



# **IT-Security (ITS) B1**

## **DIKU, E2024**



# **This is plain text**

Computing in the presence of an adversary



**This is not plain text**

Frpsxwlqj lq wkh suhvhqfh ri dq dgyhuvdub



**This is not plain text**

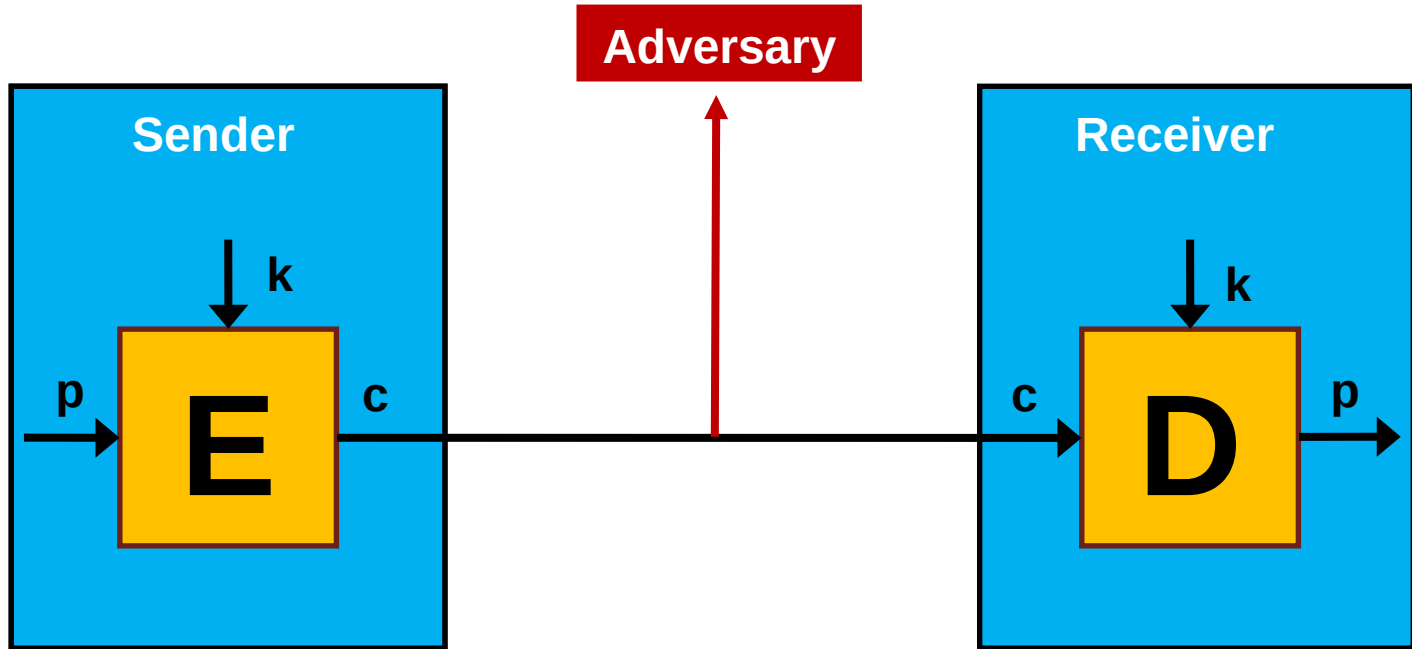
Q29tcHV0aW5nIGluIHRobzSBwcmVzZW5jZSBvZiBhbiBhZHZlcnNhcnkK



# This is not plain text

```
689b ef01 affa eb02 5618 7770 1c66 58ed  
139c 9020 8431 2ff0 e7af 0d41 b3d5 b4a6  
f222 90b3 f51a afd9 00fe e01d c509 69f4
```

# Cryptosystems





# Security goals

Confidentiality

Integrity

Authenticity

Non-repudiation



# Today's agenda

Part 1: Crypto building blocks

Part 2: More crypto building blocks

(Later: Real-world crypto protocols)





# Today's agenda

Cryptography defined

Cryptography from a historic perspective

Tools: Encryption, decryption, cryptographic hash functions, digital signatures,

Remember, cryptography is key, but hard to get right and not a panacea



# Crypto defined



# Today's agenda

Cryptology (from Ancient Greek: *kryptós* "hidden, secret"; and *graphein*, "to write"), is the practice and study of codes, both creating and breaking them

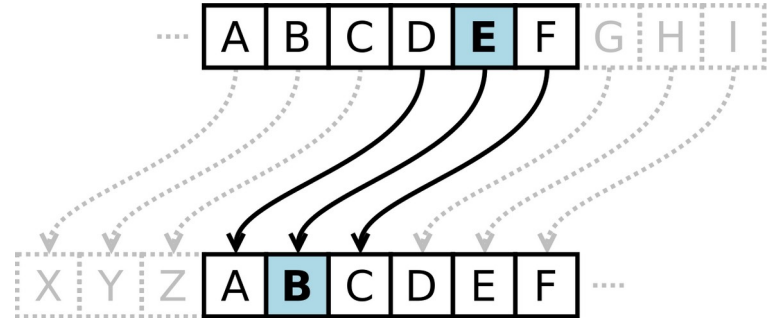
Cryptography is the art of creating codes

Cryptanalysis is the art of breaking codes



# **Crypto from a historic perspective**

# Cryptography influence world events



# Cryptography influence world events

**WESTERN UNION TELEGRAM**

GERMAN LEGATION  
MEXICO CITY

via Galveston JAN 8 9 1917

130 13042 13401 8501 115 3528 416 17214 6491 11310  
18147 18222 21560 10247 11518 23877 13605 3494 14936  
98092 5905 11311 10392 10371 0302 21290 5101 39095  
23571 17504 11295 18276 18101 0317 0228 17894 4473  
24224 22200 19452 21589 07893 5569 13918 8958 12137  
1333 4725 4458 5905 17106 13851 4458 17149 14471 0708  
13850 12224 0929 14961 7382 15857 07893 14218 36477  
5870 17553 87803 0870 5454 16102 15217 22601 17138  
21001 17388 7146 23638 18222 8719 14331 15021 23845  
3158 23552 22096 21604 4797 9497 22461 20855 4377  
23610 18140 22280 5905 13347 20420 39889 15732 50687  
0929 5275 18507 52282 1340 22049 13339 11285 22295  
10439 14814 4178 6992 8784 7632 7357 6926 52262 11287  
21100 21272 9340 9559 22464 15874 18502 18500 15857  
2188 5376 7381 98092 10127 13486 9350 9220 76036 14219  
5144 2831 17520 11347 17142 11264 7887 7762 15099 9110  
10482 97556 3509 3070

BEPASTOPFF.

