

## Dijkstra's Algorithm( lazy implementation)

- Single Source shortest path algorithm
- Works for a graph  $G(V,E)$  with non negative edge weights!
- Why: To ensure that once a node has been visited it's optimal distance cannot be improved!
- So it acts in a **greedy** way: by selecting the most promising node ever time.
- Logic:
  - > Maintain a **Priority Queue** of key-value pairs (node,distance) which tells you which node to visit next.
  - > So we maintain a **distance array** where the distance to every node =  $\infty$
  - > Kickstart the algorithm:
    - > Insert  $[s,0]$  into the Priority Queue.
    - > Loop while(P.Q. is not empty {  
pulling out the next most promising [node,distance] pair  
}
    - > for every node we visit:
      - > Iterate over all edges outwards from the current node & relax each edge appending a new [node, distance] pair to the P.Q. for every relaxation.
    - > Having duplicate key entries in P.Q. is what makes this specific implementation lazy!
- Pseudo Code:
  - # Returns an array that contains the shortest distance to every node from source(s)
  - # g: adjacency list of our weighted graph
  - # n: num of nodes in our graph
  - # s: index of source node ( $0 \leq s \leq n$ )

### dijkstra(g,n,s):

```
vis = [false, false, ...,false]; # size = n
dist = [∞, ∞, ..., ∞]; # size = n
dist[s] = 0;
init(pq); # initialize an empty priority queue
pq.insert([s,0]); # kickstart the algorithm
```

```
while(pq.size() != 0){
    index, minVal = pq.remove(); # remove most promising index, distance
    vis[index] = true; # mark index as visited
```

```
    for(edge: g[index]){
        if(vis[edge.to]){
```

```

        continue; # so we won't visit them again!
    }
    newDist = dist[index] + edge.cost;
    if(newDist < dist[edge.to]){
        dist[edge.to] = newDist;
        pq.insert([edge.to, newDist]);
    }
}
}
return dist; // return array with optimal distances

```

#### - Finding the **optimal path** :

If we want to not only find the optimal distance to a specific node but also “what sequence of nodes were taken” to get there, we need to track additional info:

<The index of the previous node> we took to get to our current node.

Prev = [null, null, ..., null]; # array of size n

```

    .
    .
    .
if(newDist < dist[edge.to]) { # Relaxation
    prev[edge.to] = index;
    dist[edge.to] = newDist;
    pq.insert([edge.to, newDist]);
}

```

#### - find the **shortest path between 2 nodes**:

# g: adjacency list of our weighted graph

# n: number of nodes in the graph

# s: index of source(starting) node

# e: index of end node

#### **findShortestPath(g, n, s, e):**

```

    dist, prev = dijkstra(g,n,s);
    path = [] ; # empty list, will be in reverse order
    if (dist[e] == oo){
        return path;
    }
    for(at=e; at!=null; at=prev[at]){
        path.add(at);
    }
    path.reverse();
    return path;

```