# Decentralized Timeline

#### LARGE SCALE DISTRIBUTED SYSTEMS

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### Introduction

Goal — Implement a decentralized timeline service

#### **Technologies**

Language — Python

Asynchronous implementation (Python's asyncio package)

DHT – kademlia (PyPI package)

Pretty print results in tabular form — tabulate (PyPI package)

#### **Bootstrapping**

All nodes are equal implementation-wise

A node can specify the addresses of other nodes it knows of when starting

There may be nodes with widely known addresses and high uptime guarantees

#### **Clock Synchronization**

Relevant for merging timelines of different users

Posts of different users not causally related

Timestamps in timelines and timeline caches handled by the owner of the timeline

Nodes' local time assumed to be reasonably synchronized with an NTP server



## Operations

#### Optional flags in all operations

- -h show help message-d show debug logs-l local port

The **userid** is a pair **ip:port** with the node's ip and port used to receive requests from other nodes. The port can be omitted, with the default being 8000. A one-to-one association between users and nodes is assumed.

#### Start — starts running the user's node

python run.py start <userid>

#### Optional flags:

- -k port used for Kademlia DHT
- -b list of addresses to bootstrap the DHT network
- -f caching frequency
- -t cache time to live
- -c max cached posts per subscription

### Group T3G14

### View — view your feed (posts made by you or by the )

python run.py view [max-posts]

#### Get — find a user's timeline

python run.py get <userid> [max-posts]

#### Post – make a new post

python run.py post <filepath>

#### Remove – delete a post you made

python run.py remove <postid>

#### Sub — subscribe to a user

python run.py sub <userid>

#### Unsub — unsubscribe from a user

python run.py unsub <userid>

#### People I may know – 2nd degree connections

python run.py people-i-may-know [max-users]

### Communication Protocol

#### DISTRIBUTED HASH TABLE

- Used to keep track of user subscriptions
- Implemented with Kademlia DHT
- By default communication happens on port 8468

#### Subscribed to users

```
key - "<userid>-subscribed"
value - [<userid>, ...]
Lists the users that a given user is subscribed to
Only that user will write to this key
```

