

Network Layer & Security

A 45 minutes walkthrough of old exam questions and some general tips

Jakob Holst Svenningsen



UNIVERSITY OF COPENHAGEN
FACULTY OF SCIENCE

Network Layer



Good things to remember about the Network Layer

- Routers forward datagrams
- Network Protocol used is **IP** (IPv4 and IPv6)
- Network control plane
 - Controls forwarding process (Routing)
- Network data plane
 - Performs actual forwarding (Forwarding)
- Routing Algorithms
 - calculates path from sender to receiver
- Network Service Model
 - Best-effort service
 - Encryption is possible at network layer (e.g. IPSec)
- In FW table longest prefix matching is performed
- Packet loss can occur in the router, when there is no memory left for incoming packets
- **Subnet** e.g. 223.1.1.0/24 means hosts here have address 223.1.1.xxx (first 24 bits are fixed)
- **DHCP** is an app-layer protocol for automatic assignment of IP-addresses
- **Network Address Translation (NAT)** solves problem with number of available IPs

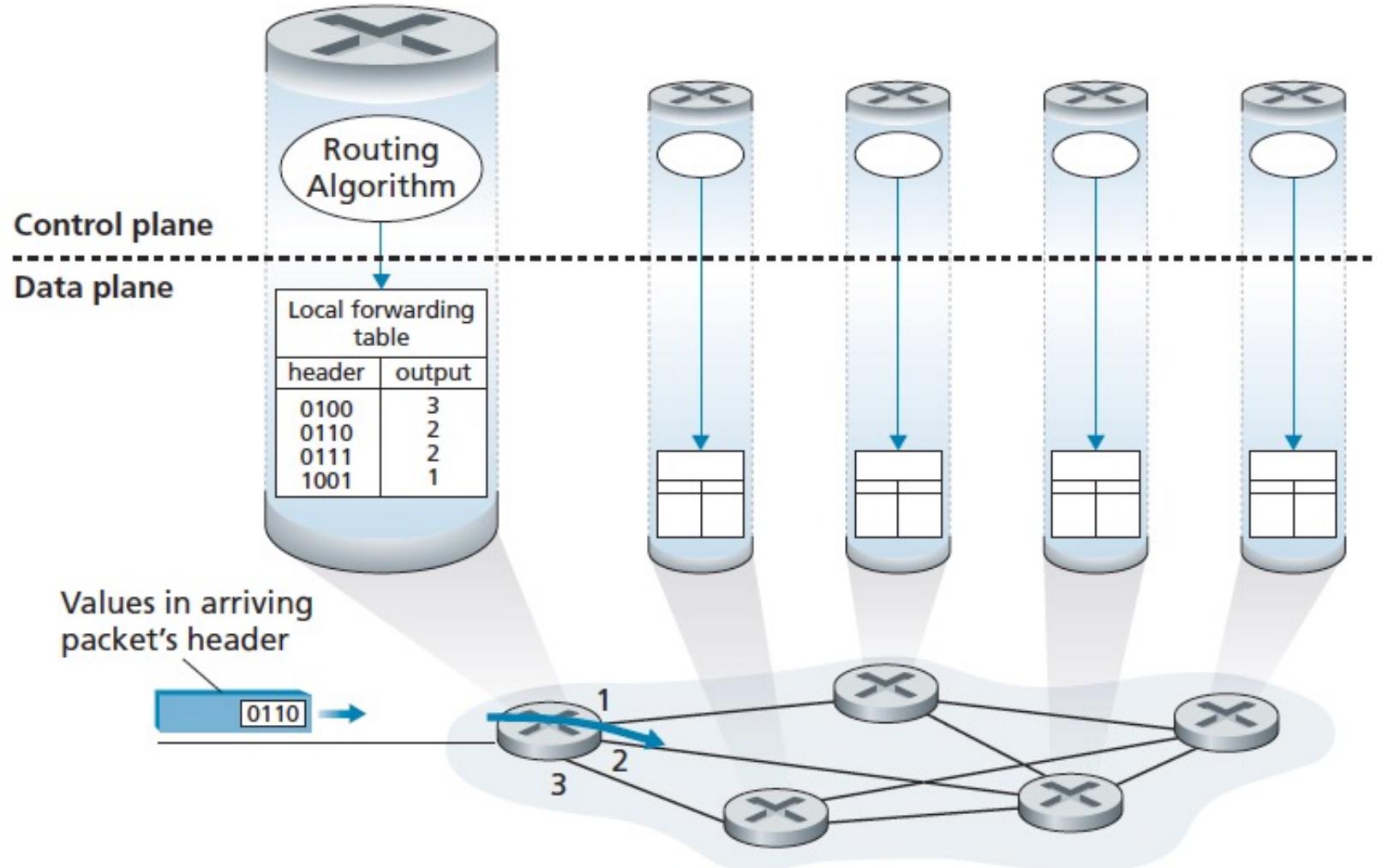


Figure 4.2 ♦ Routing algorithms determine values in forward tables

Control plane is responsible for routing, while the actual forwarding happens at the data plane

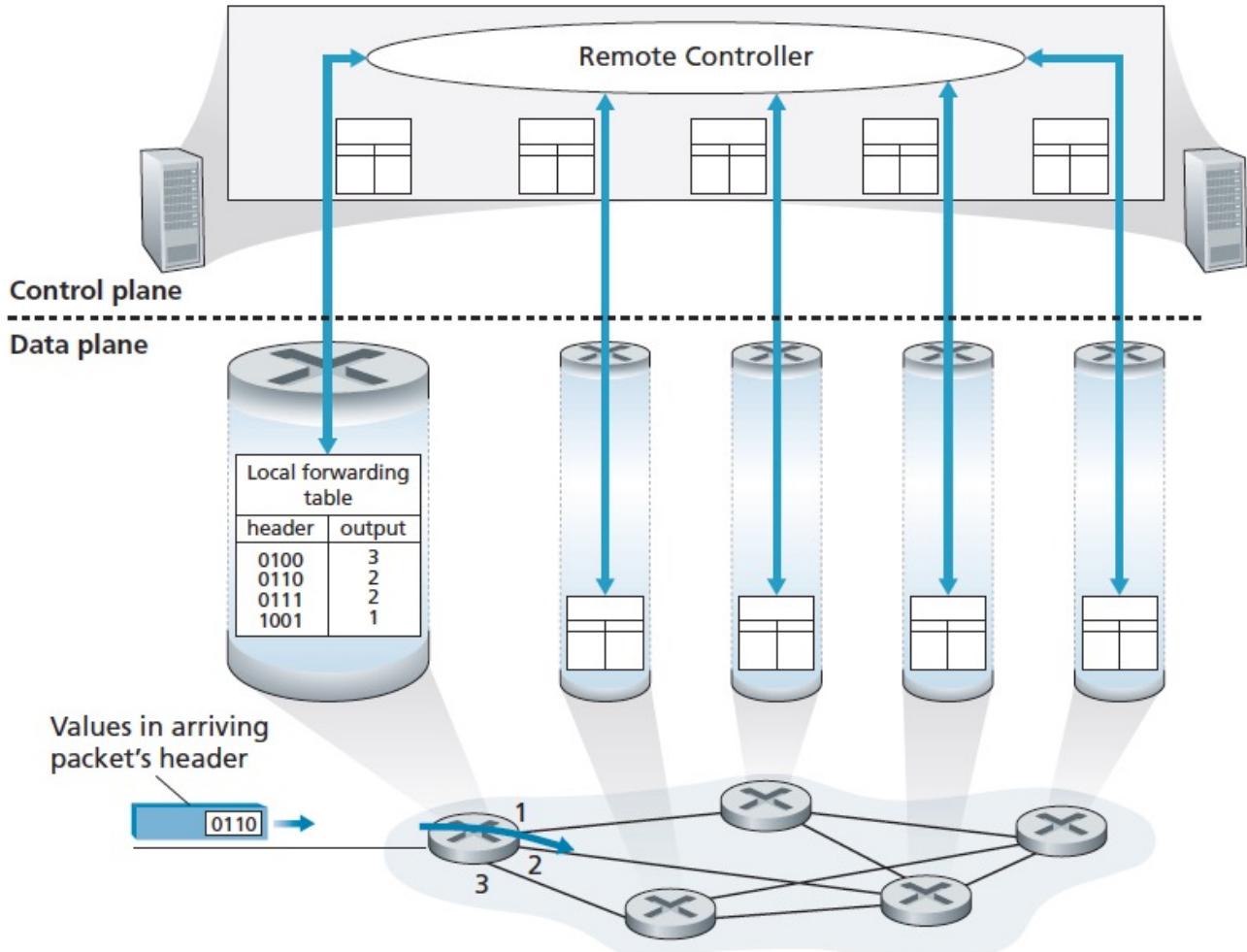


Figure 4.3 ♦ A remote controller determines and distributes values in forwarding tables

Today, the routing is often implemented in software and a centralized server (Remote Controller) tells the router how datagrams should be forwarded. It is however, possible to manually alter a forwarding table in each router.

