

Weighted Model Counting with Conditional Measures

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1 Introduction

- The Main Narrative

1. When weights are defined on literals, the measure on the free BA is fully independent.
2. This means that the BA itself must be larger (i.e., have additional ‘meaningless’ literals) to turn any probability distribution into an independent one.
3. We show how we can define conditional weights on literals, allowing us to encode any probability distribution into a Boolean algebra that’s not necessarily independent and thus can be smaller.
4. We demonstrate a specific example of this by presenting a new way to encode Bayesian networks into instances of WMC and adapting a WMC algorithm (ADDMC) to run on the new format.
5. We show that this results in significantly faster inference.
6. We show that our encoding results in asymptotically fewer literals and fewer ADDs, and thus a simpler problem.
7. (Maybe) we experimentally demonstrate a phase transition based on the number of variables per ADD.

- Potential criticism may be that this is a step backwards and doesn’t allow us to use SAT-based techniques for probabilistic inference. However, they can still be used for the ‘theory+query’ part.

- Zero-probability weights and one-probability weights can be interpreted as logical clauses. This doesn’t affect ADDMC but could be useful for other solvers.

F What are the main claims, what are the main takeaways, intuitive [??] of theorems to follow. To do this, we appeal to algebraic constructions to define the main concepts for introducing measures on Boolean algebras.

- Algorithms¹

- ADDMC [19] (rediscovered the multiplicativity of BAs in different words) (with optimal settings)
- Cachet [42]
- c2d [17]
- d4 [34] (closed source, boo!)
- miniC2D [39]

- Notable previous/related work

- Hailperin’s approach to probability logic [24]

¹<http://beyonddnp.org/pages/solvers/model-counters-exact/>

Table 1: A comparison of Boolean-algebraic (BA) and set-theoretic (ST) concepts for 2^X for some set X

BA name	BA symbol	ST symbol	ST name
bottom	\perp	\emptyset	empty set
top	\top	X	
meet, and	\wedge	\cap	intersection
join, or	\vee	\cup	union
complement, not	\neg	c	complement
	\leq	\subseteq	subset relation, set inclusion
atom			singleton, unit set

- Nilsson’s (somewhat successful) probabilistic logic [37, 38]
- Logical induction: a big paper with a good overview of previous attempts to assign probabilities to logical sentences in a sensible way [22]
- Measures on Boolean algebras
 - * On possibility and probability measures in finite Boolean algebras [7]
 - * Representation of conditional probability measures [32]
- Intuitively, a measure is just like a probability, except it’s in $\mathbb{R}_{\geq 0}$ instead of $[0, 1]$.
- Somewhere: explain things like CNF, CPT, etc.

2 Boolean Algebras and Power Sets

Let X be a set and let 2^X be its power set. We can equivalently interpret 2^X as a Boolean algebra. See Table 1 for a summary of the differences in terminology and notation. We will use both.

2.1 The Space of Functions on Boolean Algebras

Definition 1 (Operations on functions). Let $A: 2^X \rightarrow \mathbb{R}_{\geq 0}$ and $B: 2^Y \rightarrow \mathbb{R}_{\geq 0}$ be arbitrary functions, $\alpha \in \mathbb{R}_{\geq 0}$, and $x \in X$. We define the following operations:

Addition: $A + B$ is a function $A + B: 2^{X \cup Y} \rightarrow \mathbb{R}_{\geq 0}$ such that

$$(A + B)(\tau) = A(\tau \cap X) + B(\tau \cap Y)$$

for all $\tau \in 2^{X \cup Y}$.

Inverse: \bar{A} is a function $\bar{A}: 2^X \rightarrow \mathbb{R}_{\geq 0}$ such that

$$\bar{A}(\tau) = 1 - A(\tau)$$

for all $\tau \in 2^X$.

Multiplication: $A \cdot B$ is a function $A \cdot B: 2^{X \cup Y} \rightarrow \mathbb{R}_{\geq 0}$ such that

$$(A \cdot B)(\tau) = A(\tau \cap X) \cdot B(\tau \cap Y)$$

for all $\tau \in 2^{X \cup Y}$.

Scalar multiplication: αA is a function $\alpha A: 2^X \rightarrow \mathbb{R}_{\geq 0}$ such that

$$(\alpha A)(\tau) = \alpha \cdot A(\tau)$$

for all $\tau \in 2^X$.

Projection: $\exists_x A$ is a function $\exists_x A: 2^{X \setminus \{x\}} \rightarrow \mathbb{R}_{\geq 0}$ such that

$$(\exists_x A)(\tau) = A(\tau) + A(\tau \cup \{x\})$$

for all $\tau \in 2^{X \setminus \{x\}}$.

Observation 1. Let U be a set, and $\mathcal{V} = \{A: 2^X \rightarrow \mathbb{R}_{\geq 0} \mid X \subseteq U\}$. Then \mathcal{V} is a semi-vector space with three additional operations: inverse, (non-scalar) multiplication, and projection. Specifically, note that both addition and multiplication are both associative and commutative.

Definition 2 (Special functions).

- unit $1: 2^\emptyset \rightarrow \mathbb{R}_{\geq 0}$, $1(\emptyset) = 1$.
- zero $0: 2^\emptyset \rightarrow \mathbb{R}_{\geq 0}$, $0(\emptyset) = 0$.
- logical constant $[a]: 2^{\{a\}} \rightarrow \mathbb{R}_{\geq 0}$, $[a](\emptyset) = 0$, $[a](\{a\}) = 1$.

