

Peer-Assisted Content Distribution Aided by Video Popularity Evolution Model

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Abstract—Content distribution network (CDN) is widely used to efficiently deliver streaming media. The CDN with dedicated network bandwidths and hardware supports can provide high-quality streaming services but at a high cost. On the other side, the rise of peer-to-peer (P2P) networks are scalable but do not guarantee high quality streaming service due to the transient nature of peers. In this paper, we present peer-assisted CDN model, the peer-to-peer networks that can be used to help CDN to deliver streaming media. Our peer-assisted CDN model is aided by Youtube VoD views popularity model. We have comparatively evaluated our system through trace-driven simulations with synthetic workloads. Our results show that our model gives lower number of replicas while maintaining same number of peers contribution compare to previous work. We also do the significance to the number of replicas using the Kolmogorov-Smirnov statistic on two samples and we find our results are significant.

Index Terms—Internet Video on Demand (VoD), P2P, CDN, and Caching.

I. INTRODUCTION

Streaming content, especially video, represents a significant fraction of the traffic volume on the Internet, and it has become a standard practice to deliver this type of content using Content Delivery Networks (CDNs) such as Akamai and Limelight for better scaling and quality of experience for the end users. For example, YouTube uses Google cache and MTV uses Akamai in their operations.

With the spread of broadband Internet access at a reasonable flat monthly rate, users are connected to the Internet 24 hours a day and they can download and share multimedia content. P2P (peer to peer) applications are also widely deployed. In China, P2P is very popular; we see many P2P applications from China such as PPLive, PPStream, UUSe, Xunlei, etc. [1]. Some news broadcasters also rely on P2P technology to deliver popular live events. For example, CNN uses the Octoshape [2] solution that enables their broadcast to scale and offer good video quality as the number of users increases.

From the Internet provider point of view, the presence of so many always-on users suggests that it is possible to delegate a portion of computing, storage and networking tasks to the users, thus creating P2P networks where users can share files and multimedia content. Starting from file sharing protocols, P2P architectures have evolved toward video on demand and support for live events.

Alternatively, video contents can be efficiently distributed on services offered by managed network architectures and

CDN companies. The major issues of CDN are high deployment cost and good but not unlimited scalability in the long term. Given the complementary features of P2P and CDN, in recent years some hybrid solutions have been proposed and applied to the operational of CDN [3]–[5] to take the best of both approaches. In Peer assisted CDN, users can download content from CDN nodes from or other users or peers. A user may cache the content after download to serve requests from other users. Due to the complexity of the behavior of peers, the process should be done in the home gateway user where the ISP can control it.

In this work, we will revisit Guo et al., [6] work's PROP as basis to evaluate of the peer-assisted CDN and propose an improvement the model for the PROP. We also examine the characteristics of Internet VoD by investigating real-world datasets obtained from Youtube. We estimate the video popularity phase. Information of video popularity phase will be used for caching strategy In P2P assisted CDN for video on demand (VoD), most of researcher assume that catalog of video popularity rank is already established following zipf distribution. This become basis for P2P assisted CDN model in PROP [6]. Our work is quite different whereas we will use VoD view popularity to aid the PROP model. We use Youtube VoD view model for this purpose. A twofold of our contributions as follows: (1) We use the idea of VoD view popularity model to aid the PROP model. To our knowledge, the combination of PROP model and VoD view popularity model is the first. (2) From simulation-based experiments, we find that peer contributions become higher than the PROP model.

Our paper presentation as follows: (1) we describe related work in sect.II; (2) we explain detail of Youtube popularity evolution model in sect.III; (3) we explain the caching strategy for CDN and peer in sect.IV; (4) we explain our simulation design, simulator, and its evaluation in sect.V. Finally, we present our conclusions in section VI.