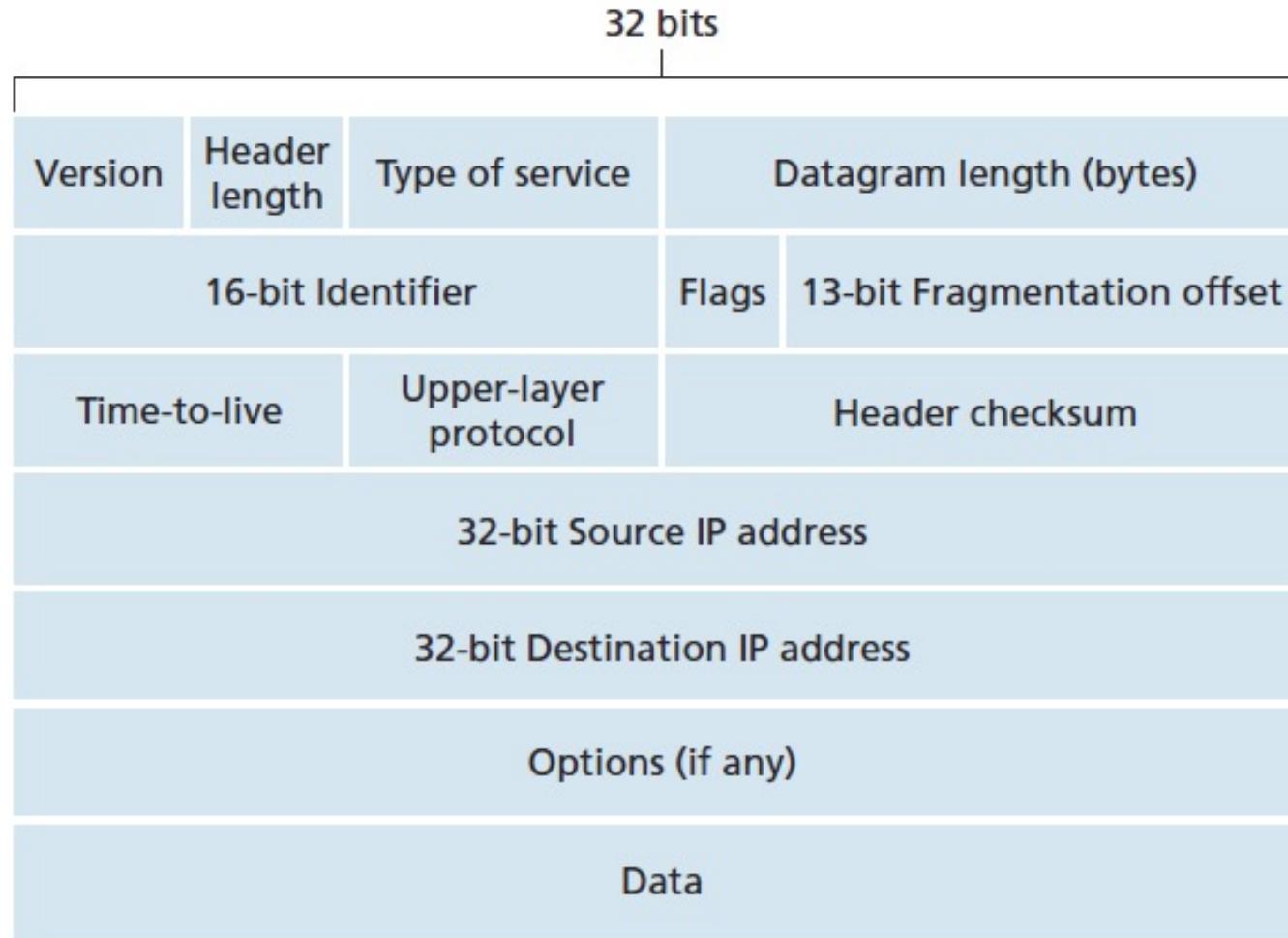


Figure 4.17 ♦ IPv4 datagram format

The IP-datatype is just yet another protocol that looks like this

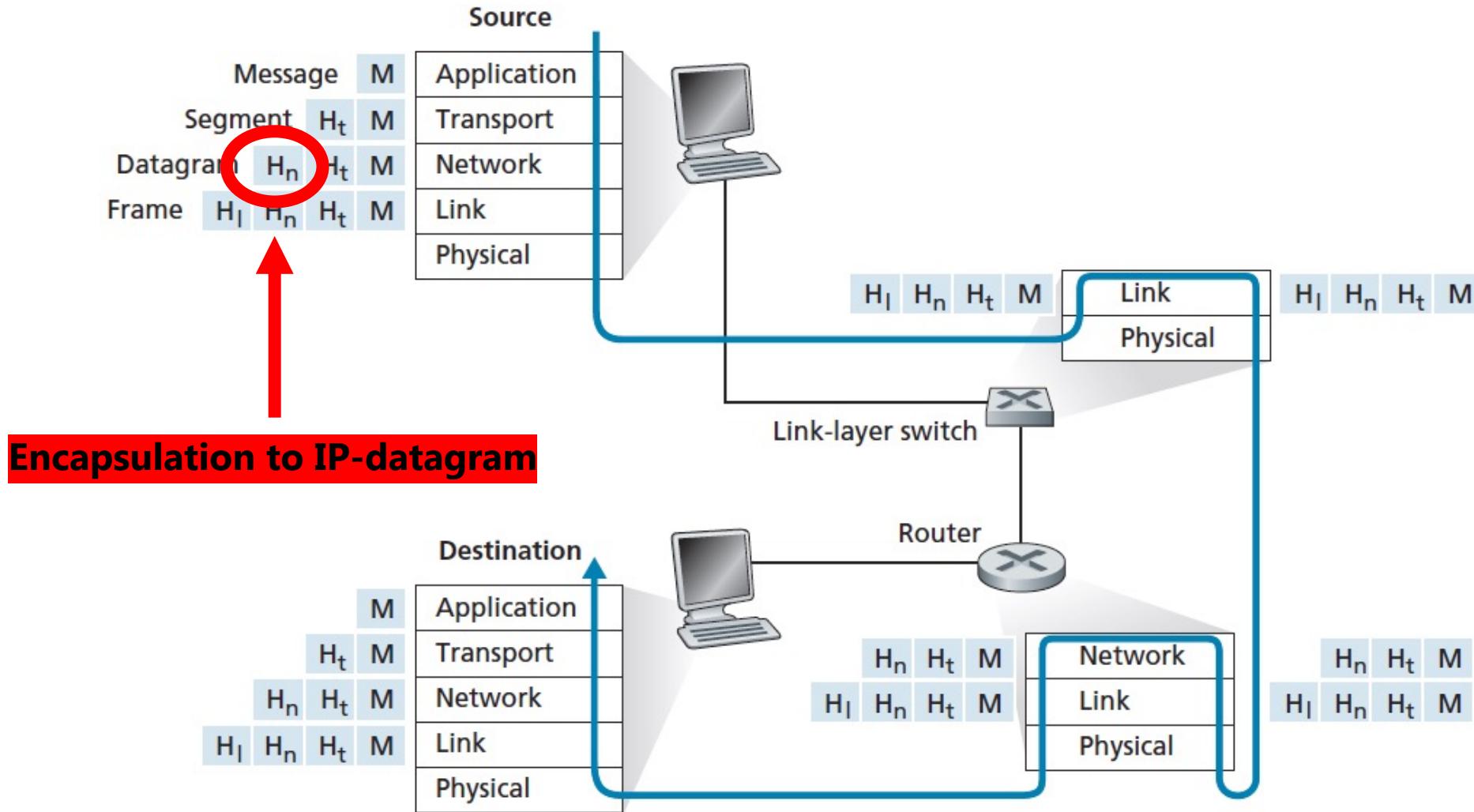


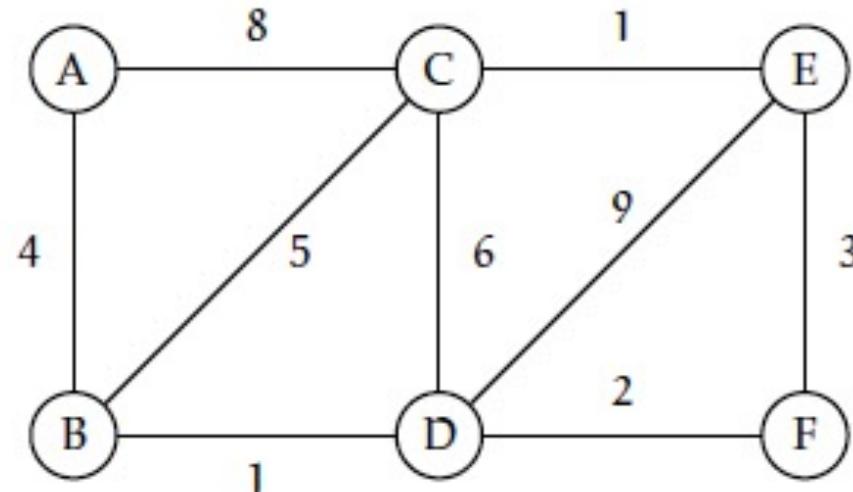
Figure 1.24 • Hosts, routers, and link-layer switches; each contains a different set of layers, reflecting their differences in functionality

In a broader context, higher level packets are encapsulated to IP-datagrams at the network layer. At the destination, we “unwrap” layer by layer. This abstraction allows arbitrary Application layer protocols to rely on IP

Link-State (LS) Routing Algorithm – Reexam 2020-21

- Centralized Routing Algorithm
- Essentially Dijkstra's shortest paths algorithm

Question 3.3.1: Consider the network topology outlined in the graph below.

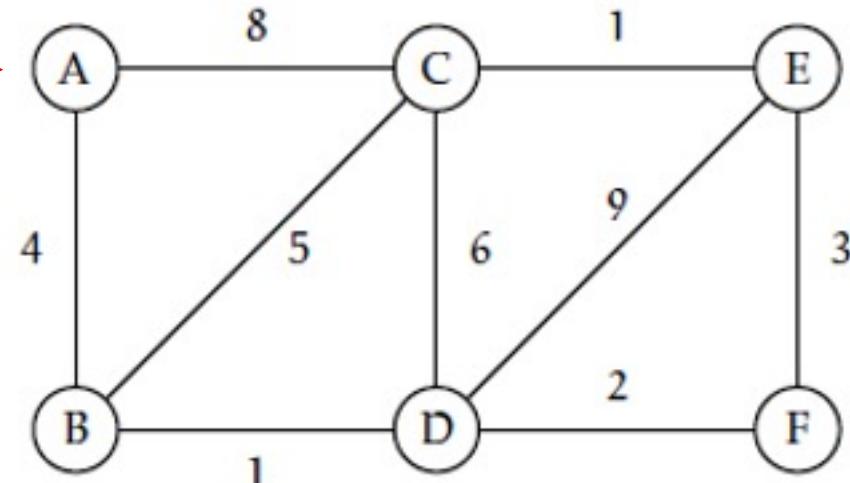


Apply the link state routing algorithm and compute the forwarding table on node A by filling out the following tables.

Solution Method

Here, **A**, is the src

- $D(v)$: Cost of least path from **u** (source) to **v** in this iteration
- $p(v)$: Previous node of **v** along least-cost path