Henceforth, for any function $A: 2^X \rightarrow \mathbb{R}_{\geq 0}$ and any set τ , we will write $A(\tau)$ to mean $A(\tau \cap X)$.

3 Weighted Model Counting as a Measure

Definition 3 (Measures and weight functions). Let U be a set.

- A *measure* is a function $M: 2^{2^U} \rightarrow \mathbb{R}_{\geq 0}$ such that $M(\perp) = 0$ and

$$M(a \vee b) = M(a) + M(b)$$

for all $a, b \in 2^{2^U}$ whenever $a \wedge b = \perp$.

- A *weight function* is a function $W: 2^U \rightarrow \mathbb{R}_{\geq 0}$.
- A weight function is *factored* if $W = \prod_{x \in U} W_x$ for some functions $W_x: 2^{\{x\}} \rightarrow \mathbb{R}_{\geq 0}$, $x \in U$.
- We say that a weight function $W: 2^U \rightarrow \mathbb{R}_{\geq 0}$ *induces* a measure $M_W: 2^{2^U} \rightarrow \mathbb{R}_{\geq 0}$ if

$$M_W(x) = \sum_{\{u\} \leq x} W(u). \quad (1)$$

- A measure $M: 2^{2^U} \rightarrow \mathbb{R}_{\geq 0}$ is *factorable* if there exists a factored weight function $W: 2^U \rightarrow \mathbb{R}_{\geq 0}$ that induces M .

Lemma 1. The function M_W , as defined by Eq. (1), is a measure.

Proof. Note that $M_W(\perp) = 0$ since there are no atoms below \perp . Let $a, b \in 2^{2^U}$ be such that $a \wedge b = \perp$. By elementary properties of Boolean algebras, all atoms below $a \vee b$ are either below a or below b . Moreover, none of them can be below both a and b because then they would have to be below $a \wedge b = \perp$. Thus

$$M_W(a \vee b) = \sum_{\{u\} \leq a \vee b} W(u) = \sum_{\{u\} \leq a} W(u) + \sum_{\{u\} \leq b} W(u) = M_W(a) + M_W(b)$$

as required. \square

In this formulation, the process of calculating the value of $M_W(x)$ for some $x \in 2^{2^U}$ with a given definition of W is known as *weighted model counting*.

Table 2: An overview of notation for a logic defined over two atoms. The elements in both columns are listed in the same order.

Name in logic	Set-theoretic notation	Boolean-algebraic notation
Atoms (elements of U)	a, b	a, b
Models (elements of 2^U)	$\emptyset, \{a\}, \{b\}, \{a, b\}$	$\neg a \wedge \neg b, a \wedge \neg b, \neg a \wedge b, a \wedge b$
	$\{\emptyset, \{a\}, \{b\}, \{a, b\}\}$	\top
	$\{\emptyset, \{a\}, \{b\}\}, \{\emptyset, \{a\}, \{a, b\}\}$	$\neg a \vee \neg b, a \rightarrow b$
	$\{\emptyset, \{b\}, \{a, b\}\}, \{\{a\}, \{b\}, \{a, b\}\}$	$b \rightarrow a, a \vee b$
Formulas (elements of 2^{2^U})	$\{\emptyset, \{a\}\}, \{\emptyset, \{b\}\}, \{\emptyset, \{a, b\}\}$	$\neg b, \neg a, a \leftrightarrow b$
	$\{\{a\}, \{b\}\}, \{\{a\}, \{a, b\}\}, \{\{b\}, \{a, b\}\}$	$(a \wedge \neg b) \vee (b \wedge \neg a), a, b$
	$\{\emptyset\}, \{\{a\}\}, \{\{b\}\}, \{\{a, b\}\}$	$\neg a \wedge \neg b, a \wedge \neg b, \neg a \wedge b, a \wedge b$
	\emptyset	\perp

3.1 Relation to the Classical (Logic-Based) View of WMC

Would it be more accurate to replace the word ‘model’ with ‘interpretation’ in most situations?

Let \mathcal{L} be a propositional logic with atoms $U = \{a, b\}$. See Table 2 for an overview of the models and formulas in \mathcal{L} as well as their set-theoretic and Boolean-algebraic representations. Note the differences in the meaning of the word ‘atom’ between the logical and the Boolean-theoretic interpretations. In the Boolean algebra 2^{2^U} , an atom is a set with a single element which corresponds to a model of \mathcal{L} (also known as an element of 2^U); whereas an atom of \mathcal{L} is an atomic formula, i.e., an element of U . We will primarily use the term ‘atom’ to mean the former.

Let $w: \{a, b, \neg a, \neg b\} \rightarrow \mathbb{R}_{\geq 0}$ be the *weight function* defined by

$$w(a) = 0.3, \quad w(\neg a) = 0.7, \quad w(b) = 0.2, \quad w(\neg b) = 0.8.$$

Let Δ be a theory in \mathcal{L} with a sole axiom a . Then Δ has two models, i.e., $\{a, b\}$ and $\{a, \neg b\}$. The *weighted model count* (WMC) [12] of Δ is then

$$\text{WMC}(\Delta) = \sum_{\omega \models \Delta} \prod_{\omega \models l} w(l) = w(a)w(b) + w(a)w(\neg b) = 0.3. \quad (2)$$

Alternatively, we can define $W_a: 2^{\{a\}} \rightarrow \mathbb{R}_{\geq 0}$ as

$$W_a(\{a\}) = 0.3, \quad W_a(\emptyset) = 0.7$$

and $W_b: 2^{\{b\}} \rightarrow \mathbb{R}_{\geq 0}$ as

$$W_b(\{b\}) = 0.2, \quad W_b(\emptyset) = 0.8.$$

Let M be the measure on 2^{2^U} induced by $W = W_a \cdot W_b$. Then, equivalently to Eq. (2), we can write

$$M(a) = W(\{a, b\}) + W(\{a\}) = W_a(\{a\})W_b(\{b\}) + W_a(\{a\})W_b(\emptyset) = 0.3.$$

Given the theory Δ , we can also compute the probability of a query b as [5, 43]

$$\Pr_{\Delta, w}(b) = \frac{\text{WMC}(\Delta \wedge b)}{\text{WMC}(\Delta)}.$$

The same thing can be accomplished using the algebraic formulation, with M replacing WMC.

