Declarative Programming

- Knowledgebases and queries in propositional logic are made up of propositions and connectives.
- Predicate logic adds the notion of predicates and variables.
- We take a non-theoretical approach to predicate logic by introducing declarative programming.
- Declarative programming is the use of mathematical logic to describe the logic of computation without describing its control flow.
- Useful for: expert systems, diagnostic systems, machine learning, parsing text, theorem proving, ...

Declarative Programming

- Programmer gives a declarative specification of the problem, using the language of logic.
- The programmer should not have to tell the computer what to do.
- To get information, the programmer simply asks a query.

Datalog

- Prolog is a declarative (logical) programming language and stands for PROgramming in LOGic
- We only look at a subset of the language which is (roughly) equivalent to Datalog.

http://www.learnprolognow.org



Basic idea of Datalog/Prolog

- Describe the situation of interest
- Ask a question
- Prolog logically deduces new facts about the situation we described
- Prolog gives us its deductions back as answers

Consequences

- Think declaratively, not procedurally
 - Challenging
 - Requires a different mindset
 - Has similarities with other programming paradigms such as functional programming.

- High-level language
 - Not as efficient
 - Useful in many Al applications

The interpreter

- Prolog has an interactive interpreter
- After starting the interpreter, it can start reading your Prolog files and answer your queries.
- To exit Prolog simply type the command halt. (note the full-stop)
- Prolog program files usually have the extension .pl
 or .pro

Where is the program written?

- Facts and Rules are stored in one or more files forming our Knowledge Base
- Files containing KB are loaded into the interpreter
- After changing these files, the files should be loaded again to be effective
- Queries are asked in the interactive mode in front of the question prompt: ?-

Reading Files

- consult (filename). Reads and compiles a Prolog source file. Example: - consult ('/home/user/prolog/sample.pl'). reconsult (filename). Reconsult a changed source files. Example reconsult ('/home/user/prolog/sample.pl'). • [\filename']. (valid but don't use it; can be confused with lists!) ['/home/user/prolog/sample.pl'].
 - make.
 - Reconsult all changed source files.

```
woman(mia).
woman(jody).
woman(yolanda).
playsAirGuitar(jody).
party.
```

```
woman(mia).
woman(jody).
woman(yolanda).
playsAirGuitar(jody).
party.
```



```
woman(mia).
woman(jody).
woman(yolanda).
playsAirGuitar(jody).
party.
```

?- woman(mia).

```
woman(mia).
woman(jody).
woman(yolanda).
playsAirGuitar(jody).
party.
```

```
?- woman(mia).
yes
?-
```

```
woman(mia).
woman(jody).
woman(yolanda).
playsAirGuitar(jody).
party.
```

```
?- woman(mia).yes?- playsAirGuitar(jody).
```

```
woman(mia).
woman(jody).
woman(yolanda).
playsAirGuitar(jody).
party.
```

```
?- woman(mia).
yes
?- playsAirGuitar(jody).
yes
?-
```

```
woman(mia).
woman(jody).
woman(yolanda).
playsAirGuitar(jody).
party.
```

```
?- woman(mia).
yes
?- playsAirGuitar(jody).
yes
?- playsAirGuitar(mia).
no
```

```
woman(mia).
woman(jody).
woman(yolanda).
playsAirGuitar(jody).
party.
```

?- tattoed(jody).

```
woman(mia).
woman(jody).
woman(yolanda).
playsAirGuitar(jody).
party.
```

```
?- tattoed(jody).
no
?-
```

```
woman(mia).
woman(jody).
woman(yolanda).
playsAirGuitar(jody).
party.
```

?- tattoed(jody).
ERROR: predicate tattoed/1 not defined.
?-

```
woman(mia).
woman(jody).
woman(yolanda).
playsAirGuitar(jody).
party.
```

?- party.

```
woman(mia).
woman(jody).
woman(yolanda).
playsAirGuitar(jody).
party.
```

```
?- party.
yes
?-
```

```
woman(mia).
woman(jody).
woman(yolanda).
playsAirGuitar(jody).
party.
```

?- rockConcert.

```
woman(mia).
woman(jody).
woman(yolanda).
playsAirGuitar(jody).
party.
```

```
?- rockConcert.
no
?-
```

```
happy(yolanda).
```

listens2music(mia).

listens2music(yolanda):- happy(yolanda).

playsAirGuitar(mia):- listens2music(mia).

playsAirGuitar(yolanda):- listens2music(yolanda).

happy(yolanda). fact

listens2music(mia).

listens2music(yolanda):- happy(yolanda).

playsAirGuitar(mia):- listens2music(mia).

playsAirGuitar(yolanda):- listens2music(yolanda).

```
happy(yolanda). fact
listens2music(mia). fact
listens2music(yolanda):- happy(yolanda).
playsAirGuitar(mia):- listens2music(mia).
playsAirGuitar(yolanda):- listens2music(yolanda).
```

happy(yolanda).

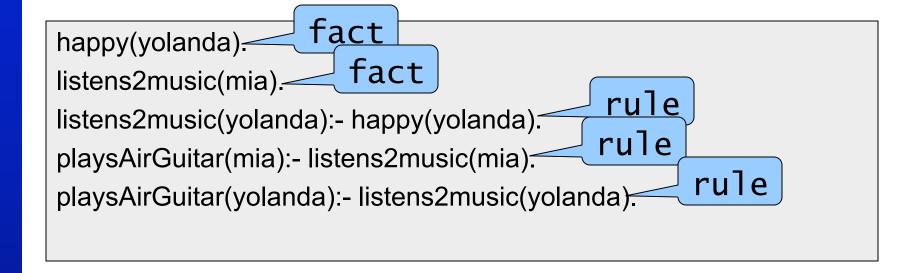
listens2music(mia).

