

Generating Random WMC Instances

An Empirical Analysis with Varying Primal Treewidth

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THE UNIVERSITY OF EDINBURGH

informatics



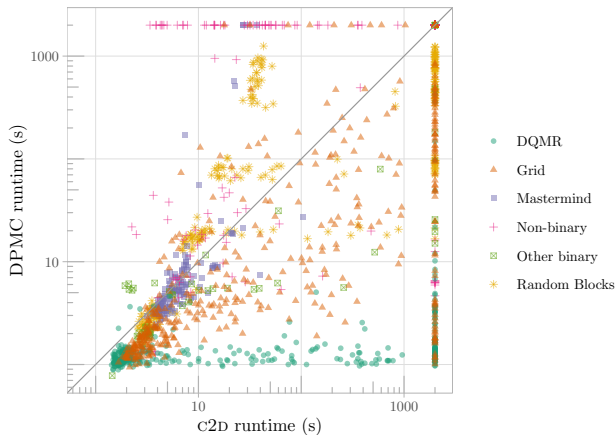
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Which Algorithm Is Better? It Depends on the Data



Data from Dilkas and Belle (2021): various Bayesian networks encoded using the method by Darwiche (2002)

The Problem: Weighted Model Counting (WMC)

- ▶ A generalisation of propositional model counting ($\#SAT$)
- ▶ Applications:
 - ▶ graphical models
 - ▶ probabilistic programming
 - ▶ neural-symbolic artificial intelligence

Example

$$w(x) = 0.3, w(\neg x) = 0.7, \\ w(y) = 0.2, w(\neg y) = 0.8$$

$$WMC(x \vee y) = w(x)w(y) + \\ w(x)w(\neg y) + w(\neg x)w(y) = 0.44$$

(Some of the) WMC Algorithms

- ▶ CACHET (Sang et al. 2004)
 - ▶ a SAT solver with clause learning and component caching
- ▶ C2D (Darwiche 2004)
 - ▶ knowledge compilation to d-DNNF
- ▶ D4 (Lagniez and Marquis 2017)
 - ▶ knowledge compilation to decision-DNNF
- ▶ MINIC2D (Oztok and Darwiche 2015)
 - ▶ knowledge compilation to decision sentential decision diagrams
- ▶ DPMC (Dudek, Phan and Vardi 2020)
 - ▶ dynamic programming with algebraic decision diagrams and tree decomposition based planning

Tree Decompositions and Primal Treewidth

Formula in CNF:

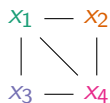
$$\phi = (x_4 \vee \neg x_3 \vee x_1) \wedge (\neg x_2 \vee x_4) \wedge (\neg x_1 \vee x_2 \vee x_4)$$

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Its primal graph:



Tree Decompositions and Primal Treewidth

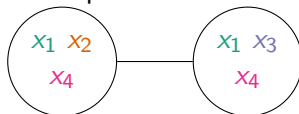
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Its primal graph:



Its minimum-width tree decomposition:



Tree Decompositions and Primal Treewidth

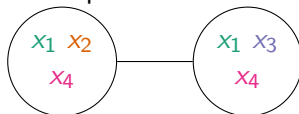
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Its primal graph:



Its minimum-width tree decomposition:



\therefore the primal treewidth of ϕ is 2

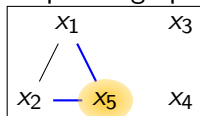
From Random SAT to Random WMC

We introduce parameter $\rho \in [0, 1]$ that biases the probability distribution towards adding variables that would introduce fewer new edges to the primal graph.