In the rest of the paper, for any set U , we will use set-theoretic notation for 2^U and Boolean-algebraic notation for 2^{2^U} , except for (Boolean) atoms in 2^{2^U} that are denoted as $\{x\}$ for some model $x \in 2^U$.

4 Limitations of Factorable Measures

- F Give a concrete example of something impossible to represent using WMC.
- F Can you say something here about factorized vs non-factorized weight function definitions? That is, factorized is when w maps literals to $R_{\geq 0}$, non-factorized is when w maps models to $R_{\geq 0}$ and
- come up with nice example when non-factorized weights are intuitive;

Lemma 2. For any measure $M: 2^{2^U} \rightarrow \mathbb{R}_{\geq 0}$ and elements $a, b \in 2^{2^U}$,

$$M(a \wedge b) = M(a)M(b) \quad (3)$$

if and only if

$$M(a \wedge b) \cdot M(\neg a \wedge \neg b) = M(a \wedge \neg b) \cdot M(\neg a \wedge b). \quad (4)$$

Proof. First, note that $a = (a \wedge b) \vee (a \wedge \neg b)$ and $(a \wedge b) \wedge (a \wedge \neg b) = 0$, so, by properties of a measure,

$$M(a) = M(a \wedge b) + M(a \wedge \neg b). \quad (5)$$

Applying Eq. (5) and the equivalent expression for $M(b)$ allows us to rewrite Eq. (3) as

$$M(a \wedge b) = [M(a \wedge b) + M(a \wedge \neg b)][M(a \wedge b) + M(\neg a \wedge b)]$$

which is equivalent to

$$M(a \wedge b)[1 - M(a \wedge b) - M(a \wedge \neg b) - M(\neg a \wedge b)] = M(a \wedge \neg b)M(\neg a \wedge b). \quad (6)$$

Since $a \wedge b$, $a \wedge \neg b$, $\neg a \wedge b$, $\neg a \wedge \neg b$ are pairwise disjoint and their supremum is 1,

$$M(a \wedge b) + M(a \wedge \neg b) + M(\neg a \wedge b) + M(\neg a \wedge \neg b) = 1,$$

and this allows us to rewrite Eq. (6) into Eq. (4). As all transformations are invertible, the two expressions are equivalent. \square

Theorem 1. A measure $M: 2^{2^U} \rightarrow \mathbb{R}_{\geq 0}$ is factorable if and only if

$$M(u \wedge v) = M(u)M(v) \quad (7)$$

for all distinct $u, v \in U \cup \{\neg w \mid w \in U\}$ such that $u \neq \neg v$.

Proof. (\Leftarrow) For each $x \in U$, let $W_x: 2^{\{x\}} \rightarrow \mathbb{R}_{\geq 0}$ be defined by $W_x(\{x\}) = M(x)$ and $W_x(\emptyset) = M(\neg x)$. Let M_W be the measure induced by

$$W = \prod_{x \in U} W_x.$$

We will show that $M = M_W$. First, note that $M_w(\perp) = 0 = M(\perp)$ by Definition 3 and Lemma 1. Second, let

$$a = \bigwedge_{u \in U} a_u$$

be an atom in 2^{2^U} such that $a_u \in \{u, \neg u\}$ for all $u \in U$. Then

$$M_W(a) = W(\{a_u \mid u \in U\}) = \prod_{u \in U} W_u(\{a_u\}) = \prod_{u \in U} M(a_u) = M\left(\bigwedge_{u \in U} a_u\right) = M(a)$$

Finally, note that if M_W and M agree on all atoms, then they must also agree on all other non-zero elements of the Boolean algebra.

(\Rightarrow) For the other direction, we are given a factored weight function

$$W = \prod_{x \in U} W_x,$$

and we want to show that its induced measure M_W satisfies Eq. (7). Let $k_u, k_v \in U \cup \{\neg w \mid w \in U\}$ be such that $k_u \in \{u, \neg u\}$, $k_v \in \{v, \neg v\}$, and $u \neq v$. We then want to show that

$$M_W(k_u \wedge k_v) = M_W(k_u)M_W(k_v) \quad (8)$$

which is equivalent to

$$M_W(k_u \wedge k_v) \cdot M_W(\neg k_u \wedge \neg k_v) = M_W(k_u \wedge \neg k_v) \cdot M_W(\neg k_u \wedge k_v) \quad (9)$$

by Lemma 2. Then

$$\begin{aligned} M_W(k_u \wedge k_v) &= \sum_{\{a\} \leq k_u \wedge k_v} W(a) = \sum_{\{a\} \leq k_u \wedge k_v} \prod_{x \in U} W_x(a) \\ &= \sum_{\{a\} \leq k_u \wedge k_v} W_u(a_u)W_v(a_v) \prod_{x \in U \setminus \{u, v\}} W_x(a) = \sum_{\{a\} \leq k_u \wedge k_v} W_u(k_u)W_v(k_v) \prod_{x \in U \setminus \{u, v\}} W_x(a) \\ &= W_u(k_u)W_v(k_v) \sum_{\{a\} \leq k_u \wedge k_v} \prod_{x \in U \setminus \{u, v\}} W_x(a) = W_u(k_u)W_v(k_v)C, \end{aligned}$$

where

$$C = \sum_{\{a\} \leq k_u \wedge k_v} \prod_{x \in U \setminus \{u, v\}} W_x(a),$$

and is the same for $M_W(\neg k_u \wedge k_v)$, $M_W(k_u \wedge \neg k_v)$, and $M_W(\neg k_u \wedge \neg k_v)$ as well. But then Eq. (9) becomes

