

Haskell - P57

standard

Monad

natural extension of applicative functor, support $\gg=$ (bind function)

$(\gg=) :: (\text{Monad } m) \Rightarrow m a \rightarrow (a \rightarrow m b) \rightarrow m b$

apply ~~function~~ function of type $a \rightarrow m b$ to value of type $m a$

$\gg=$ take one monadic value $(m a)$

one function take normal value (a) and return monadic value $(m b)$

and apply function $a \rightarrow mb$ to monadic value $m a$

$\gg=$ 作用于 Maybe 类型时，其行为类似于函数的组合

$\text{applyMaybe} :: \text{Maybe } a \rightarrow (a \rightarrow \text{Maybe } b) \rightarrow \text{Maybe } b$

$\text{applyMaybe } \text{Nothing} = \text{Nothing}$

$\text{applyMaybe } (\text{Just } x) f = f x$

so $\text{Just "smile"} \cdot \text{applyMaybe} (\lambda x \rightarrow \text{Just}(x + ":))$

$\rightarrow \text{"smile :)"}$

$\text{Nothing} \cdot \text{applyMaybe} (\lambda x \rightarrow \text{Just}(x + 1))$

$\rightarrow \text{Nothing}$

class Monad m where

$\text{return} :: a \rightarrow m a$

$(\gg=) :: m a \rightarrow (a \rightarrow m b) \rightarrow m b$

$(\gg) :: m a \rightarrow m b \rightarrow m b$

$x \gg y = x \gg= \lambda z \rightarrow y$ (order to succeed)

$\text{fail} :: \text{String} \rightarrow m a$

$\text{fail msg} = \text{error msg}$

when Haskell made, not occurred to people that applicative functor are good fit

目前 GHC.Base 中的定义： class (Applicative m) => Monad m where

$(\gg=) :: \text{forall } a b. m a \rightarrow (a \rightarrow m b) \rightarrow m b$

$(\gg) :: \text{forall } a b. m a \rightarrow m b \rightarrow m b$

$m \gg k = m \gg= \lambda _- \rightarrow k$

$\text{return} :: a \rightarrow m a$

$\text{return} = \text{pure}$

$\text{fail} :: \text{String} \rightarrow m a$

$\text{fail s} = \text{errorWithoutStackTrace s}$

Haskell - P58

Monad

`return :: (Monad m) => a -> m a`

take a value and put in minimal default context

in other word, wrap in a monad that still hold the value

do same thing as "pure" function from Applicative type class

nothing like "return" in most other languages

not end function execution or anything

`(>>) :: (Monad m) => m a -> (a -> m b) -> m b`

like function application: take normal value and feed to normal function

bind: take monadic value (a value with a context)

and feed to function that take normal value but return monadic value

`(>>) :: (Monad m) => m a -> m b -> m b`

`m >> k = m >>= _ -> k`

instance

come with default implementation and never implement when making Monad

`fail :: (Monad m) => String -> m a`

"fail" function never used explicitly in code

used by Haskell to enable failure in special syntactic construct for monad

so 将 Maybe 实现为 Monad 为 instance

instance Monad Maybe where

`return x = Just x`

`Nothing >>= f = Nothing`

`Just x >>= f = f x`

`fail _ = Nothing`

其中 `return` 等价于 `pure`, do same as in Applicative type class

`>>= (bind function)` 等价于之前定义的函数 `apply` Maybe

so `return "what" :: Maybe String -> Just "what"`

`Just 11 >>= \x -> return (x * 7) -> Just 77`

`Nothing >>= \x -> return (x * 7) -> Nothing`

Haskell - P59

考虑一个平衡问题，假定一个走钢丝的人手持一根平衡杆 (balancing pole)
在其走钢丝的过程中随机地有鸟降落 / 离开平衡杆的一侧

如果左右两侧停留的鸟的数量相差不超过 3 只，则可以保持平衡
否则走钢丝的人会失去平衡

首先定义 Birds 为 Int 的别名，Pole 是 (Birds, Birds) 的别名

type Birds = Int

type Pole = (Birds, Birds)