#### Subscribers of user

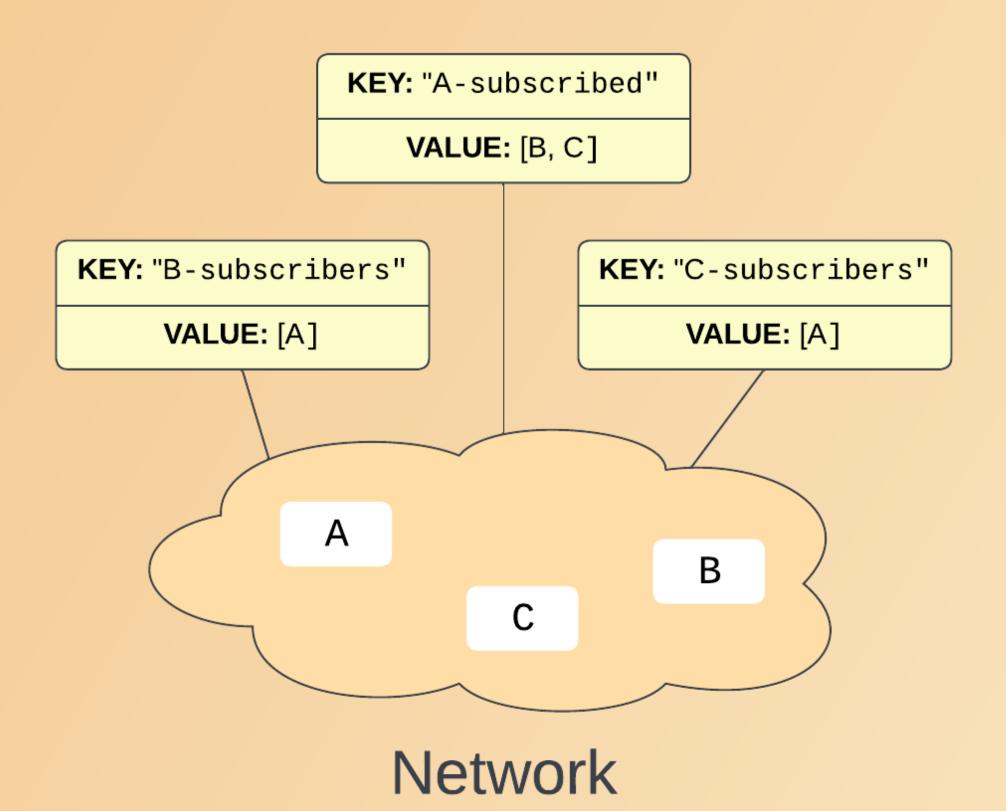


Fig. 1: Subscription information stored in the DHT when user A is subscribed to B and C.



### Communication Protocol

#### REQUESTS TO PEERS

- Used to obtain Timelines of other users
- Requests and responses sent in JSON
- Communication over TCP socket
- By default communication happens on port 8000

#### Request

```
command - "get-timeline"
userid - user who's timeline we are looking for
max-posts - maximum number of posts (optional)
```

#### Response

```
status - "ok" or "error"
error - error message (if status == "error")
timeline - timeline cache (if status == "ok")
```