Charge German Embassy.

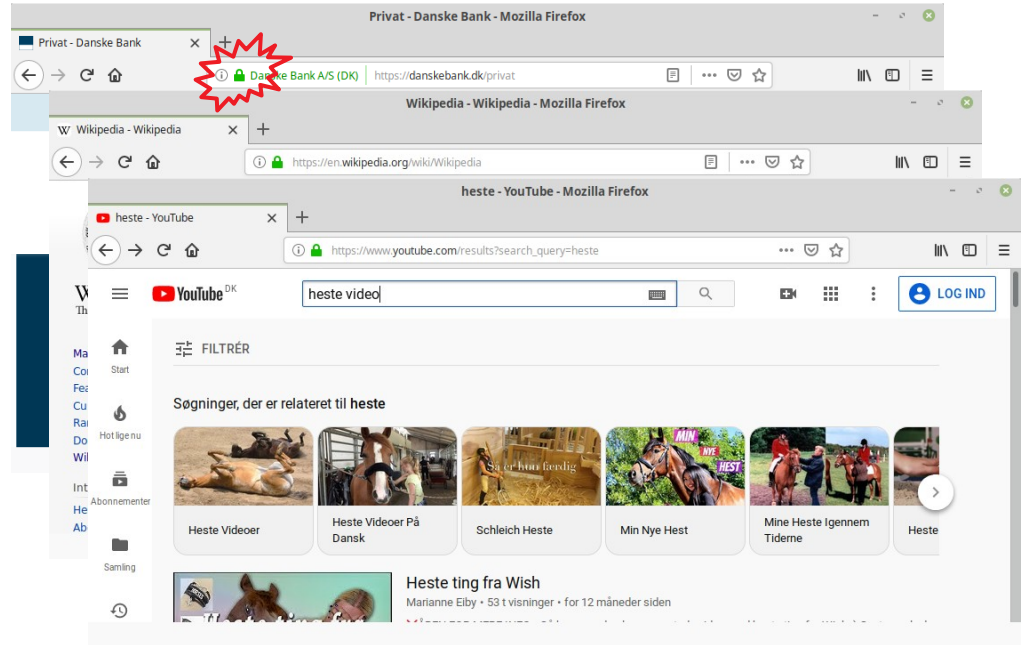


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# Cryptography influence world events



# Our goal: Secure online communication







# **Warm-up question**



# FileCrypt

“FileCrypt is a dynamic non-factor based quantum AI encryption hardware solution.

Developed by our cryptographic experts and hardwired into a tamper-resistant USB token.

Plug the token into your PC, start the program and encrypt the files you need to protect”

What problems do you see with this solution?



# Multiple concerns

#1: “Developed by our cryptographic experts”

Should we trust proprietary crypto over public peer-reviewed time-tested crypto?

#2: “Dynamic non-factor based quantum AI”

What does that mean? Are there any academic papers that discuss this concept?

#3: “Plug the token into your PC”

Can anyone do this? What if token is lost? Violates **Kerckhoffs’ Principle**

# Kerckhoffs' (2nd) Principle

The security of a cryptographic algorithm must rest solely in the secrecy of its **key**, not in the secrecy of the algorithm itself

Collaries:

- Assume attacker knows the algorithm
- Make it available for public analysis
- Protect the key!



Auguste Kerckhoffs  
(1835 – 1903)



# **Symmetric cryptosystems**



# **Symmetric cryptosystems**

## **Stream ciphers**

# One time pad

If  $k$  random,  $|k| \geq |p|$ , never reused, and kept secret, then it is impossible to decrypt or break without knowing the key (Shannon, 1949)

Key	0	1	0	1	1	1	0	0	1	0
Plaintext	1	1	0	0	0	1	1	0	0	0
Ciphertext	1	0	0	1	1	0	1	0	1	0

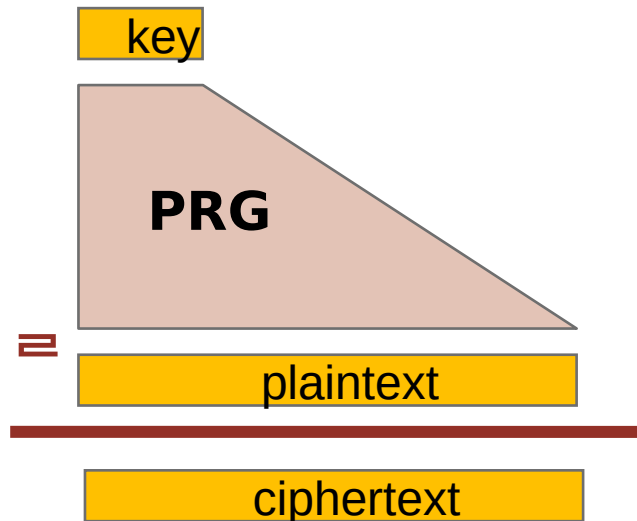
# Towards modern stream ciphers

Problem

OTP key as long as plaintext

Solution

Generate pseudo random keystream







# <sup>st</sup> 1 rule of stream ciphers

Never reuse key

$$C_1 \oplus P_1 = \text{PRG}(k)$$

$$C_2 \oplus P_2 = \text{PRG}(k)$$

$$C_1 = C_2 \quad \circ \quad P_1 = P_2$$

$$P_1 = P_2 \quad \circ \quad P_1, P_2$$



# Solution: Initialisation Vector (IV)

For each message

- Generate IV

- Mix  $k$  with IV

- Generate keystream  $\text{PRG}(k + \text{IV})$  and encrypt

- Send  $c$  and IV (in plaintext)

Change  $k$  before IVs run out

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# Stream ciphers in the wild



