

New York University, CIMS, CS, Course CSCI-GA.3140-001, Spring 2024

“Abstract Interpretation”

Ch. 15, Fixpoints

(Prerequisites to be self-studied)

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Class 3, Monday, February 5th, 2024, 4:55-6:55 PM, WWH, Room CIWW 202

These slides are available at

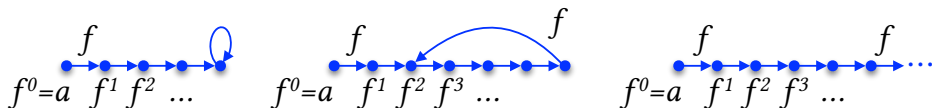
<http://cs.nyu.edu/~pcousot/courses/spring24/CSCI-GA.3140-001/slides/>

[03--2024-02-05-structural-fixpoint-prefix-trace-semantics/slides-10--15--fixpoints-AI.pdf](http://cs.nyu.edu/~pcousot/courses/spring24/CSCI-GA.3140-001/slides/10-15-fixpoints-AI.pdf)

Ch. 15, Fixpoints

Iteration

- Iterating an **operator** f on a set S (i.e. $f \in S \rightarrow S$) from an element $a \in S$ i.e. $f^0 \triangleq a$ and $f^{n+1} \triangleq f(f^n)$ can



- (1) reach a **fixpoint** $f^p = f(f^p)$;
- (2) or, reach a **cycle** and iterate for ever along that cycle $f^{p+\ell} = f^p, \ell > 0$;
- (3) or, **iterate for ever** along infinitely many different elements $\forall p \neq q . f^p \neq f^q$.

This last case is impossible if the set S is finite.

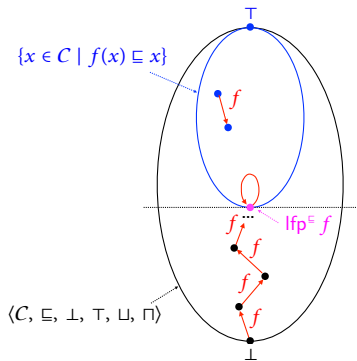
- We study sufficient conditions on S , f , and a ensuring that the iteration reaches a fixpoint.

en.wikipedia.org/wiki/Iterative_method

Fixpoints of increasing function on a complete lattice

There are two reasons why an increasing function $f \in \mathcal{C} \rightarrow \mathcal{C}$ on a complete lattice $\langle \mathcal{C}, \sqsubseteq, \perp, \top, \sqcup, \sqcap \rangle$ has a least fixpoint $\text{lfp}^{\sqsubseteq} f$:

- because it is the glb of the *post-fixpoints* $\{x \in \mathcal{C} \mid f(x) \sqsubseteq x\}$,
- because it is the limit (lub) of the iterates $\perp, f(\perp), \dots, f^n(\perp), f^{n+1}(\perp), \dots$



Fixpoint theorems

Tarski's fixpoint theorems (section 15.1)

Fixpoint

- A *fixpoint* of a function $f \in S \rightarrow S$ is an element x of S such that $f(x) = x$.
- In general a function may have no, one, or many fixpoints.

`en.wikipedia.org/wiki/Fixed_point_(mathematics)`

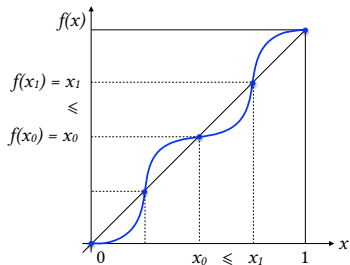
Increasing functions

- A function $f \in S \rightarrow S$ on a poset $\langle S, \sqsubseteq \rangle$ is increasing (denoted $f \in S \multimap S$) if and only if
$$\forall x, y \in S. (x \sqsubseteq y) \Rightarrow (f(x) \sqsubseteq f(y))$$
- A function f is increasing if and only if, for all $X \in \wp(S)$, when the lubs do exist,
$$\bigsqcup \{f(x) \mid x \in X\} \sqsubseteq f(\bigsqcup \{x \mid x \in X\})$$
- The composition of increasing functions is increasing

