

# Peer-Assisted Content Distribution Aided by Video Popularity Evolution Model

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**Abstract**—Generally content distribution network (CDN) have adopted two different architectures: client-server model that's is the most common architecture model and peer-assisted CDN. In client-server model, clients download content dedicated and geographically managed servers while in peer-assisted model, clients download part of content from each other client besides from CDN servers. In this paper, we develop a peer-assisted CDN model based on previous work. Our model is quite different from previous work. We add Youtube VoD views popularity model aspect to our peer-assisted CDN model. From simulation-based experiments driven by real-world Youtube datasets, we find that peers in our model can contribute to deliver contents higher than previous work on the other hand number of access to cache in peer also higher than previous work. We do significance test for both metrics (peer contributions and number of access to cache) and find that  $p$ -values are less than 1% thus the results are significant.

**Index Terms**—Internet Video on Demand (VoD), P2P, CDN, and Caching.

## I. INTRODUCTION

Streaming content, especially video, represents a significant fraction of the traffic volume on the Internet, and it has become a standard practice to deliver this type of content using Content Delivery Networks (CDNs) such as Akamai and Limelight for better scaling and quality of experience for the end users. For example, YouTube uses Google cache and MTV uses Akamai in their operations.

With the spread of broadband Internet access at a reasonable flat monthly rate, users are connected to the Internet 24 hours a day and they can download and share multimedia content. P2P (peer to peer) applications are also widely deployed. In China, P2P is very popular; we see many P2P applications from China such as PPLive, PPStream, UUSe, Xunlei, etc. [1]. Some news broadcasters also rely on P2P technology to deliver popular live events. For example, CNN uses the Octoshape [2] solution that enables their broadcast to scale and offer good video quality as the number of users increases.

From the Internet provider point of view, the presence of so many always-on users suggests that it is possible to delegate a portion of computing, storage and networking tasks to the users, thus creating P2P networks where users can share files and multimedia content. Starting from file sharing protocols, P2P architectures have evolved toward video on demand and support for live events.

Alternatively, video contents can be efficiently distributed on services offered by managed network architectures and

CDN companies. The major issues of CDN are high deployment cost and good but not unlimited scalability in the long term. Given the complementary features of P2P and CDN, in recent years some hybrid solutions have been proposed and applied to the operational of CDN [3]–[5] to take the best of both approaches. In Peer assisted CDN, users can download content from CDN nodes from or other users or peers. A user may cache the content after download to serve requests from other users. Due to the complexity of the behavior of peers, the process should be done in the home gateway user where the ISP can control it.

In this work, we will revisit Guo et al., [6] work's PROP as basis to evaluate of the peer-assisted CDN and propose an improvement the model for the PROP. We also examine the characteristics of Internet VoD by investigating real-world datasets obtained from Youtube. In P2P assisted CDN for video on demand (VoD), most of researcher assume that catalog of video popularity rank is already established following zipf distribution. This become basis for P2P assisted CDN model in PROP [6]. Our work is quite different whereas we will use VoD view popularity to aid the PROP model. We use Youtube VoD view model for this purpose. A twofold of our contributions as follows: (1) We use the idea of VoD view popularity model to aid the PROP model. To our knowledge, the combination of PROP model and VoD view popularity model is the first. (2) From simulation-based experiments, we find that peer contributions become higher than the PROP model.

Our paper presentation as follows: (1) we describe related work in sect.II; (2) we explain detail of Youtube popularity evolution model in sect.III; (3) we explain the caching strategy for CDN and peer in sect.IV; (4) we explain our simulation design, simulator, and its evaluation in sect.V. Finally, we present our conclusions in section VI.