II. RELATED WORK

Content Distribution Networks with peer assist have been successfully deployed on the Internet, such as Akamai [7], [3] and LiveSky [8]. The authors of [7] examine the risks and benefits of peer-assisted content distribution in Akamai and measure the effectiveness of its peer-assisted approach. The authors of [3] conclude from two real world traces that

hybrid CDN-P2P can significantly reduce the cost of content distribution and can scale to cope with the exponential growth of Internet video content. Yin et al. [8] described commercial operation of a peer-assisted CDN in China. LiveSky solved several challenges in the system design, such as dynamic resource scaling of P2P, low startup latency, ease of P2P integration with the existing CDN infrastructure, and network friendliness and upload fairness in the P2P operation. Xu et al. [9] used game-theory to show the right cooperative profit distribution of P2P can help the ISP to maximize the utility. Their model can easily be implemented in the context of current Internet economic settlements. Misra et al. [10] also mentioned the importance of P2P architecture to support content delivery networks. The authors use cooperative game theory to formulate simple compensation rules for users who run P2P to support content delivery networks.

The idea of telco- or ISP-managed CDN has been proposed in recent years. The complexity of the CDN business encourage telcos and ISPs to manage their own CDN, rather than allow others to run CDNs on their networks. It has been shown that it is cost effective [11] [12]. Kamiyama et al. [13] proposed optimally ISP operated CDN. Kamiyama et al. mentioned that, in order to deliver large and rich Internet content to users, ISPs need to put their CDNs in data centers. The locations are limited while the storage is large, making this solution effective, using optimum placement algorithm based on real ISP network topologies. The authors found that inserting a CDN into an ISP's ladder-type network is effective in reducing the hop count, thus reduce total link cost. Based on the author definition: Ladder-type network is a network with a maximum degree under 10. Cisco has initiated an effort to connect telco- or ISP-managed CDNs to each other, to form a CDN federation [11] using open standards [14]. They argue that the current CDN architecture is not close enough to the users and ISPs can fill this position.

The idea of utilizing the user's computation power to support ISP operation is not new. The Figaro project [15] proposed the residential gateway as an integrator of different networks and services, becoming an Internet-wide distributed content management for a proposed future Internet architecture [15]. Cha et al., [16] performed trace analysis and found that an IPTV architecture powered by P2P can handle a much larger number of channels, with lower demand for infrastructure compared to IP multicast. Jiang et al. [17] proposed scalable and adaptive content replication and request routing for CDN servers located in users' home gateways. Maki et al., [18] propose traffic engineering for peer-assisted CDN to control the behavior of clients, and present a solution for optimizing the selection of content files. Mathieu et al., [19] are using data gathered from France telecom network to calculate reduction of network load if customers are employed as peer-assisted content delivery.

Guo et al., [6] work's PROP is closest with our work. PROP uses local system (local counter) to calculate the segment popularity in peer-assisted proxy system. PROP uses popularity for proxy cache replacement strategy. In peer side, the author use utility function for cache replacement strategy. A utility function assigns numerical value to outcomes, in

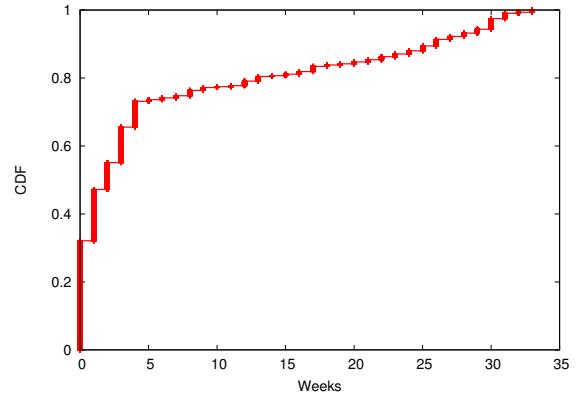


Fig. 1: Time to peak empirical distribution.

such a way that outcomes with higher utility are always preferred of outcomes with lower utilities. The utility function is also function from popularity. While the authors successfully show that the results are very good, the peer-assisted system behavior over time is not explain because the author focus on properties such as proxy cache size variations and peer cache size variations. The explanation of the optimal number of replicas is not also clear because unavailable information when the snapshot is taken. In our work, we complement Guo et al., [6] work with VoD viewing popularity evolution model and describe the behavior of the peer-assisted CDN over the time.

