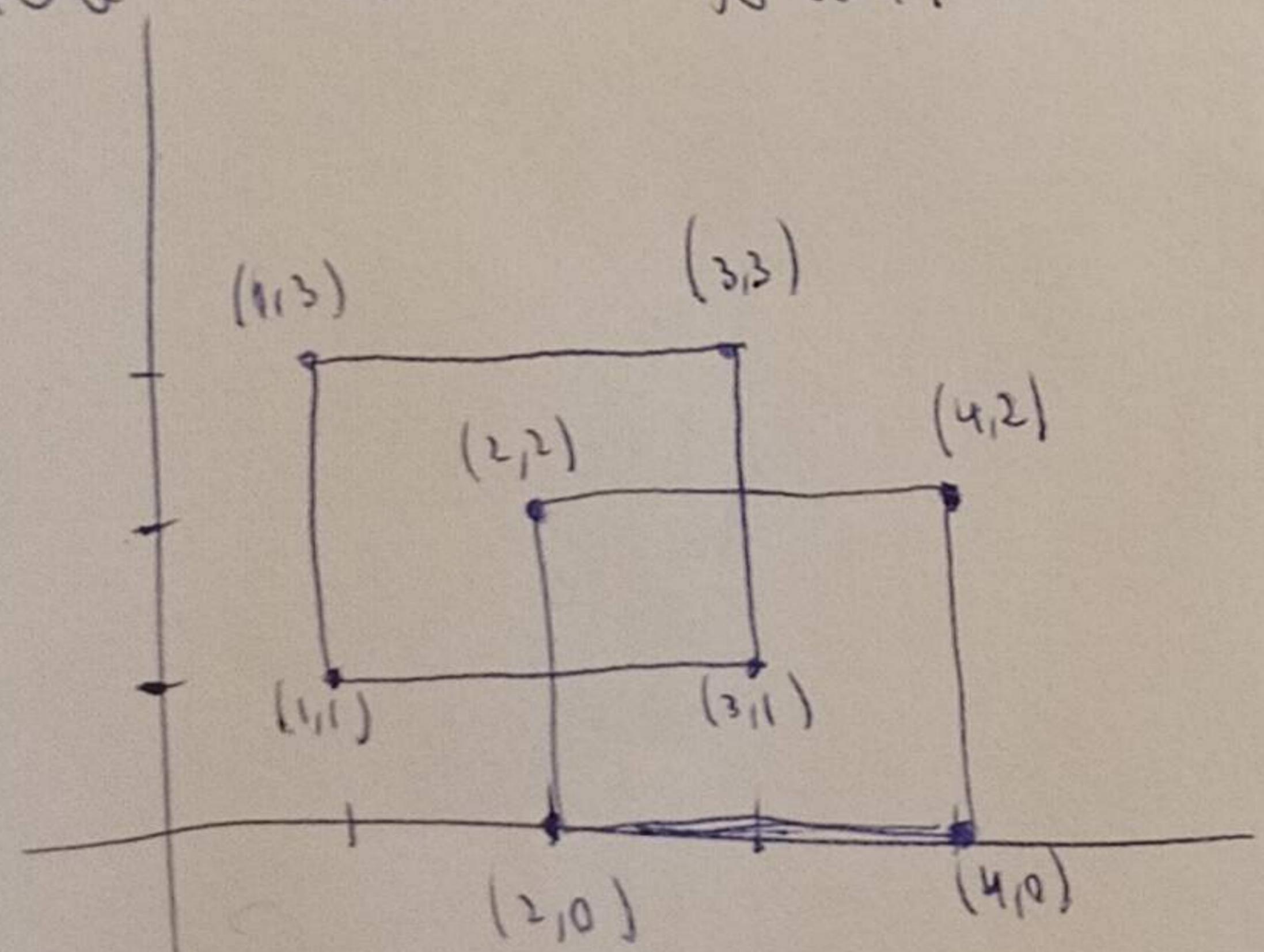


$$(a_0, a_3, a_2, a_1), (N_0, N_1, N_2, N_3) \rightarrow (e_0, e_1, e_2, e_3, e_4, e_5, e_6, e_7)$$

- intersections become points.
- points within rows get deleted.
- there must be rows!



- for each point check whether it is contained in any polygon \rightarrow Hashmap
- make a mask with ring = len(all points)
marks flagged one to be deleted

How to tell that a point is contained in a polygon?

I need now to find intersections
(and we use Hash to m^2)

~~many overlapping polygons~~

- assume one root polygon at a time (remove = "empty" (possible))

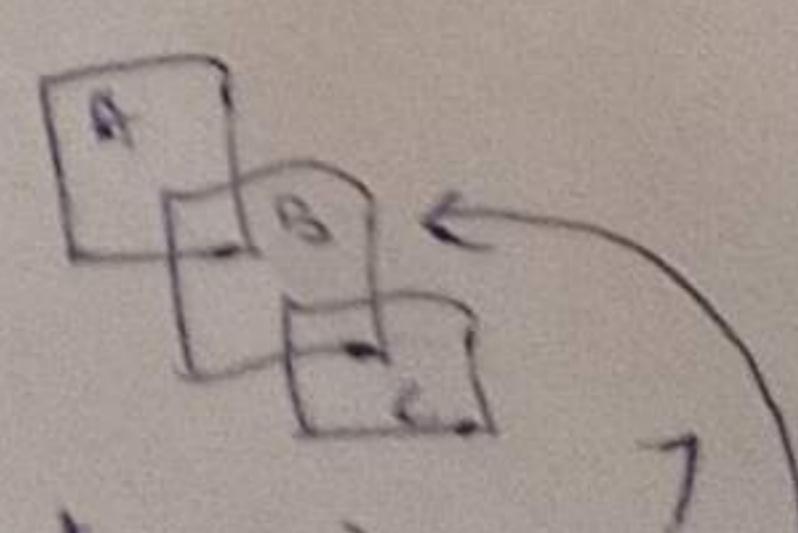
- iterate over lines of points

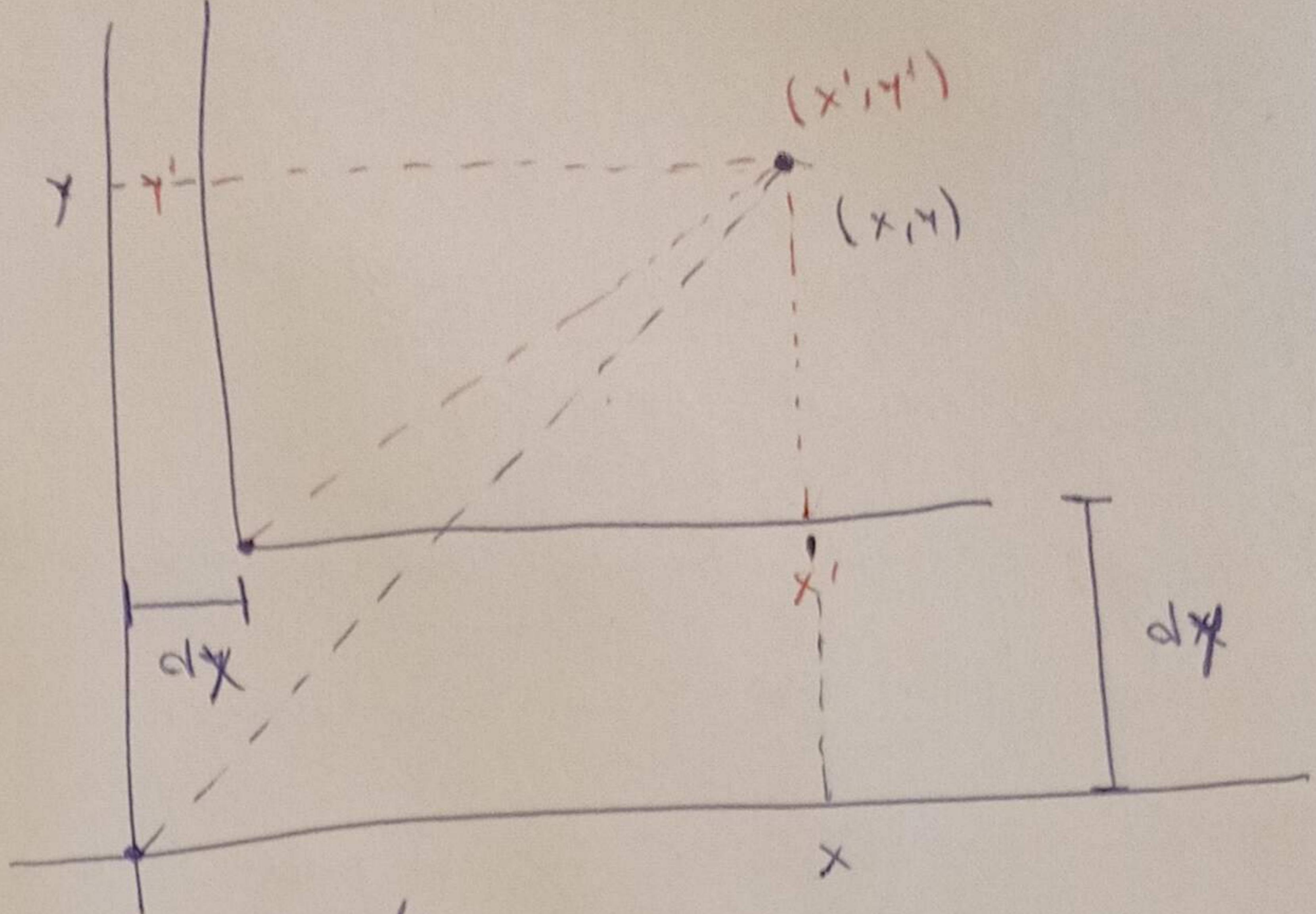
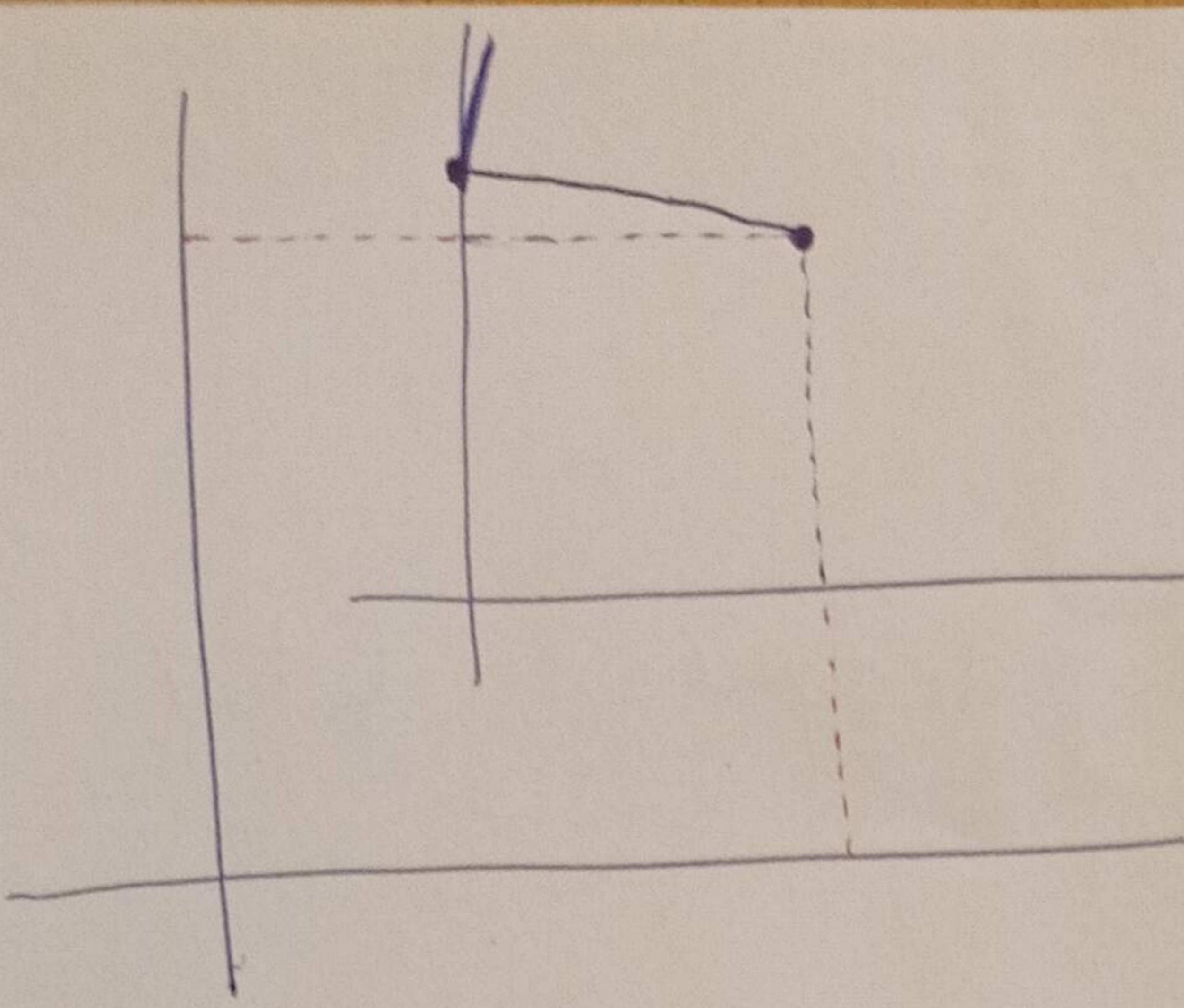
- What if polygons are disjoint?

and furthermore what if they ^(A, B) get joined together through an intermediate?

if A gets to C first and B second how would it work?

- So first always generate a polygon not within them a master polygon
- "lines" (finite) segments! illustration

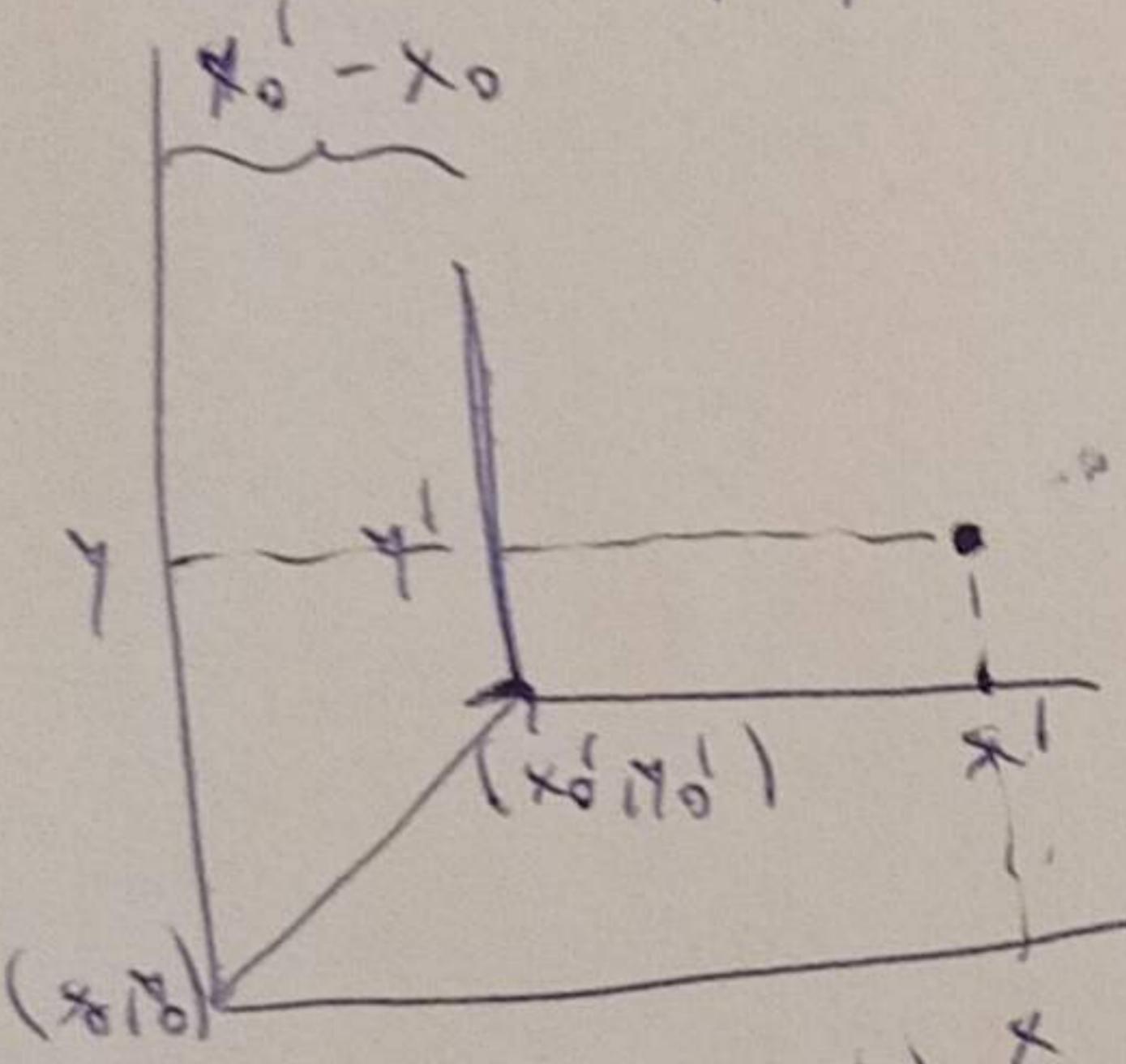




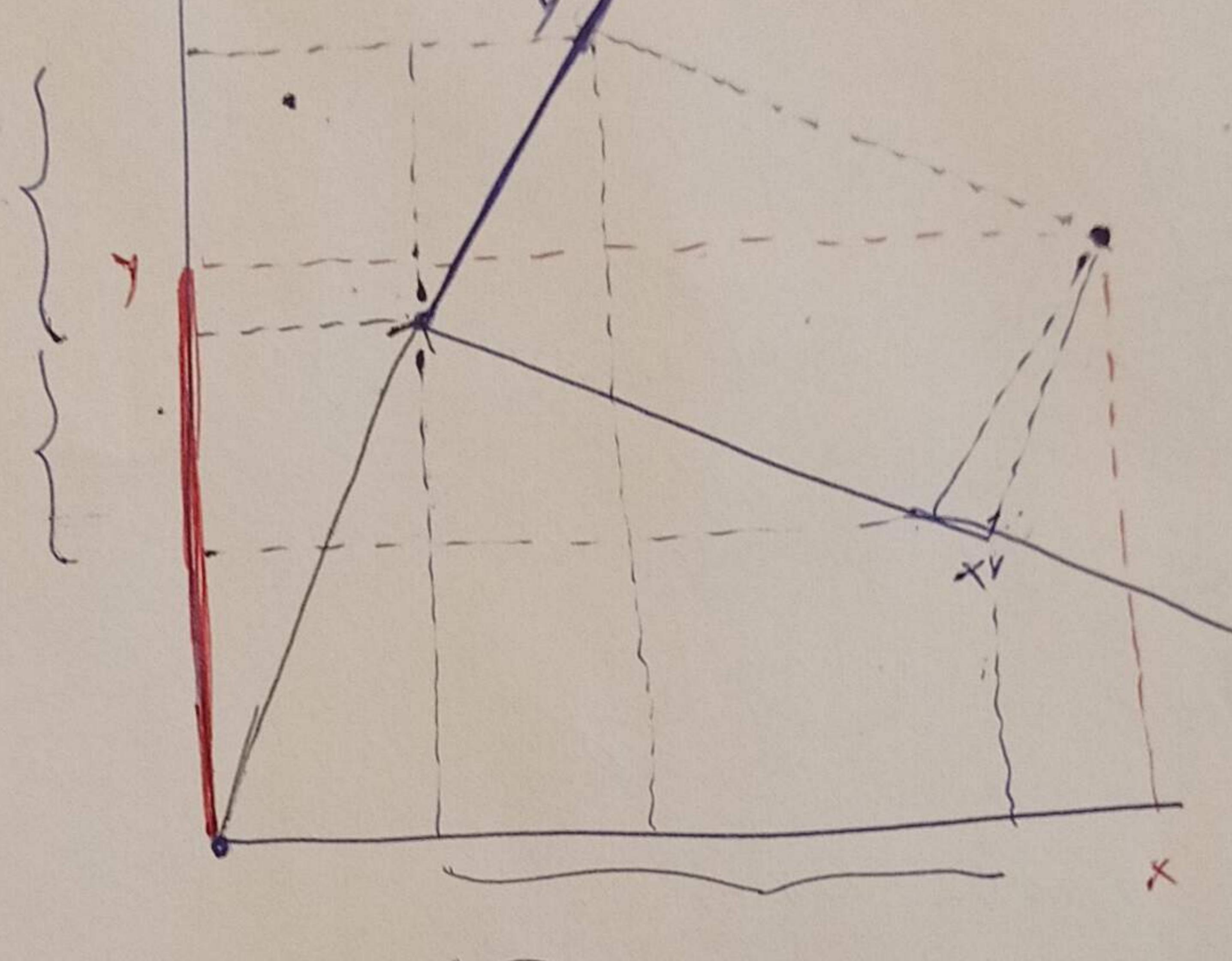
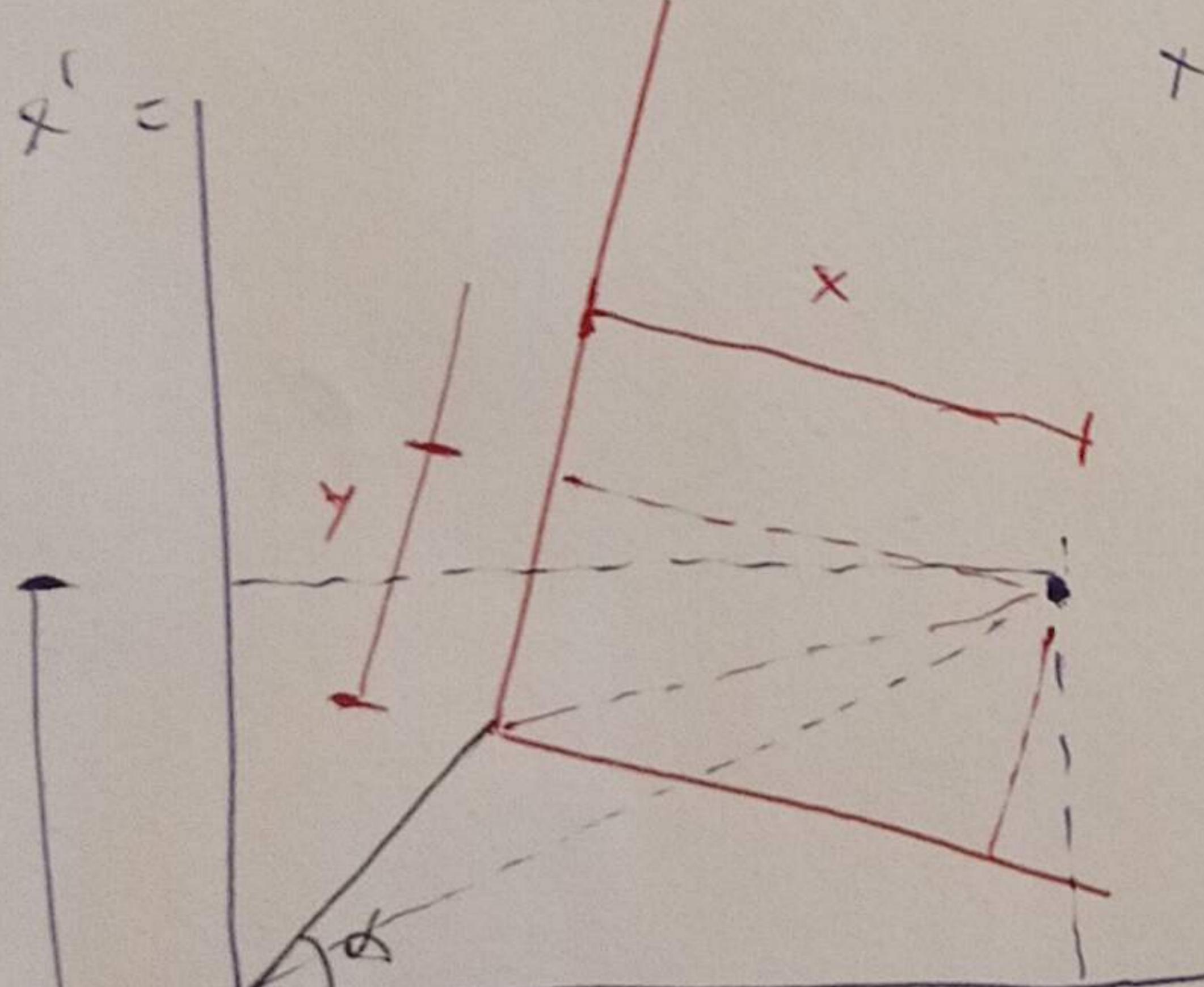
$$x = dy + y'$$

$$x = dx + x'$$

$$\gamma = \alpha x' + \beta y'$$

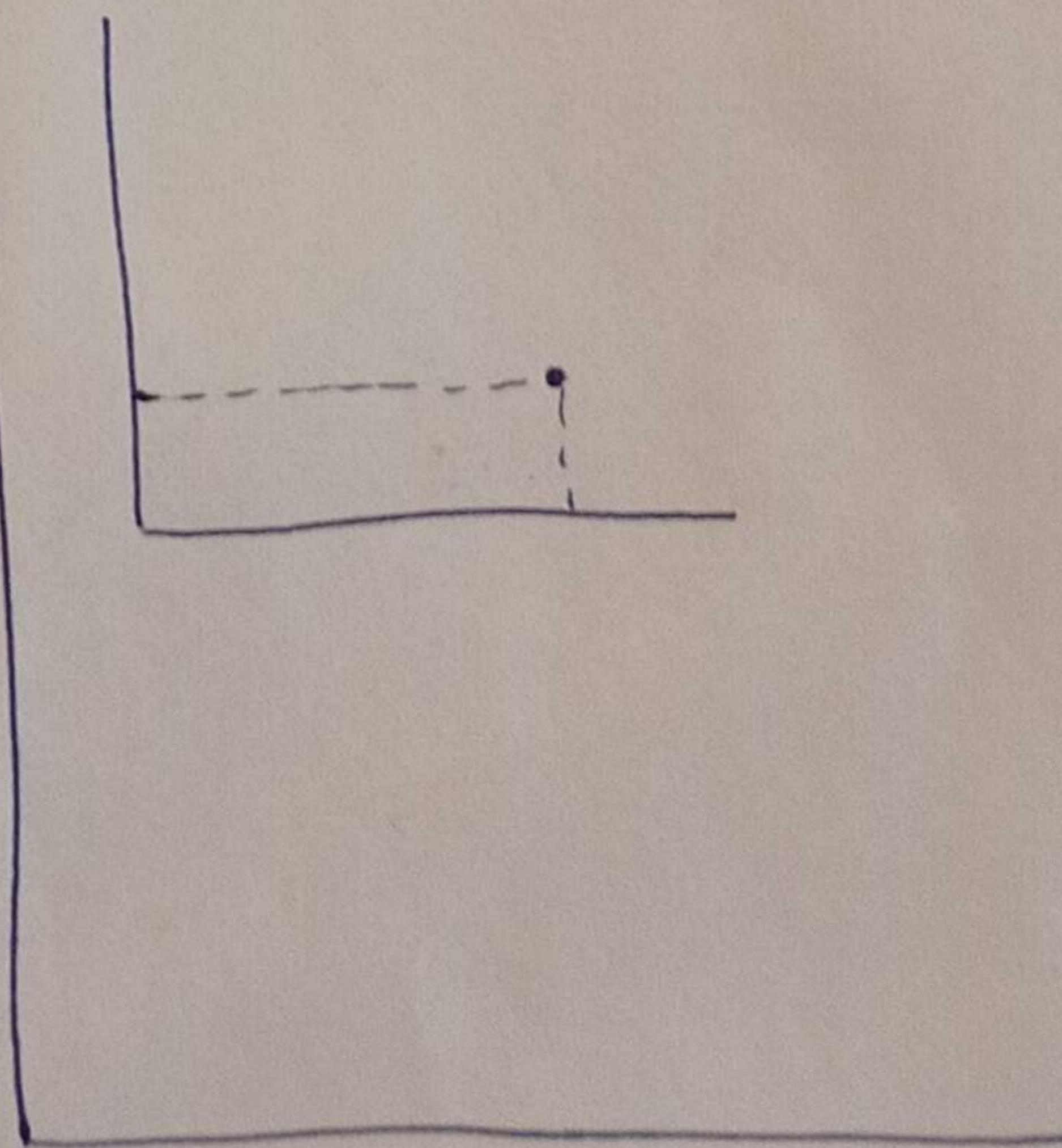
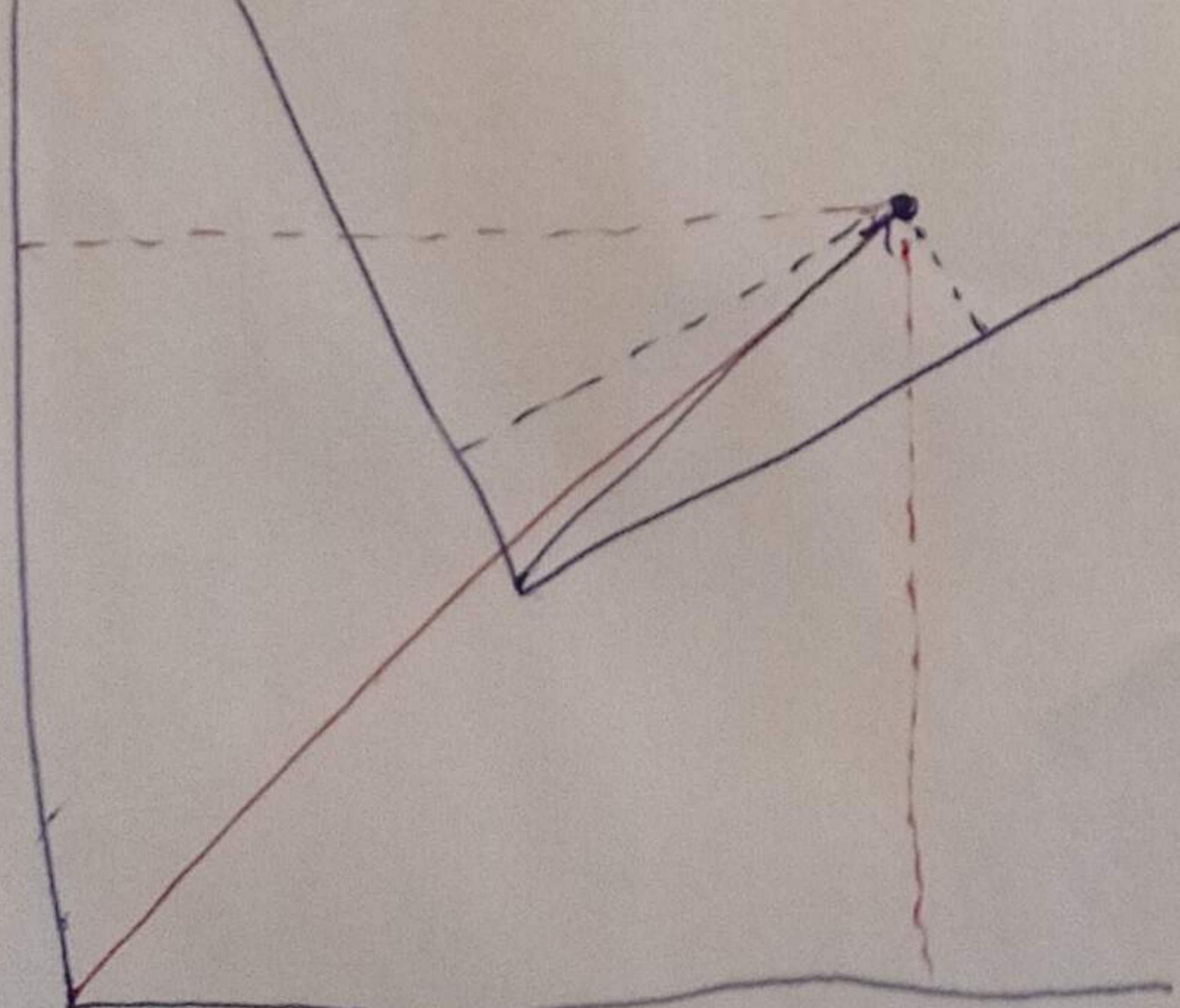
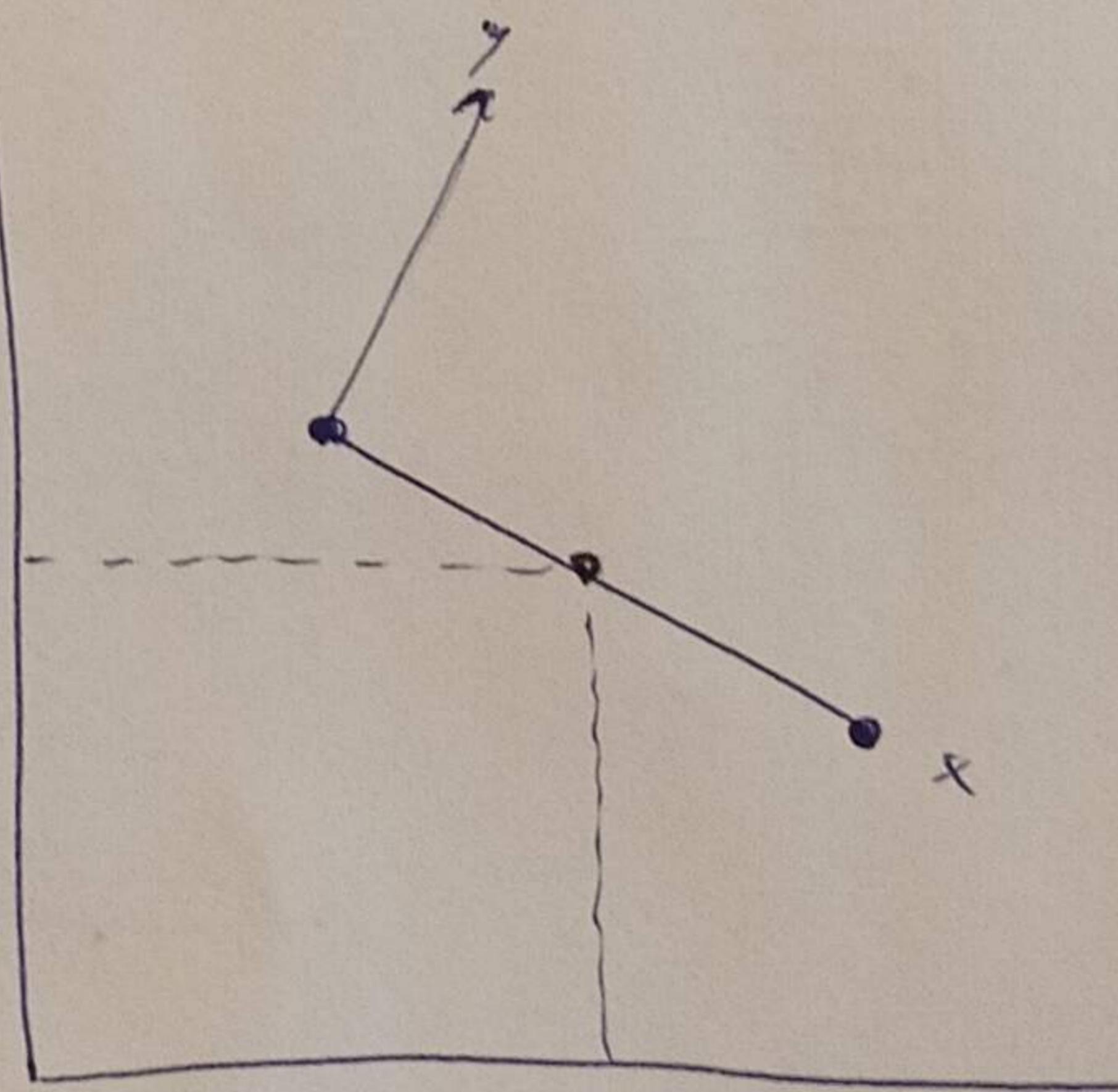


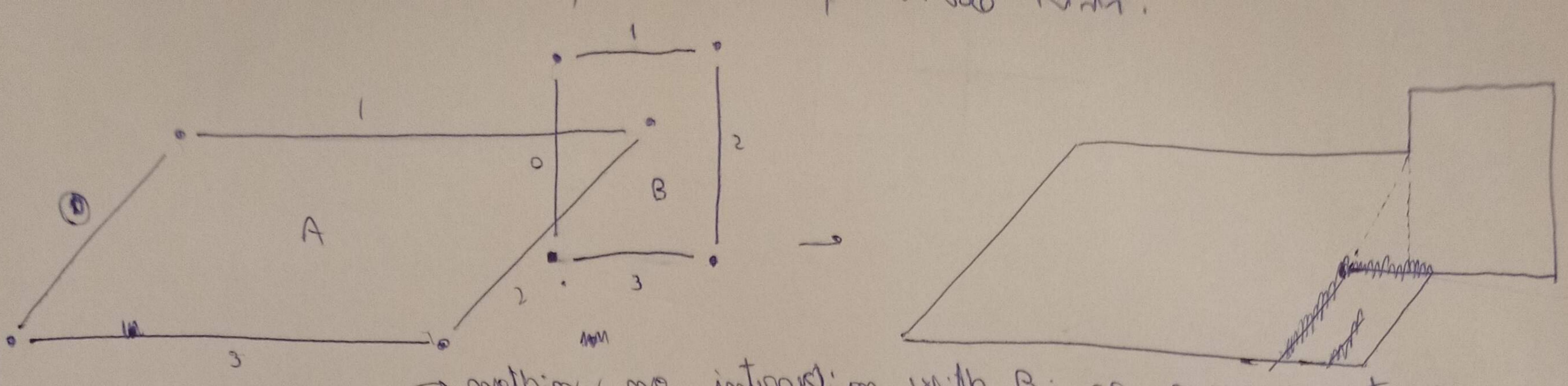
$$y^* = \gamma - (\gamma_0^* - \gamma_0)$$



$$\gamma = (\gamma_0^* - \gamma_0) + \gamma^*$$

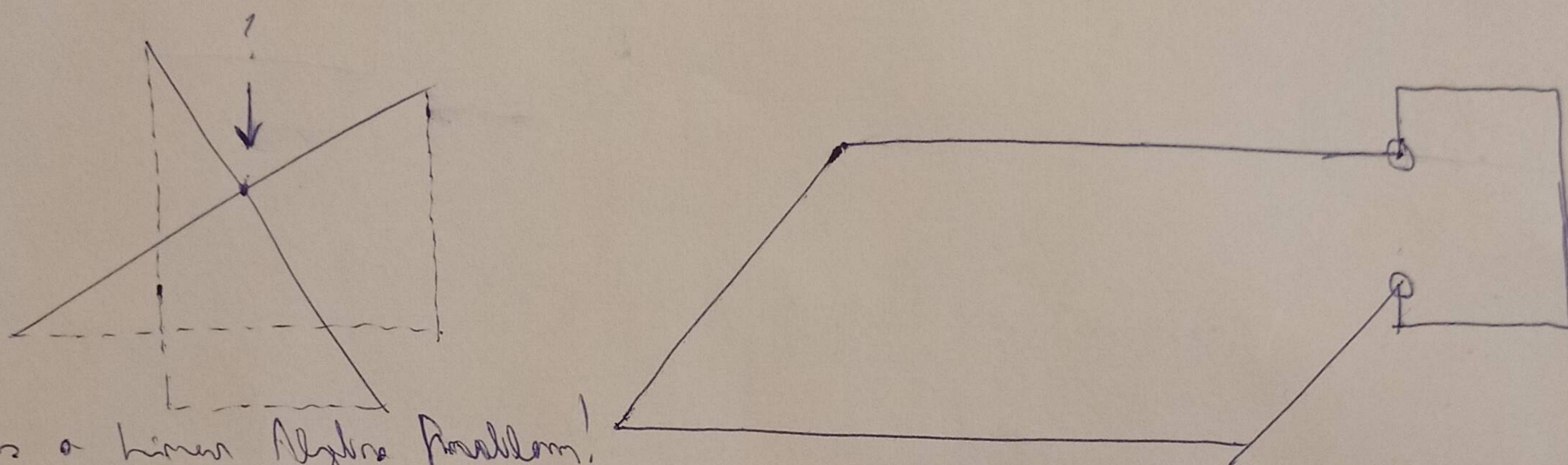
$$x = (x_0^* - x_0) + x'$$





→ nothing, no intersection with B_i so move on to A_j
 Start with segment A_0 , loop through segments of B starting with B_0 .

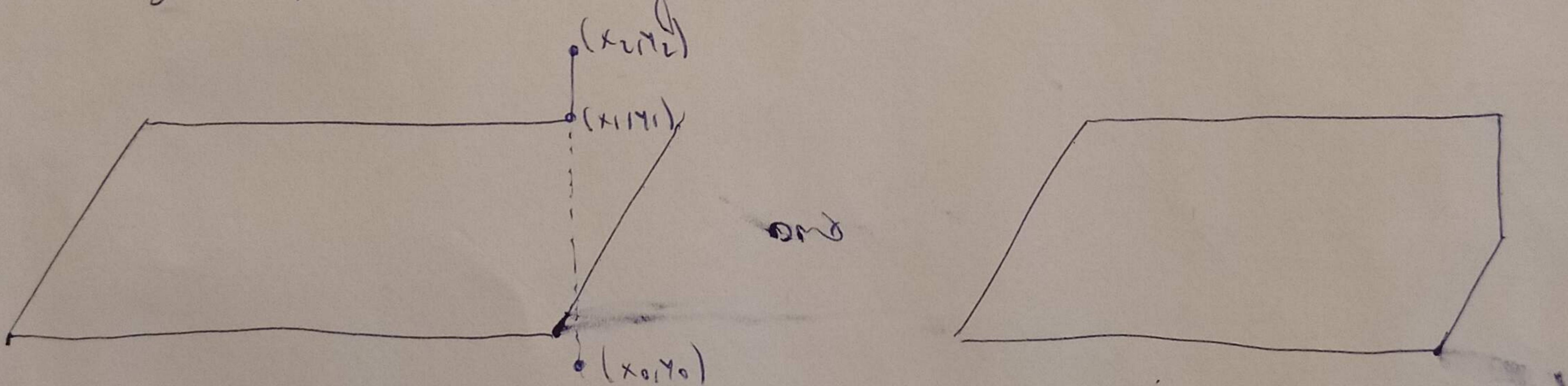
There is overlap. What is the boundary position?



