

COMPLEXITY OF US13

kruskal algorithm

```

procedure Kruskal(sortedGraphEdges: array of Edge, verticesGraph: array of
Vertex) returns Graph
    A := new Graph()
    parent := new array[verticesGraph.size()]
    rank := new array[verticesGraph.size()]

    for i := 0 to verticesGraph.size() - 1 do
        parent[i] := i
        rank[i] := 0
    end for

    nA := 0
    i := 0

    while nA < verticesGraph.size() - 1 and i < sortedGraphEdges.size() do
        e := sortedGraphEdges[i]
        u := indexOf(verticesGraph, e.getOrigin())
        v := indexOf(verticesGraph, e.getDestiny())
        rootU := find(u, parent)
        rootV := find(v, parent)

        if rootU != rootV then
            A.addEdge(e)
            union(rootU, rootV, parent, rank)
            nA := nA + 1
        end if

        i := i + 1
    end while

    return A

```

find method

```

procedure find(vertex: integer, parent: array of integer) returns integer
    if parent[vertex] != vertex then
        parent[vertex] := find(parent[vertex], parent)
    end if
    return parent[vertex]

```

addEdge* method**union* method**

```

procedure union(rootU: integer, rootV: integer, parent: array of integer,
rank: array of integer)
    if rank[rootU] > rank[rootV] then
        parent[rootV] := rootU
    else if rank[rootU] < rank[rootV] then
        parent[rootU] := rootV
    else
        parent[rootV] := rootU
        rank[rootU] := rank[rootU] + 1
    end if

```

Lines	Kruskal
1 ^a	1A
2 ^a	1A
3 ^a	1A
4 ^a	(n + 1) A
5 ^a	n * 1A
6 ^a	n * 1A
7 ^a	1A
8 ^a	1A
9 ^a	m * 2C
10 ^a	m * 1A
11 ^a	m * O(n)
12 ^a	m * O(n)
13 ^a	m * T(find) (the complexity of this algorithm is analysed below)
14 ^a	m * T(find) (the complexity of this algorithm is analysed below)
15 ^a	m * 1C
16 ^a	m * 1A
17 ^a	m * T(union) (the complexity of this algorithm is analysed below)
18 ^a	m * (1A + 1Op)
19 ^a	m * (1A m + 1Op)
20 ^a	1R
Total	9 + 2n + 3m
O estimate	O(m log m)

Subtitle:

- A: Assignment

- **n:** Number of Vertices
- **m:** Number of Edges
- **C:** Comparison
- **Op:** Arithmetic Operation
- **I:** Increment
- **R:** Return

COMPLEXITY OF US17

dijkstra algorithm

```

procedure Dijkstra(start: String, MP: String, edges: array of Edges)
    shortestPaths := new Graph()
    if MP is equal to start then
        print("The vertex is already te Metting Point!")
    else
        initialGraph := new Graph()
        for each Edge in the set of Edges do
            Add Edge to initialGraph
        end for
        vertices := get vertices from inicjalGraph
        numVertices := vertices.size()
        dist := new array[numVertices]
        prev := new array[numVertices]
        visited := new array[numVertices]
        for each index i from 0 to numVertices - 1 do
            dist[i] := -1
            prev[i] := null
            visited[i] := false
        create a PriorityQueue named Queue with a comparator based on
        the distance of vertices
        MPVertex := new Vertex(MP)
        startVertex := new Vertex(start)
        dist[vertices.indexOf(startVertex)] := 0
        queue.add(startVertex)
        while (not queue.isEmpty())
            u := queue.poll()
            uIndex := vertices.indexOf(u)
            visited[uIndex] := true

```

```

neighbors := initialGraph.getVerticesConnectedTo(u)
for each neighbor in neighbors
    vIndex := vertices.indexOf(neighbor)
    weight := initialGraph.getEdgeCost(u, neighbor)

    if (not visited[vIndex] and dist[uIndex] ≠ -1 and
    (dist[vIndex] = -1 or dist[uIndex] + weight < dist[vIndex])) then
        queue.remove(neighbor)
        dist[vIndex] := dist[uIndex] + weight
        prev[vIndex] := u
        queue.add(neighbor)
    end if
end for
end while
path := new ArrayList()
at := MPVertex
while (at ≠ null)
    path.add(at)
    at := prev[vertices.indexOf(at)]
end while
path.reverse()
i := 0
while (i < path.size() - 1)
    origin := path.get(i)
    destiny := path.get(i + 1)
    cost := initialGraph.getEdgeCost(origin, destiny)
    shortestPaths.addEdge(new Edge(origin, destiny, cost))
    i := i + 1
end while
return shortestPaths

