

# Advanced Programming

## Testing

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- Abstract data types and modules
  - Specifications
  - Testing fundamentals
  - Specification-driven testing
  - Property-based testing by input partitioning
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## Goals

Week 5 (last week):

- concurrency;
- channels and threads;
- thread-based asynchronous programming;
- client-server programming.

Week 6 (this week):

- specification-driven testing;
- randomized property based testing;
- programming with Quickcheck.

## Abstract data types and modules revisited

- An *abstract data type (ADT)* is an abstract type together with operations involving the type and a precise description of what the operations do. An ADT consists of
  - an *interface* containing the *declarations* of the type and its operations (functions),
  - their *specification* (what they do, which properties they have); and
  - *examples* illustrating correct and incorrect use of the operations.
- An *implementation* of an ADT contains
  - the *definitions* (source code) of the operations; and
  - *evidence* that it satisfies the ADT's specification (test design documentation and resulting test suite).
- An ADT typically has *multiple possible implementations*.
  - In particular, the implementation types may be different from each other.
  - Example:  $s \rightarrow (a, s)$  and  $\text{IORef } s \rightarrow \text{IO } a$  for state monads.
- *Module*: Structure with zero, one or more abstract or concrete type declarations.
  - ADT: Module with one abstract type declaration.

## Abstract data types

- Programming language support for ADTs and *modular abstraction*:
  - ML-family languages: Signatures (module types), structures (module implementations), functors (module-parameterized structures).
  - Haskell: Type classes, class instances, class constraints in functions.
  - Java-like languages: Interfaces and abstract classes, concrete classes, concrete classes with interface-parameterized constructor(s).

## Specifications

- *Specification*: Describes what a (collection of) function(s) does, not how it is implemented (its code).
  - Litmus test: Contains *everything* a user of a module *needs to know* and *nothing else*.
  - Specifications may be
    - inconsistent (contradictory): contain a contradiction;
    - incomplete (often the case): leave room for multiple implementations with different behaviors.
- *Implementation*: Code.
- *Correctness*: Evidence (e.g. test design method and resulting *test suite*) that code may satisfy its specification.
- *Requirements*: Informal description of desired functionality of a program component (what it should do and what it should not do).
  - Requirements may be unclear and leave room for interpretation.
  - A specification *resolves and documents* this ambiguity.

## Module interface specification as a contract

- Module interface specification: *Contract* between programmer *using* a module (user) and programmer *implementing* it (implementor).
  - User to implementor: “All I want to know needs to be written in the module specification. How I use your implementation is none of your business. Don’t make me look at your code.”
  - Implementor to user: “All you need to know is written in the module specification. How I implement it is none of your business. Don’t ask to look at my code. (I may change it at any time.)”
  - Implementor and user are *roles*: They *may* be the same person.
- What if specification and implementation are *inconsistent*, as evidenced by a *test failure*?
  - Change the code to meet the specification: Requires no negotiation with user(s).
  - Change the specification to meet the code: Requires negotiation with user(s) since *changes may propagate to whatever code they have written*.
  - Change both, code and specification: Natural part of *exploratory programming* during *requirements elicitation*.
  - Common problem: Most modules are *underspecified*. User tries out code and relies on unspecified behavior of implementation.

## Example: State monad specification

- How to (partially) specify an abstract data type? Typical methods:
  - Properties its functions must satisfy.
    - Usually *universal equational properties*
  - Reference implementation (executable specification, model implementation)
    - Observable results must be the same as those from a reference implementation
- Example: State monad
  - Monad laws plus

```
1 put s >> get           = return s
2 put s >> put t >> get  = put t >> get
3 get  >> get            = get
```

for all  $s, t$ .

- Coherence with purely functional implementation FState

## Example: Reference implementation by type conversion

- Idea: Provide conversion functions from/to reference type
- Example (simplified):

```

1  FState s a = s -> IO (a, s)    -- reference type
2  IState s a = IORef s -> IO a    -- new type
3
4  to :: FState s a -> IState s a
5  to fstate ref = do
6    s <- readIORef ref
7    (a, t) <- fstate s
8    writeIORef t ref
9    return a
10
11 from :: IState s a -> FState s a
12 from istate s = do
13   ref <- newIORef s
14   a <- istate ref
15   t <- readIORef ref
16   return (a, t)

```



## Example: Reference implementation by type conversion

- Idea: Check that new functions simulate reference functions
- Properties:

```

1 instance StateMonad FState where ... -- reference impl.
2 instance StateMonad IState where ... -- new impl.
3 from (to fstate) = fstate      :: FState s a
4 from get         = get         :: FState s s
5 from (put s)     = put s       :: FState s ()
6 from (return a)  = return a    :: FState s a
7 to fstate >>= f   = fstate >>= f :: FState s b

```

for all `fstate`, `s`, `a`, `f`.

- Note: Left-hand side operations from new state monad implementation; right-hand side from reference implementation.
- Problems:
  - What does `=` mean?
  - How to generate functions for testing?