type State = Maybe Pole，State 是 Maybe Pole 的别名

再定义函数 landLeft 和 landRight

landLeft :: Birds → Pole → State

landLeft n (left, right)

| abs (left + n) - right) < 4 = Just (left + n, right)

| otherwise = Nothing

landRight :: Birds → Pole → State

landRight n (left, right)

| abs (left - (right + n)) < 4 = Just (left, right + n)

| otherwise = Nothing

于是可知 landLeft Int go landRight Int 有类型 Pole → Maybe Pole

则可以用于 Monad Maybe 的实现

so return (0, 0) => landRight 2 => landLeft 2 => landRight 2

→ Just (2, 4)

return (0, 0) => landLeft 1 => landRight 5 => landLeft 4

→ Nothing

注意 Applicative operator will fetch ~~value~~ result and feed to function

in manner appropriate for each applicative

and then put final applicative value together

假定当走钢丝的人踩到香蕉皮会立刻失去平衡

banana :: Pole → State

banana _ = Nothing

so return (0, 0) => landLeft 1 => banana => landRight 1

→ Nothing

注意 $\bullet \gg = \text{banana}$ 实际上等价于 $\gg = \text{Nothing}$

so return (0, 0) => landLeft 1 => Nothing => landRight 1 → Nothing

Haskell - P60

story

1+2+9 = 12+9 = 21

do notation 考虑 $\text{Just } 11 \gg= (\lambda x \rightarrow \text{Just } "?") \gg= (\lambda y \rightarrow \text{Just } (\text{show } x ++ y))$
 $\rightarrow \text{Just } "11?"$

注意从 $\text{Just } 11$ 中提取 11 并传给参数 x 。

从 $\text{Just } "?"$ 中提取 $"?"$ 并传给参数 y 。

value with failure context, can replace with failure

$\text{Nothing} \gg= (\lambda x \rightarrow \text{Just } "?") \gg= (\lambda y \rightarrow \text{Just } (\text{show } x ++ y))$
 $\rightarrow \text{Nothing}$

$\text{Nothing} \gg= (\lambda x \rightarrow \text{Nothing}) \gg= (\lambda y \rightarrow \text{Just } (\text{show } x ++ y))$
 $\rightarrow \text{Nothing}$

可以用 do notation 代替嵌套的 lambda function

foo :: Maybe String	bar :: Maybe String
foo = do	bar = do
Just 11 >>= (\lambda x \rightarrow	x <- Just 11
Just "?" >>= (\lambda y \rightarrow	y <- Just "?")
Just (show x ++ y))	Just (show x ++ y)
)	
return Just "11?"	return Just "11?"

temporarily extract string from Maybe value

without having to check if Maybe value is Just / Nothing at every step

if any of value extract from Nothing, whole do expression result in Nothing

do expression just different syntax for chaining monadic value

every line is monadic value in do expression

use <- to inspect value

variable become String if bind with <- Maybe String

just like use >>= to feed monadic value to lambda function

foo :: Maybe Bool	bar :: Maybe Bool
foo = do	bar = do
Just 11 >>= (\lambda x \rightarrow	x <- Just 11
Just (x > 7))	Just (x > 7)
)	
return Just True	return Just True

R = 5(5) 题目

Haskell - P61

STUDY

8829 - elimination

Monad

考虑将一个过程转换为 do notation

```
routine :: Maybe Pole
routine = case landLeft 2 (0,0) of
    Nothing → Nothing
    Just first → case landRight 2 first of
        Nothing → Nothing
        Just second → landLeft 1 second
```

有 Routine → Just (3,2)

