

# STACK

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# Stack

- Non Primitive and Linear Data Structure
- LIFO [Last In First Out] principle
- Can be implemented using Array or Linked List
- A variable called TOP is used to keep track of the top of the stack.
- Stack Overflow = When we try to push a data on to the stack but the stack is already full.
- Stack Underflow = When we try to pop a data from the stack but the stack is already empty.

# Algorithm for PUSH

- Check if the stack is full as **top == MAX\_SIZE - 1**
- If the stack is full then display the appropriate message such as “Stack Overflow” and terminate the program
  - Stack Overflow
  - `exit(0)`
- Else
  - Increment the top variable as **top = top + 1**
  - Push the element to the top of the stack as `stack[top] = data`
- Terminate the program

# Algorithm for POP

- Check if the stack is empty as **top < 0**
- If the stack is empty then display the appropriate message such as “Stack Underflow” and terminate the program
  - Stack Underflow
  - `exit(0)`
- Else
  - Pop the element of the top of the stack as `data = stack[top]`
  - Decrement the top variable as **top = top - 1**
- Terminate the program

# Algorithm for PEEK

- Check if the stack is empty as **top < 0**
- If the stack is empty then display the appropriate message such as “Stack Underflow” and terminate the program
  - Stack Underflow / Stack Empty
  - `exit(0)`
- Else
  - Display the element of the top of the stack as **data = stack[top]**
- Terminate the program

# Applications of stack

- Computer internally uses stack when we implement a recursive function.
- Internet Browser uses stack to keep track of the history of websites you visited.
- Calculation of postfix expression is done using stack.

# Precedence

<b>Operators</b>	<b>Symbols</b>	<b>Priority / Precedence</b>
Exponent	$\wedge$ , \$	Highest
Division and Multiplication	/, *	Second Highest
Addition and Subtraction	+, -	Lowest

# Algorithm for Infix to Postfix Expression Conversion

- Scan the character from left to right
- If the scanned character is '('
  - PUSH TO THE STACK
- If the scanned character is an OPERAND
  - ADD TO THE POSTFIX STRING
- If the scanned character is an OPERATOR
  - Check Precedence
    - If precedence of the scanned character is greater than to the precedence of the operator in the stack then PUSH the operator to the stack
    - Else POP the operator from the stack and add the popped operator to the postfix string. Then, PUSH the scanned operator onto the stack.
- If the scanned character is ')'
  - POP ALL THE OPERATOR INSIDE PARENTHESIS “()” AND ADD IT TO THE POSTFIX STRING



Input Expression	Stack	Postfix Expression
K		K
+	+	
L	+	K L
-	-	K L +
M	-	K L + M
*	- *	K L + M
N	- *	K L + M N
+	+	K L + M N * K L + M N * -
(	+ (	K L + M N * -
O	+ (	K L + M N * - O
^	+ ( ^	K L + M N * - O
P	+ ( ^	K L + M N * - O P
)	+	K L + M N * - O P ^
*	+ *	K L + M N * - O P ^
W	+ *	K L + M N * - O P ^ W
/	+ /	K L + M N * - O P ^ W *
U	+ /	K L + M N * - O P ^ W * U
/	+ /	K L + M N * - O P ^ W * U /
V	+ /	K L + M N * - O P ^ W * U / V
*	+ *	K L + M N * - O P ^ W * U / V /
T	+ *	K L + M N * - O P ^ W * U / V / T
+	+	K L + M N * - O P ^ W * U / V / T * K L + M N * - O P ^ W * U / V / T * +
Q	+	K L + M N * - O P ^ W * U / V / T * Q
		K L + M N * - O P ^ W * U / V / T * + Q +

# Algorithm for Infix to Prefix Expression Conversion

- Scan the character from right to left.
- If the scanned character is '('
  - PUSH TO THE STACK
- If the scanned character is an OPERAND
  - ADD TO THE POSTFIX STRING
- If the scanned character is an OPERATOR
  - Check Precedence
    - If precedence of the scanned character is greater than to the precedence of the operator in the stack then PUSH the operator to the stack
    - Else POP the operator from the stack and add the popped operator to the postfix string. Then, PUSH the scanned operator onto the stack.
- If the scanned character is ')'
  - POP ALL THE OPERATOR INSIDE PARENTHESIS "()" AND ADD IT TO THE POSTFIX STRING

Finally, Reverse the result postfix to obtain prefix expression.