This is a linear Algebra Problem!

$$\vec{a} \cdot \vec{n} = \text{left boundary}$$

Once I find it, how do I go on? That is how do I choose between:



Well, I choose the one that is completely outside my polygon

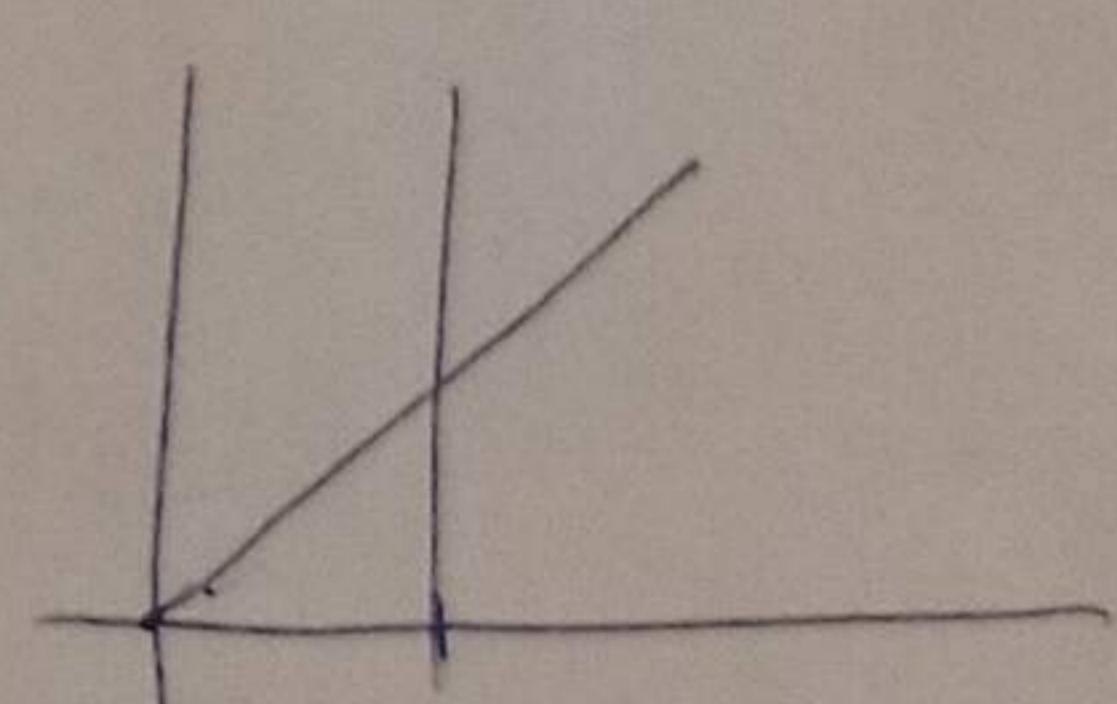
$$y = 1 \cdot x + 0$$

$$y_1 = ax_1 + b$$

I know a, b, c, d no

$$y_1 = bx_1 + d$$

$$y_1 = a \left(\frac{d-b}{a-c} \right) + b$$



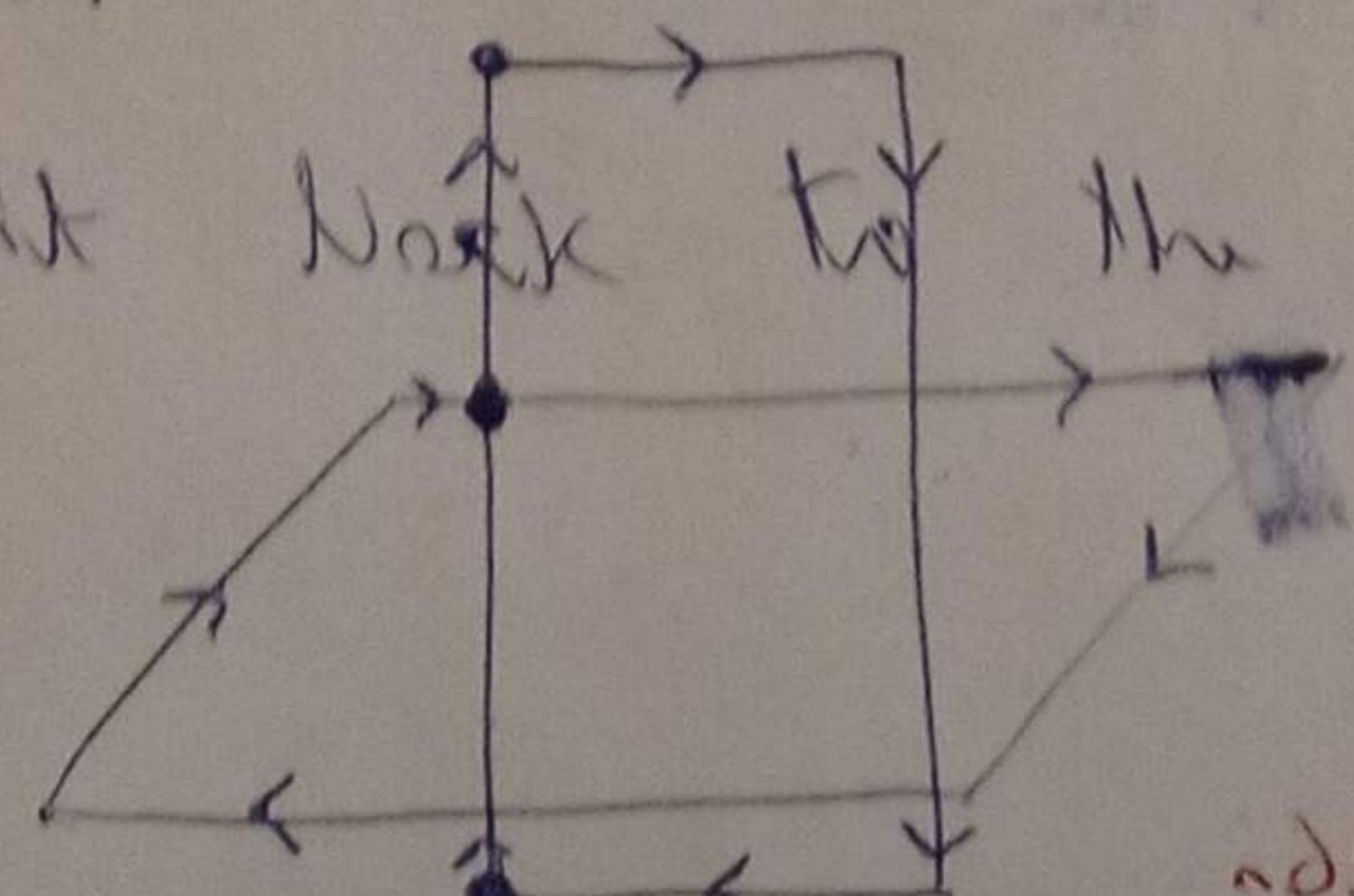
$$ax_1 + b = cx_1 + d$$

$$x_1 = \frac{d-b}{a-c}$$

(x_1, y_1) is the point of intersection!

The segment $(x_0, y_0) \rightarrow (x_1, y_1)$ intersects another off the polygon's segments no
 keep the other and move it to the queue

rule: polygon is defined
clockwise



rule: leftmost point IF POLYGONS EQUAL
 idea: polygons around from leftmost DOESN'T WORK

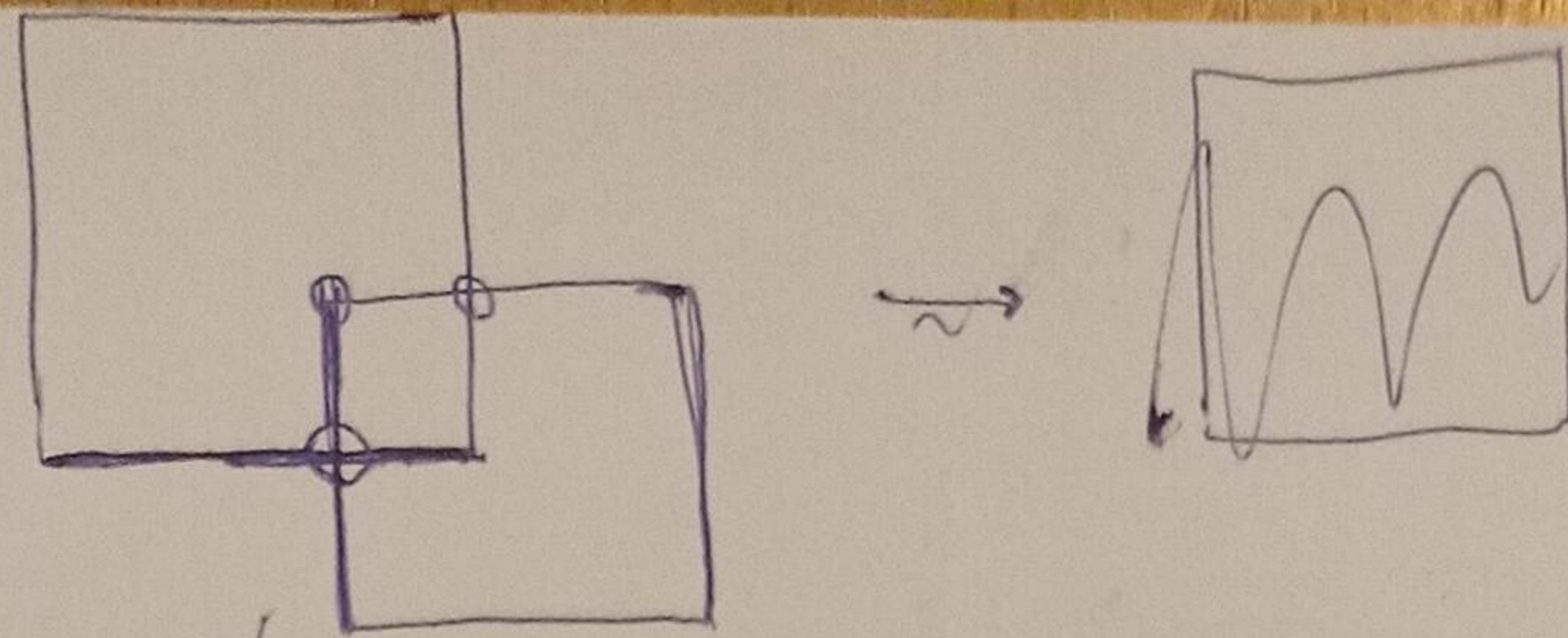
$$\gamma_0 = n\gamma_0 + N$$

$$\gamma_1 = n\gamma_1 + N$$

$$\gamma_0 - \gamma_1 = n(\gamma_0 - \gamma_1)$$

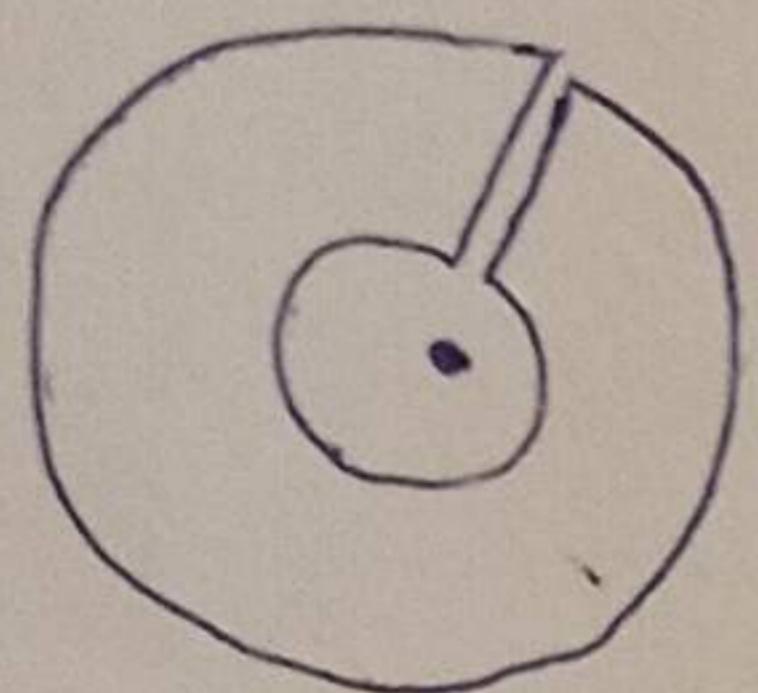
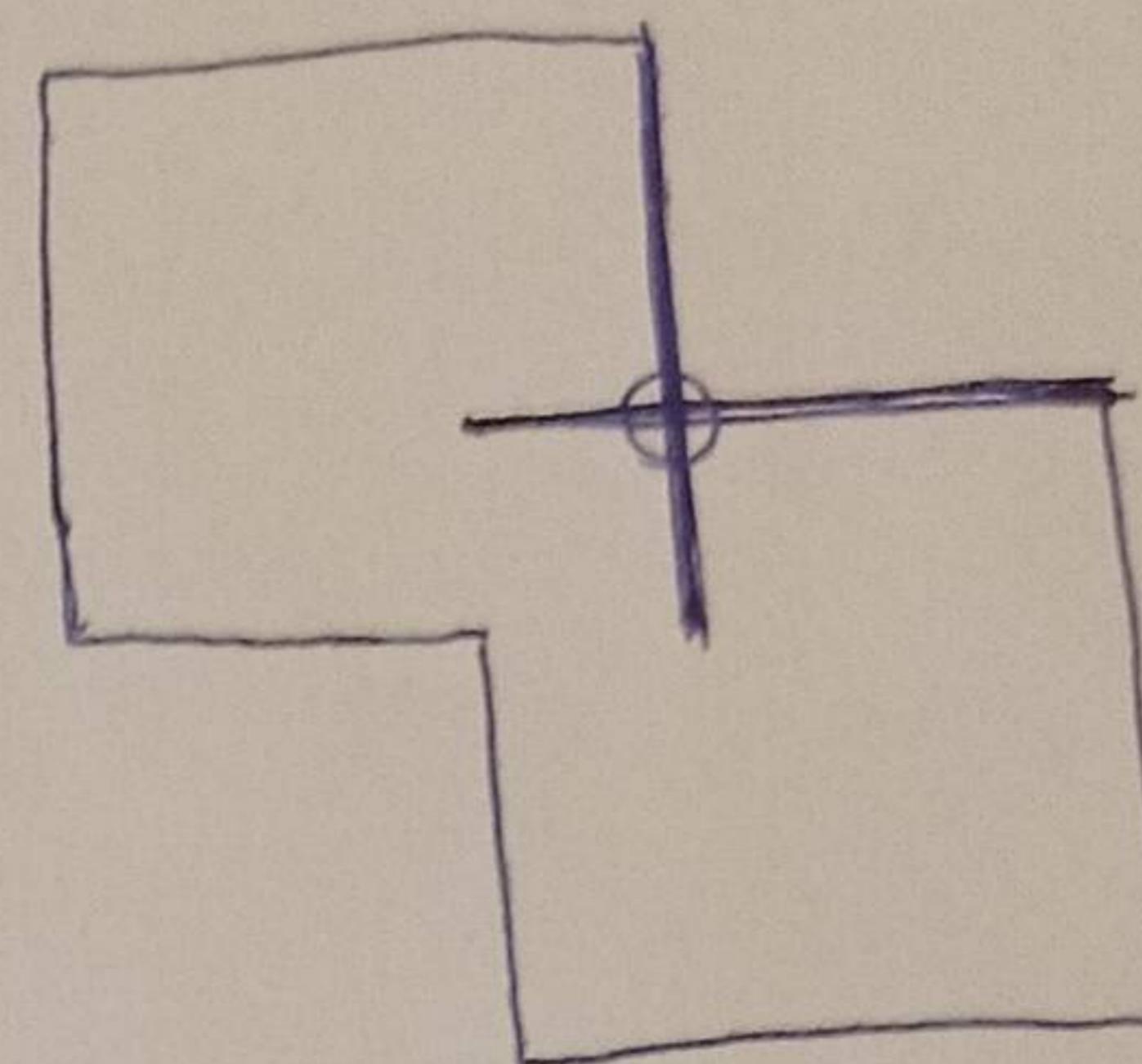
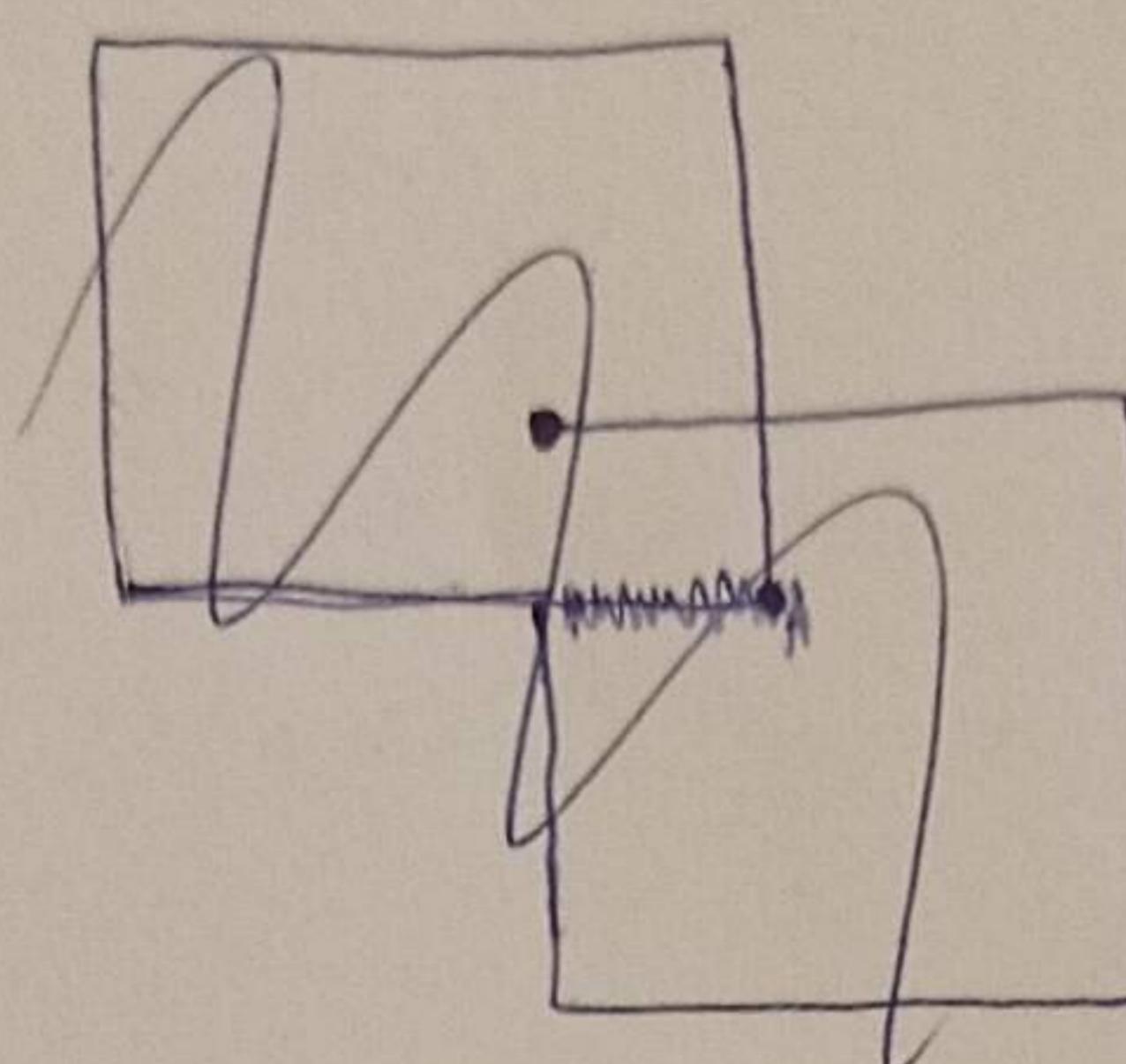
$$n = \frac{\gamma_0 - \gamma_1}{\gamma_0 - \gamma_1}$$

$$N = \gamma_0 - \left(\frac{\gamma_0 - \gamma_1}{\gamma_0 - \gamma_1} \right) \gamma_0$$



if $\gamma_0 = \gamma_1$, then $n = 0$

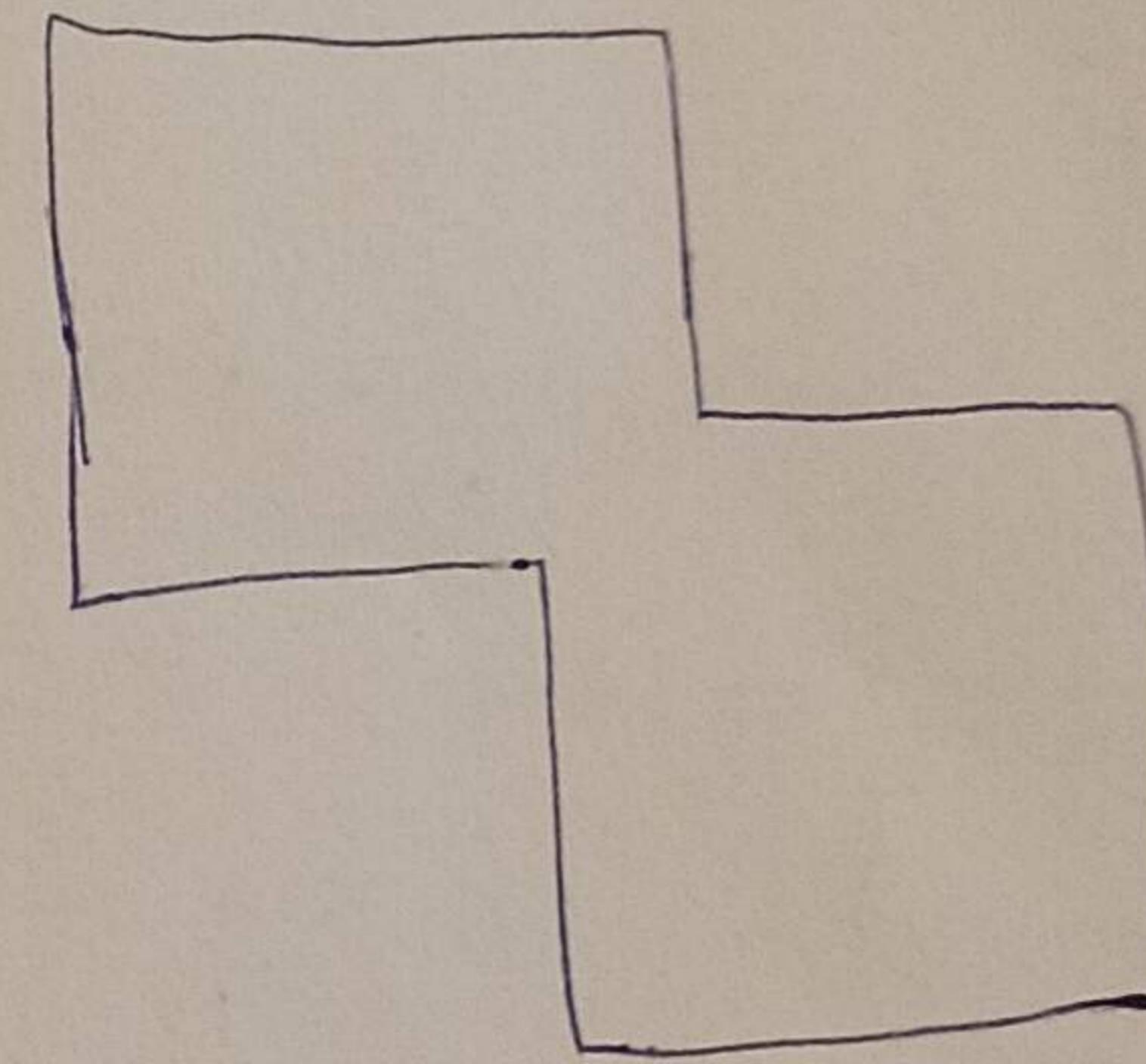
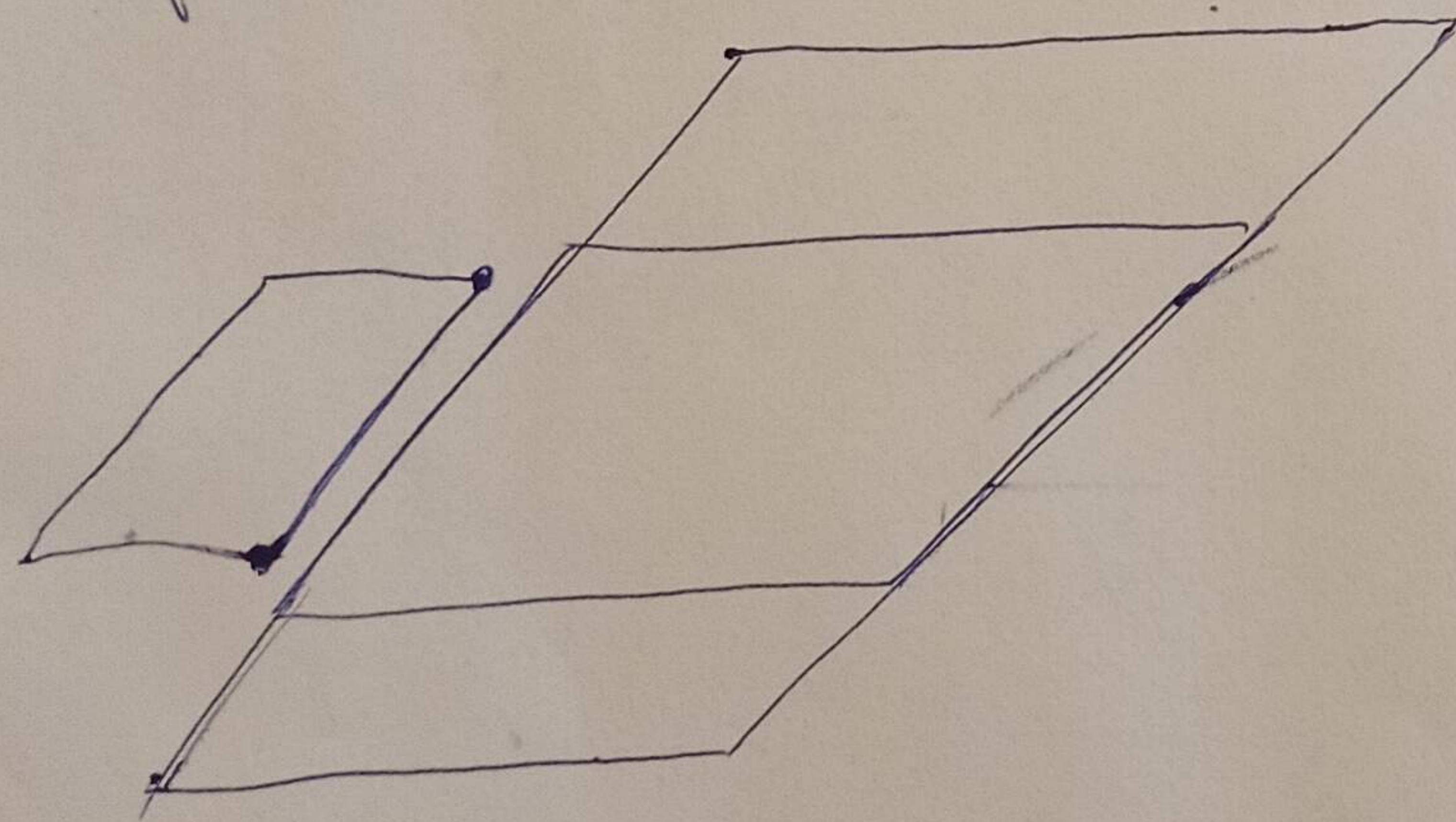
$N = \gamma_0$ which is correct for a straight line



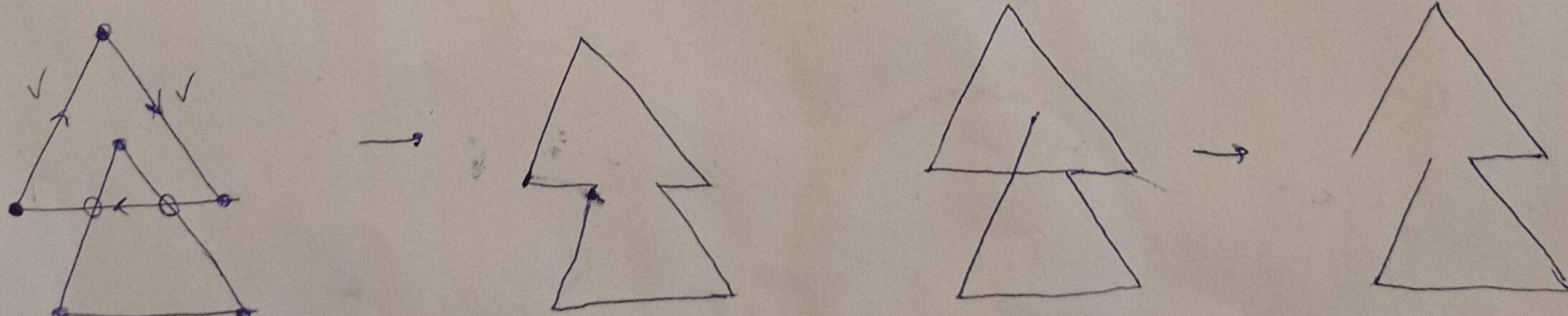
If inward segment is contained (or overlapping) just remove it

Otherwise it passes the other side, so keep only that

from outside

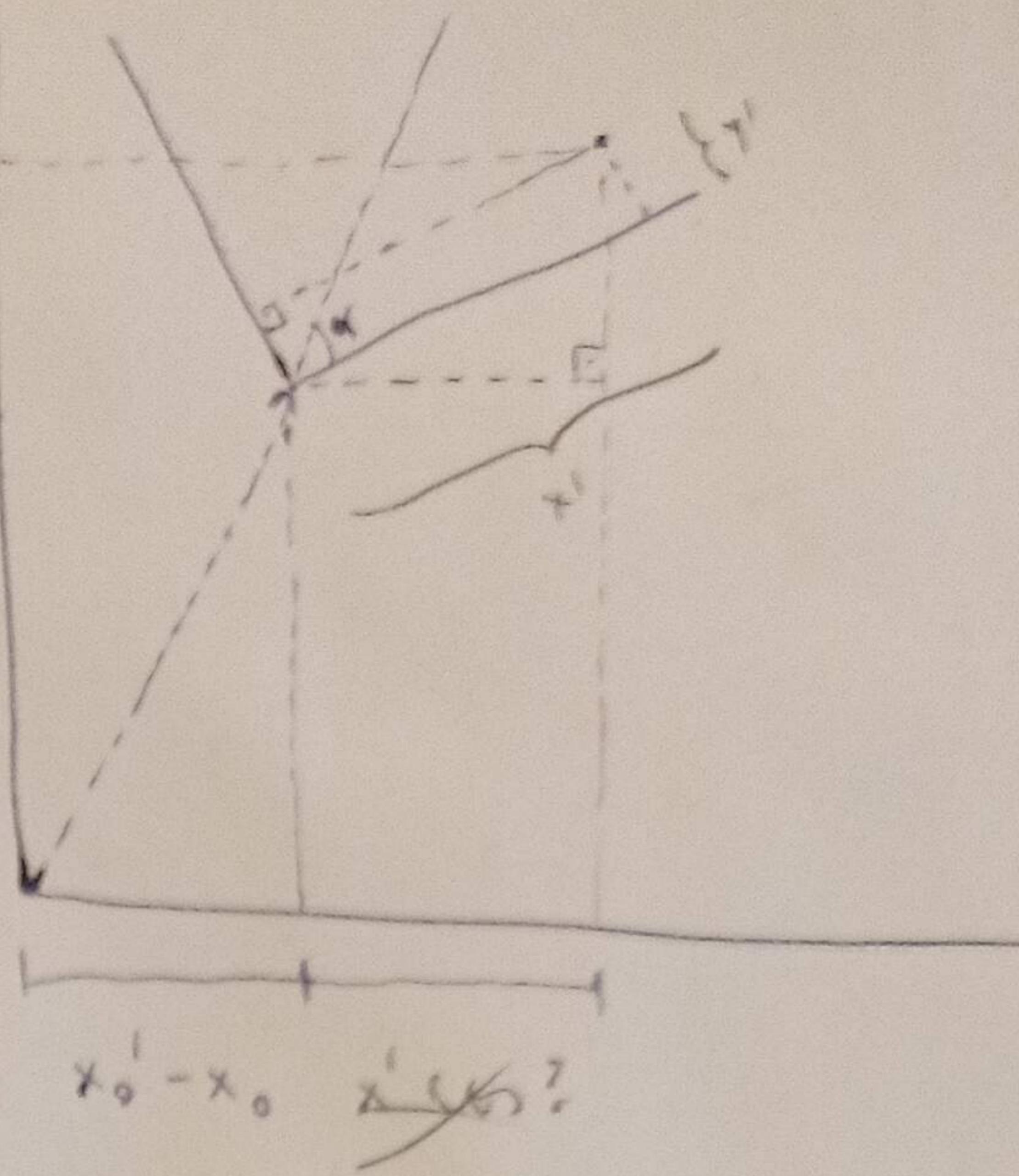
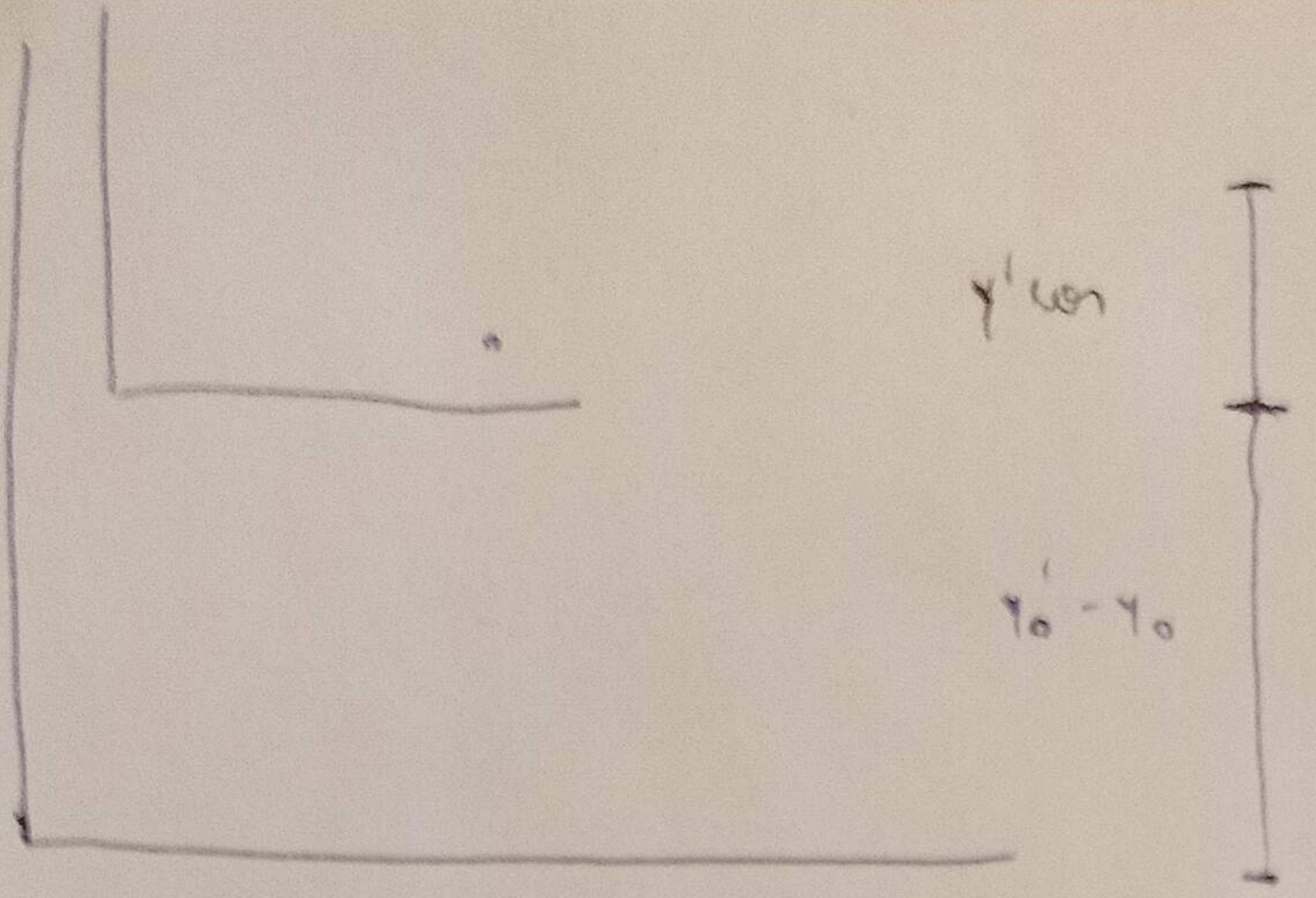


- share polygons better in having to own them into one



"anti-segments inside other polygons should be removed!"

here no seg. segment from right ??!