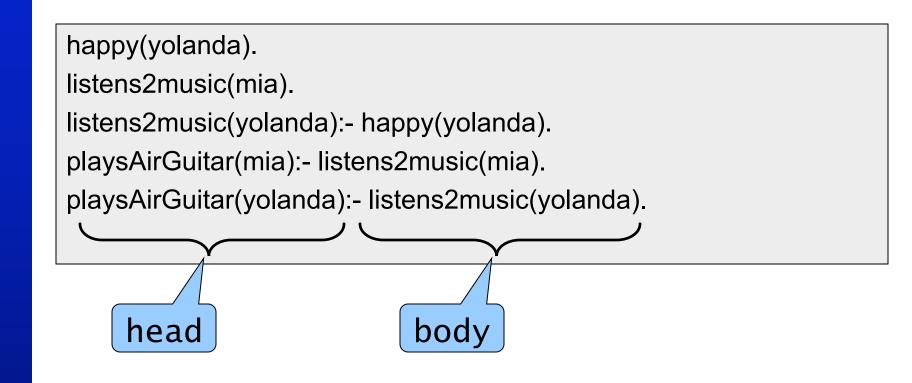
listens2music(yolanda):- happy(yolanda).

playsAirGuitar(mia):- listens2music(mia).

playsAirGuitar(yolanda):- listens2music(yolanda).

happy(yolanda):- fact
listens2music(mia):- happy(yolanda):- rule
playsAirGuitar(mia):- listens2music(mia):- rule
playsAirGuitar(yolanda):- listens2music(yolanda).





```
happy(yolanda).
listens2music(mia).
listens2music(yolanda):- happy(yolanda).
playsAirGuitar(mia):- listens2music(mia).
playsAirGuitar(yolanda):- listens2music(yolanda).
```

?-

```
happy(yolanda).
listens2music(mia).
listens2music(yolanda):- happy(yolanda).
playsAirGuitar(mia):- listens2music(mia).
playsAirGuitar(yolanda):- listens2music(yolanda).
```

```
?- playsAirGuitar(mia).
yes
?-
```

```
happy(yolanda).
listens2music(mia).
listens2music(yolanda):- happy(yolanda).
playsAirGuitar(mia):- listens2music(mia).
playsAirGuitar(yolanda):- listens2music(yolanda).
```

```
?- playsAirGuitar(mia).yes?- playsAirGuitar(yolanda).yes
```

Clauses

happy(yolanda).

listens2music(mia).

listens2music(yolanda):- happy(yolanda).

playsAirGuitar(mia):- listens2music(mia).

playsAirGuitar(yolanda):- listens2music(yolanda).

There are five clauses in this knowledge base: two facts, and three rules.

The end of a clause is marked with a full stop.

Predicates

happy(yolanda).

listens2music(mia).

listens2music(yolanda):- happy(yolanda).

playsAirGuitar(mia):- listens2music(mia).

playsAirGuitar(yolanda):- listens2music(yolanda).

There are three **predicates**in this knowledge base:
happy, listens2music, and playsAirGuitar

```
happy(vincent).
listens2music(butch).
playsAirGuitar(vincent):- listens2music(vincent), happy(vincent).
playsAirGuitar(butch):- happy(butch).
playsAirGuitar(butch):- listens2music(butch).
```

Expressing Conjunction

```
happy(vincent).
```

listens2music(butch).

playsAirGuitar(vincent):- listens2music(vincent), happy(vincent).

playsAirGuitar(butch):- happy(butch).

playsAirGuitar(butch):- listens2music(butch).

The comma "," expresses conjunction in Prolog

```
happy(vincent).
listens2music(butch).
playsAirGuitar(vincent):- listens2music(vincent), happy(vincent).
playsAirGuitar(butch):- happy(butch).
playsAirGuitar(butch):- listens2music(butch).
```

```
?- playsAirGuitar(vincent).
no
?-
```

```
happy(vincent).
listens2music(butch).
playsAirGuitar(vincent):- listens2music(vincent), happy(vincent).
playsAirGuitar(butch):- happy(butch).
playsAirGuitar(butch):- listens2music(butch).
```

```
?- playsAirGuitar(butch).
yes
?-
```

Expressing Disjunction

```
happy(vincent).
listens2music(butch).
playsAirGuitar(vincent):- listens2music(vincent), happy(vincent).
playsAirGuitar(butch):- happy(butch).
playsAirGuitar(butch):- listens2music(butch).
```

```
happy(vincent).
listens2music(butch).
playsAirGuitar(vincent):- listens2music(vincent), happy(vincent).
playsAirGuitar(butch):- happy(butch); listens2music(butch).
```

Prolog and Logic

- Clearly Prolog has something to do with logic
- Operators
 - Implication :-
 - Conjunction ,
 - Disjunction ;
- Use of modus ponens

```
woman(mia).
woman(jody).
woman(yolanda).

loves(vincent, mia).
loves(marsellus, mia).
loves(pumpkin, honey_bunny).
loves(honey_bunny, pumpkin).
```

Prolog Variables

```
woman(mia).
woman(jody).
woman(yolanda).

loves(vincent, mia).
loves(marsellus, mia).
loves(pumpkin, honey_bunny).
loves(honey_bunny, pumpkin).
```

?- woman(X).