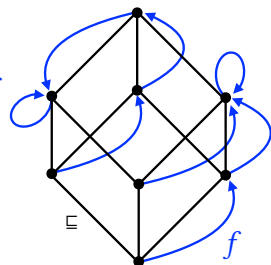
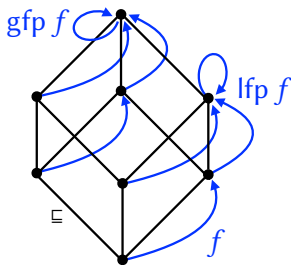
en.wikipedia.org/wiki/Monotonic_function

Fixpoints of increasing functions

- As shown by Alfred Tarski [Tarski, 1955], an increasing function on a complete lattice has at least one fixpoint and has a least one.



increasing function f



non-increasing function f

Least and greatest fixpoint

- $\text{lfp}^\sqsubseteq f$ denotes the least fixpoint of $f \in L \rightarrow L$ for the partial order $\langle L, \sqsubseteq \rangle$, if any.
- If the partial order \sqsubseteq is understood from the context, we write $\text{lfp } f$.
- $\text{lfp}_a^\sqsubseteq f$ is the least fixpoint of f for \sqsubseteq , which is \sqsubseteq -greater than or equal to $a \in L$, if any i.e.
 $a \sqsubseteq \text{lfp}_a^\sqsubseteq f = f(\text{lfp}_a^\sqsubseteq f)$ and if $a \sqsubseteq x = f(x)$ then $\text{lfp}_a^\sqsubseteq f \sqsubseteq x$
- Dually for $\text{gfp}_a^\sqsubseteq f$, $\text{gfp}^\sqsubseteq f$, and $\text{gfp } f$

en.wikipedia.org/wiki/Least_fixed_point

Tarski fixpoint theorem

Theorem 15.6 An increasing function $f \in L \multimap L$ on a complete lattice $\langle L, \sqsubseteq, \perp, \top, \sqcap, \sqcup \rangle$ has a least fixpoint $\text{lfp} \sqsubseteq f = \sqcap \{x \in L \mid f(x) \sqsubseteq x\}$.

Tarski fixpoint theorem

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- By duality, f has a greatest fixpoint $\text{gfp}^{\sqsubseteq} f = \sqcup \{x \in L \mid x \sqsubseteq f(x)\}$
- The set of fixpoints of f ordered by \sqsubseteq form a complete lattice
- This lattice of fixpoints is not necessary a sublattice of L meaning that the join and meet in this complete lattice of fixpoints may not be \sqcup and \sqcap
- Tarski's theorem 15.6 was generalized to cpos [Markowsky, 1976, theorem 9] and [Patariaia, 1997] as reported by [Escardó, 2003, Corollary 2.1.]

en.wikipedia.org/wiki/Alfred_Tarski
en.wikipedia.org/wiki/Knaster-Tarski_theorem

Proof of theorem 15.6 Let $a \triangleq \bigcap \{x \in L \mid f(x) \sqsubseteq x\}$ where the greatest lower bound \bigcap is well-defined in a complete lattice.