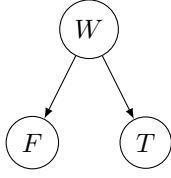
$$W_u(k_u)W_v(k_v)W_u(\neg k_u)W_v(\neg k_v)C^2 = W_u(k_u)W_v(\neg k_v)W_u(\neg k_u)W_v(k_v)C^2$$

which is trivially true. \square

5 Previous Work

5.1 Using WMC to Perform Inference on Bayesian Networks

Hitherto, four techniques have been proposed for encoding Bayesian networks into instances of WMC. We will identify them based on the initials of authors as well as publications years: **d02** [16], **sbk05** [43], **cd05** [9], and **cd06** [10]. Below we friendly summarise the observed performance differences among them. Sang et al. [43] claim that **sbk05** is a smaller encoding than **d02** with respect to both the number of clauses and the number of variables but provide no experimental comparison. Chavira and Darwiche [9] compare **cd05** with **d02** by measuring the time it takes to compile either encoding into an arithmetic circuit (but do not measure inference time). The more recent encoding **cd05** always compiles faster and results in a smaller arithmetic circuit (as measured by the number of edges). In their subsequent paper, the same authors perform two sets of experiments (that are relevant to this summary) [10]. First, they compile **cd05** and **cd06** encodings into d-DNNF (i.e., deterministic decomposable negation normal form [15]), measuring both compilation time and numbers of edges in the d-DNNF diagram. The results are mostly in favour of **cd06**. Second, they compare the inference time of **sbk05** run with Cachet [42] with the compile times of **cd05** and **cd06**, but only on five (types of) instances. In these experiments, **cd06** is always faster than **cd05**, while the comparison with **sbk05** is mixed. Sometimes **cd06** is orders of magnitude faster than **sbk05**, sometimes slightly slower. The



w	$\Pr(W = w)$	w	f	$\Pr(F = f \mid W = w)$	w	t	$\Pr(T = t \mid W = w)$
1	0.5	1	1	0.6	1	l	0.2
0	0.5	1	0	0.4	1	m	0.4
		0	1	0.1	1	h	0.4
		0	0	0.9	0	l	0.6
					0	m	0.3
					0	h	0.1

Figure 1: An example Bayesian network with its CPTs

performance difference between **sbk05** and **cd05** is even harder to judge: **sbk05** is better on three out of five instances and worse on the remaining two. Based on this description, one would expect **cd06** to be faster than both **cd05** and **sbk05**, both of which should be faster than **d02**. The experiments in Section 7, however, strongly disagree with this prediction, showing that the quality of an encoding depends strongly on the underlying search algorithm or compilation technique.

5.2 Algebraic Decision Diagrams and Their Use in Probabilistic Inference

Cover:

- ADDs provide an efficient way to manipulate functions from BAs/power sets [2].
- background reading
 - Compiling Bayesian Networks Using Variable Elimination (Chavira and Darwiche) [11]
 - On the Relationship between Sum-Product Networks and Bayesian Networks [49]
 - Using ROBDDs for Inference in Bayesian Networks with Troubleshooting as an Example [36]
 - SPUDD: Stochastic Planning using Decision Diagrams [25]

6 Encoding Bayesian Networks Using Conditional Weights

Let V denote the set of random variables in a Bayesian network. For any random variable $X \in V$, let $\text{pa}(X)$ denote the set of parents of X and $\text{im } X$ denote the set of possible values. We will use the Bayesian network in Fig. 1 (loosely based on an example by Sang et al. [43]) as a running example throughout this section. For the Bayesian network in Fig. 1, we have:

$$\begin{aligned}
V &= \{W, F, T\}, \\
\text{pa}(W) &= \emptyset, \\
\text{pa}(F) &= \text{pa}(T) = \{W\}, \\
\text{im } W &= \text{im } F = \{0, 1\}, \\
\text{im } T &= \{l, m, h\}.
\end{aligned}$$

Definition 4 (Indicator variables). Let $X \in V$ be a random variable. If X is binary (i.e., $|\text{im } X| = 2$), we can arbitrarily identify one of the values as 1 and the other one as 0 (i.e., $\text{im } X \cong \{0, 1\}$). Then X can be represented by a single *indicator variable* $\lambda_{X=1}$. For notational simplicity, for any set S , whenever we write $\lambda_{X=0} \in S$ or $S = \{\lambda_{X=0}, \dots\}$, we actually mean $\lambda_{X=1} \notin S$.

On the other hand, if X is not binary, we represent X with $|\text{im } X|$ indicator variables, one for each value. We let

$$E(X) = \begin{cases} \{\lambda_{X=1}\} & \text{if } |\text{im } X| = 2 \\ \{\lambda_{X=x} \mid x \in \text{im } X\} & \text{otherwise.} \end{cases}$$

denote the set of indicator variables for X and

$$E^*(X) = E(X) \cup \bigcup_{Y \in \text{pa}(X)} E(Y).$$

denote the set of indicator variables for X and its parents in the Bayesian network. Finally, let

$$U = \bigcup_{X \in V} E(X)$$

denote the set of all indicator variables for all random variables in the Bayesian network.

