

# MONTY HALL PROBLEM

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Abstract: an exposition on Monty Hall problem.

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## 1. THE GAME

Consider the following game. A prize is placed behind door 1, door 2 or door 3 with equal probability. So if the random variable  $X$  denotes where the prize is then

$$P(X = 1) = P(X = 2) = P(X = 3) = \frac{1}{3}$$

A player then chooses one of the doors with equal probability. If the door he chooses is where the prize is, he wins the prize. So if the random variable  $Y$  denotes his choice then

$$P(Y = 1) = P(Y = 2) = P(Y = 3) = \frac{1}{3}$$

After the player has made his choice, the host will go behind the doors. If the prize is behind the chosen door, that is  $X = Y$  then he will open one of the remaining two doors with equal probability. If the prize is not behind the chosen door, that is  $X \neq Y$  then he will open the remaining door. So if the random variable  $Z$  denotes his choice then

$$P(Z = 1 | X = 1, Y = 1) = 0$$

$$P(Z = 2 | X = 1, Y = 1) = \frac{1}{2}$$

$$P(Z = 3 | X = 1, Y = 1) = \frac{1}{2}$$

$$P(Z = 1 | X = 1, Y = 2) = 0$$

$$P(Z = 2 | X = 1, Y = 2) = 0$$

$$P(Z = 3 | X = 1, Y = 2) = 1$$

$\vdots$

Gathering all this together, we have the trivariate random variable  $(X, Y, Z)$  with probability mass function  $p_{X,Y,Z}(x, y, z)$  where

$$p_{X,Y,Z}(1, 1, 1) = 0, p_{X,Y,Z}(1, 1, 2) = \frac{1}{18}, p_{X,Y,Z}(1, 1, 3) = \frac{1}{18}$$



$$\begin{aligned}
P(\text{winning by switching}) &= P(X = W) \\
&= \sum_{x=w} p_{X,Y,Z,W}(x, y, z, w) \\
&= \sum_{x \neq y \text{ and } x \neq z} p_{X,Y,Z}(x, y, z) \\
&= \frac{1}{9} + \frac{1}{9} + \frac{1}{9} + \frac{1}{9} + \frac{1}{9} + \frac{1}{9} \\
&= \frac{2}{3}
\end{aligned}$$

Alternatively, one can reason that

$$\begin{aligned}
P(X = W) &= P(X \neq Y \text{ and } X \neq Z) \\
&= P(X \neq Y) \\
&= 1 - P(X = Y) \\
&= 1 - \frac{1}{3} = \frac{2}{3}
\end{aligned}$$

**Example 2.2.** We can compute the distribution of  $Z$

$$\begin{aligned}
P(Z = 1) &= \sum_{z=1} p_{X,Y,Z}(x, y, z) \\
&= \frac{1}{18} + \frac{1}{9} + \frac{1}{9} + \frac{1}{18} \\
&= \frac{1}{3}
\end{aligned}$$

Similarly,  $P(Z = 2) = P(Z = 3) = \frac{1}{3}$ .

**Example 2.3.** We can compute the distribution of  $W$

$$\begin{aligned}
P(W = 1) &= \sum_{w=1} p_{X,Y,Z,W}(x, y, z, w) \\
&= \sum_{y \neq 1 \text{ and } z \neq 1} p_{X,Y,Z}(x, y, z) \\
&= \frac{1}{9} + \frac{1}{9} + \frac{1}{18} + \frac{1}{18} \\
&= \frac{1}{3}
\end{aligned}$$

Similarly,  $P(W = 2) = P(W = 3) = \frac{1}{3}$ .

All  $X, Y, Z, W$  are identically distributed.

### 3. INFORMATION AND ENTROPY

If we think the host's action makes a difference in this game, can we quantify it? That brings up information and entropy. We limit ourselves to discrete random variables.

**Definition 3.1.** For each random variable  $(\Omega, \mathcal{F}, P) \xrightarrow{X} (\mathbb{R}, \mathcal{B}(\mathbb{R}))$ , we define its information to be the random variable

$$(\mathbb{R}, \mathcal{B}(\mathbb{R}), X_*(P)) \xrightarrow{I_X} (\mathbb{R}, \mathcal{B}(\mathbb{R}))$$

$$x \mapsto \begin{cases} 0 & \text{if } p_X(x) = 0 \\ \log_2\left(\frac{1}{p_X(x)}\right) & \text{otherwise} \end{cases}$$

Note that  $p_X(x) = X_*(P)(x)$ . Here the base is chosen to be 2, and the unit of  $I_X$  is called *bit*. Other popular bases are 10 with unit *nat* and  $e$  with unit *hartley*. All results then differ by constant scalars  $\log_{10} 2$  and  $\log_e 2$  respectively.