For all $x \in L$,

$$f(x) \sqsubseteq x$$

$$\Rightarrow a \sqsubseteq x$$

$\{ \text{since } x \in \{x \in L \mid f(x) \sqsubseteq x\} \text{ and the greatest lower bound } a \text{ is smaller than or equal to any element of the set} \}$

$$\Rightarrow f(a) \sqsubseteq f(x) \sqsubseteq x$$

$\{ f \text{ is increasing and } f(x) \sqsubseteq x \}$

$$\Rightarrow f(a) \sqsubseteq \bigcap \{x \in L \mid f(x) \sqsubseteq x\} = a$$

$\{ \text{since } f(a) \text{ is a lower bound of } \{x \in L \mid f(x) \sqsubseteq x\} \text{ hence smaller than or equal to its greatest lower bound } a \}$

With $a \triangleq \bigcap \{x \in L \mid f(x) \sqsubseteq x\}$, we have proved that

$$f(a) \sqsubseteq a$$

$$\Rightarrow f(f(a)) \sqsubseteq f(a) \quad \text{\textit{(since } } f \text{ is increasing\textit{)}}$$

$$\Rightarrow f(a) \in \{x \in L \mid f(x) \sqsubseteq x\} \quad \text{\textit{(def. } \in \textit{)}}$$

$$\Rightarrow a = \bigcap \{x \in L \mid f(x) \sqsubseteq x\} \sqsubseteq f(a)$$

$$\quad \text{\textit{(def. greatest lower bound smaller than or equal to any element of } } \{x \in L \mid f(x) \sqsubseteq x\} \textit{)}}$$

$$\Rightarrow a = f(a) \quad \text{\textit{(by antisymmetry)}}$$

So a is a fixpoint of f .

Let us prove that it is the least one.

- Let $x = f(x)$ be a fixpoint of f .
- Then $x \in \{x' \in L \mid f(x') \sqsubseteq x'\}$ by reflexivity.
- It follows that $a = \bigsqcap \{x' \in L \mid f(x') \sqsubseteq x'\} \sqsubseteq x$ by def. greatest lower bound.
- Therefore $a = \text{lfp} \sqsubseteq f$ is the least fixpoint of f . □

Optimization problem

- Tarski's theorem 15.6 shows that computing the least fixpoint $\text{lfp}^{\sqsubseteq} f$ can be understood as solving the *optimization problem* [Boyd and Vandenberghe, 2004]

$$\begin{array}{ll}\underset{x}{\text{minimize}} & f(x) \\ \text{subject to} & f(x) \sqsubseteq x\end{array}$$

where f is the objective function and $f(x) \sqsubseteq x$ is an inequality constraint.

- The optimization problem may be solvable for some complete lattices L [Cousot, 2005; Liberti and Marinelli, 2014] (e.g. finite ones)
- The optimization problem is, in general, uncomputable [Liberti, 2019].

en.wikipedia.org/wiki/Optimization_problem

Properties of the fixpoints of increasing functions

Let $f \in L \multimap L$ be an increasing function on a complete lattice $\langle L, \sqsubseteq, \perp, \top, \sqcap, \sqcup \rangle$

- $\text{lfp}^\sqsubseteq f = \sqcap \{x \in L \mid f(x) = x\}$
- $\text{lfp}^\sqsubseteq f = \text{lfp}^\sqsubseteq x \mapsto x \sqcup f(x)$.
- Given a prefixpoint $a \in L$ such that $a \sqsubseteq f(a)$, $\text{lfp}_a^\sqsubseteq f = \sqcap \{x \in L \mid a \sqsubseteq x \wedge f(x) \sqsubseteq x\}$ is the least fixpoint of f greater than or equal to a .
- For all $a \in L$ such that $a \sqsubseteq f(a)$, $\text{lfp}^\sqsubseteq x \mapsto a \sqcup f(x) = \text{lfp}_a^\sqsubseteq f$ where $\text{lfp}_a^\sqsubseteq f$ is the least fixpoint of f greater than or equal to $a \in L$.

