

Compsys Recap

10/01/24

Understanding, races and Virtual address translation
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Agenda

- OS in general
- Races
- Virtual memory and virtual address translation
- Exam exercises



OS in general

- process vs program
 - program is the “dead” code
 - process is the “live” instance running the program
- Kernel
 - “Always resident code that services request from the hardware and manages processes”
 - Unprivileged, must make calls to kernel to switch to privileged state (interrupts)
 - The kernel handles: hardware, system memory, File I/O and context switching



Processes, Threads and races

- Processes
 - duplicate of parent, but with own address space (changes are not reflected)
 - Child processes must be reaped, adopted but init process if termination of parent. Processes that never terminates are “zombie children”
- Threads
 - run in same address as the calling process (changes are reflected)
 - have own thread context: ID, SP, PC, general purpose registers, condition codes
 - can access “critical memory” must be handled with semaphores (mutexes) and / or condition variables
- Races
 - code depend on order of execution

Fork example from exam:

```
1  #include <stdio.h>
2  #include <unistd.h>
3  #include <sys/types.h>
4  #include <sys/wait.h>
5  void main () {
6      if (fork() == 0) {
7          printf("1");
8          if (fork() == 0) {
9              printf("2");
10             }
11         } else {
12             pid_t pid = fork();
13             if (waitpid(pid, NULL, 0) > 0) {
14                 if (fork() == 0) {
15                     printf("3");
16                 } else {
17                     printf("4");
18                 }
19             }
20             printf("5");
21         }
22     }
```



Virtual addresses

- Why is virtual memory useful?:
 - caching
 - memory management
 - memory protection
- Components of virtual address:
 - VPN, VPO, Tag og Index
 - TLB
 - PPN and PPO (physical addresses)
 - Page table (fully associative)



Translation of virtual addresses

- split the address in its components VPO, VPN, and then split the VPN in index and tag
- bits of the VPO = $\log_2(\text{pageSize})$
- bits of VPN = rest
- bits of set = $\log_2(\text{\#sets})$ bits of VPN
- bits of Tag = rest of vpn
- Good idea to mark this
- Translate hex-adresses to bit addresses (each digit corresponds to 4 bits)



Translation cont. (algorithm for translating)

- read description to find page size and #sets
- translate the hex to binary
- Find VPN, index and set
- look up in TLB with index and set
 - if the tag exists AND the valid bit is set it is a hit, otherwise a miss.
- if no hit in TLB look in page table , with VPN
 - if VPN exists AND the valid bit is set hit, otherwise page fault and a new page must be brought from memory
- If hit, copy ppn and ppo = vpo



Exam set:

Exam-set 2022- 23

if time

Re-Exam-set 2022-23