## II. RELATED WORK

Content Distribution Networks with peer assist have been successfully deployed on the Internet, such as Akamai [7], [3] and LiveSky [8]. The authors of [7] examine the risks and benefits of peer-assisted content distribution in Akamai and measure the effectiveness of its peer-assisted approach. The authors of [3] conclude from two real world traces that hybrid CDN-P2P can significantly reduce the cost of content distribution and can scale to cope with the exponential growth

of Internet video content. Yin et al. [8] described commercial operation of a peer-assisted CDN in China. LiveSky solved several challenges in the system design, such as dynamic resource scaling of P2P, low startup latency, ease of P2P integration with the existing CDN infrastructure, and network friendliness and upload fairness in the P2P operation. Xu et al. [9] used game-theory to show the right cooperative profit distribution of P2P can help the ISP to maximize the utility. Their model can easily be implemented in the context of current Internet economic settlements. Misra et al. [10] also mentioned the importance of P2P architecture to support content delivery networks. The authors use cooperative game theory to formulate simple compensation rules for users who run P2P to support content delivery networks.

The idea of telco- or ISP-managed CDN has been proposed in recent years. The complexity of the CDN business encourage telcos and ISPs to manage their own CDN, rather than allow others to run CDNs on their networks. It has been shown that it is cost effective [11] [12]. Kamiyama et al. [13] proposed optimally ISP operated CDN. Kamiyama et al. mentioned that, in order to deliver large and rich Internet content to users, ISPs need to put their CDNs in data centers. The locations are limited while the storage is large, making this solution effective, using optimum placement algorithm based on real ISP network topologies. The authors found that inserting a CDN into an ISP's ladder-type network is effective in reducing the hop count, thus reduce total link cost. Based on the author definition: Ladder-type network is a network with a maximum degree under 10. Cisco has initiated an effort to connect telco- or ISP-managed CDNs to each other, to form a CDN federation [11] using open standards [14]. They argue that the current CDN architecture is not close enough to the users and ISPs can fill this position.

The idea of utilizing the user's computation power to support ISP operation is not new. The Figaro project [15] proposed the residential gateway as an integrator of different networks and services, becoming an Internet-wide distributed content management for a proposed future Internet architecture [15]. Cha et al., [16] performed trace analysis and found that an IPTV architecture powered by P2P can handle a much larger number of channels, with lower demand for infrastructure compared to IP multicast. Jiang et al. [17] proposed scalable and adaptive content replication and request routing for CDN servers located in users' home gateways. Maki et al., [18] propose traffic engineering for peer-assisted CDN to control the behavior of clients, and present a solution for optimizing the selection of content files. Mathieu et al., [19] are using data gathered from France telecom network to calculate reduction of network load if customers are employed as peer-assisted content delivery.

Guo et al., [6] work's PROP is closest with our work. PROP uses local system (local counter) to calculate the segment popularity in peer-assisted proxy system. PROP uses popularity for proxy cache replacement strategy. In peer side, the author use utility function for cache replacement strategy. A utility function assigns numerical value to outcomes, in such a way that outcomes with higher utility are always preferred over outcomes with lower utilities. The utility function

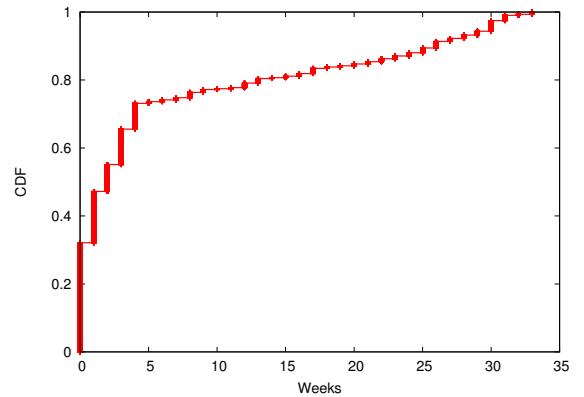


Fig. 1: Time to peak empirical distribution.

is also function from popularity. While the authors successfully show that the results are very good, the peer-assisted system behavior over time is not explain because the author focus on properties such as proxy cache size variations and peer cache size variations. The explanation of the optimal number of replicas is not also clear because unavailable information when the snapshot is taken. In our work, we complement Guo et al., [6] work with VoD viewing popularity evolution model and describe the behavior of the peer-assisted CDN over the time.