Fixpoint induction

- Let $f \in L \multimap L$ be an increasing function on the complete lattice $\langle L, \sqsubseteq, \perp, \top, \sqcap, \sqcup \rangle$ and $x \in L$. We have

$$\text{lfp}^{\sqsubseteq} f \sqsubseteq x \text{ if and only if } \exists y \in L . f(y) \sqsubseteq y \wedge y \sqsubseteq x$$

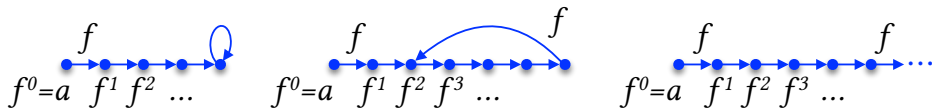
(This is the basis for the main program verification methods)

- Let $x \in L$. $\text{lfp}^{\sqsubseteq} f \sqsubseteq x$ if and only if $\text{lfp}^{\sqsubseteq} f \sqcap f(x) \sqsubseteq x$ ([Park, 1969, (2.2)]).

Iterative fixpoint theorems (section 15.2)

Iteration

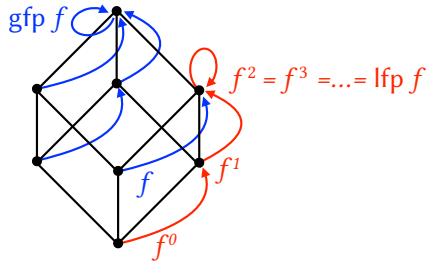
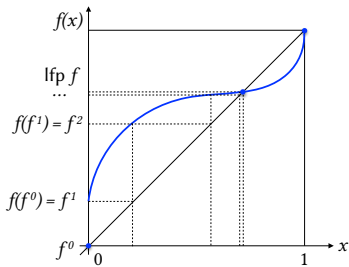
- In general, the iteration $f^0 = a$ and $f^{n+1} = f(f^n)$ for $n \in \mathbb{N}$ of a function $f \in S \rightarrow S$ on a set S from an element $a \in S$ either reaches a fixpoint, a cycle, or loops for ever.



- In the case of an increasing function on a poset $\langle P, \sqsubseteq \rangle$, the iterates from $a \in S$ such that $a \sqsubseteq f(a)$ are increasing so reaching a cycle is impossible by antisymmetry.
- If limits exist and have a continuity property then a fixpoint is always reached.

en.wikipedia.org/wiki/Fixed-point_theorem

Examples of iteration of increasing functions



Iterates

Definition 15.20 Let $f \in P \rightarrow P$ be a function on a poset $\langle P, \sqsubseteq, \perp, \sqcup \rangle$ with infimum \perp and lub \sqcup .

The (partially defined) *iterates* of f from $a \in P$ are $\langle f^n, n \in \mathbb{N} \cup \{\omega\} \rangle$ such that

- $f^0 \triangleq a$,
- $f^{n+1} \triangleq f(f^n)$ for $n \in \mathbb{N}$
- $f^\omega \triangleq \bigsqcup \{f^n \mid n \in \mathbb{N}\}$

(where, to denote the limit of infinite iterations, the naturals are extended by the infinity ω such that $\forall n \in \mathbb{N} . n < \omega \leq \omega$).

By default, $a = \perp$.

We write $\langle f^n(a), n \in \mathbb{N} \cup \{\omega\} \rangle$ to make clear that the iterates start from $a \neq \perp$.

The iterates *converge* in finitely many steps or are *ultimately stationary* at $\ell \in \mathbb{N}$ if and only if $\forall n \geq \ell . f^n = f^\ell$.

Tarski-Kantorovich iterative fixpoint

Theorem 15.21 Let $f \in P \rightarrow P$ be an increasing function on a poset $\langle P, \sqsubseteq, \perp, \sqcup \rangle$ with infimum \perp and lub \sqcup .

Let $\langle f^n, n \in \mathbb{N} \cup \{\omega\} \rangle$ be the iterates of f .