**Example 3.2.** The joint random variable  $(X, Y)$  has information

$$(\mathbb{R}^2, \mathcal{B}(\mathbb{R}^2), (X, Y)_*(P)) \xrightarrow{I_{X,Y}} (\mathbb{R}, \mathcal{B}(\mathbb{R}))$$

$$(x, y) \mapsto \begin{cases} 0 & \text{if } p_{X,Y}(x, y) = 0 \\ \log_2\left(\frac{1}{p_{X,Y}(x, y)}\right) & \text{otherwise} \end{cases}$$

Note that  $p_{X,Y}(x, y) = (X, Y)_*(P)(x, y)$ . We will use this in example 3.10.

**Definition 3.3.** For each random variable  $X$ , we define its entropy  $H(X)$  to be the expected value  $E(I_X)$  of  $I_X$ .

Explicitly, one has

$$H(X) = \sum_{x_i} p_X(x_i) I_X(x_i)$$

$$= \sum_{x_i} p_X(x_i) \log_2\left(\frac{1}{p_X(x_i)}\right)$$

For each pair of random variables  $X, Y$ , given event  $Y = y$ , we can think of

- conditional probability measure  $P(-|Y = y)$
- its pushforward  $X_*(P(-|Y = y))$
- random variable  $(\mathbb{R}, \mathcal{B}(\mathbb{R}), X_*(P(-|Y = y))) \xrightarrow{X|Y=y} (\mathbb{R}, \mathcal{B}(\mathbb{R}))$
- its information  $I_{X|Y=y}$

Note that  $p_{X|Y=y}(x) = X_*(P(-|Y = y))(x)$ .

**Definition 3.4.** For each pair of random variables  $X, Y$ , we define the entropy  $H(X|Y = y)$  of  $X$  given an event  $Y = y$  to be the expected value  $E(I_{X|Y=y})$  of  $I_{X|Y=y}$ .

Explicitly, one has

$$H(X|Y = y) = \sum_{x_i} p_{X|Y=y}(x_i) I_{X|Y=y}(x_i)$$

$$= \sum_{x_i} p_{X|Y=y}(x_i) \log_2\left(\frac{1}{p_{X|Y=y}(x_i)}\right)$$

If we view  $(\mathbb{R}, \mathcal{B}(\mathbb{R}), Y_*(P)) \xrightarrow{H(X|Y=-)} (\mathbb{R}, \mathcal{B}(\mathbb{R})), y \mapsto H(X|Y=y)$  as a random variable then it has expected value

$$\begin{aligned}
E(H(X|Y=-)) &= \sum_{y_j} p_Y(y_j) H(X|Y=y_j) \\
&= \sum_{y_j} p_Y(y_j) \sum_{x_i} p_{X|Y=y_j}(x_i) \log_2 \left( \frac{1}{p_{X|Y=y_j}(x_i)} \right) \\
&= \sum_{x_i, y_j} p_Y(y_j) p_{X|Y=y_j}(x_i) \log_2 \left( \frac{1}{p_{X|Y=y_j}(x_i)} \right) \\
&= \sum_{x_i, y_j} p_Y(y_j) \frac{p_{X,Y}(x_i, y_j)}{p_Y(y_j)} \log_2 \left( \frac{1}{p_{X|Y=y_j}(x_i)} \right) \\
&= \sum_{x_i, y_j} p_{X,Y}(x_i, y_j) \log_2 \left( \frac{p_Y(y_j)}{p_{X,Y}(x_i, y_j)} \right)
\end{aligned}$$

**Definition 3.5.** For each pair of random variables  $X, Y$  we define the entropy  $H(X|Y)$  of  $X$  given  $Y$  to be the expected value  $E(H(X|Y=-))$  of  $H(X|Y=-)$  above.

The difference  $H(X) - H(X|Y)$  is called mutual information between  $X$  and  $Y$  and denoted by  $I(X, Y)$ . It also equals the mutual information  $I(Y, X) = H(Y) - H(Y|X)$  between  $Y$  and  $X$  and has explicit formula

$$I(X, Y) = \sum_{i,j} p_{X,Y}(x_i, y_j) \log_2 \left( \frac{p_{X,Y}(x_i, y_j)}{p_X(x_i)p_Y(y_j)} \right)$$

Now we can compute all informations and entropies in the game.

