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P2PCDN

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SUMMARY
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1. Introduction

In recent years due to the widespread of broadband Internet connections, user achieved the possibility of downloading and sharing their multimedia contents. Futhermore, more advanced hardware with a reduction of computer costs allowed users to own a high performance PC to perform complex task. More people are watching video on their PC using specific applications or plugins that integrated into browsers. Common technical solution just rely on a client-server architecture. This solution is not viable since it can not scale if the content becomes popular. Some content companies have agreement with CDN (content delivery network) providers such as Akamai, Limelight, or Level3 to distributed the contents better scale and better QoE (quality of experience) to end users. For example, Youtube uses Google cache, MTV uses Akamai, etc. P2P (peer to peer) applications are also widely deployed and their success is quite important. In China, P2P is very popular. We see many P2P applications from China such as PPLive, PPStream, UUSe, Xunlei, etc. Some news broadcaster are also rely on P2P technology to deliver popular live events. For example, CNN uses Octoshape solution that enable to scale and offer good video quality while the number of users increase. From Internet provider point of view, the diffusion of powerful computers in users side suggested that it could be possible to delegate a portion of computing or storage tasks to the users thus creating P2P networks where users can share files and multimedia contents. Starting from file sharing only protocol, P2P architecture evolved toward video on demand and live events support.

A P2P based architecture usually requires a sufficient number of nodes supplying the data (termed as seeders) to start the distribution process among the joining peers. A peer usually offers low outbound streaming rate due to traditional asymmectrical DSL home connections) and hence

multiple peers must jointly stream a media data to a requesting peer. The decentralized, uncoordinated operation implies that scaling to high number of peers comes with side effects. Typical problem of a P2P streaming architecture are the low stream quality with undesirable disruptions, resources unfairness due to heteregeneous peer resources and high startup delay. Moreover, current P2P applications are not aware of the underlying network and may conflict with the ISP routing policies and business model.

Nowadays, a number of P2P streaming applications have been designed, analyzed and deployed attracting a significant amount of users. Research studies and deployement experiences have both demostated that P2P is a promising solution in terms of scalability and deployment costs. On the other hand, the heteregeneous nature and unstable behavior of the peers contributing bandwidth and computational resources along with the networking issues are limit the commercial success of P2P video streaming applications. Apart from P2P protocols, video contents can be efficiently distributed leveraging on services offered by managed networks architectures and CDN companies. The major issue of CDN are high deployment cost and limited scalability in the long term. Given the complementary features of P2P and CDN, in the last years some hybrid solutions have been proposed [1–3] to take the best of both approaches.

Nowadays, xDSL networks are deployed worldwide and broadband acess network helps the P2P application to perform better. In the coming years, FTTH will be massively deployed by network operators in order to offer to their customers and even larger acces bandwidth. As bandwidth increase, more the P2P systems is efficient since a peer can contribute much more.

P2P architecture offers great potential for content providers to cost effectively distribute content by capitalizing network resources of end users. It is economical superiority to the traditional client-server architecture and it has been demonstrated by lots of academic work and successful commercial systems [3]. We believe more content providers will adopt P2P technology in the future.

In the mentioned work, each end user device is involved directly in the P2P swarm. In such cases, the installation of a software for the P2P engine in every end device is necessary and the user itself is directly involved in the content's swarm. In such condition peers may result unstable and the swarm can be affected by rapid and frequent disconnections. Furthermore, user's devices usually can contribute to the swarm with limited upload bandwidth and techniques
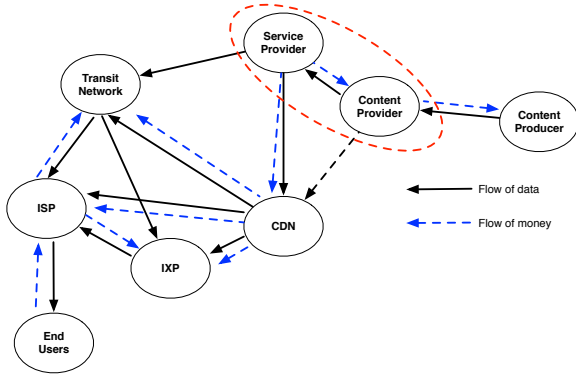
Fig. 1: Complex relationship of entities in Internet.The data flow from content producer to content provider. From content provider data can flow to service provider then goes to CDN. Depends on routing and peering policy, data from CDN can flow directly to ISP or flow to ISP via IXP (Internet Exchange) or via transit network. Flow of data noted by straight arrow and flow of money noted by dash arrow line.