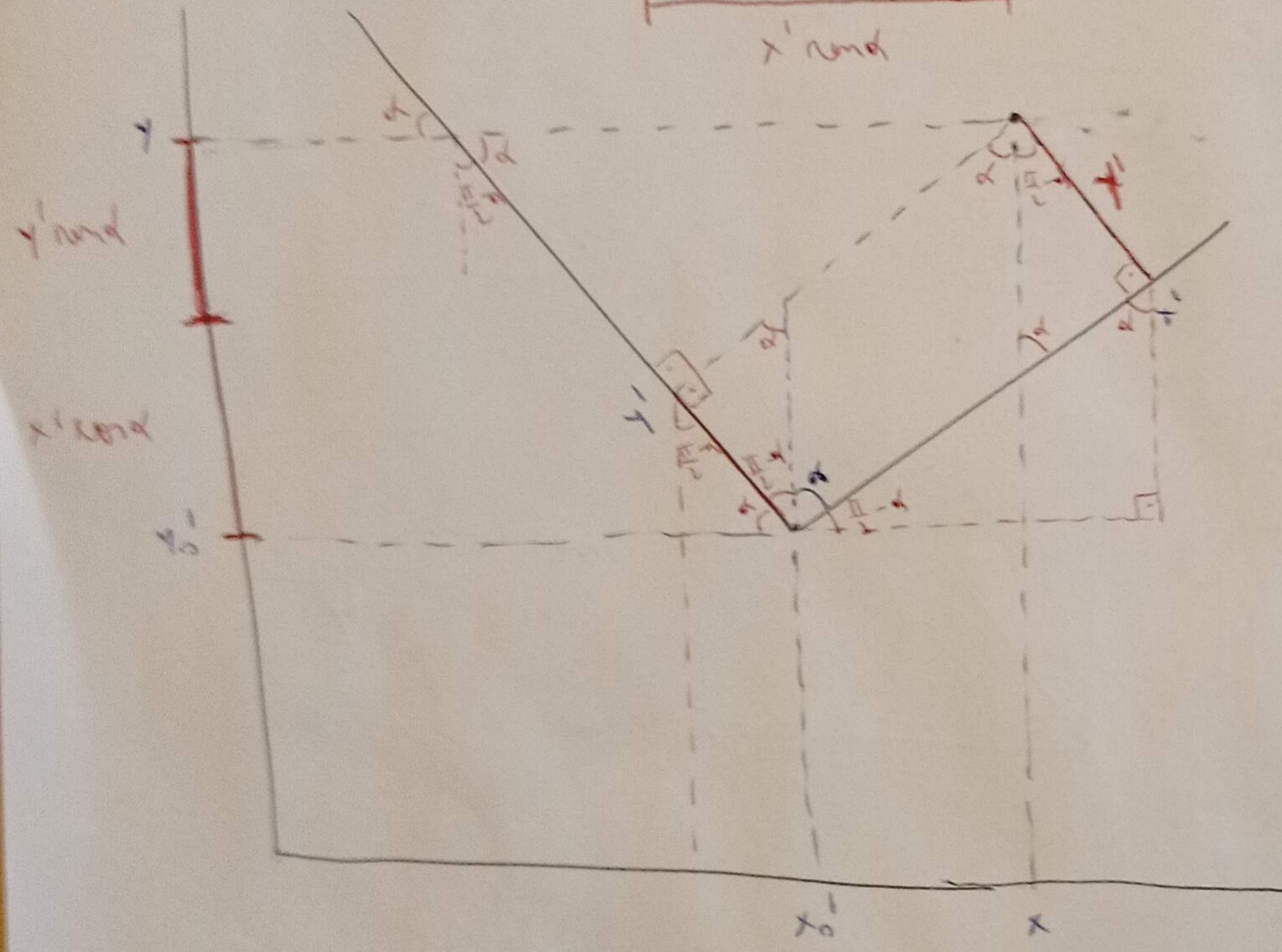
would make a translation

followed by a rotation

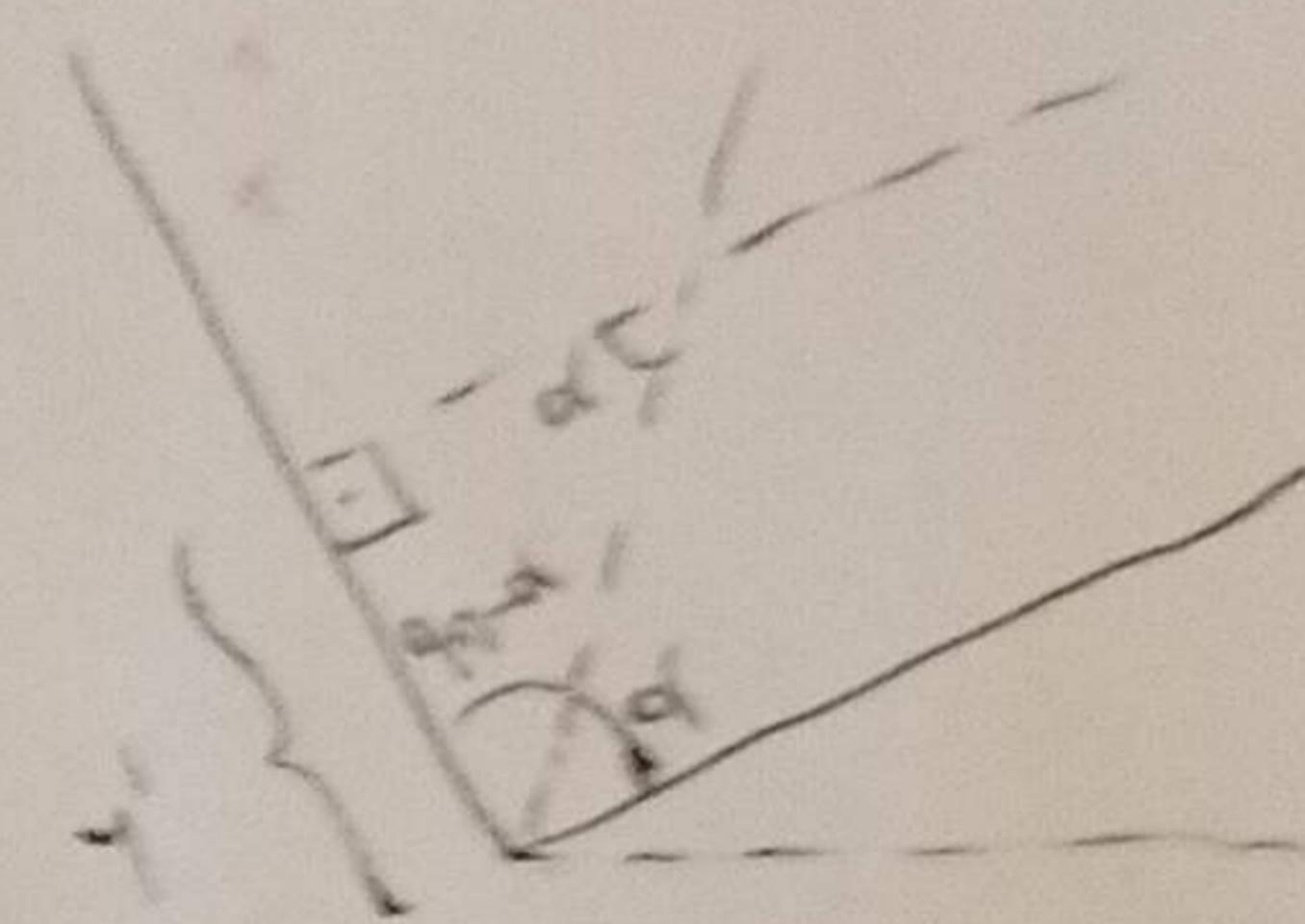
define rotation

$$y'_{wind}$$

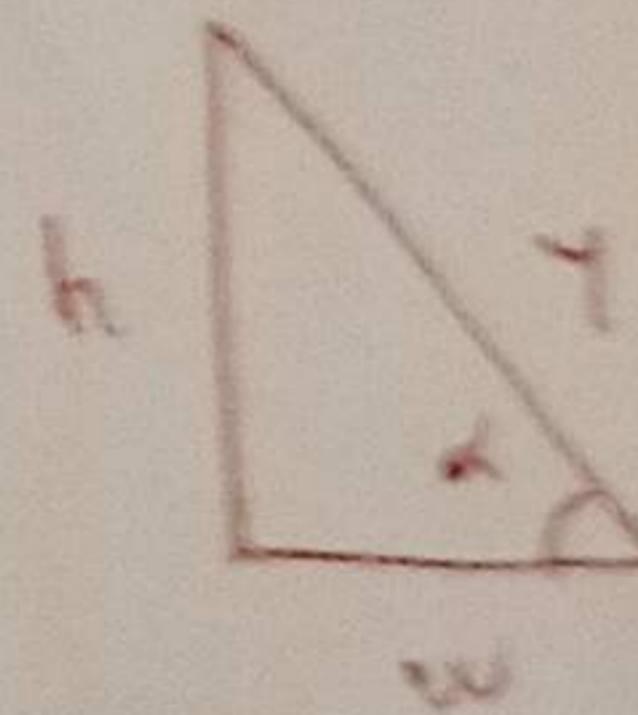
$$x'_{wind}$$



F



const =



$$\text{wind} = \frac{h}{\theta}$$

$$\text{wind} = \frac{h}{\gamma'}$$

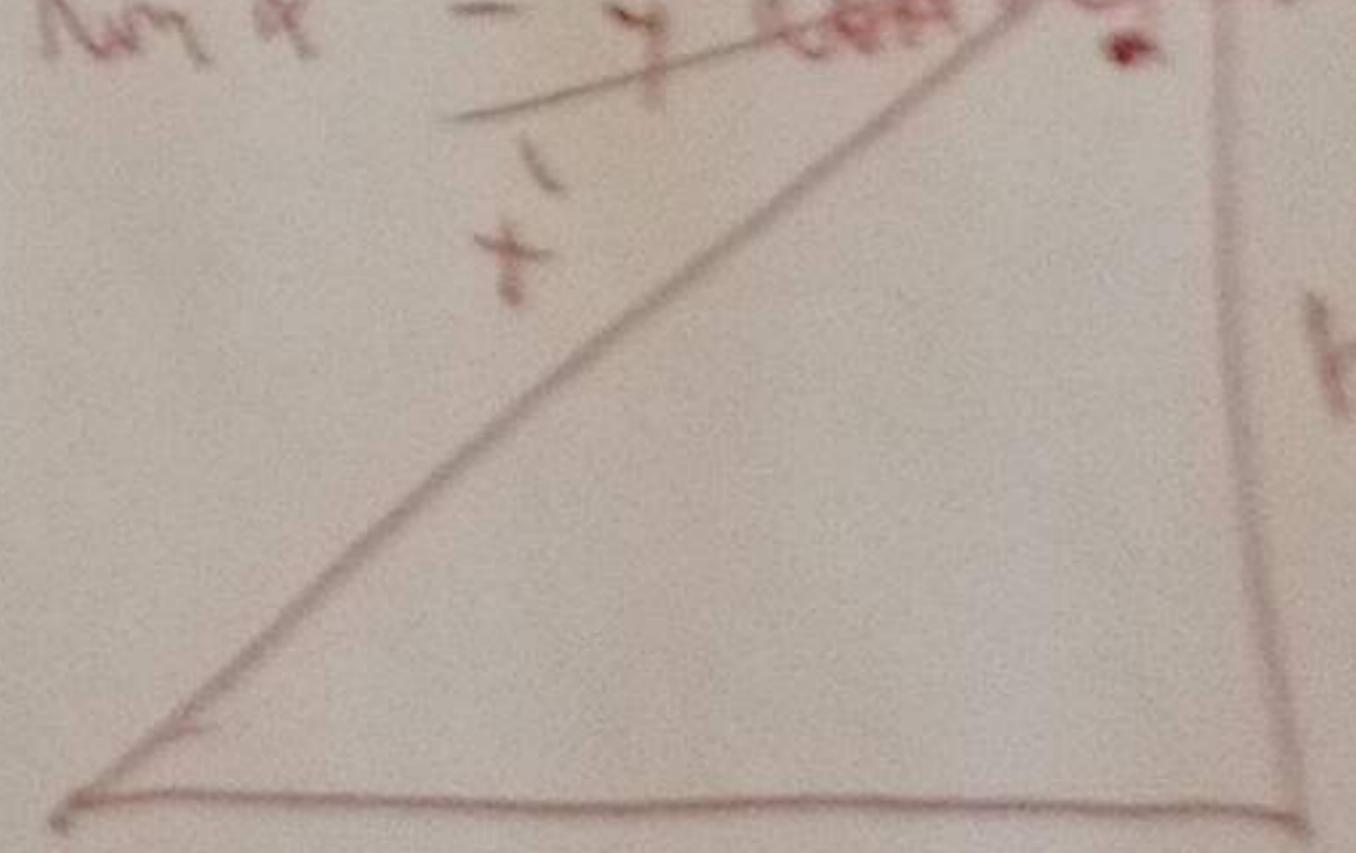
$$\text{const} = \frac{h}{x'}$$

$$y = y'_0 + x' \text{const} + y'_{wind}$$

$$y_{\text{total}} = y'_0 \text{wind} + x' \text{wind}^2 + y'_{\text{wind}}$$

$$x = x'_0 + x' \text{const} - y'_{wind}$$

$$x_{\text{total}} = x'_0 \text{wind} + x' \text{wind}^2 - y'_{\text{wind}}$$



$$y_{wind} + x_{wind} = y'_0 \text{wind} + x'_0 \text{wind} + x'$$

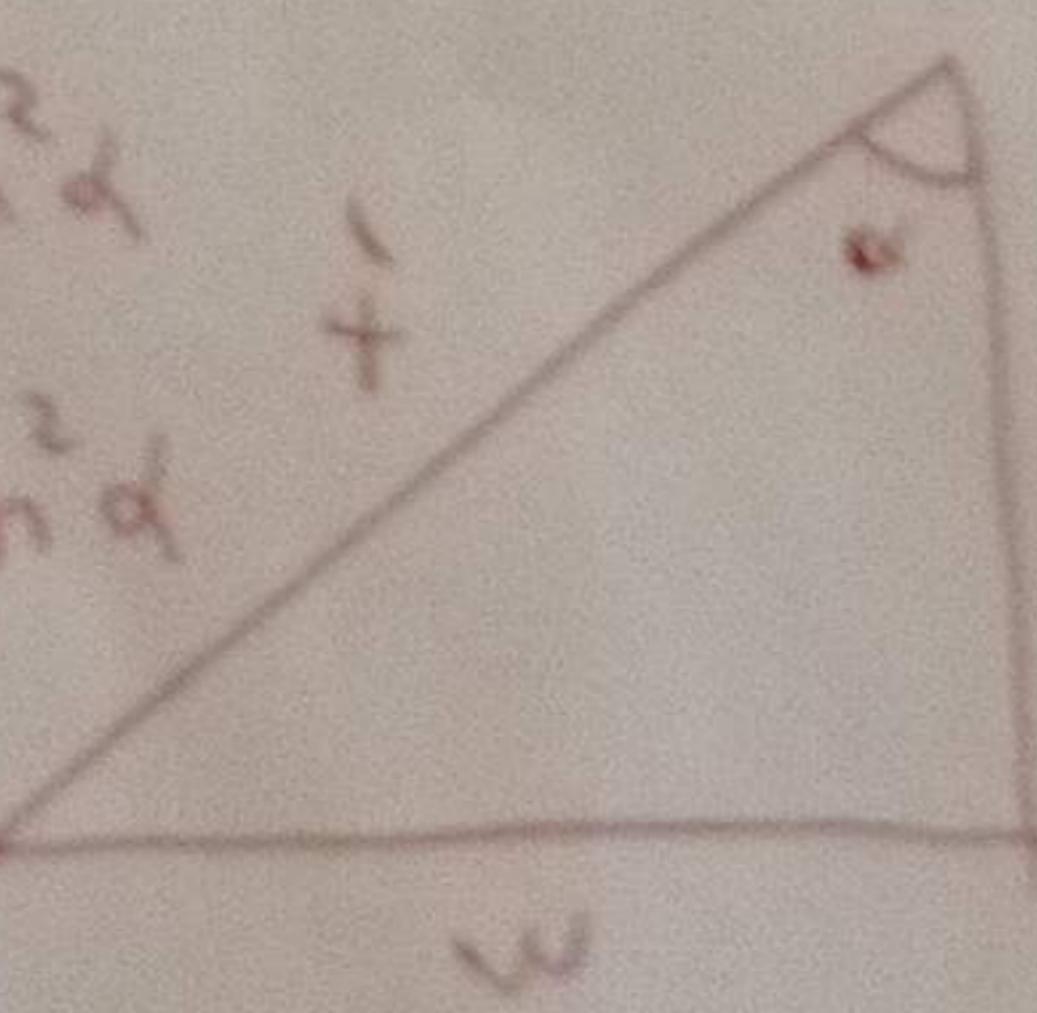
$$x' = (y - y'_0) \text{wind} + (x - x'_0) \text{wind}$$

$$y_{wind} = y'_0 \text{wind} + x' \text{wind}^2 + y'_{wind}$$

$$x_{wind} = x'_0 \text{wind} + x' \text{wind}^2 - y'_{wind}$$

$$y_{wind} - x_{wind} = y'_0 \text{wind} - x'_0 \text{wind} + y'$$

$$y' = (y - y'_0) \text{wind} - (x - x'_0) \text{wind}$$



$$\text{wind} = \frac{w}{\gamma'}$$



$$\text{const} = \frac{w}{x'}$$

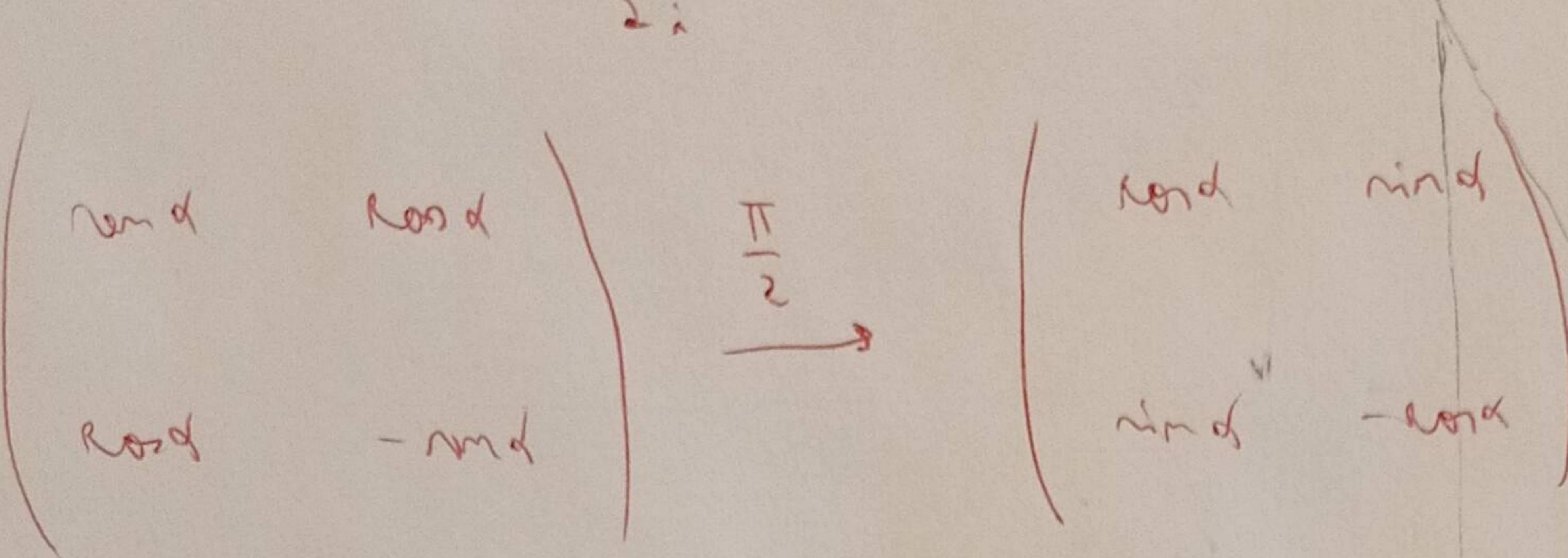
$$\begin{pmatrix} x' \\ y' \end{pmatrix} = \begin{pmatrix} \text{wind} & \text{wind} \\ \text{wind} & -\text{wind} \end{pmatrix} \begin{pmatrix} x - x'_0 \\ y - y'_0 \end{pmatrix}$$

$$\sin\left(\frac{\pi}{2} - \alpha\right) = \frac{e^{i(\frac{\pi}{2} + \alpha)} - e^{-i(\frac{\pi}{2} + \alpha)}}{2i} =$$

NAIT

$$\frac{ie^{i\alpha} + ie^{-i\alpha}}{2i} = \cos\alpha$$

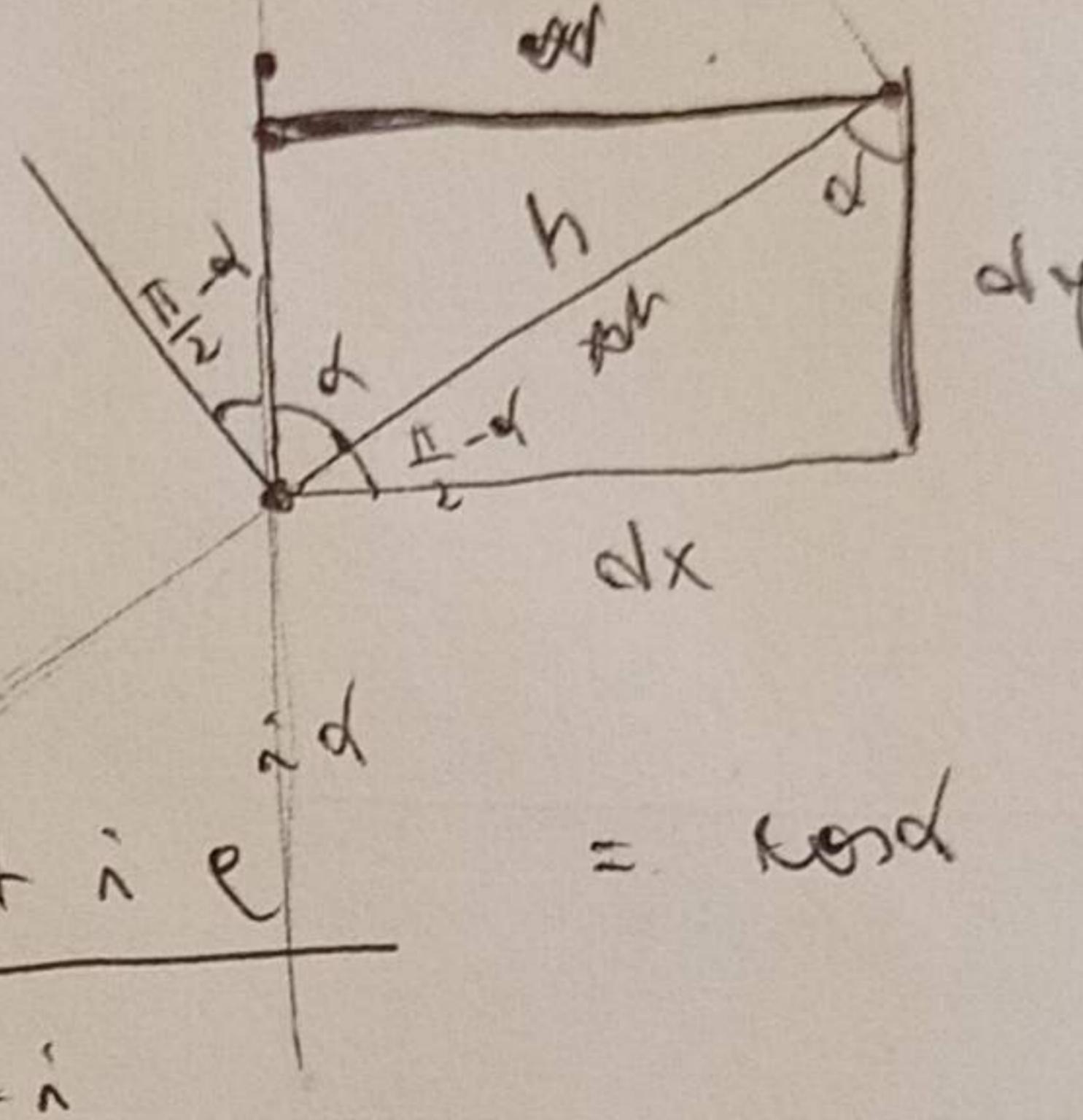
$$e^{i\frac{\pi}{2}} = \cos\frac{\pi}{2} + i\sin\frac{\pi}{2} = i$$



How do I find α ? Not present

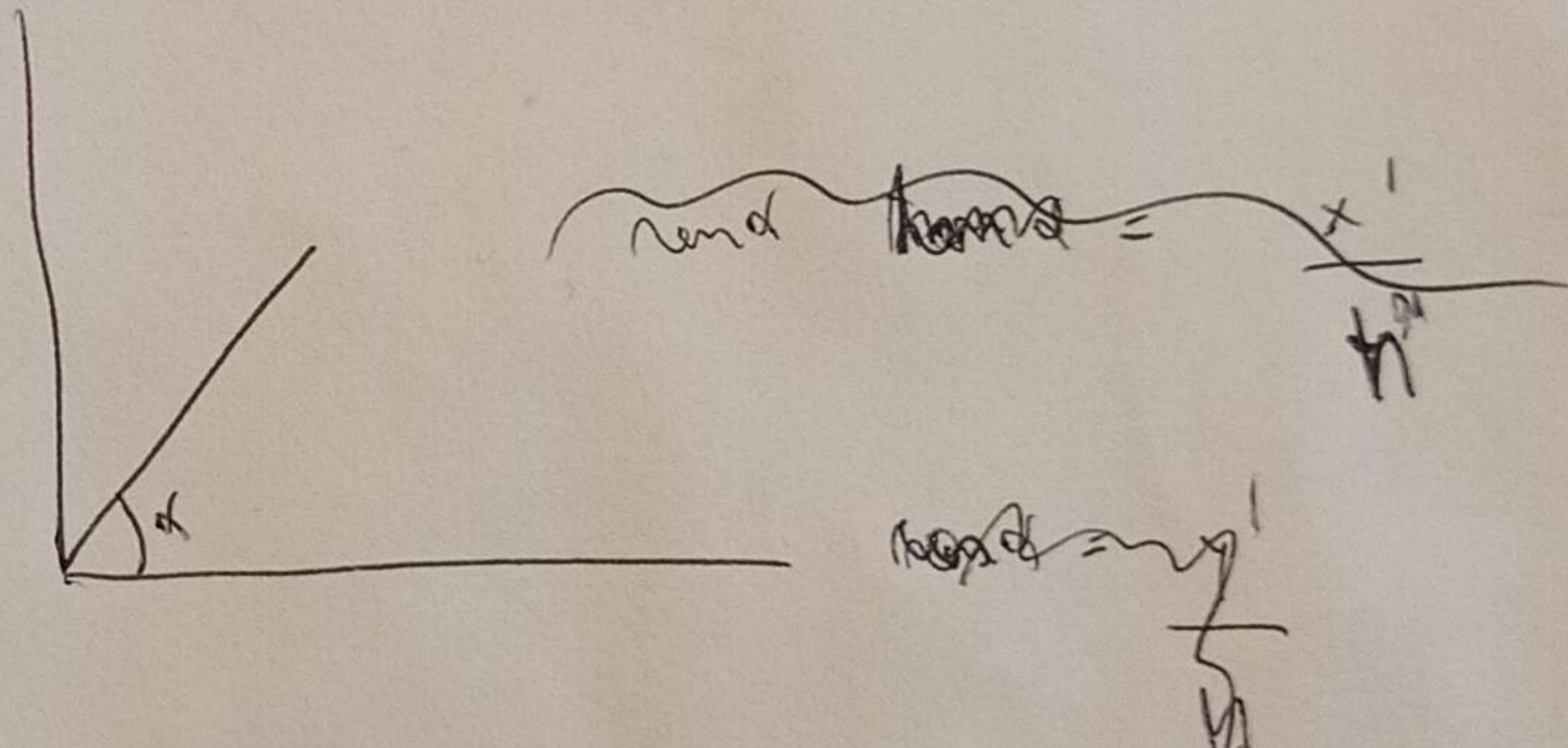
$$\sin\left(\frac{\pi}{2} - \alpha\right) = \frac{e^{i(\frac{\pi}{2} - \alpha)} - e^{-i(\frac{\pi}{2} - \alpha)}}{2i} = \frac{-i\alpha + i\alpha}{2i} = \cos\alpha$$

$$\cos\left(\frac{\pi}{2} - \alpha\right) = \frac{e^{i(\frac{\pi}{2} - \alpha)} + e^{-i(\frac{\pi}{2} - \alpha)}}{2} = \frac{i\alpha - i\alpha}{2} = \frac{-i\alpha + i\alpha}{2} = \cos\alpha$$



$$|\vec{a} \cdot \vec{b}| = |\alpha| |\beta| \cos\alpha$$

$$\cos\alpha = \frac{|\alpha| |\beta|}{|\vec{a} \cdot \vec{b}|}$$



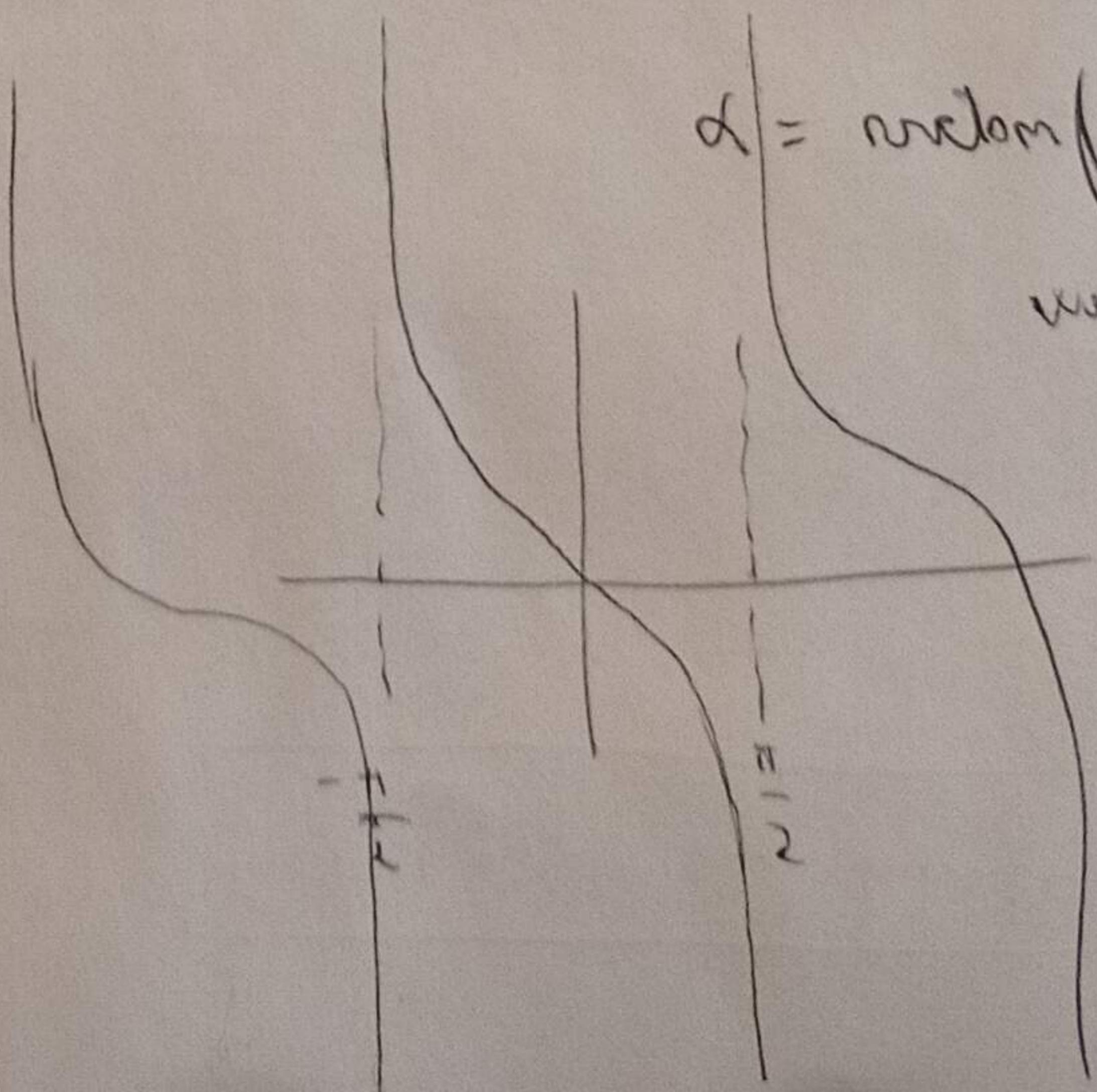
$$\text{real} = \frac{dx}{h}$$

$$\text{real} = \frac{dy}{h}$$


$$\tan\alpha = \frac{dx}{dy}$$

$$\alpha = \arctan\left(\frac{dx}{dy}\right)$$

well defined for $\frac{dx}{dy} \neq 0$



You're given a list of segments, how to order them into a polygon, efficiently?
 Intuitively: Start with a segment, iterate over all other segments to find the
 joint one that whose end equals this segment's start

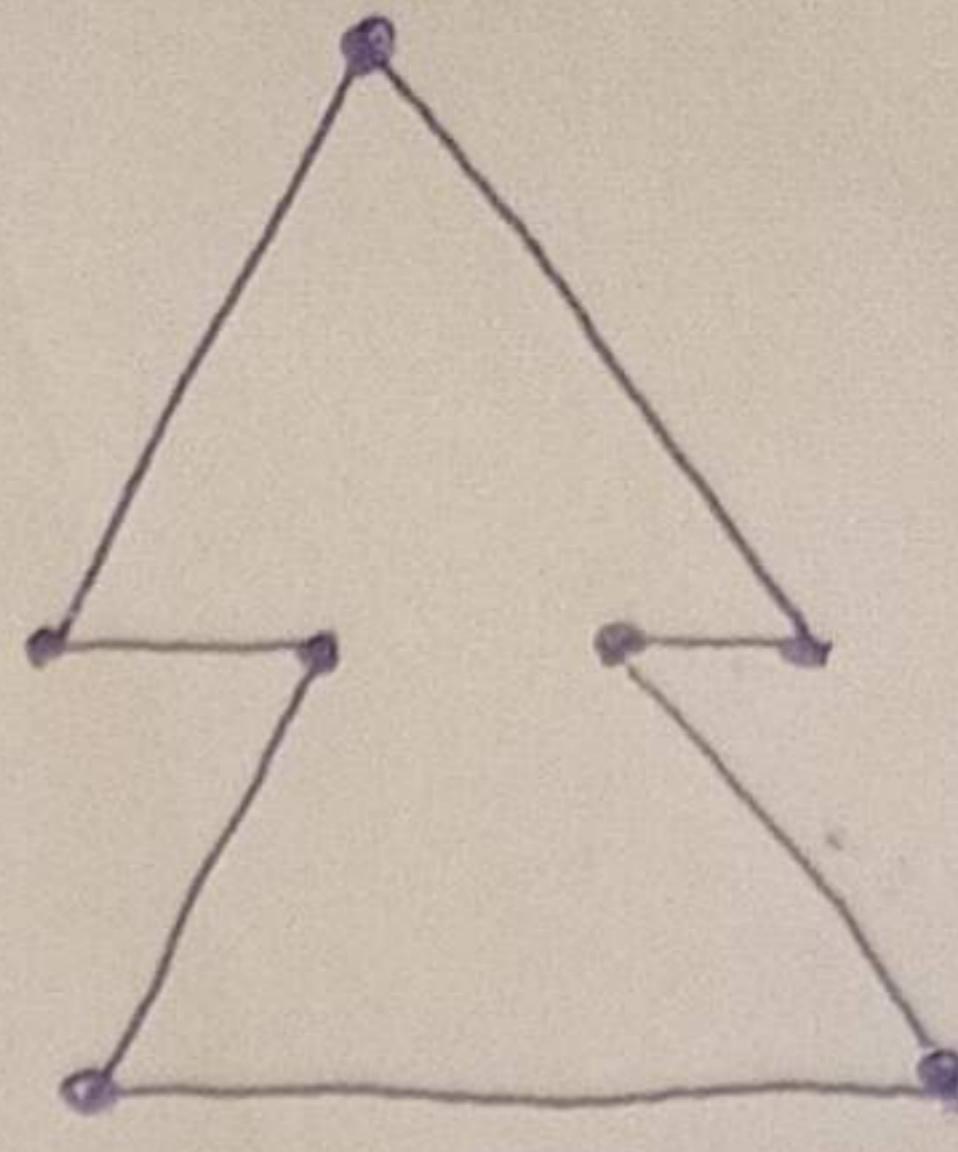
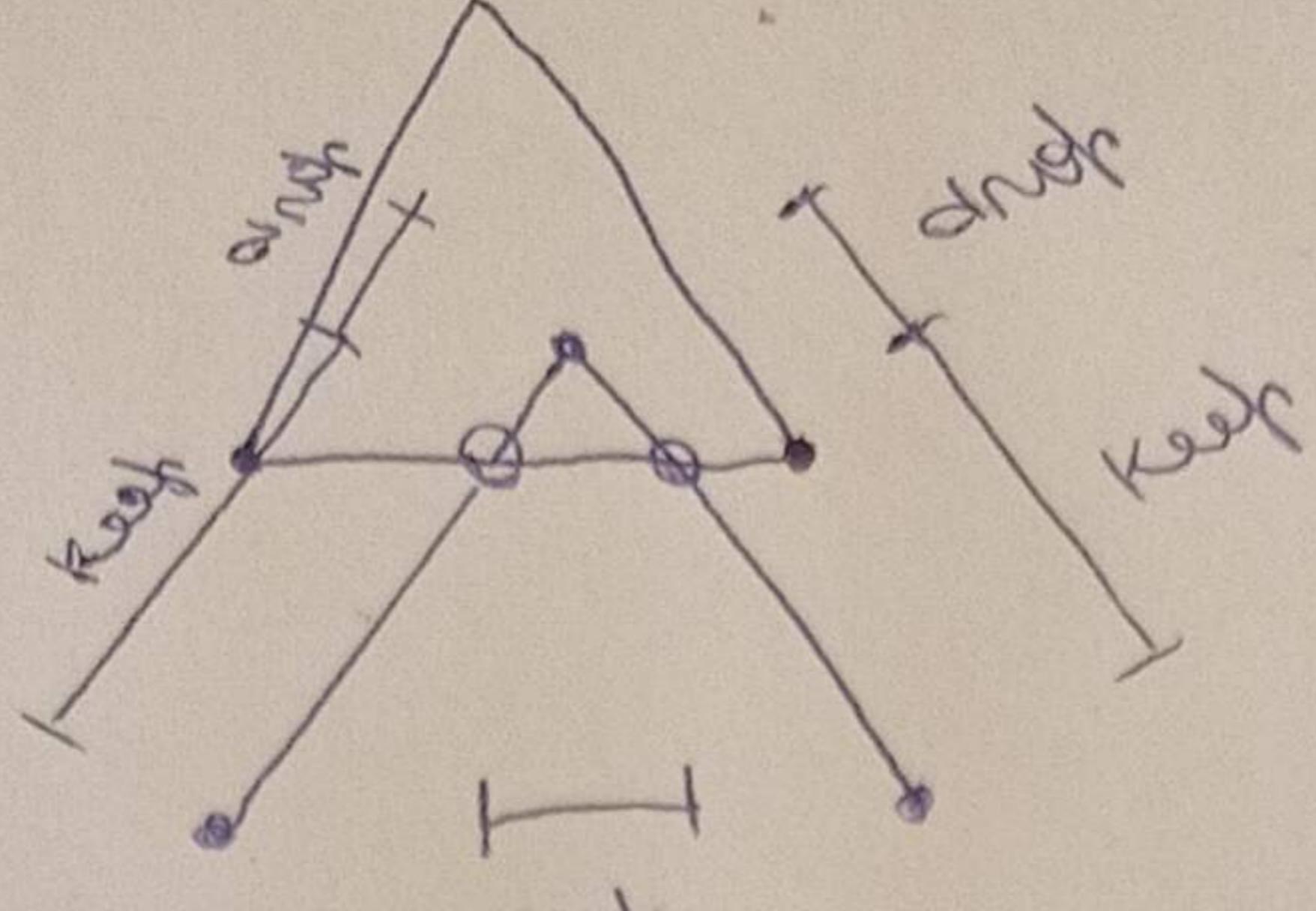
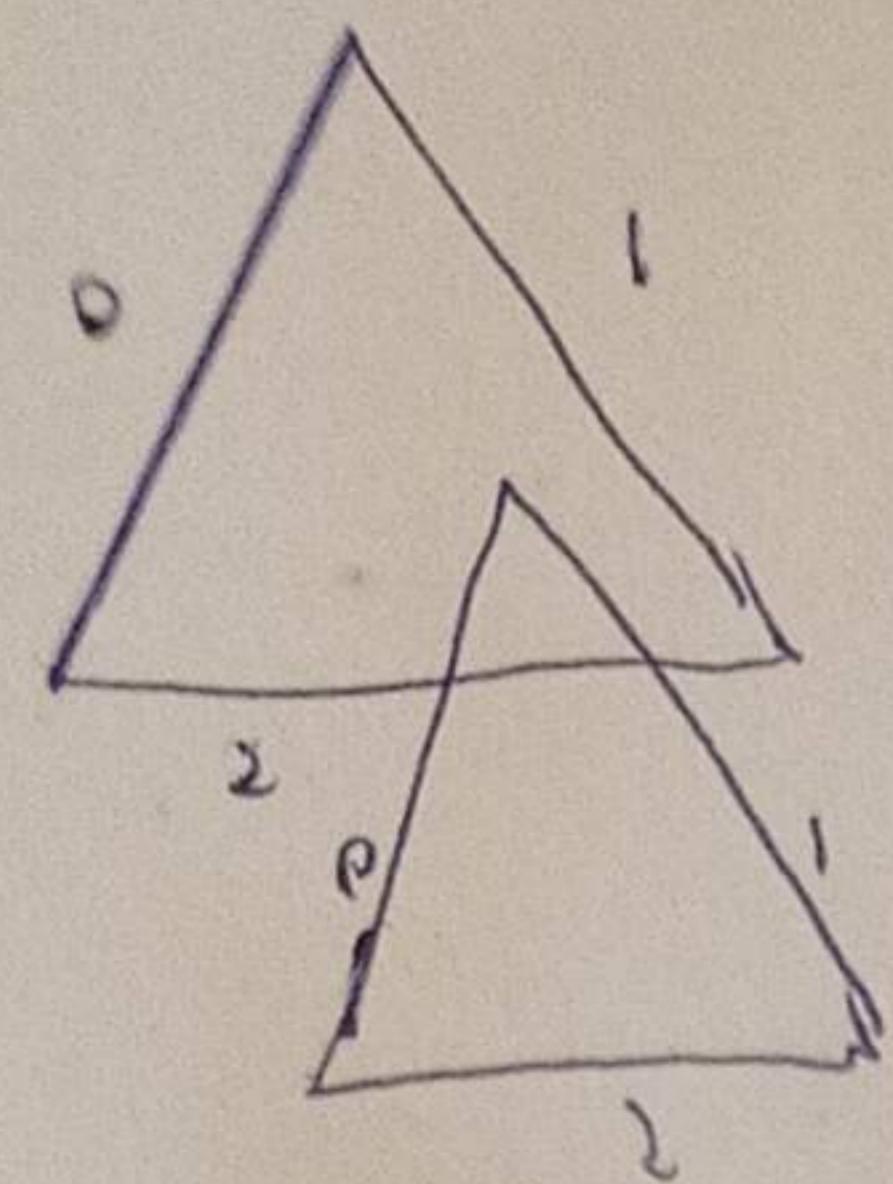


TODO: deal with disjoint polygons

TODO: fix point intersections no fix vs needs

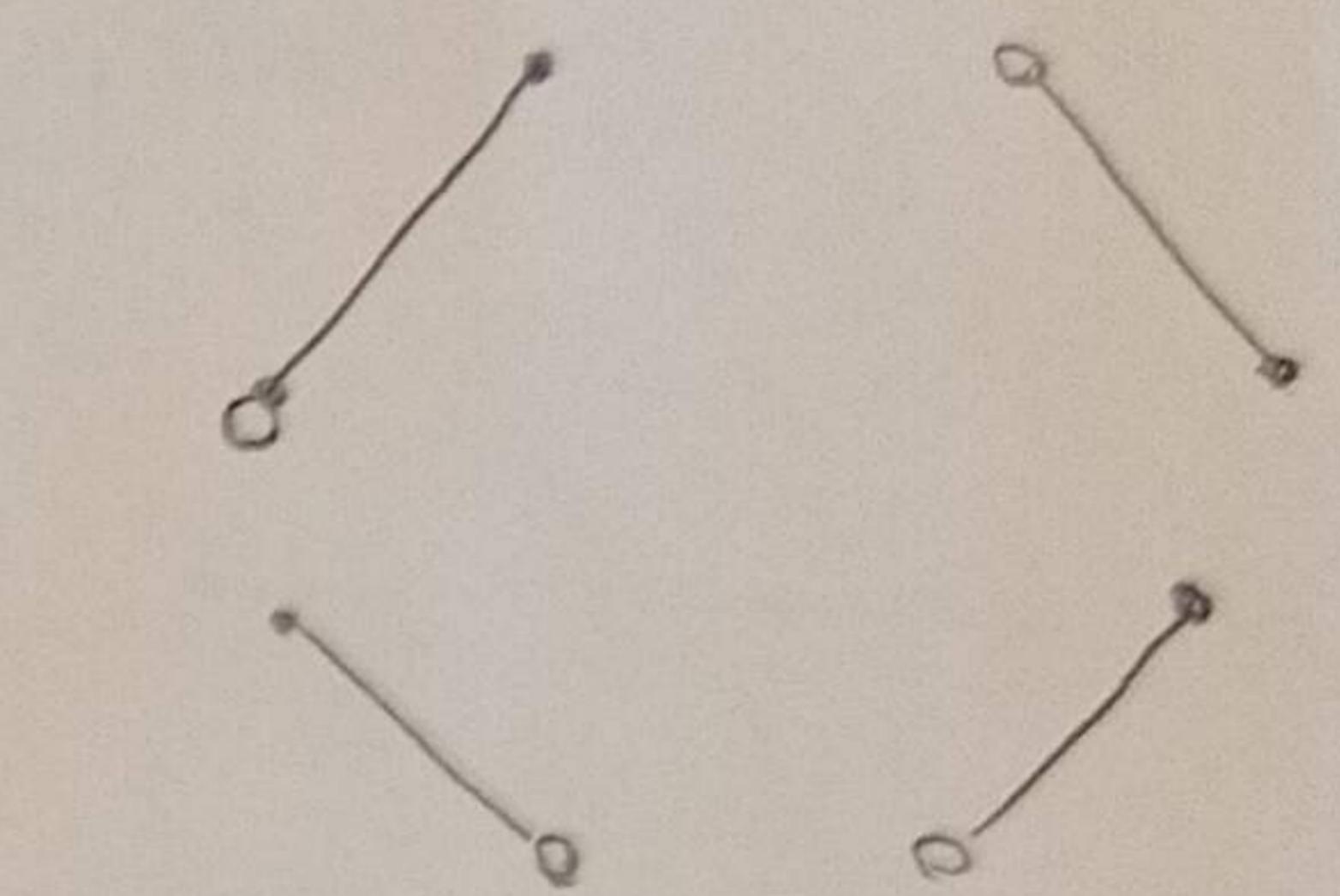
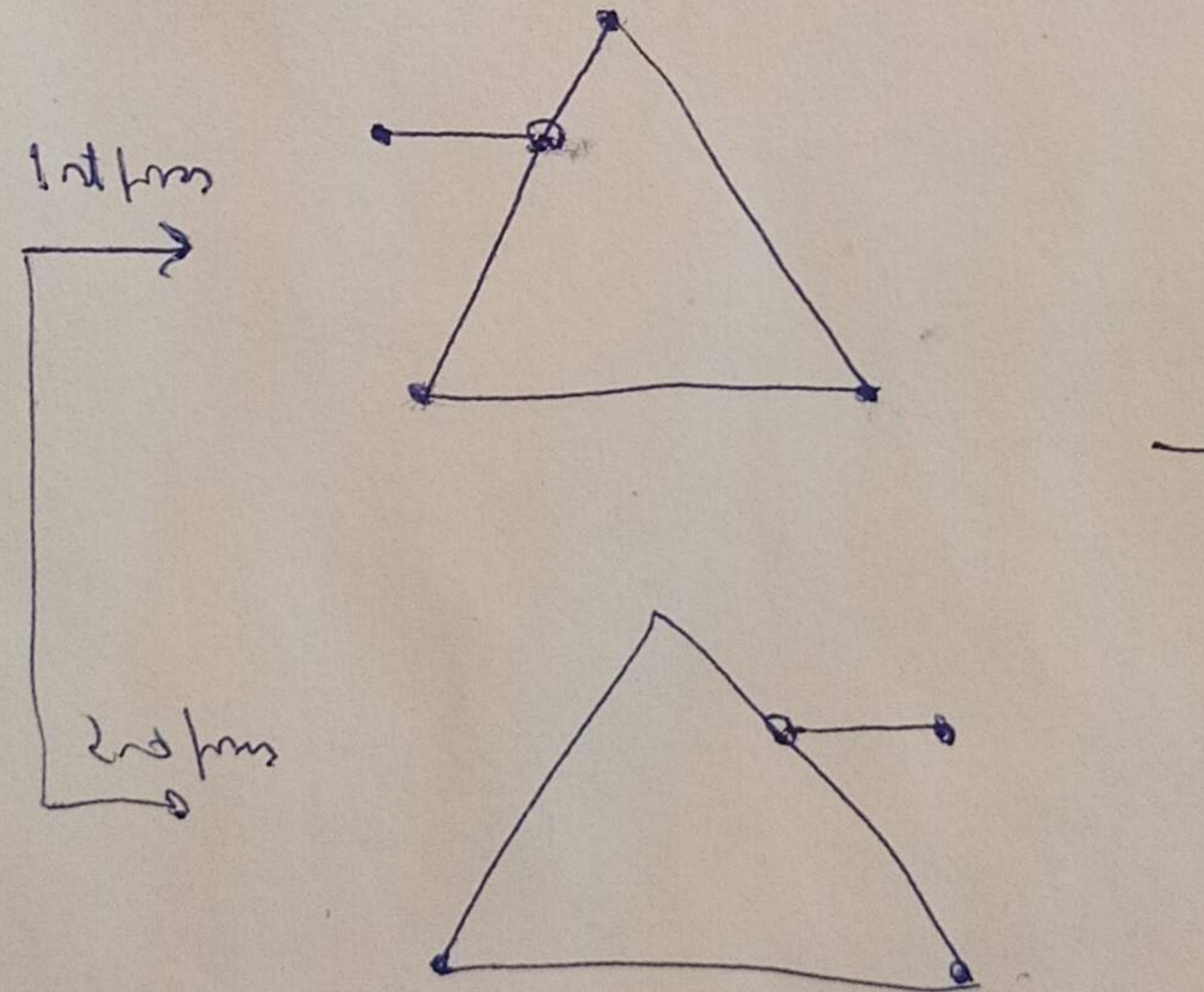
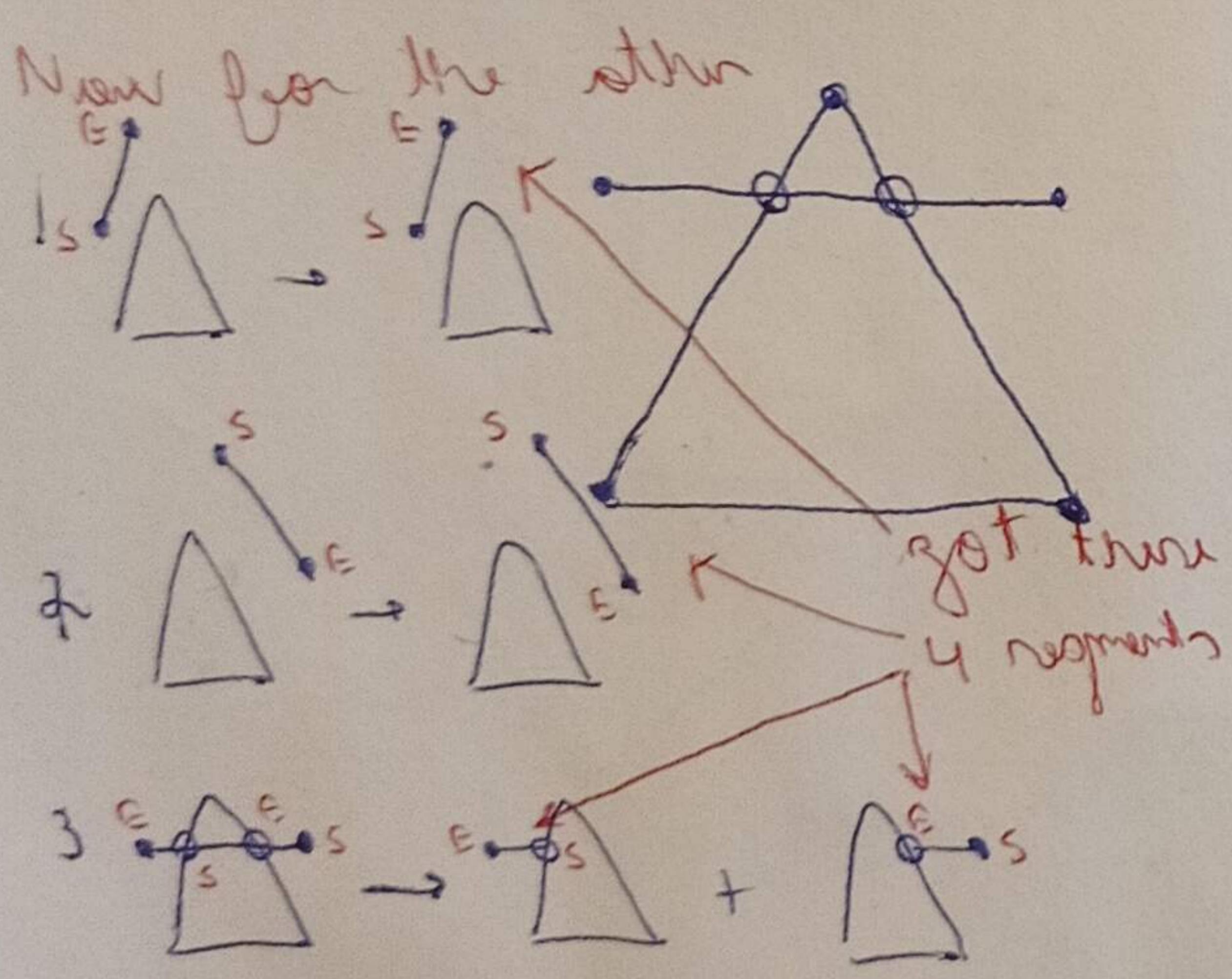
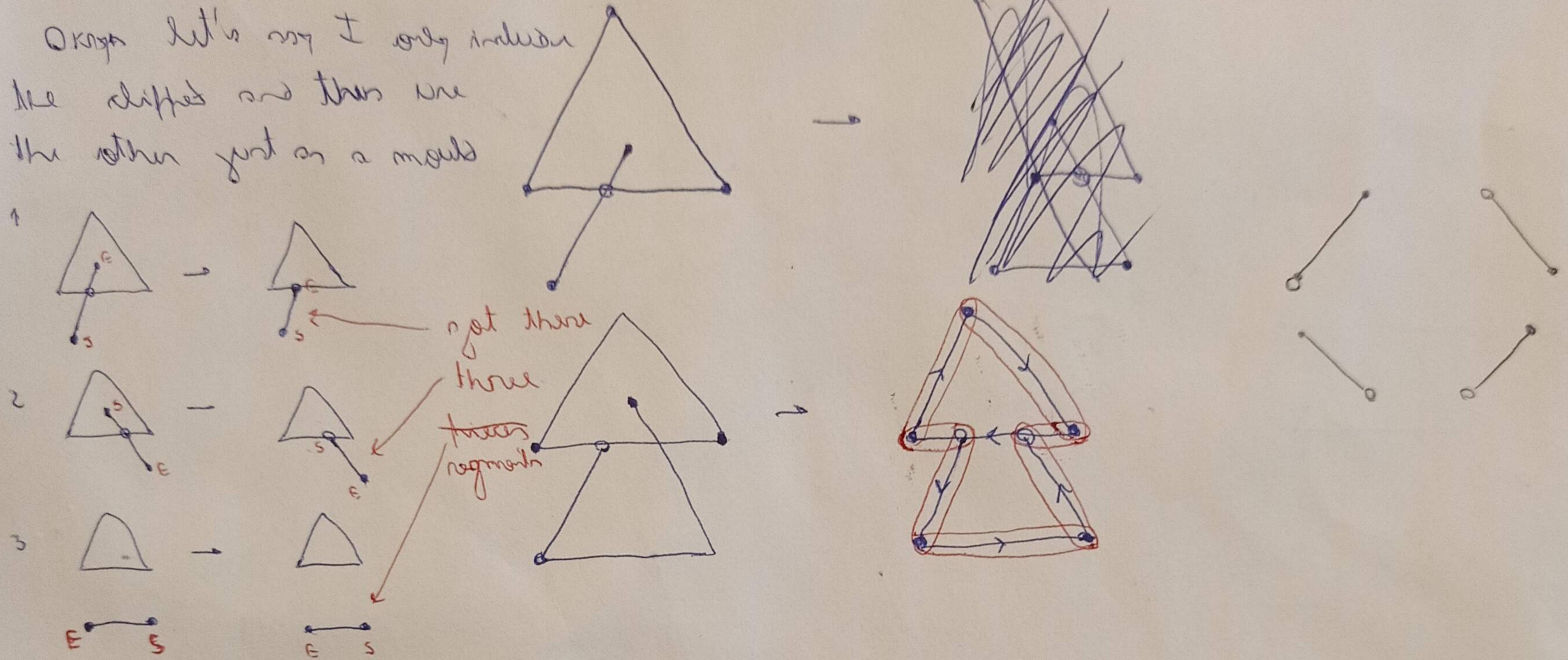
TODO: handle many polygons

counter clockwise!