III. CHARACTERIZING INTERNET VOD POPULARITY

The objective of estimating Internet VoD popularity phase is to get popularity state of a video whether a video is at before-peal popularity phase or at peak popularity phase, or at after-peal popularity phase. We use YouTube as an example of VoD service where we get YouTube content popularity from Borghol et al., [20]. The datasets itself measure the evolution of content popularity in 36 weeks by recording video count statistics of YouTube.

In YouTube datasets, we have one-week spacing between consecutive snapshots. We can get how many times the video was view during the one-week period since last week or since snapshot ($i - 1$). Borghol et al., [20] define time-to-peak for a video as its age (time since upload) at which its weekly viewing rate is the highest during measurement (from the first week until end of measurement).

The time-to-peak distributions is shown in fig.1. Figure 1 shows Borghol et al., [20] work that around three-quarters of a large fraction videos peak within the first six weeks since their upload and beyond six weeks we have uniform distribution thus the time-to-peak is exponentially distributed mixture with uniform distribution. Because we know the peak time (at-peak phase) of every video, we can also know before-peak phase and after-phase of every videos.

To estimate the the rate parameter of exponential part of time-to-peak distribution, we use Maximum Likelihood Estimation (MLE) [21]. Using MLE method, we can get exponential parameter $\lambda = 0.59$. For weekly views distribution,

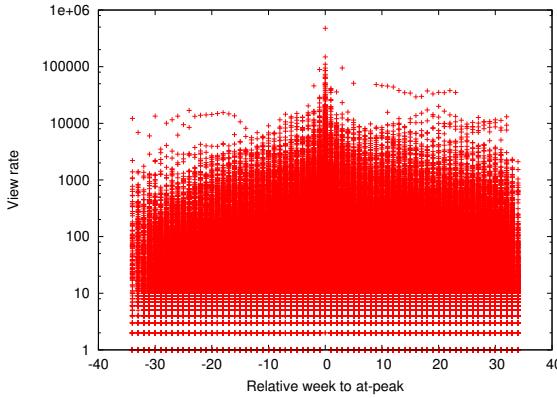


Fig. 2: View rate distribution versus week relative to at-peak phase week for every videos, where y-axis in logscale. Every points lie in negative x-axis mean view rate of every videos in before-peak phase. Every points lie in x-axis= 0 mean view rate of every videos at-peak phase. Every points lie in positive x-axis mean view rate of every videos in after-peak phase.

A	view rate	1	1	1	10	1	1	1
	relative week to peak	-3	-2	-1	0	1	2	3
↓								
B	view rate	1	1	1	10	1	1	1
	relative week to peak	-3	-2	-1	0	1	2	3
	number of week	0	1	2	3	4	5	6

Fig. 3: Transformation of view rate distribution. We add number week and make it as x -axis, View rate as y -axis, and relative week to peak as z -axis.

Borghol et al., [20] found that beta distribution is a good model to explain video views popularity evolution thus we follow Borghol et al., [20] for weekly views distribution model. To reveal data distribution of view rate for every video, we plot view rate versus week where we shift week of view rate at-peak phase to zero. Therefore we can get view rate distribution relative to at-peak week as shown in fig. 2

We use estimation of video popularity phase in peer caching strategy side. How we estimate the video popularity phase is shown in fig.3 and fig.4. In fig. 3 part A, we have view rate (y-axis) and relative week to peak (x-axis) which is view rate distribution versus week relative to at-peak phase as also shown completely in fig. 2. We transform these numbers by adding number of week and make number of week as x -axis, view rate as y -axis, and relative week to peak as z -axis fig. 3 part B. This transformation is shown in fig. 4 denote as diamond points. We want to estimate what is the phase of that video. Is the video in at-peak phase, before-phase, or after phase. We can estimate the that video phase by averaging