### III. CHARACTERIZING INTERNET VOD POPULARITY

Before analyzing the system description and video caching, we first examine the popularity characteristics of Internet VoD services. We use YouTube as example of VoD service. The studies of content popularity evolution are mostly considered in short time periods. Borghol et al., [20] measure the evolution of content popularity in long periods (36 weeks, from 3 August 2008 until 29 March 2009) in which view count statistics of YouTube.

In datasets, we have one-week spacing between consecutive snapshots. We can get how many times the video was viewed during the one-week period since last week or since snapshot  $(i - 1)$ . Borghol et al., [20] define time-to-peak for a video as its age (time since upload) at which its weekly viewing rate is the highest during measurement (from the first week until end of measurement).

The time-to-peak distributions is shown in fig.1. Figure 1 shows Borghol et al., [20] work that around three-quarters of a large fraction videos peak within the first six weeks since their upload and beyond six weeks we have uniform distribution thus the time-to-peak is exponentially distributed mixture with uniform distribution. Because we know the peak time (at-peak phase) of every video, we can also know before-peak phase and after-phase of every videos.

To estimate the the rate parameter of exponential part of time-to-peak distribution, we use Maximum Likelihood Estimation (MLE) [21]. Using MLE method, we can get exponential parameter  $\lambda = 0.59$ . For weekly views distribution, Borghol et al., [20] found that beta distribution is a good model to explain video views popularity evolution thus we follow

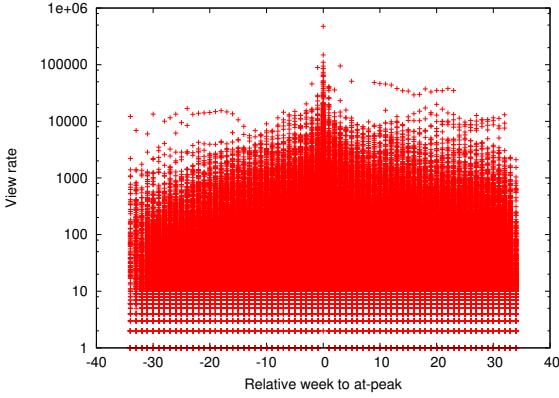


Fig. 2: View rate distribution versus week relative to at-peak phase week for every videos, where y-axis in logscale. Every points lie in negative x-axis mean view rate of every videos in before-peak phase. Every points lie in x-axis= 0 mean view rate of every videos at-peak phase. Every points lie in positive x-axis mean view rate of every videos in after-peak phase.

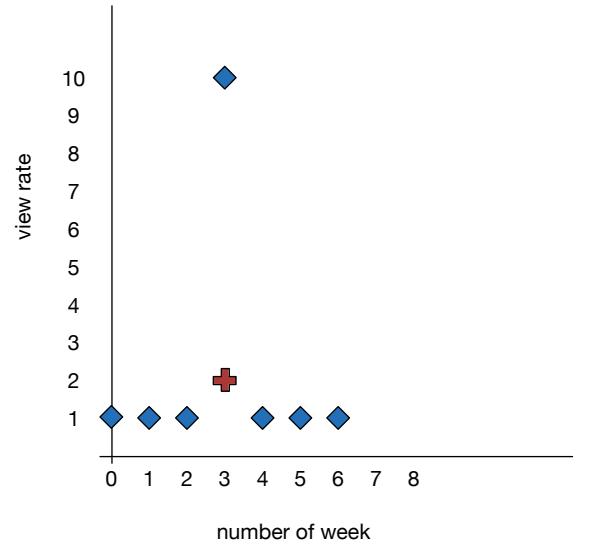


Fig. 4: graph of view rate distributio and red

A	view rate	1	1	1	10	1	1	1
	relative week to peak	-3	-2	-1	0	1	2	3

B	view rate	1	1	1	10	1	1	1
	relative week to peak	-3	-2	-1	0	1	2	3
	number of week	0	1	2	3	4	5	6

Fig. 3: view rate distribution example.

