Fast-Reg Memory Registration in Mellanox OFED 3.3 By Austin Pohlmann

I Introduction

Mellanox's RDMA drivers provide both user level and kernel level APIs for users to develop RDMA applications. However, the kernel level API is severely lacking in documentation, mostly due to all the changes between versions. The kernel API provides a few methods of registering memory for use with rdma transfers, and the most recent drivers (4.0+) have changed the way users do this. This paper is to serve as a reference to anyone who wishes to use version 3.3 (and possibly 3.4) of Mellanox's OFED drivers. All of this is the way I perceived the usage of the fast reg memory system, and could possibly be incorrect.

The fast reg system provides a fast and asynchronous way to prepare and register memory regions for use with rdma operations. Most of this information has been inferred after looking through the source code for the drivers and reading through the appropriate header files. However, with all the assumptions I made, I was able to successfully able to register and use memory with the methods described here. All bolded words are ones that will most likely show up when trying to use this registration.

II Parameters

The functions referenced are in the **ib_verbs.h** file in directory where the driver was installed to. The use must first allocate the **ib_fast_reg_page_list** structure via a call to

ib_alloc_fast_reg_page_list(struct ib_device
 *device, int page_list_len)

The device can be found using a previously allocated protection domain as follows:

pd->device

The page_list_length should be set to the number of pages covered by a contiguous buffer that was previously allocated with kmalloc. The returned structure will have the following members:

struct ib_device *device; u64 *page_list; unsigned int max_page_list_len;

max_page_list_len is the length of the array page_list, which will be at LEAST the size specified upon allocation. The page list is to be populated by the DMA address that are returned from the ib_dma_* functions. Also, the page list must contain the address of the PAGES. A simple way to do this is to map the entire kmalloc'ed buffer using ib_dma_map_single, which will give you the DMA address of the buffer. With this address, you can loop through the page list buffer size/page size times, setting each entry to the sum of the previous entry and the page size.

NOTE: the original DMA address might not be on a boundry, so be sure logical AND the DMA address with the **PAGE_MASK** macro supplied by the kernel. This will effectively set any offset bits to zero.

Alongside allocating the page list, you must also allocate the memory region to be associated with the page list. A simple call to

will return the **ib_mr** structure needed to proceed. Note that the **max_page_list_length** should be the value that was stored in the page list structure earlier.

Before continuing, be sure that both sides are connected. On the server, this is accomplished after calling **rdma_accept**, and on the client, this is accomplished after the event handler

processes the **RDMA_CM_EVENT_ESTABLISHED** event.

The final step before the buffer is registered is the **fast_reg** work request. The **ib_send_wr** structure's union, **wr**, contains the **fast_reg** portion that we are concerned with.

iova_start

This value is the start of your buffer, and should be set to the DMA address returned from ib_dma_map_single

page_list

This is the same **page_list** found in the structure we allocated earlier.

page_shift

This is set to the **PAGE_SHIFT** macro provided by the kernel. Its value is the number of bits that are used as the offset in the address you provided in the page list.

page_list_length

This is the value that YOU specified to the first function call mentioned earlier, not max_page_list_length. This is how many of the entries in page_list that you populated.

length

This is the length of the buffer that starts at the address **iova_start**.

access_flags

The access rights for the memory being registered. These are the same flags used when registering a DMA mr.

rkey

This is the **rkey** contained in the **mr** structure allocated earlier

After the work request has been filled out, a call to **ib_post_send** will asynchronously complete

the registration (don't forget to set the send_flags to IB_SEND_SIGNALED in the ib_send_wr). When it is complete, completion queue will be notified and the completion's opcode will be IB_WC_FAST_REG_MR. When the currently running code tries to use the registered memory for the first time, it should sleep if the request has not yet been completed. The buffer should not be touched until after the work is completed. When the memory is no longer needed, please free everything in the following order:

ib_dma_unmap_single ib_dereg_mr ib_free_fast_reg_page_list kfree

Final notes: if you wish to use the memory without having to reregister it, please look into the <code>ib_update_fast_reg_key</code> function and the <code>invalidate_rkey</code> value in <code>ib_send_wr</code>. You must then post another <code>fast_reg</code> work request with the new <code>rkey</code> to notify the rdma device of the change.