Come up with a less repetitive wording.

For the Bayesian network in Fig. 1, we have:

$$\begin{aligned} E(W) &= \{\lambda_{W=1}\}, \\ E(F) &= \{\lambda_{F=1}\}, \\ E(T) &= \{\lambda_{T=l}, \lambda_{T=m}, \lambda_{T=h}\}, \\ E^*(W) &= \{\lambda_{W=1}\}, \\ E^*(F) &= \{\lambda_{F=1}, \lambda_{W=1}\}, \\ E^*(T) &= \{\lambda_{T=l}, \lambda_{T=m}, \lambda_{T=h}, \lambda_{W=1}\}. \end{aligned}$$

Algorithm 1: Encoding a Bayesian network as a function $2^U \rightarrow \mathbb{R}_{\geq 0}$

Data: a Bayesian network with vertices V and probability distribution Pr

Result: a function $\phi: 2^U \rightarrow \mathbb{R}_{\geq 0}$

$\phi \leftarrow 1;$

for $X \in V$ **do**

let $\text{pa}(X) = \{Y_1, \dots, Y_n\};$

$\text{CPT}_X \leftarrow 0;$

if $|\text{im } X| = 2$ **then**

for $(y_1, \dots, y_n) \in \prod_{i=1}^n \text{im } Y_i$ **do**

$p_1 \leftarrow \text{Pr}(X = 1 \mid Y_1 = y_1, \dots, Y_n = y_n);$

$p_0 \leftarrow \text{Pr}(X \neq 1 \mid Y_1 = y_1, \dots, Y_n = y_n);$

$\text{CPT}_X \leftarrow \text{CPT}_X + p_1[\lambda_{X=1}] \cdot \prod_{i=1}^n [\lambda_{Y_i=y_i}] + p_0[\overline{\lambda_{X=1}}] \cdot \prod_{i=1}^n [\lambda_{Y_i=y_i}];$

else

let $\text{im } X = \{x_1, \dots, x_m\};$

for $x \in \text{im } X$ **and** $(y_1, \dots, y_n) \in \prod_{i=1}^n \text{im } Y_i$ **do**

$p_x \leftarrow \text{Pr}(X = x \mid Y_1 = y_1, \dots, Y_n = y_n);$

$\text{CPT}_X \leftarrow \text{CPT}_X + p_x[\lambda_{X=x}] \cdot \prod_{i=1}^n [\lambda_{Y_i=y_i}] + [\overline{\lambda_{X=x}}] \cdot \prod_{i=1}^n [\lambda_{Y_i=y_i}];$

$\text{CPT}_X \leftarrow \text{CPT}_X \cdot (\sum_{i=1}^m [\lambda_{X=x_i}]) \cdot \prod_{i=1}^m \prod_{j=i+1}^m ([\overline{\lambda_{X=x_i}}] + [\overline{\lambda_{X=x_j}}]);$

$\phi \leftarrow \phi \cdot \text{CPT}_X;$

return $\phi;$

Ditto.

For the Bayesian network in Fig. 1, we have:

$$\text{CPT}_W = 0.5[\lambda_{W=1}] + 0.5[\overline{\lambda_{W=1}}] = 0.5 \cdot 1,$$

$$\begin{aligned} \text{CPT}_F &= 0.6[\lambda_{F=1}] \cdot [\lambda_{W=1}] + 0.4[\lambda_{F=0}] \cdot [\lambda_{W=1}] + 0.1[\lambda_{F=1}] \cdot [\lambda_{W=0}] + 0.9[\lambda_{F=0}] \cdot [\lambda_{W=0}] \\ &= 0.6[\lambda_{F=1}] \cdot [\lambda_{W=1}] + 0.4[\overline{\lambda_{F=1}}] \cdot [\lambda_{W=1}] + 0.1[\lambda_{F=1}] \cdot [\overline{\lambda_{W=1}}] + 0.9[\overline{\lambda_{F=1}}] \cdot [\overline{\lambda_{W=1}}], \end{aligned}$$

$$\text{CPT}_T = ([\lambda_{T=l}] + [\lambda_{T=m}] + [\lambda_{T=h}]) \cdot ([\overline{\lambda_{T=l}}] + [\overline{\lambda_{T=m}}]) \cdot ([\overline{\lambda_{T=l}}] + [\overline{\lambda_{T=h}}]) \cdot ([\overline{\lambda_{T=m}}] + [\overline{\lambda_{T=h}}]) \cdot (\dots).$$

Describe Algorithm 1

6.1 Proof of Correctness

Lemma 3. *Let $X \in V$ be a random variable with parents $\text{pa}(X) = \{Y_1, \dots, Y_n\}$. Then $\text{CPT}_X : 2^{E^*(X)} \rightarrow \mathbb{R}_{\geq 0}$ is such that for any $x \in \text{im } X$ and $(y_1, \dots, y_n) \in \prod_{i=1}^n \text{im } Y_i$,*

$$\text{CPT}_X(\{\lambda_{X=x}\} \cup \{\lambda_{Y_i=y_i} \mid i = 1, \dots, n\}) = \Pr(X = x \mid Y_1 = y_1, \dots, Y_n = y_n).$$

Proof. Let $\tau = \{\lambda_{X=x}\} \cup \{\lambda_{Y_i=y_i} \mid i = 1, \dots, n\}$. If X is binary, then CPT_X is a sum of $2 \prod_{i=1}^n |\text{im } Y_i|$ terms, one for each possible assignment of values to variables X, Y_1, \dots, Y_n . Exactly one of these terms is nonzero when applied to τ , and it is equal to $\Pr(X = x \mid Y_1 = y_1, \dots, Y_n = y_n)$ by definition.

