Weighted Model Counting Without Parameter Variables

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The Computational Problem of Probabilistic Inference

ProbLog

```
0.001 :: burglary.
0.002 :: earthquake.
0.95 :: alarm
                  :- burglary, earthquake.
0.94 :: alarm :- burglary, \+ earthquake.
0.29 :: alarm :- \+ burglary, earthquake.
0.001 :: alarm
              :- \+ burglary . \+ earthquake .
0.9
     :: johnCalls :- alarm.
0.05
     :: johnCalls :- \+ alarm.
0.7
     :: marvCalls :- alarm.
0.01
     :: maryCalls :- \+ alarm.
```

Bayesian Network



BLOG

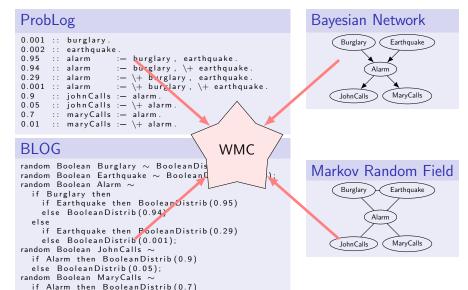
```
random Boolean Burglary ~ BooleanDistrib(0.001);
random Boolean Earthquake ~ BooleanDistrib(0.002);
random Boolean Alarm ~

if Burglary then
 if Earthquake then BooleanDistrib(0.95)
 else BooleanDistrib(0.94)
 else
 if Earthquake then BooleanDistrib(0.29)
 else BooleanDistrib(0.001);
random Boolean JohnCalls ~
 if Alarm then BooleanDistrib(0.9)
 else BooleanDistrib(0.05);
random Boolean MaryCalls ~
 if Alarm then BooleanDistrib(0.7)
 else BooleanDistrib(0.01);
```

Markov Random Field



The Computational Problem of Probabilistic Inference



else Boolean Distrib (0.01);

Weighted Model Counting (WMC)

- Generalises propositional model counting (#SAT)
- Applications:
 - graphical models
 - probabilistic programming
 - neural-symbolic artificial intelligence
- ► Main types of algorithms:
 - using knowledge compilation
 - using a SAT solver
 - manipulating pseudo-Boolean functions

$$w(x) = 0.3, w(\neg x) = 0.7,$$

 $w(y) = 0.2, w(\neg y) = 0.8$

$$WMC(x \lor y) = w(x)w(y) + w(x)w(\neg y) + w(\neg x)w(y) = 0.44$$

The Problem with Assigning Weights to Literals

A Simple Bayesian Network



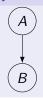
- ► from 2 binary variables
- ▶ to 8 variables and 17 clauses
- with lots of redundancy

Its WMC Encoding

```
p cnf 8 17
-2 -1 0
1 2 0
-3 1 0
-1 \ 3 \ 0
-5 -1 0
-5 -4 0
1 4 5 0
-6 -1 0
-6 4 0
-4160
-7 1 0
-7 -4 0
-1 4 7 0
-810
-8 \ 4 \ 0
-4 -1 8 0
-40
c weights 1.0 \ 1.0 \ 0.5 \ 1.0 \ \setminus
0.5 1.0 1.0 1.0 0.6 1.0 \
0.4 1.0 0.1 1.0 0.9 1.0
```

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- ► from 2 binary variables
- ▶ to 8 variables and 17 clauses
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Its WMC Encoding

```
p cnf 8 17
-2 -1 0
                     \neg x_1 \Leftrightarrow x_2
1 2 0
-310
                       x_1 \Leftrightarrow x_3
-1 3 0
-5 -1 0
-5 -4 0
                    \neg x_1 \land \neg x_4 \Leftrightarrow x_5
1 4 5 0
-6 -1 0
-640
                      \neg x_1 \land x_4 \Leftrightarrow x_6
-4160
-7 -4 0
                     X_1 \land \neg X_4 \Leftrightarrow X_7
-1 4 7 0
-840
                       x_1 \wedge x_4 \Leftrightarrow x_8
-4 -1 8 0
                        \neg x_4
-40
c weights 1.0 1.0 0.5 1.0 \
0.5 1.0 1.0 1.0 0.6 1.0 \
041001100910
```

Outline

A More Expressive Alternative

When Does This Transformation Work?