Variable Instantiation

```
woman(mia).
woman(jody).
woman(yolanda).

loves(vincent, mia).
loves(marsellus, mia).
loves(pumpkin, honey_bunny).
loves(honey_bunny, pumpkin).
```

```
?- woman(X).
X=mia
```

```
woman(mia).
woman(jody).
woman(yolanda).

loves(vincent, mia).
loves(marsellus, mia).
loves(pumpkin, honey_bunny).
loves(honey_bunny, pumpkin).
```

```
?- woman(X).
X=mia;
```

```
woman(mia).
woman(jody).
woman(yolanda).

loves(vincent, mia).
loves(marsellus, mia).
loves(pumpkin, honey_bunny).
loves(honey_bunny, pumpkin).
```

```
?- woman(X).
X=mia;
X=jody
```

```
woman(mia).
woman(jody).
woman(yolanda).

loves(vincent, mia).
loves(marsellus, mia).
loves(pumpkin, honey_bunny).
loves(honey_bunny, pumpkin).
```

```
?- woman(X).
X=mia;
X=jody;
X=yolanda
```

```
woman(mia).
woman(jody).
woman(yolanda).

loves(vincent, mia).
loves(marsellus, mia).
loves(pumpkin, honey_bunny).
loves(honey_bunny, pumpkin).
```

```
?- woman(X).
X=mia;
X=jody;
X=yolanda;
no
```

```
woman(mia).
woman(jody).
woman(yolanda).

loves(vincent, mia).
loves(marsellus, mia).
loves(pumpkin, honey_bunny).
loves(honey_bunny, pumpkin).
```

?- loves(marsellus,X), woman(X).

```
woman(mia).
woman(jody).
woman(yolanda).

loves(vincent, mia).
loves(marsellus, mia).
loves(pumpkin, honey_bunny).
loves(honey_bunny, pumpkin).
```

```
?- loves(marsellus,X), woman(X).
X=mia
?-
```

```
woman(mia).
woman(jody).
woman(yolanda).

loves(vincent, mia).
loves(marsellus, mia).
loves(pumpkin, honey_bunny).
loves(honey_bunny, pumpkin).
```

?- loves(pumpkin,X), woman(X).

```
woman(mia).
woman(jody).
woman(yolanda).

loves(vincent, mia).
loves(marsellus, mia).
loves(pumpkin, honey_bunny).
loves(honey_bunny, pumpkin).
```

```
?- loves(pumpkin,X), woman(X).
no
?-
```

loves(vincent,mia).
loves(marsellus,mia).
loves(pumpkin, honey_bunny).

loves(honey_bunny, pumpkin).

jealous(X,Y):-loves(X,Z), loves(Y,Z).

```
loves(vincent,mia).
loves(marsellus,mia).
loves(pumpkin, honey_bunny).
loves(honey_bunny, pumpkin).

jealous(X,Y):- loves(X,Z), loves(Y,Z).
```

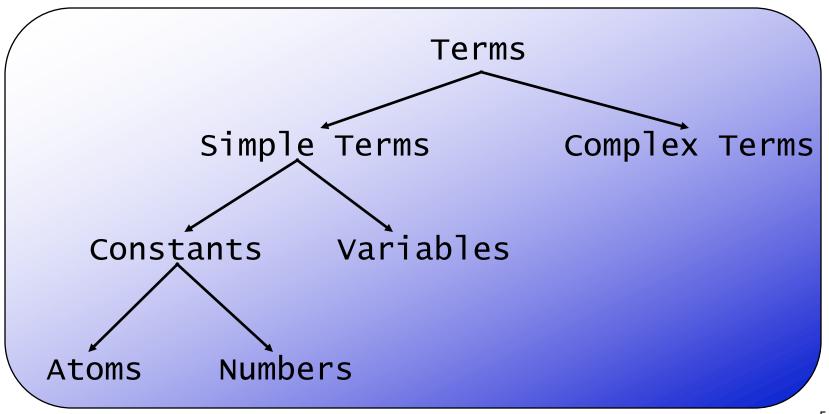
?- jealous(marsellus,W).

```
loves(vincent,mia).
loves(marsellus,mia).
loves(pumpkin, honey_bunny).
loves(honey_bunny, pumpkin).
jealous(X,Y):- loves(X,Z), loves(Y,Z).
```

```
?- jealous(marsellus,W).
W=vincent
?-
```

Prolog Syntax

 What exactly are facts, rules and queries built out of?



Atoms

- A sequence of characters of upper-case letters, lower-case letters, digits, or underscore, starting with a lowercase letter
 - Examples: butch, big_kahuna_burger, playGuitar
- An arbitrary sequence of characters enclosed in single quotes
 - Examples: 'Vincent', 'Five dollar shake', '@\$%'
- A sequence of special characters
 - Examples: : , ; . :-

Numbers

• Integers: 12, -34, 22342

• Floats: 34573.3234

Variables

 A sequence of characters of upper-case letters, lower-case letters, digits, or underscore, starting with either an uppercase letter or an underscore

Examples:

X, Y, Variable, Vincent, _tag

Complex Terms

- Atoms, numbers and variables are building blocks for complex terms
- Complex terms are built out of a functor directly followed by a sequence of arguments
- Arguments are put in round brackets, separated by commas
- The functor must be an atom