If X is not binary, then

$$\left(\sum_{i=1}^m [\lambda_{X=x_i}] \right) (\tau) = 1,$$

and

$$\left(\prod_{i=1}^m \prod_{j=i+1}^m ([\overline{\lambda_{X=x_i}}] + [\overline{\lambda_{X=x_j}}]) \right) (\tau) = 1,$$

so, by a similar argument as before,

$$\text{CPT}_X(\tau) = \Pr(X = x \mid Y_1 = y_1, \dots, Y_n = y_n).$$

□

Proposition 1. $\phi : 2^U \rightarrow \mathbb{R}_{\geq 0}$ represents the full probability distribution of the Bayesian network, i.e., if $V = \{X_1, \dots, X_n\}$, then

$$\phi(\tau) = \begin{cases} \Pr(X_1 = x_1, \dots, X_n = x_n) & \text{if } \tau = \{\lambda_{X_i=x_i} \mid i = 1, \dots, n\} \text{ for some } (x_1, \dots, x_n) \in \prod_{i=1}^n \text{im } X_i \\ 0 & \text{otherwise,} \end{cases}$$

for all $\tau \in 2^U$.

Proof. If $\tau = \{\lambda_{X=v_X} \mid X \in V\}$ for some $(v_X)_{X \in V} \in \prod_{X \in V} \text{im } X$, then

$$\phi(\tau) = \prod_{X \in V} \Pr \left(X = v_X \mid \bigwedge_{Y \in \text{pa}(X)} Y = v_Y \right) = \Pr \left(\bigwedge_{X \in V} X = v_X \right)$$

by Lemma 3 and the definition of a Bayesian network. Otherwise there must be some non-binary random variable $X \in V$ such that $|E(X) \cap \tau| \neq 1$. If $E(X) \cap \tau = \emptyset$, then

$$\left(\sum_{i=1}^m [\lambda_{X=x_i}] \right) (\tau) = 0,$$

and so $\text{CPT}_X(\tau) = 0$, and $\phi(\tau) = 0$. If $|E(X) \cap \tau| > 1$, then we must have two different values $x_1, x_2 \in \text{im } X$ such that $\{\lambda_{X=x_1}, \lambda_{X=x_2}\} \subseteq \tau$ which means that

$$(\overline{[\lambda_{X=x_1}]} + \overline{[\lambda_{X=x_2}]}) (\tau) = 0,$$

and so, again, $\text{CPT}_X(\tau) = 0$, and $\phi(\tau) = 0$. □

Theorem 2. *Let $\phi: 2^U \rightarrow \mathbb{R}_{\geq 0}$ be a function generated by the algorithm. Then*

$$(\exists_U(\phi \cdot [\lambda_{X=x}])(\emptyset) = \Pr(X = x).$$

Proof. Let $V = \{X, Y_1, \dots, Y_n\}$. Then

$$\begin{aligned} (\exists_U(\phi \cdot [\lambda_{X=x}])(\emptyset) &= \sum_{\tau \in 2^U} (\phi \cdot [\lambda_{X=x}])(\tau) = \sum_{\lambda_{X=x} \in \tau \in 2^U} \phi(\tau) = \sum_{\lambda_{X=x} \in \tau \in 2^U} \left(\prod_{Y \in V} \text{CPT}_Y \right) (\tau) \\ &= \sum_{(y_1, \dots, y_n) \in \prod_{i=1}^n \text{im } Y_i} \Pr(X = x, Y_1 = y_1, \dots, Y_n = y_n) = \Pr(X = x) \end{aligned}$$

by the following arguments:

- the proof of Theorem 1 in the ADDMC paper [19];
 - if $\lambda_{X=x} \notin \tau \in 2^U$, then $(\phi \cdot [\lambda_{X=x}])(\tau) = \phi(\tau) \cdot [\lambda_{X=x}](\tau \cap \{\lambda_{X=x}\}) = \phi(\tau) \cdot 0 = 0$;
 - Proposition 1;
 - marginalisation of a probability distribution.
-

6.2 Textual Representation

The Bayesian network in Fig. 1 can be represented in a textual format as

$\lambda_{T=l}$	$\lambda_{T=m}$	$\lambda_{T=h}$	0	
	$-\lambda_{T=l}$	$-\lambda_{T=m}$	0	
	$-\lambda_{T=l}$	$-\lambda_{T=h}$	0	
	$-\lambda_{T=m}$	$-\lambda_{T=h}$	0	
w	$\lambda_{W=1}$		0.5	0.5
w	$\lambda_{F=1}$	$\lambda_{W=1}$	0.6	0.4
w	$\lambda_{F=1}$	$-\lambda_{W=1}$	0.1	0.9
w	$\lambda_{T=l}$	$\lambda_{W=1}$	0.2	1
w	$\lambda_{T=m}$	$\lambda_{W=1}$	0.4	1
w	$\lambda_{T=h}$	$\lambda_{W=1}$	0.4	1
w	$\lambda_{T=l}$	$\lambda_{W=0}$	0.6	1
w	$\lambda_{T=m}$	$\lambda_{W=0}$	0.3	1
w	$\lambda_{T=h}$	$\lambda_{W=0}$	0.1	1

with each λ replaced with a unique positive integer. This format is based on the format used by the Cachet solver [42] to encode WMC problems, which extends the DIMACS format for CNF formulas with weight clauses. Subsequently, we extend it in two ways:

- a single weight clause now supports an arbitrary number of literals,
- and each weight clause has two probabilities instead of one (i.e., we no longer assume that $\Pr(v) + \Pr(\neg v) = 1$ for all variables $v \in U$).