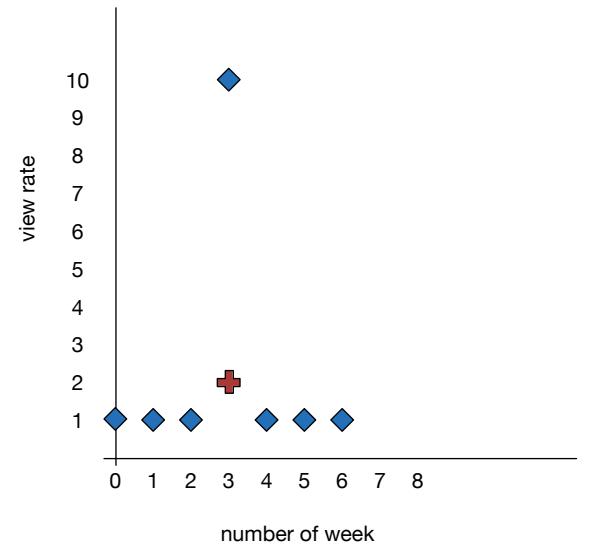


Fig. 4: Final view rate distribution after transformation where x -axis is number of week, y -axis is view rate.

relative week to peak numbers (the points at z-axis) of the nearest point from datasets. If the average value less than 0 we estimate the video is at before-peak phase, if the average value equal to 0 we estimate the video is at at-peak phase, and if the average value more than 0 we estimate the video is at after-peak phase.

For example: there is a peer that requests a video where the position of video is in fourth week with the last week view rate $vr = 2$ (we can get as this data from CDN) shown in fig. 4 denote as cross. In this case, the nearest points are the point at third week $(2, 1, -1)$ and the point at fifth week $(4, 1, 1)$. By averaging the points at z-axis of the nearest points $(-1 + 1)/2 = 0$, we can get estimate that video is in at-peak phase.

IV. SYSTEM DESCRIPTION

In our work, we use Youtube VoD view model to aid our work that based from PROP. The Youtube VoD view model will be used in peer-caching strategy side to exploits the video popularity.

A. Peer caching strategy

Since we can estimate before-peak week, at-peak week, and after-peak week of video, we modified the original utility function from PROP by adding a weight as follows:

$$u = \frac{(f(p) - f(p_{min}))(f(p_{max}) - f(p))}{r^{\alpha+\beta}} + z \quad (1)$$

Following [6] we can calculate p as follows:

$$p = \min\left(\frac{n_i^r}{t_i^r - t_a^i}, \frac{1}{t_{cur} - t_i^r}\right) \quad (2)$$

where n_i^r is number of requested video, t_i^r is last time the video is requested, t_a^i is the uploaded time of the video. z is proportion of view count that we get from Youtube datasets.

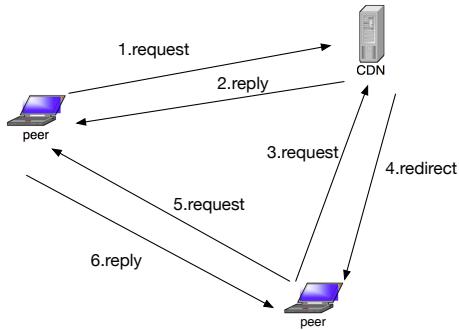


Fig. 5: Peer interaction in simulator.

In before-peak week, we get $z = 0.149538787758$, in at-peak week, we get $z = 0.470040393021$, and in after-peak week, we get $z = 0.380420819221$. In PROP's utility function, the difference between very popular videos and unpopular video is very difficult to differentiate. p represents popularity of the video, p_{min} represents estimation of minimum popularity in P2P system, p_{max} represents estimation of maximum popularity in P2P system, r represents the number of replicas of the video in the system, and $f(p)$ is monotonic non-decreasing function. α and β are the adjustment factor. The CDN can calculate p_{min} and p_{max} then propagate to the P2P system. To able to track the simulation, we use default value from PROP for $\alpha = \beta = 1$ and $f(p) = \log(p)$. We choose the video with the smallest u value as the candidate to be replaced when a peer's cache capacity is full. For unpopular video $f(p)$ will be very close to $f(p_{min})$ thus $f(p) - f(p_{min})$ will be very close to 0 then utility function become very small. For very popular video $f(p)$ will be very close to $f(p_{max})$ thus $f(p_{max}) - f(p)$ will be very close to 0 then utility function become very small. Linear addition of z factor here can help the differentiate the value of utility function.