Examples of complex terms

- Examples we have seen before:
 - playsAirGuitar(jody)
 - loves(vincent, mia)
 - jealous(marsellus, W)

- Complex terms inside complex terms:
 - hide(X,father(father(father(butch))))

Arity

 The number of arguments a complex term has is called its <u>arity</u>

Examples:

woman(mia) is a term with arity 1
loves(vincent,mia) has arity 2
father(father(butch)) arity 1

Arity is important

- In Prolog you can define two predicates with the same functor but with different arity
- Prolog would treat this as two different predicates
- In Prolog documentation arity of a predicate is usually indicated with the suffix "/" followed by a number to indicate the arity

Example of Arity

```
happy(yolanda).
```

listens2music(mia).

listens2music(yolanda):- happy(yolanda).

playsAirGuitar(mia):- listens2music(mia).

playsAirGuitar(yolanda):- listens2music(yolanda).

- This knowledge base defines
 - happy/1
 - listens2music/1
 - playsAirGuitar/1

- Working definition:
 - Two terms unify if they are the same term or if they contain variables that can be uniformly instantiated with terms in such a way that the resulting terms are equal.

- This means that:
 - mia and mia unify
 - 42 and 42 unify
 - woman(mia) and woman(mia) unify

- This also means that:
 - vincent and mia do not unify
 - woman(mia) and woman(jody) do not unify

- What about the terms:
 - mia and X

- What about the terms:
 - mia and X
 - woman(Z) and woman(mia)

- What about the terms:
 - mia and X
 - woman(Z) and woman(mia)
 - loves(mia,X) and loves(X,vincent)

Instantiations

- When Prolog unifies two terms it performs all the necessary instantiations, so that the terms are equal afterwards
- This makes unification a powerful programming mechanism

Revised Definition 1/3

If T₁ and T₂ are constants, then
 T₁ and T₂ unify if they are the same atom, or the same number.

Revised Definition 2/3

- If T₁ and T₂ are constants, then
 T₁ and T₂ unify if they are the same atom, or the same number.
- 2. If T₁ is a variable and T₂ is any type of term, then T₁ and T₂ unify, and T₁ is instantiated to T₂. (and vice versa)

Revised Definition 3/3

- If T₁ and T₂ are constants, then
 T₁ and T₂ unify if they are the same atom, or
 the same number.
- 2. If T₁ is a variable and T₂ is any type of term, then T₁ and T₂ unify, and T₁ is instantiated to T₂. (and vice versa)
- 3. If T₁ and T₂ are complex terms then they unify if:
 - They have the same functor and arity, and
 - 2. all their corresponding arguments unify, and
 - 3. the variable instantiations are compatible.

Prolog unification: =/2

?- mia = mia.

yes ?-

Prolog unification: =/2

```
?- mia = mia.
```

yes

?- mia = vincent.

no

?-

Prolog unification: =/2

?- mia = X. X=mia yes ?-

How will Prolog respond?

?- X=mia, X=vincent.

How will Prolog respond?

?- X=mia, X=vincent.

no

?_

Why? After working through the first goal, Prolog has instantiated X with **mia**, so that it cannot unify it with **vincent** anymore. Hence the second goal fails.

?-
$$k(s(g),Y) = k(X,t(k))$$
.

```
?- k(s(g),Y) = k(X,t(k)).
X=s(g)
Y=t(k)
yes
?-
```

?-
$$k(s(g),t(k)) = k(X,t(Y))$$
.

```
?- k(s(g),t(k)) = k(X,t(Y)).
X=s(g)
Y=k
yes
?-
```

One last example

?-loves(X,X) = loves(marsellus,mia).

?-

```
?- vertical(line(point(1,1),point(1,3))).
yes
?-
```

```
?- vertical(line(point(1,1),point(1,3))).

yes
?- vertical(line(point(1,1),point(3,2))).

no
?-
```

```
?- horizontal(line(point(1,1),point(1,Y))).
Y = 1;
no
?-
```

```
?- horizontal(line(point(2,3),Point)).
Point = point(_554,3);
no
?-
```

Proof Search

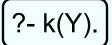
- Now that we know about unification, we are in a position to learn how Prolog searches a knowledge base to see if a query is satisfied.
- In other words: we are ready to learn about <u>proof</u> search
- Prolog has a specific way of answering queries:
 - Search knowledge base from top to bottom
 - Processes clauses from left to right
 - Backtracking to recover from bad choices

Example

```
f(a).
f(b).
g(a).
g(b).
h(b).
k(X):- f(X), g(X), h(X).
```

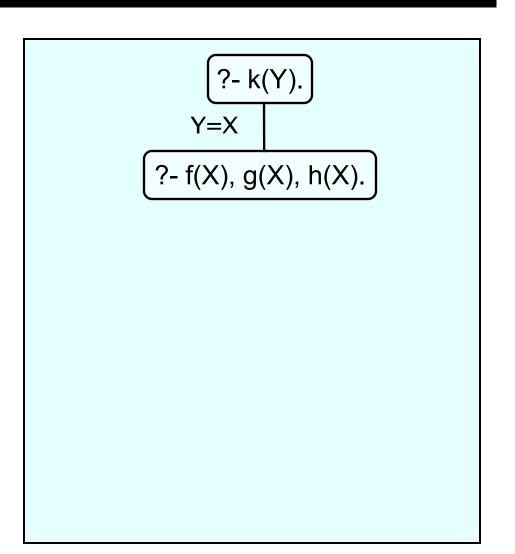
```
?- k(Y).
```

```
f(a).
f(b).
g(a).
g(b).
h(b).
k(X):- f(X), g(X), h(X).
```