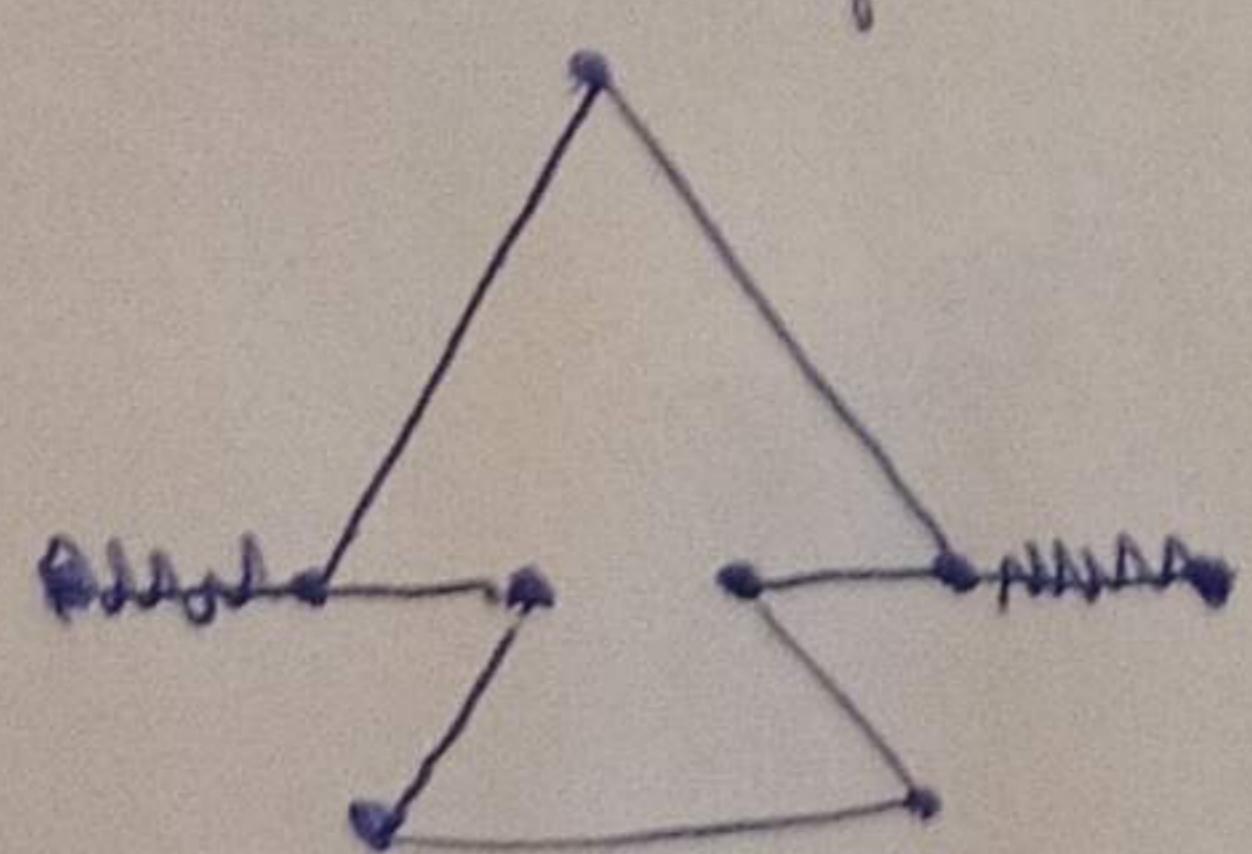
I know how to produce the correct segments. Now I do not know how to join them after the first

Okay let's say I only include the slanted and then use the other just as a mould



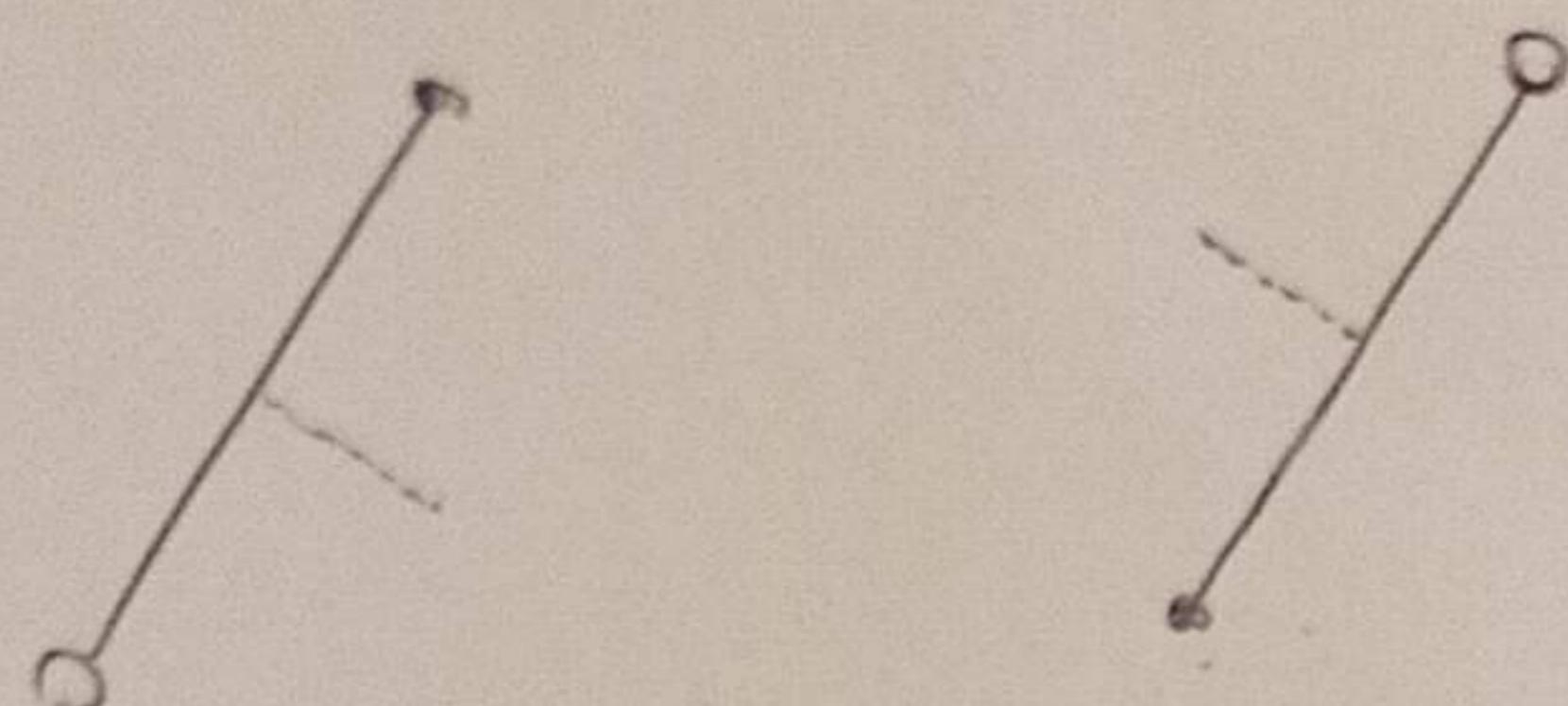
Putting it all together:

and if we start off the other! To determine order...



It works! And is highly malleable!

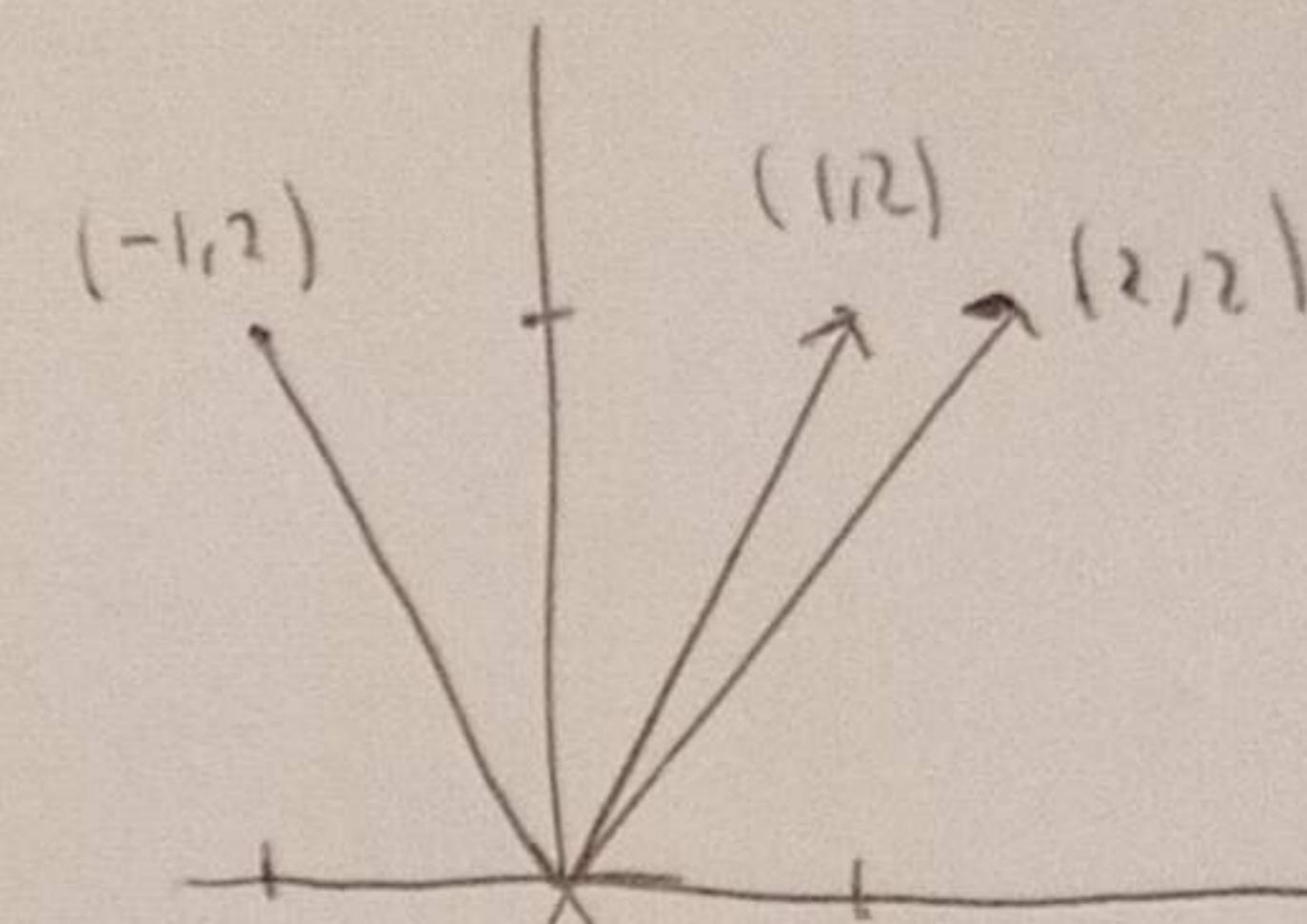
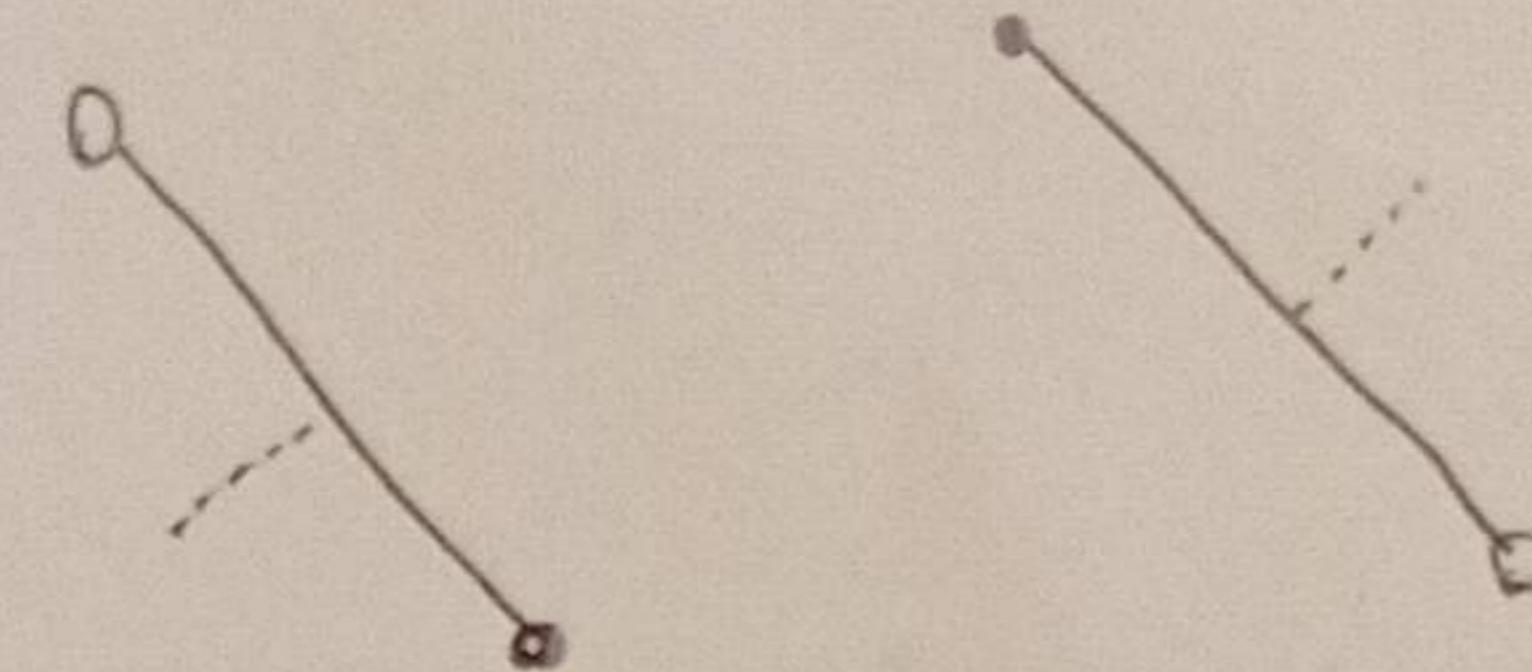
Normal vector construction



$$2 \cdot 1 + 2 \cdot 1$$

$$2 \cdot 1 + 2 \cdot -1$$

$$\sqrt{-1 \cdot 2 + 2 \cdot 2} > 0$$

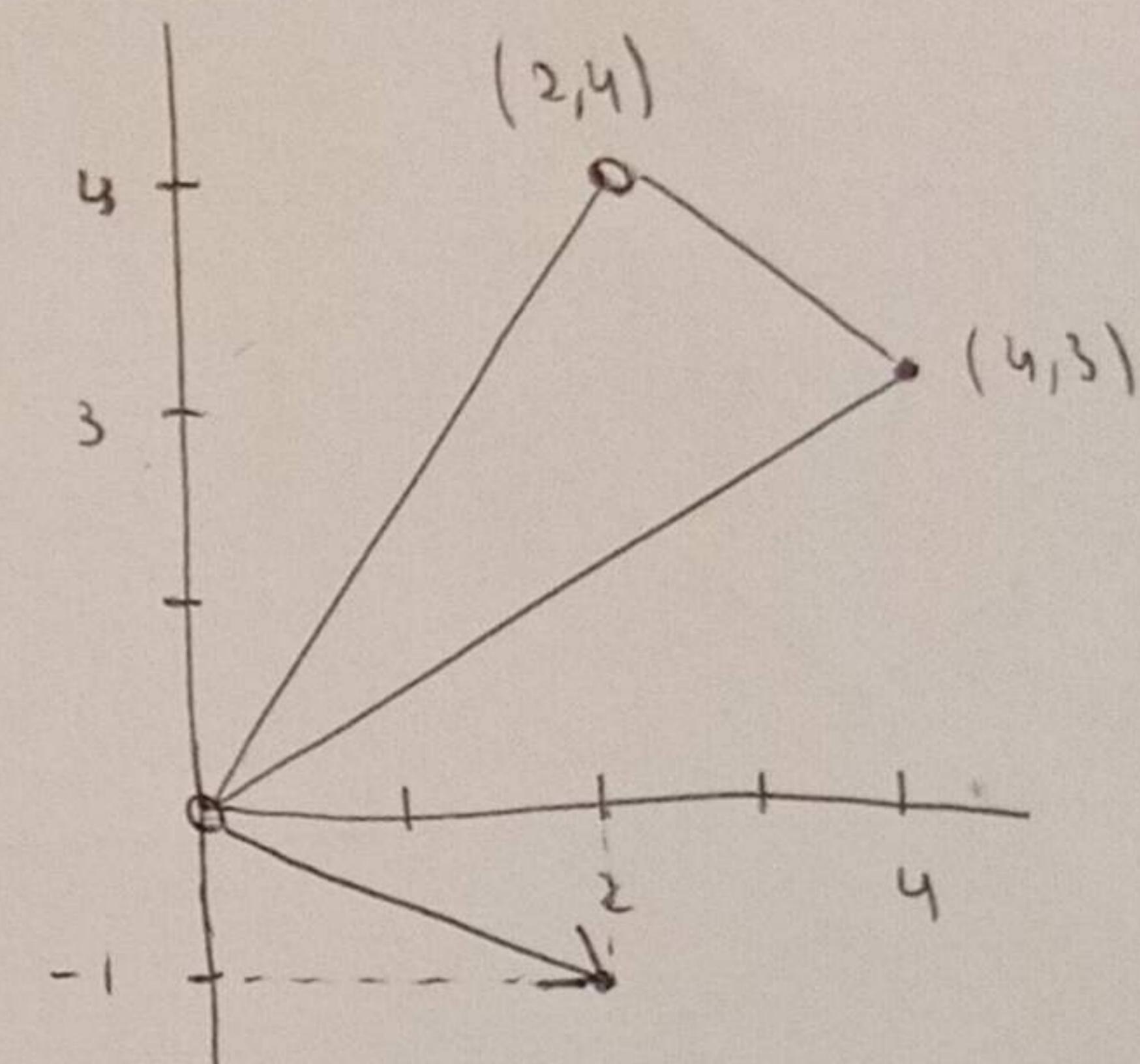
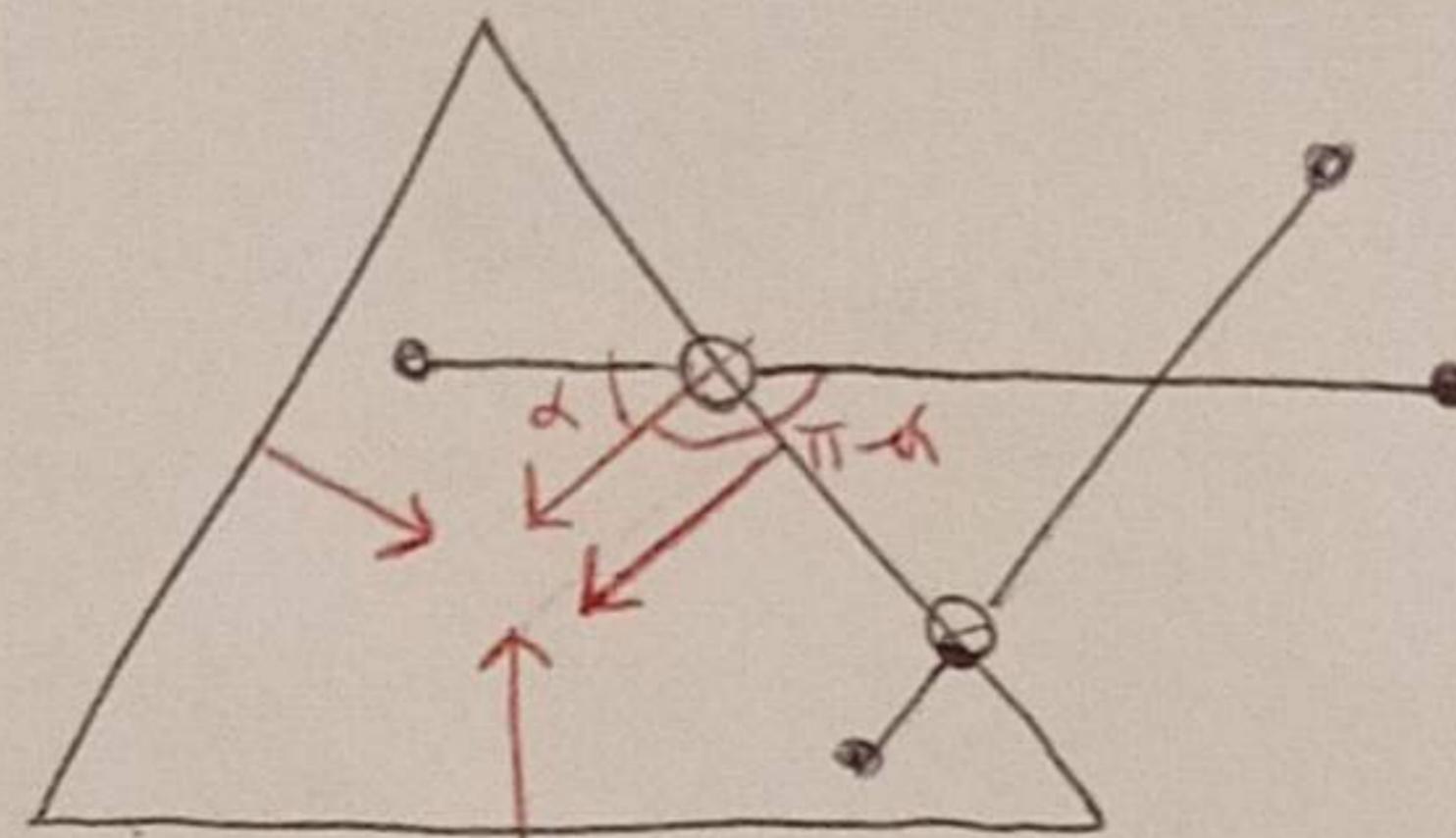
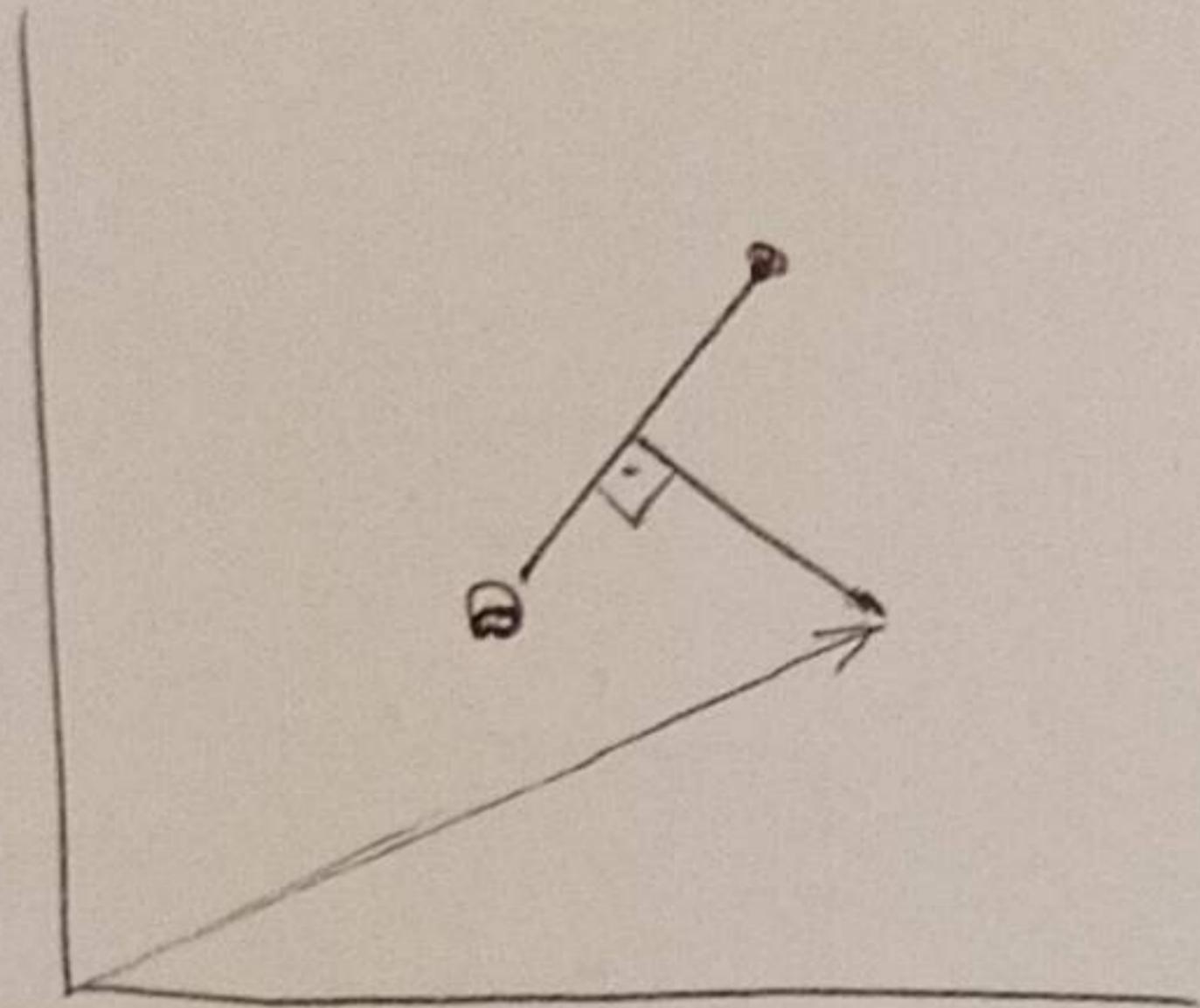


$$\sqrt{2 \cdot 1 + 2 \cdot 2}$$

$$(-1,2) \quad (1,2) \quad (2,2)$$

$$(1,1,-2) \quad \sqrt{2 \cdot 1 - 2 \cdot 2}$$

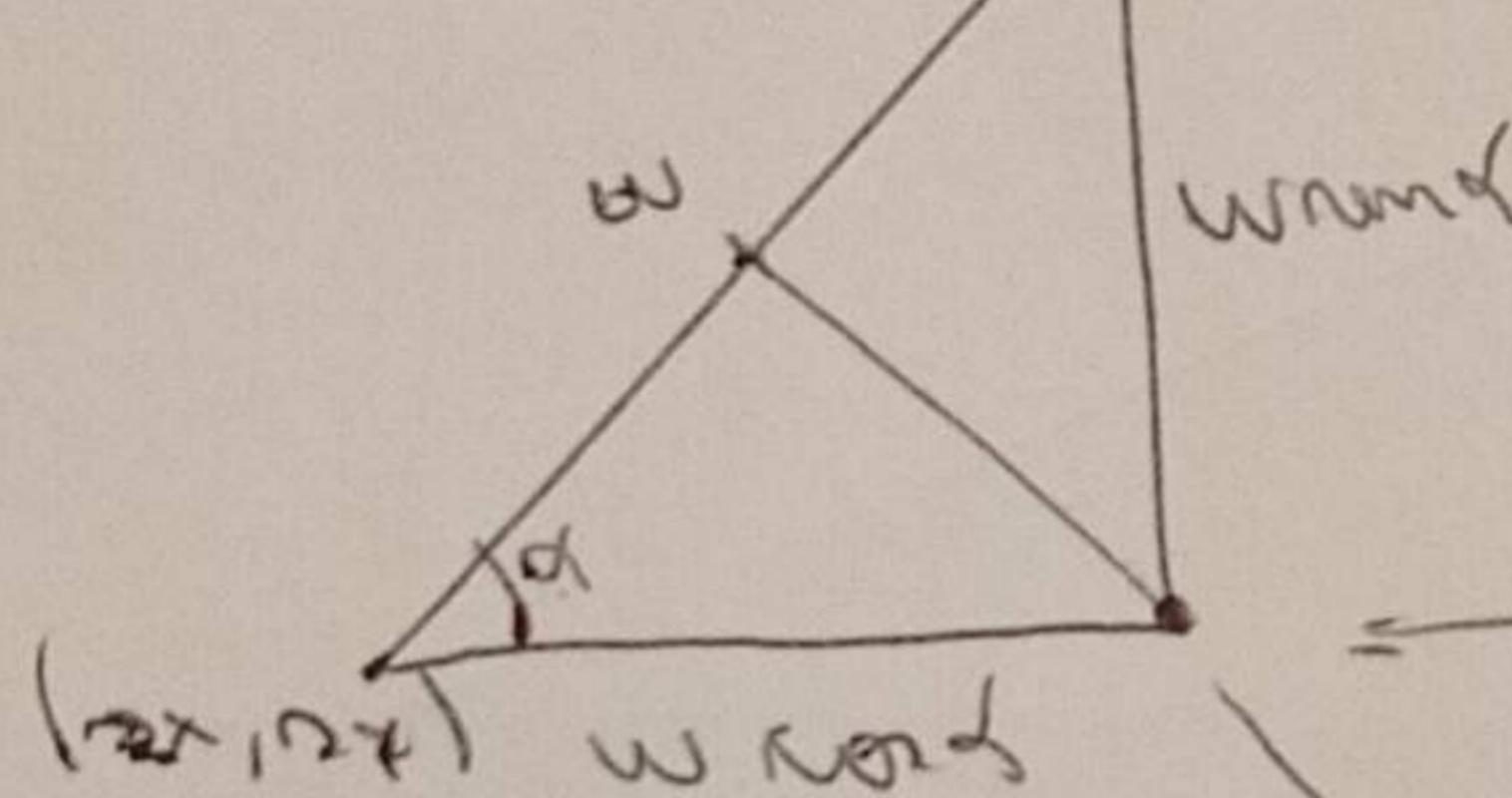
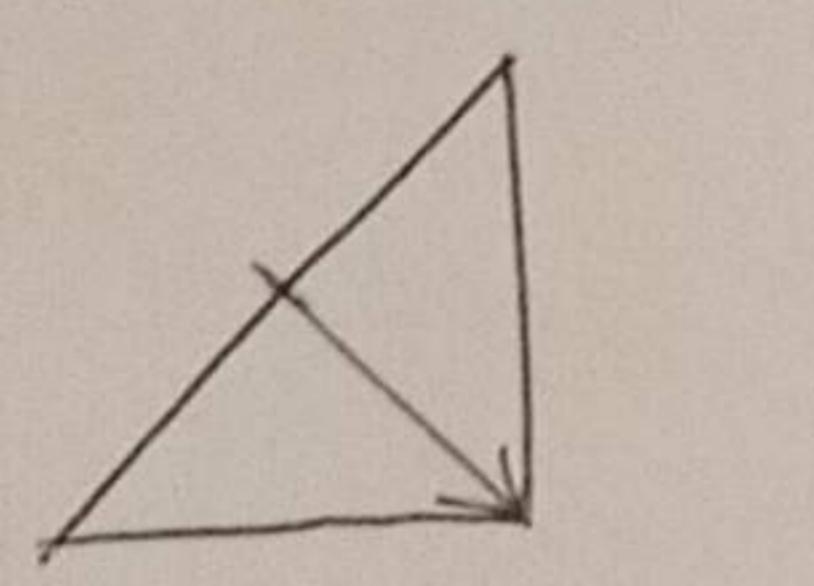
$$(1,1,-2) \quad \sqrt{2 \cdot 1 - 2 \cdot 2}$$



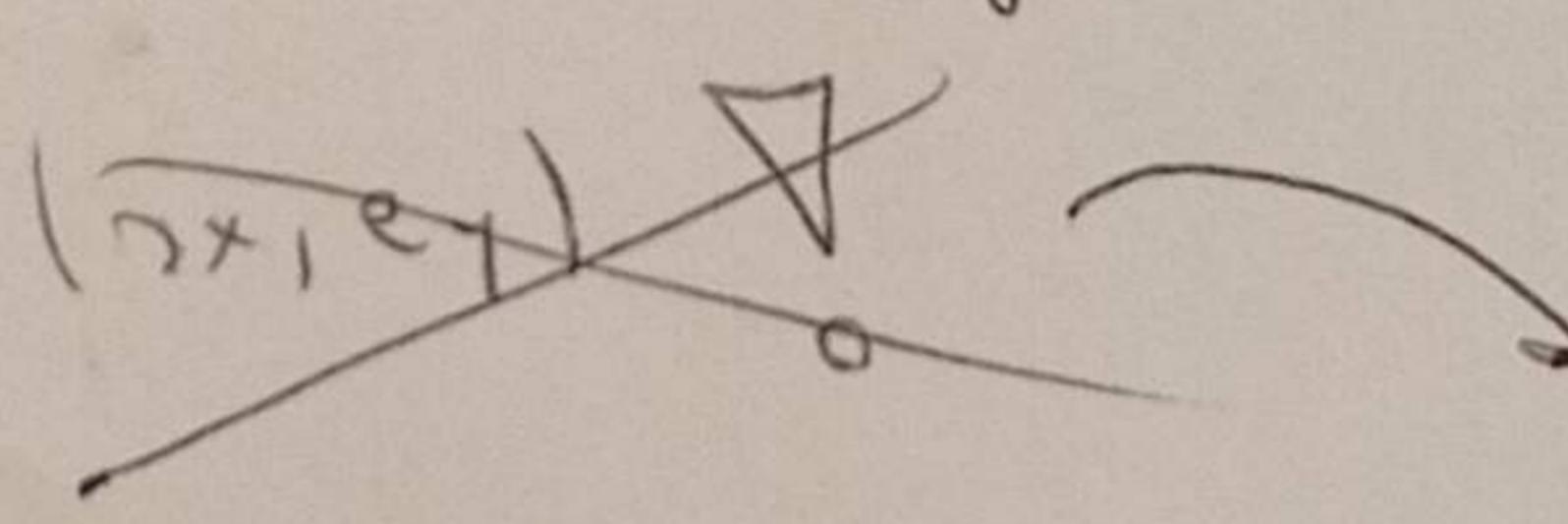
$$(4,3) - (2,4) = (2,1)$$

$$\vec{o} \cdot \vec{n} = 0 \text{ if } \cos \theta = 0$$

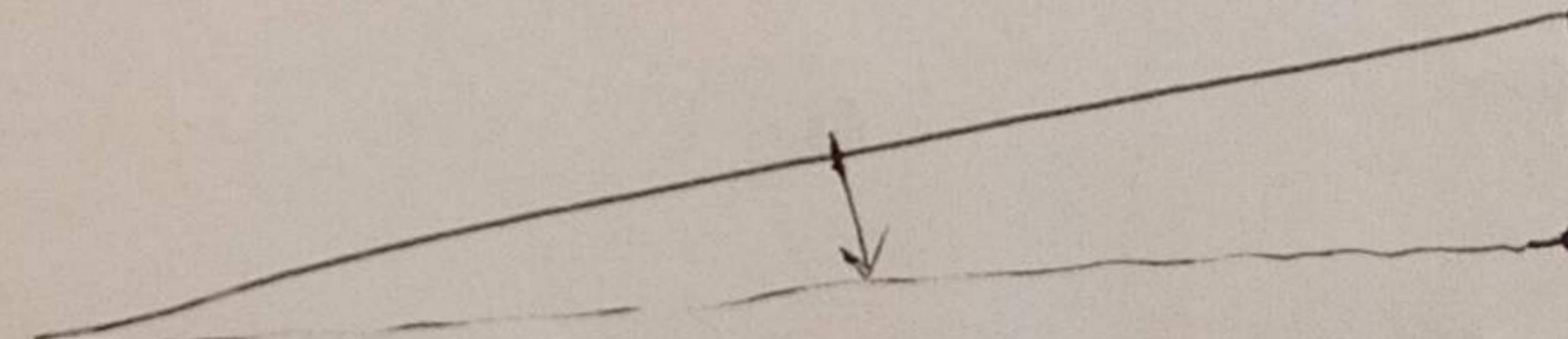
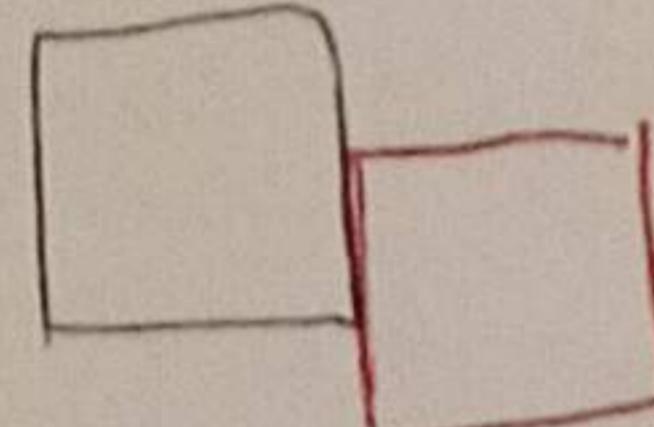
$$(e_x, e_y)$$



intersection of $(w_x, 0)$ with $(0, w_y)$

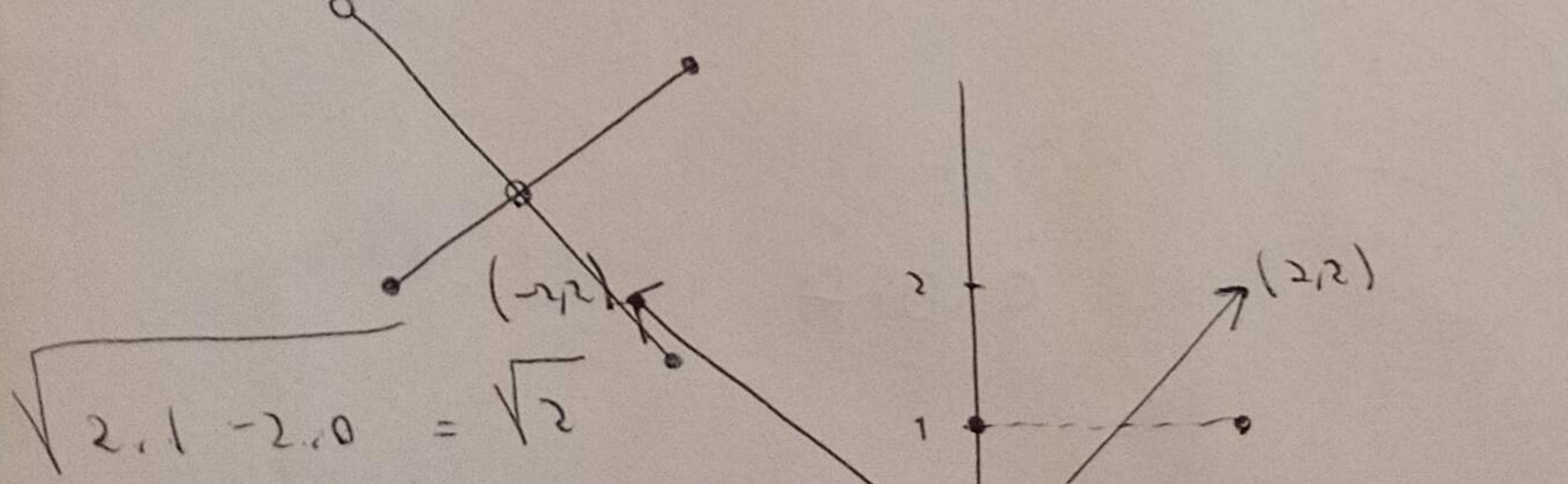
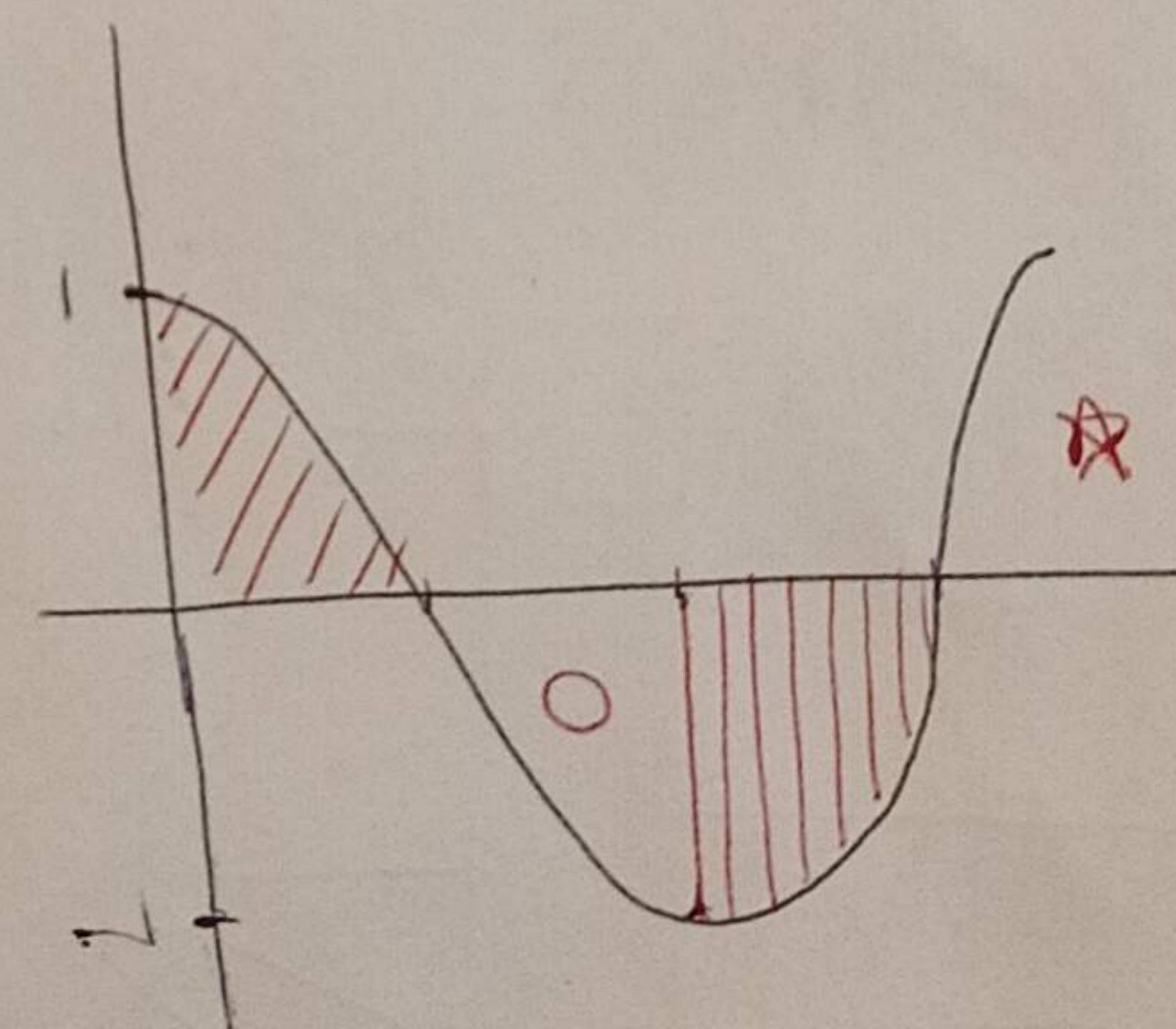
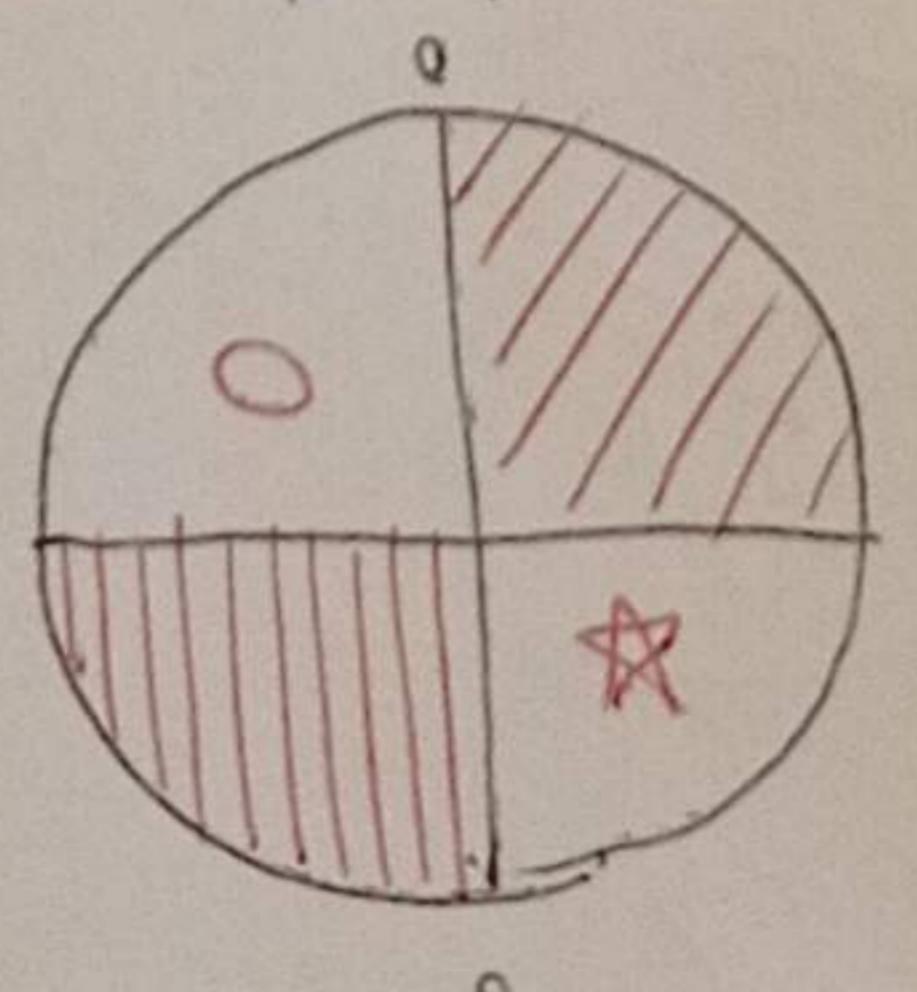