- N' : subset of nodes, that are part of shortest path



Steps of the algorithm:

Step	N'	$D(B), p(B)$	$D(C), p(C)$	$D(D), p(D)$	$D(E), p(E)$	$D(F), p(F)$
0	A	4, A	8, A	inf	inf	inf
1						
2						
3						
4						
5						

Initialization:

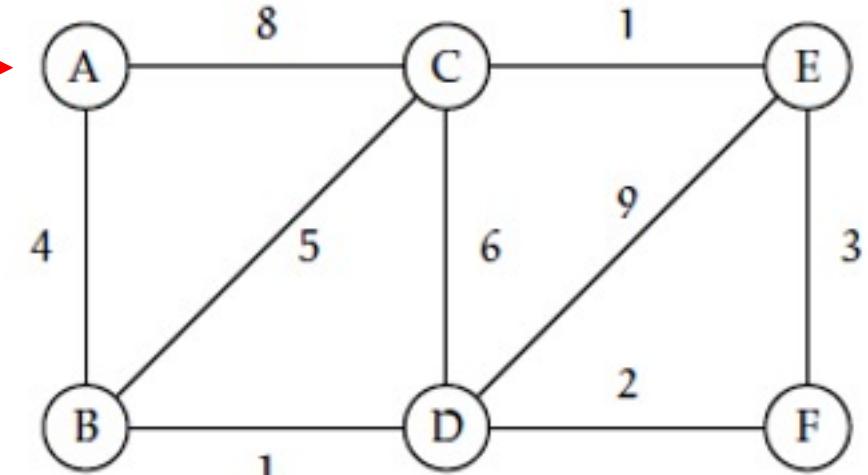
```

 $N' = \{u\}$ 
for all nodes  $v$ 
  if  $v$  is a neighbor of  $u$ 
    then  $D(v) = c(u, v)$ 
  else  $D(v) = \infty$ 
  
```

Execution table on node A.

Solution Method

Here, A, is the src



Loop

find w not in N' such that $D(w)$ is a minimum
add w to N'

update $D(v)$ for each neighbor v of w and not in N' :

$$D(v) = \min(D(v), D(w) + c(w, v))$$

/* new cost to v is either old cost to v or known
least path cost to w plus cost from w to v */

until $N' = N$

1: If we choose B as w, then $D(w)$ is minimum

2: Iterate over all neighbours of w (which is B here) that are not in the subset N'

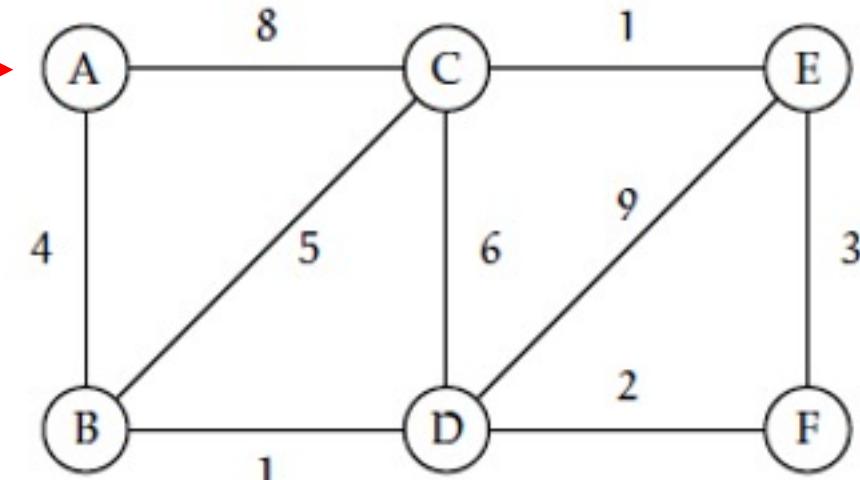
Steps of the algorithm:

Step	N'	$D(B), p(B)$	$D(C), p(C)$	$D(D), p(D)$	$D(E), p(E)$	$D(F), p(F)$
0	A	4,A	8,A	inf	inf	inf
1	A,B	4,A	8,A	5,B	inf	inf
2						
3						
4						
5						

Execution table on node A.