Example partially-filled formula:

$$(\neg x_5 \vee x_2 \vee x_1) \wedge (x_5 \vee ? \vee ?)$$

Its primal graph:



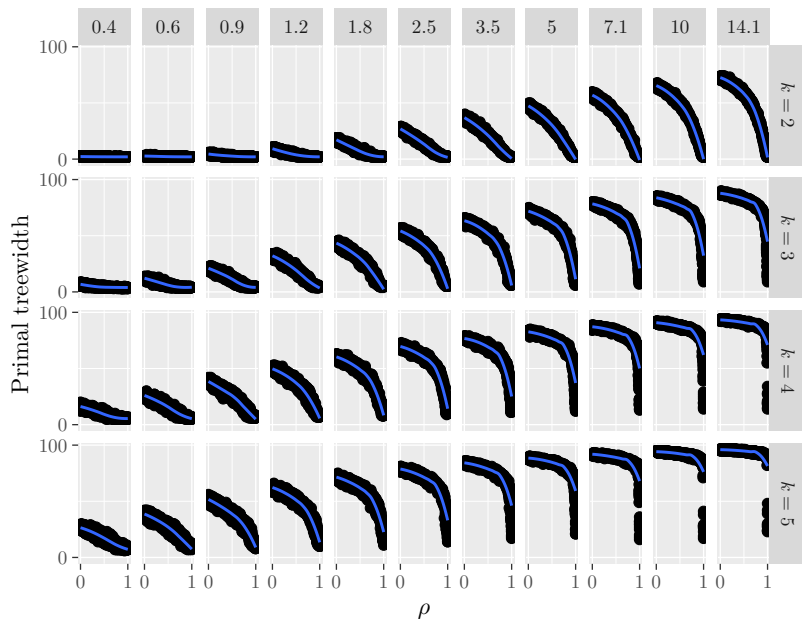
The probability distribution for the next variable

Base probability of each variable being chosen:

$$\frac{1 - \rho}{4}.$$

Both x_1 and x_2 get a bonus probability of $\rho/2$ for each being the endpoint of **one** out of the **two** neighbourhood edges.

The Relationship Between ρ and Primal Treewidth



Peak Hardness w.r.t. Density

Let μ denote the **density**, i.e., the number of clauses divided by the number of variables.

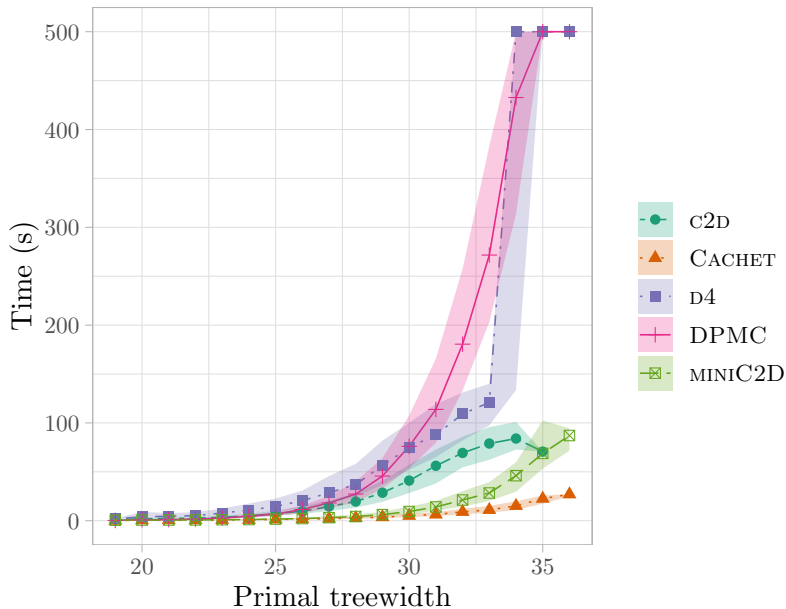
- ▶ CACHET is known to peak at $\mu = 1.8$ (Sang et al. 2004)
- ▶ Bayardo Jr. and Pehoushek (2000) show some #SAT algorithms to peak at $\mu = 1.2$ and $\mu = 1.9$