developed in P2P to select neighboring peers are often unfriendly with ISP's routing policies. Different topologies have been proposed in literature, where the collaborative mechanism for content distribution is created among more stable devices such as residential gateways which are still geographically close to the end users but directly owned and managed by ISPs. In such conditions, running on more stable and powerful device, each peer can contribute with more bandwidth to the content swarm with respect to traditional end-user P2P. Peer selection procedures can also managed directly by ISP, avoiding the traversal of multiple nodes across ISP boundaries.

In some recent works, the residential gateway (i.e a home gateway placed in home premise, serving several terminals within the home network but directly managed by ISP) is considered the central entity of such managed P2P infrastructure as describe in [4, 5]. Since P2P traffic now decreasing and moving to cloud [6], there is room from ISP to use user's home gateway as peer assisted CDN since users usually do more download rather than upload. The user home gateway is always on 24 hours therefore this is perfect device to run peer assisted application. ISP can give discount price to user's who wants their home gateway to be used as peer assisted by ISP and ISP can get benefit from peer assisted running in their customer home gateway to help their CDN. With the growing interest to do interconnect of CDN [7,8] this architecture can give good benefit for ISP.

In this paper we study, the economic incentive for ISP who adopt peer assisted CDN and shows that the advantages to adopt peer assisted CDN.

2. CDN Complexity

As shown in Fig.1, how content in Internet delivered to end-users. Content producer such as movies companies send their movies to content provider for example: Netflix and

Hulu. Content providers then deliver the movies using CDN or they also can use their upstream provider. Depends on peering policies, CDN can reach to customer directly via ISP or via IXP then reach ISP or goes to transit networks then reach ISP. In current modern Internet topology, content provider and service provider can be merged in on entity, for example: Google. Labovitz et al., [6] mentioned that the hypergiant entities such as Google doing massive peering to IXP in order to be closed to ISP. Although CDN is placed close to ISP network, it does not guarantee end-users can get good quality video stream [9]. Other than technical complexity as mentioned before, CDN also faces economic complexity. Therefore, the future of CDN business is likely to live deeper into ISP networks, more integrated into and interleaved with ISP infrastructures.

3. System Description

Fig.2 shows the model of system archictecture. That's model offers the opportunity to connect the P2P swarm not directly but through an intermediate later of proxies. This function has potential benefit: boosting download performances (by exposing through proxy devices higher upload bandwidth than end users) and increasing privacy (by reducing swarm exposure for end users). In Fig.2 two level CDN is defined. The first level L1 consists high capacity servers strategically placed to increase network backbone capacity. L1 is more like current CDN network. In L1 level, CDN node can communicate each other by using common interconnect interface eventhough the CDN node are owned by different companies or different ISPs [7]. L0 consists CDN proxies are owned or to be managed by ISP. These nodes are placed closer to end users. CDN proxies can communicate each other to react to the content requests from users. L0 are the natural service points for the users and they can offer the best quality of service if the are able to efficiently cache the requested content. The state of the L0 and L1 caches is stored in DHT. The DHT is used to keep an updated record for each content in scalable manner. This different caching levels are used to offer an adaptive, flexible, and scalable service to the users.

Other architecture, CDN federation proposed by Cisco and some ISPs to interconnect the ISP's CDN [10]. The goal this architecture is to interconnect ISPs owned CDN thus reduce Capex and Opex for content delivery. This architecture does not involved P2P. Some peering or transit might be needed by ISP to join this architecture.

