## Towards Practical First-Order Model Counting

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### Motivation

## **Example Setting**

- ▶ Let  $\triangle$  be a set of cardinality n
- ▶ Suppose we want to count all  $P \subseteq \Delta^2$  (as a function of n) that are:
  - functions,
  - bijections,
  - partial orders,
  - symmetric,
  - transitive,
  - etc.

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  - etc.
- Propositional model counting (#SAT) is #P-complete
- But many of these counting problems have efficient solutions
- ► And we can find them using first-order model counting
  - ▶ i.e., reasoning about sets, subsets, and arbitrary elements without grounding them

# First-Order Model Counting

### The Problem with CRANE

## A Solution Produced for the Bijection-Counting Problem

$$f(m,n) = \sum_{l=0}^{n} {n \choose l} (-1)^{n-l} g(l,m),$$
  
$$g(l,m) = g(l-1,m) + mg(l-1,m-1)$$

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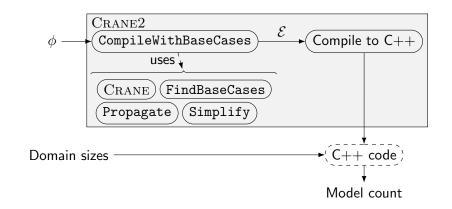
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#### Issues

Completeness: what are the base cases of g?

Usability: how do I compute, e.g., f(7,7)?

## Knowledge Compilation Workflow



# Knowledge Compilation Workflow (1/2)

- 1. Use Crane to compile  $\phi$  into a set of equations  $\mathcal E$
- 2. Simplify them, e.g.,

$$g(l,m) = \sum_{k=0}^{m} [0 \le k \le 1] {m \choose k} g(l-1, m-k)$$

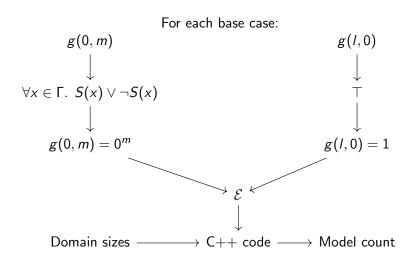
becomes

$$g(l,m) = g(l-1,m) + mg(l-1,m-1)$$

3. Identify a sufficient set of base cases

• e.g., 
$$\{g(0,m),g(1,0)\}$$

# Knowledge Compilation Workflow (2/2)



### **Benchmarks**

► Friends & Smokers

$$(\forall x, y \in \Delta. \ S(x) \land F(x, y) \to S(y)) \land (\forall x \in \Delta. \ S(x) \to C(x))$$

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Functions

$$(\forall x \in \Gamma. \exists y \in \Delta. P(x,y)) \land (\forall x \in \Gamma. \forall y, z \in \Delta. P(x,y) \land P(x,z) \rightarrow y = z)$$

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Bijections

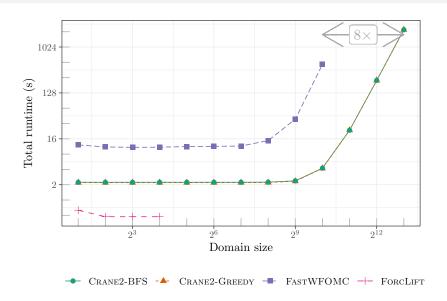
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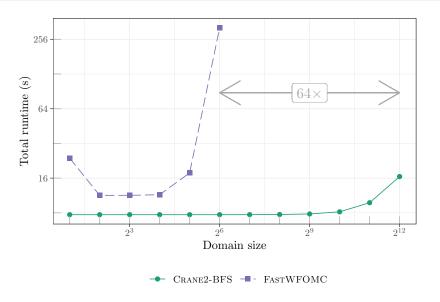
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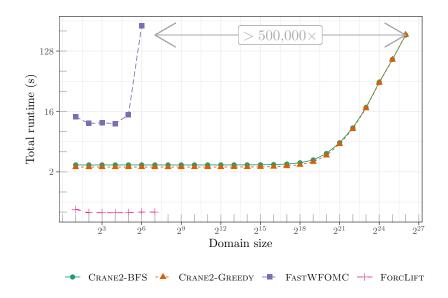
## Friends & Smokers



# Bijections



## **Functions**



# Summary

TODO: and future work