Infix to prefix :  $K + L - M * N + (O^P) * W/U/V * T + Q$

Input expression	Stack	Prefix expression
Q		Q
+	+	Q
T	+	QT
*	+*	QT
V	+*	QTV
/	+*/	QTV
U	+*/	QTVU
/	+*//	QTVU
W	+*//	QTVUW
*	+*//*	QTVUW
)	+*//*)	QTVUW
P	+*//*)	QTVUWP
^	+*//*)^	QTVUWP
O	+*//*)^	QTVUWPO
{	+*//*)	QTVUWPO^
+	++	QTVUWPO^*//*
N	++	QTVUWPO^*//*N
*	++*	QTVUWPO^*//*N
M	++*	QTVUWPO^*//*NM
-	++-	QTVUWPO^*//*NM*
L	++-	QTVUWPO^*//*NM*L
+	+++	QTVUWPO^*//*NM*L
K	+++	QTVUWPO^*//*NM*LK
		QTVUWPO^*//*NM*LK+--+

Thus, The reverse of QTVUWPO^\*//\*NM\*LK+--+ is **++-+KL\*MN\*//\*^OPWUVTQ** which is the required prefix expression.

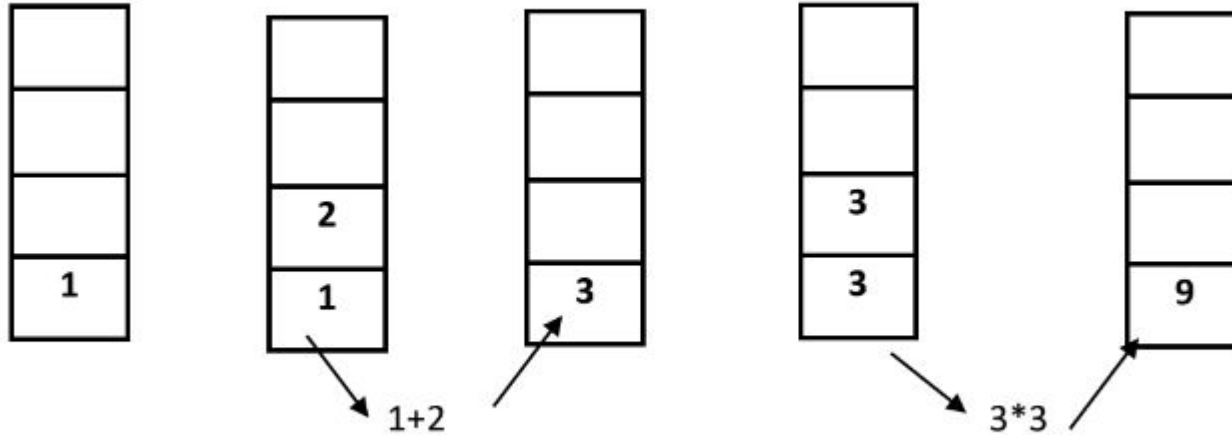
# Evaluation of postfix expression

Algorithm to evaluate the postfix expression:

1. Start
2. If character is number  
    Push on the stack
3. If character is operator(op)
  - a. Val1=pop
  - b. Val2=pop
  - c. Perform  $\text{result} = \text{val2 op val1}$
  - d. Push the result into stack
4. Repeat 2 to 3 until postfix expression is scanned.
5. Output the result
6. Stop

# Evaluate $AB+C^*$

Let  $A=1$ ,  $B=2$  and  $C=3$



Ans: 9

# Evaluate $ABC^*DEF^{\wedge}/G^*-H^*+$

Let A= 2, B= 3, C=9, D=8, E=1, F=4, G= 2, H=7

						4	
					1	1	1
		9		8	8	8	8
	3	3	27	27	27	27	27
2	2	2	2	2	2	2	2

	2					
8	8	16		7		
27	27	27	11	11	77	
2	2	2	2	2	2	79

Ans: 79