类似于 use >>= to chain successive step

```
return (0,0) >>= landLeft 2 >>= landRight 2 >>= landLeft 1
→ Just (3,2)
```

each step rely on previous step and have added context of possible failure

可以等价地改写为 do notation 的形式

```
routineDo :: Maybe Pole
routineDo = do
    start ← return (0,0)
    first ← landLeft 2 start
    second ← landRight 2 first
    landLeft 1 second
```

each line of do expression must feature monadic value

do notation look like imperative code (指令式编程)

because written line by line

do notation is sequential

each value in each line rely on result of previous ones

along with context, whether succeeded or failed

在 do notation 也有 pattern match

so justH :: Maybe Char

```
justH = do
```

```
    (x:xs) ← Just "hello"
```

```
    ← return x
```

```
justH → Just 'h'
```

输出是 "Just 'h'"

Haskell - P62

Brackets

2859 - Combination

Monad

考虑 $\text{return } (0,0) \gg= \text{landLeft } 2 \gg= \text{Nothing}, \gg= \text{landRight } 2 \gg= \text{landLeft } 1$

$\rightarrow \text{Nothing}$

可以转换成函数形式 $(0,0)$ 被 $\text{landLeft } 2$ 处理

$\text{routineNothing} :: \text{Maybe Pole}$

$\text{routineNothing} = \text{case landLeft } 2 (0,0) \text{ of}$

$\text{Nothing} \rightarrow \text{Nothing}$

$\text{Just first} \rightarrow \text{case banana first of}$

$\text{Nothing} \rightarrow \text{Nothing}$

$\text{Just something} \rightarrow \text{case landRight } 2 \text{ something of}$

$\text{Nothing} \rightarrow \text{Nothing}$

$\text{Just second} \rightarrow \text{landLeft } 1 \text{ second}$

只有 $\text{routineNothing} \rightarrow \text{Nothing}$

可以进一步地改写成 do notation 的形式

$\text{routineDoNothing} :: \text{Maybe Pole}$

$\text{routineDoNothing} = \text{do}$

$\text{Start} \leftarrow \text{return } (0,0)$

$\text{first} \leftarrow \text{landLeft } 2 \text{ Start}$

$\rightarrow \text{Nothing}$ equivalent to $_ \leftarrow \text{Nothing}$

$\text{second} \leftarrow \text{landRight } 2 \text{ first}$

$\text{landLeft } 1 \text{ second}$

只有 $\text{routineDoNothing} \rightarrow \text{Nothing}$

write line in do notation without binding monadic value with \leftarrow

just like put \gg after monadic value whose result want to ignore

if matching fall through all patterns for given function

error thrown and program crash

dailed pattern matching in let expression result in error produced right away

while in do expression, call fail function

in context

part of Monad type class, enable failed pattern matching result in failure

so $\text{justN} :: \text{Maybe Char}$

$\text{fail} :: (\text{Monad m}) \Rightarrow \text{String} \rightarrow \text{m a}$

$\text{justN} = \text{do}$

$\text{fail msg} = \text{error msg}$

$(x :: \text{xs}) \leftarrow \text{Just }$

$x = \text{fail "No Char"}$

$\text{return } x$

而 instance Monad Maybe 实现中

$\text{justN} \rightarrow \text{Nothing}$

$\text{fail_} = \text{Nothing}$

Haskell - P63

+ 9 - Haskell



Monad

when deal with non-determinism, can make many choices

try all the choices, result non-determinism has many more values

so Monad type class 既是 list 的实现

instance Monad [] where

return $x = [x]$

$xs >>= f = concat (map f xs)$

fail _ = []

return : do the same thing as pure

make a list has only input one value

useful when wrap normal value into list

so can interact with non-deterministic value

$>>= :: [a] \rightarrow (a \rightarrow [b]) \rightarrow [b]$

take all possible values in [a]