The way we use this encoding, it is always the case that either both probabilities sum to one, or the second probability (i.e., the probability for the complement of the variable) is equal to one.

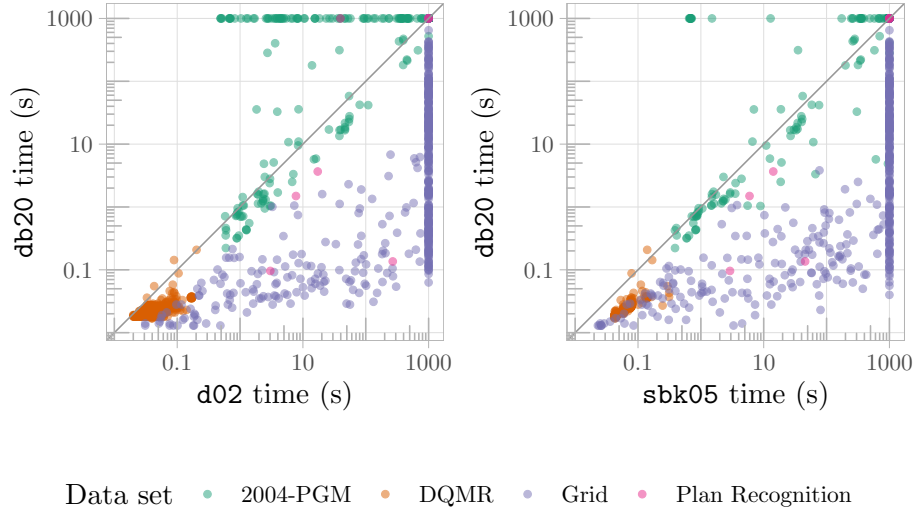


Figure 2: Scatter plots

6.3 Changes to ADDMC

ADDMC constructs the *Gaifman graph* [21] of the input CNF formula as an aid for the algorithm’s heuristics. This graph has as vertices the variables of the formula, and there is an edge between two variables u and v if there is a clause in the formula that contains both u and v . We extend this definition to functions on Boolean algebras, i.e., the factors of ϕ . For any pair of distinct variables $u, v \in U$, we draw an edge between them in the Gaifman graph if there is a function $A: 2^X \rightarrow \mathbb{R}_{\geq 0}$ that is a factor of ϕ such that $u \in X$ and $v \in X$. For instance, a factor such as CPT_X will enable edges between all distinct pairs of variables in $E^*(X)$.

This needs rewriting.

ADDMC multiplies the final answer by $w(u) + w(-u)$ for every $u \in U$ that was ‘simplified out’ of the ADD for ϕ and so is not featured in the \exists_U projection. In our format, for every such $u \in U$, we multiply the answer by two. Each such $u \in U$ must satisfy two properties:

- all the probabilities/weights associated with u are equal to 0.5 (otherwise the corresponding CPT would depend on u),
- and all other CPTs are independent of u (or they may have a trivial dependence, where the probability stays the same if u is replaced with its complement).

Thus, the CPT that corresponds to u still multiplies every model by 0.5, but the number of models being considered by the ADDMC is halved. To correct for this, we need to perform this multiplication.

7 Experimental Comparison

- We don’t compare ‘compile times’ because our encoding time is linear, so we would easily beat everyone else.
- When the Bayesian network has an evidence file, we compute the probability of evidence. Otherwise, let X denote the last-mentioned node in the Bayesian network. If `true` is a valid value of X , we compute the marginal probability of $X = \text{true}$. Otherwise, we pick the first value of X and calculate

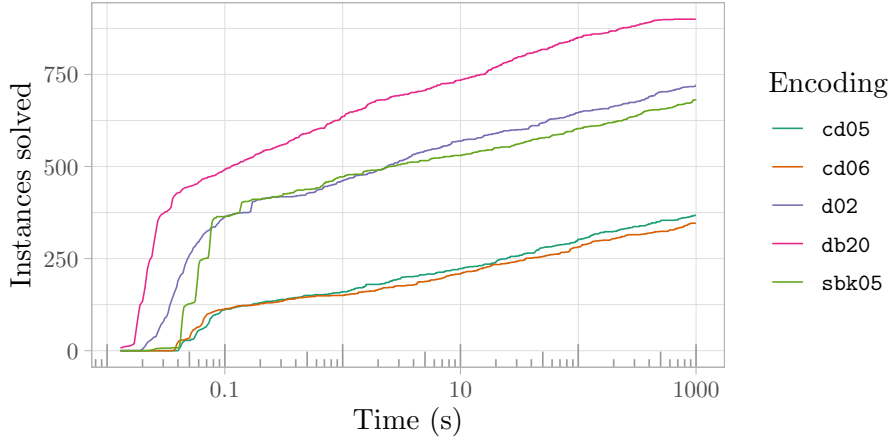


Figure 3: Cumulative plot

Table 3: Out of 1216 instances

Encoding	Unique	Fastest	Solved
cd05	2	3	372
cd06	0	1	351
d02	41	99	726
db20	228	871	901
sbk05	6	36	687

its marginal probability. This applies to the Grid data set (as intended) and also to two instances of Plan Reconstruction and roughly half of the instances from 2004-PGM that have empty evidence files.