Borghol et al., [20] for weekly views distribution model. To reveal data distribution of view rate for every video, we plot view rate versus week where we shift week of view rate at-peak phase to zero. Therefore we can get view rate distribution relative to at-peak week as shown in fig. 2

#### IV. SYSTEM DESCRIPTION

In our work, we use Youtube VoD view model to aid our work that based from PROP. The Youtube VoD view model will be used in peer-caching strategy side to exploits the video popularity.

##### A. Peer caching strategy

As we mentioned before, we use Youtube VoD view model in peer caching strategy side. Our utility function is different from PROP. Our utility function need the estimation of video position whether the requested video is in before-peak phase, at-peak phase, or after-peak phase. How we estimate the video position is shown in fig. 3 and fig. 4. In fig. 3 part A, we have view rate (y-axis) and relative week to peak (x-axis) which is

view rate distribution versus week relative to at-peak phase as also shown completely in fig. 2. We transform these numbers by adding number of week and make number of week as x-axis, view rate as y-axis, and relative week to peak as z-axis fig. 3 part B. This transformation is shown in fig. 4 as diamond points. Assume there is a peer that request a video where the position of video in third week with the last week view rate  $vr = 2$  (we can get as this data from CDN) shown in fig. 4 as red cross. We want to estimate what is the position of that video. Is the video in at-peak phase, before-phase, or after phase. We can estimate the that video position by averaging relative week to peak numbers (the points at z-axis) of the nearest point from datasets. If the average value less than 0 we estimate the video position at before-peak phase, if the average value equal to 0 we estimate the video position at at-peak phase, and if the average value more than 0 we estimate the video position at after-peak phase. In this case, the nearest points are the point at third week (2, 1, -1) and the point at fifth week (4, 1, 1). By averaging the points at z-axis of the nearest points  $(-1+1)/2 = 0$ , we can get estimate that video is in at-peak phase.

Since we can estimate before-peak week, at-peak week, and after-peak week of video, we modified the original utility function from PROP as follows:

- In before-peak and after-peak phases, we assume that requests to video are low thus we only consider minimum popularity.

$$u = \frac{f(p) - f(p_{min})}{r^{\alpha+\beta}}. \quad (1)$$

- In at-peak phase, we assume that request to video are high thus we only consider maximum popularity.

$$u = \frac{f(p_{max}) - f(p)}{r^{\alpha+\beta}}. \quad (2)$$

$p$  represents popularity of the video,  $p_{min}$  represents estimation of minimum popularity in P2P system,  $p_{max}$  represents

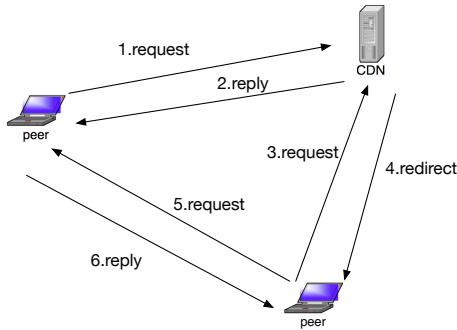


Fig. 5: Peer interaction in simulator.

estimation of maximum popularity in P2P system,  $r$  represents the number of replicas of the video in the system, and  $f(p)$  is monotonic non-decreasing function.  $\alpha$  and  $\beta$  are the adjustment factor. The CDN can calculate  $p_{min}$  and  $p_{max}$  then propagate to the P2P system. To able to track the simulation, we use default value from PROP for  $\alpha = \beta = 1$  and  $f(p) = \log p$ . We choose the video with the smallest  $u$  value as the candidate to be replaced when a peer's cache capacity is full.