seed into function, take normal value and return monadic value

即类型为 $a \rightarrow [b]$

最后用过 $concat :: (Foldable t) \Rightarrow t [a] \rightarrow [a]$

将 $[[b]]$ 组装为 $[b]$

so $[1, 2, 3] >>= \lambda x \rightarrow [x, -x]$

$\rightarrow [1, -1, 2, -2, 3, -3]$

jail : non-determinism also support for failure

empty list [] 类似于 Maybe 中的 Nothing

signify absence of result

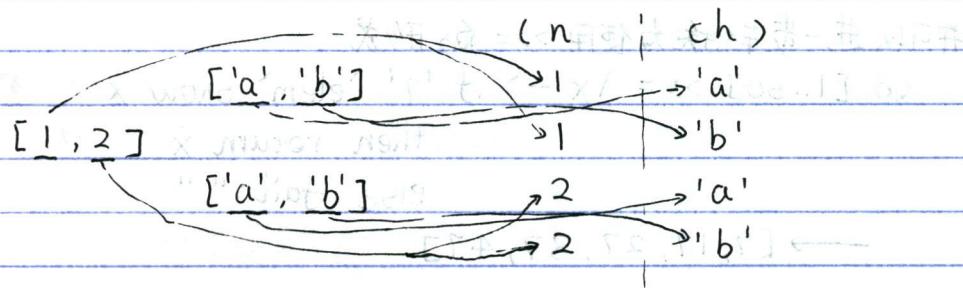
so $[] >>= \lambda x \rightarrow "bad"$

$[1, 2, 3] >>= \lambda x \rightarrow []$

与 Maybe 类似, 可以连续地使用 $>>=$

so $[1, 2] >>= \lambda n \rightarrow ['a', 'b'] >>= \lambda ch \rightarrow return (n, ch)$

$\rightarrow [(1, 'a')), (1, 'b')), (2, 'a')), (2, 'b'))$



C Haskell - P64

Monad

对于 $[1..2] >>= \lambda n \rightarrow [a', b'] >>= \lambda ch \rightarrow \text{return } (n, ch)$

$\rightarrow [(1, 'a'), (1, 'b'), (2, 'a'), (2, 'b')]$

转换为 do notation 的形式有

$\text{listOfTuple} :: [(Int, Char)]$

$\text{listOfTuple} = \text{do}$

$n <- [1..2]$

$ch <- ['a', 'b']$

$\text{return } (n, ch)$

$\text{listOfTuple} \rightarrow [(1, 'a'), (1, 'b'), (2, 'a'), (2, 'b')]$

注意到这个 do notation 类似于列表生成式 (list comprehension)

即有 $[(n, ch) | n <- [1..2], ch <- ['a', 'b']]$

$\rightarrow [(1, 'a'), (1, 'b'), (2, 'a'), (2, 'b')]]$

实际上列表生成式是使用 Monad [] 的语法糖

syntactic sugar for using list as monad

对于列表生成式中使用筛选条件 (filter)

如 $[x | x <- [1..50], '7' `elem` show x]$

$\rightarrow [7, 17, 27, 37, 47]$

同样可以转换为 do notation 的形式

注意对于不满足筛选条件的情况，可以调用 fail 函数以生成 []

$\text{listOfSeven} :: [Int]$

$\text{listOfSeven} = \text{do}$

$n <- [1..50] - "fail" <- x / = << []$

$\text{if } '7' `elem` show n \text{ then } = << []$

$\text{return } n$

else

$= << \text{用 } fail \text{ 替换 } n$

$(n, n) - "fail" <- "b" / = << ['d', 'b'] - n / = << []$

$\text{listOfSeven} \rightarrow [7, 17, 27, 37, 47]$

并可以进一步转换为使用 >>= 的形式

如 $[1..50] >>= \lambda x \rightarrow \text{if } '7' `elem` show x$

$\text{then return } x$

else fail " "

$\rightarrow [7, 17, 27, 37, 47]$

Haskell - P65

STLSet

P659 - <http://hackage.haskell.org/>

MonadPlus 在 Haskell 中提供 3 MonadPlus type class

for Monad can also act as Monoid

在 GHC.Base 处理 type class MonadPlus 的定义为

~~infixl 3 <|>~~

class (Applicative f) => Alternative f where

empty :: f a default value from identity of '<|>'