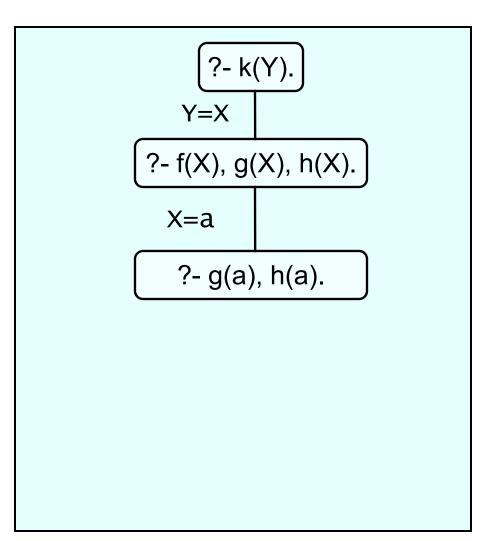


```
f(a).
f(b).
g(a).
g(b).
h(b).
k(X):- f(X), g(X), h(X).
```

```
?- k(Y).
```

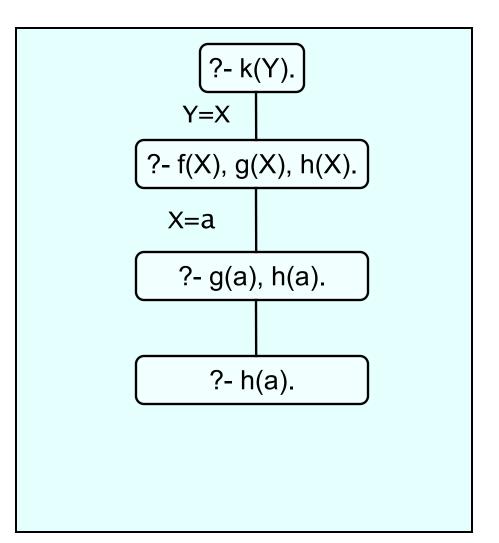


```
f(a).
f(b).
g(a).
g(b).
h(b).
k(X):- f(X), g(X), h(X).
```

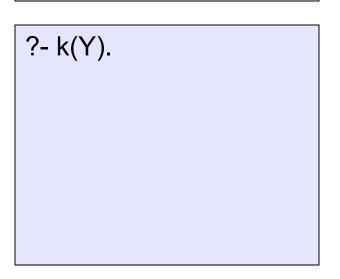


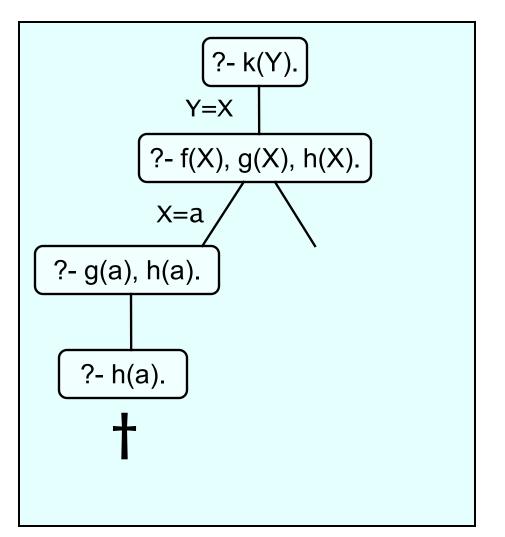
```
f(a).
f(b).
g(a).
g(b).
h(b).
k(X):- f(X), g(X), h(X).
```



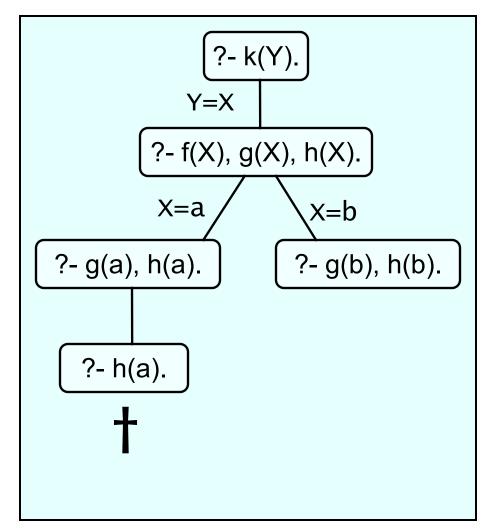


```
f(a).
f(b).
g(a).
g(b).
h(b).
k(X):- f(X), g(X), h(X).
```