Solution Method

Here, A, is the src



Loop

find w not in N' such that $D(w)$ is a minimum
add w to N'

update $D(v)$ for each neighbor v of w and not in N' :

$$D(v) = \min(D(v), D(w) + c(w, v))$$

/* new cost to v is either old cost to v or known
least path cost to w plus cost from w to v */

until $N' = N$

1: If we choose D as w, then $D(w)$ is minimum

2: Iterate over all neighbours of w (which is D here) that are not in the subset N'

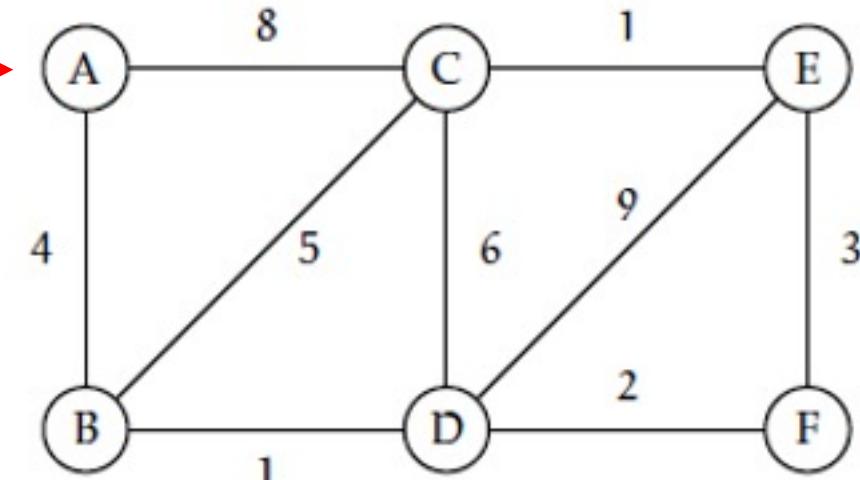
Steps of the algorithm:

Step	N'	$D(B), p(B)$	$D(C), p(C)$	$D(D), p(D)$	$D(E), p(E)$	$D(F), p(F)$
0	A	4,A	8,A	inf	inf	inf
1	A,B	4,A	8,A	5,B	inf	inf
2	A,B,D	4,A	8,A	5,B	14,D	7,D
3						
4						
5						

Exchanging table on node A.

Solution Method

Here, A, is the src



Loop

find w not in N' such that $D(w)$ is a minimum
add w to N'

update $D(v)$ for each neighbor v of w and not in N' :

$$D(v) = \min(D(v), D(w) + c(w, v))$$

/* new cost to v is either old cost to v or known
least path cost to w plus cost from w to v */

until $N' = N$

**1: If we choose F as w,
then $D(w)$ is
minimum**

**2: Iterate over all
neighbours of w
(which is F here) that
are not in the subset
 N' .**

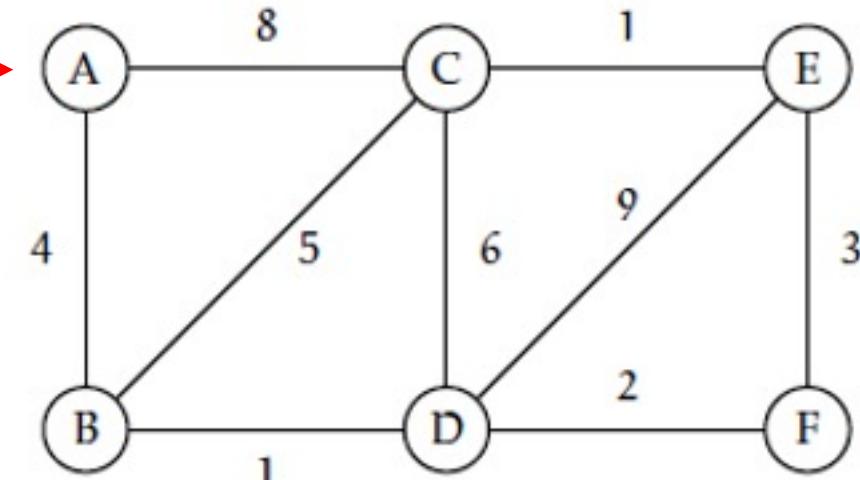
Steps of the algorithm:

Step	N'	$D(B), p(B)$	$D(C), p(C)$	$D(D), p(D)$	$D(E), p(E)$	$D(F), p(F)$
0	A	4,A	8,A	inf	inf	inf
1	A,B	4,A	8,A	5,B	inf	inf
2	A,B,D	4,A	8,A	5,B	14,D	7,D
3	A,B,D,F	4,A	8,A	5,B	10,F	7,D
4						
5						

Experiments table on node A.

Solution Method

Here, A, is the src



Loop

find w not in N' such that $D(w)$ is a minimum
add w to N'

update $D(v)$ for each neighbor v of w and not in N' :

$$D(v) = \min(D(v), D(w) + c(w, v))$$

/* new cost to v is either old cost to v or known
least path cost to w plus cost from w to v */

until $N' = N$

1: If we choose C as w, then $D(w)$ is minimum

2: Iterate over all neighbours of w (which is C here) that are not in the subset N' .

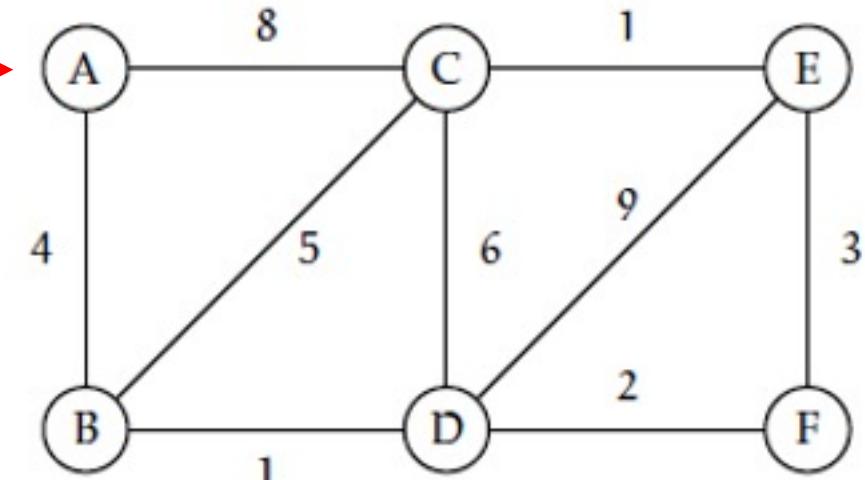
Steps of the algorithm:

Step	N'	$D(B), p(B)$	$D(C), p(C)$	$D(D), p(D)$	$D(E), p(E)$	$D(F), p(F)$
0	A	4,A	8,A	inf	inf	inf
1	A,B	4,A	8,A	5,B	inf	inf
2	A,B,D	4,A	8,A	5,B	14,D	7,D
3	A,B,D,F	4,A	8,A	5,B	10,F	7,D
4	A,B,D,F,C	4,A	8,A	5,B	9,C	7,D
5						

Exchanging table on node A.

Solution Method

Here, A, is the src



Loop

find w not in N' such that $D(w)$ is a minimum
add w to N'

update $D(v)$ for each neighbor v of w and not in N' :

$$D(v) = \min(D(v), D(w) + c(w, v))$$

/* new cost to v is either old cost to v or known
least path cost to w plus cost from w to v */

until $N' = N$

**1: If we choose E as w,
then $D(w)$ is
minimum**

**2: Iterate over all
neighbours of w
(which is E here) that
are not in the subset
 N' .**

Steps of the algorithm:

Step	N'	$D(B), p(B)$	$D(C), p(C)$	$D(D), p(D)$	$D(E), p(E)$	$D(F), p(F)$
0	A	4,A	8,A	inf	inf	inf
1	A,B	4,A	8,A	5,B	inf	inf
2	A,B,D	4,A	8,A	5,B	14,D	7,D
3	A,B,D,F	4,A	8,A	5,B	10,F	7,D
4	A,B,D,F,C	4,A	8,A	5,B	9,C	7,D
5	A,B,D,F,C,E	4,A	8,A	5,B	9,C	7,D

Exampling table on node A.