How Good Is It?

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WMC, Formally

Definition

A WMC instance is a tuple (ϕ, X_I, X_P, w) , where

- \triangleright X_I is the set of indicator variables,
- ▶ X_P is the set of parameter variables (with $X_I \cap X_P = \emptyset$),
- ϕ is a propositional formula in CNF over $X_I \cup X_P$,
- ▶ $w: X_I \cup X_P \cup \{\neg x \mid x \in X_I \cup X_P\} \rightarrow \mathbb{R}$ is a weight function
 - ▶ such that $w(x) = w(\neg x) = 1$ for all $x \in X_I$.

Definition

Let ϕ be a formula over a set of variables X. Then $Y \subseteq X$ is a minimum-cardinality model of ϕ if

- \triangleright $Y \models \phi$,
- ▶ and $|Y| \le |Z|$ for all $Z \models \phi$.

WMC and Minimum-Cardinality WMC

The goal of WMC is to compute

$$\sum_{Y \models \phi} \prod_{Y \models I} w(I)$$

whereas the goal of minimum-cardinality WMC is to compute

$$\sum_{Y \models \phi, \ |Y| = k} \prod_{Y \models I} w(I),$$

where

$$k = \min_{Y \models \phi} |Y|.$$

A More Expressive Alternative

For any propositional formula ϕ over a set of variables X and $p,q\in\mathbb{R}$, let $[\phi]_q^p\colon 2^X\to\mathbb{R}$ be the pseudo-Boolean function defined as

$$[\phi]_q^p(Y) := egin{cases} p & ext{if } Y \models \phi \ q & ext{otherwise} \end{cases}$$

for any $Y \subseteq X$.

Definition (Pseudo-Boolean Projection (PBP))

A PBP instance is a tuple (F, X, ω) , where X is the set of variables, F is a set of two-valued pseudo-Boolean functions $2^X \to \mathbb{R}$, and $\omega \in \mathbb{R}$ is the scaling factor.

Example

- ► Indicator variable: x
- ► Parameter variables: p, q
- ▶ Weights: w(p) = 0.2, w(q) = 0.8, and $w(\neg p) = w(\neg q) = 1$

WMC Clause

```
    \begin{array}{l}
      \neg x \Rightarrow p \\
      p \Rightarrow \neg x \\
      x \Rightarrow q
    \end{array}
```

 $q \Rightarrow x$ $\neg x$

- ► Indicator variable: x
- ► Parameter variables: p, q
- ► Weights: w(p) = 0.2, w(q) = 0.8, and $w(\neg p) = w(\neg q) = 1$

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WMC Clause	In CNF	Pseudo-Boolean Function
$\neg x \Rightarrow p$	$x \lor p$	$[\neg x]_1^{0.2}$
$p \Rightarrow \neg x$	$\neg x \lor \neg p$	0.0
$x \Rightarrow q$	$\neg x \lor q$	$[x]_1^{0.8}$
$q \Rightarrow x$	$x \vee \neg q$	1
$\neg x$	$\neg \chi$	$[\neg x]_0^1$

- ► Indicator variable: x
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WMC Clause	In CNF	Pseudo-Boolean Function	
$\neg x \Rightarrow p$	$x \lor p$	$[\neg x]_1^{0.2}$	0.0
$p \Rightarrow \neg x$	$\neg x \lor \neg p$. 10.0	$[x]_{0.2}^{0.8}$
$x \Rightarrow q$	$\neg x \lor q$	$[x]_1^{0.8}$	
$q \Rightarrow x$	$x \vee \neg q$	1	
¬ <i>x</i>	$\neg X$	$[\neg x]_0^1$	$[\neg x]_0^1$

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Correctness Conditions for WMC (1/2)

For each parameter variable $p \in X_P$,

- \blacktriangleright $w(\neg p) = 1$,
- \triangleright and the set of clauses that mention p or $\neg p$ is

$$\left\{p \vee \bigvee_{i=1}^{n} \neg I_{i}\right\} \cup \left\{I_{i} \vee \neg p \mid i = 1, \ldots, n\right\}$$

for some non-empty family of indicator literals $(I_i)_{i=1}^n$.