Topo:



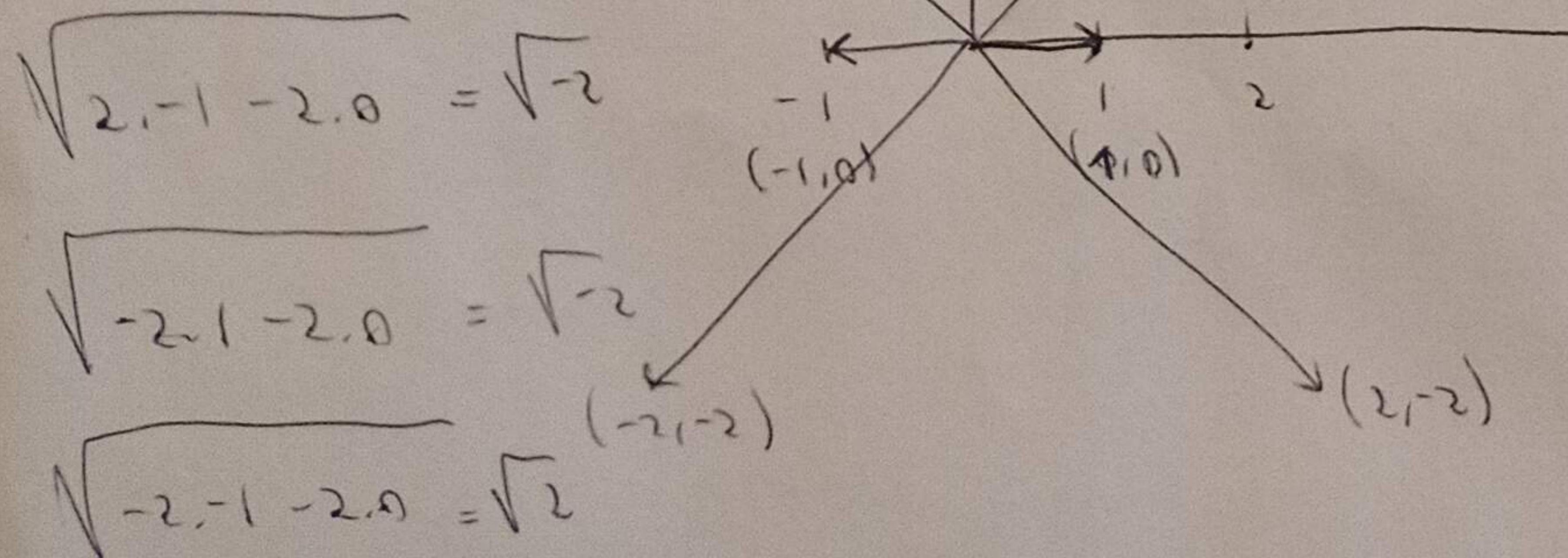
$$\vec{m} \cdot \vec{n} = \sqrt{M_x M_x + N_y N_y} = |\vec{n}|$$

$$\cos \theta = \frac{|\vec{o} \cdot \vec{n}|}{|\vec{o}| |\vec{n}|} \in [0,1]$$



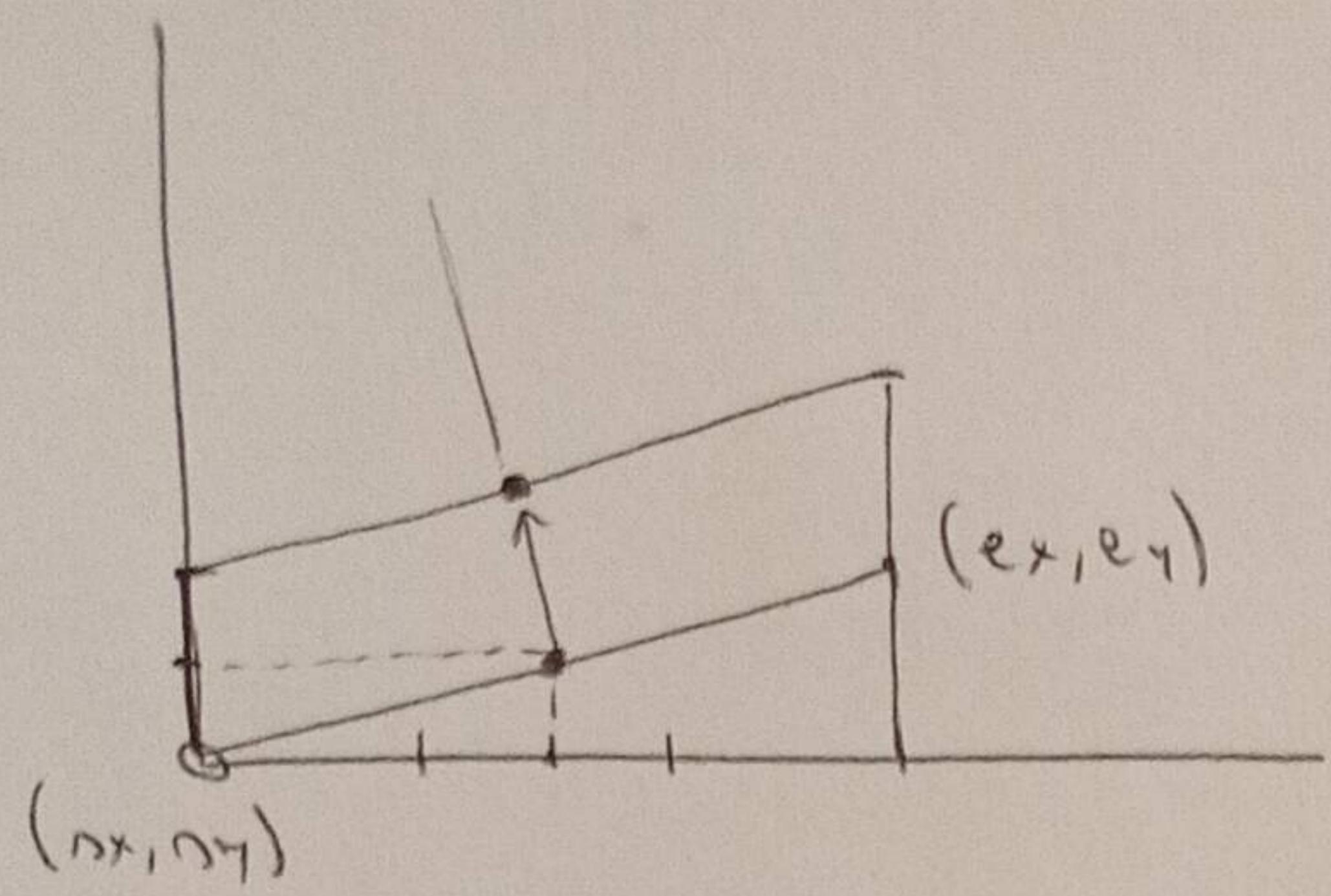
$$\sqrt{2 \cdot 1 - 2 \cdot 0} = \sqrt{2}$$

$$\sqrt{-1 \cdot 1 + 0 \cdot 0} = \sqrt{-1}$$



$$\sqrt{-2 \cdot 1 + 2 \cdot 0} = \sqrt{-2}$$

$$\sqrt{-2 \cdot 1 + 2 \cdot 0} = \sqrt{2}$$



$$\vec{R}_{\text{box}} = \left(\frac{ex - nx}{2}, \frac{ey - ny}{2} \right)$$

\vec{R}_{box}

$$\vec{R}_{\text{box}} = \left(\frac{nx + ex}{2}, \frac{ny + ey}{2} \right)$$

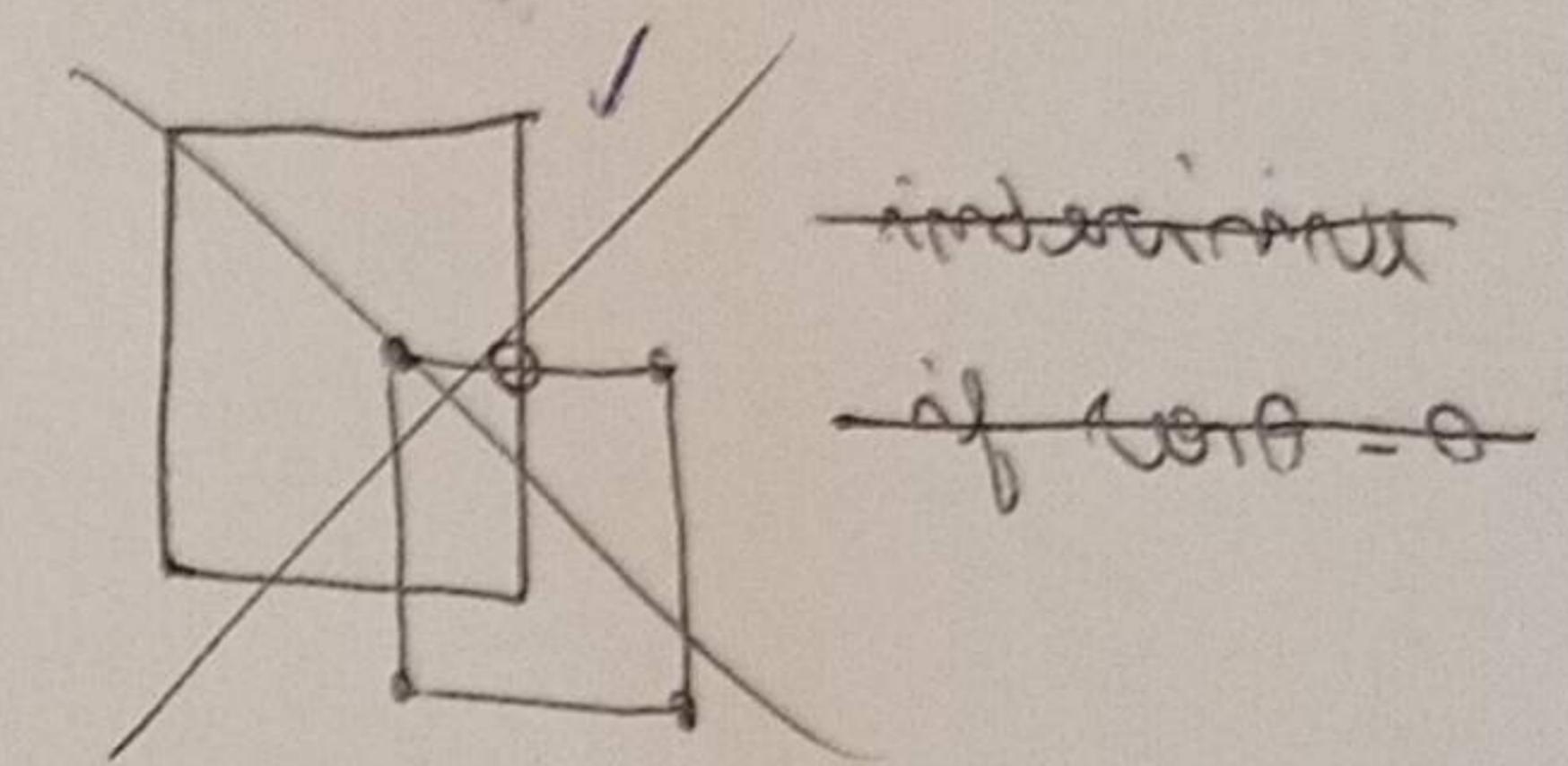
$$\vec{R}_{\text{box}} = \left(\frac{nx + ex}{2}, \frac{ny + ey}{2} \right)$$

How do I find \vec{e}_c ?

$$\vec{a} \times \vec{n} \text{ or } \vec{n} \times \vec{a}$$

$$\vec{a} = (\vec{R}_{\text{box}}, \vec{R}_{\text{box}}, 0)$$

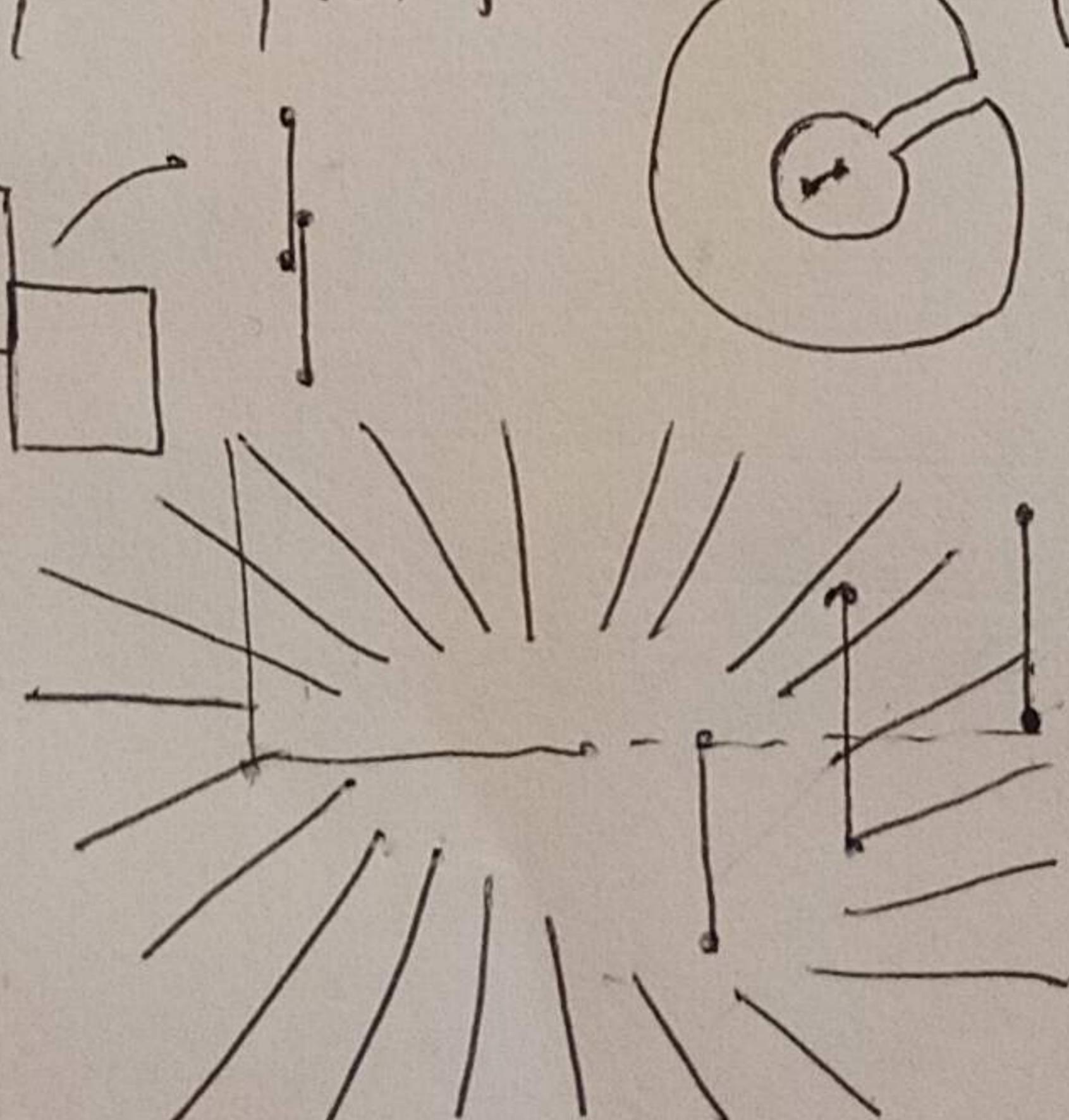
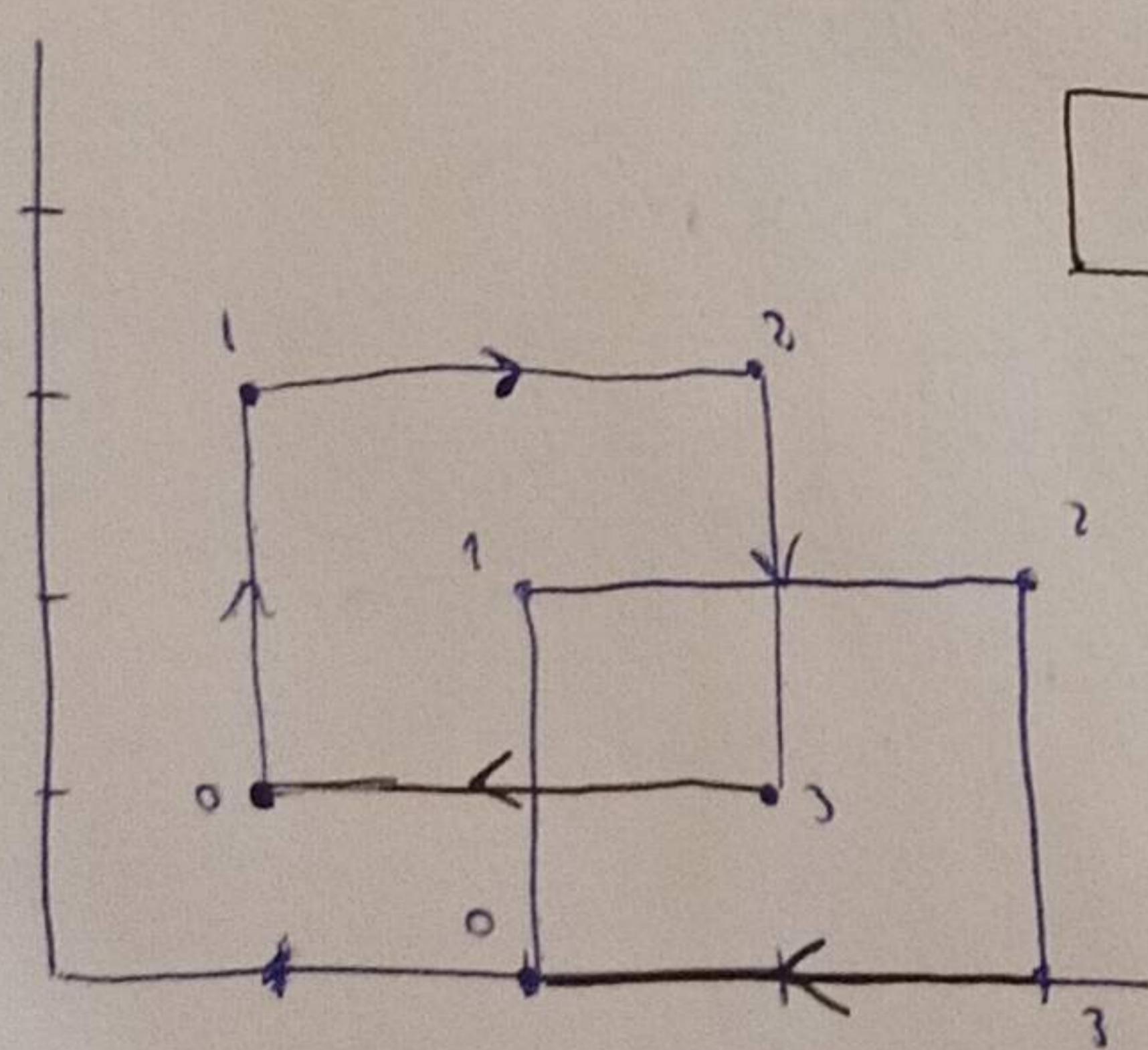
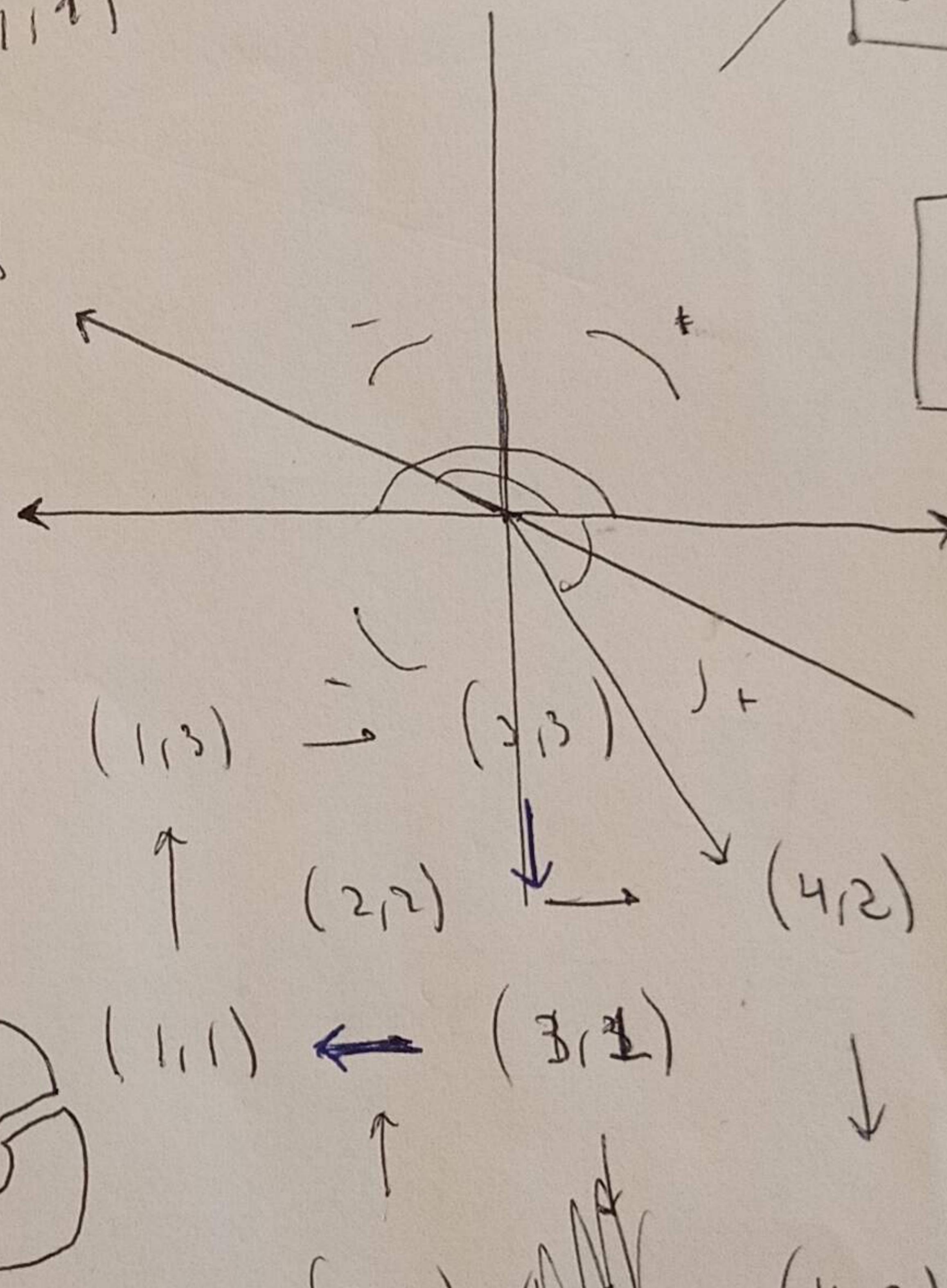
$$\vec{n} = (\vec{R}_{\text{box}}, \vec{R}_{\text{box}}, 1)$$



$$\vec{n} \cdot \vec{a} = a_x n_x + a_y n_y + a_z n_z, \text{ little messy}$$

$$\vec{n} \cdot \vec{a} = |\vec{a}| |\vec{n}| \cos \theta$$

$$\begin{vmatrix} \hat{i} & \hat{j} & \hat{k} \\ a_x & a_y & a_z \\ n_x & n_y & n_z \end{vmatrix} = (a_y n_z - a_z n_y) \hat{i} - (a_x n_z - a_z n_x) \hat{j} + (a_x n_y - a_y n_x) \hat{k}$$



$$A, B, C, D$$

$$A, B, E, D, C$$

$$y = ax + b$$

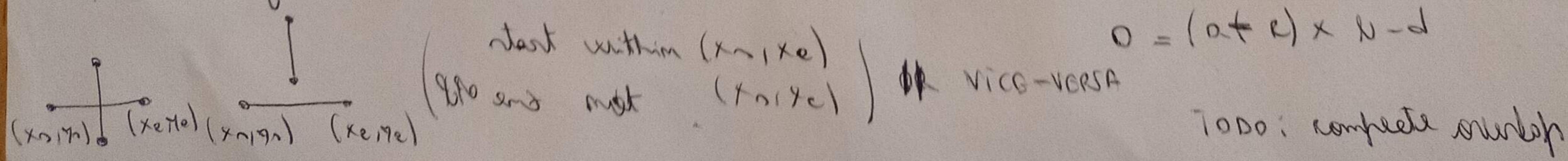
$$y = cx + d$$

There is now an intersection with, problem is that there is $\pm\infty$ here is there another way to detect intersection between two segments?

I could forget one segment on the other's turn

$$a = \frac{(x_2 - x_1)}{x_2 - x_1}$$

Does $x_1 < x_2$ contains $x_3 < x_4$



$$O = (a + c) \times N - d$$

To do: complete workup

How to know in advance which polygons will make NP disjoint first
polygons

Why I want construct a "polygon" in a list of polygons?

- time where would it fail?

As polygon size went up it becomes more expensive to clip

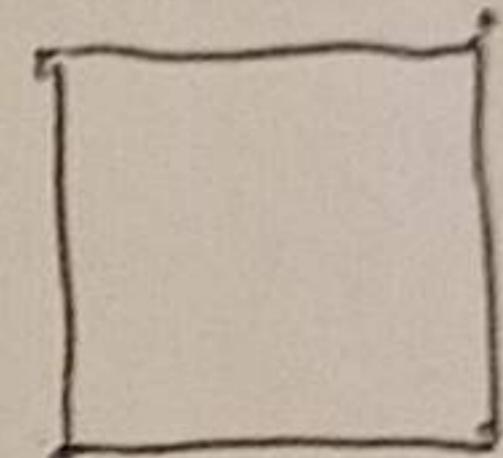
- it would fail in drawing but ^{before play} there I could split them in

Polygon :: from - numbered - segments (...) → Vect< Polygon >

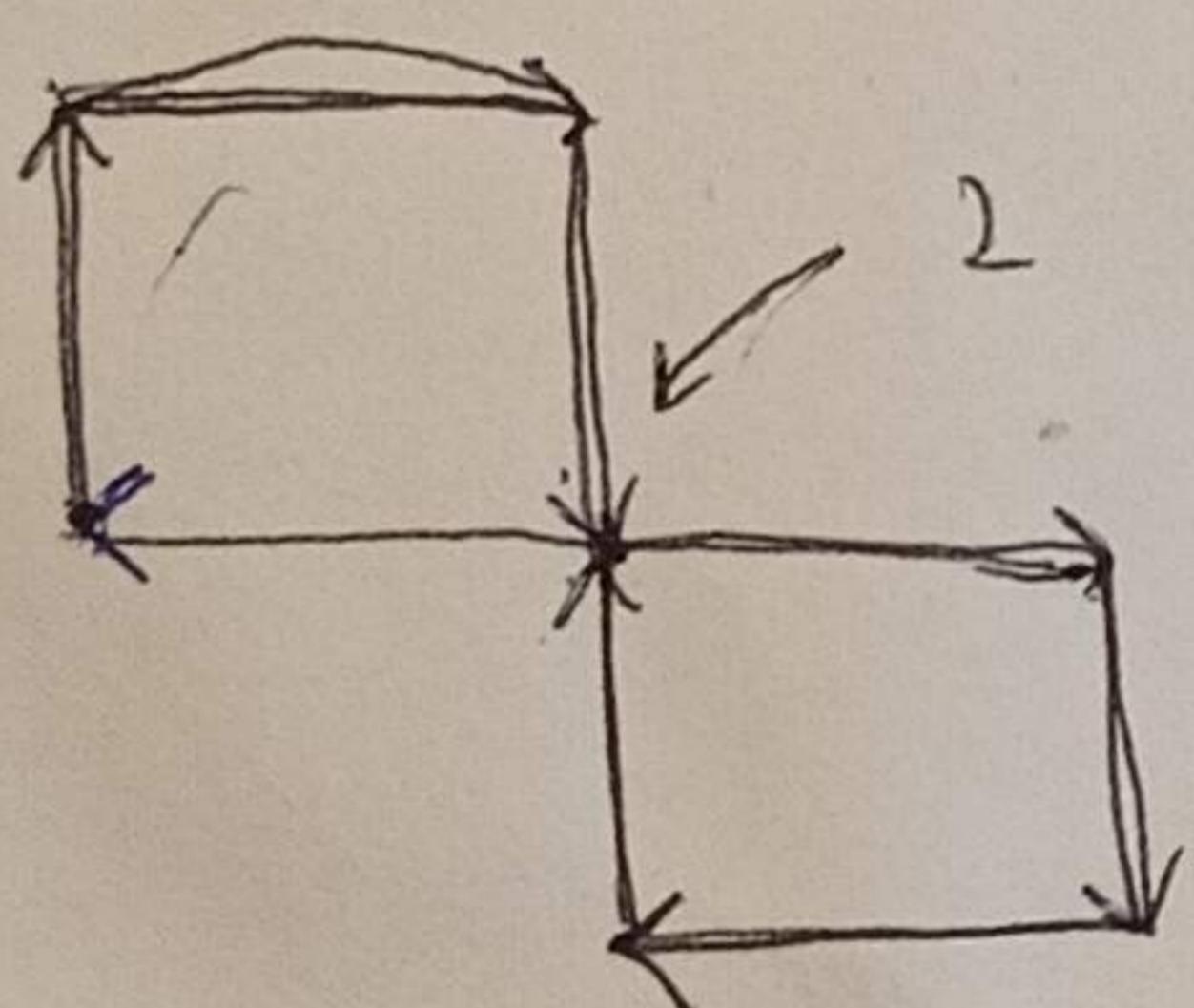
clip should take and return segments

first def retN

num returns list of polygons



edge case:

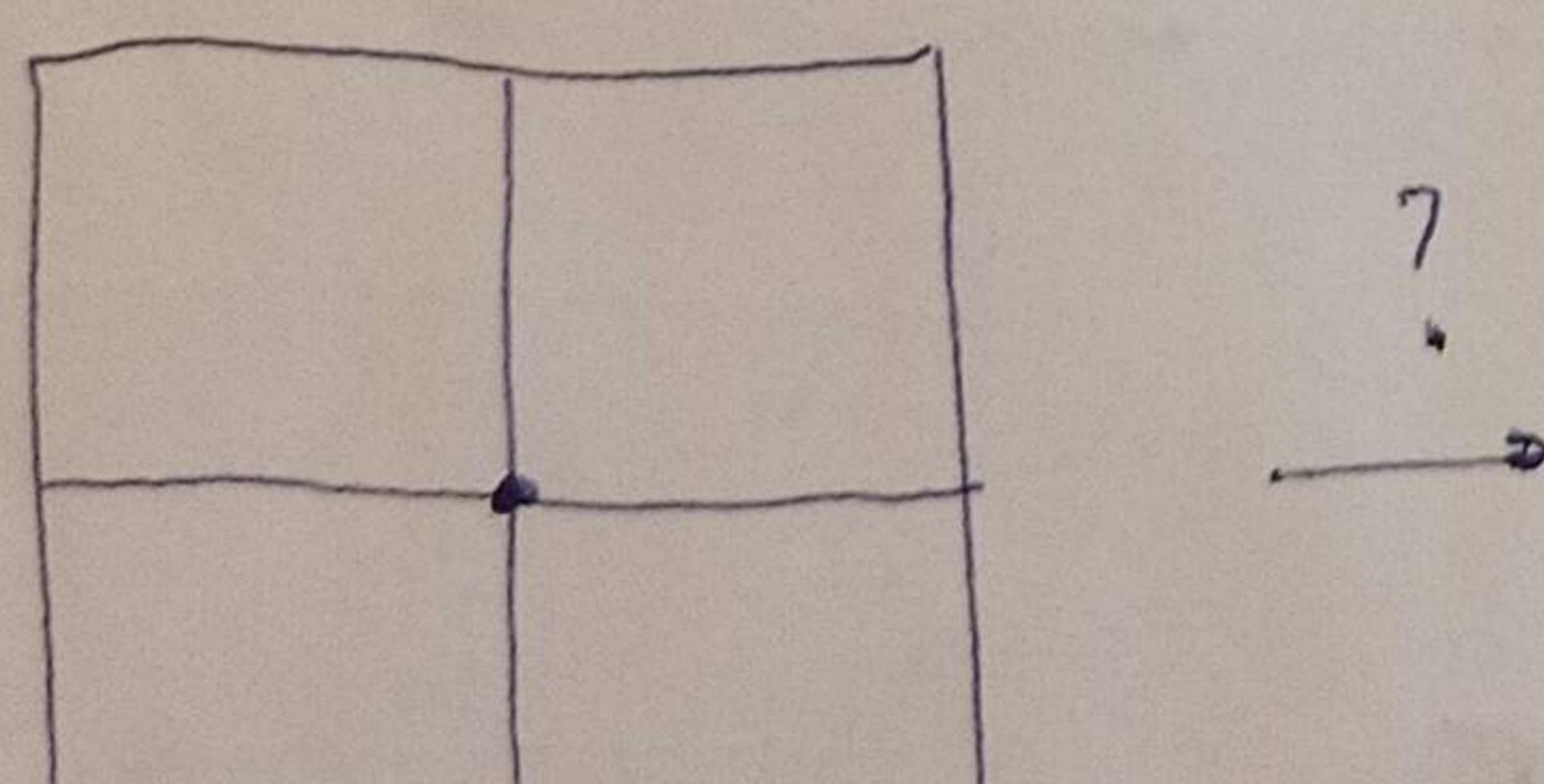


I need to

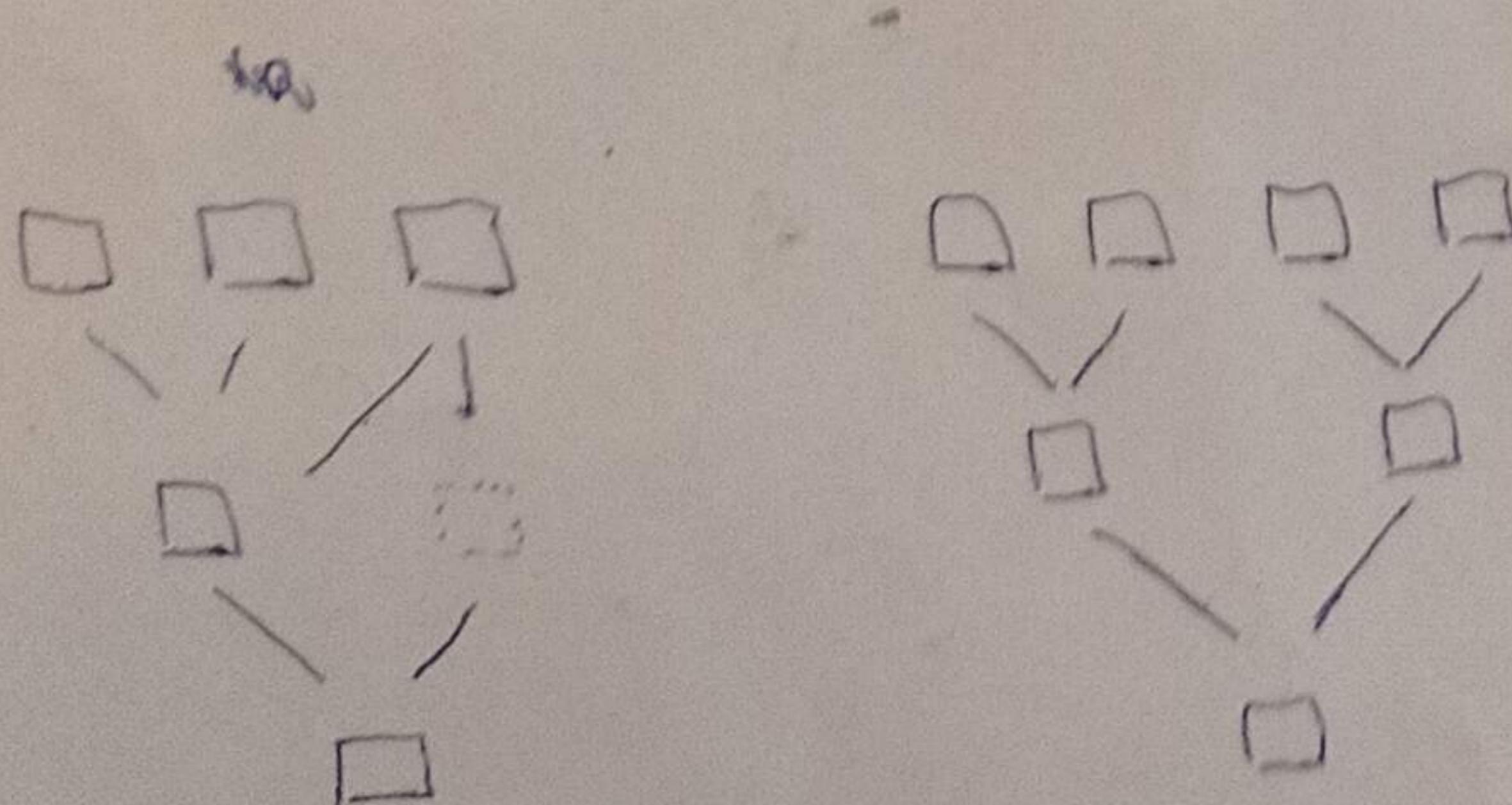
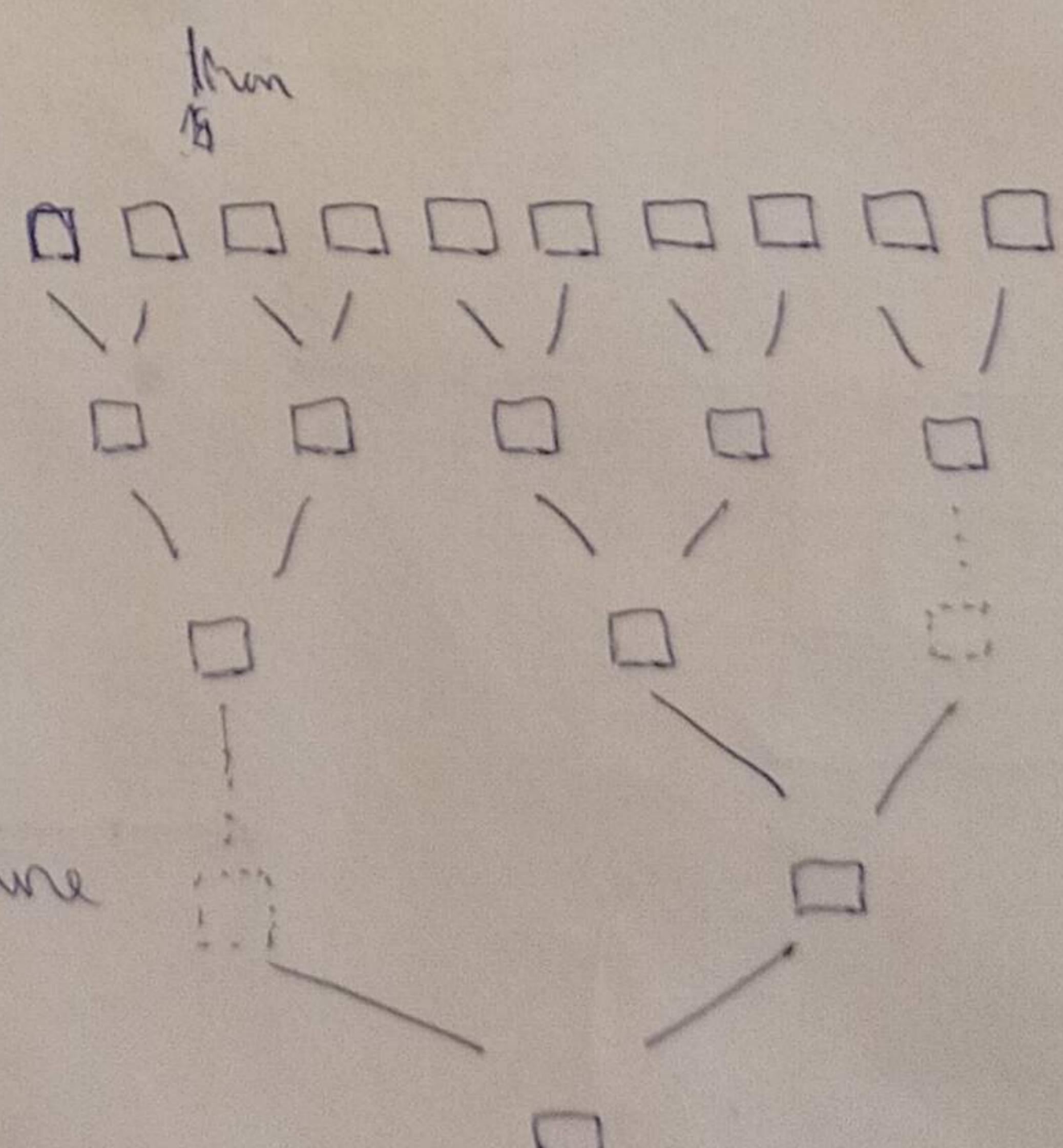
remove won't work because I need to
handle for thin edge case