- After the experiments are finished, note the processor, memory per thread.
- The experiments were run on Intel Xeon Gold 6138 processor with an 8 GB memory limit.
- I will call my encoding **db20**.
- All other encodings are implemented in Ace 3.0² and were compiled with `-encodeOnly` (i.e., don't compile the CNF into an AC) and `-noEclause` (i.e., only use standard syntax) flags.
- We don't use Ace's evidence-encoding mechanism because it has multiple significant bugs (and the code is closed-source). Ace's implementation (depending on the exact encoding) is supposed to perform some simplifications when encoding evidence. We just do the trivial thing—no simplifications.
- **cd05** relaxes the encoding so much that extra models become possible. They are supposed to be filtered out by the algorithm, but mine can't do that because it doesn't deal with models. Same for **cd06** because it's based on **cd05**.

Data. For experiments, we use the Bayesian networks available with Ace³ and Cachet⁴, most of which happen to be binary. We classify them into the following seven categories: • DQMR and • grid networks as described by Sang et al. [43], • friends and smokers, • mastermind, and • random blocks from the work

²<http://reasoning.cs.ucla.edu/ace/>

³<http://reasoning.cs.ucla.edu/ace/>

⁴<https://www.cs.rochester.edu/u/kautz/Cachet/>

of Chavira et al. [13], • remaining binary Bayesian networks that include plan recognition [43], students and professors [13], and `tcc4f`, and • non-binary classic Bayesian networks (`alarm`, `diabetes`, `hailfinder`, `mildew`, `munin1-4`, `pathfinder`, `pigs`, `water`).

8 Explaining The Performance Benefits

- `d02` has

$$\sum_{X \in V} |\text{im } X| + |\text{im } X| \prod_{Y \in \text{pa}(X)} |\text{im } Y|$$

variables and

$$\sum_{X \in V} 1 + \binom{|\text{im } X|}{2} + |\text{im } X|(2 + |\text{pa}(X)|) \prod_{Y \in \text{pa}(X)} |\text{im } Y|$$

clauses (along with one ADD per variable to encode the weights).

- `sbk05` is a bit harder to evaluate due to a handful of small optimisations in the encoding. Could find an upper bound anyway.
- `db20` (my encoding) has

$$\sum_{X \in V} |\text{im } X|$$

variables (less for binary) and

$$\sum_{X \in V} |\text{im } X| + 1 + \binom{|\text{im } X|}{2}$$

ADDs.

- Let:
 - $N = |V|$ (i.e., the number of nodes in the Bayesian network),
 - $D = \max_{X \in V} |\text{pa}(X)|$ (i.e., the maximum in-degree or the number of parents),
 - $V = \max_{X \in V} |\text{im } X|$ (i.e., the maximum number of values per variables).
- Then my encoding has $\mathcal{O}(NV)$ variables and $\mathcal{O}(NV^2)$ ADDs while `d02` has $\mathcal{O}(NV^{D+1})$ variables and $\mathcal{O}(NDV^{D+1})$ ADDs.

Calculate `numVariables/numClauses` (or the other way around) for each instance and plot this ratio vs runtime (for each encoding, or at least mine and `D02`). The new CP paper kind of beat me to it...

9 Conclusion and Future Work

- Bayesian networks and ADDMC are only particular examples. This should also work with Cachet.
- Extra benefit: one does not need to come up with a way to turn some probability distribution to into a fully independent one.
- Important future work: replacing ADDs with AADDs⁵ [44] is likely to bring performance benefits. Other extensions:
 - FOADDs can represent first order statements;

⁵<https://github.com/ssanner/dd-inference>

- XADDs can replace WMI for continuous variables;
- ADDs with intervals can do approximations.
- Filtering out ADDs that have nothing to do with the answer helps tremendously, but I’m purposefully not doing that. Perhaps a heuristic could do the same thing?
- Encodings for everything else
 - probabilistic programs [26]
 - ProbLog [20]
 - * For the ProbLog to WMC conversion, check out this guy: <https://users.ics.aalto.fi/ttj/>.
 - * proof-based [35]
 - * rule-based [28]
 - * For ground ProbLog, we can encode a program


```
p :: a :- b
q :: a :- c
```

 into $P(a \mid b) = p, P(a \mid c) = q$ instead of having clauses $b \Rightarrow a, c \Rightarrow a$. Some logical structure is likely to remain.
- Bayesian networks are often solved in a compile once, query many times fashion. This can be achieved using ADDMC by selecting a subset S of variable we may want to query over and running ADDMC while excluding S from variable elimination/projection/ \exists .
- More references
 - Measures on/in Boolean algebras: Horn and Tarski [27], Jech [29]
 - On Boolean algebras and their role in analysis [47]
 - Infinite domains
 - * Markov Logic in Infinite Domains (Singla and Domingos) [45]
 - * Objective Bayesian probabilistic logic (Williamson) [46]
 - * Unifying Logic and Probability (Russell) [41]
 - Logical induction [22]
 - Quantum probabilistic logic programming [3]
 - WMC
 - * algebraic model counting [31]
 - * Explanation-Based Approximate Weighted Model Counting for Probabilistic Logics [40]
 - * OUWMC [4]
 - * Formula-Based Probabilistic Inference [23]
 - * Parallel Probabilistic Inference by WMC [14]
 - * Semiring Programming [6]
 - * theoretical extension: WMC beyond two-variable logic [33]
 - * from weighted to unweighted model counting [8]
 - * theory behind WMC algorithms: solving $\#SAT$ and Bayesian inference with backtracking search [1]

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