## Testing

- *Testing* is (partial) verification by
  - executing code on inputs and checking whether the result is evidence of a specification violation;
  - *systematically* constructing the inputs to *maximize their likelihood of demonstrating a specification violation*.
- Terminology:
  - **Formal verification:** Mathematical proof that there exists *no* valid input that results in a specification violation.
  - **Testing:** The *systematic* discipline of finding “nasty” inputs that maximize the likelihood of finding specification violations.
    - Not being able to “break the code” is taken as evidence (not proof) that the code may be correct with respect to its given specification.
  - **Trying out, illustrating, exemplifying:** Applications of the code showing that it satisfies its specification for *some* inputs (if there is a specification) or that it does something useful (if there is no specification, only informal requirements).
- Don't use the term “test” unless you have *both* specification and code and you are *systematically* looking for specification violations.
  - Use “illustrate” or “exemplify” or “try out” instead.

## Specification-driven testing

- Recall: All forms of testing require *both* specification and code.
  - No specification, no testing.
- *Specification-driven function testing*: Analyze specification and systematically
  - find *valid* data with maximal likelihood to expose a specification violation of the code;
  - determine corresponding *expected* outputs *from the specification*.
    - Don't look at the code in this process.
- Test-driven development (TDD):
  - Develop the test suite *before* you have any code.
  - Continuously run the test suite to drive and check the code as you develop it.
- Specification-driven testing is also called *specification-based testing*, *external testing*, and *black-box testing*.
  - “Black-box testing” is not used here since much of questionable value is unfortunately written on the web on that.

## Properties

- A *logical statement* is a statement that is either true or false.
- A *(logical) property* of one or more functions is a true logical statement involving the function(s).
- A *universally quantified property* has the form

$$\forall x, y, \dots Q(x, y, \dots)$$

or

$$\forall x, y, \dots (P(x, y, \dots) \Rightarrow Q(x, y, \dots)).$$

In words: For all (valid)  $x, y, \dots$  the property  $Q(x, y, \dots)$  is true.

- Example:

$$\forall x \in \text{Int} ((x \geq 0) \Rightarrow \text{fib}(x) = fib(x))$$

where `fib` is some code and *fib* is the “true” mathematically defined Fibonacci function.

- A *(partial) formal specification* is often a conjunction of universally quantified properties.
- *Property-based testing*: Systematic design of *finite* subset of inputs aimed at falsifying a property.

## Input partitioning

- *Input partitioning* is a design method for constructing *test data* by inspecting the *properties to be checked*.
- Given a universally quantified property, partition its valid inputs into a *finite* set of partitions, that is into pairwise disjoint subsets whose union is the entire set of valid inputs.
- From each partition,
  - choose a *typical* element and one or more *boundary elements*;
  - determine the *expected* output for each element according to the function specification—*not* the code.
- Test suite: The chosen input/expected output pairs. The boundary elements increase the likelihood of finding a specification violation.
- Test execution: Run the function on the inputs in the test suite; check that it produces the corresponding expected outputs.
- Input partitioning is also called *equivalence partitioning* or *equivalence class partitioning*.

## Input partitioning: Numbers

- Generally useful partitioning of integral and floating-point numbers.
  - `int` (32-bit 2's-complement integers):
    - Partitions: Negative numbers, `{0}`, positive numbers.
    - Values: `-2147483648`, `-137`, `-1`, `0`, `1`, `2377`, `2147483647`.
  - `float` (64-bit IEEE 754 floating-point numbers):
    - Partitions: `{ Negative infinity }`, negative numbers, `{0}`, positive numbers, `{ positive infinity }`, `{ NaN }` ("not-a-number").
    - Values: `System.Double.NegativeInfinity`, `2.2250738585072014e-308`, `-1.0`; `-System.Double.Epsilon`, `0.0`, `System.Double.Epsilon`, `1.0`, `47771354343.98989988`, `1.7976931348623158e+308`, `System.Double.PositiveInfinity`, `System.Double.NaN`.
  - Note: F# syntax used!
- Note that most test values are boundary values.

## Summary: Programming = Specification + coding + verification

- *Programming* consists of

- 1 *Specification*: What is computed? Which functions, what are their *properties*?
- 2 *Coding*: How is it computed? How are the functions implemented (“code”)?
- 3 *Verification*: Which explicit *evidence* is there for the program code meeting its specification (correctness)?

- Verification always requires *both*, program code *and* specification.

- Program code may be correct with respect to one specification and incorrect with respect to another specification.

- Avoid statements such as “My code works”.

- With respect to which specification?

- Code is trivially correct with respect to whatever it happens to do—there is nothing to test; it does what it does. But what *does* it do that a code user can rely on?

- Does it *always* work and can you provide conclusive evidence for that, for *all* valid inputs and in *all* conceivable contexts, or just on *some* data that you happen to have tried out?

- Next time: Randomized property testing and QuickCheck.

## Summary

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- Specifications
- Testing fundamentals
- Specification-driven testing
- Property-based testing by input partitioning