**https://**

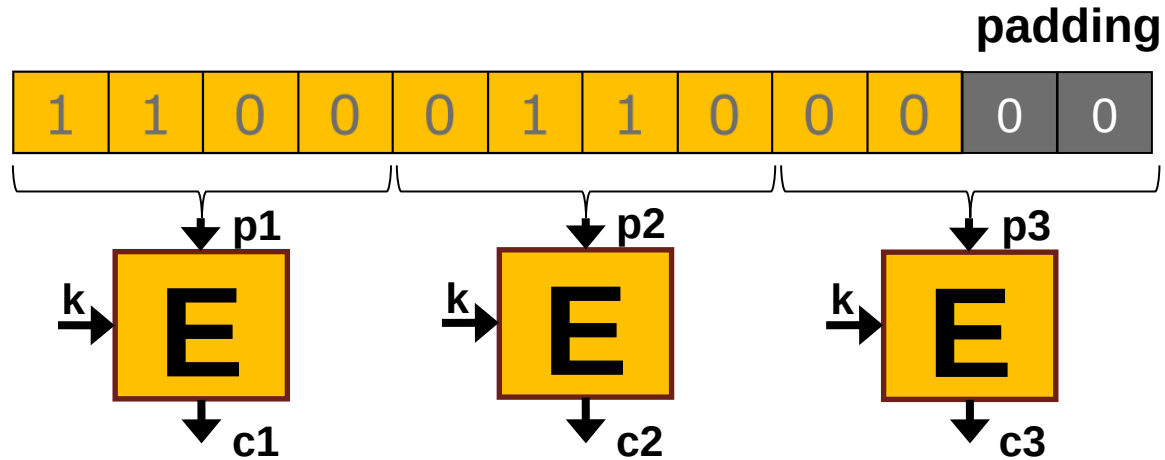




# Block ciphers

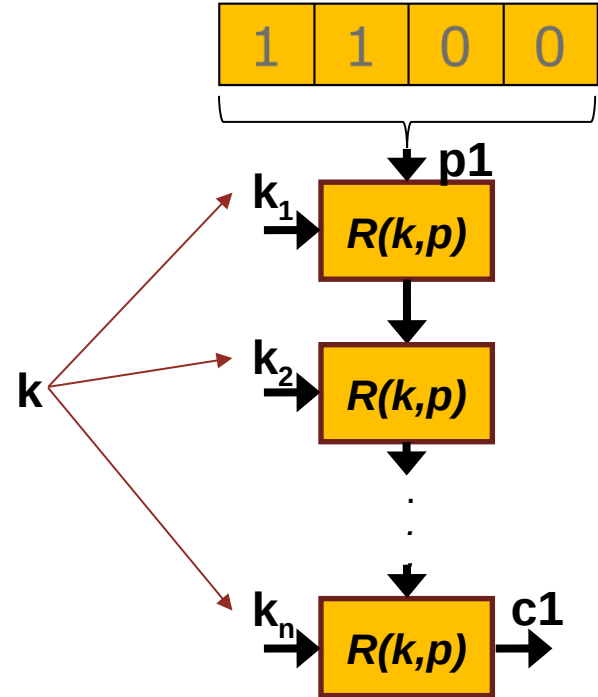
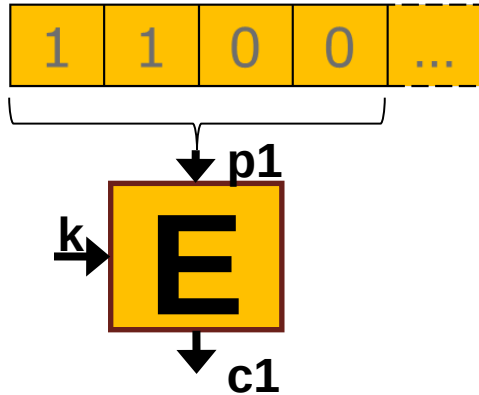
# Block ciphers

One block at a time – as opposed to one bit at a time



# One block at a time

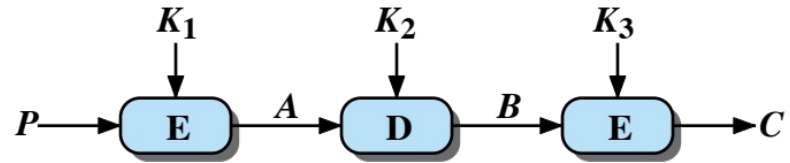
Blocks, rounds function, key schedule, iterations