Correctness Conditions for WMC (2/2)

For each parameter variable $p \in X_P$,

- \blacktriangleright $w(p) + w(\neg p) = 1$,
- each clause has at most one parameter variable,
- ▶ there is no clause $c \in \phi$ such that $\neg p \in c$,
- ▶ if $\{p\} \in \phi$, then this is the only clause that mentions p,
- ▶ and for any $c, d \in \phi$ such that $c \neq d$, $p \in c$ and $p \in d$,

$$\bigwedge_{I \in c \setminus \{p\}} \neg I \wedge \bigwedge_{I \in d \setminus \{p\}} \neg I$$

is false.

Additional Conditions for Minimum-Cardinality WMC

- ▶ All models of $\{c \in \phi \mid c \cap X_P = \emptyset\}$ have the same number of positive indicator literals,
- and

$$\min_{Z\subseteq X_P} |Z|$$
 s.t. $Y \cup Z \models \phi$

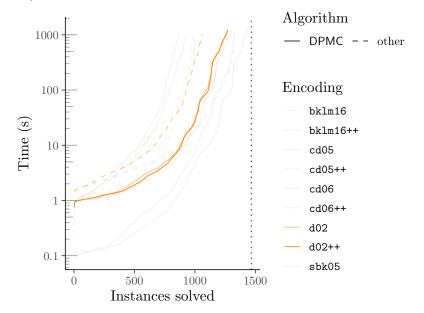
is the same for all $Y \models \{c \in \phi \mid c \cap X_P = \emptyset\}$.

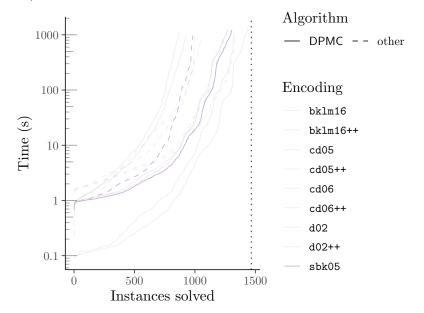
Outline

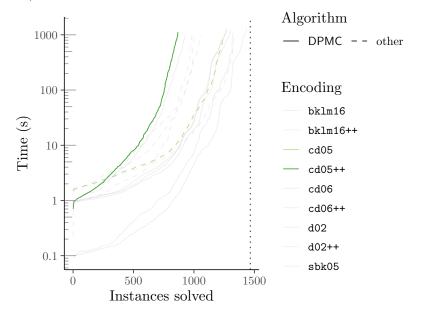
A More Expressive Alternative

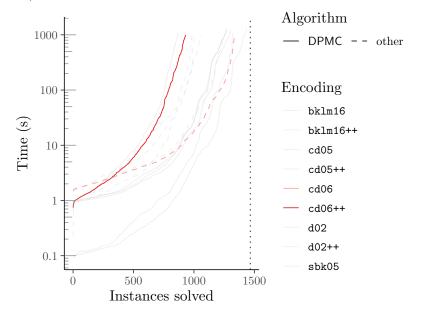
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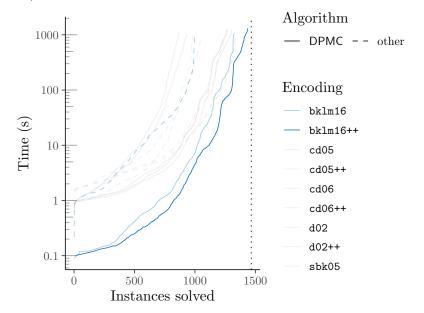
How Good Is It?



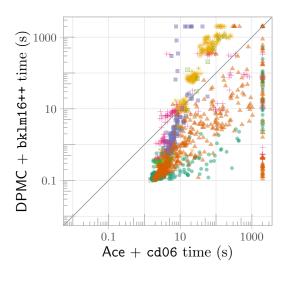






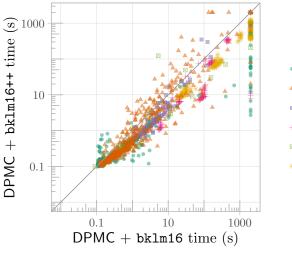


Compared to the Previous State of the Art



- DQMR
- Grid
- Mastermind
 - Non-binary
- Other binary
- Random Blocks

The Best Encoding for DPMC: Before and After



- DQMR
- Grid
- Mastermind
- Non-binary
- Other binary
- * Random Blocks

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