**Example 3.6.** Without the host's action, the prize location  $X$  has entropy

$$\begin{aligned}
H(X) &= \frac{1}{3} \log_2 \left( \frac{1}{\frac{1}{3}} \right) + \frac{1}{3} \log_2 \left( \frac{1}{\frac{1}{3}} \right) + \frac{1}{3} \log_2 \left( \frac{1}{\frac{1}{3}} \right) \\
&= \log_2 3
\end{aligned}$$

**Example 3.7.** With the host's action of opening door 3, the prize location  $X|Z=3$  has entropy

$$\begin{aligned}
H(X|Z=3) &= p_{X|Z=3}(1) \log_2 \left( \frac{1}{p_{X|Z=3}(1)} \right) + p_{X|Z=3}(2) \log_2 \left( \frac{1}{p_{X|Z=3}(2)} \right) + p_{X|Z=3}(3) \log_2 \left( \frac{1}{p_{X|Z=3}(3)} \right) \\
&= \frac{1}{2} \log_2 \left( \frac{1}{\frac{1}{2}} \right) + \frac{1}{2} \log_2 \left( \frac{1}{\frac{1}{2}} \right) + 0 \\
&= 1
\end{aligned}$$

With the host's action of opening door 3, the entropy of  $X$  decreases. Similarly  $H(X|Z=1) = H(X|Z=2) = 1$ .

**Example 3.8.** With the host's action, the prize location has expected entropy

$$\begin{aligned}
H(X|Z) &= p_Z(1)H(X|Z=1) + p_Z(2)H(X|Z=2) + p_Z(3)H(X|Z=3) \\
&= \frac{1}{3} \cdot 1 + \frac{1}{3} \cdot 1 + \frac{1}{3} \cdot 1 \\
&= 1
\end{aligned}$$

One can verify this directly with the explicit formula in definition 3.5 as well.

**Example 3.9.** The information the host provides about the prize location is

$$\begin{aligned}
I(X, Z) &= 0 + \left(\frac{1}{18} + 0 + \frac{1}{9}\right) \log_2 \left(\frac{\frac{1}{18} + 0 + \frac{1}{9}}{\frac{1}{3} \cdot \frac{1}{3}}\right) + \left(\frac{1}{18} + \frac{1}{9} + 0\right) \log_2 \left(\frac{\frac{1}{18} + \frac{1}{9} + 0}{\frac{1}{3} \cdot \frac{1}{3}}\right) \\
&\quad + \left(0 + \frac{1}{18} + \frac{1}{9}\right) \log_2 \left(\frac{0 + \frac{1}{18} + \frac{1}{9}}{\frac{1}{3} \cdot \frac{1}{3}}\right) + 0 + \left(\frac{1}{9} + \frac{1}{18} + 0\right) \log_2 \left(\frac{\frac{1}{9} + \frac{1}{18} + 0}{\frac{1}{3} \cdot \frac{1}{3}}\right) \\
&\quad + \left(0 + \frac{1}{9} + \frac{1}{18}\right) \log_2 \left(\frac{0 + \frac{1}{9} + \frac{1}{18}}{\frac{1}{3} \cdot \frac{1}{3}}\right) + \left(\frac{1}{9} + 0 + \frac{1}{18}\right) \log_2 \left(\frac{\frac{1}{9} + 0 + \frac{1}{18}}{\frac{1}{3} \cdot \frac{1}{3}}\right) + 0 \\
&= \log_2 \frac{3}{2} \\
&= \log_2 3 - 1
\end{aligned}$$

One can see that  $I(X, Z) = H(X) - H(X|Z)$ . This is how much the entropy of  $X$  decreases by.

**Example 3.10.** The joint random variable  $(X, Z)$  has entropy

$$\begin{aligned}
H(X, Z) &= E(I_{X,Z}) \\
&= \sum_{x_i, z_k} p_{X,Z}(x_i, z_k) \log_2 \left(\frac{1}{p_{X,Z}(x_i, z_k)}\right) \\
&= \frac{1}{6} \log_2 \left(\frac{1}{\frac{1}{6}}\right) + \frac{1}{6} \log_2 \left(\frac{1}{\frac{1}{6}}\right) + \frac{1}{6} \log_2 \left(\frac{1}{\frac{1}{6}}\right) + \frac{1}{6} \log_2 \left(\frac{1}{\frac{1}{6}}\right) + \frac{1}{6} \log_2 \left(\frac{1}{\frac{1}{6}}\right) + \frac{1}{6} \log_2 \left(\frac{1}{\frac{1}{6}}\right) \\
&= \log_2 6 \\
&= \log_2 3 + 1
\end{aligned}$$

One can see that  $H(X, Z) = H(X) + H(Z) - I(X, Z) = H(X|Z) + I(X, Z) + H(Z|X)$ . Draw a Venn diagram.

**Exercise 3.11.** Compute  $H(Y)$ ,  $H(Y|Z = 3)$ ,  $H(Y|Z)$ ,  $I(Y, Z)$ ,  $H(Y, Z)$ . Hint: swap prize and player.

#### 4. SIMULATION

See Python code and games at [github.com/dinhuun/probability\\_statistics](https://github.com/dinhuun/probability_statistics).