Assume that the least upper bound $f^\omega = \sqcup \{f^n \mid n \in \mathbb{N}\}$ exists and $f(\sqcup \{f^n \mid n \in \mathbb{N}\}) = \sqcup \{f(f^n) \mid n \in \mathbb{N}\}$ then

- $\langle f^n, n \in \mathbb{N} \cup \{\omega\} \rangle$ is an increasing chain
- f has a least fixpoint $\text{lfp}^\sqsubseteq f = \sqcup \{f^n \mid n \in \mathbb{N}\}$.

en.wikipedia.org/wiki/Kleene_fixed-point_theorem

The iterative fixpoint theorem is often proved with stronger hypotheses, see below.

Proof of theorem 15.21 Fixpoint:

- Observe that

$$\begin{aligned} & \bigsqcup \{f^n \mid n \in \mathbb{N}\} \\ &= \bigsqcup (\{f^{n+1} \mid n \in \mathbb{N}\} \cup \{f^0\}) && \text{\textit{(extract case } } n = 0 \text{)}} \\ &= (\bigsqcup \{f^{n+1} \mid n \in \mathbb{N}\}) \sqcup \perp && \text{\textit{(since } } f^0 = \perp \text{ is the infimum)}} \\ &= \bigsqcup \{f(f^n) \mid n \in \mathbb{N}\} && \text{\textit{(since } } f^{n+1} = f(f^n) \text{ for } n \in \mathbb{N}\text{)}} \end{aligned}$$

- Then

$$\begin{aligned} & f(\bigsqcup \{f^n \mid n \in \mathbb{N}\}) \\ &= \bigsqcup \{f(f^n) \mid n \in \mathbb{N}\} && \text{\textit{(by hypothesis)}} \\ &= \bigsqcup \{f^n \mid n \in \mathbb{N}\} \end{aligned}$$

- It follows that $\bigsqcup \{f^n \mid n \in \mathbb{N}\}$ is a fixpoint of f . □

Least fixpoint:

- Let $x = f(x)$ be a fixpoint of f . $f^0 = \perp \sqsubseteq x$ since \perp is the infimum (smallest element) of P .
- Assume that $f^n \sqsubseteq x$ by induction hypothesis.
- Then $f^{n+1} = f(f^n) \sqsubseteq f(x) = x$ by def. of the iterates, f increasing, and fixpoint property.
- By recurrence $\forall n \in \mathbb{N} . f^n \sqsubseteq x$.
- So x is an upper bound of the $\{f^n \mid n \in \mathbb{N}\}$
- Hence x is larger than or equal to the least upper bound $\bigsqcup \{f^n \mid n \in \mathbb{N}\}$,
- It follows that $\bigsqcup \{f^n \mid n \in \mathbb{N}\} = \text{lfp}^{\sqsubseteq} f$ is the least fixpoint of f . □

Increasing chain:

- We have $f^0 = \perp \sqsubseteq f^1$ by def. infimum.
 - Assume $f^n \sqsubseteq f^{n+1}$.
 - By definition 15.20 of the iterates and f increasing, we have $f^{n+1} = f(f^n) \sqsubseteq f(f^{n+1}) = f^{n+2}$, proving $\langle f^n, n \in \mathbb{N} \rangle$ to be increasing.
 - The lub $f^\omega = \bigsqcup \{f^n \mid n \in \mathbb{N}\}$ is assumed to exist so $\langle f^n, n \in \mathbb{N} \cup \{\omega\} \rangle$ is an increasing chain.
-

Chain (recall)

- A subset $\{x_i \mid i \in \mathbb{N}\} \subseteq P$ of a poset $\langle P, \sqsubseteq \rangle$ is an (enumerable) increasing chain if and only if $\forall i, j \in \mathbb{N} . (i \leq j) \Rightarrow (x_i \sqsubseteq x_j)$.
- The family $x_i, i \in \mathbb{N}$ considered as a function of i is therefore increasing.