② channel Works as long as there's a border alive

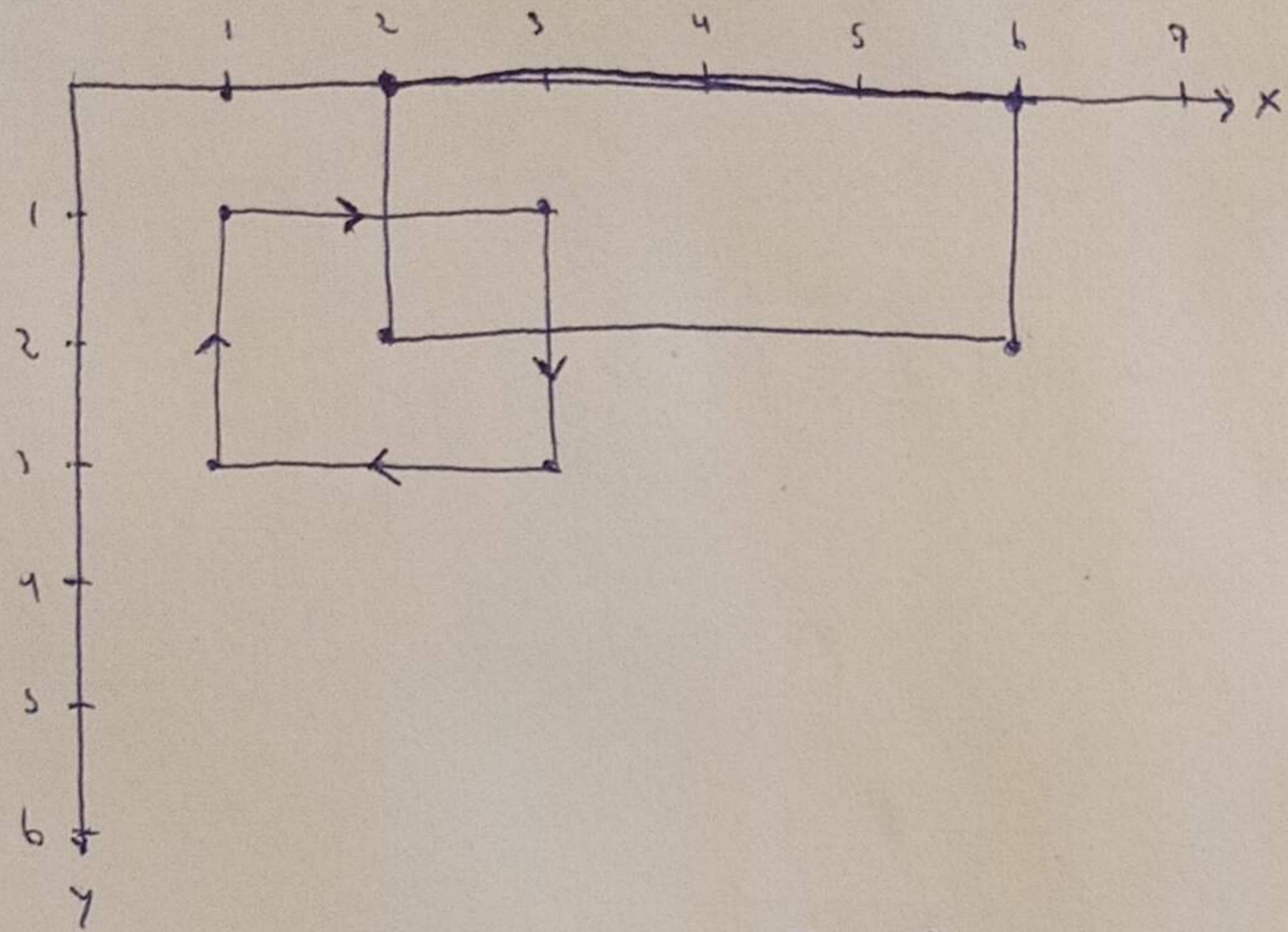
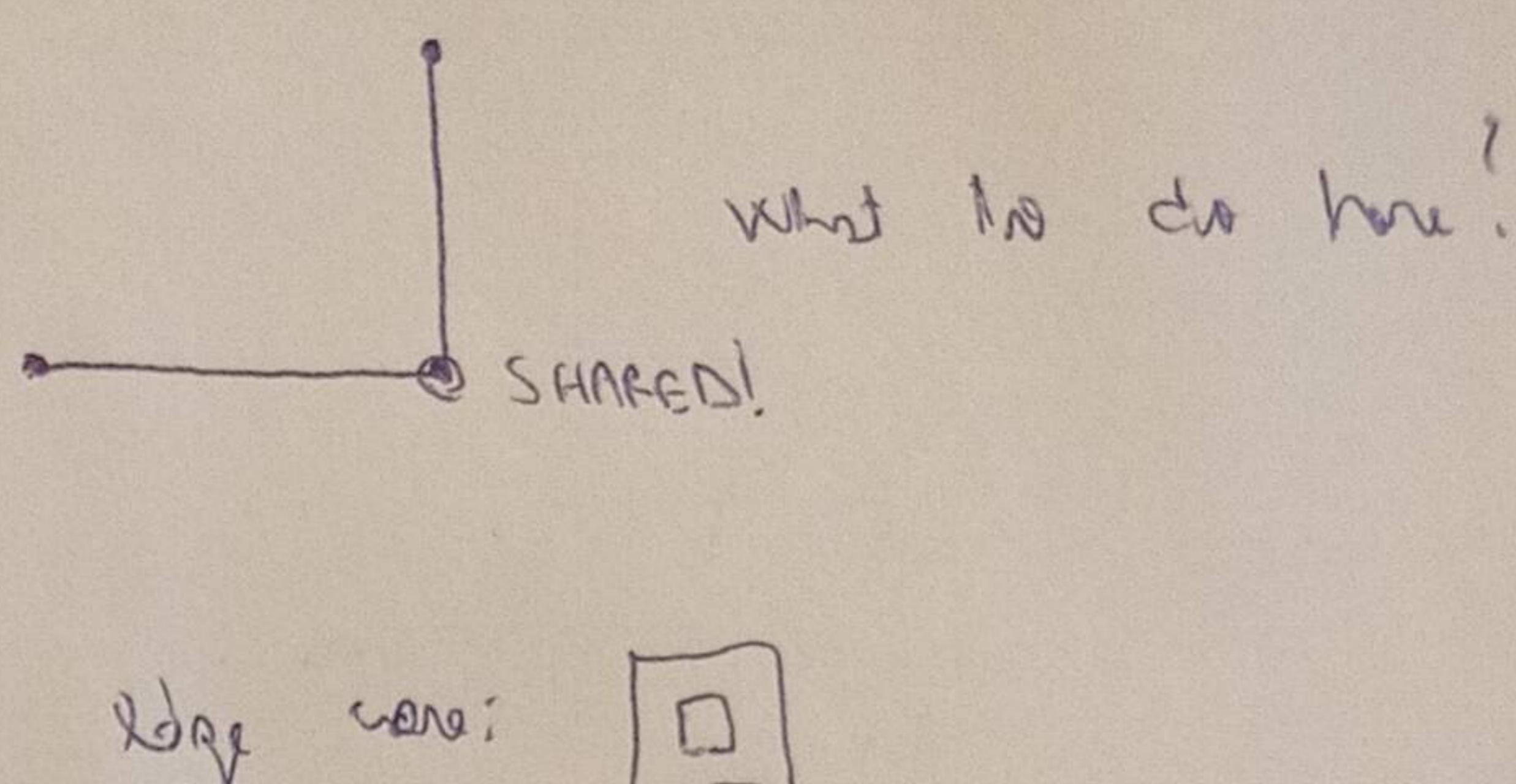
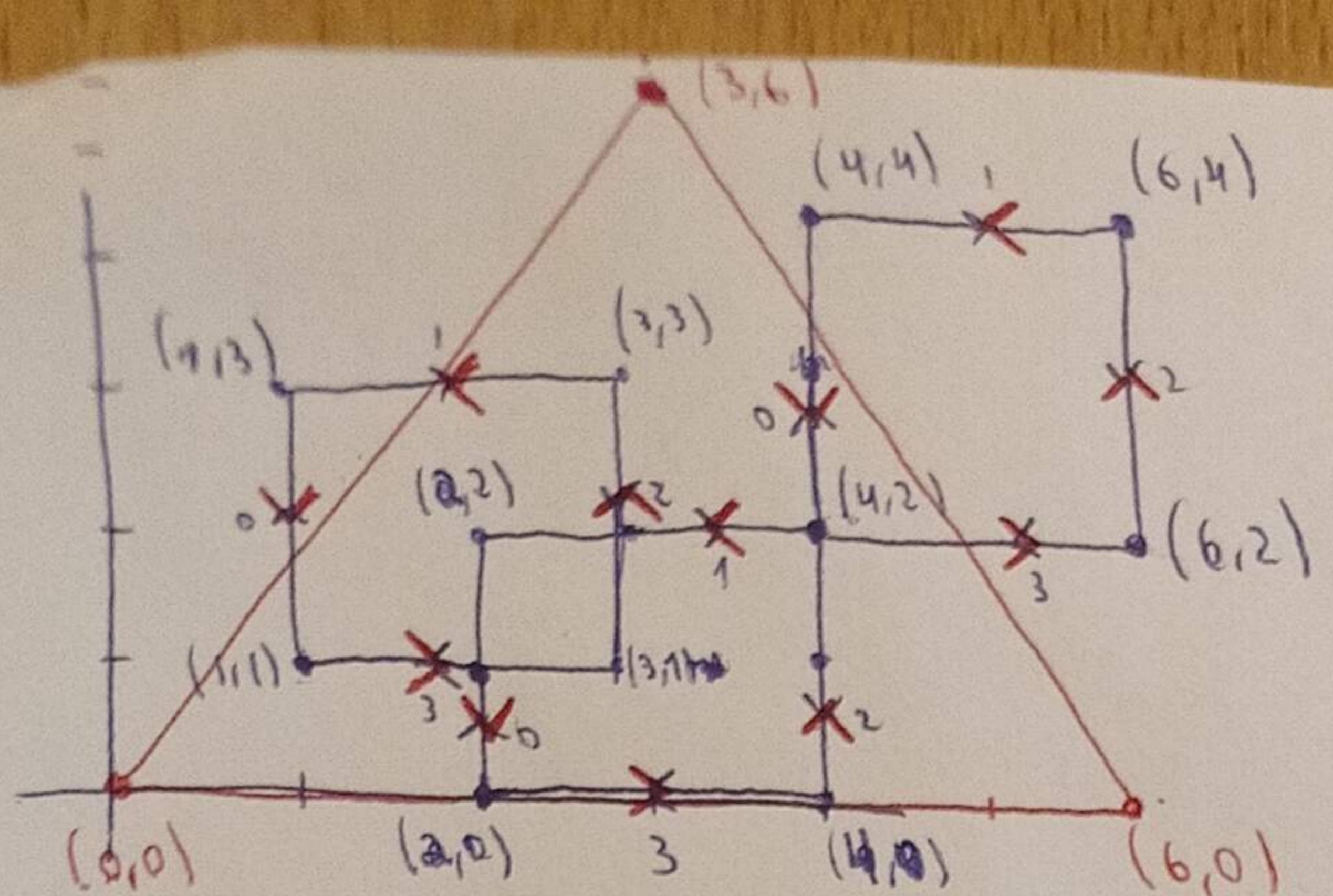
I'd like to take 2 entries from the channel and read it up to a threshold
for processing. At the end of processing a Vect< Segment > is sent back to the channel



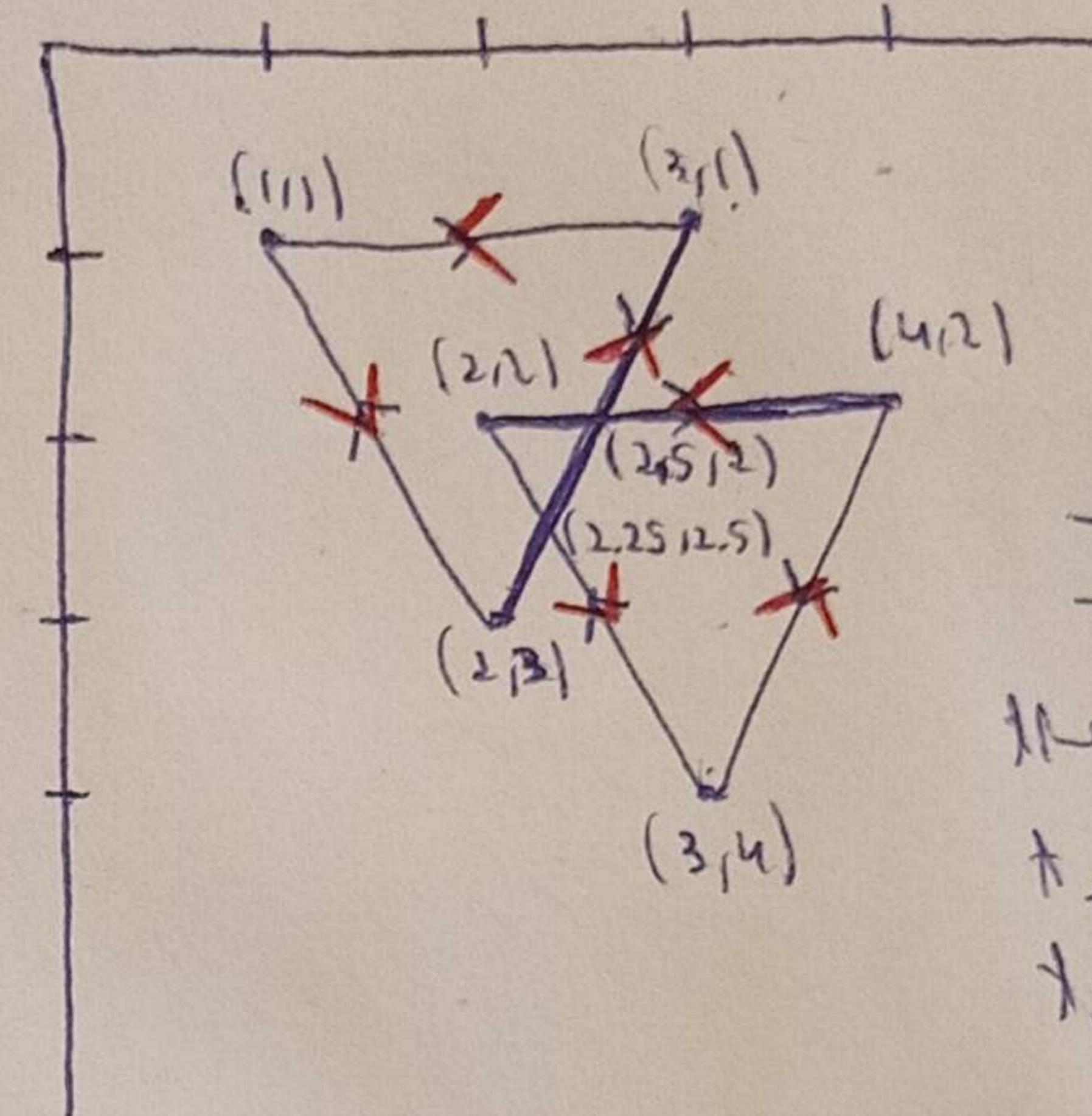
10
5 ready



removal is not there



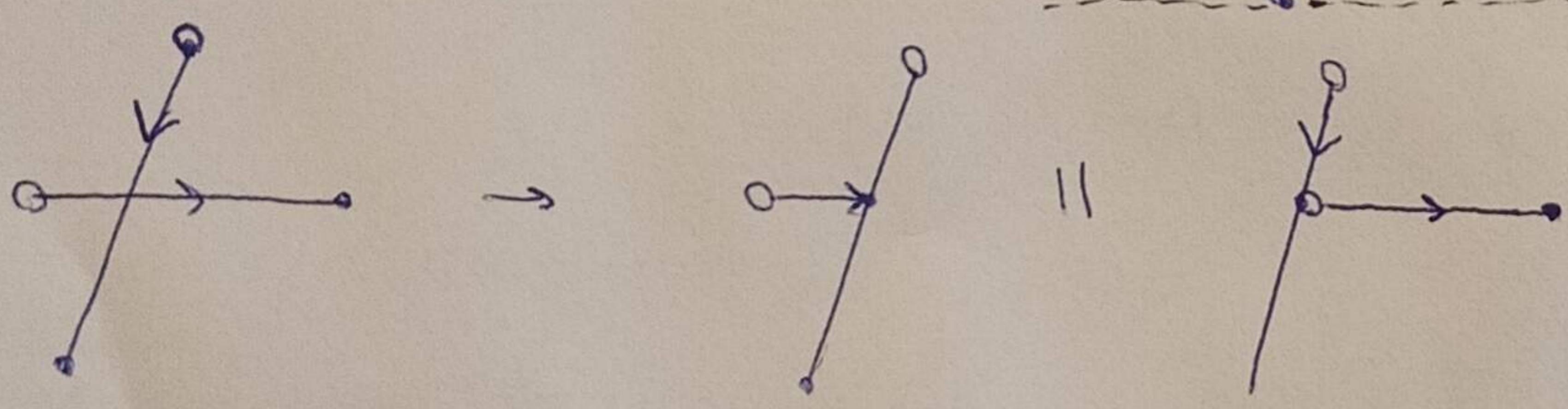
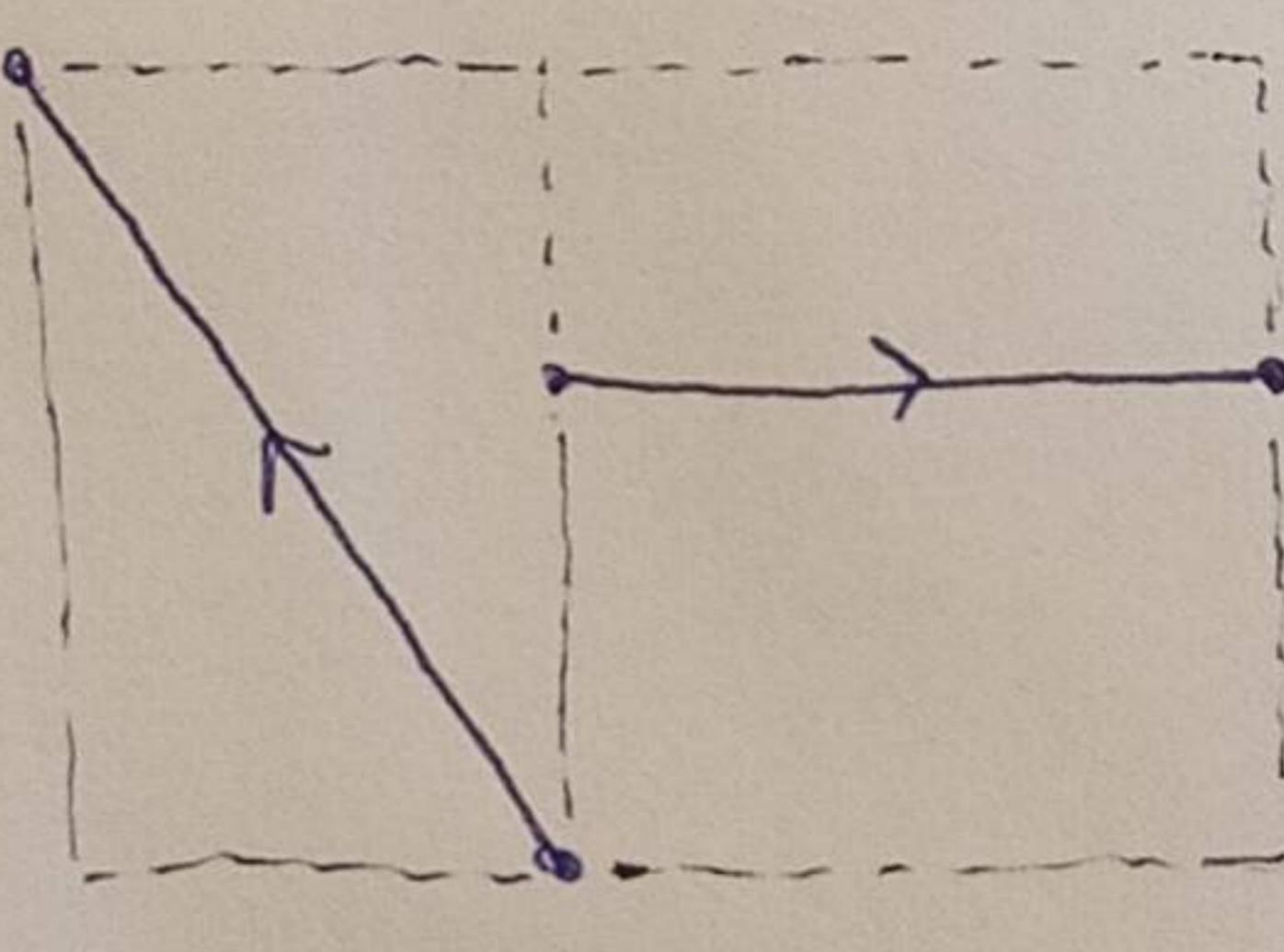
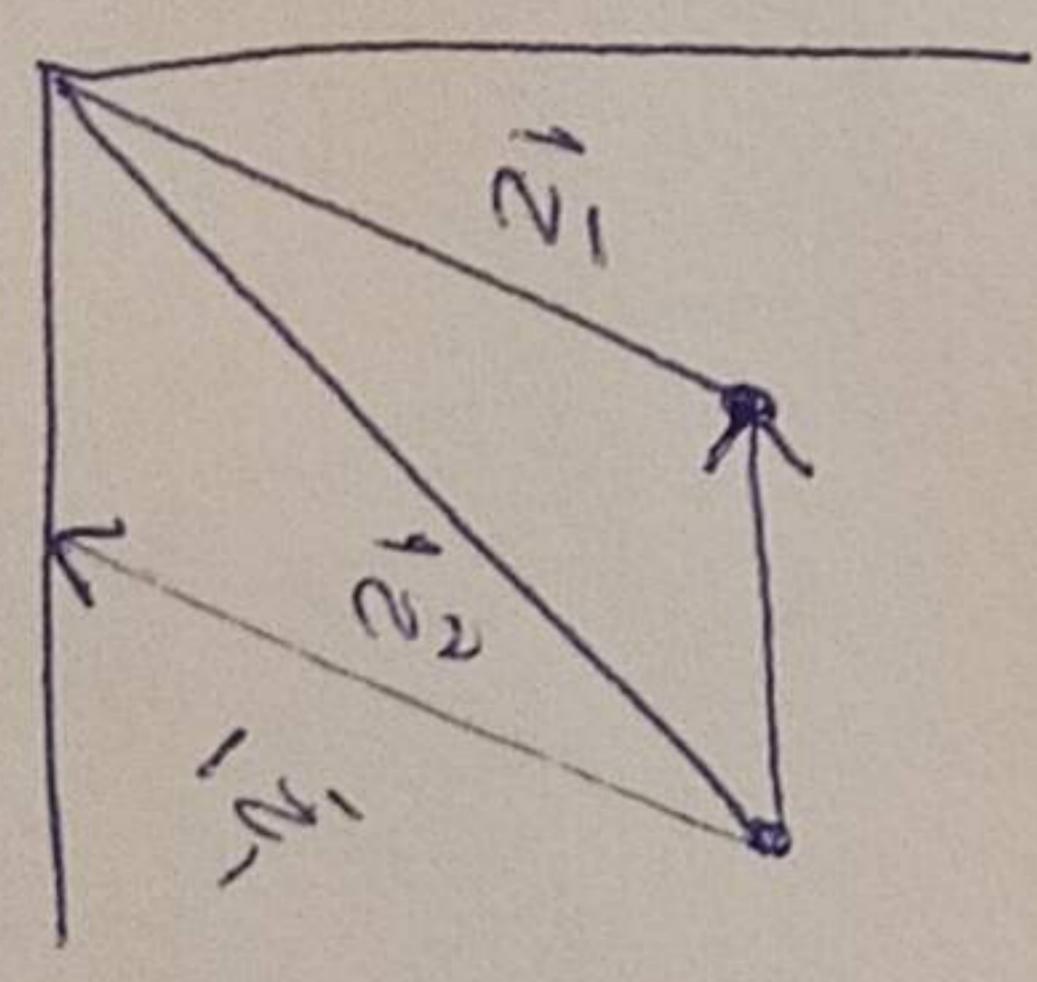
$$\vec{v}_2 - \vec{v}_1$$



I think I should correct the creation of quadrilaterals to keep the correct orientation. It already is

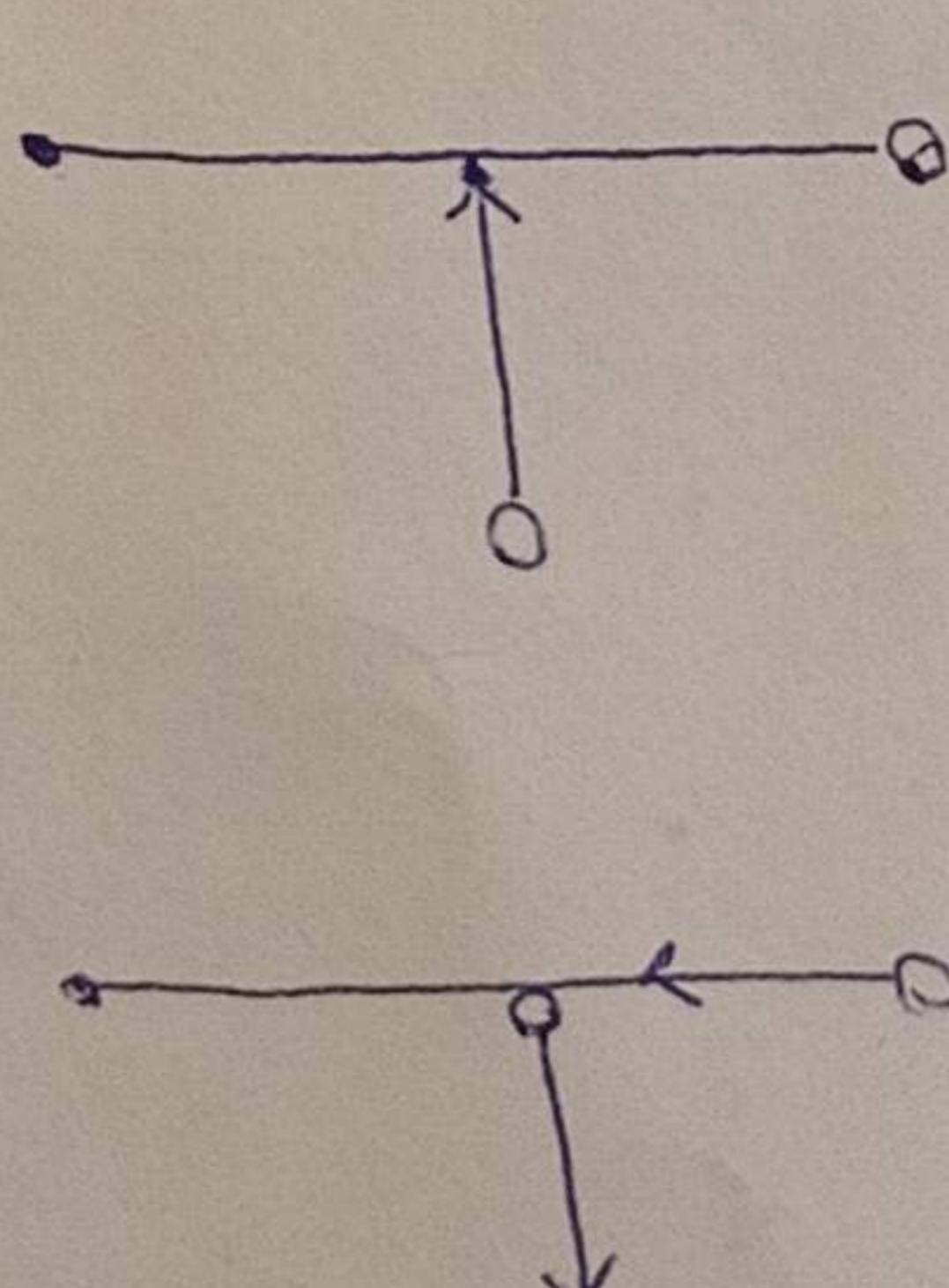
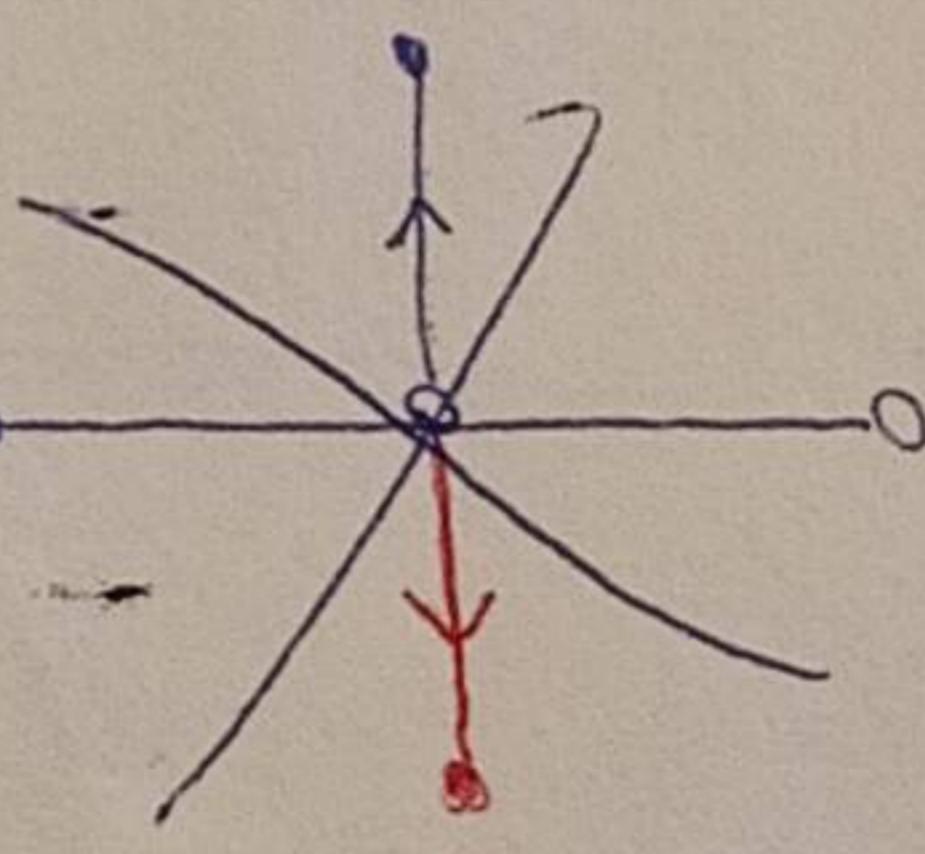
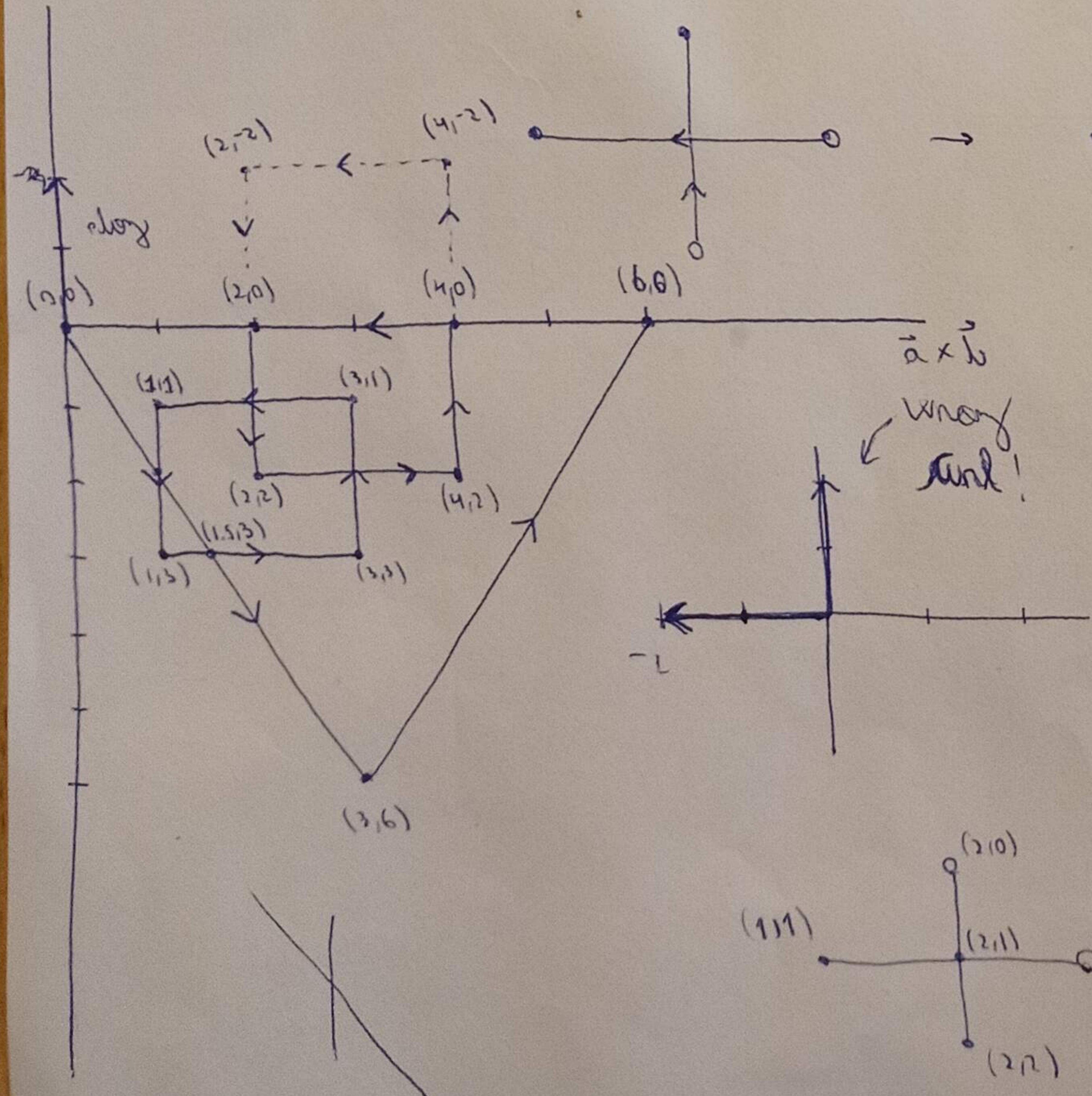
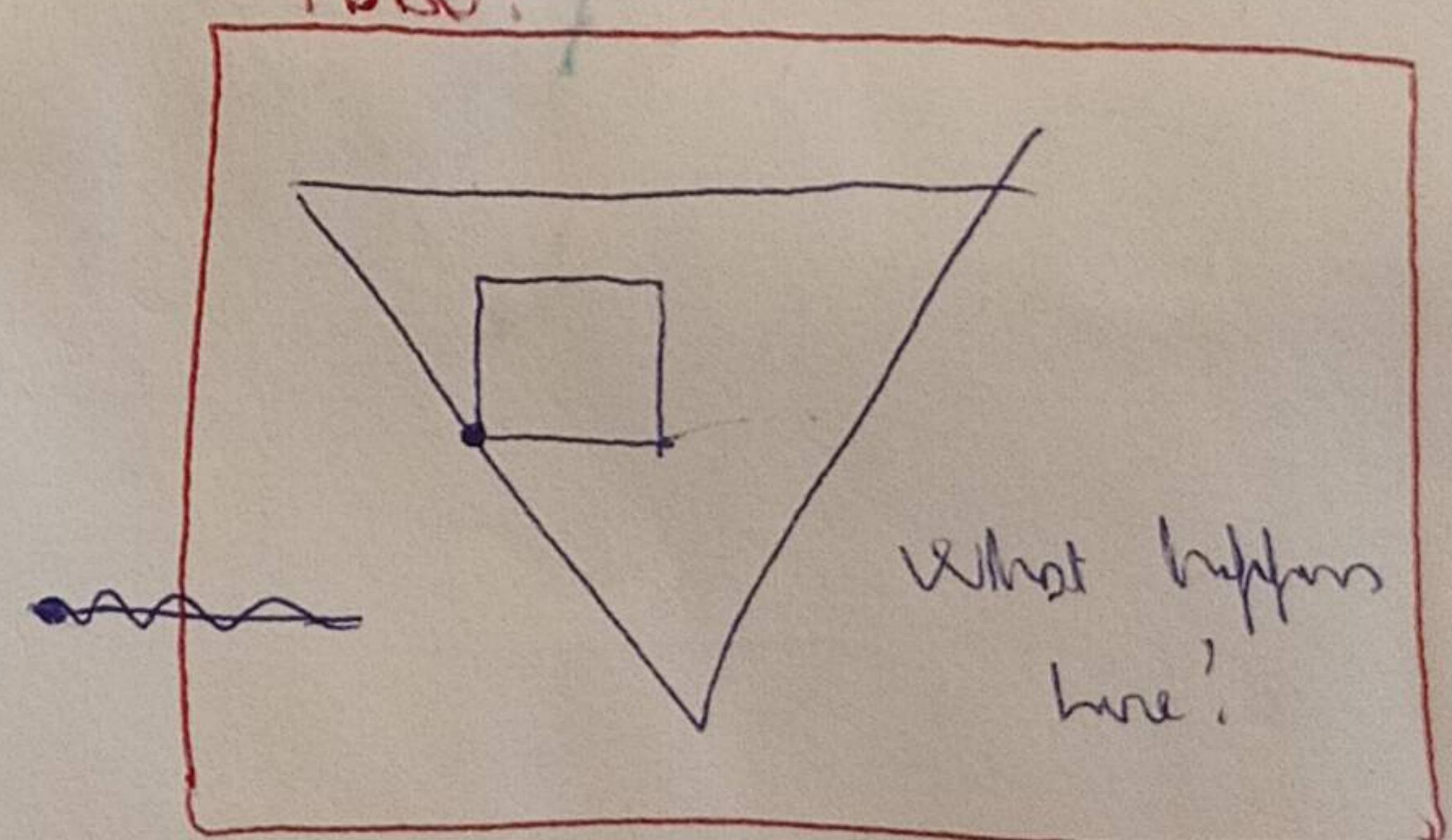
- getting no clipping? Why?
- means wrong segments were kept

now 6 different routes



BOTH POINT OUTWARDS!!!