4. System Model

We present a cooperative model between P2P and CDN using simple a single bitrate for video streaming as shown in Fig.3. The model is a stochastic fluid model similar to [11]. While the authors in [11] focus on probabiity of degraded service on small and large systems, our work focus on lower bound of peers that can support intended bit rate. The knowledge of lower bound is very import for ISP
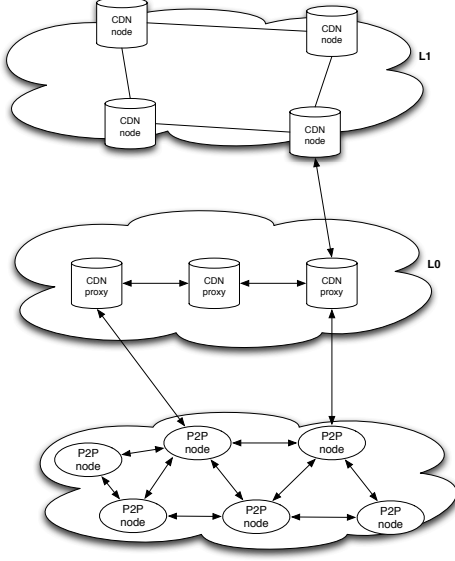
Fig. 2: Two tier CDN-P2P topology. The CDN proxy maintained by ISP. ISP also has CDN node. CDN node can exchange contents with other CDN node owned by other CDN companies.P2P nodes are ISP's customers that run P2P software.
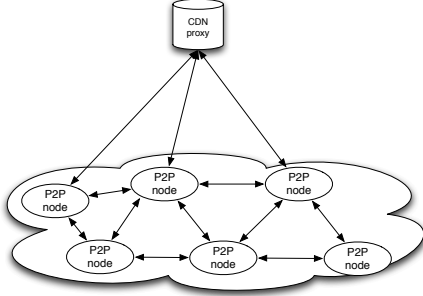

Fig. 3: System architecture on L0. P2P nodes that run on user homegateway communicates to CDN proxy.

to estimate number of customers that can join to P2P. We aim to find out how many peers will be sufficient to provide a stream to a certain number of leechers. Throughout the rest of this section we assume that CDN proxy has bit rate r and the upload capacity of the CDN proxy is C_{proxy}.

We analyze the system for the theoretical unconstrained case when peers can have unlimited number of connections with other peers. We also present the analysis more realistic case for constrained case when peers can only have a limited number of incoming and outgoing connections. Moreover, for both case above we also present the analysis for churnless and churn systems.

For churn case, we assume peers join and leave at random times, stay in for a random period, then leave the system. Assuming poisson process with rate λ. Peers stay in the system for a period of time that follows a general probability distribution with mean $\frac{1}{\gamma}$. Define $N(t)$ as the number of peers in the system at time t, $N(t)$ can be represented as the number of peers in a $M/G/\infty$ queue.

4.1 Unconst rained System

In unconstrained churnless system, denote n_l is the number of leechers and n_s is the number of seeders in the system. Also denote u_i^l , u_j^s as upload rate of leecher i and seeder j. Following two condition must be holded: $C_{proxy} \geq n_s r$ and $\sum_{j=1}^{n_s} u_j^s \geq r$. First inequality express server capacity constraint and second inequality express the set of seeders have minimum upload capacity. The maximum achievable streaming rate r_{max} is given by

$$r_{max} = min\{U_s, \frac{U_s + \sum_{i=1}^{n_l} u_i^l}{n_l}\} \tag{1}$$

U_s is total update rate of all seeders and $U_s = \sum_{j=1}^{n_s} u_j^s$ for simplicity we assume all leechers have the same upload rate of u_l and all seeders have the same upload rate of u_s thus u_l and u_s can be considered as the average upload rates over all leechers and seeders. we assume $r \geq u_l$. it means that the average upload rate of single leecher is not enough to support sharing the whole stream with other peers. in that case r_{max} can be reduced to $r_{max} = min\{n_s u_s, \frac{n_s u_s + n_l u_l}{n_l}\}$. when the system in the steady state condition we can assume that $n_s u_s \geq \frac{n_s u_s + u_l u_l}{n_l}$ then we can get that a churnless system can support a streaming rate of:

$$r \leq \frac{n_s u_s + n_l u_l}{n_l} \tag{2}$$

we can derive for theoretical lower bound $n_s \geq \frac{n_l(r - u_l)}{u_s}$ on the number of seeders to support bitrate r.