(<|>) :: f a -> f a -> f a associative binary operation

{some :: f a -> f [a]

one or more | some v = Some_v

where many_v = Some_v <|> pure [] not

some_v = liftA2 (:) v many_v

{many :: f a -> f [a]

zero or more | many v = many_v

where many_v = Some_v <|> pure [] not

Some_v = liftA2 (:) v many_v

class (Alternative m, Monad m) => MonadPlus m where

mzero :: m a default value] default definition of

mzero = empty identity of 'mplus'

mplus :: m a -> m a -> m a default definition of

mplus = (<|>) associative operation

其中 mzero synonymous to mempty for Monoid type class

represent non-deterministic computation has no result at all (failed)

mplus synonymous to mappend for Monoid type class

join two non-deterministic values into one

Maybe 类型有 Alternative 的实现为

instance Alternative Maybe where] instance MonadPlus Maybe

empty = Nothing] (采用在 Alternative 中的实现)

Nothing <|> r = r] (采用在 Alternative 中的实现)

Nothing <|> _ = Nothing] (采用在 Alternative 中的实现)

empty = []] (采用在 Alternative 中的实现)

(<|>) = (++)] (采用在 Alternative 中的实现)

Haskell - P66

MonadPlus 在 Haskell 中，Control.Monad 提供了与 Monad 相关的类和函数

do guard :: (MonadPlus m) => Bool -> m ()

guard True = return ()

guard False = mzero

当传入 Bool 值为 True 时，将空元组 () 装入 minimal default context

而传入 Bool 值为 False 时，返回 failed monadic value

do guard (5 > 1) :: Maybe () -> Just ()

guard (5 > 11) :: Maybe () -> Nothing

guard (5 > 1) :: [()] -> [()]

guard (5 > 11) :: [()] -> []

当 guard 与 >> 组合使用时，>> ignore empty tuple () and present other result

如果 guard fail，则其后的操作也会失败

即向 >>= 传入一个 failed monadic value 也会返回一个 failed monadic value

如果 guard succeed，则 make successful value has dummy result of () inside
all allow computation to continue

do guard (5 > 1) -> return "TRUE" :: [String] -> ["TRUE"]

guard (5 > 11) -> return "TRUE" :: [String] -> []

再考虑列表生成式中使用筛选条件，可以使用 guard 进行实现

如 [x | x <- [1..50], '7' `elem` show x] -> [7, 17, 27, 37, 47]

[1..50] -> [7, 17, 27, 37, 47]

[1..50] -> (\x -> guard ('7' `elem` show x) >> return x)

可以进一步地用 guard 实现 do notation 的形式

do sevenOnly :: [Int]

sevenOnly = do

x <- [1..50]

guard ('7' `elem` show x)

return x

sevenOnly -> [7, 17, 27, 37, 47]

注意在这个过程中，由于 guard 函数抛出的 failed monadic value 阻断计算

所以实际上相当于一个 if - then - else 的结构

并在失败时调用 fail 函数

Haskell - P67

Knight's quest, 考虑一个国际象棋棋盘(chess board)且仅有一个骑士(knight)

对于给定 knight 的起始位置和一个目标位置 s

考虑是否可以用恰好3次移动到达

令 type Knight Pos = (Int, Int)

其中第一个元素表示 knight 所在的列

第二个元素表示 knight 所在的行

可以通常地定义一个使用 filter 的函数

moveKnight :: Knight Pos → [Knight Pos]

moveKnight (c, r) = filter onBoard

[(c+2, r+1), (c+1, r+2), (c-1, r+2), (c-2, r+1),

[(c-2, r-1), (c-1, r-2), (c+1, r-2), (c+2, r-1)]

where onBoard (c', r') = c' `elem` [1..8] && r' `elem` [1..8]