Solution Method

- To arrive at a final solution, we see that if we stand at node **A** then:
 - Shortest path from **A** to **B** is by taking edge **(A,B)** and here the total cost is 4
 - Shortest path from **A** to **C** is by taking edge **(A,C)** and here the total cost is cost 8
 - Shortest path from **A** to **D** is by taking edge **(A,B)** and here the total cost is cost 5
 - Shortest path from **A** to **E** is by taking edge **(A,C)** and here the total cost is cost 9
 - Shortest path from **A** to **F** is of cost 7. We see that the previous node on this path is **D** and we know that the shortest path to **D** is by taking edge **(A,B)**. Hence shortest path from **A** to **F** is by taking edge **(A,B)**.

Forwarding table on node **A**:

Destination node	Edge
B	(A,B)
C	(A,C)
D	(A,B)
E	(A,C)
F	(A,B)

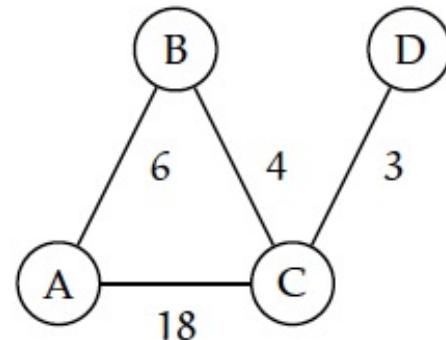


Remember that 'edge' here means standing at **A** in which direction should you go such that you end up taking the shortest path to the destination node

Distance-Vector (DV) Routing Algorithm – Exam 2021-22

- De-centralized Routing Algorithm
- An iterative, asynchronous, and distributed algorithm
- Information propagates in the network
 - Hence, we need to iterate, until information has propagated

Consider the network topology outlined in the graph below



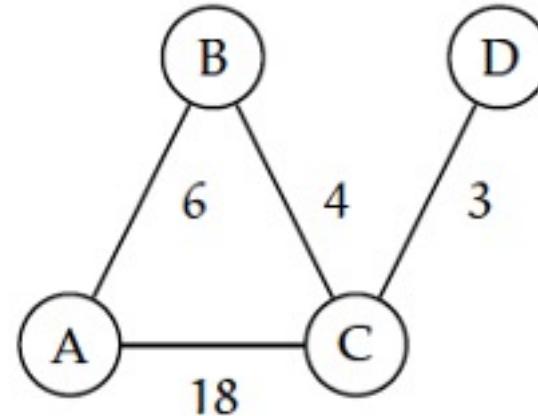
Question 3.3.1: Describe how to apply the Distance Vector Algorithm for Node A in the above diagram, so that it gets a complete routing table to all other Nodes. You do not need to describe each individual step on each Node in detail, but should be able to provide a generalised description that could apply to any Node. For this answer you can assume no connection costs change.

Node A table				
	A	B	D	D
A				
B				
C				
D				

Node B table				
	A	B	D	D
A				
B				
C				
D				

Node C table				
	A	B	D	D
A				
B				
C				
D				

Node D table				
	A	B	D	D
A				
B				
C				
D				



Initialization:

```

for all destinations y in N:
     $D_x(y) = c(x, y)$  /* if y is not a neighbor then  $c(x, y) = \infty$  */
for each neighbor w
     $D_w(y) = ?$  for all destinations y in N
for each neighbor w
    send distance vector  $\mathbf{D}_x = [D_x(y) : y \text{ in } N]$  to w
  
```

- **Initialize distance vector for each node with distance to reachable neighbours**
- **Distance to self is simply 0**
- **If a neighbour is unreachable, initialize to infinity**

Node A table

	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node B table

	A	B	C	D
A				
B				
C				
D				

Node C table

	A	B	C	D
A				
B				
C				
D				

Node D table

	A	B	C	D
A				
B				
C				
D				

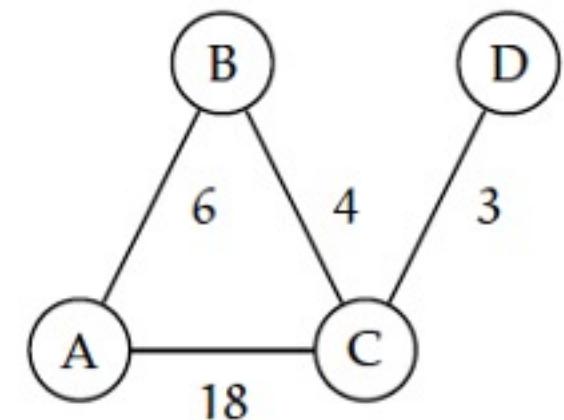
Initializing distance vector for A

- From A we can immediately reach B (cost 6) and C (cost 18)
- D is not immediately reachable, and hence initialized to inf (Abbreviated IF)
- At this point we do not know what the Distance Vectors of A's neighbours are, so these are simply initialized to IF
- Remember that the algorithm is run for each node asynchronously
 - Therefore this procedure has to be repeated for both A,B,C and D

**Initialization:**

```

for all destinations y in N:
   $D_x(y) = c(x,y)/*$  if y is not a neighbor then  $c(x,y) = \infty */$ 
  for each neighbor w
     $D_w(y) = ?$  for all destinations y in N
  for each neighbor w
    send distance vector  $D_x = [D_x(y) : y \in N]$  to w
  
```



- Initialize distance vector for each node with distance to reachable neighbours
- Distance to self is simply 0
- If a neighbour is unreachable, initialize to infinity

Node A table (Init)

	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node B table (Init)

	A	B	C	D
A	IF	IF	IF	IF
B	6	0	4	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node C table (Init)

	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node D table (Init)

	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	3	0

We initialize B, D and C in the same manner

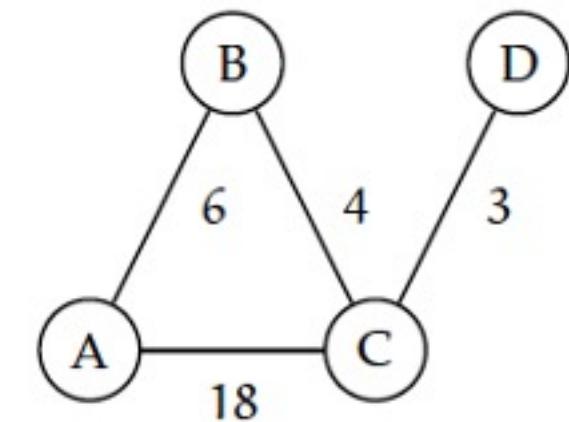
- From B we can reach A (cost 6) and C (cost 4)
- From C we can reach A (cost 18), B (cost 4) and D (cost 3)
- From D we can reach C (cost 3)

Initialization:

```

for all destinations y in N:
     $D_x(y) = c(x,y)/*$  if y is not a neighbor then  $c(x,y) = \infty */$ 
for each neighbor w
     $D_w(y) = ?$  for all destinations y in N
for each neighbor w
    send distance vector  $D_x = [D_x(y) : y \in N]$  to w

```



- Initialize distance vector for each node with distance to reachable neighbours**
- Distance to self is simply 0**
- If a neighbour is unreachable, initialize to infinity**

Node A table (Init)				
	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

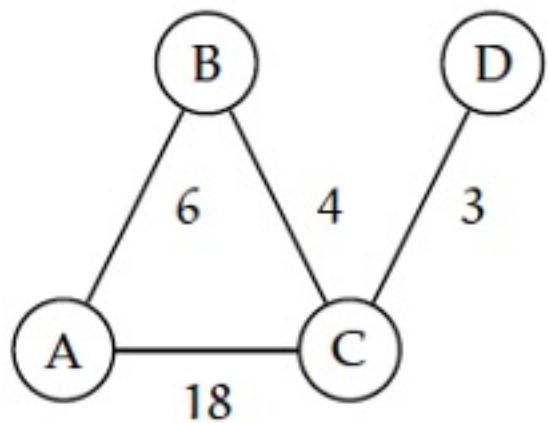
Node B table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	6	0	4	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node C table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node D table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	3	0