TODO:



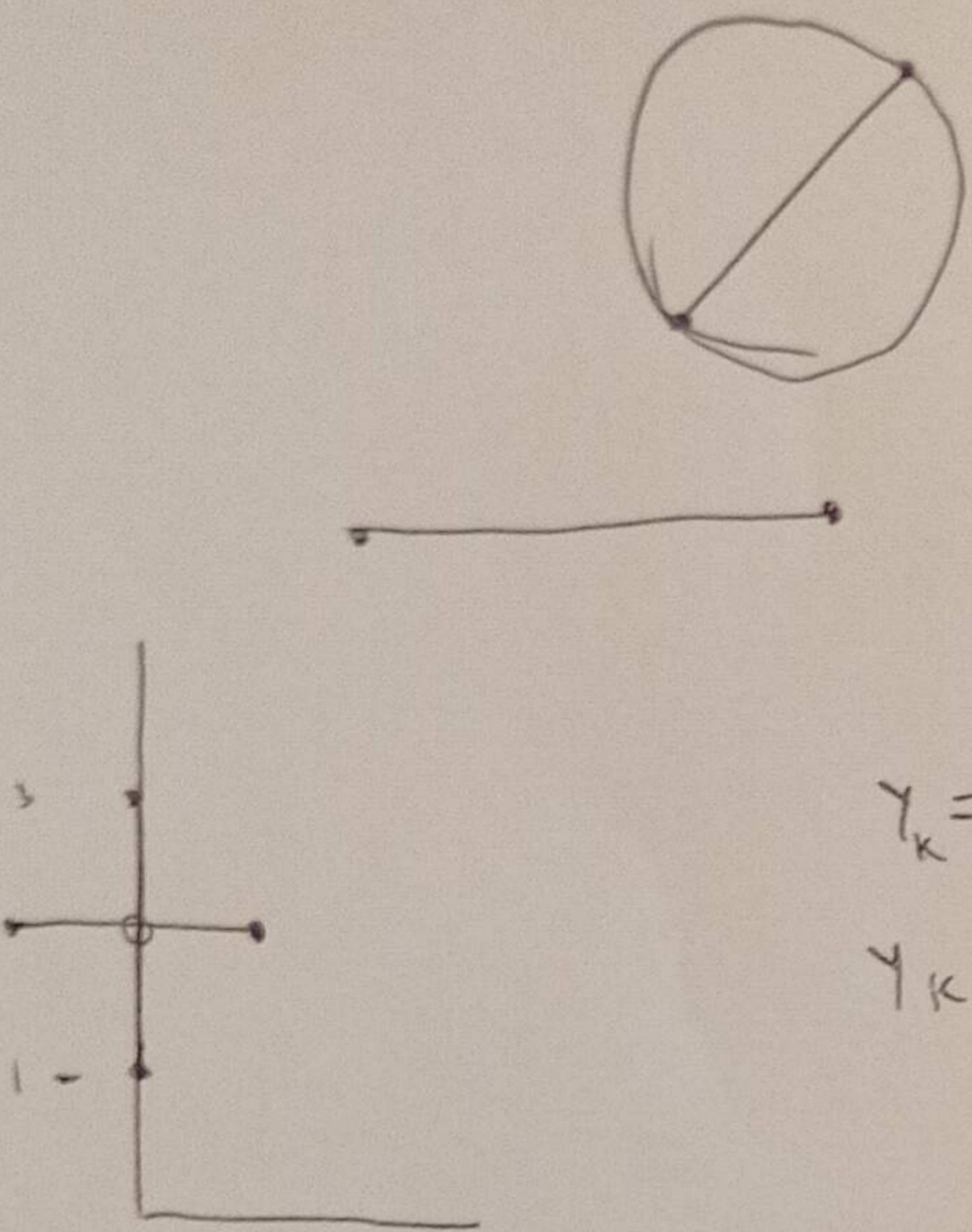
BOTH POINT INWARDS!!!

$$(2,1) \rightarrow (2,0)$$

$$(2,-2), (0,-1)$$

$$(0, -1)$$

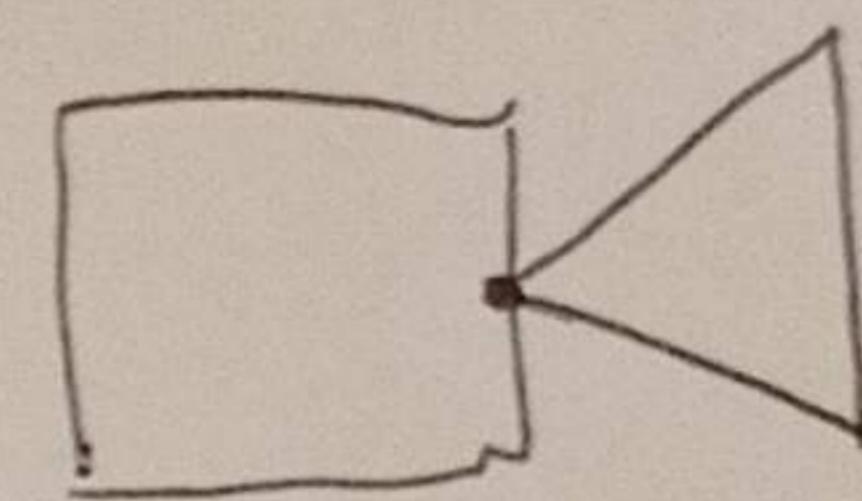
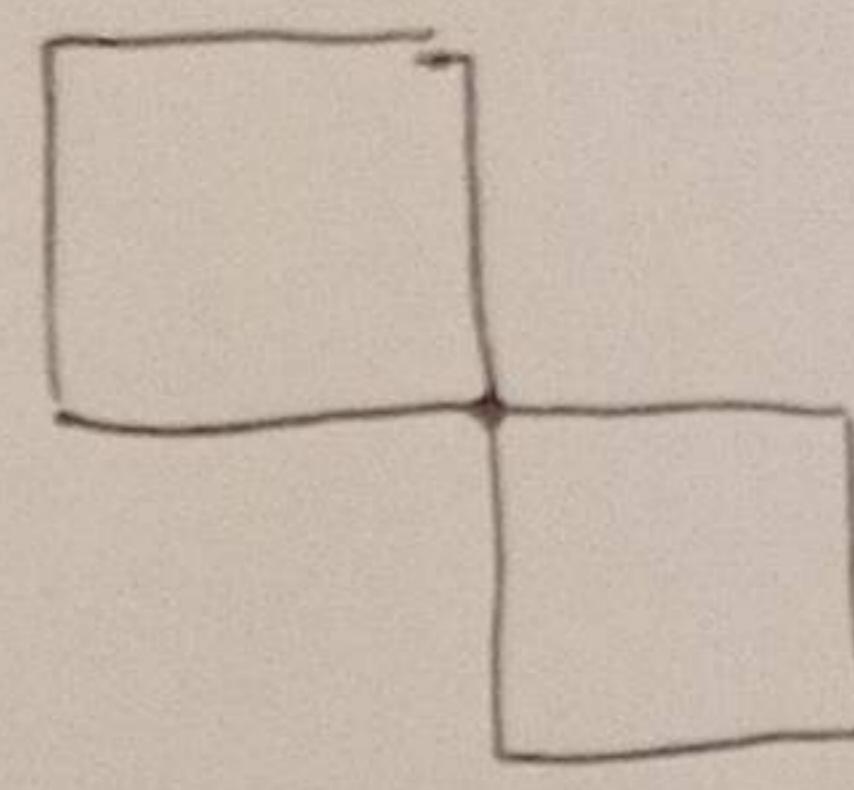
$$(2,0) - (2,1) = (0, -1)$$



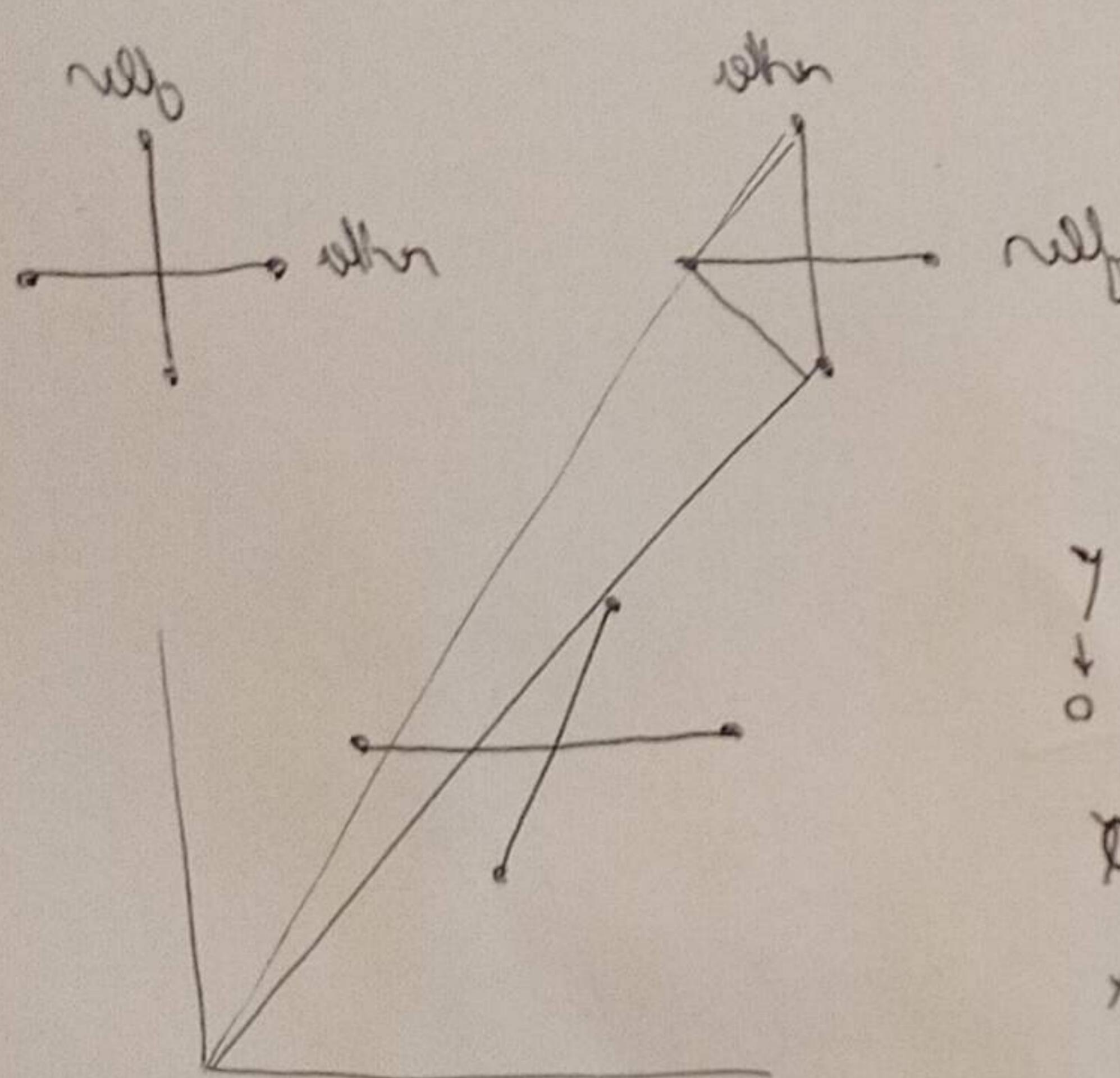
$$\gamma_k = \alpha x_k + b$$

$$\gamma_k = \alpha x_k + b$$

↓ 0 ↓ 2



$$0 = (\alpha - c)x_k + (b - d)$$



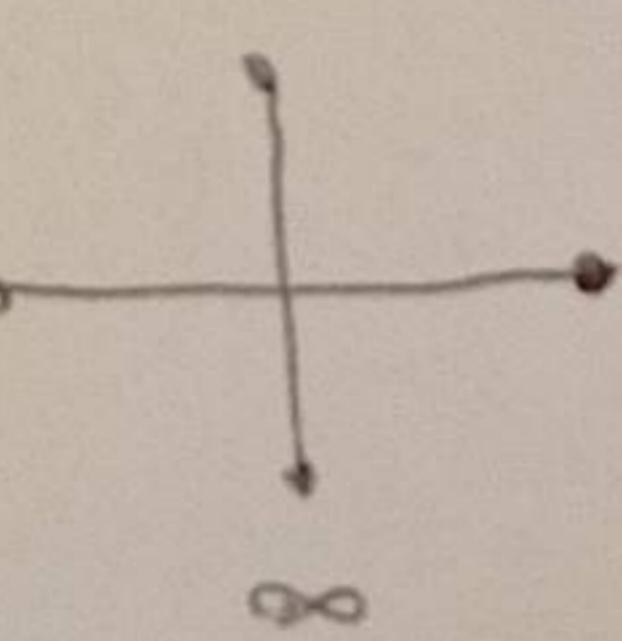
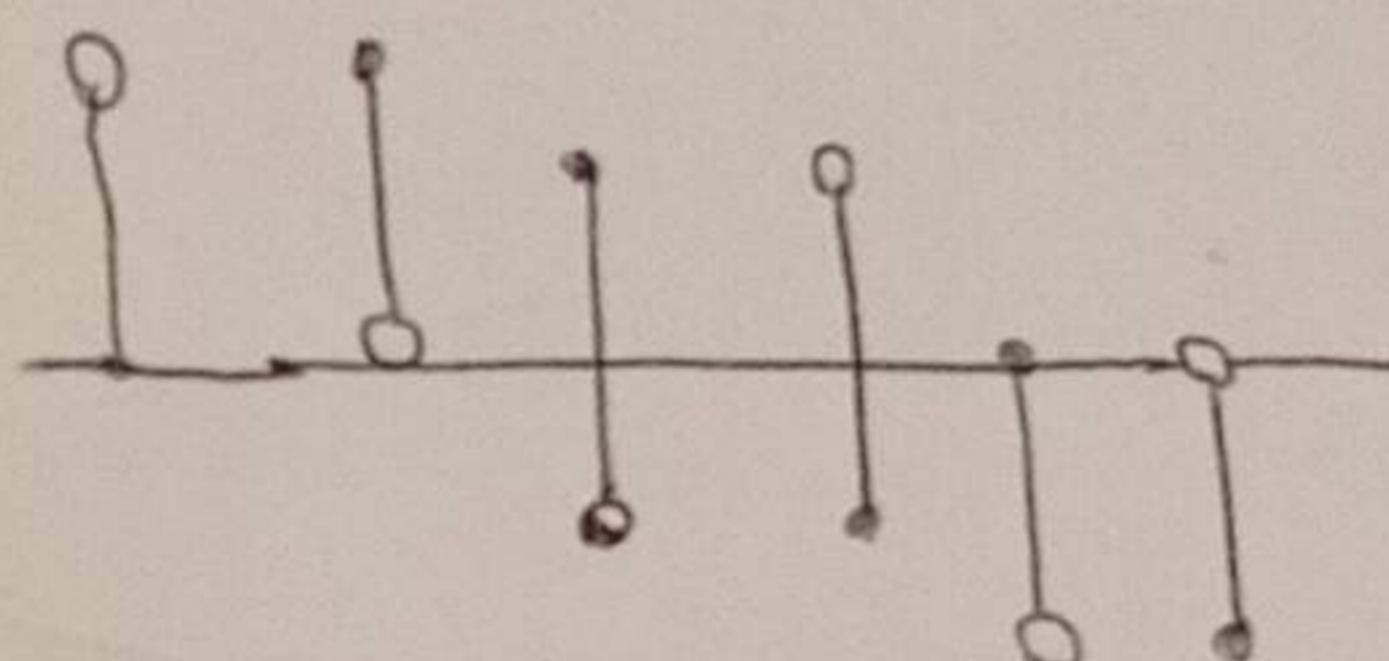
$$x_k = \frac{b-d}{\alpha - c} = \frac{-\infty + 2}{+\infty + 0}$$

$$\gamma = \alpha x + b$$

↓ 0 ↓ 2

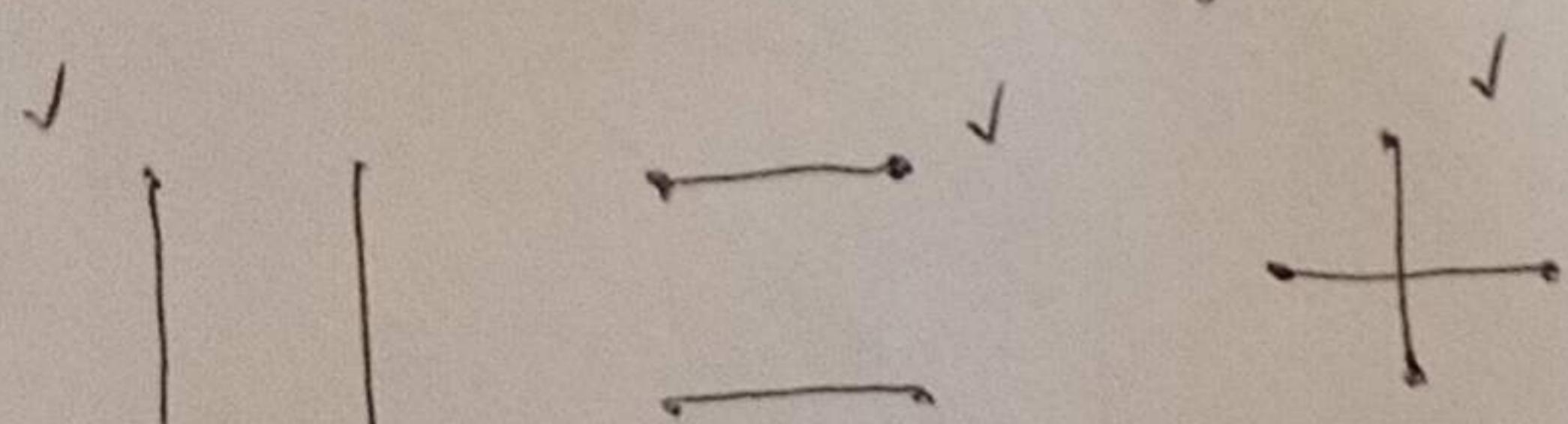
$$0 = \alpha x + b$$

$$x = -\frac{b}{\alpha}$$



intersections removed we solve from two different places bc of 64 with more the hypothesis

I think my drawing of Norms is not correct!

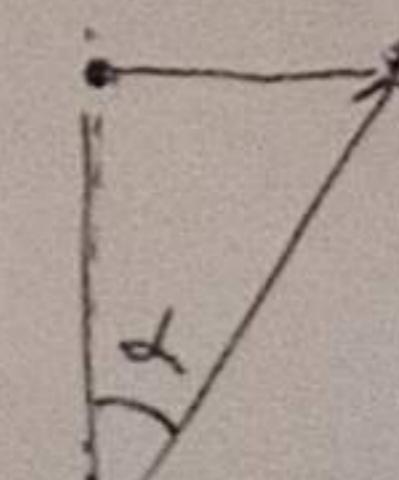
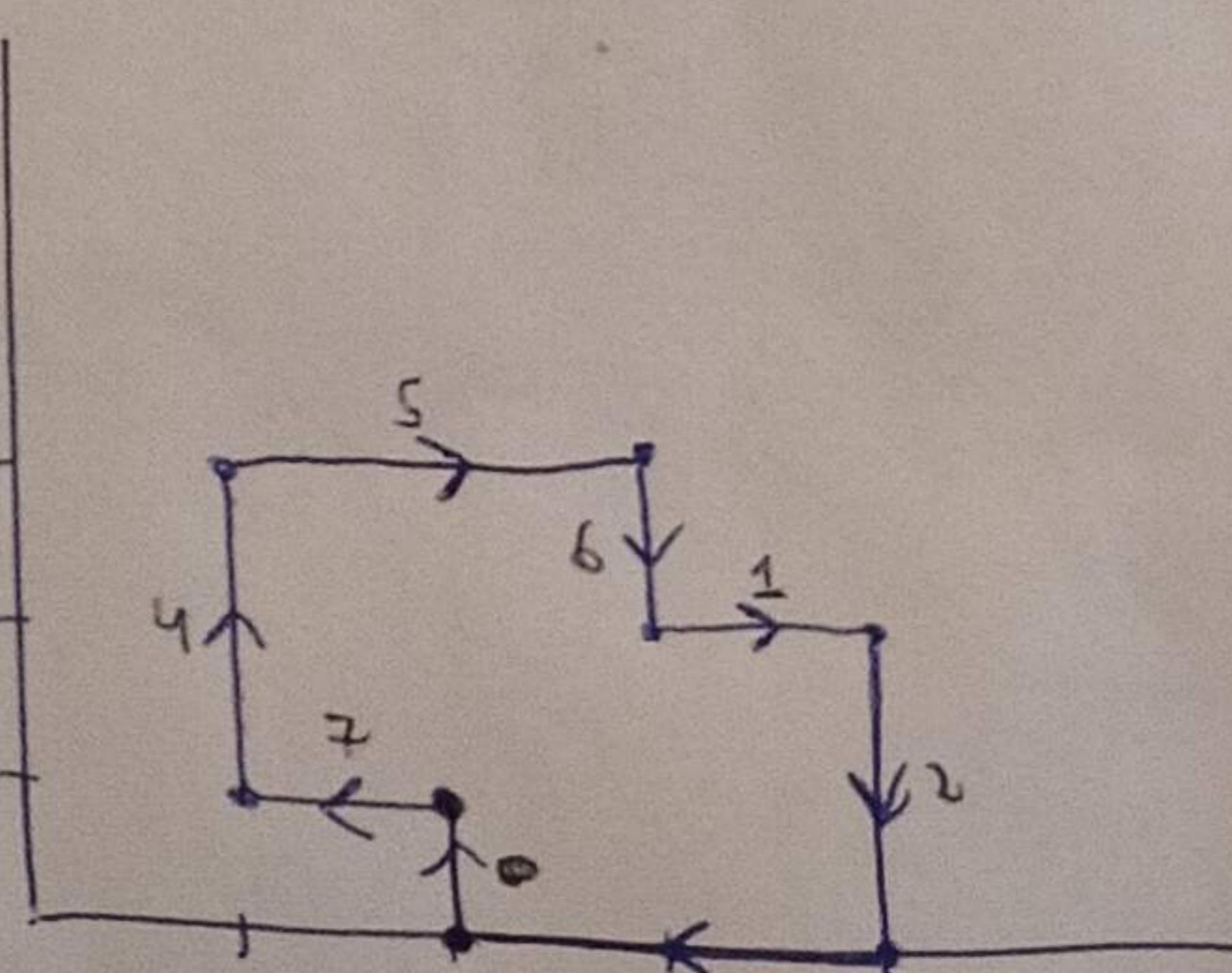
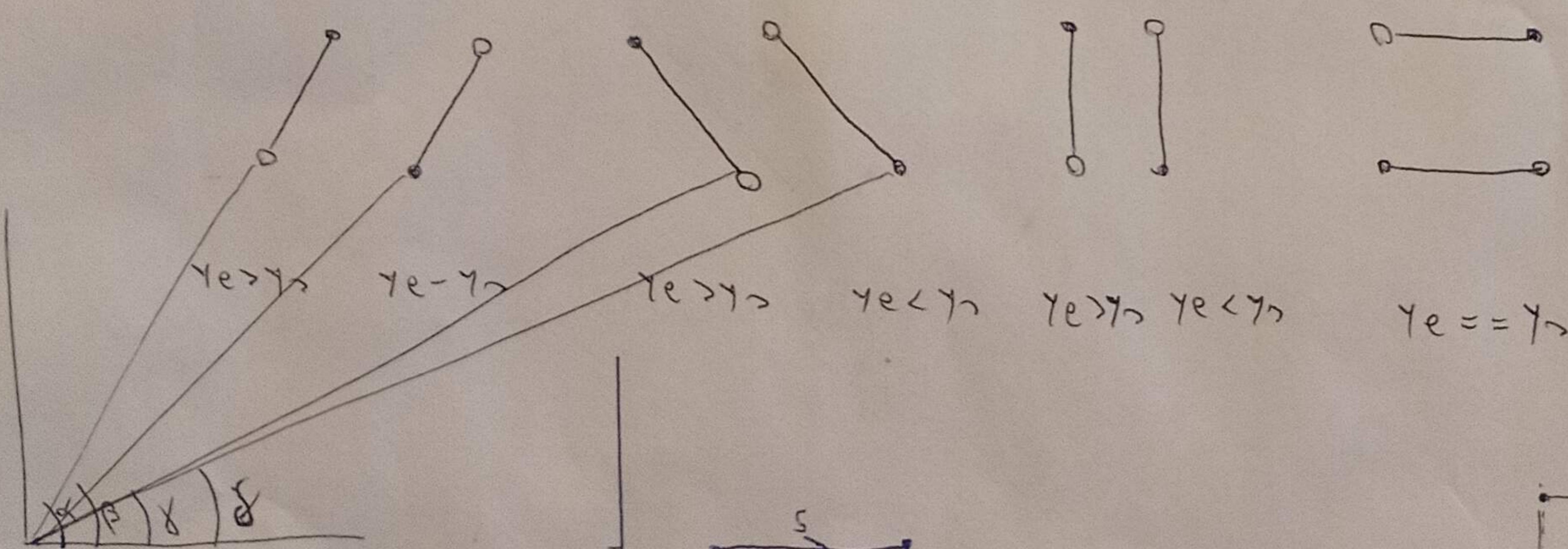


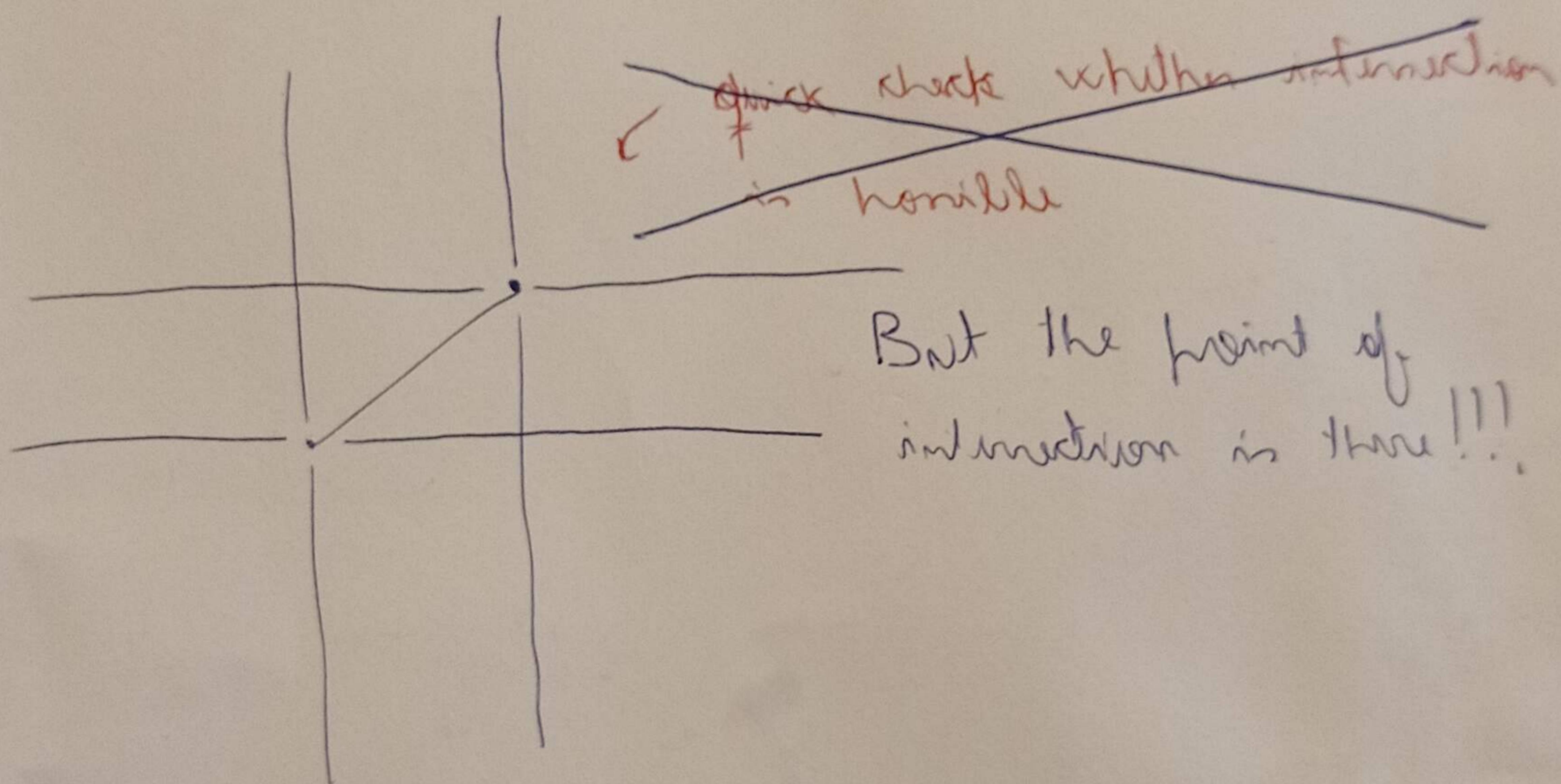
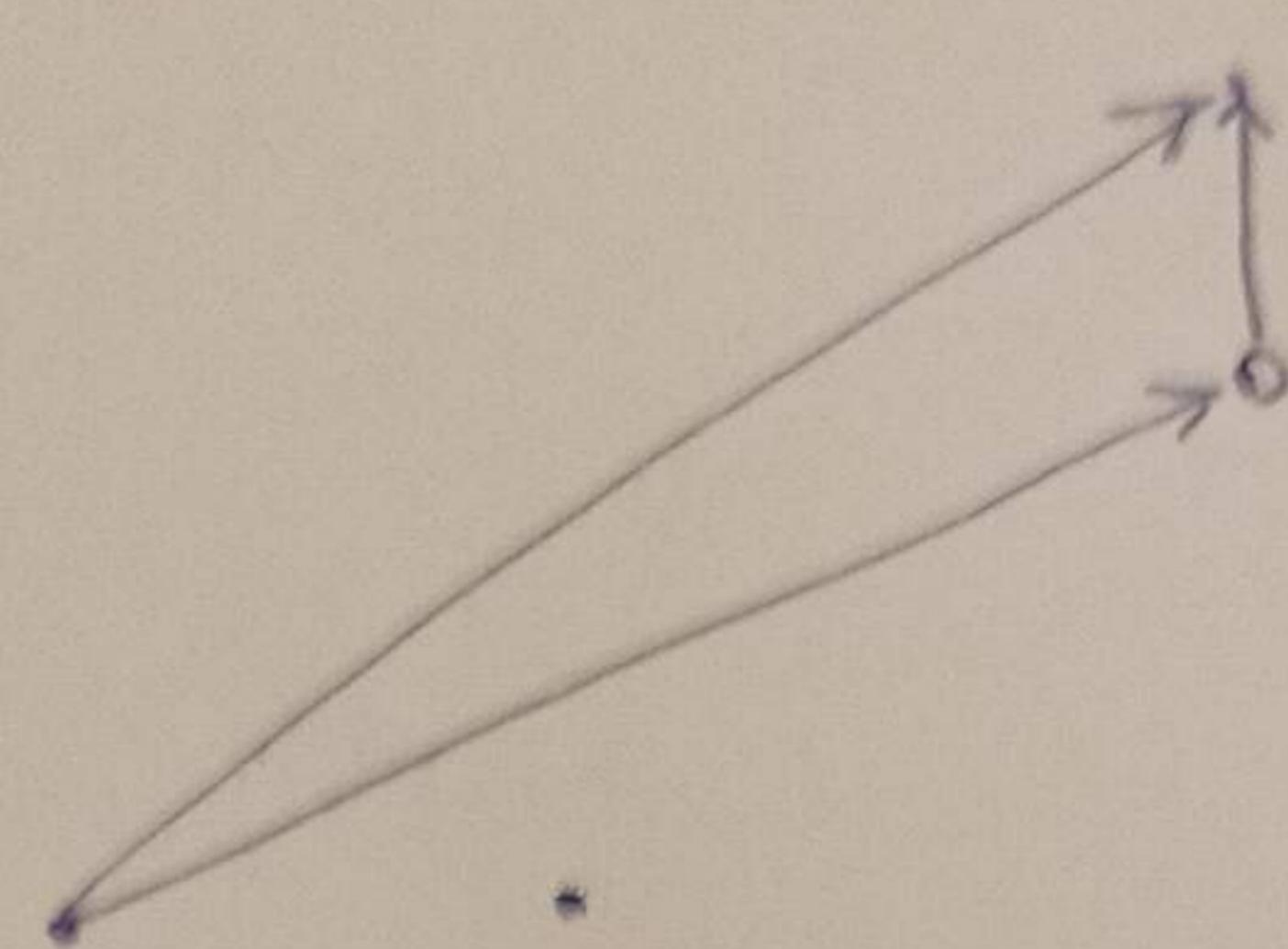
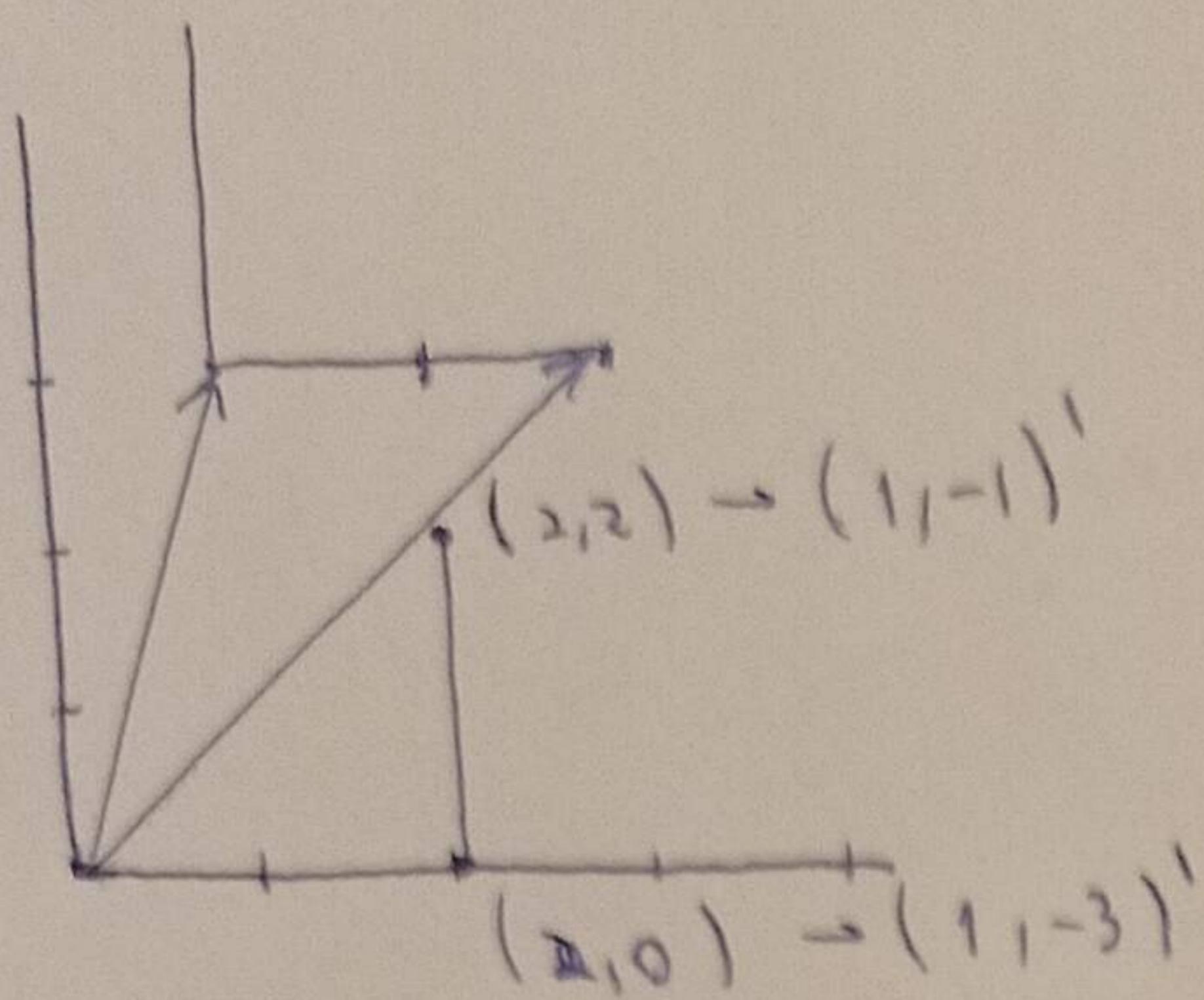
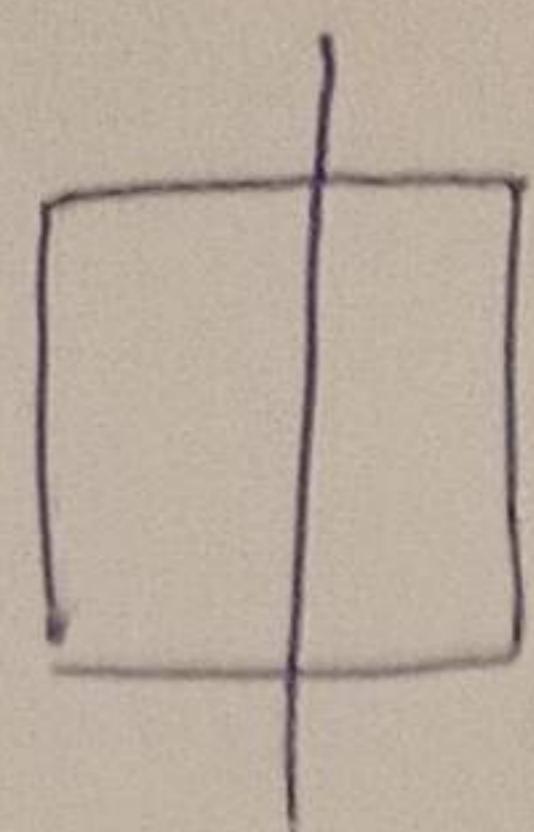
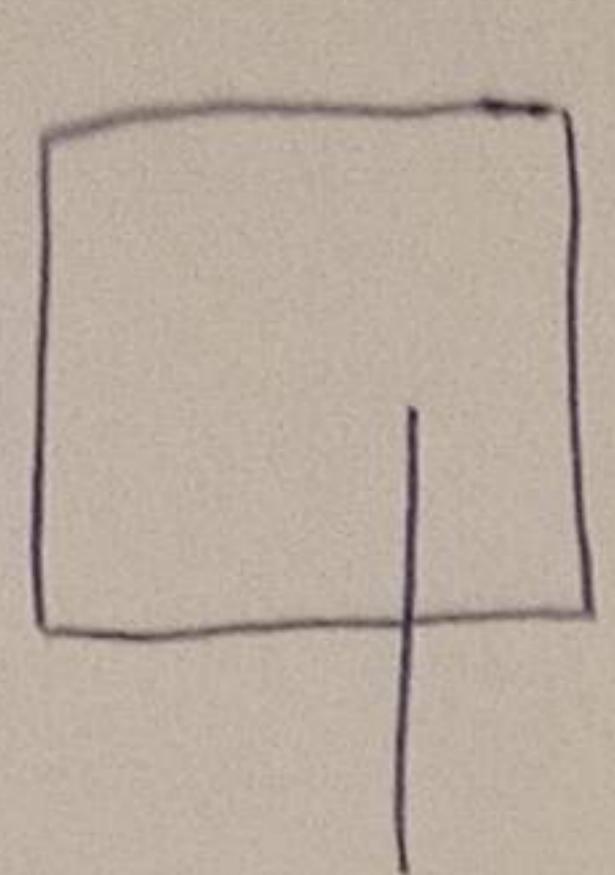
$$\gamma = \alpha x + b$$

$$\gamma = \beta x + d$$

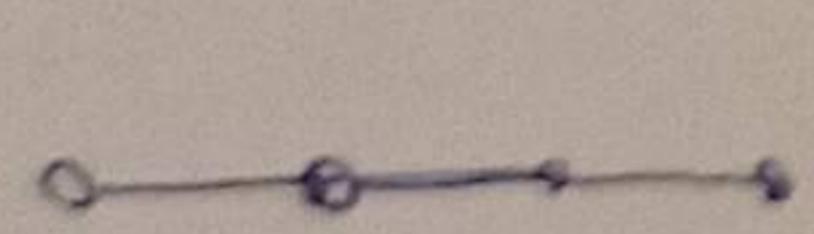
$$0 = b - d$$

if they are
parallel



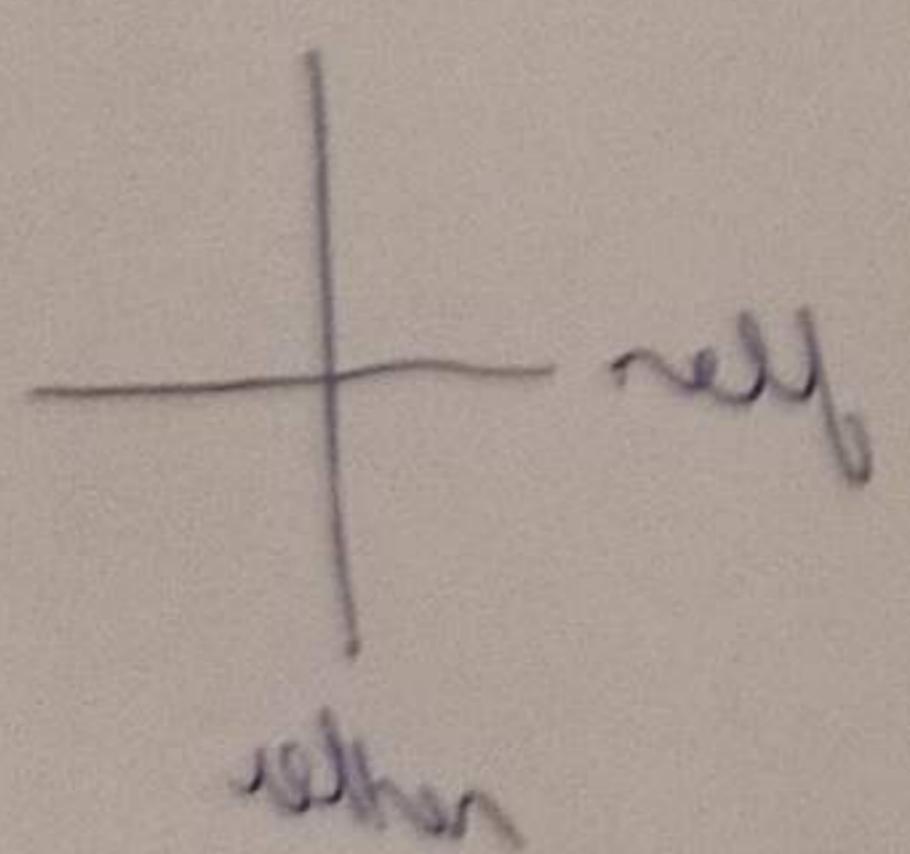


They are either parallel, perpendicular or neither



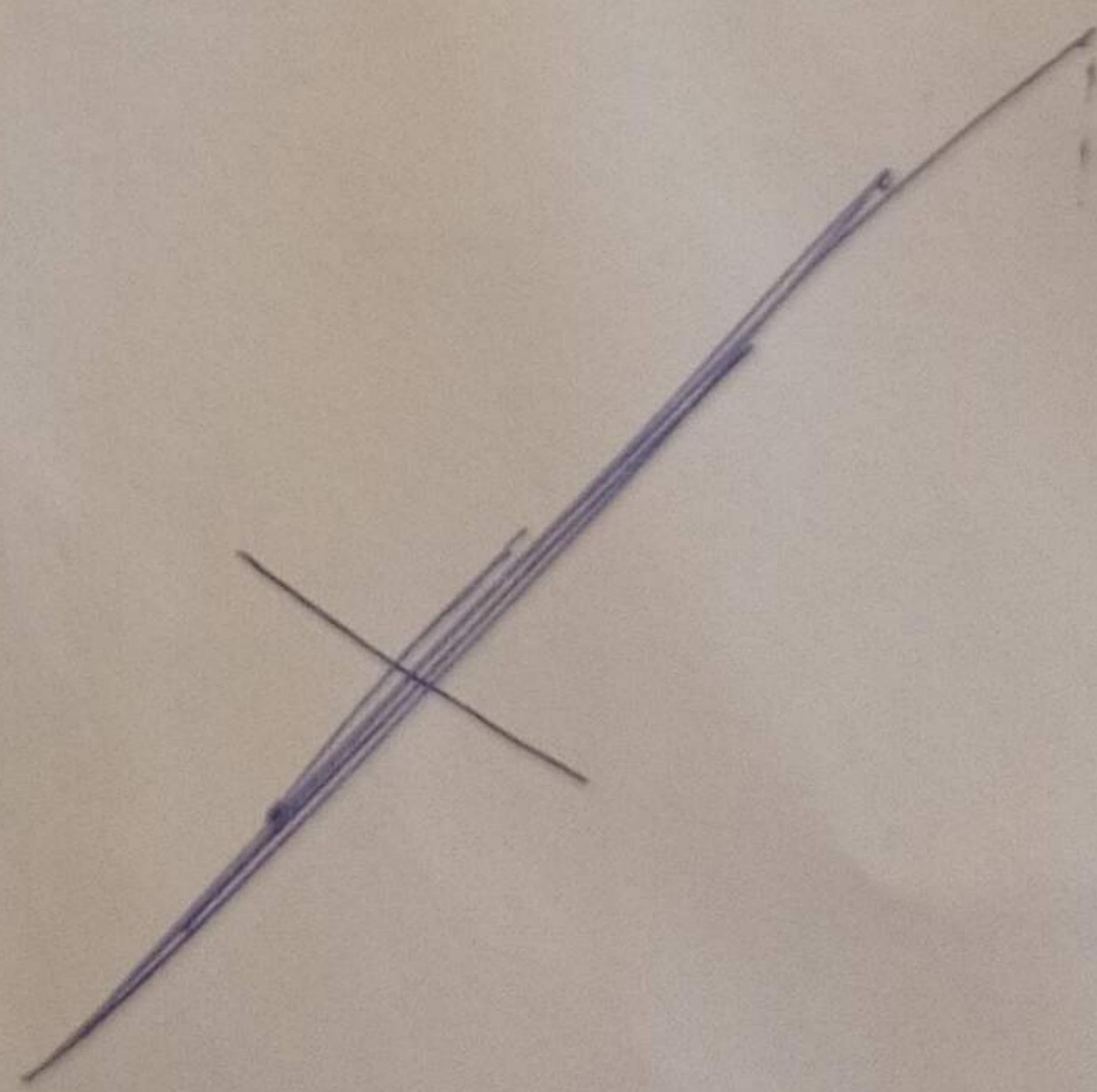
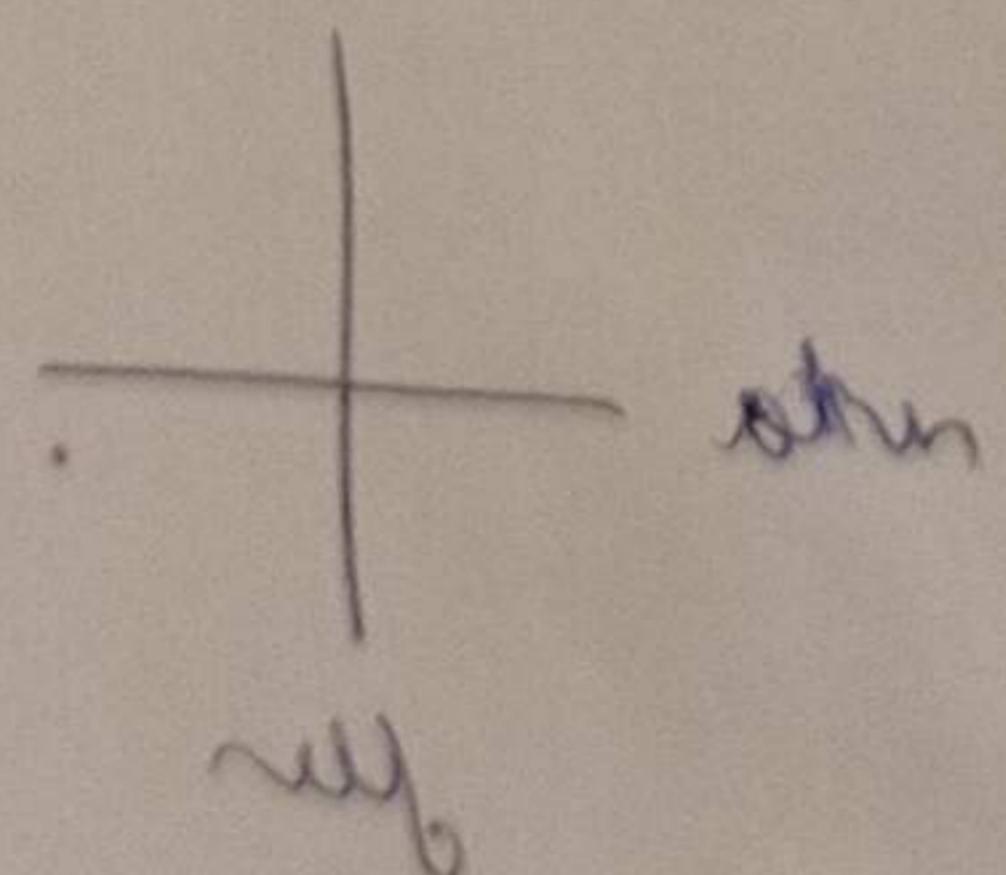
$$y = ax + b$$

