

Wireshark IP

ip-ethereal-trace-1

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*Apply a display filter ... <Ctrl>/

No.	Time	Source	Destination	Protocol	Length	Info
10.000000		Telebit 73:8d:ce	Broadcast	ARP	60	Who has 192.168.1.11? Tell 192.168.1.104
24.866867		192.168.1.100	192.168.1.1	SSDP	174	M-SEARCH * HTTP/1.1
34.868147		192.168.1.100	192.168.1.1	SSDP	175	M-SEARCH * HTTP/1.1
45.363536		192.168.1.100	192.168.1.1	SSDP	174	M-SEARCH * HTTP/1.1
55.364799		192.168.1.100	192.168.1.1	SSDP	175	M-SEARCH * HTTP/1.1
65.864428		192.168.1.100	192.168.1.1	SSDP	174	M-SEARCH * HTTP/1.1
75.865461		192.168.1.100	192.168.1.1	SSDP	175	M-SEARCH * HTTP/1.1
86.163045		192.168.1.102	128.59.23.100	ICMP	98	Echo (ping) request id=0x0300, seq=20483/848, ttl=1 (no response found!)
96.176826		192.168.1.102	128.59.23.100	ICMP	98	Time-to-live exceeded (Time to live exceeded in transit)
106.188629		192.168.1.102	128.59.23.100	ICMP	98	Echo (ping) request id=0x0300, seq=20739/849, ttl=2 (no response found!)
116.202957		192.168.1.102	128.59.23.100	ICMP	98	Time-to-live exceeded (Time to live exceeded in transit)
126.208997		192.168.1.102	128.59.23.100	ICMP	98	Echo (ping) request id=0x0300, seq=20995/850, ttl=3 (no response found!)
136.234505		192.168.1.102	128.59.23.100	ICMP	98	Time-to-live exceeded (Time to live exceeded in transit)
146.238695		192.168.1.102	128.59.23.100	ICMP	98	Echo (ping) request id=0x0300, seq=21251/851, ttl=4 (no response found!)
156.252722		192.168.1.102	128.59.23.100	ICMP	98	Time-to-live exceeded (Time to live exceeded in transit)

> Frame 8: 98 bytes on wire (784 bits), 98 bytes captured (784 bits)

> Ethernet II, Src: Actionte_8a:70:1a (00:20:e0:8a:70:1a), Dst: LinksysG_da:af:73 (00:06:25:da:af:73)

> Internet Protocol Version 4, Src: 192.168.1.102, Dst: 128.59.23.100

0100 = Version: 4

.... 0101 = Header Length: 20 bytes (5)

> Differentiated Services Field: 0x00 (DSCP: CS0, ECN: Not-ECT)

Total Length: 84

Identification: 0x32d0 (13008)

> Flags: 0x00

Fragment offset: 0

> Time to live: 1

Protocol: ICMP (1)

Header checksum: 0x2d2c [validation disabled]

[Header checksum status: Unverified]

0000 00 06 25 da af 73 00 20 e0 8a 70 1a 08 00 45 00 ...S...p...E

0010 00 54 32 d0 00 01 01 2d 2c c0 a8 01 66 80 3b ...T2....f...

0020 17 c4 08 00 f7 ca 03 00 50 03 37 32 20 aa aa aa ...d.....P.72...

0030 aa aa aa aa aa aa aa aa aa aa aa aa aa aa aa aa

0040 aa aa aa aa aa aa aa aa aa aa aa aa aa aa aa aa

0050 aa aa aa aa aa aa aa aa aa aa aa aa aa aa aa aa

0060 aa aa

Frame (frame), 98 bytes

Packets: 380 - Displayed: 380 (100.0%) - Load time: 0:0.5

Profile: Def

1. Select the first ICMP Echo Request message sent by your computer, and expand the Internet Protocol part of the packet in the packet details window. What is the IP address of your computer?

192.168.1.102

2. Within the IP packet header, what is the value in the upper layer protocol field?
3. How many bytes are in the IP header? How many bytes are in the payload of the IP datagram? Explain how you determined the number of payload bytes.

.... 0101 = Header Length: 20 bytes (5)

Header length: 20 bytes. There are 20 bytes in the IP header which leaves 36 bytes for the payload of the IP datagram because we were sending a packet of length 56 bytes.

4. Has this IP datagram been fragmented? Explain how you determined whether or not the datagram has been fragmented

Identification: 0x32d0 (13008)

> Flags: 0x00

0... = Reserved bit: Not set

.0.. = Don't fragment: Not set

..0. = More fragments: Not set

Fragment offset: 0

.... 0101 = Header Length: 20 bytes (5) No. Fragment offset: 0

5. Which fields in the IP datagram always change from one datagram to the next within this series of ICMP messages sent by your computer?

The header checksum, time to live and the identification changes from each datagram to the next.

6. Which fields stay constant? Which of the fields must stay constant? Which fields must change? Why?

Fields that stay constant are the same as fields must stay constant: Version (Ipv4...), length of the header (ICMP packets), source IP (sending the same place), destination IP (contacting the same

place), upper layer protocol (because of ICMP), Differentiated Services (since all packets are ICMP they use the same Type of Service class).

Fields that must change: Header checksum (Header changes), identification (Each packets must have different IP), time to live (traceroute increments each subsequent packet)

7. Describe the pattern you see in the values in the Identification field of the IP datagram.

The pattern is that the IP header Identification fields increment with each ICMP Echo request.

8. What is the value in the Identification field and the TTL field?

Identification: 0x32d0 (13008)

> Flags: 0x00

Fragment offset: 0

> Time to live: 1

Identification: 0x32d0 (13008)

TTL: 1

9. Do these values remain unchanged for all of the ICMP TTL-exceeded replies sent to your computer by the nearest (first hop) router? Why?

The identification field changes for all the ICMP TTL-exceeded replies because the identification field is a unique value. When two or more IP datagrams have the same identification value, then it means that these IP datagrams are fragments of a single large IP datagram. The TTL field remains unchanged because the TTL for the first hop router is always the same.

10. Find the first ICMP Echo Request message that was sent by your computer after you changed the Packet Size in pingplotter to be 2000. Has that message been fragmented across more than one IP datagram?

Yes

11. Print out the first fragment of the fragmented IP datagram. What information in the IP header indicates that the datagram been fragmented? What information in the IP header indicates whether this is the first fragment versus a latter fragment? How long is this IP datagram?

The datagram has been fragmented because the flags bit for more fragment is set. The fragment offset is 0 so that this is the first fragment.

The first datagram has the length of 1500.