V. EVALUATION

In order to evaluate the proposed peer-caching strategy using before-peak, at-peak, and after-peak information from Youtube VoD view model, we have to compare our model to PROP model. We evaluate three metrics which are peer contribution to delivery contents during simulation, access frequency of cache during simulation, and number of replicas. Peer contribution metric related to byte-hit-ratio. Byte-hit-ratio is defined as the total bytes contents served by peers normalized by the total bytes of video all peers and CDN consume. It means more peer contributions, more byte-hit-ratio because peer can get content from another peers. Access frequency of cache reflects the storage utilization. More access means more peer storage utilization. Number of replicas is also related to peer storage utilization. However, too many replicas will waste the storage resources. To evaluate these metrics, we developed a peer-assisted CDN simulator.

A. Simulation Design

An event driven simulator is developed using Python for this purpose. In our simulator, time is divided into rounds. During

a round, a peer request a video.

In fig.5, we describe the process of a peer that requests a video in simulator which derived from PROP. A peer and a CDN are implemented in object oriented model. When a peer requests a video, it always goes to a CDN server (step 1). The CDN provides the videos to the peer (step 2). If there is another peer request same video, that request will go to CDN (step 3). A CDN will check its record to see if there are some peers cache that requested video. If there are some peers cache that requested video, a CDN will reply with redirect message that asking a peer to download requested video from other peer (step 4). If there are no peers have requested video, a CDN will serve the video. A peer then can request the video to other peer and get the video (step 5 and step 6). From above description, we can see that deploying peer-assisted CDN can save some traffic since the clients which form P2P network can sharing the contents or videos.

1) Catalog Generator: The goal for catalog generator is to create a catalog video that consist video-id, time when a video is uploaded, a video size, view count terminus, and progress of videos popularity. We assume that a video is uploaded to server following Poisson process with mean rate $\lambda = 1$ thus we can get the time when a video is uploaded. The view count terminus for every video is assigned randomly uniform from YouTube datasets and video size for every video is assigned randomly uniform between 1 and 200MB. Finally, we have a catalog that consists of: video-id, time when a video is uploaded, view count terminus, and video size.

2) Peer Request Generator: In catalog generator, we assume peer request a video to CDN following poisson process with a mean rate $\lambda = 1$ [22]. We assume that a peer choose a video based on a preference from view count and view rate that we can get from catalog generator. Because peer requests are generate from catalog, the video request will follow global popularity video from YouTube

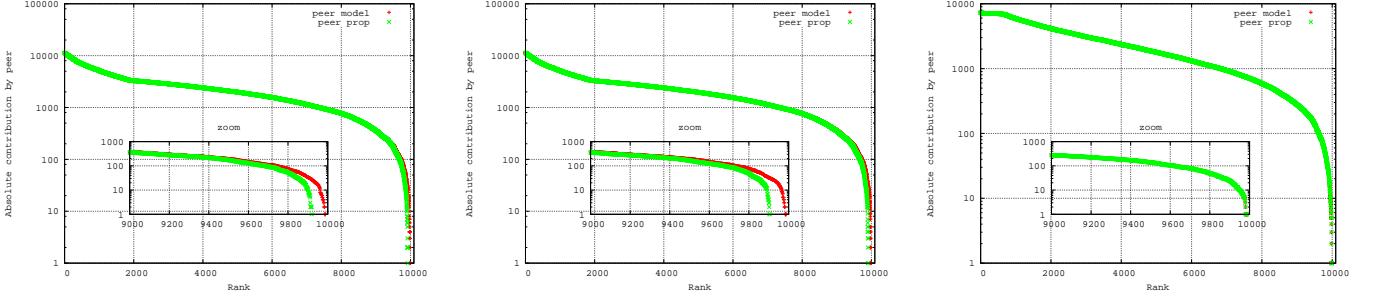
3) Simulation Parameters and Scenarios: The simulation parameters are follows:

- Length: 360 days.
- Video size: uniform random between 1MB and 200MB.
- Peer capacity: 500MB.
- CDN capacity: 10000MB.
- Number of peers: 100000.
- Number of videos: 10000.

