

Algorithms for Data Science

Optimization: Integer Programming

Integer Programming

Integer programming (IP) is an extension of linear programming where some or all variables are restricted to integer values.

Standard Form:

Maximize or Minimize $c^T x$ s.t. $Ax \leq b, x \in \mathbb{Z}^n$

Key Differences from LP:

LP allows for continuous variables, IP requires discrete variables. Integer constraints make the problem combinatorial.









Types of Integer Programming

Pure IP

- All decision variables are integers.
- E.g. Assigning workers to tasks where partial assignments are not allowed.
- Use Case: Resource Allocation

Mixed IP (MIP)

- Some variables are integers, others continuous.
- E.g. Optimizing costs with integer constraints but continuous costs.
- Use Case: Portfolio Optimization

Binary IP

- Variables are restricted to 0 or 1 (binary).
- E.g. Selecting projects to fund given budget.
- Use Case: Project Selction



Methods for Solving IP Problems

Branch & Bound

 Systematically divides the problem into smaller subproblems to exclude infeasible solutions.

Process:

- Solve the relaxed LP problem.
- Identify fractional variables.
- Prune subproblems

Cutting Plane

 Iteratively refines the feasible region by adding linear constraints to exclude non-integer solutions.

Process:

- Solve the relaxed LP problem.
- Identify violated integer constraints.
- Add cutting planes.

Heuristic Methods

 Approximation techniques that provide near-optimal solutions quickly.

• Examples:

- Genetic Algorithms
- Simulated Annealing



Integer Programming Example

- **Objective:** Assign tasks to workers while minimizing the total costs.
- **Constraints:**
 - Each task must be assigned to exactly one coworker.
 - Each worker can handle only a limited number of tasks.
 - Assignments must be binary.
- **Mathematical Formulation:**
 - \circ Let x_{ij} represent whether task j is assigned to worker i or not.

$$\circ$$
 Objective Function: Minimize $\sum_{i=1}^m \sum_{j=1}^n c_{ij} x_{ij}$

$$\circ$$
 Constraints: $\sum_{iar{n}^1}^m x_{ij} = 1, orall j = 1$ (Each task assigned to one worker)

$$\sum_{i=1}^n x_{ij} \leq k_i, orall i$$
 (Worker capacity constraint)



$$x_{ij} \in \{0,1\}, orall i, j$$