## V. EVALUATION

In order to evaluate the proposed peer-caching strategy using before-peak, at-peak, and after-peak information from Youtube VoD view model, we have to compare our model to PROP model. We evaluate two metrics which is peer contribution to delivery contents during simulation and access frequency of cache during simulation. Peer contribution metric related to byte-hit-ratio. Byte-hit-ratio is defined as the total bytes contents served by peers normalized by the total bytes of video all peers and CDN consume. It means more peer contributions, more byte-hit-ratio because peer can get content from another peers. Access frequency of cache reflects the storage utilization. More access means more peer storage utilization. To evaluate these metrics, we developed peer-assisted CDN simulator.

### A. Simulation Design

An event driven simulator is developed using Python for this purpose. In our simulator, time is divided into rounds. During a round, a peer request a video.

In fig.5, we describe the process of a peer that requests a video in simulator. A peer and a CDN are implemented in object oriented model. When a peer requests a video, it always goes to a CDN server (step 1). The CDN provides the videos to the peer (step 2). If there is another peer request same video, that request will go to CDN (step 3). A CDN will check its record to see if there are some peers cache that requested video. If there are some peers cache that requested video, a CDN will reply with redirect message that asking a peer to download requested video from other peer (step 4). If there are no peers have requested video, a CDN will serve the video. A peer then can request the video to other peer

and get the video (step 5 and step 6). From this description, we can see that deploying peer-assisted CDN can save some traffic since the clients which form P2P network can sharing the contents or videos.

**1) Catalog Generator:** In catalog generator, we assume peer request a video to CDN following poisson process with a mean rate  $\lambda = 1.1$  [22] and we made it 3600 videos per hour, finally we generate video request for 360 days of simulation thus we have 31104000 requests by peers.

We calculate the number of videos at-peak time as follows: sample  $N$  value from the time-to-peak distribution and determine the number of videos  $n_j^{at}$  that peak at week  $j$ . Total number of video  $N = n_j^{before} + n_j^{at} + n_j^{after}$ .

Next, we determine view count terminus which the number of final view count of video. In view count terminus, we assume that a video will not get big additional view after at-peak phase. We assign view count terminus randomly from datasets. After determining view count terminus, we assign beta distribution parameter for every video.

Since we can estimate the time of at-peak phase for each video, we know the mode of beta distribution value and we can calculate  $\alpha$  and  $\beta$  value using the mode of distribution formula:  $m = \frac{\alpha-1}{\alpha+\beta-2}$ . We assign  $\alpha$  value randomly between 1 and 2 thus we can calculate  $\beta$  value. With the knowledge of beta distribution of every video and its view count terminus, we can know the view count and view rate of every video as function from time. In last step, we assign file size of video randomly between 1MB and 200MB. Finally, we have a catalog that consists of: video-id to be chosen, time when uploaded, view count terminus, at-peak week, and video size.

**2) Simulation Parameters:** The simulation parameters are follows:

- Length: 360 days.
- Video size: random between 1MB and 200MB.
- Peer capacity: [500MB,1000MB].
- CDN capacity: 10000MB and 20000MB.
- Number of peers: 10000.
- Number of videos: 10000.

We compare our proposed improvement of PROP to original PROP [6] implementation.

### B. Result and Discussion

Figure 6 shows CDF peer contribution for our model compare to PROP for peer capacity 500MB and 1000MB, with CDN capacity 20GB. One dot in a figure means how many percentage a video delivered by peers during simulation. Moreover fig. 6a shows the comparison of CDF of peer contributions between our model and PROP for peer capacity 500MB. Figure 6b shows the comparison of CDF of peer contribution between our model and PROP for peer capacity 1000MB. Figure 6c we compare our model between peer capacity 500MB and peer capacity 1000MB. Peer capacity 1000MB gives higher peer contribution because additional space makes a peer can cache more videos. From fig. 6b and fig. 6c, We can see that our model can gain higher peer contributions than PROP. We do significance statistical testing if our model has significantly different from PROP. We use