There are three scenarios in our simulations. First, peers choose a video that has a popularity following from YouTube data sets that we already explained in V-A1. Second, peers choose a video that has a popularity following from YouTube data sets and we shift the requests time four weeks. Third, peers choose a video that has a popularity following zipf distribution with rate= 0.9 [23]. We compare our results to original PROP [6] implementation.

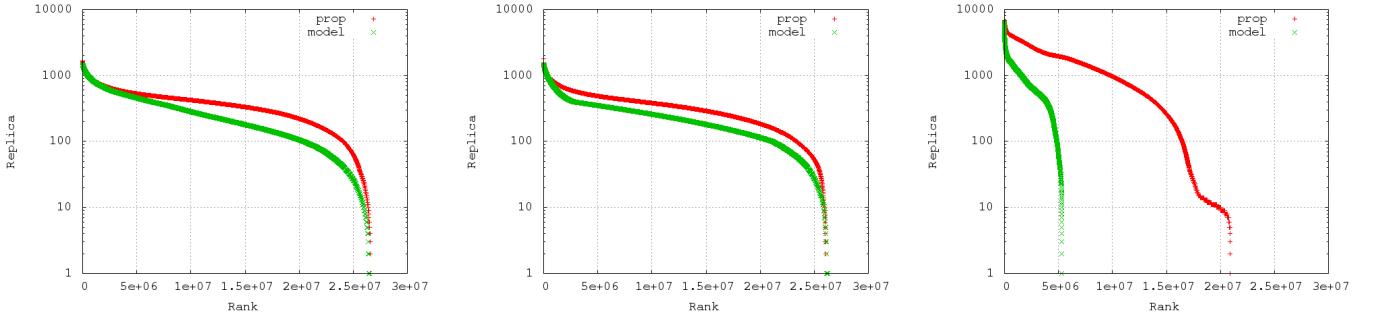
B. Result and Discussion

Figure 6a, 6b, and 6c show the absolute peer contribution to deliver videos compared between model and prop. Figure 6a and fig.6b show same pattern. The peers give more contribution in the tail while in the third scenario the peer contribution



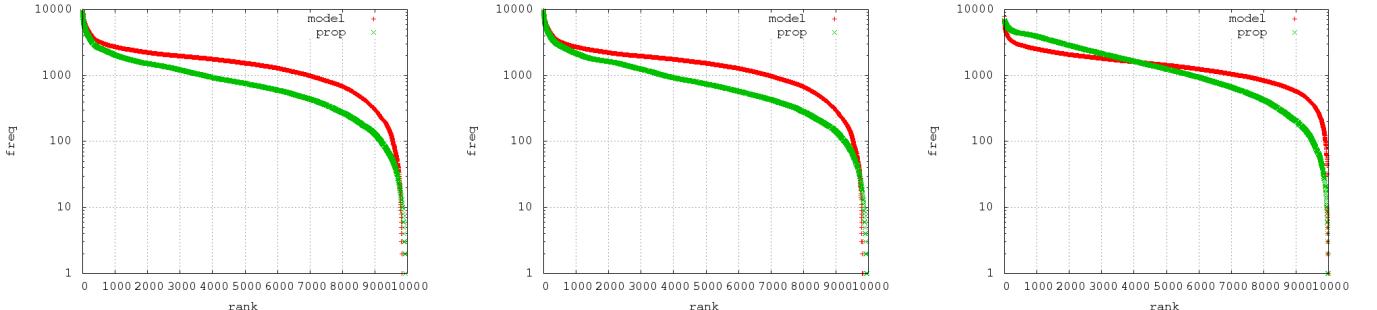
(a) Absolute of contribution of peer for the first scenario where y -axis in log-scale.
(b) Absolute of contribution of peer for the second scenario where y -axis in log-scale.
(c) Absolute contribution of peers for the third scenario where y -axis in log-scale.