In unconstrained with churn system, denote $N(t)$ is the total number of peers in the system at time t (including seeders and leechers) and N follows a Poisson distribution with rate $\rho = \frac{\lambda}{\gamma}$. for computing the probability that the system will be able to support a streaming rate r in case of node churn, we assume that node churn only happens in leecher nodes. This means number of seeders in the system are assumed to be constant. Only leechers arrive and leave the system. The probability that we can support a bitrate r as following:

$$P(r) = P(r \leq \frac{n_s u_s + N u_l}{N})$$

$$= P(N \leq \frac{n_s u_s}{r - u_l})$$

$$= F(\frac{n_s u_s}{r - u_l})$$

$F(w) = \sum_{x=0}^{w} \frac{e^\rho \rho^x}{x!}$ for large system we can assume $\rho \rightarrow \infty$ thus we can approximate the Poisson distribution with a Gaussian distribution with mean $\mu = \rho$ and standard deviation $\sigma = \sqrt{\rho}$ as $\frac{1}{\sigma}\phi(\frac{x-\mu}{\sigma})$. Hence, we can get:

$$P(r) = P(\frac{N - \rho}{\sqrt{\rho}} \leq \frac{\frac{n_s u_s}{r - u_l} - \rho}{\sqrt{\rho}})$$

$$= \Phi(\frac{\frac{n_s u_s}{r - u_l} - \rho}{\rho})$$

$\Phi(z)$ is the cumulative distribution function of the standard normal random variable. Let $\phi_{1-\alpha} = 1 - \alpha$ where $\phi_{1-\alpha}$ is a positive real number. The confidence $(1 - \alpha)x100\%$ that peers can support bitrate r is $(\frac{n_s u_s}{r-u_l} - \rho)/\sqrt{\rho} \geq \phi_{1-\alpha}$. Finally we can derive lower bound on the number of peers:

$$n_s \geq \frac{(\phi_{1-\alpha}\sqrt{\rho} + \rho)(r - u_l)}{u_s} \qquad (3)$$

4.2 Constrained System

In constrained system, each peer as a limited number of inbound and outbound connections. We defined the limit as following:

- S_{in} is the maximum number of incoming connections a seeder can accept.
- Y_{in} is the maximum number of incoming connections a leecher can accept.
- Y_{out} is the number connections a leecher can initiate. We assume that there is no limit on the number of connections a leecher can initiate. It means the bottleneck is in the number of connections that actually established. This number is mainly controlled by parameter S_{in} and Y_{in}

We define η as the efficiency of the P2P protocol which can be computed as the probability of any leecher finding new content at other leechers when they establish a connection [12]. It is shown that Bittorrent efficiency can exceed 0.9 if the file has more than ten pieces [13]. P2P protocol efficiency means that a leecher has an effective upload rate of ηu_l. Denote d as the average download rate for any leecher in the swarm, then d can be computed as:

$$d = \sum_x E[d]|\text{leecheer connected to x seeders} \times P_r\{x\}$$
$$= \sum_x (\frac{xu_s}{S_{in}} + \frac{(Y_{out} - x)\eta u_l}{Y_{in}})P_r\{x\}$$
$$= \frac{Y_{out}\eta u_l}{Y_{in}} + (\frac{u_s}{S_{in}} - \frac{\eta u_l}{Y_{in}})\sum_x xPr\{x\}$$

Note that $\sum_x sPrx$ is the average number of seeders connected to each leecher which can be approximated by the value $n_s S_{in}/n_l$. Note also that when $n_l \geq n_s$, we can calculate an approximate value for the number of connections each leecher can establis by $Y_{out} = (n_s S_{in} + n_l Y_{in})/n_l$. Subtituting these two equations in equation above , then d can be reduced to:

$$d = \frac{n_s u_s + \eta n_l u_l}{n_s} \qquad (4)$$

Since d can be considered the average bitrate that can be supported by the system, then the number of seeders sufficient to support the bitrate is $n_s = n_l(r - \eta u_l)/u_s$.