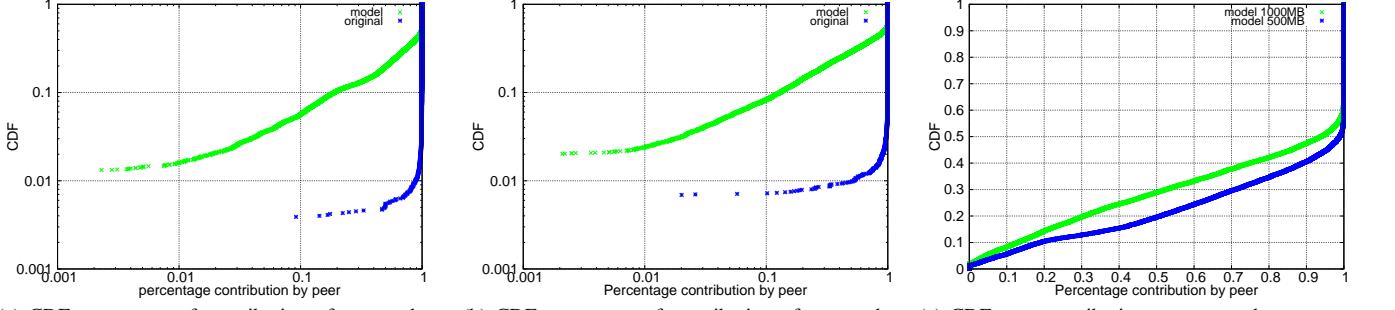


Fig. 6: CDF peer contribution for peer capacity 500MB and 1000MB, CDN capacity 10000MB. Model refers to our work and original refer to PROP [6].

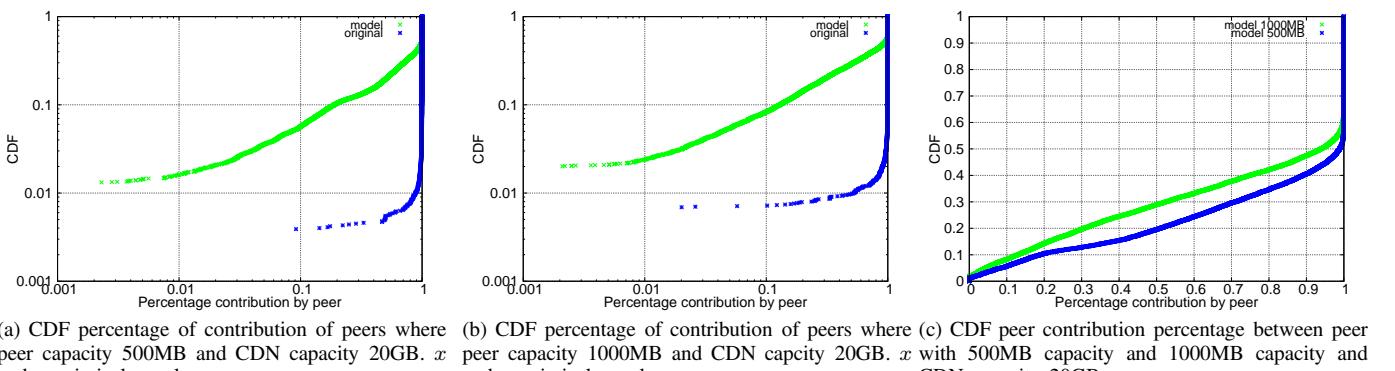


Fig. 7: CDF peer contribution for peer capacity 500MB and 1000MB, CDN capacity 20000MB. Model refers to our work and original refer to PROP [6].