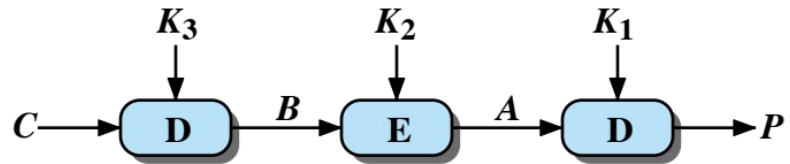
# DES

DES

Key 64, block 64, rounds 16



(a) Encryption



(b) Decryption

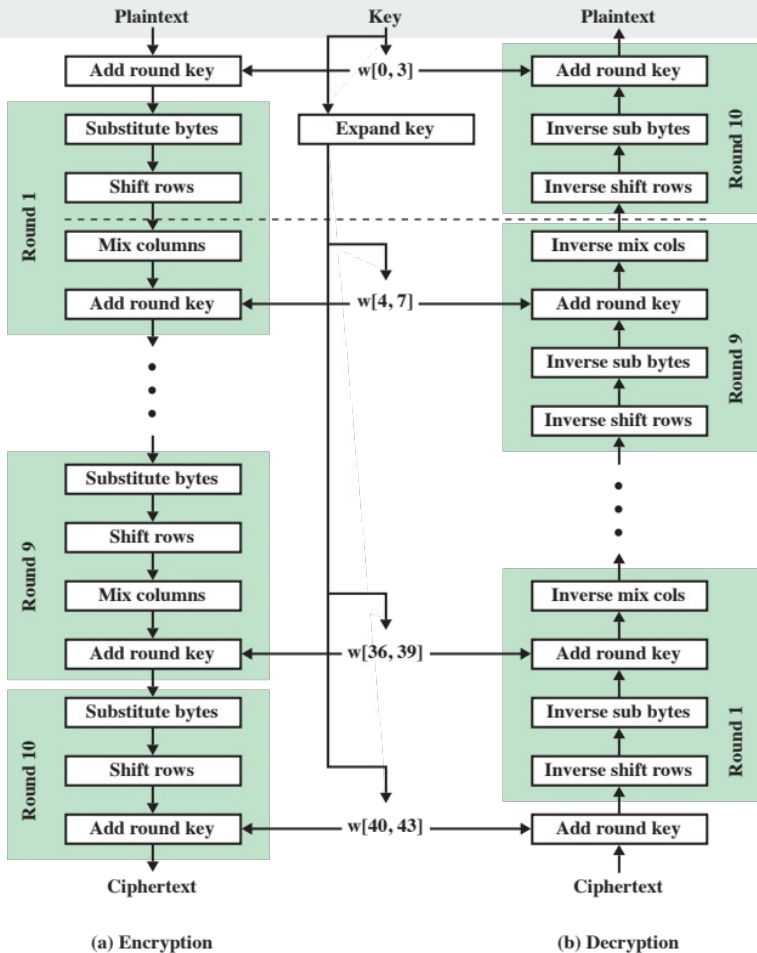
# AES

AES

Keys 128/192/256

Block 128

Rounds 10/12/14

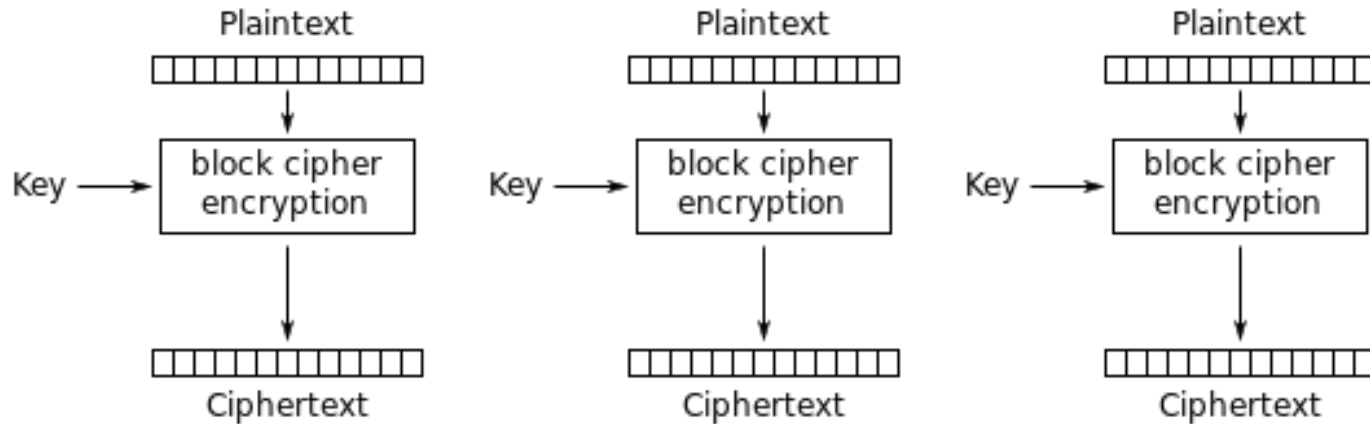






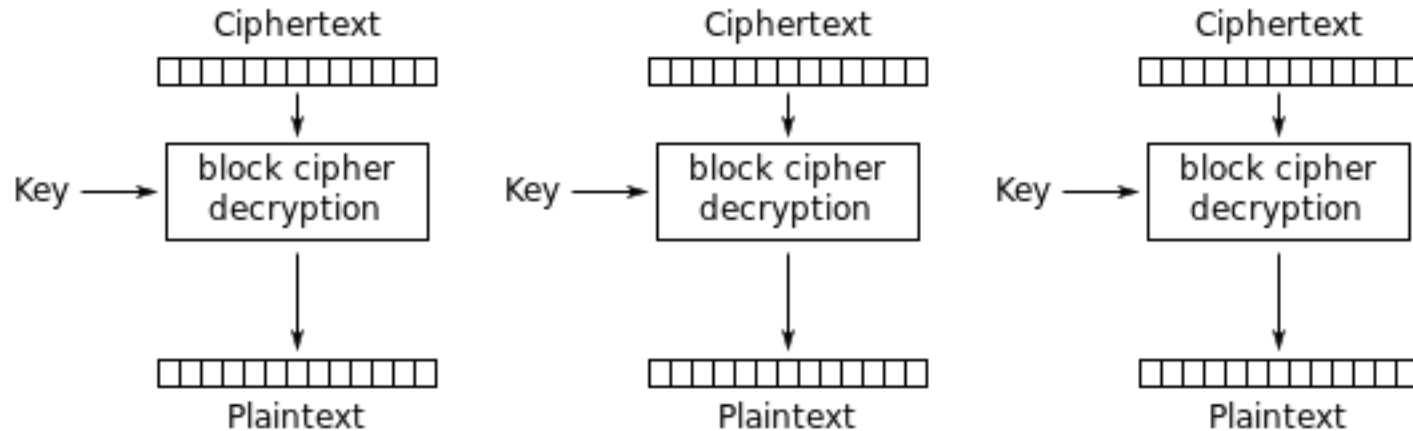
# Modes of operation

# Electronic Codebook (ECB)



Electronic Codebook (ECB) mode encryption

# ECB decryption



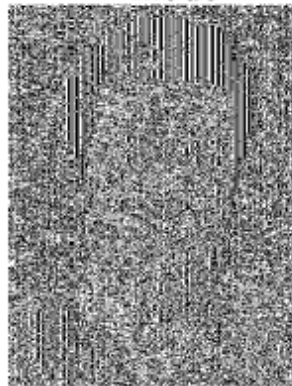
Electronic Codebook (ECB) mode decryption