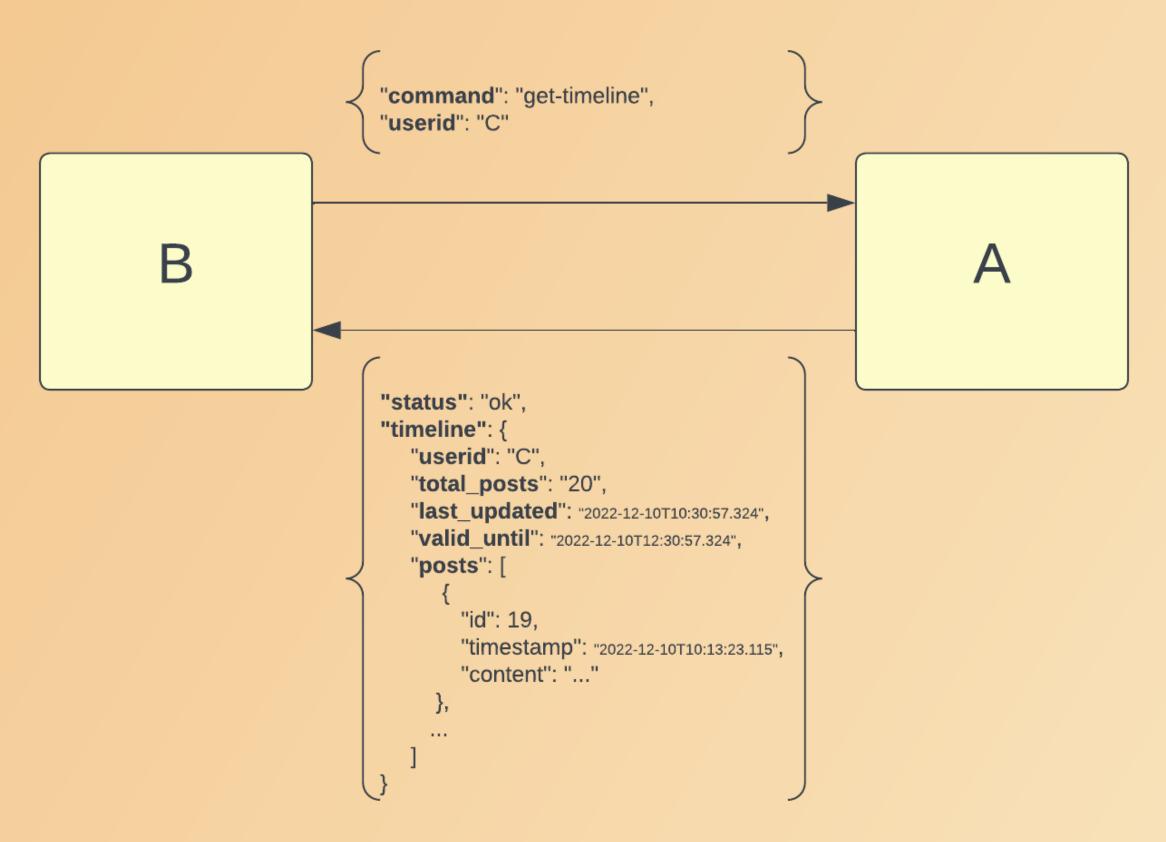


Fig. 2: Example of a request from user B to user A and the respective response.



### Communication Protocol

#### LOCAL CONNECTION

- Used to execute operations in the user's running node
- Requests and responses sent in JSON
- Communication over TCP socket
- By default communication happens on port 8600

#### Request

#### Response

```
status - "ok" or "error"
error - error message (if status == "error")
timeline, users, warnings - depends on the command
```

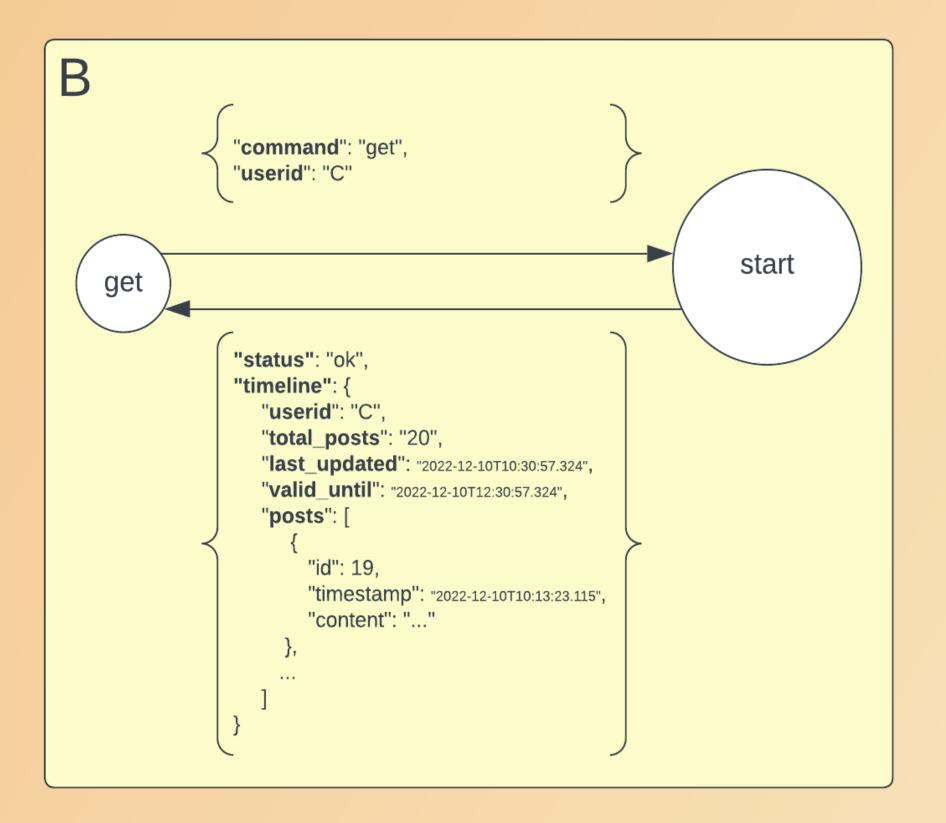


Fig. 3: Example of the communication between the processes "run.py get" and "run.py start" in the computer of the user B in order to execute the **get** command.