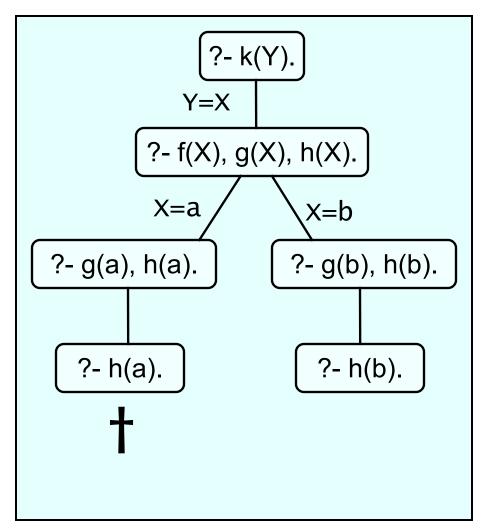




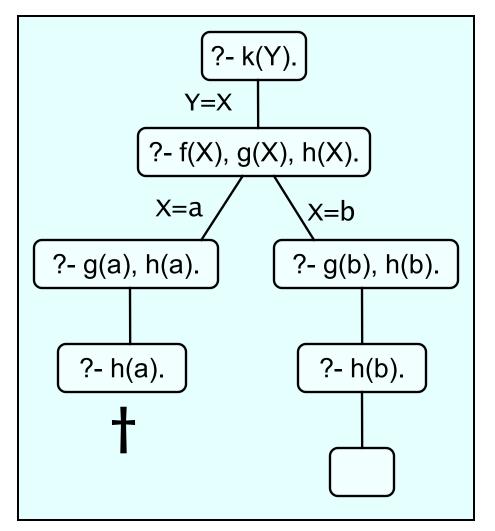
```
f(a).
f(b).
g(a).
g(b).
h(b).
k(X):- f(X), g(X), h(X).
```



```
f(a).
f(b).
g(a).
g(b).
h(b).
k(X):- f(X), g(X), h(X).
```

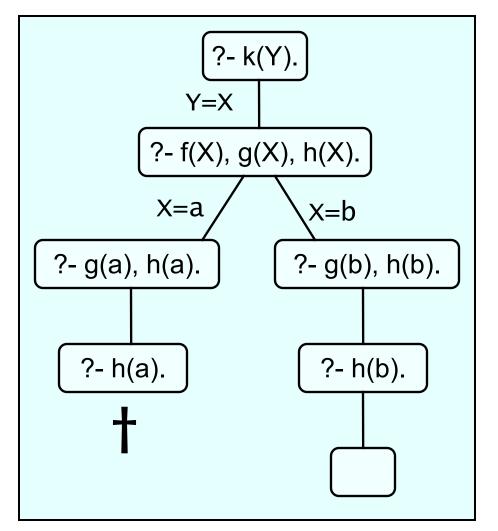


```
f(a).
f(b).
g(a).
g(b).
h(b).
k(X):- f(X), g(X), h(X).
```



```
f(a).
f(b).
g(a).
g(b).
h(b).
k(X):- f(X), g(X), h(X).
```

```
?- k(Y).
Y=b;
no
?-
```



loves(vincent,mia). loves(marsellus,mia).

jealous(A,B):loves(A,C), loves(B,C).

?- jealous(X,Y).

loves(vincent,mia). loves(marsellus,mia).

jealous(A,B):loves(A,C), loves(B,C).

?- jealous(X,Y).

?- jealous(X,Y).

loves(vincent,mia). loves(marsellus,mia).

jealous(A,B):loves(A,C), loves(B,C).

?- jealous(X,Y).

?- jealous(X,Y).

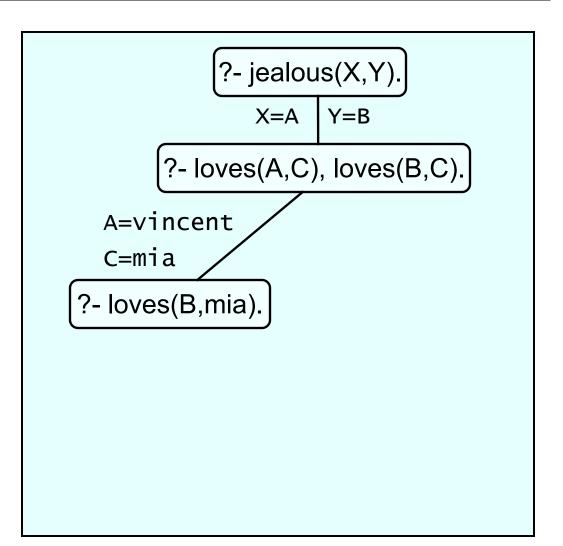
X=A Y=B

?- loves(A,C), loves(B,C).

loves(vincent,mia). loves(marsellus,mia).

jealous(A,B):loves(A,C), loves(B,C).

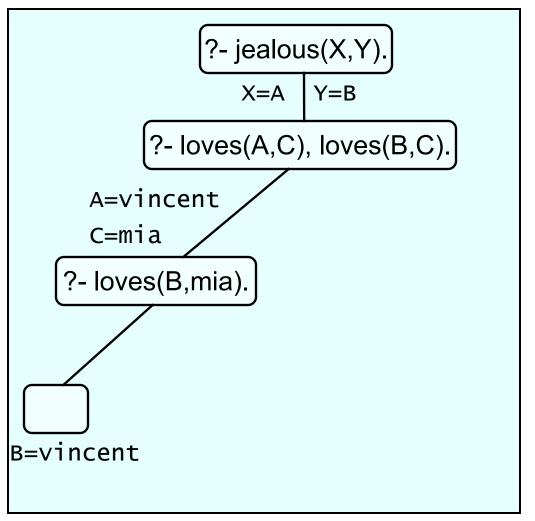
?- jealous(X,Y).



loves(vincent,mia). loves(marsellus,mia).

jealous(A,B):loves(A,C), loves(B,C).

?- jealous(X,Y). X=vincent Y=vincent



loves(vincent,mia). loves(marsellus,mia).

jealous(A,B):loves(A,C),

loves(B,C).

