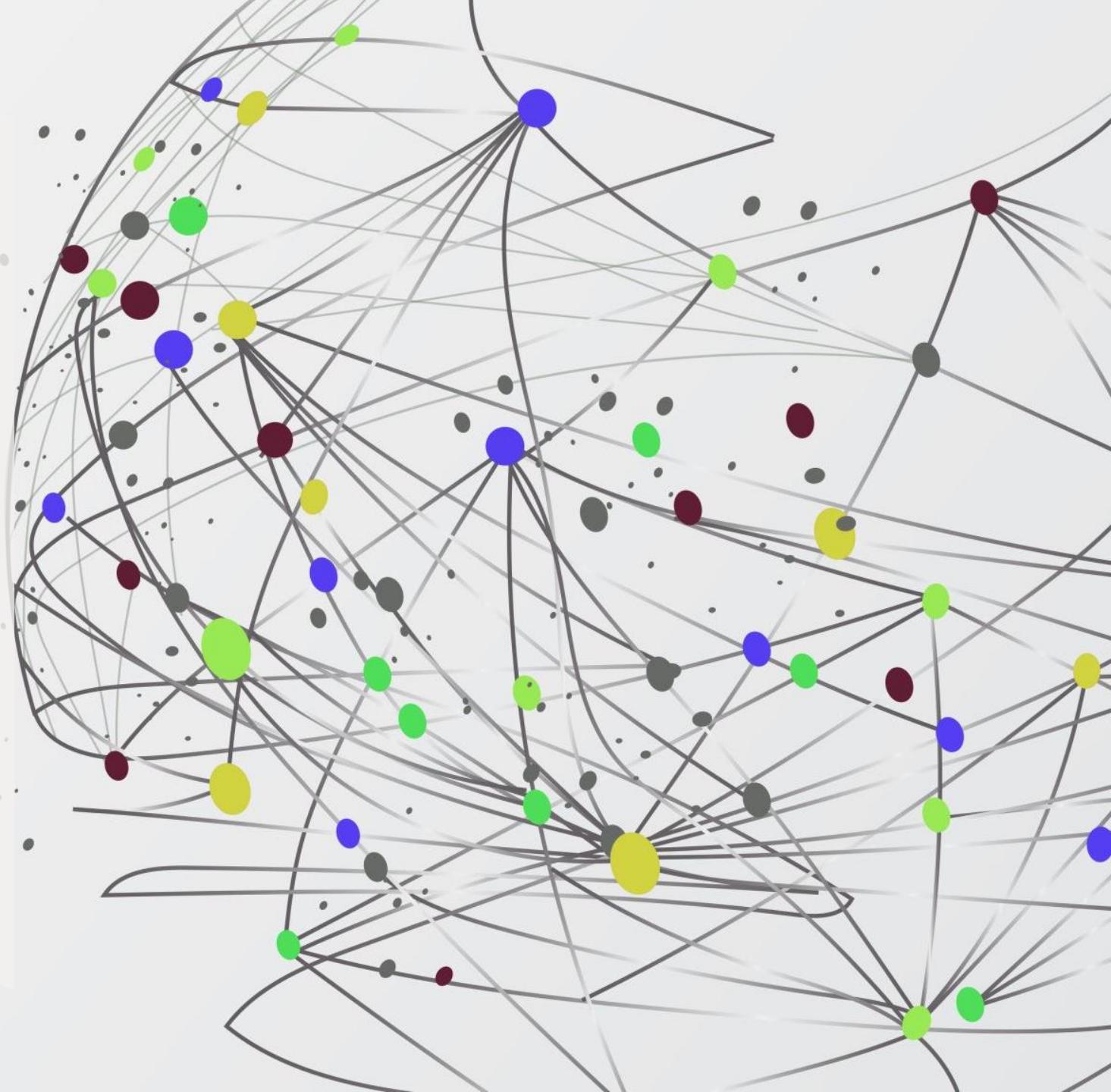


# GRAPH COLOURING PROBLEM



OPTIMIZATION METHODS - II

-Dikshant Joshi

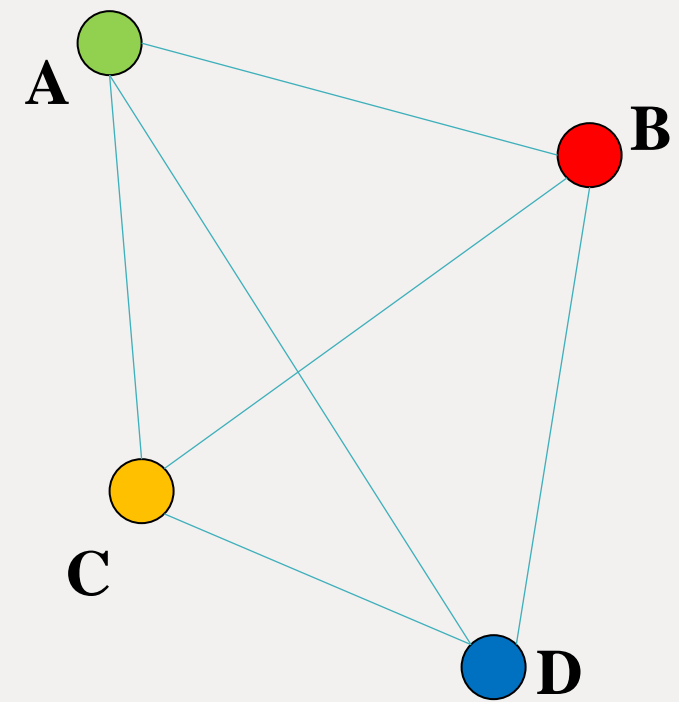


# Agenda



- Overview
- Application of Graph Coloring Problem
- Problem Statement
- Algebraic Formulation
- Modelling using Python
- Modelling using Excel
- Results

# OVERVIEW



- The graph coloring problem is a well-known problem in computer science and graph theory.
- It is the problem of assigning colors to the vertices of a graph such that no two adjacent vertices have the same color, while using the fewest number of colors possible

# APPLICATION OF GCP



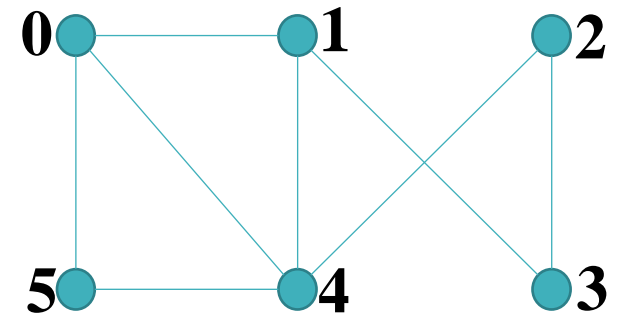
- Timetabling
- Register allocation
- Airline Scheduling

**Note:** Graph coloring problem is NP Hard .This means that there is no known algorithm that can solve it in polynomial time for all instances



# ? PROBLEM STATEMENT

- There are 6 nodes in a graph given and we need to colour these node in a way that adjacent nodes do not have same colour.
- The adjacency matrix is shown in the below graph which tells if two nodes are connected (1) or not(0).
- Based on the given problem we need to find optimum number of colors required to color the nodes of the graph



	0	1	2	3	4	5
0	0	1	0	0	1	1
1	1	0	0	1	1	0
2	0	0	0	1	1	0
3	0	1	1	0	0	0
4	1	1	1	0	0	1
5	1	0	0	0	1	0



# ALGEBRAIC FORMULATION

## Decision Variables

- »  $Y_j$  : a binary variable that indicates whether color  $j$  has been used or not (1 if used , 0 otherwise)
- »  $X_{ij}$  : a binary variable that indicates whether node  $i$  is colored with color  $j$  or not

$Y_j$	0	0	0	0	0	0
-------	---	---	---	---	---	---

## Objective Function:

- » Minimize  $Z: \sum_i Y_i$   $i \in N$ ,

Where  $N$  is the set of all nodes.