(Upper) continuity

Definition 15.22 A function $f \in P \rightarrow P$ on a poset $\langle P, \sqsubseteq \rangle$ is upper continuous (denoted $f \in P \xrightarrow{uc} P$) if and only if

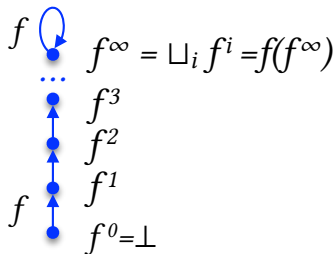
- for any enumerable increasing chain $\{x_i \mid i \in \mathbb{N}\} \subseteq P$ of P which has a least upper bound $\bigsqcup \{x_i \mid i \in \mathbb{N}\} \in P$ then $\bigsqcup \{f(x_i) \mid i \in \mathbb{N}\}$ exists in P
- and $f(\bigsqcup \{x_i \mid i \in \mathbb{N}\}) = \bigsqcup \{f(x_i) \mid i \in \mathbb{N}\}$.

The dual is lower-continuity (denoted $f \in P \xrightarrow{lc} P$).

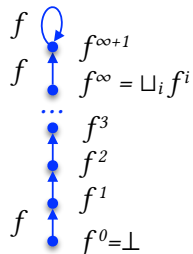
Observe that join preservation implies upper continuity.

en.wikipedia.org/wiki/Scott_continuity

Examples of upper-continuous functions



continuous



not continuous

Examples of upper-continuous functions

- An upper continuous function on a poset is **increasing** (but not all increasing functions are upper continuous).
- The **composition** of continuous functions is continuous.

CPO (recall)

- A complete partial order (cpo) is a poset $\langle P, \sqsubseteq, \perp \rangle$ with infimum \perp such that any increasing chain $\{x_i \mid i \in \mathbb{N}\}$ has a least upper bound $\bigsqcup \{x_i \mid i \in \mathbb{N}\} \in P$.

en.wikipedia.org/wiki/Chain-complete_partial_order

en.wikipedia.org/wiki/Complete_partial_order

Kleene/Scott iterative fixpoint theorem

Theorem 15.26 If $f \in L \xrightarrow{uc} L$ be an upper continuous function on a cpo $\langle L, \sqsubseteq, \perp, \sqcup \rangle$ then f has a least fixpoint $\text{lfp}^\sqsubseteq f = \bigsqcup_{n \in \mathbb{N}} f^n(\perp)$.

en.wikipedia.org/wiki/Kleene_fixed-point_theorem

The proof is similar to that of theorem 15.21.

Continuity of the least fixpoint of continuous functions

Theorem 15.32 If $F \in L \xrightarrow{uc} L \xrightarrow{uc} L$ be an upper continuous function on a cpo $\langle L, \sqsubseteq, \perp \rangle$ then $R \in L \mapsto \text{lfp}^{\sqsubseteq} F(R) \in L \xrightarrow{uc} L$.

Proof of theorem 15.32

- $F(R) \in L \xrightarrow{uc} L$ so $\text{lfp}^\sqsubseteq F(R) = \bigsqcup_{n \in \mathbb{N}} F(R)^n(\perp)$ exists by theorem 15.26.
- Let us prove, by recurrence, that $R \in L \mapsto F(R)^n(\perp) \in L \xrightarrow{uc} L$.
 - For the basis, $R \in L \mapsto F(R)^0(\perp) = R \in L \mapsto \perp$ is constant so continuous.
 - For the induction, $R \in L \mapsto F(R)^n(\perp)$ is continuous by induction hypothesis so $R \in L \mapsto F(R)^{n+1}(\perp) = F(R) \circ R \in L \mapsto F(R)^n(\perp)$ is continuous since $F(R) \in L \xrightarrow{uc} L$ and,
the composition of continuous functions is continuous. □

- For $F \in L \xrightarrow{uc} L \xrightarrow{uc} L, R \in L \mapsto \text{lfp}^\sqsubseteq F(R)$ is continuous. Let $\langle R_i, i \in \mathbb{N} \rangle$ be a \sqsubseteq -increasing chain.