Fig. 6: Peer contributions compared between model and PROP.



(a) Number of a video replicas when a peer request a video for the first scenario where y axis in log-scale.
(b) Number of a video replicas when a peer request a video for the second scenario where y axis in log-scale.
(c) Number of a video replicas when a peer request a video for the third scenario where y axis in log-scale.

Fig. 7: Comparison of available replicas between model and prop when a peer requests a video.



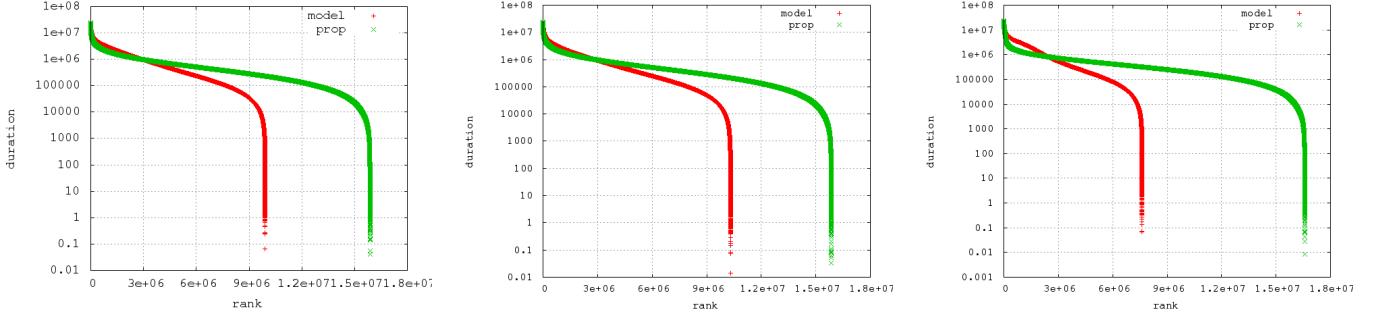
(a) Frequency a video in peers for the first scenario.
(b) Frequency a video in peers for the second scenario.
(c) Frequency a video in peers for the third scenario.

Fig. 8: Frequency a video in peers compared between model and prop.

is mostly same between model and PROP. A peers can give more contribution because a video has longer duration than other videos in a peer's cache thus other peer's requests are served by the peer. A video has longer duration than other videos in peer's cache because that a video has bigger utility function than other videos for example a video that will enter the cache.

Figure 7a, 7b, and 7c show the number of a video replicas available in system when a peer requests a video. As we can

see from all figures, the model gives us lower number of replicas than PROP. The model gives lower number of replicas than PROP because when a peer requests a video, that peer does not cached the video. We can see the proportion of cached and not-cached event in table.I and detail of the video phase breakdown is shown in table.II. In model, not-cached events take around 65% from all events and majority of video phase is in after-peak for both cached events and not-cached events. Because the majority of video phase is in after-peak for both



(a) Cache duration in peers for the first scenario. (b) Cache duration in peers for the second scenario. (c) Cache duration in peers for the third scenario.

Fig. 9: Duration compared between model and prop.

TABLE I: Percentage of Cached events and Not-Cached events in Model and PROP

Scenario	Type	Cached	Not-Cached
Scenario 1	Model	33.5%	66.5%
	PROP	52%	48%
Scenario 2	Model	34.8%	65.2%
	PROP	52.6%	47.4%
Scenario 3	Model	32.4%	67.6%
	PROP	67.7%	32.3%

cached events and not-cached events, In PROP, cached events take around 52% from all events for the first scenario and the second scenario, while for the third scenario is 67.7%. In model not-cached events are bigger than PROP, means peers do not cached the videos thus we get lower replicas number than PROP.