		0	1	2	3	4	5
0	0	0	0	0	0	0	0
1	0	0	0	0	0	0	0
2	0	0	0	0	0	0	0
3	0	0	0	0	0	0	0
4	0	0	0	0	0	0	0
5	0	0	0	0	0	0	0



# ALGEBRAIC FORMULATION

- **Constraints**

- » Each node should have only one color:

$$\sum_j X_{ij} = 1 \text{ for all } i$$

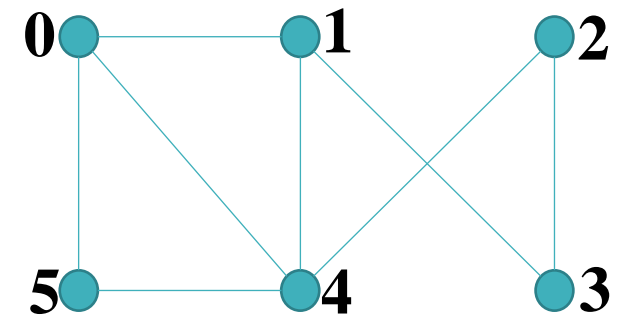


# ALGEBRAIC FORMULATION

- **Constraints**

- » Each node should have only one color:

$$\sum_j X_{ij} = 1 \text{ for all } i$$





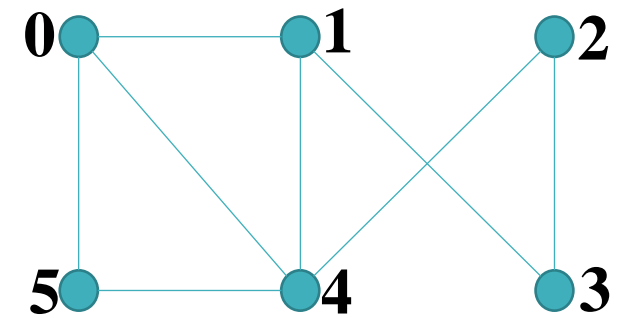


# ALGEBRAIC FORMULATION

- **Constraints**

- » Each node should have only one color:

$$\sum_j X_{ij} = 1 \text{ for all } i$$



	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	0	0	0	0	0
2	0	0	0	0	0	0
3	0	0	0	0	0	0
4	0	0	0	0	0	0
5	0	0	0	0	0	0

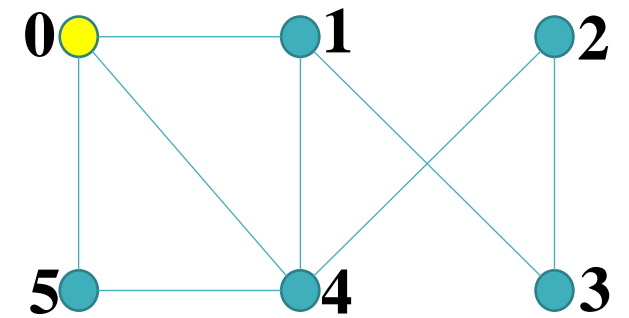


# ALGEBRAIC FORMULATION

## ■ Constraints

» Each node should have only one color:

$$\sum_j X_{ij} = 1 \text{ for all } i$$



	0	1	2	3	4	5
0	1	0	0	0	0	0
1	0	0	0	0	0	0
2	0	0	0	0	0	0
3	0	0	0	0	0	0
4	0	0	0	0	0	0
5	0	0	0	0	0	0



# ALGEBRAIC FORMULATION

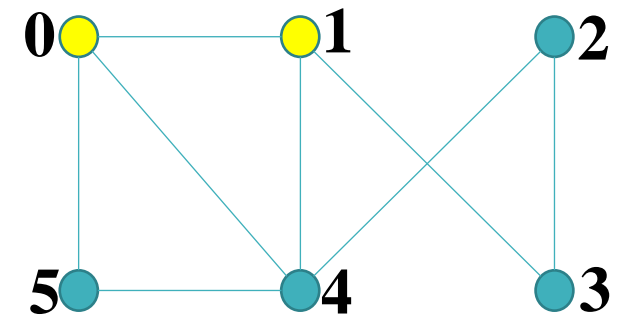
## ■ Constraints

- » Each node should have only one color:

$$\sum_j X_{ij} = 1 \text{ for all } i$$

- » Adjacent nodes must have different color:

$$X_{ik} + X_{jk} \leq 1 \text{ for all } A_{ij} = 1, k \in N \text{ where } A \text{ is adjacency matrix}$$