$$\begin{aligned}
& (R \in L \mapsto \text{lfp}^\sqsubseteq F(R))(\bigsqcup_{i \in \mathbb{N}} R_i) \\
= & (R \in L \mapsto \bigsqcup_{n \in \mathbb{N}} F(R)^n(\perp))(\bigsqcup_{i \in \mathbb{N}} R_i) && \{ \text{theorem 15.26} \} \\
= & \bigsqcup_{n \in \mathbb{N}} F(\bigsqcup_{i \in \mathbb{N}} R_i)^n(\perp) && \{ \text{function application} \} \\
= & \bigsqcup_{n \in \mathbb{N}} \bigsqcup_{i \in \mathbb{N}} F(R_i)^n(\perp) && \{ R \in L \mapsto F(R)^n(\perp) \in L \xrightarrow{uc} L \text{ is continuous} \} \\
= & \bigsqcup_{i \in \mathbb{N}} \bigsqcup_{n \in \mathbb{N}} F(R_i)^n(\perp) && \{ \text{def. } \sqcup \} \\
= & \bigsqcup_{i \in \mathbb{N}} (R \in L \mapsto \bigsqcup_{n \in \mathbb{N}} F(R)^n(\perp))(R_i) && \{ \text{def. function application} \} \\
= & \bigsqcup_{i \in \mathbb{N}} (R \in L \mapsto \text{lfp}^\sqsubseteq F(R))(R_i) && \{ \text{theorem 15.26} \}
\end{aligned}$$

Conjugate fixpoint theorem

David Park conjugate fixpoint theorem [Park, 1969], [Park, 1979, (4.1.2)]

Theorem 15.33 Let S be a set and $f \in \wp(S) \rightarrow \wp(S)$ be \subseteq -increasing on the complete lattice $\langle \wp(S), \subseteq, \emptyset, S, \cap, \cup \rangle$. Let $\neg X \triangleq S \setminus X$ be the set complement. Define $\tilde{f} \triangleq X \mapsto \neg f(\neg X)$. Then $\text{gfp}^\subseteq f = \neg \text{lfp}^\subseteq \tilde{f}$.

- $\text{lfp}^\subseteq f \cap \text{lfp}^\subseteq \tilde{f} = \emptyset$ [Park, 1969, (2.2.3)]
- $\text{lfp}^\subseteq f \cup \text{lfp}^\subseteq \tilde{f} = S$ if and only if $\text{lfp}^\subseteq f$ is the unique fixpoint of f [Park, 1969, (2.2.4)]
- This generalizes to complete boolean lattices [Monk and Bonnet, 1983].

en.wikipedia.org/wiki/Complete_Boolean_algebra

Proof of theorem 15.33

$$\begin{aligned} & \neg \text{lfp}^{\subseteq} \tilde{f} \\ &= \neg \text{lfp}^{\subseteq} X \mapsto \neg f(\neg X) && \{ \text{def. } \tilde{f} \} \\ &= \neg \bigcap \{ X \mid \neg f(\neg X) \subseteq X \} \\ & && \{ f \text{ is increasing so } X \mapsto \neg f(\neg X) \text{ is increasing and theorem 15.6} \} \\ &= \bigcup \{ \neg X \mid f(\neg X) \supseteq \neg X \} && \{ \text{De Morgan laws} \} \\ &= \bigcup \{ Y \mid Y \subseteq f(Y) \} && \{ \text{letting } Y \triangleq \neg X \} \\ &= \text{gfp}^{\subseteq} f && \{ \text{dual of theorem 15.6} \} \quad \square \end{aligned}$$

Conclusion

Conclusion

- Fixpoint theorems are the mathematical foundations of program semantics, verification, and static analysis;
- Fixpoint theorems will be extensively used in the following chapters of the book;
- They deserve to be studied carefully, e.g. by trying to redo the proofs by yourself.

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The End, Thank you