其中 onBoard 函数用于保证输出位置均在棋盘内部

则 moveKnight 也可以转换为等价的 do notation 形式

moveKnight :: Knight Pos → [Knight Pos]

moveKnight (c, r) = do

(c', r') ← [(c+2, r+1), (c+1, r+2), (c-1, r+2), (c-2, r+1),

[(c-2, r-1), (c-1, r-2), (c+1, r-2), (c+2, r-1)]

guard (c' `elem` [1..8] && r' `elem` [1..8])

return (c', r')

to moveKnight (7, 7) → [(5, 8), (5, 6), (6, 5), (8, 5)]

可以定义函数 at3 从返回给定起始位置可以用恰好3次移动到达的位置

at3 :: Knight Pos → [Knight Pos]

at3 start = do

first ← moveKnight start

second ← moveKnight first

moveKnight second

等价于 at3 start = return start >>= moveKnight >>= moveKnight

再定义函数 canReachAt3 用于判断

canReachAt3 :: Knight Pos → Knight Pos → Bool

canReachAt3 start end = end `elem` at3 start

to (8, 1) `canReachAt3` (7, 7) → True

Haskell - P68

Monad

与 functor 类似, Monad 也需要符合一些规则, 即 monad law. 注意在 Haskell 中, 对于 Monad type class 的 instance 并不要求满足 monad law, 所以需要在实现 instance Monad 的时候特别注意.

left identity: $\text{return } x >>= f \text{ 等价于 } f x$

feed monadic value to function $f :: (\text{Monad } m) \Rightarrow a \rightarrow m b$

not different than apply function to normal value

如 return "wow" $>>= (\lambda x \rightarrow [x, x, x])$

$\rightarrow ["\text{wow}", "\text{wow}", "\text{wow}"]$

$(\lambda x \rightarrow [x, x, x]) = "w\text{ow}" = (\lambda x \rightarrow [x, x, x]) "w\text{ow}"$

$\rightarrow ["\text{wow}", "\text{wow}", "\text{wow}"]$

right identity: $m >>= \text{return } m \text{ 等价于 } m$

注意 $\text{return} :: (\text{Monad } m) \Rightarrow a \rightarrow m a = a$

如对于 instance Monad [] 中的 $>>=$ 的实现 ($\text{concat } (\text{map } f \text{ xs})$)

有 $\text{xs} >>= f = \text{concat } (\text{map } f \text{ xs}) = ["\text{"}]$

于是 $[1, 2, 3] >>= \text{return} \rightarrow \text{concat } [[1], [2], [3]] \rightarrow [1, 2, 3]$

associativity: $(m >>= f) >>= g \text{ 等价于 } m >>= (\lambda x \rightarrow f x >>= g)$

"chain of monadic function application with $>>=$ not matter how nested"

类似于函数复合运算符, $(.) :: (b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow (a \rightarrow c)$

$f . g = (\lambda x \rightarrow f (g x))$

有 $(<= <) :: (\text{Monad } m) \Rightarrow (b \rightarrow m c) \rightarrow (a \rightarrow m b) \rightarrow (a \rightarrow m c)$

$f <= < g = (\lambda x \rightarrow g x >>= f)$

如 let $f x = [x, -x]$

let $g x = [x * 3, x * 2]$

let $h = f <= < g$

$h 1 \rightarrow [3, -3, 2, -2]$

结合 Monad law 则有: 对于任意 monadic function $f :: (\text{Monad } m) \Rightarrow a \rightarrow m b$

有 $f <= < \text{return}$ 等价于 f

$\text{return} <= < f$ 等价于 f

$(f <= < g) <= < h$ 等价于 $f <= < (g <= < h)$

其中 $g :: (\text{Monad } m) \Rightarrow c \rightarrow m a$, $h :: (\text{Monad } m) \Rightarrow d \rightarrow m c$