	0	1	2	3	4	5
0	1	0	0	0	0	0
1	1	0	0	0	0	0
2	0	0	0	0	0	0
3	0	0	0	0	0	0
4	0	0	0	0	0	0
5	0	0	0	0	0	0



# ALGEBRAIC FORMULATION

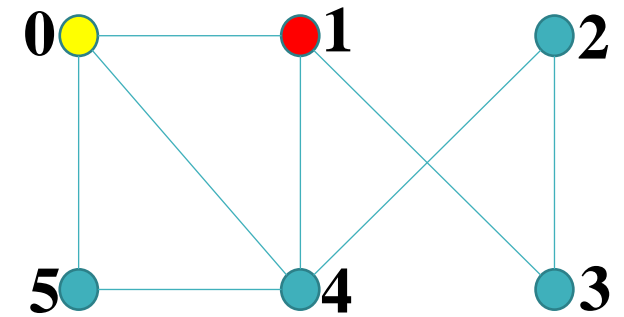
## ■ Constraints

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	0	1	2	3	4	5
0	1	0	0	0	0	0
1	0	1	0	0	0	0
2	0	0	0	0	0	0
3	0	0	0	0	0	0
4	0	0	0	0	0	0
5	0	0	0	0	0	0



# ALGEBRAIC FORMULATION

## ■ Constraints

- » Each node should have only one color:

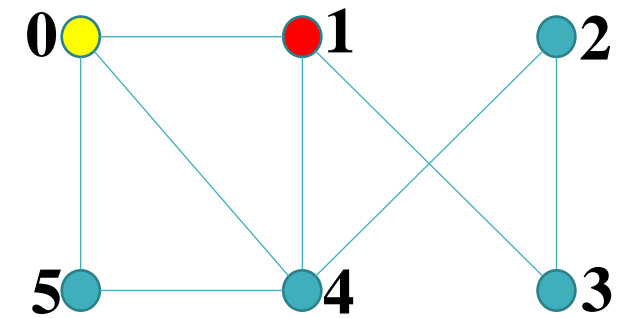
$$\sum_j X_{ij} = 1 \text{ for all } i$$

- » Adjacent nodes must have different color:

$$X_{ik} + X_{jk} \leq 1 \text{ for all } A_{ij} = 1, k \in N \text{ where } A \text{ is adjacency matrix}$$

- » The colour should be used only n number of times:

$$\sum_i X_{ij} \leq n \text{ for all } j$$



	0	1	2	3	4	5
0	1	0	0	0	0	0
1	0	1	0	0	0	0
2	0	0	0	0	0	0
3	0	0	0	0	0	0
4	0	0	0	0	0	0
5	0	0	0	0	0	0



# ALGEBRAIC FORMULATION

## ■ Constraints

- » Each node should have only one color:

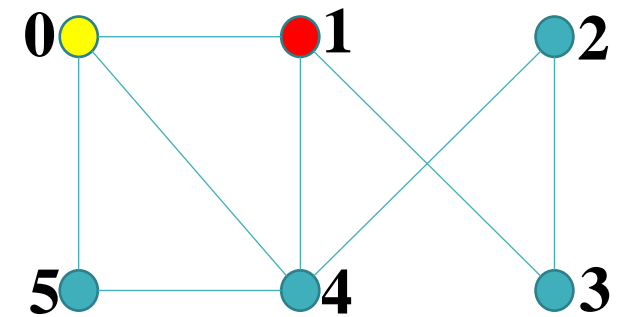
$$\sum_j X_{ij} = 1 \text{ for all } i$$

- » Adjacent nodes must have different color:

$$X_{ik} + X_{jk} \leq 1 \text{ for all } A_{ij} = 1, k \in N \text{ where } A \text{ is adjacency matrix}$$