Denote u_{dl} is the minimum utility function for a video inside the cache and u_{ms} is utility function for a video that will enter the cache, p_{dl} is the popularity for a video inside the cache and p_{ms} is the popularity for a video that will enter the cache. In order a requested video is cached by a peer, the utility function for u_{dl} must be lower than the utility function for u_{ms} .

$$u_{dl} < u_{ms} \quad (3)$$

$$\frac{(f(p_{dl}) - f(p_{min}))(f(p_{max}) - f(p_{dl}))}{r_{dl}^{\alpha+\beta}} + z_{dl} < \frac{(f(p_{ms}) - f(p_{min}))(f(p_{max}) - f(p_{ms}))}{r_{ms}^{\alpha+\beta}} + z_{ms} \quad (4)$$

We assume that number of replicas are same, thus:

$$(f(p_{dl}) - f(p_{min}))(f(p_{max}) - f(p_{dl})) - (f(p_{ms}) - f(p_{min}))(f(p_{max}) - f(p_{ms})) < z_{ms} - z_{dl} \quad (5)$$

Since p_{min} and p_{max} are same, we can arrange the equation become:

$$f(p_{ms}) - f(p_{dl}) > z_{dl} - z_{ms} \quad (6)$$

TABLE II: Percentage of Video Phase for Model in cached and not-cached events

Scenario	Type/Events	Before-Peak	At-Peak	After-Peak
Scenario 1	Model-Cached	8.2%	1.2%	24.1%
	Model-Not-Cached	11.2%	0.8%	54.5%
Scenario 2	Model-Cached	6.2%	1.2%	29.8%
	Model-Not-Cached	5.2%	0.8%	56.8%
Scenario 3	Model-Cached	8.0%	1.8%	22.7%
	Model-Not-Cached	15.1%	0.8%	51.6%

and comparing with PROP:

$$f(p_{ms}) - f(p_{dl}) > 0 \quad (7)$$

$$f(p_{ms}) - f(p_{dl}) > z_{dl} - z_{ms} \quad (8)$$

As we know from table.II that the majority of a requested video phase is in after-peak and a requested video phase that is in at-peak phase is very small portion, then we can see that $z_{dl} - z_{ms}$ majority will be in negative term. While comparing to PROP, the left term $f(p_{ms}) - f(p_{dl})$ will be equal to 0 thus as result in our model a peer rejects the a requested video to be cached.

Figure 8a, 8b, and 8c show the frequency a video stay in peers compared between model and prop. As all figure show the model has higher frequency than prop to stay in peers except for the beginning rank of data where the model has same frequency with prop in first and second scenario. In the third scenario, in the beginning rank of data the model has lower frequency than prop, then around rank 1000 the model has higher frequency than prop until the end of data. The frequency a video stay in a video can also be viewed in fig 9a, 9b, and 9c, where in the model some videos have longer cache duration than prop, while others have shorter cache duration than prop. Thus, we can see the relationship between cache duration and frequency a video stays in peers.

VI. CONCLUSION AND FUTURE WORK

This paper presents a scheme for peer-to-peer network can help CDN to deliver the content over the Internet. We show that by introducing the weight to utility function we can increase the peer contribution to deliver the content while decreasing required replicas. We found that there are no much different between the first scenario, the second scenario and

the third scenario in peer contribution to deliver a video. We found that in the all scenarios, the model gives lower replica than PROP. This is because in the model, we found that not-cached events are higher than cached events. Therefore, in the model the number of available replicas are lower than PROP. We also did the significance test to the number of replicas using the Kolmogorov-Smirnov statistic on two samples and we find that for all scenarios the p -values are less than 1% thus the results are significant.

Some areas of improvement that we have identified for future are: The energy trade off this peer-assisted CDN architecture in order to know how much energy saving by ISP and how much increase of energy at users home gateway side in this architecture since we have higher peer contribution. More numerical experiments for other zipf shape parameters.

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