Send distance vector to neighbours

- A informs B & C
- B informs A & C
- C informs A, B & D
- D informs C



Initialization:

```

for all destinations y in N:
   $D_x(y) = c(x, y) /* if y is not a neighbor then  $c(x, y) = \infty */$ 
for each neighbor w
   $D_w(y) = ?$  for all destinations y in N
for each neighbor w
  send distance vector  $D_x = [D_x(y) : y \in N]$  to w$ 
```

Node A table (Init)				
	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node B table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	6	0	4	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node C table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node D table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	3	0

Node A table (Iter 1)				
	A	B	C	D
A	0	6	18	IF
B	6	0	4	IF
C	18	4	0	3
D	IF	IF	IF	IF

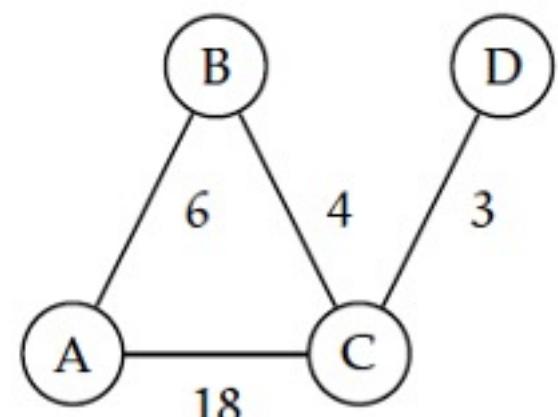
Node B table (Iter 1)				
	A	B	C	D
A	0	6	18	IF
B	6	0	4	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 1)				
	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	3	0

Node D table (Iter 1)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	3	0

Now DVs have been sent to neighbours and we enter the loop

- Both A,B,C and D have indeed received a DV from some neighbour, so therefore we have to execute the **for loop** on both A,B,C and D



```

loop
  wait (until I see a link cost change to some neighbor w or
        until I receive a distance vector from some neighbor w)

    for each y in N:
       $D_x(y) = \min_v \{ c(x, v) + D_v(y) \}$ 

    if  $D_x(y)$  changed for any destination y
      send distance vector  $D_x = [D_x(y) : y \in N]$  to all neighbors

  forever

```

Node A table (Init)				
	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node B table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	6	0	4	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node C table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node D table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	3	0

Node A table (Iter 1)				
	A	B	C	D
A	0	6	18	IF
B	6	0	4	IF
C	18	4	0	3
D	IF	IF	IF	IF

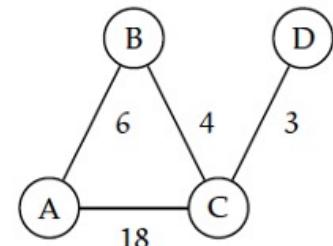
Node B table (Iter 1)				
	A	B	C	D
A	0	6	18	IF
B	6	0	4	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 1)				
	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	3	0

Node D table (Iter 1)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	3	0

We execute the for loop for Node A table

- Here we check if there is a shorter path from A to B
- $D_A(B) = \min_v\{c(A, v) + D_v(B)\}$
= $\min\{c(A, B) + D_B(B), c(A, C) + D_C(B)\} = \min\{6 + 0, 18 + 4\} = 6$
- In the above, we are essentially checking if we by going through one of our neighbours can obtain a shorter path to B
- $D_A(B)$ did **not** change as the shortest path from A to B was already 6
- We also check if there is a shorter path from A to C and D:



```

loop
  wait (until I see a link cost change to some neighbor w or
        until I receive a distance vector from some neighbor w)

    for each y in N:
       $D_x(y) = \min_v\{c(x, v) + D_v(y)\}$ 

    if  $D_x(y)$  changed for any destination y
      send distance vector  $D_x = [D_x(y) : y \in N]$  to all neighbors

  forever

```

Node A table (Init)				
	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node B table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	6	0	4	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node C table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node D table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	3	0

Node A table (Iter 1)				
	A	B	C	D
A	0	6	10	21
B	6	0	4	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node B table (Iter 1)				
	A	B	C	D
A	0	6	18	IF
B	6	0	4	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 1)				
	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	3	0

Node D table (Iter 1)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	3	0

We execute the for loop for Node A table

- **Checking if there is a shorter path from A to C**
- $D_A(C) = \min_v\{c(A, v) + D_v(C)\}$
= $\min\{c(A, B) + D_B(C), c(A, C) + D_C(C)\} = \min\{6 + 4, 18 + 0\} = 10$
- Great, we saw that we can go from A to C via B at cost 10!
- **Checking if there is a shorter path from A to D**
- $D_A(D) = \min_v\{c(A, v) + D_v(D)\}$
= $\min\{c(A, B) + D_B(D), c(A, C) + D_C(D)\} = \min\{6 + 4, 18 + 3\} = 21$
- Great, we can go from A to D at cost 21!
- As there were changes, we **send our updated DV** to our neighbors

```

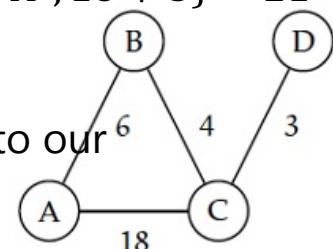
loop
    wait (until I see a link cost change to some neighbor w or
          until I receive a distance vector from some neighbor w)

    for each y in N:
         $D_x(y) = \min_v\{c(x, v) + D_v(y)\}$ 

    if  $D_x(y)$  changed for any destination y
        send distance vector  $D_x = [D_x(y) : y \in N]$  to all neighbors

```

forever



Node A table (Init)				
	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node B table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	6	0	4	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node C table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node D table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	3	0

Node A table (Iter 1)				
	A	B	C	D
A	0	6	10	21
B	6	0	4	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node B table (Iter 1)				
	A	B	C	D
A	0	6	18	IF
B	6	0	4	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 1)				
	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	3	0

Node D table (Iter 1)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	3	0

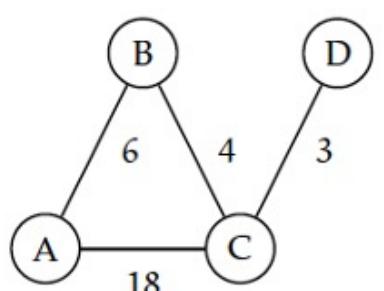
Node A table (Iter 2)				
	A	B	C	D
A	0	6	10	21
B	6	0	4	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node B table (Iter 2)				
	A	B	C	D
A	0	6	10	21
B	6	0	4	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 2)				
	A	B	C	D
A	A	0	6	10
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	3	0

Node D table (Iter 2)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	3	0

But remember! The algorithm also needs to be run for the B, C and D tables!



Node A table (Init)

	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node A table (Iter 1)

	A	B	C	D
A	0	6	18	IF
B	6	0	4	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node B table (Init)

	A	B	C	D
A	IF	IF	IF	IF
B	6	0	4	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node B table (Iter 1)

	A	B	C	D
A	0	6	18	IF
B	6	0	4	7
C	18	4	0	3
D	IF	IF	IF	IF

Node C table (Init)

	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 1)

	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	3	0

Node D table (Init)

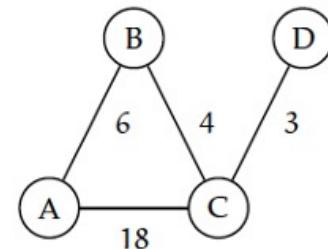
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	3	0

Node D table (Iter 1)

	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	3	0

Remember that we also have to run the algorithm for the B,C and D tables:

- **Checking if there is a shorter path from B to A**
 - There is not
- **Checking if there is a shorter path from B to C**
 - There is not
- **Checking if there is a shorter path from B to D**
 - There is: By going via C, the cost is 7
- Hence, we need to send our DV to neighbors



Node A table (Init)				
	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node B table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	6	0	4	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node C table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node D table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	3	0

Node A table (Iter 1)				
	A	B	C	D
A	0	6	10	21
B	6	0	4	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node B table (Iter 1)				
	A	B	C	D
A	0	6	18	IF
B	6	0	4	7
C	18	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 1)				
	A	B	C	D
A	0	6	18	IF
B	6	0	4	IF
C	18	4	0	3
D	IF	IF	3	0

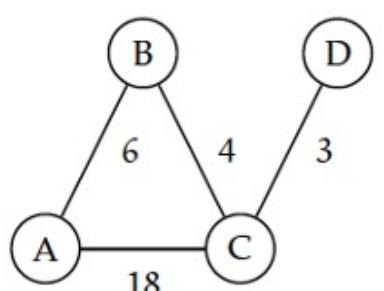
Node D table (Iter 1)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	3	0

Node A table (Iter 2)				
	A	B	C	D
A	0	6	10	21
B	6	0	4	7
C	18	4	0	3
D	IF	IF	IF	IF

Node B table (Iter 2)				
	A	B	C	D
A	0	6	10	21
B	6	0	4	7
C	18	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 2)				
	A	B	C	D
A	A	0	6	10
B	6	0	4	7
C	18	4	0	3
D	IF	IF	3	0

Node D table (Iter 2)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	3	0



Node A table (Init)

	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node B table (Init)

	A	B	C	D
A	IF	IF	IF	IF
B	6	0	4	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node C table (Init)

	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node D table (Init)

	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	3	0

Red arrows indicate transitions from the initial tables to the first iteration tables for each node.

