

COMPLEXITY OF US13

kruskal algorithm

```

procedure Kruskal(sortedGraphEdges: array of Edge, verticesGraph: array of
Vertex) returns Graph
  A := new Graph()
  parent := new array[verticesGraph.size()]
  rank := new array[verticesGraph.size()]

  for i := 0 to verticesGraph.size() - 1 do
    parent[i] := i
    rank[i] := 0
  end for

  nA := 0
  i := 0

  while nA < verticesGraph.size() - 1 and i < sortedGraphEdges.size() do
    e := sortedGraphEdges[i]
    u := indexOf(verticesGraph, e.getOrigin())
    v := indexOf(verticesGraph, e.getDestiny())
    rootU := find(u, parent)
    rootV := find(v, parent)

    if rootU != rootV then
      A.addEdge(e)
      union(rootU, rootV, parent, rank)
      nA := nA + 1
    end if

    i := i + 1
  end while

  return A

```

find method

```

procedure find(vertex: integer, parent: array of integer) returns integer
  if parent[vertex] != vertex then
    parent[vertex] := find(parent[vertex], parent)
  end if
  return parent[vertex]

```

addEdge method**union method**

```

procedure union(rootU: integer, rootV: integer, parent: array of integer,
rank: array of integer)
  if rank[rootU] > rank[rootV] then
    parent[rootV] := rootU
  else if rank[rootU] < rank[rootV] then
    parent[rootU] := rootV
  else
    parent[rootV] := rootU
    rank[rootU] := rank[rootU] + 1
  end if

```

Lines	Kruskal
1 ^a	1A
2 ^a	1A
3 ^a	1A
4 ^a	$(n + 1) A$
5 ^a	$n * 1A$
6 ^a	$n * 1A$
7 ^a	1A
8 ^a	1A
9 ^a	$m * 2C$
10 ^a	$m * 1A$
11 ^a	$m * O(n)$
12 ^a	$m * O(n)$
13 ^a	$m * T(\text{find})$ (the complexity of this algorithm is analysed below)
14 ^a	$m * T(\text{find})$ (the complexity of this algorithm is analysed below)
15 ^a	$m * 1C$
16 ^a	$m * 1A$
17 ^a	$m * T(\text{union})$ (the complexity of this algorithm is analysed below)
18 ^a	$m * (1A + 1Op)$
19 ^a	$m * (1A m + 1Op)$
20 ^a	1R
Total	$9 + 2n + 3m$
O estimate	$O(m \log m)$

Subtitle:

- A: Assignment

- ***n***: Number of Vertices
- ***m***: Number of Edges
- **C**: Comparison
- **Op**: Arithmetic Operation
- **I**: Increment
- **R**: Return

COMPLEXITY OF US17

dijkstra algorithm

```

procedure Dijkstra(start: String, MP: String, edges: array of Edges)
    shortestPaths := new Graph()
    if MP is equal to start then
        print("The vertex is already te Metting Point!")
    else
        initialGraph := new Graph()
        for each Edge in the set of Edges do
            Add Edge to initialGraph
        end for
        vertices := get vertices from inicalGraph
        numVertices := vertices.size()
        dist := new array[numVertices]
        prev := new array[numVertices]
        visited := new array[numVertices]
        for each index i from 0 to numVertices - 1 do
            dist[i] := -1
            prev[i] := null
            visited[i] := false

            create a PriorityQueue named Queue with a comparator based on
the distance of vertices

            MPVertex := new Vertex(MP)
            startVertex := new Vertex(start)
            dist[vertices.indexOf(startVertex)] := 0
            queue.add(startVertex)
            while (not queue.isEmpty())
                u := queue.poll()
                uIndex := vertices.indexOf(u)
                visited[uIndex] := true

```

```

        neighbors := initialGraph.getVerticesConnectedTo(u)
        for each neighbor in neighbors
            vIndex := vertices.indexOf(neighbor)
            weight := initialGraph.getEdgeCost(u, neighbor)

            if (not visited[vIndex] and dist[uIndex] ≠ -1 and
(dist[vIndex] = -1 or dist[uIndex] + weight < dist[vIndex])) then
                queue.remove(neighbor)
                dist[vIndex] := dist[uIndex] + weight
                prev[vIndex] := u
                queue.add(neighbor)
            end if
        end for
    end while
    path := new ArrayList()
    at := MPVertex
    while (at ≠ null)
        path.add(at)
        at := prev[vertices.indexOf(at)]
    end while
    path.reverse()
    i := 0
    while (i < path.size() - 1)
        origin := path.get(i)
        destiny := path.get(i + 1)
        cost := initialGraph.getEdgeCost(origin, destiny)
        shortestPaths.addEdge(new Edge(origin, destiny, cost))
        i := i + 1
    end while
return shortestPaths