```

Lines	Dijkstra
1 ^a	1A
2 ^a	1C
3 ^a	1A
4 ^a	1A
5 ^a	1A
6 ^a	$m * 1A$
7 ^a	$m * 1A$
8 ^a	$n * 1A$
9 ^a	1A
10 ^a	$n * 1A$
11 ^a	$n * 1A$
12 ^a	$n * 1A$
13 ^a	$n * 1C$
14 ^a	$n * 1A$
15 ^a	$n * 1A$
16 ^a	$n * 1A$
17 ^a	1A
18 ^a	1A
19 ^a	1A
20 ^a	1A
21 ^a	$1A + \log(n)$
22 ^a	$n * \log(n)$
23 ^a	$n * \log(n)$
24 ^a	$n * 1A$
25 ^a	$n * 1A$
26 ^a	$n * 1A$
27 ^a	$m * 1C$
28 ^a	$m * 1A$
29 ^a	$m * 1A$
30 ^a	$m * 3C$
31 ^a	$m * \log(n)$
32 ^a	$m * 1A$
33 ^a	$m * 1A$
34 ^a	$m * \log(n)$
35 ^a	$m * 1A$
36 ^a	$m * 1A$
37 ^a	$n * \log(n)$
38 ^a	1A
39 ^a	1A
40 ^a	$n * 1C$
41 ^a	$n * 1A$
42 ^a	$n * 1A$
43 ^a	$n * 1A$

44 ^a	n * 1A
45 ^a	1A
46 ^a	n * 1C
47 ^a	n * 1A
48 ^a	n * 1A
49 ^a	n * 1A
50 ^a	n * 1A
51 ^a	n * 1A
52 ^a	1A
53 ^a	1A
54 ^a	1R
Total	$9+2n+3m+2m\log(n)+6n+5m+5n+1$
O estimate	$O(m \log(n))$

Subtitle:

- **A:** Assignment
- **n:** Number of Vertices
- **m:** Number of Edges
- **C:** Comparison
- **R:** Return

COMPLEXITY OF US18

```

neighbors := initialGraph.getVerticesConnectedTo(u)

for each neighbor in neighbors do

    vIndex := vertices.indexOf(neighbor)

    weight := initialGraph.getEdgeCost(u, neighbor)

    if not visited[vIndex] and dist[uIndex] != -1 and
(dist[vIndex] = -1 or dist[uIndex] + weight < dist[vIndex]) then

        queue.remove(neighbor)

        dist[vIndex] := dist[uIndex] + weight

        prev[vIndex] := u

        queue.add(neighbor)

    end if

end for

end while

path := new ArrayList<>()

for Vertex at := MPVertex; at != null; at :=
prev[vertices.indexOf(at)] do

    path.add(at)

end for

path.reverse()

shortestPath := new Graph()

for each index i from 0 to path.size() - 2 do

    Vertex origin := path.get(i)

    Vertex destiny := path.get(i + 1)

    int cost := initialGraph.getEdgeCost(origin, destiny)

    shortestPath.addEdge(new Edge(origin, destiny, cost))

end for

shortestPaths.add(shortestPath)

end if

end for

return shortestPaths

```

Lines	Dijkstra
1 ^a	1A
2 ^a	k * 1C
3 ^a	k * 1C
4 ^a	k * 1A
5 ^a	k * 1A
6 ^a	k * 1A
7 ^a	k * m * 1A
8 ^a	k * m * 1A
9 ^a	k * n * 1A
10 ^a	k * 1A
11 ^a	k * n * 1A
12 ^a	k * n * 1A
13 ^a	k * n * 1A
14 ^a	k * n * 1C
15 ^a	k * n * 1A
16 ^a	k * n * 1A
17 ^a	k * n * 1A
18 ^a	k * 1A
19 ^a	k * 1A
20 ^a	k * 1A
21 ^a	k * 1A
22 ^a	k * (1A + log(n))
23 ^a	k * n * log(n)
24 ^a	k * n * log(n)
25 ^a	k * n * 1A
26 ^a	k * n * 1A
27 ^a	k * n * 1A
28 ^a	k * m * 1C
29 ^a	k * m * 1A
30 ^a	k * m * 1A
31 ^a	k * n * 3C
32 ^a	k * m * log(n)
33 ^a	k * m * 1A
34 ^a	k * m * 1A
35 ^a	k * m * log(n)
36 ^a	k * m * 1A
37 ^a	k * m * 1A
38 ^a	k * n * log(n)
39 ^a	k * 1A
40 ^a	k * 1A

41 ^a	k * n * 1C
42 ^a	k * n * 1A
43 ^a	k * n * 1A
44 ^a	k * n * 1A
45 ^a	k * n * 1A
46 ^a	k * 1A
47 ^a	k * n * 1C
48 ^a	k * n * 1A
49 ^a	k * n * 1A
50 ^a	k * n * 1A
51 ^a	k * n * 1A
52 ^a	k * 1A
53 ^a	k * n * 1A
54 ^a	k * 1A
55 ^a	k * 1A
56 ^a	1R
Total	$k * (18 + 4m + 11n + (n*m)) * \log(n)$
O estimate	$O(k * n^2 * \log(n))$

Subtitle:

- **A:** Assignment
- **k:** Meeting points.
- **n:** Number of Vertices
- **m:** Number of Edges
- **C:** Comparison
- **R:** Return