Node A table (Iter 1)

	A	B	C	D
A	0	6	18	IF
B	6	0	4	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node B table (Iter 1)

	A	B	C	D
A	0	6	18	IF
B	6	0	4	7
C	18	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 1)

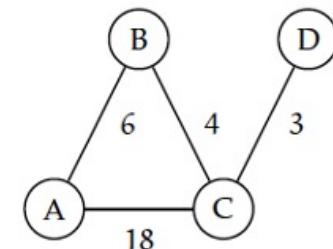
	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	10	4	0	3
D	IF	IF	3	0

Node D table (Iter 1)

	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	3	0

Remember that we also have to run the algorithm for the B,C and D tables:

- **Checking if there is a shorter path from C to A**
 - There is, by going via B, we get cost $4 + 6 = 10$
 - Hence, we need to send our DV to neighbors
- **Checking if there is a shorter path from C to B**
 - There is not
- **Checking if there is a shorter path from C to D**
 - There is not



Node A table (Init)				
	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node B table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	6	0	4	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node C table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node D table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	3	0

Node A table (Iter 1)				
	A	B	C	D
A	0	6	10	21
B	6	0	4	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node B table (Iter 1)				
	A	B	C	D
A	0	6	18	IF
B	6	0	4	7
C	18	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 1)				
	A	B	C	D
A	0	6	18	IF
B	6	0	4	IF
C	10	4	0	3
D	IF	IF	3	0

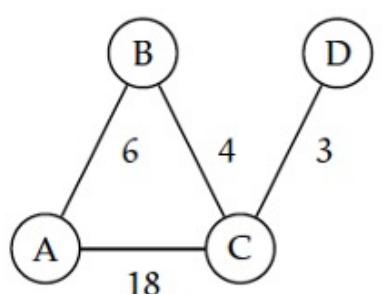
Node D table (Iter 1)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	3	0

Node A table (Iter 2)				
	A	B	C	D
A	0	6	10	21
B	6	0	4	7
C	10	4	0	3
D	IF	IF	IF	IF

Node B table (Iter 2)				
	A	B	C	D
A	0	6	10	21
B	6	0	4	7
C	10	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 2)				
	A	B	C	D
A	A	0	6	10
B	6	0	4	7
C	10	4	0	3
D	IF	IF	3	0

Node D table (Iter 2)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	10	4	0	3
D	IF	IF	3	0



Node A table (Init)

	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node A table (Iter 1)

	A	B	C	D
A	0	6	18	IF
B	6	0	4	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node B table (Init)

	A	B	C	D
A	IF	IF	IF	IF
B	6	0	4	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node B table (Iter 1)

	A	B	C	D
A	0	6	18	IF
B	6	0	4	7
C	18	4	0	3
D	IF	IF	IF	IF

Node C table (Init)

	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 1)

	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	10	4	0	3
D	IF	IF	3	0

Node D table (Init)

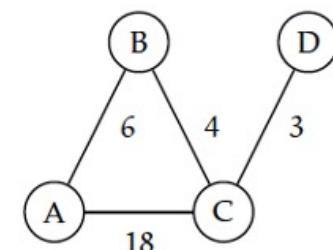
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	3	0

Node D table (Iter 1)

	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	21	7	3	0

Remember that we also have to run the algorithm for the B,C and D tables:

- **Checking if there is a shorter path from D to A**
 - There is, by going from D to A via C ($3+18 = 21$)
- **Checking if there is a shorter path from D to B**
 - There is by going via C ($3+4 = 7$)
- **Checking if there is a shorter path from D to C**
 - There is not
- Once again, update DV and send to neighbors



Node A table (Init)				
	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node B table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	6	0	4	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node C table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node D table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	3	0

Node A table (Iter 1)				
	A	B	C	D
A	0	6	10	21
B	6	0	4	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node B table (Iter 1)				
	A	B	C	D
A	0	6	18	IF
B	6	0	4	7
C	18	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 1)				
	A	B	C	D
A	0	6	18	IF
B	6	0	4	IF
C	10	4	0	3
D	IF	IF	3	0

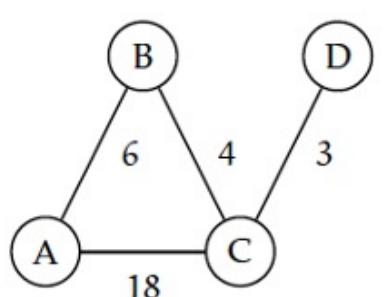
Node D table (Iter 1)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	21	7	3	0

Node A table (Iter 2)				
	A	B	C	D
A	0	6	10	21
B	6	0	4	7
C	10	4	0	3
D	IF	IF	IF	IF

Node B table (Iter 2)				
	A	B	C	D
A	0	6	10	21
B	6	0	4	7
C	10	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 2)				
	A	B	C	D
A	A	0	6	10
B	6	0	4	7
C	10	4	0	3
D	21	7	3	0

Node D table (Iter 2)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	10	4	0	3
D	21	7	3	0



Get ready for yet another
iteration...

Node A table (Init)				
	A	B	C	D
A	0	6	18	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node B table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	6	0	4	IF
C	IF	IF	IF	IF
D	IF	IF	IF	IF

Node C table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node D table (Init)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	IF	IF	IF	IF
D	IF	IF	3	0

Node A table (Iter 1)				
	A	B	C	D
A	0	6	10	21
B	6	0	4	IF
C	18	4	0	3
D	IF	IF	IF	IF

Node B table (Iter 1)				
	A	B	C	D
A	0	6	18	IF
B	6	0	4	7
C	18	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 1)				
	A	B	C	D
A	0	6	18	IF
B	6	0	4	IF
C	10	4	0	3
D	IF	IF	IF	0

Node D table (Iter 1)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	18	4	0	3
D	21	7	3	0

Node A table (Iter 2)				
	A	B	C	D
A	0	6	10	21
B	6	0	4	7
C	10	4	0	3
D	IF	IF	IF	IF

Node B table (Iter 2)				
	A	B	C	D
A	0	6	10	21
B	6	0	4	7
C	10	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 2)				
	A	B	C	D
A	A	0	6	10
B	6	0	4	7
C	10	4	0	3
D	21	7	3	0

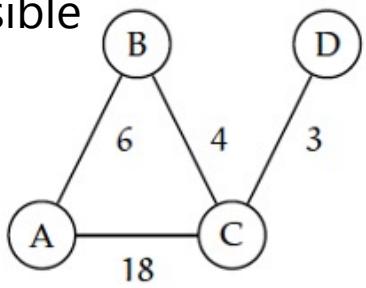
Node D table (Iter 2)				
	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	10	4	0	3
D	21	7	3	0

For Node A table:

We can get to D via B, hence A to D cost $6+7 = 13$
Otherwise no optimizations

For Node B table:

No optimizations possible



For Node C table:

No optimizations possible

For Node D table:

We can get to A via C, hence D to A cost $3+10 = 13$

Node A table (Iter 2)

	A	B	C	D
A	0	6	10	13
B	6	0	4	7
C	10	4	0	3
D	IF	IF	IF	IF

Node B table (Iter 2)

	A	B	C	D
A	0	6	10	21
B	6	0	4	7
C	10	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 2)