- » The colour should be used only n number of times:

$$\sum_i X_{ij} \leq n * Y_j \text{ for all } j$$



	0	1	2	3	4	5
0	1	0	0	0	0	0
1	0	1	0	0	0	0
2	0	0	0	0	0	0
3	0	0	0	0	0	0
4	0	0	0	0	0	0
5	0	0	0	0	0	0





# ALGEBRAIC FORMULATION

## ■ Constraints

- » Each node should have only one color:

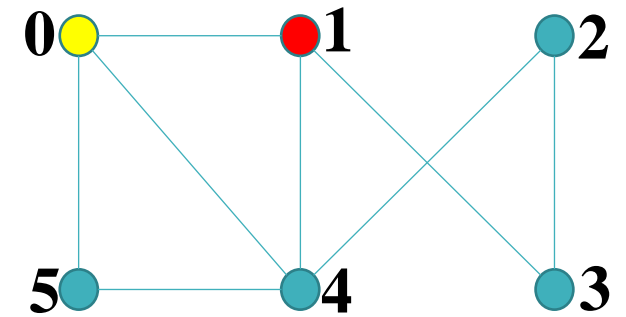
$$\sum_j X_{ij} = 1 \text{ for all } i$$

- » Adjacent nodes must have different color:

$$X_{ik} + X_{jk} \leq 1 \text{ for all } A_{ij} = 1, k \in N \text{ where } A \text{ is adjacency matrix}$$

- » The colour should be used only n number of times:

$$\sum_i X_{ij} \leq n \text{ for all } j$$



	0	1	2	3	4	5
0	1	0	0	0	0	0
1	0	1	0	0	0	0
2	0	0	0	0	0	0
3	0	0	0	0	0	0
4	0	0	0	0	0	0
5	0	0	0	0	0	0





# ALGEBRAIC FORMULATION

## ■ Constraints

- » Each node should have only one color:

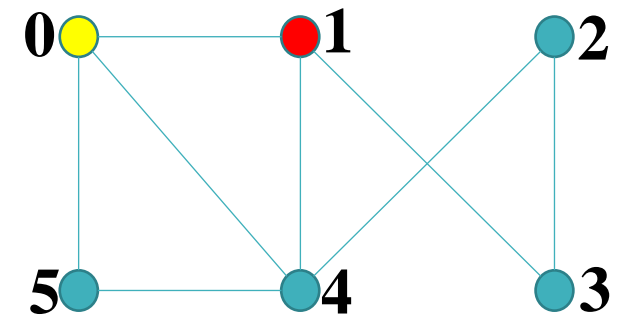
$$\sum_j X_{ij} = 1 \text{ for all } i$$

- » Adjacent nodes must have different color:

$$X_{ik} + X_{jk} \leq 1 \text{ for all } A_{ij} = 1, k \in N \text{ where } A \text{ is adjacency matrix}$$

- » The colour should be used only n number of times:

$$\sum_i X_{ij} \leq n * Y_j \text{ for all } j$$



	0	1	2	3	4	5
0	1	0	0	0	0	0
1	0	1	0	0	0	0
2	0	0	0	0	0	0
3	0	0	0	0	0	0
4	0	0	0	0	0	0
5	0	0	0	0	0	0