III Example

The following code is a functioning example using the Mellanox OFED 3.3 drivers on centos 6.9 with the 3.10.87 kernel. The structs and global variables used in the code are shown in the beginning. This function would take place after a connection has already been established, and relies on the completion handler function (defined when creating the completion queue) to sleep until conditions are met. The fast reg region is created, and the relevant information for using it is sent to the remote side. What is not shown is the deregistration process, as it happens outside of the function. The error markers at the end can be examined in order to see what the deregistration would look like in case of an error at any point of the process

```
enum crdma state {
               // state is set to WAITING before going to sleep
     CONNECT REQUEST, // appears when a client requests to connect to a server
     CONNECTED, // appears when a client and server successfully connect
     ROUTE RESOLVED, // client has successfully resolved the route to server
     ADDR RESOLVED,
                     // client has resolved the host's address
     FRMR COMPLETE,
                     // the fast reg mr has been successfully registered
     RDMA WRITE COMPLETE,
                           // These a used when a completion for the
     RDMA SEND COMPLETE,
                          // relevant operations have been processed
     RDMA RECV COMPLETE,
     RDMA READ COMPLETE,
     DISCONNECT, // the remote side has disconnected
     ERROR // an error has occured
} state;
struct crdma cb {
     u8 addr[4]; //holds the address of the server (ipv4)
     uint16 t port; // the port that is being used to communicate
     struct rdma cm id *cmid, *child; // the cm ids, child is used for the
client's connection on the server's side and cmid is for self reference
     struct ib cq *cq; // the completion queue
     struct ib pd *pd; // the protection domain
     struct ib qp *qp; // the queue pair
     struct ib mr *mr; // the mr that I used before setting up the fast reg mr
     struct ib mr *frmr; // the mr used for fast req
     struct ib fast reg page list *frpl; // the page list for fast reg
     struct ib send wr send wr; // the wr used for any type of send
     struct ib send wr fast wr; // unused, I used send wr instead of this
     struct ib recv wr recv wr; // the wr used for any type of receive
     struct ib send wr *bad send; // supplied as an argument to ib post send
     struct ib recv wr *bad recv; // supplied to ib post recv
     struct ib sge test sge[2]; // 0 is recieve, 1 is send
```

```
u64 dma[3]; // array of dma addresses
     uint32 t remote rkey; // the rkey for the remote mr
     uint64 t remote addr; // the remote dma address of the mr
     uint32 t remote len; // the length of the remote mr
     };
static int crdma fr(struct crdma cb *cb){
     int ret=0,i;
                     // return variable and for loop variable
     // allocate the page list structure
     cb->frpl = ib alloc fast reg page list(cb->pd->device, 8);
     pr info("Page list: %p\nMax page list length: %u\n", cb->frpl, cb->frpl-
>max page list len);
     // see if we got a valid address. If not, there was an error
     if(IS ERR(cb->frpl)) {
           pr err("Failed to allocate the page list!!!\n");
           return PTR ERR(cb->frpl);
     // Allocate the fast reg mr and see if we got a valid address
     cb->frmr = ib alloc fast reg mr(cb->pd,
                            cb->frpl->max page list len);
     if (IS ERR(cb->frmr)) {
           pr_err("fast_reg_mr failed\n");
           ret = PTR ERR(cb->frmr);
           goto error0;
     }
     pr info("Fast reg rkey: %lu\n", (long unsigned)cb->frmr->rkey);
     // get the memory we want to register and map it to a dma address
     cb->buffs[2] = kmalloc(8*4096, GFP KERNEL);
     cb->dma[2] = ib dma map single(cb->pd->device, cb->buffs[2], 8*4096,
DMA BIDIRECTIONAL);
     // make sure everything mapped okay
     if(ib dma mapping error(cb->pd->device, cb->dma[2])){
           pr err("Error mapping fast reg buffer\n");
```

```
ret = 1;
            goto error1;
      }
      // fill the page list by removing the offset values from every page in
our buffer
      for (i=0; i<7; i++) {
            cb->frpl->page list[i] = (cb->dma[2] + i*4096) & PAGE MASK;
      pr info("Page mask: %llx\n", PAGE MASK);
      // fill out the fast reg work request
      cb->send wr.opcode = IB WR FAST REG MR;
      cb->send wr.num sge = 0;
      cb->send wr.sg list =NULL;
      cb->send wr.next = NULL;
      cb->send wr.send flags = IB SEND SIGNALED;
      cb->send_wr.wr.fast_reg.access_flags =
            IB_ACCESS LOCAL WRITE
                                            IB ACCESS REMOTE READ
IB ACCESS REMOTE WRITE;
      cb->send wr.wr.fast reg.page list = cb->frpl;
      cb->send wr.wr.fast reg.rkey = cb->frmr->rkey;
      cb->send wr.wr.fast reg.page shift = PAGE SHIFT;
      cb->send wr.wr.fast reg.page list len = 8;
      cb->send wr.wr.fast reg.iova start=cb->dma[2];
      cb->send wr.wr.fast reg.length = 8*4096;
      // sleep until the work request finishes/fails
      state = WAITING;
      if(ib post send(cb->qp, &cb->send wr, &cb->bad send)){
            pr err("Failed to post work request to send queue!!!\n");
            goto error2;
      if(wait event interruptible(cb->wqueue, state >=FRMR COMPLETE)){
            pr info("Interrupted\n");
      }
```

```
pr info("State after wakikng: %u\n", state);
      if(state >= DISCONNECT)
           goto error2;
      state = WAITING;
      // state was set to waiting for our upcoming sleep
      // fill out the work requests
      cb->test sge[0].lkey = cb->frmr->lkey;
      cb->test sge[1].lkey = cb->frmr->lkey;
      cb->test sge[0].addr = cb->dma[2];
      cb->test sge[1].addr = cb->dma[2] + 4096;
      cb->test sge[0].length = 4*1024;
      cb->test sge[1].length = 4*1024;
      // I mess with the recv here because I used it AFTER this function. It
can be ignored
      cb->recv wr.next = NULL;
      cb->recv wr.num sge = 1;
      cb->recv wr.sg list = cb->test sge;
      // copy the dma addr and rkey to be send to the remote side
      memcpy(cb->buffs[2] + 4096, &cb->dma[2], sizeof(cb->dma[2]));
     memcpy(cb->buffs[2] + 4096+sizeof(cb->dma[2]), &cb->frmr->rkey,
sizeof(cb->frmr->rkey));
      pr info("Address: %llx\n",(long long unsigned)cb->dma[2]);
      cb->send wr.next = NULL;
      cb->send wr.sg list = &cb->test sge[1];
      cb->send wr.num sge = 1;
      cb->send wr.opcode = IB WR SEND;
      cb->send wr.send flags = IB SEND SIGNALED;
      if(ib post send(cb->qp, &cb->send wr, &cb->bad send)){
           pr err("Failed to post work request to send queue!!!\n");
           goto error2;
      if(wait event interruptible(cb->wqueue, state >=RDMA_SEND_COMPLETE)){
           pr info("Interrupted\n");
```