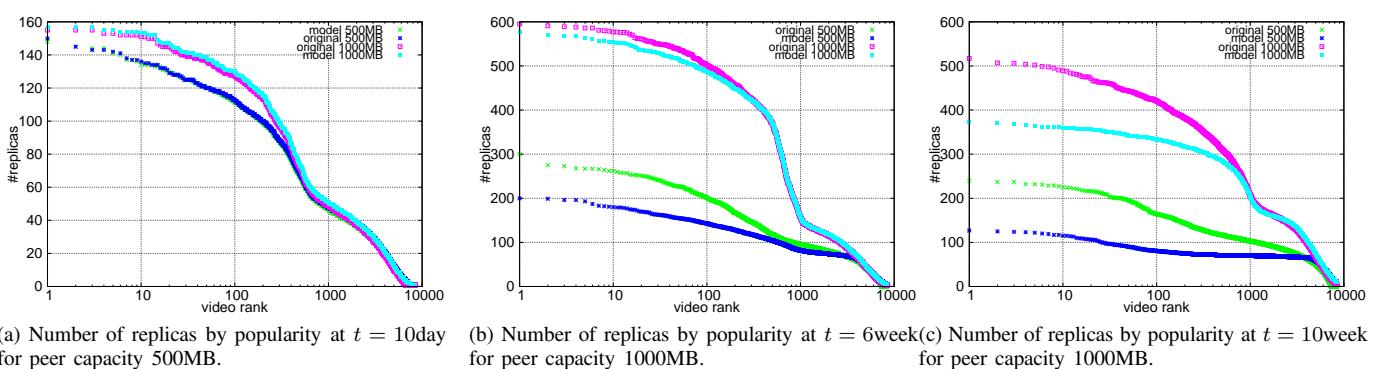
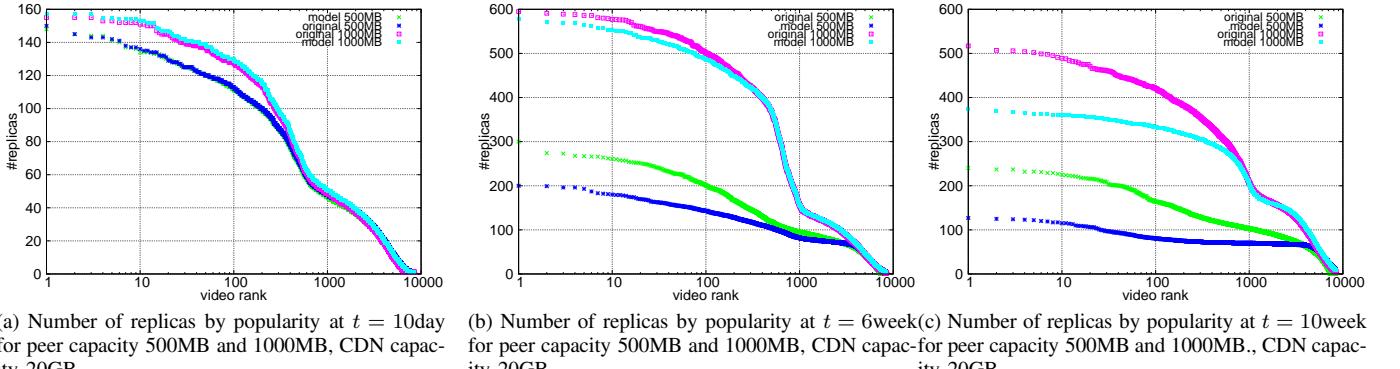


Fig. 8: Distribution number of video replicas at snapshot  $t = 10$ day,  $t = 6$ week, and  $t = 10$ week for peer capacity 500MB and 1000MB, CDN capacity 10GB.

Kolmogorov-Smirnov (KS) test for this purpose. KS-test tries to determine if two datasets differ significantly. It has the advantage of making no assumption about the distribution of data. KS-test reject the null hypothesis of no difference between datasets if  $p$ -value is small. For peer capacity 500MB when we compare our model and PROP, we get  $p$ -value  $0.5e - 005$ . In peer capacity 1000MB case, we get  $p$ -value

$0.46e - 005$ . Because both  $p$ -values are below 1%, we can reject null hypothesis that both data are the same thus our results are significant.