# Persistent Storage

- Keeps data stored in disk to withstand crashes
- Stores own timeline, subscribed timeline caches, list of users it is subscribed to and counter for unique post ids

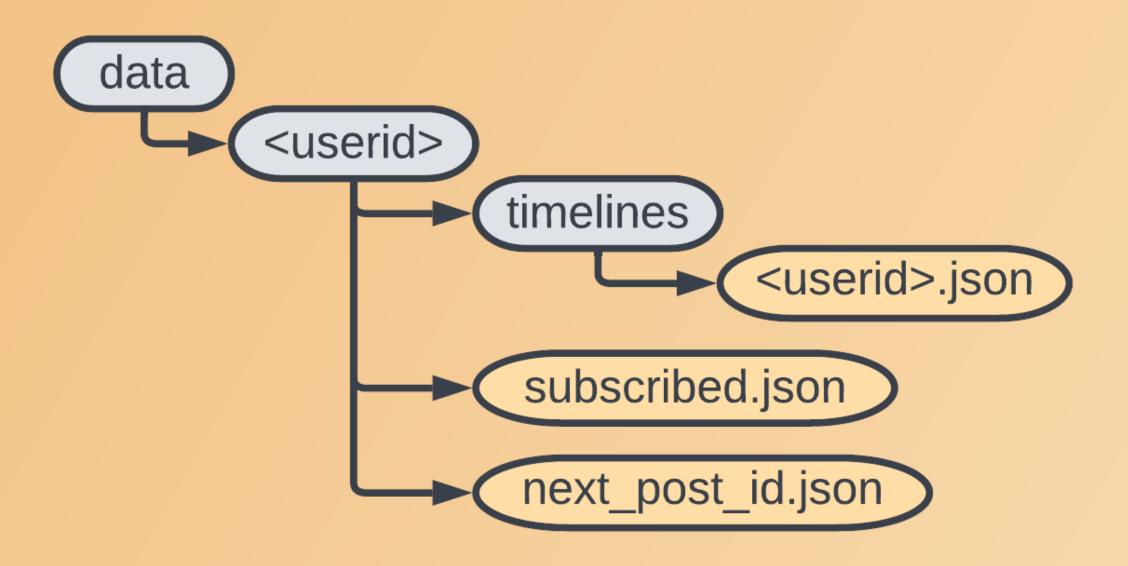


Fig. 4: File structure used to persistently store relevant data.

# Periodic Caching

- Every 120 seconds (this period can be configured), the node will try to update the cache of subscribed timelines
- A cache may become invalid after a period of time (configurable)
- A cache only contains the 20 most recent posts (amount of posts is configurable)

#### For each subscription:

Ensure the user appears as subscribed in the DHT

Tries to retrieve the timeline from the owner

IF FAILS

Tries to retrieve cache from other subscribers

this recovers from race conditions that may have happened

only accepts caches that are more recent than its own



### Retrieve Timelines

#### 1. Locally:

- Local get command
- **get-timeline** command from other node (Note: if a non-subscribed timeline is requested, checks DHT to recover from race conditions that may have happened after unsubscribing)

#### if our timeline was requested

Make cache of owned timeline

OR

#### if other user's timeline was requested

Return cached timeline from local storage (only if it exists and has not expired)

#### 2. From peers:

- Periodic caching
- Local get command (if failed to get locally)

Request timeline directly to owner



Request timeline from another subscriber



IF FETCHED MOST RECENT TIMELINE OR WITH PROBABILITY p ( $p \leftarrow p/2$ )

### Conclusions and Future Work

- The goals of the assignment were accomplished
- Our implementation is robust and we correctly implemented all intended features given the assumptions specified
- The communication between processes using the local TCP socket permits an easy extension with a graphical user interface, which is left as a future work
- Another feature to implement in the future could be protection from malicious writes to the DHT and changes to a timeline cache