**If  $p1 = p2$ , then  $c1 = c2$**

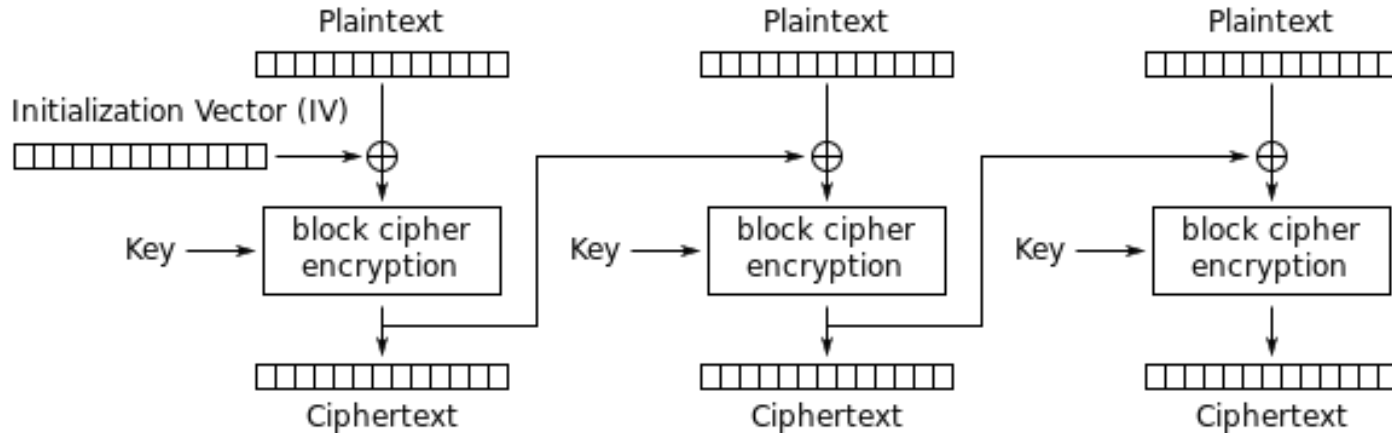
An example plaintext



Encrypted with AES in ECB mode

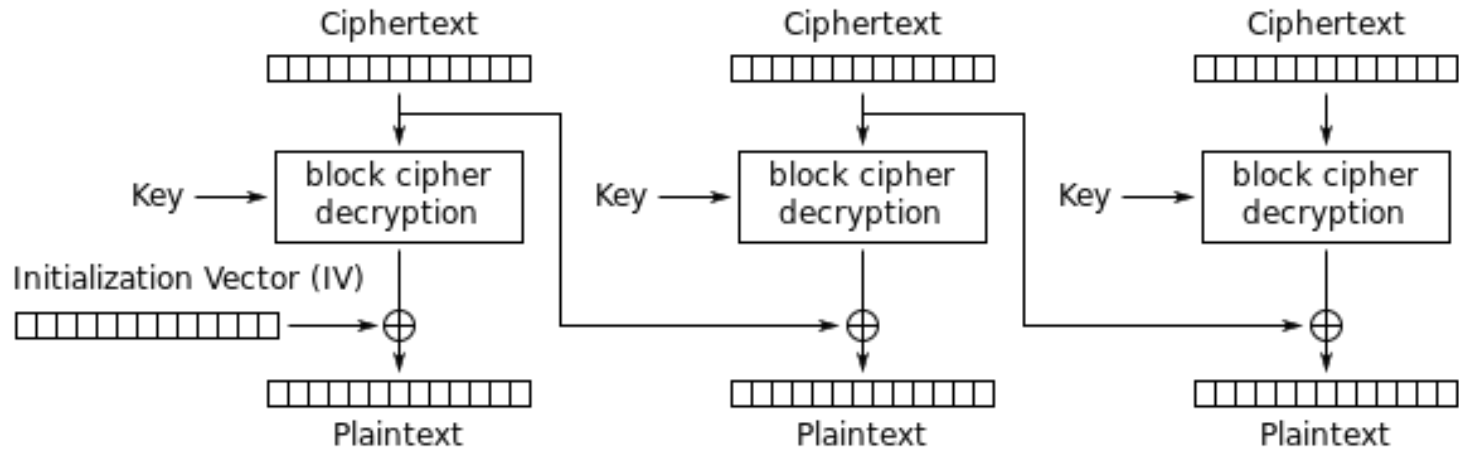


# Cipher Block Chaining



Cipher Block Chaining (CBC) mode encryption

# CBC decryption



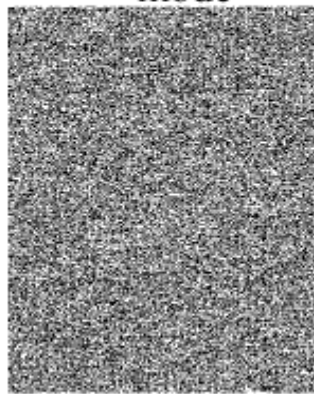
Cipher Block Chaining (CBC) mode decryption

# Better

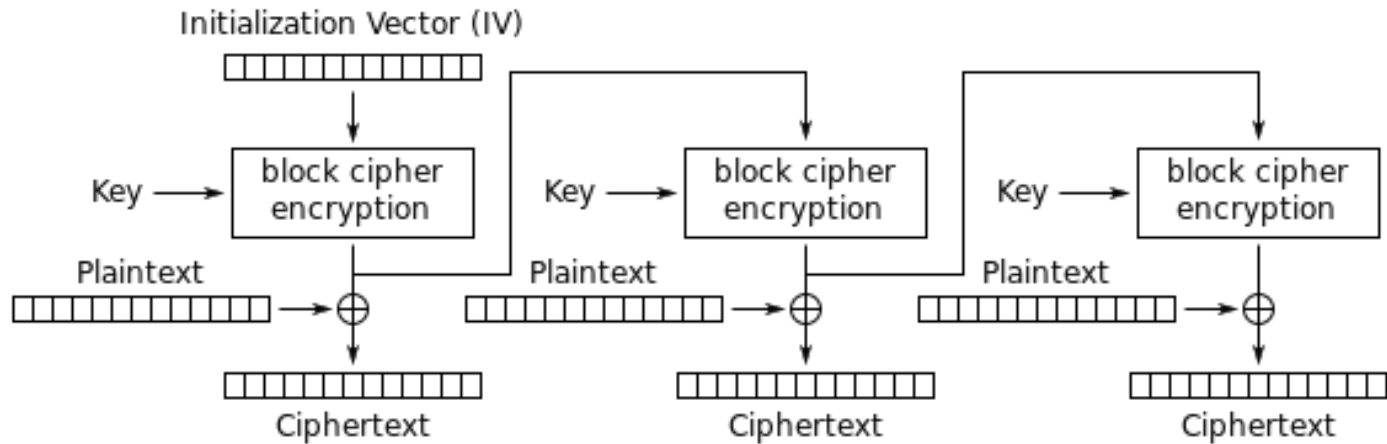
An example plaintext



Encrypted with AES in CBC  
mode



# Output Feedback



Output Feedback (OFB) mode encryption





# Security goals revisited

“Susceptibility to malicious insertions and modifications. Because each symbol is separately enciphered, an active interceptor who has broken the code can splice together pieces of previous messages and transmit a spurious new message that may look authentic.” - Phleeger & Phleeger in Security in Computing, Pearson, 2003

*Is this a disadvantage of stream cipher? Why, why not?*

**Security goal of encryption: Confidentiality**



# Status

*Confidentiality: Check!*

*Integrity: Missing*



# **Message authentication code (MAC)**



# Message authentication code

Goal: Provide integrity

Process

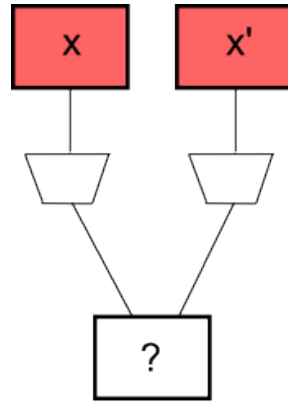
Choose a cryptographic hash function  $h : \{0,1\}^x \rightarrow \{0,1\}^n$

Sender: Send  $h(m), m$

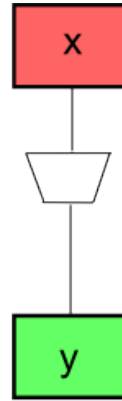
Receiver: Calculate  $h(m)$  and verify it matches  $h(m)$

Examples MD5 ( $n = 128$ ), SHA-256 ( $n = 256$ )