In constrained system with churn, we can compute the probability of supporting a range of bitrates r in a system as follows:

$$P(r) = P(\frac{n_s u_s + \eta N_u l}{N} - \delta \leq r \leq \frac{n_s u_s + \eta N u_l}{N} + \delta)$$
$$P(r) = P(\frac{n_s u_s}{r - \eta u_l + \delta} \leq N \leq \frac{n_s u_s}{r - \eta u_l - \delta})$$

As we did earlier, N can be approximate by a Gaussian random variable. With $(1 - \alpha)100\%$ confidence interval for N we can set:

$$\frac{\frac{n_s u_s}{r - \eta u_l + \delta} - \rho}{\sqrt{\rho}} \leq -\phi_{1-\alpha/2}, \frac{\frac{n_s u_s}{r - \eta u_l - \delta} - \rho}{\sqrt{\rho}} \geq \phi_{1-\alpha/2}$$

We define $\hat{\phi} = \phi_{1-\alpha/2}$ then we can get following confidence interval for N

$$\frac{(\rho + \hat{\phi}\sqrt{\rho})(r - \eta u_l - \delta)}{u_s} \leq n_s \leq \frac{(\rho - \hat{\phi}\sqrt{\rho})(r - \eta u_l + \delta)}{u_s}$$

This inequality provides valid intervals only for δ values that satisfy the condition $\delta \geq \frac{\phi_{1-\alpha/2}(r-\eta u_l)}{\sqrt{\rho}}$. We can see that δ is inversely proportional to ρ. It means that the higher peer arrival rates and the longer peers stay in the system, the lower δ becomes. Lower δ mean a smaller interval in inequality which yiedls a higher guarantee that the number of peers given by above interval will be sufficient for supporting the bitrate r.

Figure 4 shows numerical results of lower bound for seeders agaitns ρ for constrained system with churn and Fig.5 shows numerical results of lower bound for seeders agaitns ρ for unconstrained system with churn. Both numerical results have 95% confidence interval on the number of seeders sufficient to support bitrate r. We use $1/\gamma = 0.5$ hour thus we suppose that peers sojourn in the system on average for 30 minutes and we fix the arrival rate of the peers to $\lambda = 100$ per hour. We increase arrival rate λ until 160 per hour. We can see that the difference in the number of seeder sufficient to support different bitrates is not large. We also see that that lower bound number for seeders is very easy to fullfil by ISP to employ system because ISP's customer number much higher that that lower bound.

4.3 ISP Strategies: Game Theory Approach

In this section, we present simple game theory approach for analysis of ISP strategies. We define the pure strategy space S_{isp} to include two strategies combination choice $s \in \{No, P2P\}$ and corresponding subscription free P^s. Futhermore, we denote the strategies as:

- $s_1 = (No, P^{s_1})$: the ISP decides to do nothing (use CDN only) and thus keeps charging the initial subscription free P^{s_1}.
- $s_2 = (P2P, P^{s_2})$: the ISP decides recommends P2P and offers a subscription fee $P^{s_2} \leq P^{s_1}$.

Where the ISP strategy space $S_{isp} = \{s1, s2\}$. After defined the strategy space of ISP, now we will describe user reaction. One obvious user action is to accept ISP strategy. When user receives offer from ISP whether user want to join P2P
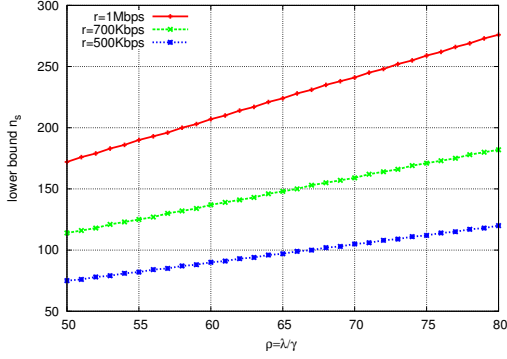
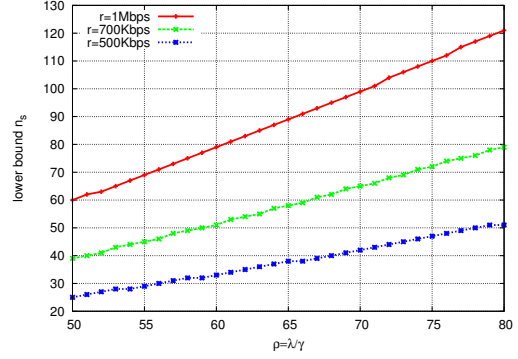Fig. 4: Constrained system with churn, 95% confidence