1	1	0	0	0	0
---	---	---	---	---	---





# MODELLING USING PYTHON

```
# Define the adjacency matrix for the graph with 5 nodes
```

```
adj_matrix = np.array([[0, 1, 0, 0, 1, 1],  
                        [1, 0, 0, 1, 1, 0],  
                        [0, 0, 0, 1, 1, 0],  
                        [0, 1, 1, 0, 0, 0],  
                        [1, 1, 1, 0, 0, 1],  
                        [1, 0, 0, 0, 1, 0]])
```

```
# Define the number of nodes
```

```
num_nodes = 6
```

```
# Define the decision variable
```

```
# Boolean variables y[j] that colour j is used or not
```

```
# Assumption: we have n number of colours available where n = nodes
```

```
y = cp.Variable(num_nodes, boolean=True)
```

```
# Boolean variable array X[i,j] that node i is coloured with colour j or not
```

```
x = cp.Variable((6,num_nodes), integer=True)
```

```
# Define the objective function
```

```
obj = cp.Minimize(cp.sum(y))
```

```
# Define the constraints
```

```
# This allows to have only one color per node
```

```
constraints = []
```

```
for i in range(num_nodes):
```

```
    for j in range(num_nodes):
```

```
        constraints.append(x[i, j] >= 0)
```

```
    constraints.append(cp.sum(x[i, :]) == 1)
```

```
# This allows the adjacent node to have different colors
```

```
for i in range(num_nodes):
```

```
    for j in range(num_nodes):
```

```
        if adj_matrix[i, j] == 1:
```

```
            for k in range(num_nodes):
```

```
                constraints.append(x[i, k] + x[j, k] <= 1)
```

```
# This constraint is indirectly ensuring that if a colour is used to colour a node than it should reflect in the array y.
```

```
constraints.append(cp.sum(x, axis=0) <= num_nodes * y)
```

```
# Solve the problem
```

```
problem = cp.Problem(obj, constraints)
```

```
problem.solve(solver=cp.GLPK_MI)
```

```
# Print the solution
```

```
print("Minimum number of colors:", int(obj.value))
```

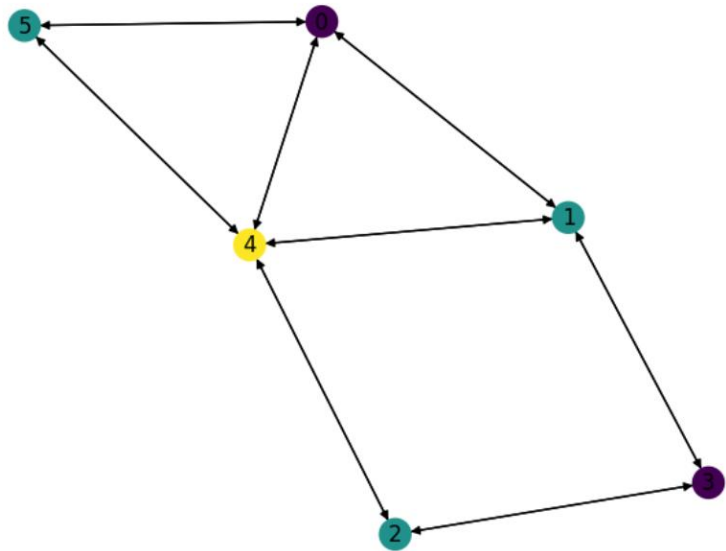
```
print("Node colors:")
```

```
print(np.argmax(x.value, axis=1))
```



## RESULT

Minimum number of colors: 3  
Node colors:  
[2 3 3 2 4 3]



The Optimum Number of colours for this Graph  
colouring problem is : 3  
With colours as:

0	1	2	3	4	5
2	3	3	2	4	3



Set Objective:

To: ☐ Max ☒ Min ☐ Value Of:

By Changing Variable Cells:

Subject to the Constraints:

☒ Make Unconstrained Variables Non-Negative

Select a Solving Method:

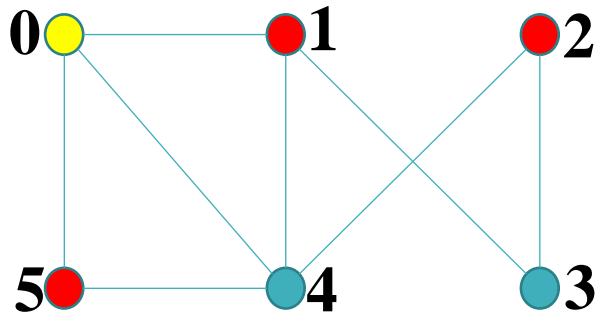


## RESULT

Nodes	0	1	2	3	4	5
Color	2	0	0	1	1	0
No. of colors used	3					

The Optimum Number of colours for this Graph colouring problem is : 3  
With colours as:

0	1	2	3	4	5
2	0	0	1	1	0





THANK YOU



QUESTIONS ?