	A	B	C	D
A	A	0	6	10
B	6	0	4	7
C	10	4	0	3
D	21	7	3	0

Node D table (Iter 2)

	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	10	4	0	3
D	13	7	3	0

For Node A table:

We can get to D via B, hence A to D cost $6+7 = 13$
Otherwise no optimizations

Node A table (Iter 3)

	A	B	C	D
A	0	6	10	13
B	6	0	4	7
C	10	4	0	3
D	IF	IF	IF	IF

Node B table (Iter 3)

	A	B	C	D
A	0	6	10	13
B	6	0	4	7
C	10	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 3)

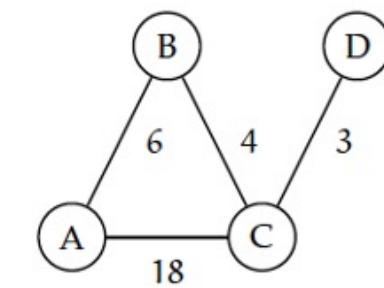
	A	B	C	D
A	0	6	10	13
B	6	0	4	7
C	10	4	0	3
D	13	7	3	0

Node D table (Iter 3)

	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	10	4	0	3
D	13	7	3	0

For Node D table:

We can get to A via C, hence D to A cost $3+10 = 13$



Node A table (Iter 2)

	A	B	C	D
A	0	6	10	13
B	6	0	4	7
C	10	4	0	3
D	IF	IF	IF	IF

Node B table (Iter 2)

	A	B	C	D
A	0	6	10	21
B	6	0	4	7
C	10	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 2)

	A	B	C	D
A	A	0	6	10
B	6	0	4	7
C	10	4	0	3
D	21	7	3	0

Node D table (Iter 2)

	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	10	4	0	3
D	13	7	3	0

For Node A table:

We can get to D via B, hence A to D cost $6+7 = 13$
Otherwise no optimizations

Node A table (Iter 3)

	A	B	C	D
A	0	6	10	13
B	6	0	4	7
C	10	4	0	3
D	IF	IF	IF	IF

Node B table (Iter 3)

	A	B	C	D
A	0	6	10	13
B	6	0	4	7
C	10	4	0	3
D	IF	IF	IF	IF

Node C table (Iter 3)

	A	B	C	D
A	0	6	10	13
B	6	0	4	7
C	10	4	0	3
D	13	7	3	0

Node D table (Iter 3)

	A	B	C	D
A	IF	IF	IF	IF
B	IF	IF	IF	IF
C	10	4	0	3
D	13	7	3	0

For Node A table:

No optimizations possible

For Node C table:

No optimizations possible

For Node D table:

We can get to A via C, hence D to A cost $3+10 = 13$

For Node C table:

No optimizations possible

For Node D table:

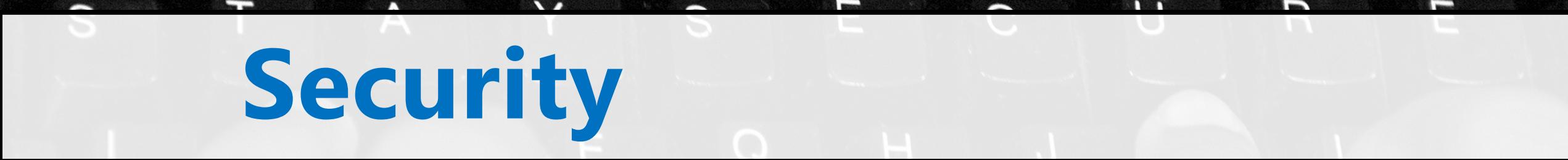
No optimizations possible

As there are no more optimizations, changes have propagated and we stop

Final table (Composed of DVs from A,B,C and D)				
	A	B	C	D
A	0	6	10	13
B	6	0	4	7
C	10	4	0	3
D	13	7	3	0



Security



Good things to remember about Security

- Questions within security are often open-ended and you need to know the **theory** well
- Here, I will briefly give you a brush up on Confidentiality and Integrity
- **Learning goal:**
 - **Kompetencer i at** "Ræsonnere omkring enkle sikkerhedsegenskaber for et givet system"
 - **Viden om** "Basal anvendt kryptografi, og operationelle metoder til at sikre styresystemsmekanismer, netværksprotokoller og applikationer."
 - As an example, that does **NOT** mean being able to generate arbitrary RSA keys by hand, but rather understand when and why RSA can be beneficial and when it is not.

Confidentiality

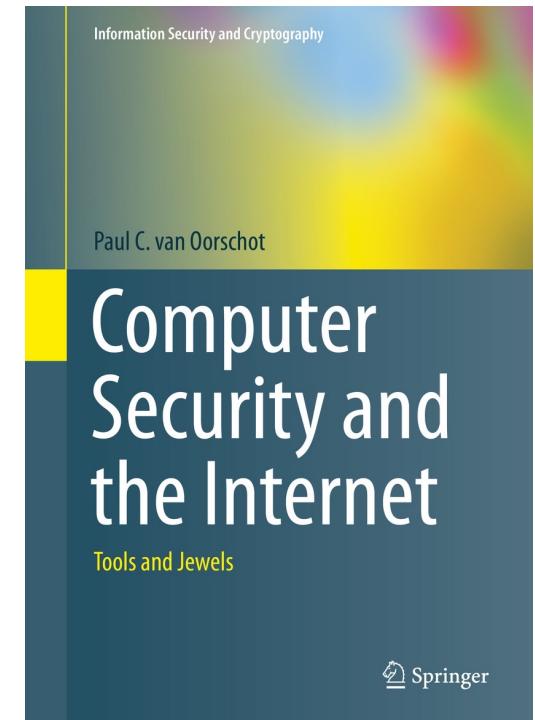
- A typical aim in IT-security is to establish confidential communication between Alice and Bob
 - Use encryption to ensure **confidentiality**
 - RSA is asymmetric (private,public), i.e. keypair
 - However encrypting each message through RSA is computationally expensive
so typically RSA is used to encrypt a symmetric session key
 - Further communication is then encrypted using a session key
 - This is essentially what happens in the **TLS**, remember **TLS** is what makes **HTTP -> HTTPS**
 - **Note:** Just because something is confidential, it does not guarantee protection against e.g. replay attacks.
 - Also in Databases, passwords are not stored in plaintext
 - Instead, a Hash of (Password+Salt) is stored – the salt helps guard against Rainbow Tables and dictionary attacks

Integrity

- Another aim is for Alice and Bob to ensure that messages are not altered while in transit
 - Therefore, a **checksum** can be used, just as we did in A3 and A4
 - A checksum is a hash of the message, using a **cryptographic hashing algorithm** such as SHA-256, SHA-1 or MD5. If this is sent along the message, the receiver can generate his own hash and compare just as in A3 and A4.
- **Integrity** does not by itself guarantee the identity of the sender. For that some kind of **authentication** is required
 - Hash(message + secret) is called a **Message Authentication Code (MAC)**
 - If only Alice and Bob know the **secret**, authentication proves identity of sender
 - Another tool is the use of **Digital Signatures** where a **private key** is used to sign (encrypt) the hash of a message. Receiver can use senders **public key** to decrypt the hash of the message and compare with his own hash of the message

If you are interested in IT-Security

- CompSys gives only a very brief introduction to IT-Security
- If you are interested in the field, I recommend taking the elective Bsc course IT-Security in blok 1 → See <https://kurser.ku.dk/course/ndaa09025u/>
- The course gives you a theoretical foundation and you get to try out some practical things through the assignments



Postludium

- Thank you for now, it has been a pleasure teaching you
- **Remember** you have come a long way! In week 1 of compSys, 99% of you didn't know what a pointer or malloc was. Also, **C** was just a letter in the alphabet.
- Good luck at the compSys exam (and also good luck at MAD/SS, etc.)
- I will see you around DIKU
- And for those of you taking the course **Participatory design af informationssystemer**, I will see you there ☺

**And remember:
Even though it
may be tempting,
don't answer this
at the exam...**

2 Short Answer Questions

11. [10 points] Name and describe the five key phases of software development.

- ~~1. denial~~
~~2. bargaining~~
~~3. Anger~~
~~4. depression~~
~~5. acceptance~~