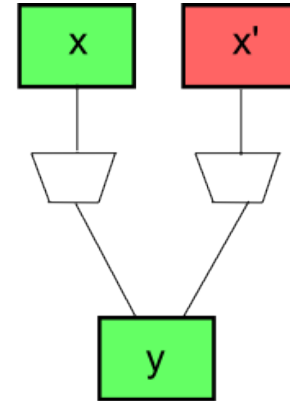
# Cryptographic hash functions



Finding  
Collision



Finding  
Inversion



Finding  
2nd Pre-image



# Hash-based MAC (HMAC)

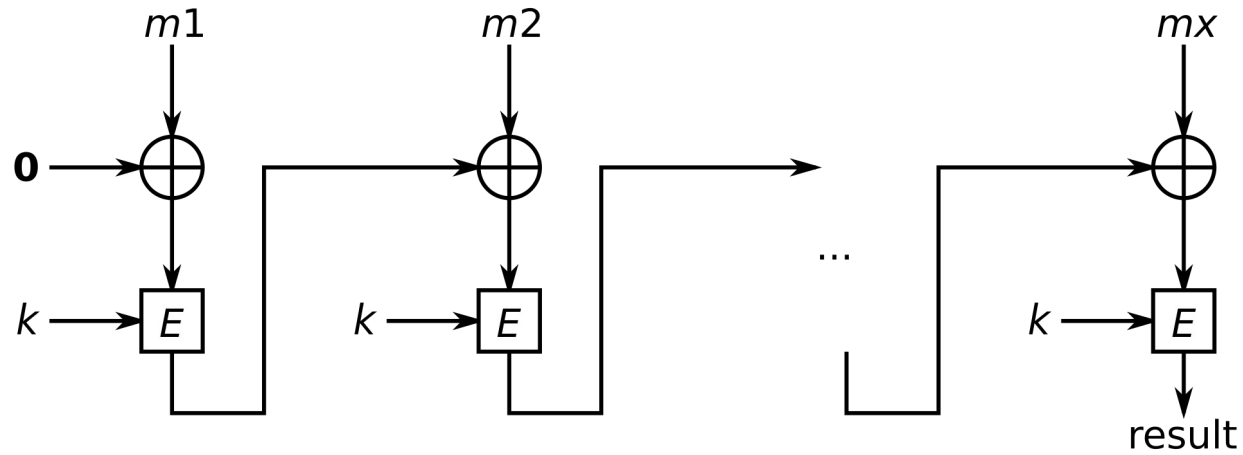
RFC2104: Hash-based MAC

$\text{HMAC}(h,k,m) =$

$$h((k \oplus \text{opad}) \parallel h((k \oplus \text{ipad}) \parallel m))$$

HMAC provides integrity and authenticity

# CBC-MAC





# Car keys

Your car key sends the code for "open the door", together with a MAC, to the car whenever you press the button.

*What could go wrong?*

Replay attack: attacker records message and replays it later

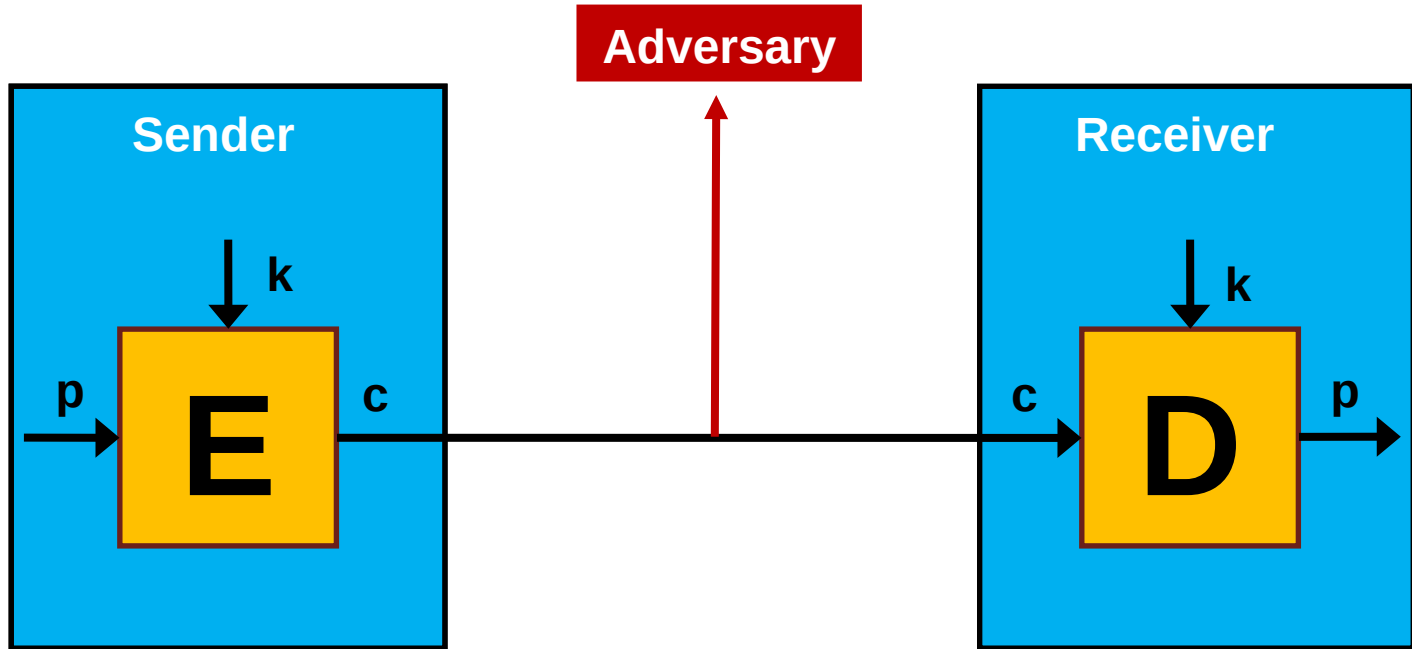
We need some freshness: a timestamp or nonce