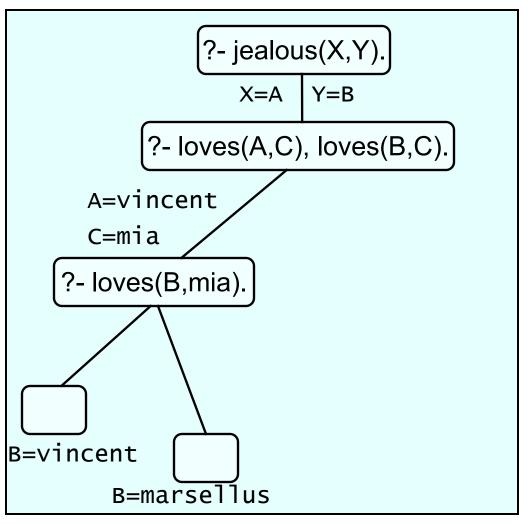
?- jealous(X,Y).

X=vincent

Y=vincent;

X=vincent

Y=marsellus



loves(vincent,mia). loves(marsellus,mia).

jealous(A,B):-

loves(A,C), loves(B,C).

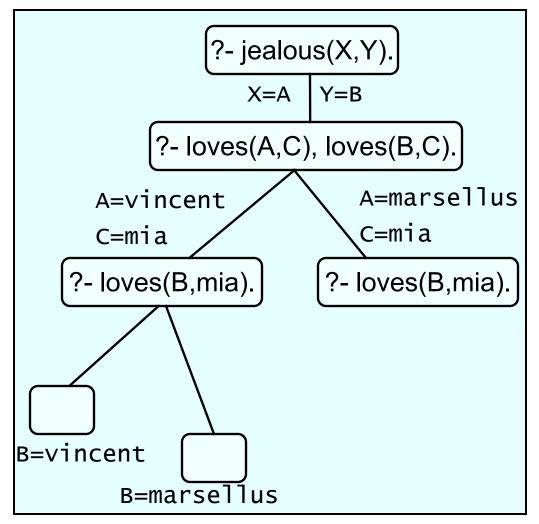
?- jealous(X,Y).

X=vincent

Y=vincent;

X=vincent

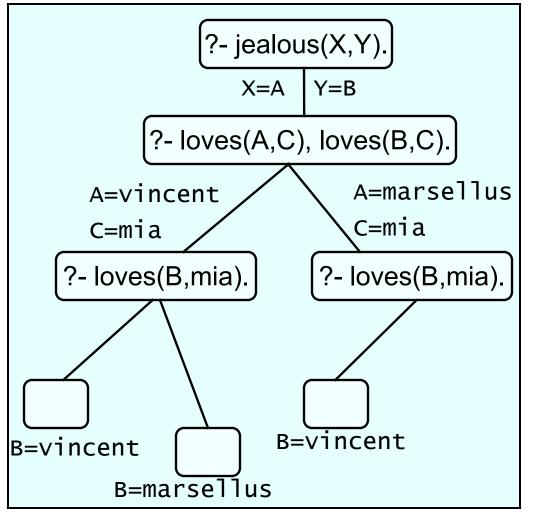
Y=marsellus;



loves(vincent,mia). loves(marsellus,mia).

jealous(A,B):loves(A,C), loves(B,C).

X=vincent
Y=marsellus;
X=marsellus
Y=vincent



loves(vincent,mia). loves(marsellus,mia).

jealous(A,B):loves(A,C),

loves(B,C).

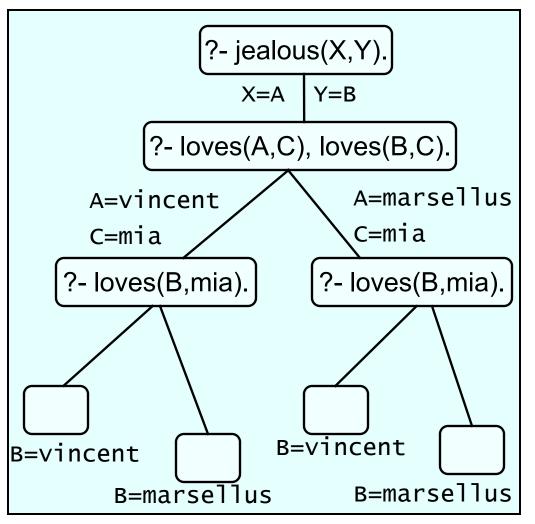
. . . .

X=marsellus

Y=vincent;

X=marsellus

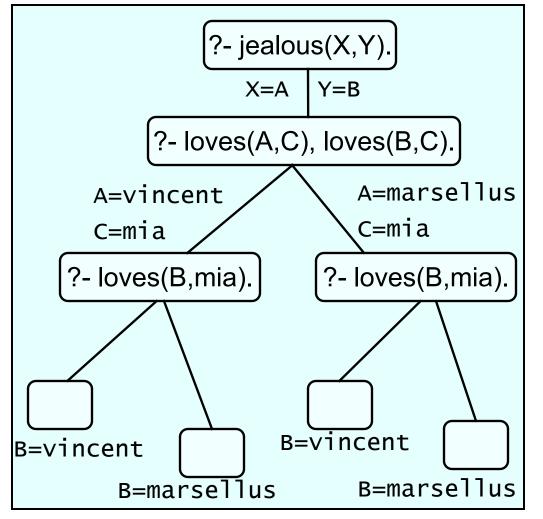
Y=marsellus



Another example

loves(vincent,mia).
loves(marsellus,mia).
jealous(A,B):loves(A,C),
loves(B,C).

