

13/09/2022

13 September 2022 19:01

input data: \rightarrow 2 3 5 6 4 8 1 7 9 0 order

arr \rightarrow (0, 1, 2, 3, 9, 5, 6, 1, 2, 8, 1, 7)

output data: 2 3 5 6 8 1 1 9 0

let order = [2, 3, 5, 6, 4, 8, 1, 7, 9, 0]

let op = []
for i =

order.forEach((i) => T

let cut = 0

{ init how many times i is there
tr and push }
qnn.foreach((j) => T

T + (j == i)
cut++

i }

push i
(cut times order+1)
S while (cut > 0){
op.push(i)
cut--1

going through all elements

return op:

return of:

$$\left\{ \begin{array}{l} A \xrightarrow{x} 1 \text{ltr} \rightarrow K \\ B \xrightarrow{y} 1 \text{ltr} \rightarrow L \end{array} \right\} X \xrightarrow{z} \{ q_1, q_2, q_3, q_4, \dots, q_n \}$$

input data

$$\begin{aligned} & A \rightarrow X, 1 \text{ltr}, \quad 1 \text{ltr} = k \rightarrow k \\ & B \rightarrow Y, 1 \text{ltr}, \quad 1 \text{ltr} = l \rightarrow l \\ & q = [] - \underbrace{\text{valts}}_{c} \quad (k, l, x, y) \end{aligned}$$

outp & desc: no of holds crossed
by size for each bin

$$\frac{S(n)}{2} \times K \times K = \boxed{k^n}, \quad k^n$$

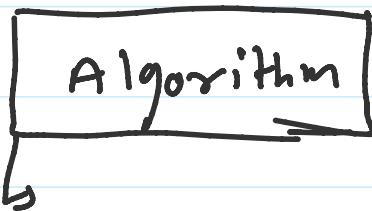
$$q = [q_1, q_2, q_3, \dots, q_n]$$

for each (i) \Rightarrow {
if ($k^n - i \geq 0$)
 $c_{n-i} +$ }

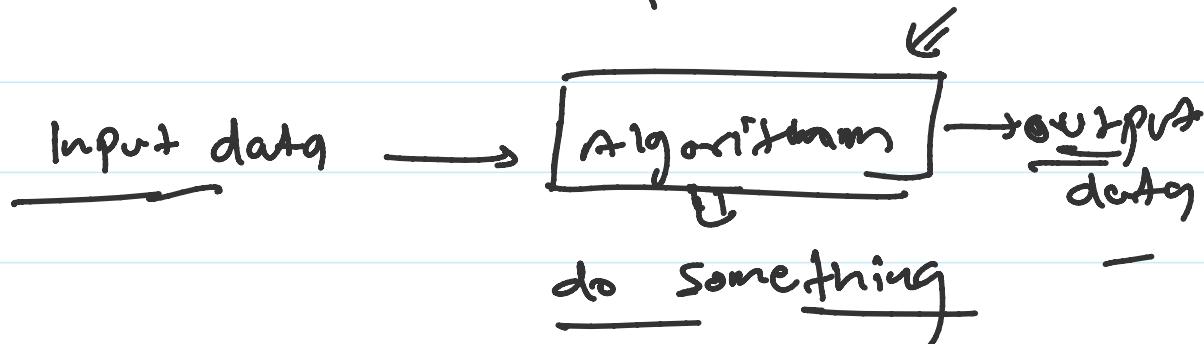
$\text{if } (x_n - 1 \geq 0)$
 $\text{cut}++$
 $k_n = k_n - i$

return cut;

Time Complexity:



{ a formalized way to {
solve problem.



input data: $a[0, 1, 2, \dots, n]$

$$\sum_{i=0}^{n-1} a[i]$$

$i=0$

Sum = 0

for $i = 0 \dots n-1$:
 $\dots \rightarrow i$

for $i = 0 \dots n-1$

 sum += i;

 print(sum);

- 1) Is it correct?
- 2) What is time complexity?
- 3) Can we do better?

Time Complexity:

→ time taken by your processor to solve problem?

in physics 1 sec, 2sec, 1 msec

1 sec

? i5 9th generation → my laptop process
? i3 10th generation → what?

2sec

So what we can do now? How to measure time complexity?