```

Lines	Dijkstra
1 ^a	1A
2 ^a	1C
3 ^a	1A
4 ^a	1A
5 ^a	1A
6 ^a	$m * 1A$
7 ^a	$m * 1A$
8 ^a	$n * 1A$
9 ^a	1A
10 ^a	$n * 1A$
11 ^a	$n * 1A$
12 ^a	$n * 1A$
13 ^a	$n * 1C$
14 ^a	$n * 1A$
15 ^a	$n * 1A$
16 ^a	$n * 1A$
17 ^a	1A
18 ^a	1A
19 ^a	1A
20 ^a	1A
21 ^a	$1A + \log(n)$
22 ^a	$n * \log(n)$
23 ^a	$n * \log(n)$
24 ^a	$n * 1A$
25 ^a	$n * 1A$
26 ^a	$n * 1A$
27 ^a	$m * 1C$
28 ^a	$m * 1A$
29 ^a	$m * 1A$
30 ^a	$m * 3C$
31 ^a	$m * \log(n)$
32 ^a	$m * 1A$
33 ^a	$m * 1A$
34 ^a	$m * \log(n)$
35 ^a	$m * 1A$
36 ^a	$m * 1A$
37 ^a	$n * \log(n)$
38 ^a	1A
39 ^a	1A
40 ^a	$n * 1C$
41 ^a	$n * 1A$
42 ^a	$n * 1A$
43 ^a	$n * 1A$

44 ^a	$n * 1A$
45 ^a	$1A$
46 ^a	$n * 1C$
47 ^a	$n * 1A$
48 ^a	$n * 1A$
49 ^a	$n * 1A$
50 ^a	$n * 1A$
51 ^a	$n * 1A$
52 ^a	$1A$
53 ^a	$1A$
54 ^a	$1R$
Total	$9+2n+3m+2m\log(n)+6n+5m+5n+1$
<i>O</i> estimate	$O(m \log(n))$

Subtitle:

- **A:** Assignment
- **n:** Number of Vertices
- **m:** Number of Edges
- **C:** Comparison
- **R:** Return

COMPLEXITY OF US18

```

        neighbors := initialGraph.getVerticesConnectedTo(u)
        for each neighbor in neighbors do
            vIndex := vertices.indexOf(neighbor)
            weight := initialGraph.getEdgeCost(u, neighbor)
            if not visited[vIndex] and dist[uIndex] != -1 and
(dist[vIndex] = -1 or dist[uIndex] + weight < dist[vIndex]) then
                queue.remove(neighbor)
                dist[vIndex] := dist[uIndex] + weight
                prev[vIndex] := u
                queue.add(neighbor)
            end if
        end for
    end while
    path := new ArrayList<>()
    for Vertex at := MPVertex; at != null; at :=
prev[vertices.indexOf(at)] do
        path.add(at)
    end for
    path.reverse()
    shortestPath := new Graph()
    for each index i from 0 to path.size() - 2 do
        Vertex origin := path.get(i)
        Vertex destiny := path.get(i + 1)
        int cost := initialGraph.getEdgeCost(origin, destiny)
        shortestPath.addEdge(new Edge(origin, destiny, cost))
    end for
    shortestPaths.add(shortestPath)
end if
end for
return shortestPaths

```


Lines	Dijkstra
1 ^a	1A
2 ^a	$k * 1C$
3 ^a	$k * 1C$
4 ^a	$k * 1A$
5 ^a	$k * 1A$
6 ^a	$k * 1A$
7 ^a	$k * m * 1A$
8 ^a	$k * m * 1A$
9 ^a	$k * n * 1A$
10 ^a	$k * 1A$
11 ^a	$k * n * 1A$
12 ^a	$k * n * 1A$
13 ^a	$k * n * 1A$
14 ^a	$k * n * 1C$
15 ^a	$k * n * 1A$
16 ^a	$k * n * 1A$
17 ^a	$k * n * 1A$
18 ^a	$k * 1A$
19 ^a	$k * 1A$
20 ^a	$k * 1A$
21 ^a	$k * 1A$
22 ^a	$k * (1A + \log(n))$
23 ^a	$k * n * \log(n)$
24 ^a	$k * n * \log(n)$
25 ^a	$k * n * 1A$
26 ^a	$k * n * 1A$
27 ^a	$k * n * 1A$
28 ^a	$k * m * 1C$
29 ^a	$k * m * 1A$
30 ^a	$k * m * 1A$
31 ^a	$k * n * 3C$
32 ^a	$k * m * \log(n)$
33 ^a	$k * m * 1A$
34 ^a	$k * m * 1A$
35 ^a	$k * m * \log(n)$
36 ^a	$k * m * 1A$
37 ^a	$k * m * 1A$
38 ^a	$k * n * \log(n)$
39 ^a	$k * 1A$
40 ^a	$k * 1A$

41 ^a	$k * n * 1C$
42 ^a	$k * n * 1A$
43 ^a	$k * n * 1A$
44 ^a	$k * n * 1A$
45 ^a	$k * n * 1A$
46 ^a	$k * 1A$
47 ^a	$k * n * 1C$
48 ^a	$k * n * 1A$
49 ^a	$k * n * 1A$
50 ^a	$k * n * 1A$
51 ^a	$k * n * 1A$
52 ^a	$k * 1A$
53 ^a	$k * n * 1A$
54 ^a	$k * 1A$
55 ^a	$k * 1A$
56 ^a	$1R$
Total	$k * (18 + 4m + 11n + (n*m)) * \log(n)$
<i>O</i> estimate	$O(k * n^2 * \log(n))$

Subtitle:

- **A:** Assignment
- **k:** Meeting points.
- **n:** Number of Vertices
- **m:** Number of Edges
- **C:** Comparison
- **R:** Return