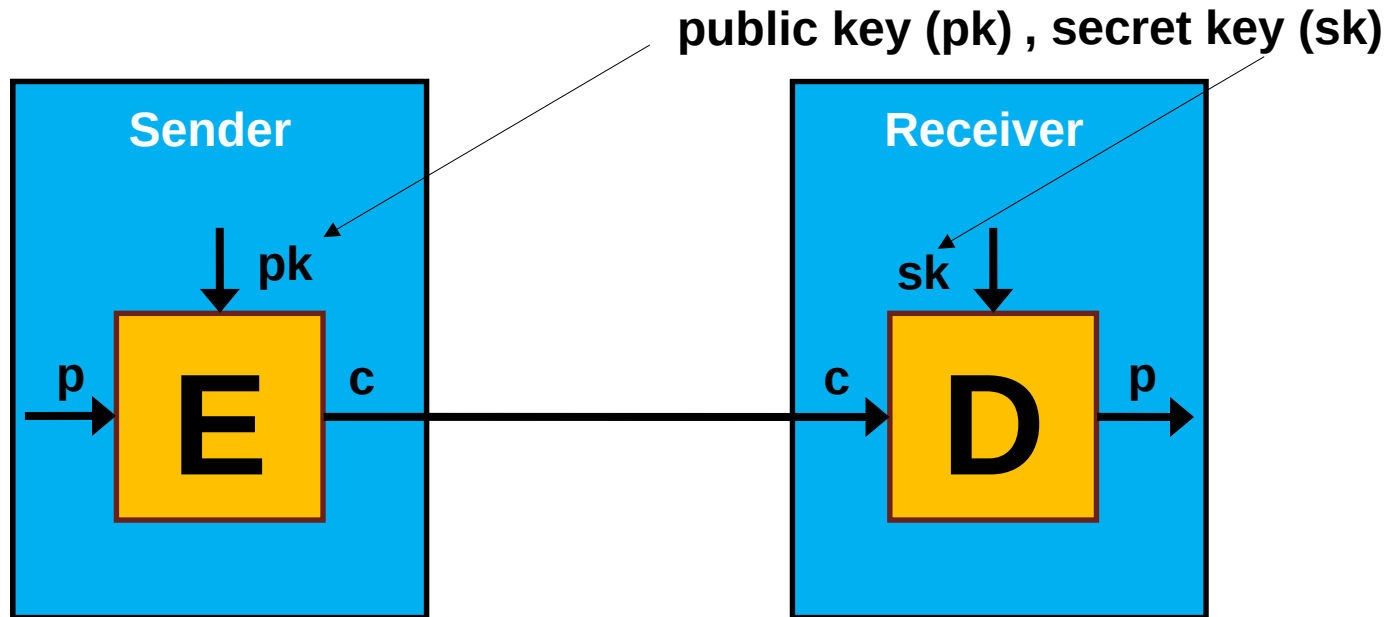


# Non-repudiation

# Cryptosystems



# Enter: Asymmetric encryption



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# Analogy: Combination locks

Bob sends out locks with combination he only knows

Alice picks one of Bob's locks, places her message in a box and locks it with Bob's lock

Bob is the only one who can open the box now





# No pre-shared key!

Bob

Publish public key, protect private key

Alice

Encrypt message with Bob's public key

Bob

Decrypts with his private key



# **Rivest Shamir Adleman (RSA), 1978**

First asymmetric cryptosystem



# RSA encryption and decryption

Public key (N,e), private key (d)

$$C = M^e \pmod{N}$$

$$M = C^d \pmod{N}$$

Asymmetric encryption: Yes! But what about non-repudiation?



# Reverse

Public key  $(N, e)$ , private key  $(d)$

Signature  $\text{sig}(M) = M^d \pmod{N}$

Verify  $\text{ver}(M, \text{sig}(M)) = \text{true}$  iff  $M = (M^d)^e \pmod{N}$





# Wrap-up



# Security goals achieved

Confidentiality

Integrity

Authenticity

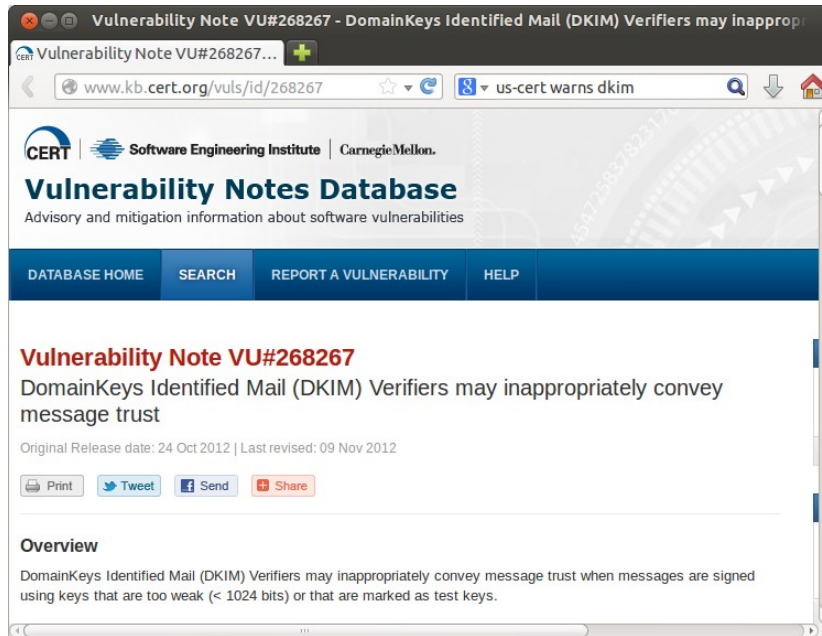
Non-repudiation

CHECK!



**But crypto can still fail**

# Small keys fail



The screenshot shows a web browser window with the title "Vulnerability Note VU#268267 - DomainKeys Identified Mail (DKIM) Verifiers may Inappropriately convey message trust". The address bar shows the URL "www.kb.cert.org/vuls/id/268267". The page header includes the CERT logo and the text "Software Engineering Institute | Carnegie Mellon". The main heading is "Vulnerability Notes Database" with the subtitle "Advisory and mitigation information about software vulnerabilities". A navigation bar contains links for "DATABASE HOME", "SEARCH", "REPORT A VULNERABILITY", and "HELP". The main content area features the title "Vulnerability Note VU#268267" in red, followed by the subtitle "DomainKeys Identified Mail (DKIM) Verifiers may inappropriately convey message trust". Below this, it states "Original Release date: 24 Oct 2012 | Last revised: 09 Nov 2012". There are buttons for "Print", "Tweet", "Send", and "Share". The "Overview" section begins with the text: "DomainKeys Identified Mail (DKIM) Verifiers may inappropriately convey message trust when messages are signed using keys that are too weak (< 1024 bits) or that are marked as test keys."

Vulnerability Note VU#268267 - DomainKeys Identified Mail (DKIM) Verifiers may Inappropriately convey message trust

Vulnerability Note VU#268267

DomainKeys Identified Mail (DKIM) Verifiers may inappropriately convey message trust

Original Release date: 24 Oct 2012 | Last revised: 09 Nov 2012

Print Tweet Send Share

**Overview**

DomainKeys Identified Mail (DKIM) Verifiers may inappropriately convey message trust when messages are signed using keys that are too weak (< 1024 bits) or that are marked as test keys.