X=marsellus
Y=vincent;
X=marsellus
Y=marsellus;
no



Recursive Definitions

- Prolog predicates can be defined recursively
- A predicate is recursively defined if one or more rules in its definition refers to itself

?-

Example: Descendant

```
child(anna,bridget).
child(bridget,caroline).
child(caroline,donna).
child(donna,emily).

descend(X,Y):- child(X,Y).
descend(X,Y):- child(X,Z), descend(Z,Y).
```

Example: Descendant

```
child(anna,bridget).
child(bridget,caroline).
child(caroline,donna).
child(donna,emily).

descend(X,Y):- child(X,Y).
descend(X,Y):- child(X,Z), descend(Z,Y).
```

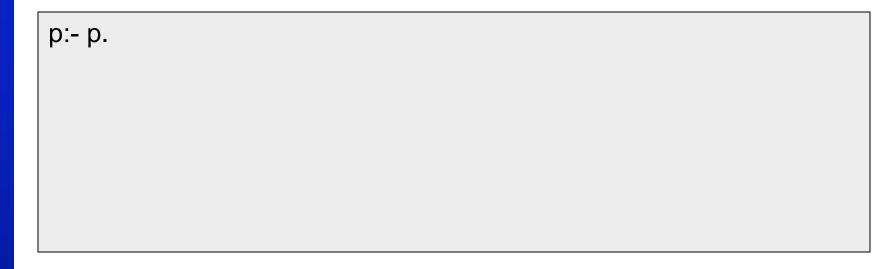
?- descend(anna,donna). yes

Another recursive definition



?- p.

Another recursive definition



?- p.
ERROR: out of memory

- Suppose we use the following way to write numerals:
 - 0 is a numeral.
 - 2. If X is a numeral, then so is succ(X).

numeral(0).

numeral(succ(X)):- numeral(X).

```
numeral(0).
numeral(succ(X)):- numeral(X).
```

```
?- numeral(succ(succ(succ(0)))).
yes
?-
```

numeral(0). numeral(succ(X)):- numeral(X).

?- numeral(X).

```
numeral(0).
numeral(succ(X)):- numeral(X).
```

```
?- numeral(X).
X=0;
X = succ(0);
X=succ(succ(0));
X=succ(succ(succ(0)));
X=succ(succ(succ(0))))
```

Example: Addition

?- add(succ(succ(0)),succ(succ(succ(0))), Result).
Result=succ(succ(succ(succ(succ(0)))))
yes

Example: Addition

add(0,X,X). %%% base clause

?- add(succ(succ(0)),succ(succ(succ(0))), Result).
Result=succ(succ(succ(succ(succ(0)))))
yes

Example: Addition

```
add(0,X,X). %%% base clause

add(succ(X),Y,succ(Z)):- %%% recursive clause

add(X,Y,Z).
```

```
?- add(succ(succ(0)),succ(succ(succ(0))), Result).
Result=succ(succ(succ(succ(succ(0)))))
yes
```

Examples from the book

descend1.pl

```
child(anna,bridget).
child(bridget,caroline).
child(caroline,donna).
child(donna,emily).

descend(X,Y):- child(X,Y).
descend(X,Y):- child(X,Z), descend(Z,Y).
```

?- descend(A,B).

A=anna

B=bridget

descend2.pl

```
child(anna,bridget).
child(bridget,caroline).
child(caroline,donna).
child(donna,emily).

descend(X,Y):- child(X,Z), descend(Z,Y).
descend(X,Y):- child(X,Y).
```

?- descend(A,B).

A=anna

B=emily

descend3.pl

```
child(anna,bridget).
child(bridget,caroline).
child(caroline,donna).
child(donna,emily).

descend(X,Y):- descend(Z,Y), child(X,Z).
descend(X,Y):- child(X,Y).
```

?- descend(A,B).

ERROR: OUT OF LOCAL STACK

descend4.pl

```
child(anna,bridget).
child(bridget,caroline).
child(caroline,donna).
child(donna,emily).

descend(X,Y):- child(X,Y).
descend(X,Y):- descend(Z,Y), child(X,Z).
```

?- descend(A,B).

Using dif predicate

```
mother(X, Y) :- parent(X, Y), female(X).
sister(X, Y) :-
    parent(Z, X),
    parent(Z, Y),
    female(X).
```

- What is wrong with this rule?
 - Any female is her own sister
- Solution?
 - Use dif ($\=$) predicate: ..., $X \setminus = Y$.

Comments

• Multi-line:

```
/* This is a comment
   This is another comment */
```

• Short:

% This is also a comment

Notation of Predicate Descriptions

+ Argument must be fully instantiated.

Think of the argument as *input*.

- Argument must be unbound.

Think of the argument as *output*.

? Either instantiated or unbound

Think of the argument as either *input* or *output*

Examples:

Write a predicate sum (+List, ?Sum) so that ...

Write a predicate append (?List1, ?List2, ?List1AndList2) such that ...

Side effects

• Some built-in "Predicates" may have side effects.

• Example: The built-in predicate write (+Term)

```
print if positive(X) :- X > 0, write(X).
```

Logical quantification

 Variables that appear in the head of a rule are universally quantified.

 Variables that <u>only</u> appear in the body of a rule are existentially quantified.

```
path(X,Y):- edge(X,Z), path(Z,Y).
```

For all nodes X and Y, there is a path from X to Y if there exists a node Z such that there is an edge from X to Z and there is path from Z to Y.