$$y = cx + d$$



$$y = \alpha x + b$$

$$y = \beta x + c$$



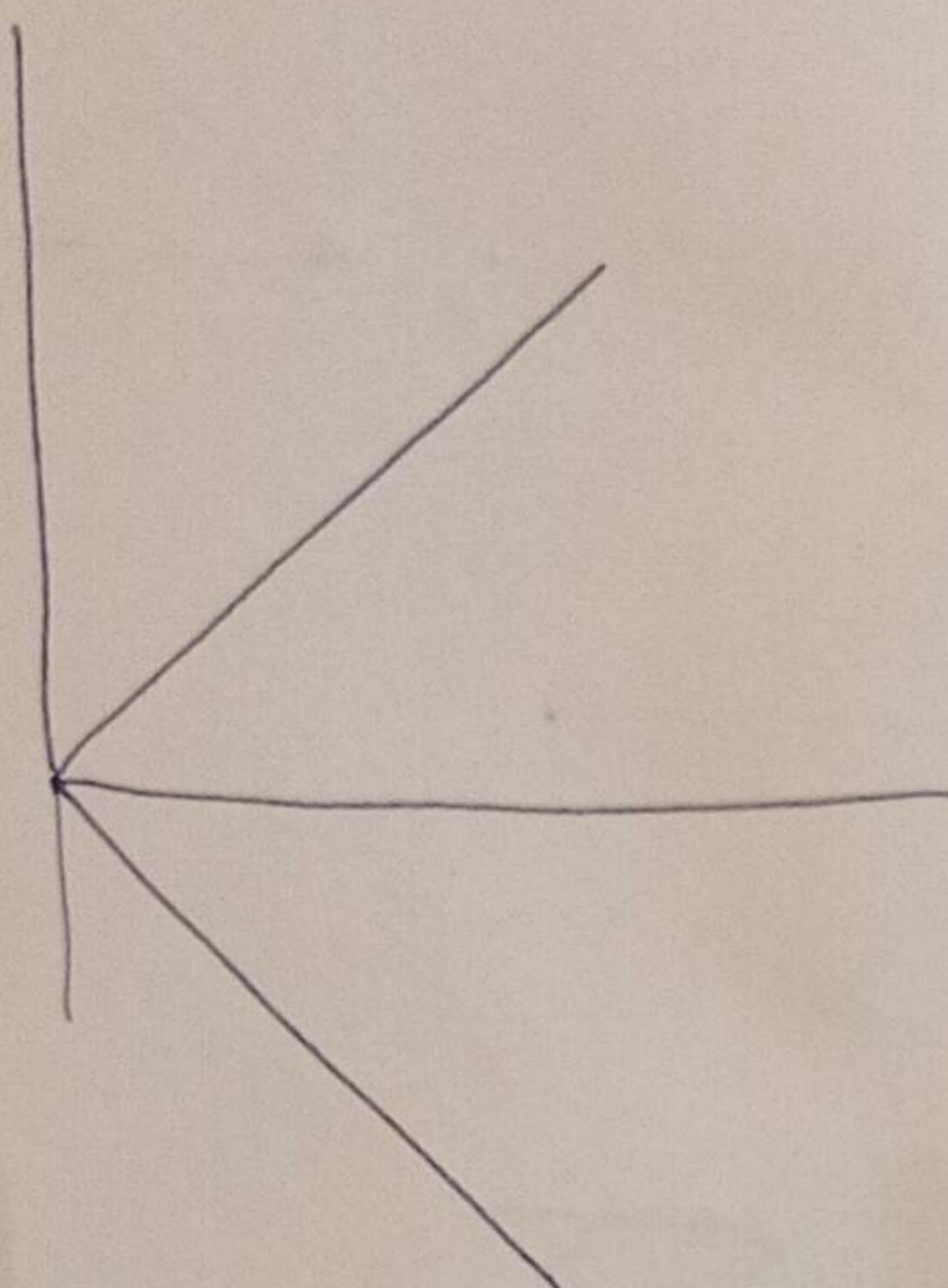
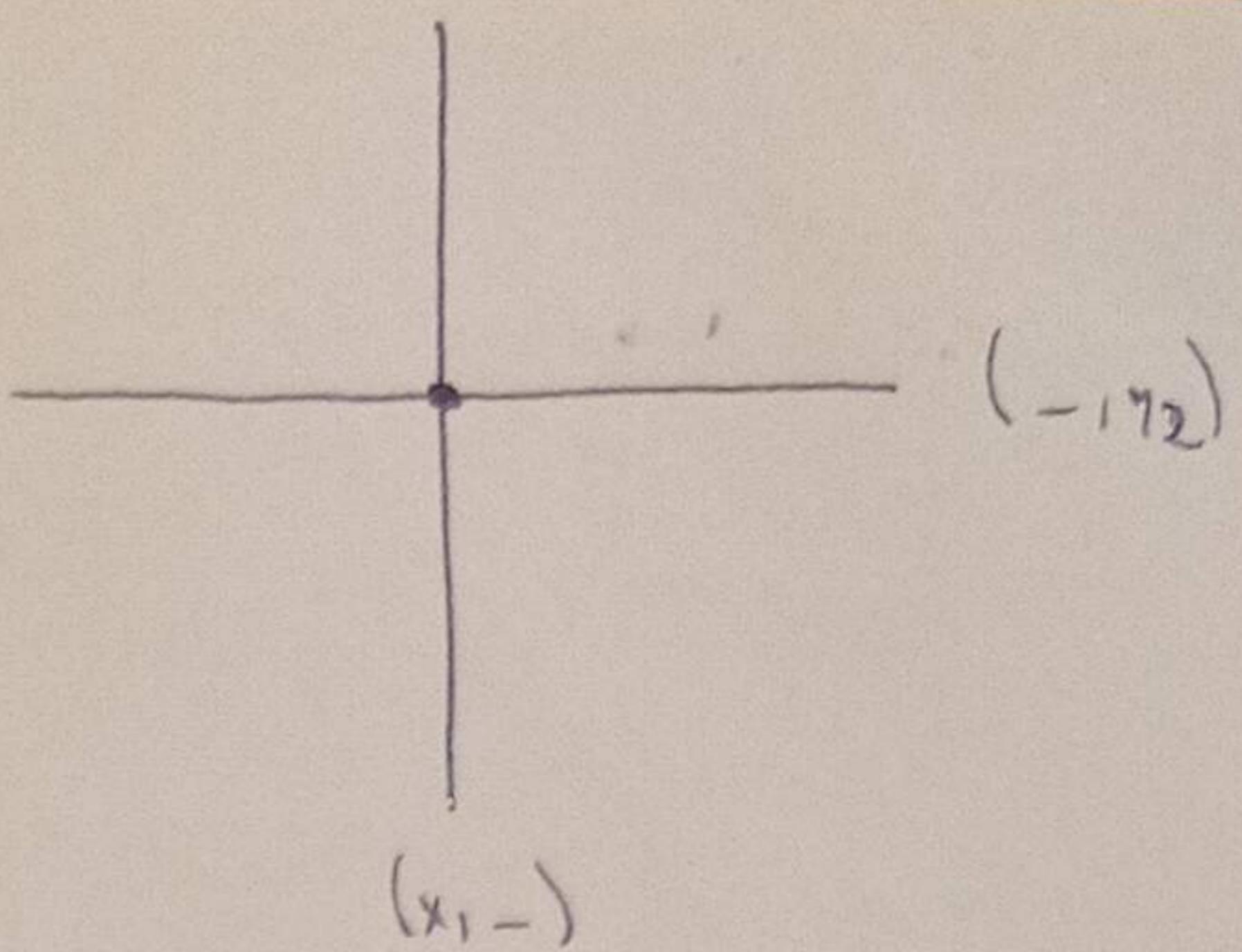
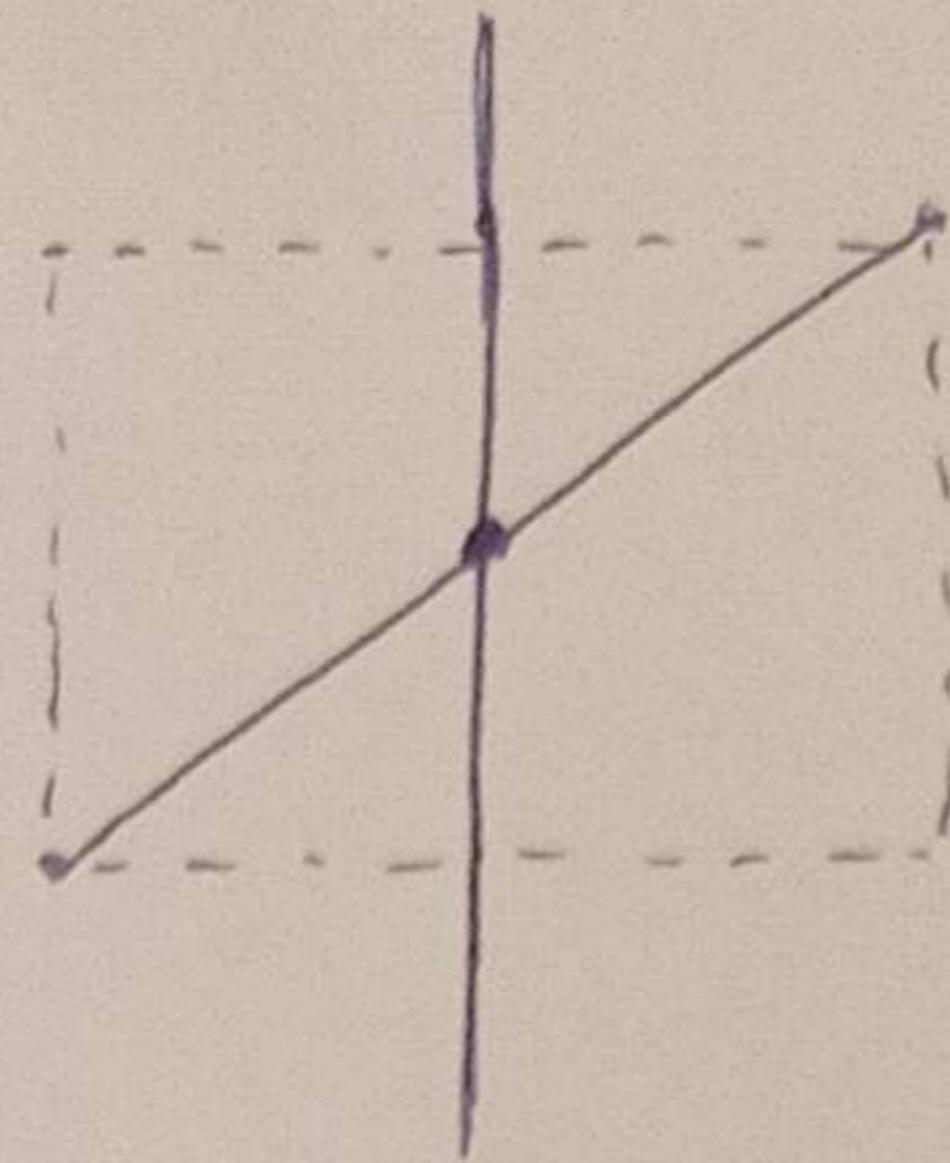
A slant of lines refers the lines with parallelism and perpendicularity

$$\gamma = \infty x + N$$

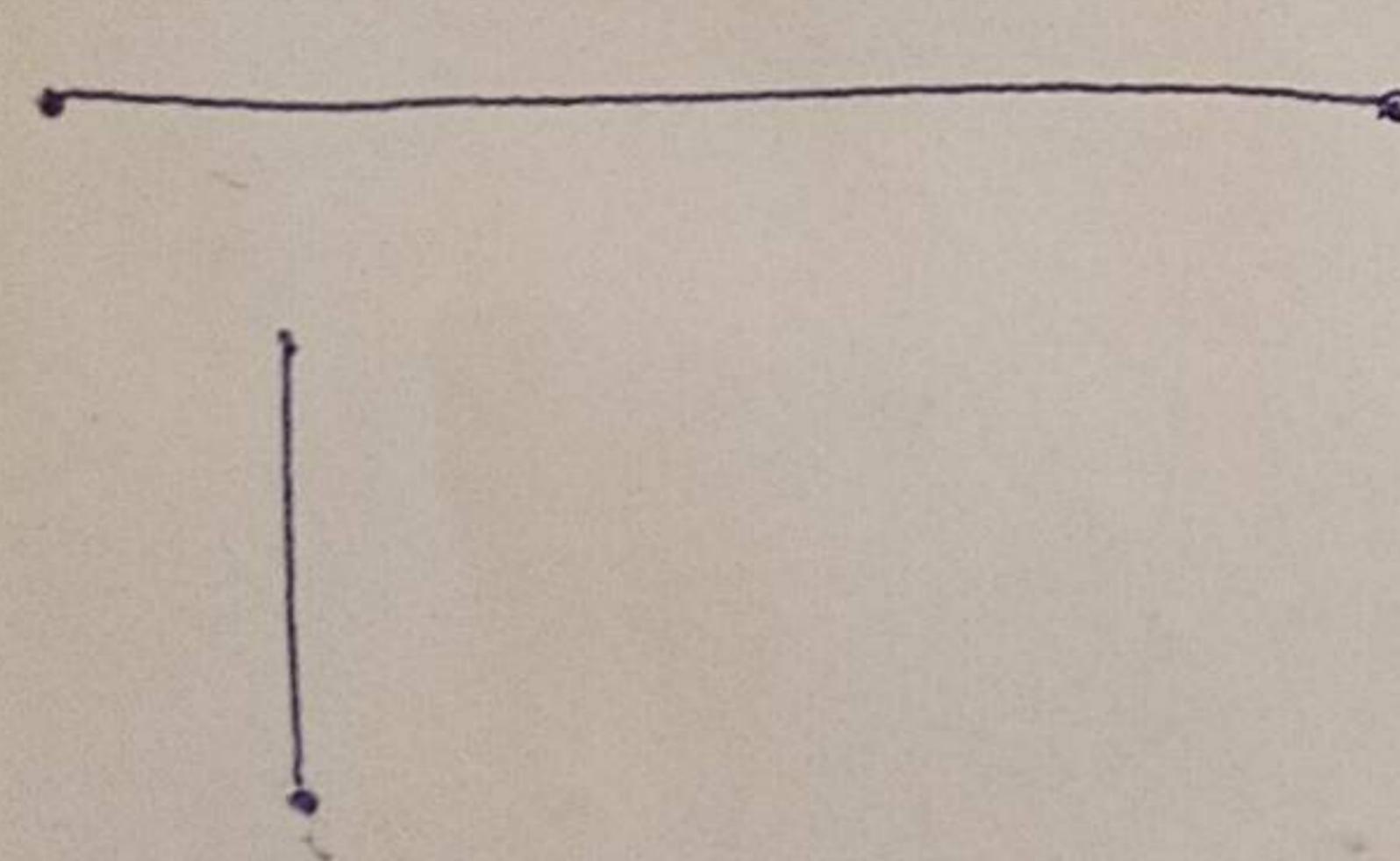
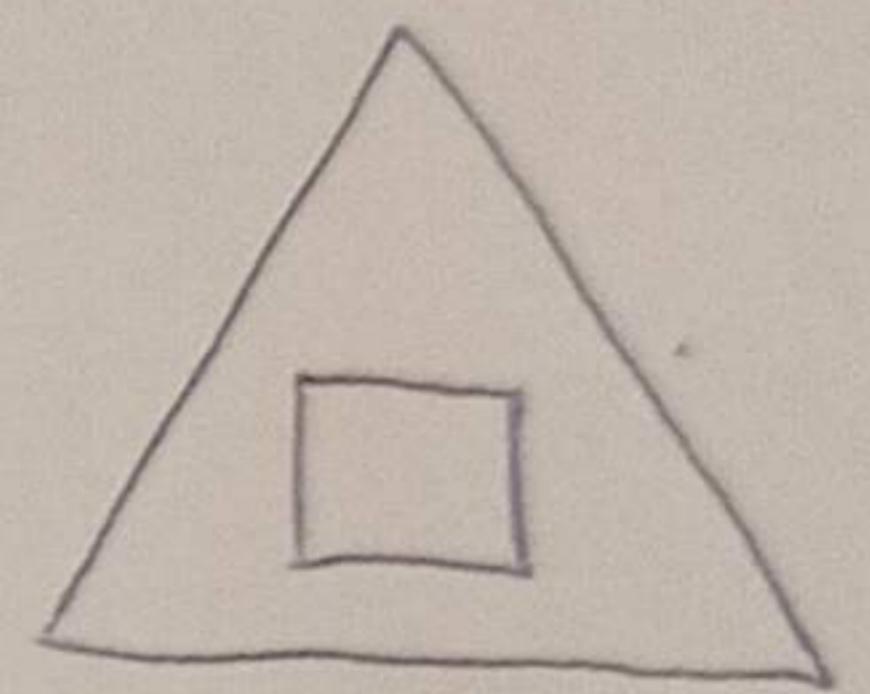
$$\gamma = 0 \cdot x + \gamma_0$$

$$0 = (\infty - \alpha) x + (N - \alpha)$$

undefined



Or say more let's deal with a polygon completely containing a segment



Next move:

a poly might be inside another
poly and be divided, but that doesn't
necessarily hold for a third poly

algo does not distinguish

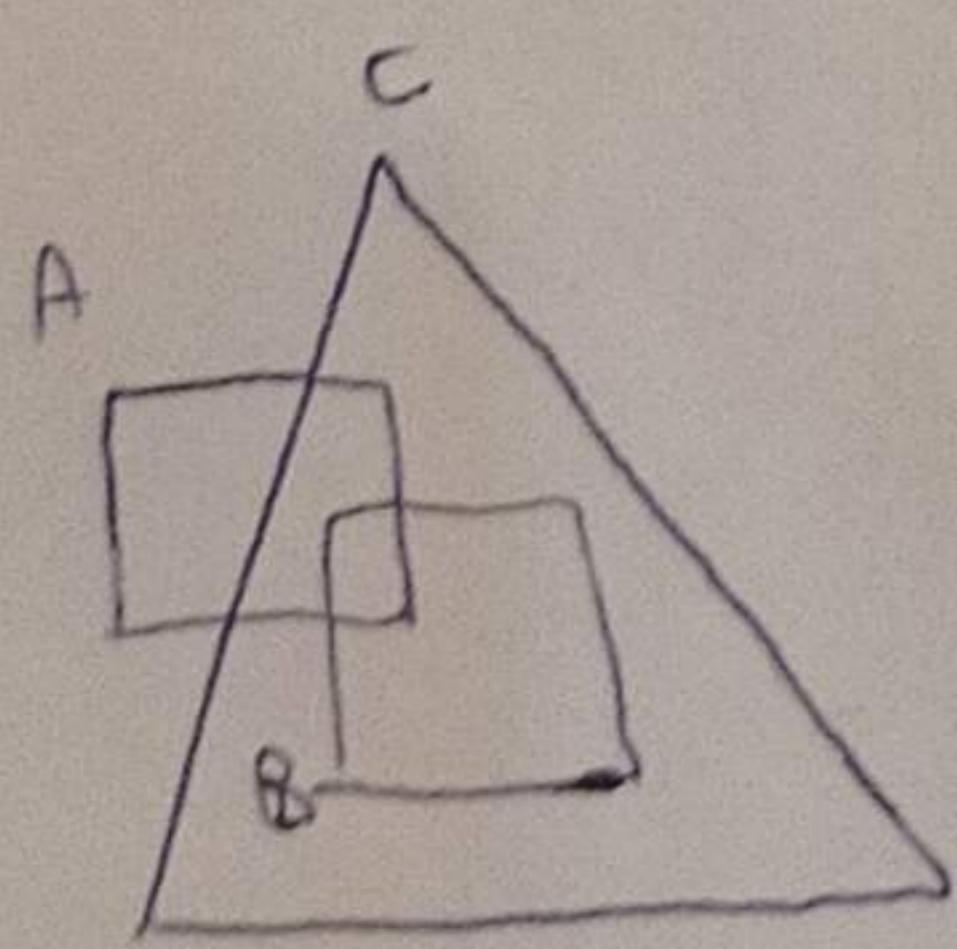
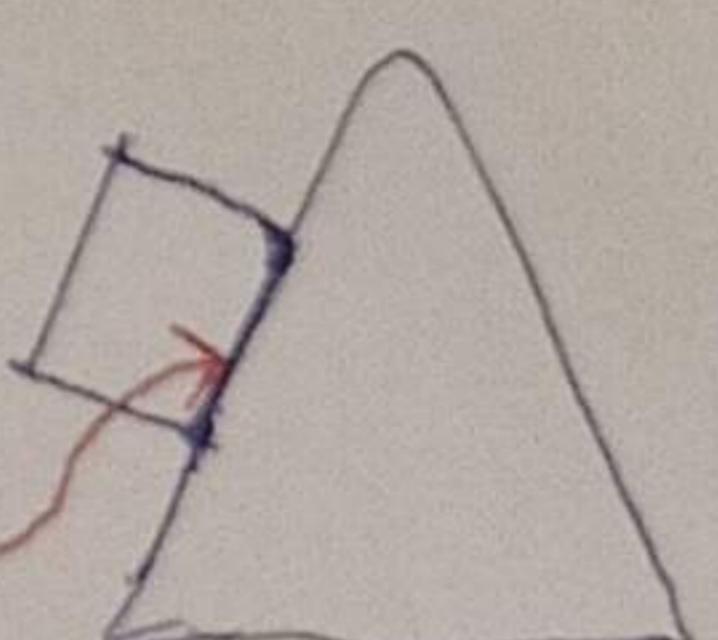
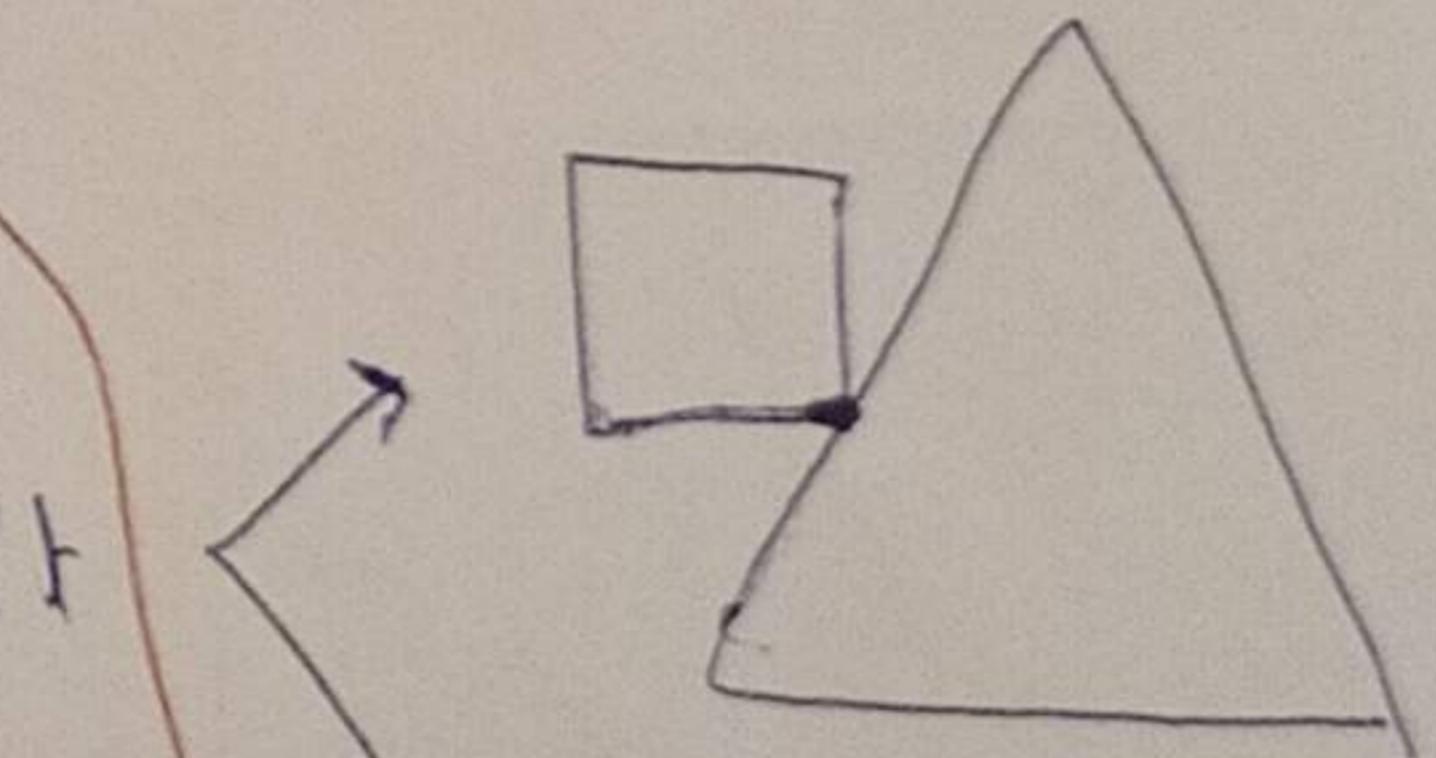
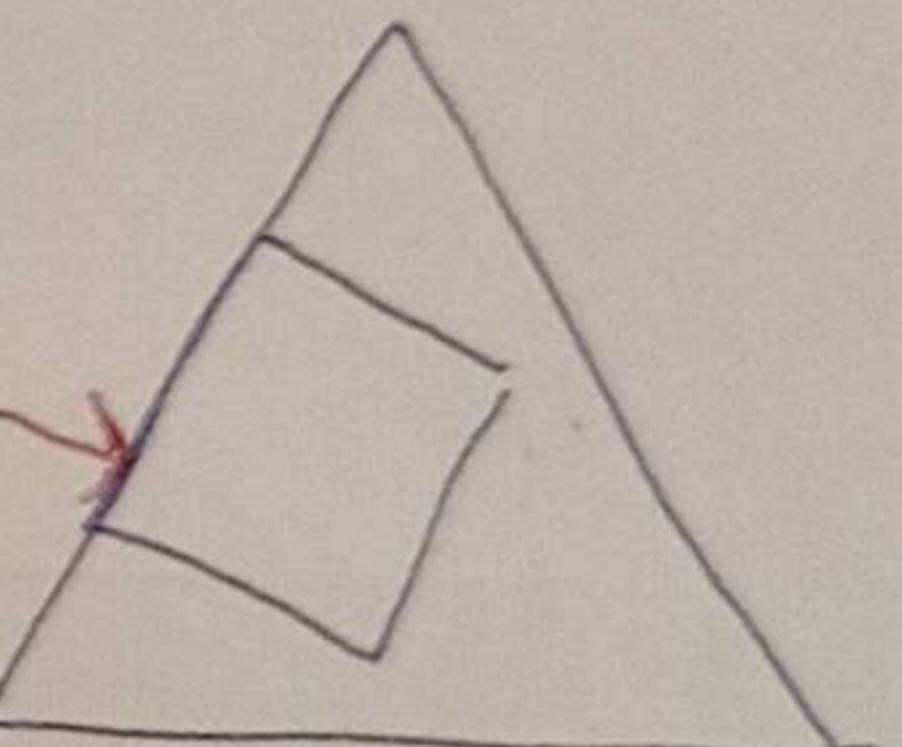
between this segment

and this segment

but these don't

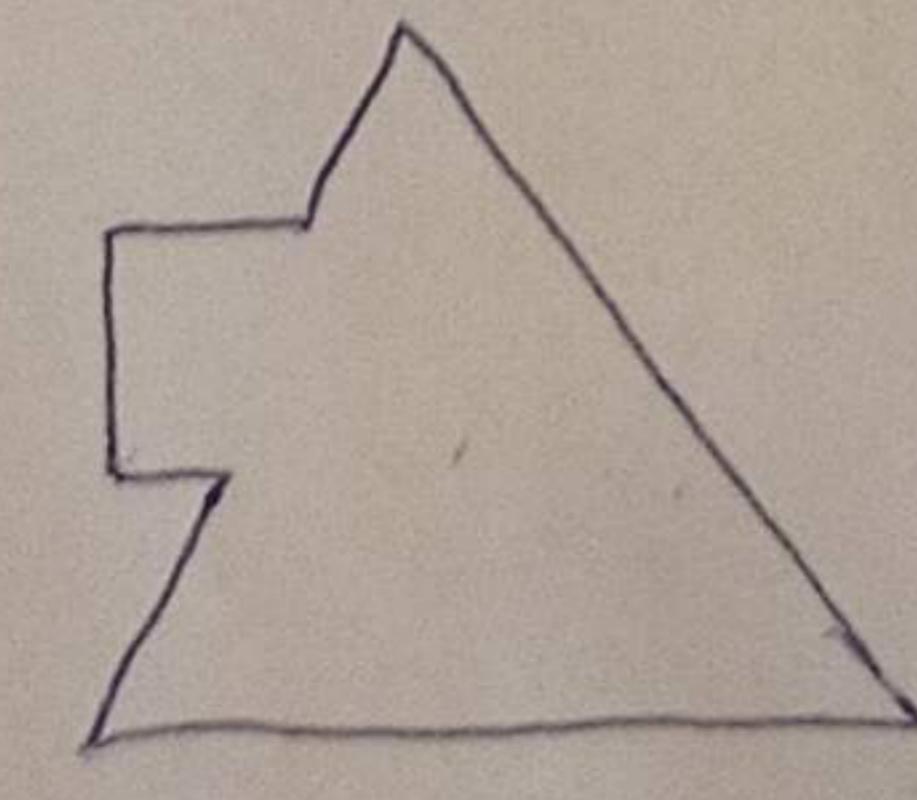
and the latter must

be cut

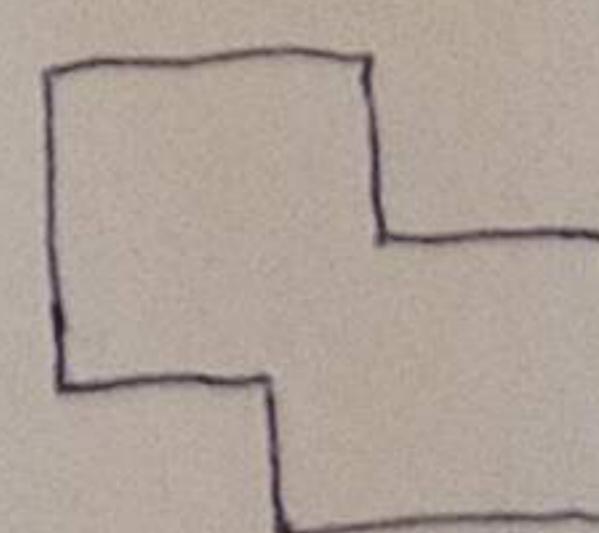


ABC

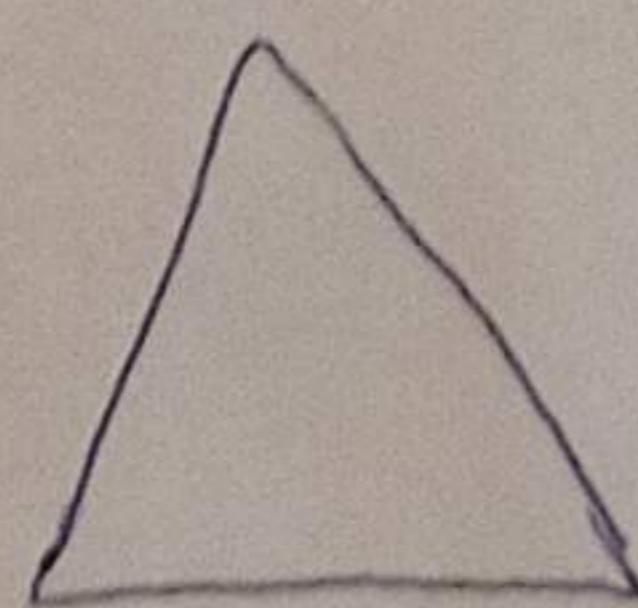
A;C =

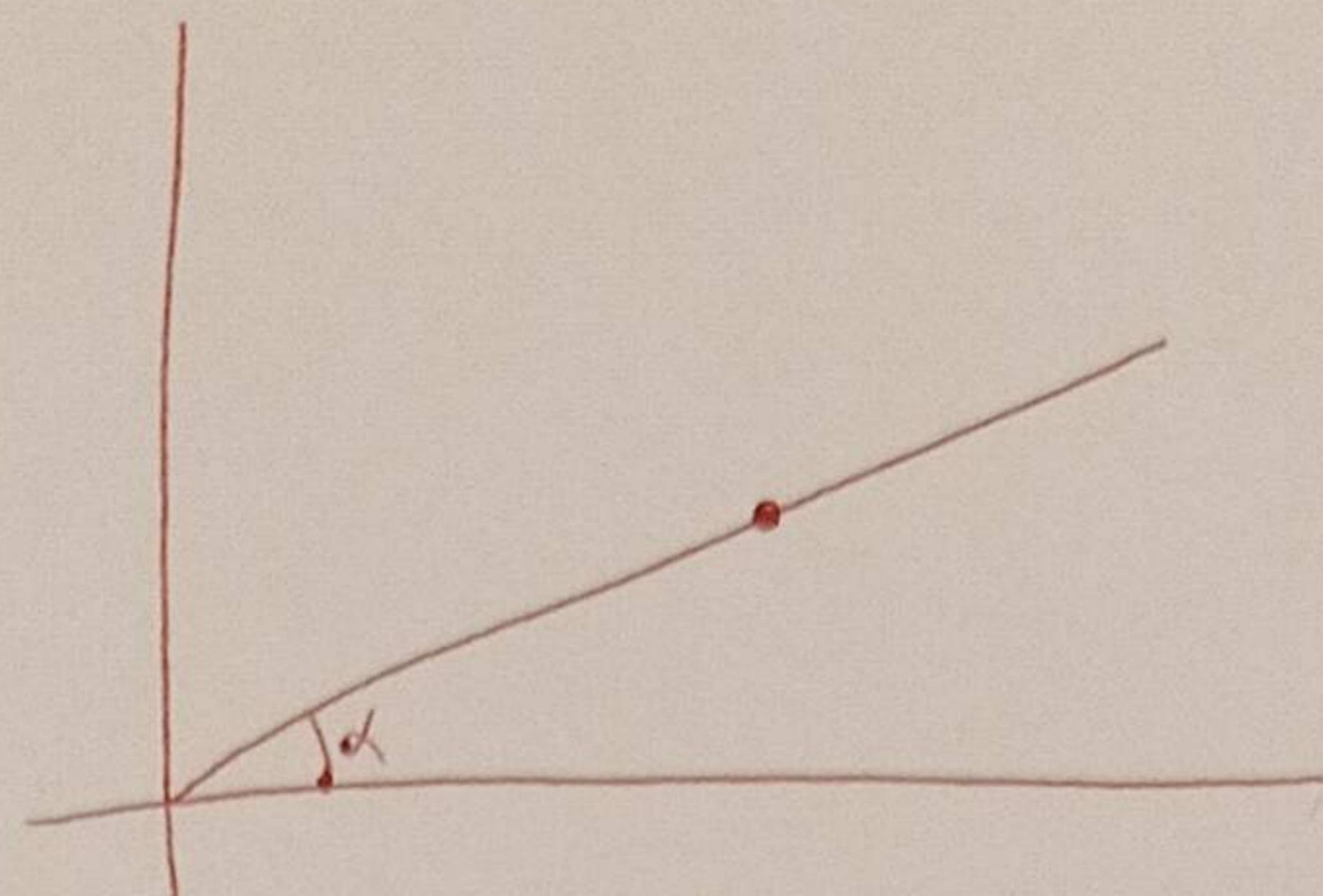
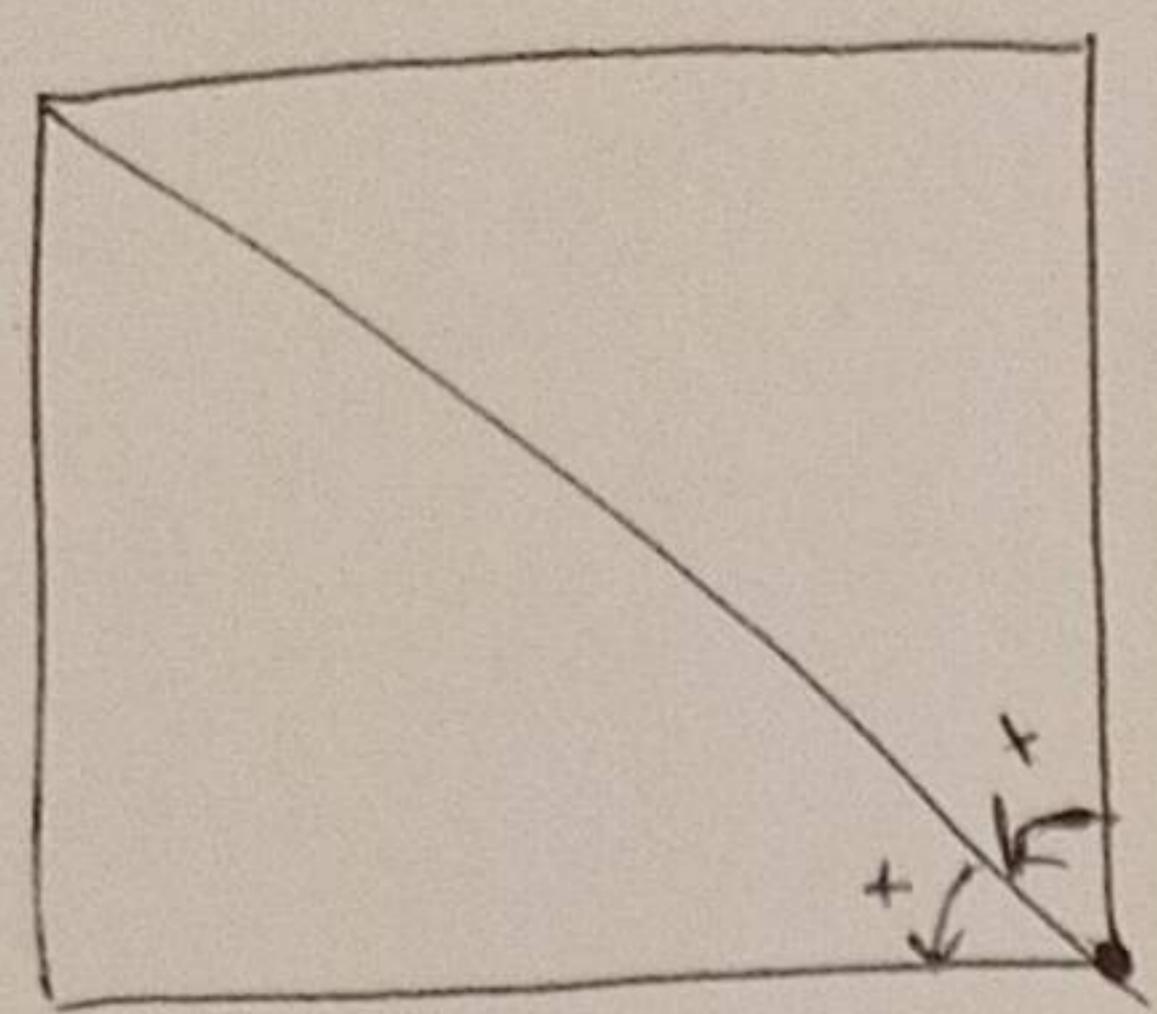
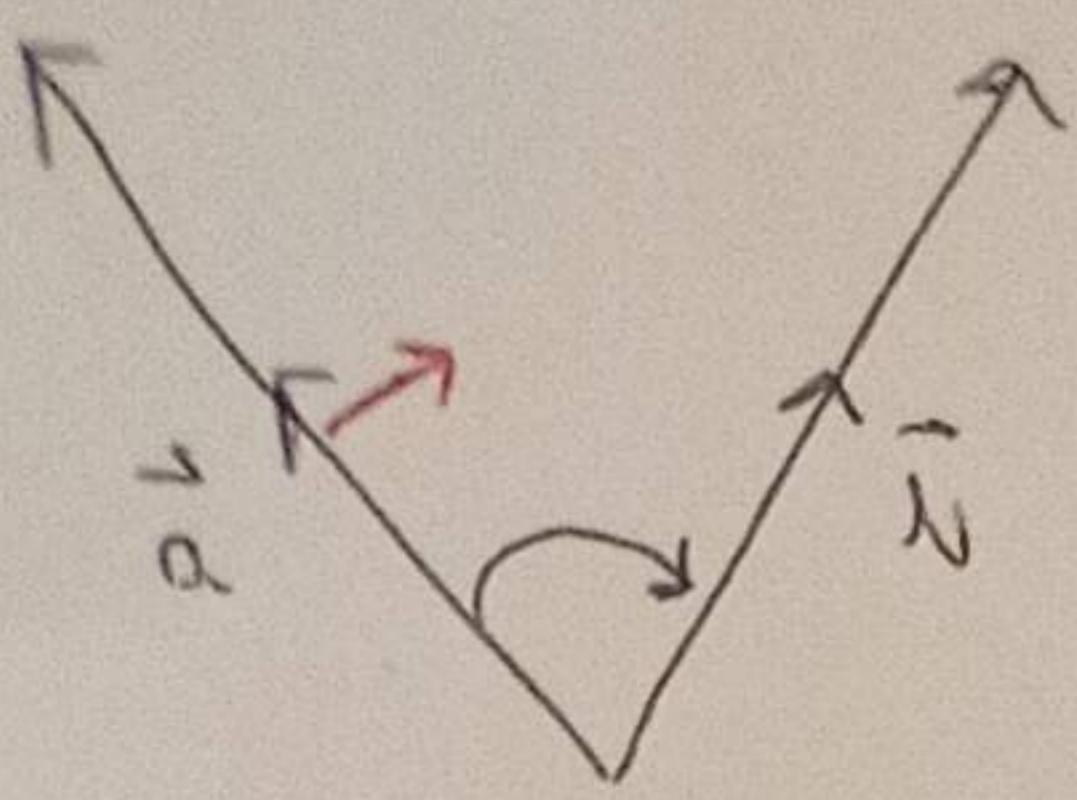


A;B =



B;C =





A B C D E

E D C B A

from C

shifted: [A, B]

forw trait IntroIteration {

type Item;

type IntroIter: Iteration<Item = Self::Item>;

fn intro_iter(Self) → Self::IntroIter;

{

C
D
E
x B
x A

impl (I: Iteration) IntroIteration for I
type Item = I::Item;

type IntroIter = I;

fn intro_iter(Self) → I
Self

{

impl ('a, T) IntroIteration for &'a Vec<T> {

type Item = &'a T;

type IntroIter: slice::Iter<'a, T>;

fn intro_iter(Self) → Self::IntroIter {

Self::iter()

{

{

ABC → CAB → BCA → ABC

A B C
↑ ↑ ↑

C A B
↑ ↑ ↑