We measure time complexity in terms
of "no. of operations"

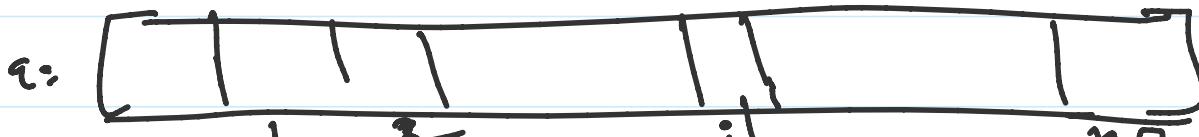
Computational mode)

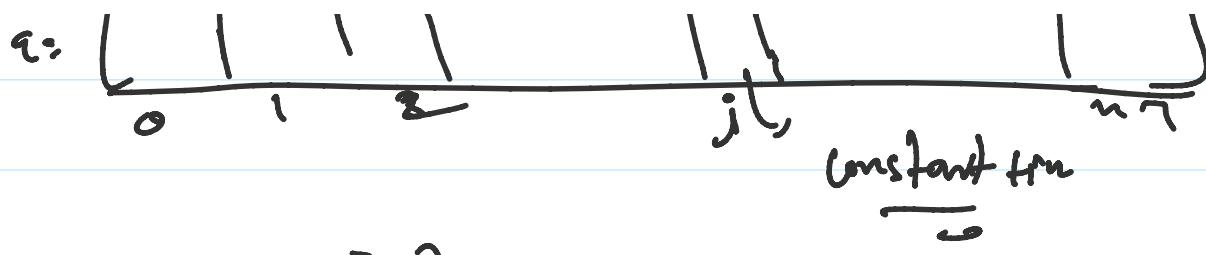
↳ a mathematical model which
kind of simulate our computer
processor.

- ↳ RAM-machine | mode |
- ↳ pointer-machine | mode |

RAM-machine | mode | :

(random access memory)





$a[i]$ -

↳ what are the operations?

set of of operations :-

↳ (+, -, *, /, %) →

↳ (declaring a variable)

↳ (assigning a value to variable)

↳ access any element of array ($a[i]$)

↳ changing value of memory at location

$T[i] = v$

Time complexity :- "no. of operations"

2 ^b $\text{Sum} = 0$ for (for loop iteration)

for i=0 --- n = 2

 \sum sum += i; $\frac{2x}{2} = \sum + i$

 print (sum);

- - - - - ← -

$$\text{Time complexity} = \frac{5n^2 + 3}{3n + 2}$$

operation depends upon size of input

$$T.C. = \frac{5n^2 + 3}{3n + 2} \quad \geq 1000$$

it $n > 1000$ (we can ignore)

$T.C. = \frac{5n^2}{3n}$ is not depends on Algorithm

depends upon your programming skills

$$T.C. = n \quad O(n)$$

notation | asymptotic

$O()$ $\Omega()$ $\Theta()$	$\xrightarrow{\text{upper bound}}$ $\xrightarrow{\text{lower bound}}$ $\rightarrow \text{fixed}$	At most At least fixed
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omeg-(Θ) \rightarrow fixed / fixed
 theta- $\Theta()$ \subseteq

Big-O ($O(D)$) //

$$\underline{f = O(g)} \Rightarrow \boxed{\underline{f \leq g}}$$

$\forall g, c, n_0$

$$\underline{h \geq n_0} \quad \underline{f \leq cg + c}$$

$$f = cn \quad \underline{out}$$

$$g = \underline{c_1 n^k} \quad \underline{or 2^n}$$

$$\begin{matrix} n & \geq & c_1 n^k \\ \cancel{n} & \cancel{\geq} & \cancel{n^k} \\ \cancel{n} & \cancel{\geq} & \cancel{n^k} \end{matrix}$$

$$\underline{f = O(g)}$$

O - upper bound / not more than g

\Rightarrow how much time your algo. will take

if it won't take more than usec

Up to it won't take more than $\underline{\underline{n}}$ sec

$$\Theta(f + O(g)) = f \geq g$$

$$T = \underline{5n+3} = \underline{\underline{O(n)}}$$

$$\begin{aligned} n &\rightarrow \text{algo will take at least} \\ &= 1 \text{ sec} \\ T &= \underline{5n+3} = \underline{\underline{O(n)}} \end{aligned}$$

when O & $\underline{\underline{O}}$ are same \rightarrow treat

$$\underline{\underline{O(n)}}$$

$n^{(m+2)}$

for $i=0 \dots m-2$

un [for $j=0 \dots n-1 \dots -2$

$x_{m+j} = a[i][j] \dots -L$

$$\begin{aligned} T_C &2 & O(n^m) & (m+2)n \\ \in & \in & \in & n^2 + 2n \end{aligned}$$

\approx

$\overline{\approx}$

$\frac{nh^2 + 2h}{\cancel{h}}$

$O(n^2)$

$\overline{\approx}$