Figure 7 shows CDF peer contribution for our model compare to PROP for peer capacity 500MB and 1000MB, with CDN capacity 20GB. Figure 7a shows the comparison of CDF of peer contributions between our model and PROP



(a) Number of replicas by popularity at  $t = 10\text{day}$  for peer capacity 500MB and 1000MB, CDN capacity 20GB.

(b) Number of replicas by popularity at  $t = 6\text{week}$  for peer capacity 500MB and 1000MB, CDN capacity 20GB.

(c) Number of replicas by popularity at  $t = 10\text{week}$  for peer capacity 500MB and 1000MB, CDN capacity 20GB.

Fig. 9: Distribution number of replicas at snapshot  $t = 10\text{day}$ ,  $t = 6\text{week}$ , and  $t = 10\text{week}$  for peer capacity 500MB and 1000MB, CDN capacity 20GB.

for peer capacity 500MB. Figure 7b shows the comparison of CDF of peer contributions between our model and PROP for peer capacity 1000MB. From fig. 7 and fig. 7a show that our model can gain higher peer contributions than PROP. We also do significance statistical testing for both cases. We get  $p$ -value  $0.5e - 005$  for peer capacity 500MB case and  $p$ -value  $0.47e - 005$  for peer capacity 1000MB case. Because the  $p$ -values less than 1%, therefore the results are significant. Finally, fig. 7c shows the comparison between peer capacity 500MB and peer capacity 1000MB in our model. This figure shows that peer capacity 1000MB gives more contributions than peer capacity 500MB because additional space in peer makes a peer can cache more videos.

Figure 8 shows distribution of the number of replicas at snapshot  $t = 10\text{day}$ ,  $t = 6\text{week}$ , and  $t = 10\text{week}$  ranked by video popularity for peer capacity 500MB and 1000MB, CDN capacity 10GB between our model and PROP. Figure 8a shows distribution of the number of replicas comparison for peer capacity 500MB and 1000MB and CDN capacity 10GB sorted by popularity rank at snapshot  $t = 10\text{day}$ . This figure shows that on snapshot  $t = 10\text{day}$  there are no many different for peer capacity 500MB between our model and PROP except for peer capacity 1000MB PROP has lower number of replicas than our model for popular video rank between 1 until 1000 beyond that the number of replicas is same. Figure 8b shows distribution of the number of replicas comparison for peer capacity 500MB and 1000MB and CDN capacity sorted by popularity rank at snapshot  $t = 6\text{week}$ . Both peer capacity 500MB and 1000MB show that our model has lower number of replicas compare to PROP for popular video rank between 1 and 1000 for peer capacity 500MB while for peer capacity 1000MB our model gives lower replicas for popular video rank between 1 and 400. Peer capacity 1000MB gives higher replicas than 500MB because additional capacity make more room for peer to cache videos. Figure 8c shows distribution of the number of replicas comparison for peer capacity 500MB and 1000MB and CDN capacity 10GB sorted by popularity rank at snapshot  $t = 10\text{week}$ . Again, our model gives lower number of replicas compare to PROP. Snapshot  $t = 10\text{ weeks}$  gives lower number of replicas than snapshot at  $t = 6\text{week}$  because in our model the utility function knows that  $t = 10\text{week}$  is at after-peak

phase thus the model only considering minimum popularity for all videos.

## VI. CONCLUSION AND FUTURE WORK

This paper presents a scheme for a ISP managed peer-assisted CDN model that Some areas of improvement that we have identified for future are: We are also very interested to include energy trade off this peer-assisted CDN architecture in order to know how much energy saving by ISP and how much increase of energy at users home gateway side in this architecture.

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