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- ▶ Bayardo Jr. and Pehoushek (2000) show some #SAT algorithms to peak at $\mu = 1.2$ and $\mu = 1.9$
- ▶ In our experiments:
 - ▶ DPMC peaks at $\mu = 2.2$
 - ▶ all other algorithms peak at $\mu = 1.9$

Hardness w.r.t. Primal Treewidth (when $\mu = 1.9$)

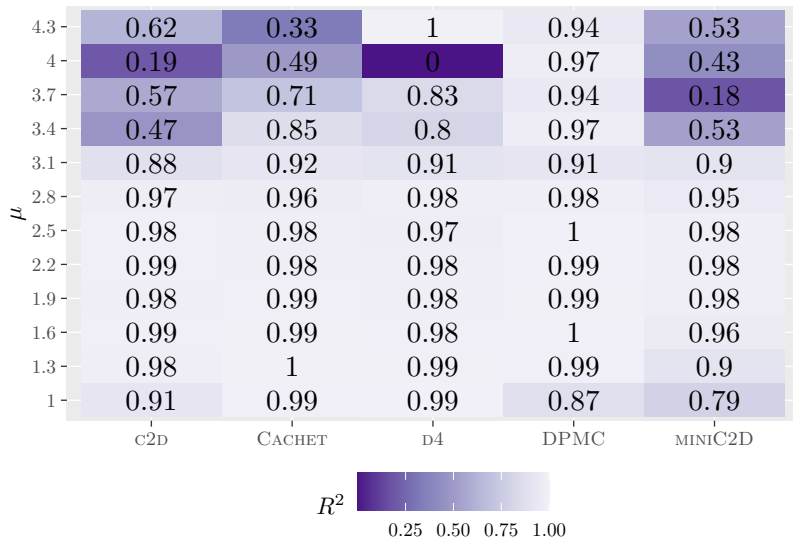


Is The Relationship Exponential?

Let us fit the model $\ln t \sim \alpha w + \beta$, i.e., $t \sim e^\beta (e^\alpha)^w$, where t is runtime, and w is primal treewidth

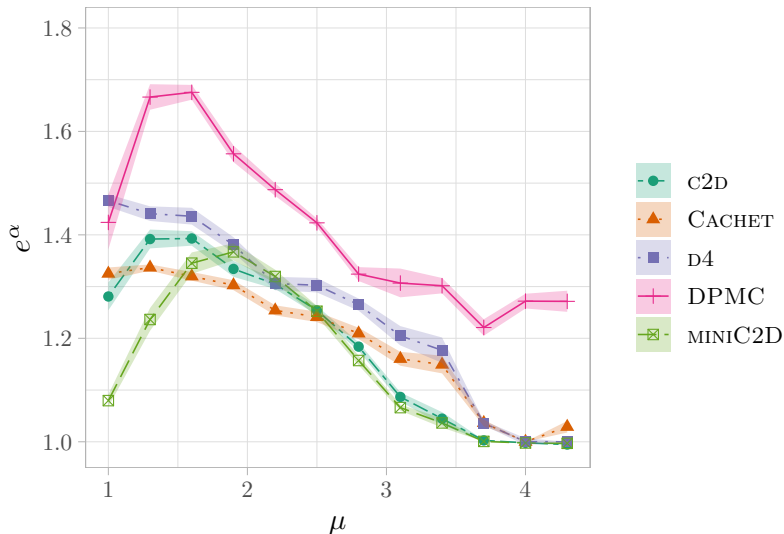
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Summary

- ▶ Introduced a random model for WMC instances
- ▶ Observations:
 - ▶ All algorithms scale exponentially w.r.t. primal treewidth
 - ▶ The running time of DPMC peaks at a higher density and scales worse w.r.t. primal treewidth
- ▶ Future work:
 - ▶ A theoretical relationship between ρ and primal treewidth
 - ▶ Non- k -CNF instances
 - ▶ Algorithm portfolios for WMC