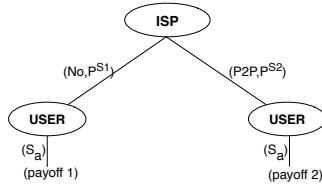Fig. 5: Unconstrained system with churn, 95% confidence

Fig. 6: Extensive form of game between ISP and users

strategy or not to join P2P strategy, user can only accept one of that strategy. If user decides to join P2P strategy from ISP then it means user accept s_2 of ISP. If user decides not to join P2P strategy from ISP then it means user accept s_1 of ISP. Therefore whatever strategy offered by ISP, the user can only accept one of the ISP strategy. We describe that situation as extensive form in Fig.6.

4.3.1 ISP Payoff

Next, we describe payoff function for ISP in our model. We assume a profit maximaxing ISP that gains higher utility proportionally with increasing profits. A simplified business model that determining ISP payoff are mainly: user generated revenue, traffif cost from inter ISP traffic (transit and peering), and ISP preferences.

For revenue model, we assume that the ISP collects revenue solely by charging an initial flat rate subscription fee P^{s_1} to its N homogeneous users who are purchasing internet access with equal and fixed quality of service. ISPs are often price discriminatory towards its customers and thus operates with different price levels for different speed or QoS. ISPs are also get revenue from other business area such as email hosting, web hosting, etc. In our model, we do not include such that revenue. In order to keep simple, we found it reasonable only to focus on users that buy the same internet access product. Given above simplifications, the ISP collects a total revenue (when deciding on a strategy s) of $R = NP^{(s)}$.

For cost model, we identified that traffic costs including (transit traffic and paid peering) are the most relevant aspect. We define the transit cost $\tau^{(s)}$ for a given ISP strategy s to be denoted as

$$\tau_t^{(s)} = \alpha^{(s)}\beta T^{(s_1)}p_t \tag{5}$$

where:

- $T^{(s_1)}$ represent the average inter-ISP traffic volume in Mbps.
- $\alpha^{(s)}$ represent the expected reduction factor of inter-ISP traffic for strategy s such that: $\alpha^{(s)} \in [0, 1]$ for $s \in \{s2\}$; $\alpha^{(s_1)} = 1$.
- β represent the ratio of inter-ISP traffic volume $T^{(s)}$ that traverse transit links such that: $\beta \in [0, 1]$.
- p_t represent the price of transit traffic.

We define the paid peering traffic cost $\tau^{(s)}$ in a similar way to transit traffic. Paid peering costs are subject to a peering traffic price p_p. Futhermore, as β is the ratio of the inter-ISP traffic $T^{(s)}$ that traverses transit links, $(1-\beta)$ express the ratio that traverses paid peering links. We can denote the paid peering traffic cost as:

$$\tau_p^{(s)} = \alpha^{(s)}(1-\beta)T^{(s_1)}p_p \tag{6}$$

Given transit cost and paid peering cost in Eq.5 and Eq.6, we can define total inter-ISP traffic cost $\tau^{(s)}$ as:

$$\tau^{(s)} = \tau_t^{(s)} + \tau_p^{(s)} \tag{7}$$

$$\tau^{(s)} = \alpha^{(s)}\beta T^{(s_1)}p_t + \alpha^{(s)}(1-\beta)T^{(s_1)}p_p \tag{8}$$

In addition to the traffic costs, we also assume that ISP needs investment cost for peer-assisted CDN. More specifically refer to Fig.3 ISP required to setup their L0 or CDN proxy. We define the investment cost $C_i^{(s)}$ to be a monthly amortization of a larger implementation related cost for the case when the ISP decides to implement P2P. We denote:

$$C_i^{(s)} > 0; \text{ for } s = s_2 \tag{9}$$

$$= 0; \text{ for } s \in \{s_1\} \tag{10}$$

5. Result and Discussion

6. Related Work

7. Conclusion and Future Work

This section conclude the papers.

Acknowledgements

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