# Collision fail



The image is a screenshot of the top portion of an Ars Technica web page. At the top left is the Ars Technica logo, consisting of an orange circle with the word 'ars' in white and 'technica' in white text to its right. To the right of the logo is an orange rectangular box containing the text 'See what Accuweather built for Windows' in white. Below these elements is a dark grey navigation bar with a home icon, 'MAIN MENU', 'MY STORIES: 25', 'FORUMS', 'SUBSCRIBE', and 'VIDEO'. Below the navigation bar is a large grey header with the text 'RISK ASSESSMENT / SECURITY & HACKTIVISM'. Underneath this is the article title 'Crypto breakthrough shows Flame was designed by world-class scientists' in a large, bold, dark font. Below the title is a subtitle in a smaller, lighter font: 'The spy malware achieved an attack unlike any cryptographers have seen before.' At the bottom left of the article header is the byline 'by Dan Goodin - June 7 2012, 8:20pm -200'. At the bottom right are two orange tags, 'BLACK HAT' and 'NATIONAL SECURITY', followed by a grey box containing the number '161'.

ars technica See what Accuweather built for Windows

MAIN MENU MY STORIES: 25 FORUMS SUBSCRIBE VIDEO

RISK ASSESSMENT / SECURITY & HACKTIVISM

**Crypto breakthrough shows Flame was designed by world-class scientists**

The spy malware achieved an attack unlike any cryptographers have seen before.

by Dan Goodin - June 7 2012, 8:20pm -200

BLACK HAT NATIONAL SECURITY 161

# Impressive fail

## New attack steals e-mail decryption keys by capturing computer sounds

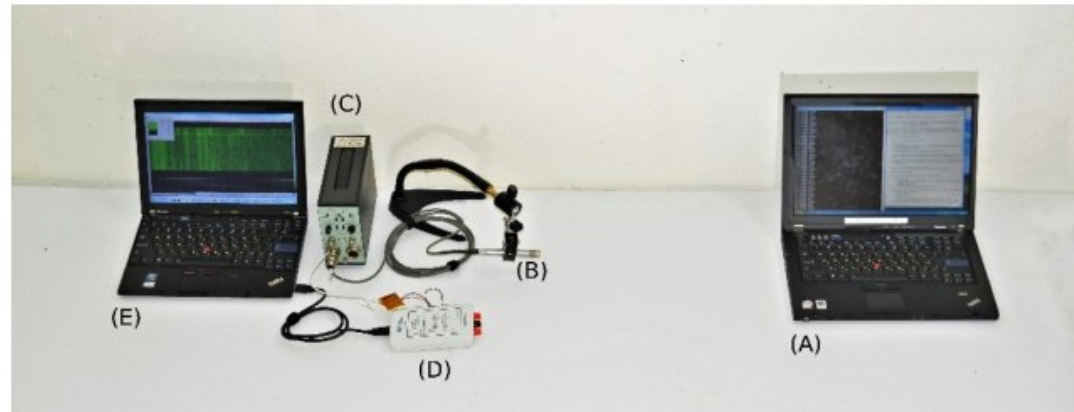
Scientists use smartphone to extract secret key of nearby PC running PGP app.

by Dan Goodin - Dec 18, 2013 11:25 pm UTC

Share

Tweet

108





# Bad choice fail

## IRS Encourages Poor Cryptography

Buried in one of the [documents](#) are the [rules for encryption](#):

While performing AES encryption, there are several settings and options depending on the tool used to perform encryption. IRS recommended settings should be used to maintain compatibility:

- Cipher Mode: ECB (Electronic Code Book).
- Salt: No salt value
- Initialization Vector: No Initialization Vector (IV). If an IV is present, set to all zeros to avoid affecting the encryption.
- Key Size: 256 bits / 32 bytes Key size should be verified and moving the key across operating systems can affect the key size.
- Encoding: There can be no special encoding. The file will contain only the raw encrypted bytes.
- Padding: PKCS#7 or PKCS#5.

ECB? Are they [serious](#)?

# DIY fail

**Smart grid security WORSE than we thought**



OSGP's DIY MAC is a JOKE





# Backdoor fail

Topic: *Security*

Follow via:  

## NIST finally dumps NSA-tainted random number algorithm

**Summary:** *Many years since a backdoor was discovered, probably planted by the NSA, public pressure finally forces NIST to formally remove Dual\_EC\_DRBG from their recommendations.*



By [Larry Seltzer](#) for [Zero Day](#) | April 23, 2014 -- 14:04 GMT (07:04 PDT)

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# Supply chain fail

## Schneier on Security

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### Crypto AG Was Owned by the CIA

The Swiss cryptography firm Crypto AG sold equipment to governments and militaries around the world for decades after World War II. They were owned by the CIA:

But what none of its customers ever knew was that Crypto AG was secretly owned by the CIA in a highly classified partnership with West German intelligence. These spy agencies rigged the company's devices so they could easily break the codes that countries used to send encrypted messages.

This isn't really news. We have long known that Crypto AG was backdooring crypto equipment for the Americans. What is new is the formerly classified documents describing the details:

The decades-long arrangement, among the most closely guarded secrets of the Cold War, is laid bare in a classified, comprehensive CIA history of the operation obtained by The Washington Post and ZDF, a German public broadcaster, in a joint reporting project.

The account identifies the CIA officers who ran the program and the

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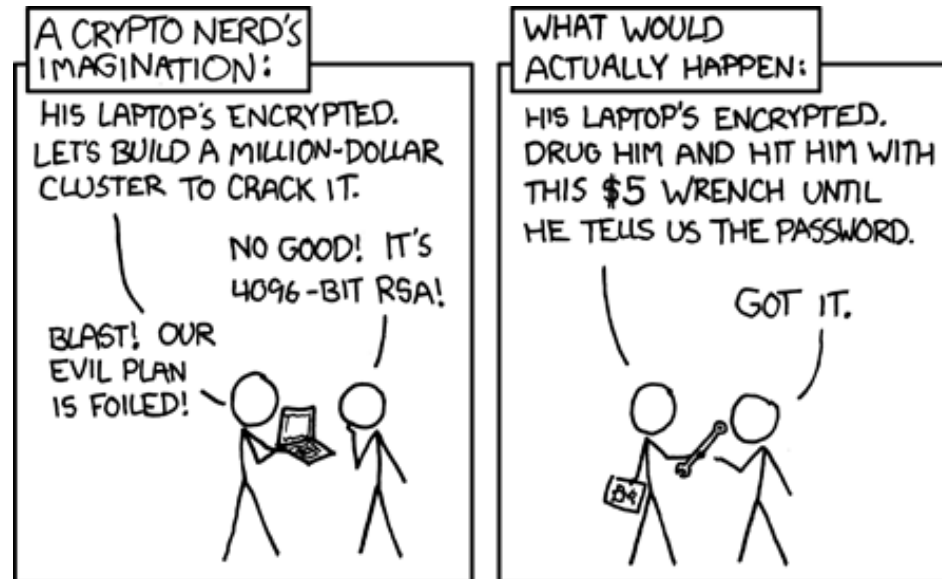
#### About Bruce Schneier



# Malware fail



# Real-world fail



# Suggested reading

