

*Great minds discuss ideas;  
Average minds discuss events;  
Small minds discuss people.*

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# TCS-503: Design and Analysis of Algorithms

## Heapsort

## Unit I: Syllabus

### ○ Introduction:

- Algorithms
- Analysis of Algorithms
- Growth of Functions
- Master's Theorem
- Designing of Algorithms

# Unit I: Syllabus

- **Sorting and Order Statistics**
  - **Heap Sort**
  - Quick Sort
  - Sorting in Linear Time
    - Counting Sort
    - Bucket Sort
    - Radix Sort
  - Medians and Order Statistics

## Sorting Revisited

### Bubble Sort

Design approach: incremental

Sorts in place: Yes

Running time:  $\Theta(n^2)$

### Selection sort

Design approach: incremental

Sorts in place: Yes

Running time:  $\Theta(n^2)$

### Insertion sort

Design approach: incremental

Sorts in place: Yes

Running time:  $\Theta(n^2)$

### Merge Sort

Design approach: Divide & Conquer

Sorts in place: No

Running time:  $\Theta(n \lg n)$

## Sorting Revisited

What is the advantage of merge sort?

Answer: good worst-case running time  $O(n \lg n)$

Conceptually easy, Divide-and-Conquer

What is the advantage of insertion sort?

Answer: sorts in place-

only a constant number of array elements are stored outside the input array at any time.

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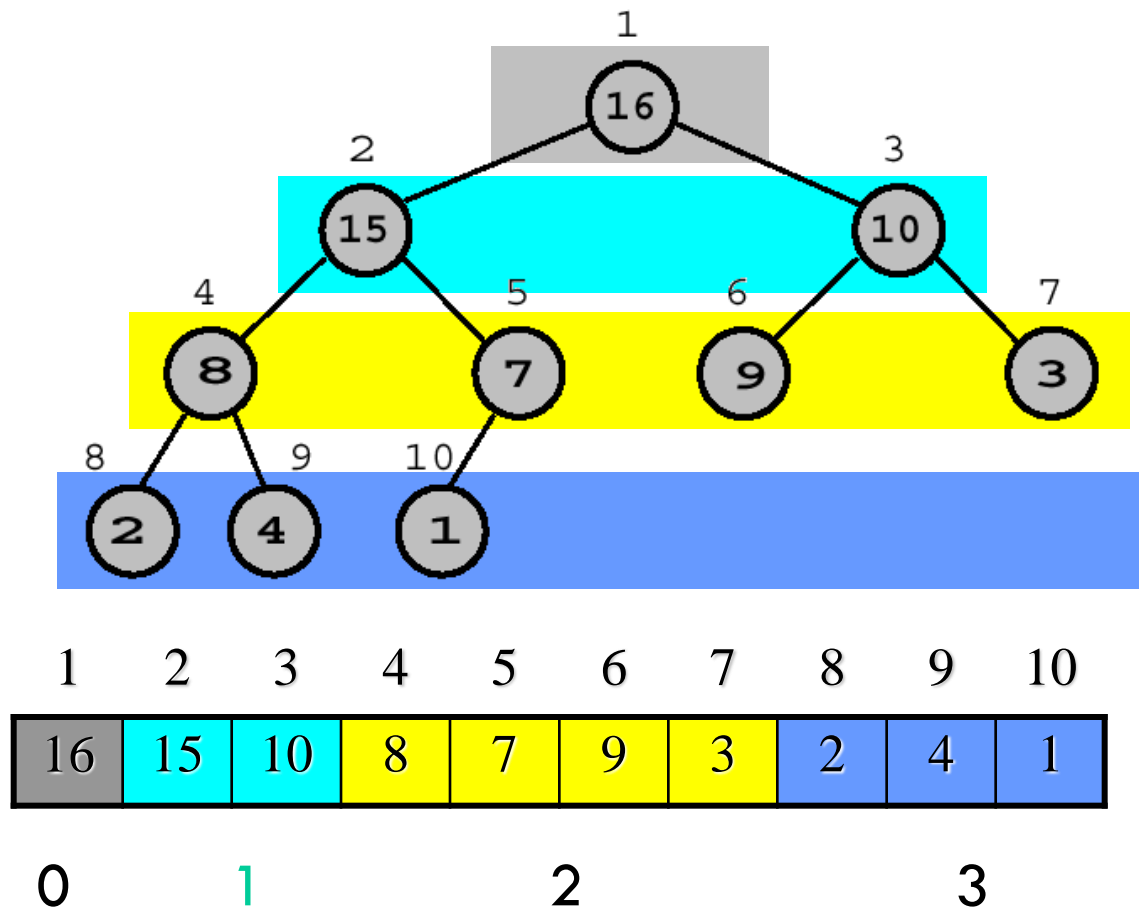
Next on the agenda:

*Heap sort*

Combines advantages of both previous algorithms

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### Heap





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## Heap

A heap can be seen as a complete binary tree with the following properties:

If a Node is at position  $i$ , then

Its left child will be at position  $2i$

Its right child will be at position  $2i+1$

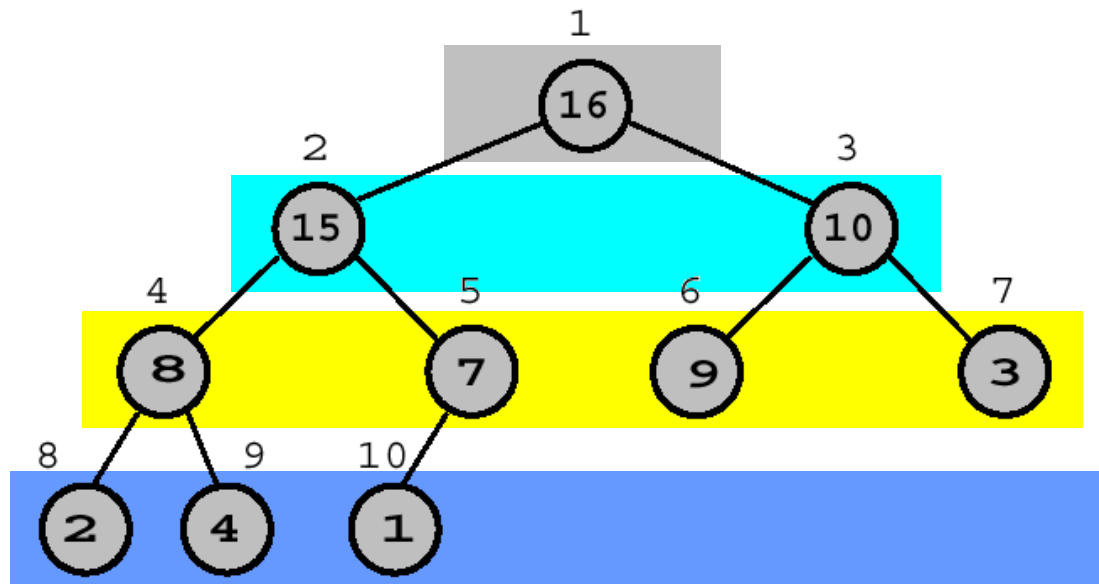
And

Its parent will be at position  $i/2$

## Heap Property

Heaps also satisfy the *heap property*:

$$A[\text{Parent}(i)] \geq A[i] \text{ for all nodes } i > 1$$



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## Heap Sort Algorithm

has

two supporting algorithms:

**MAXHEAPIFY()**

and

**BUILD-MAX-HEAP()**

## MAX-HEAPIFY() Algorithm

MAX-HEAPIFY(A, i)

l=LEFT(i)

r=RIGHT(i)

If  $l \leq \text{heap-size}[A]$  and  $A[l] > A[i]$  then

largest  $\leftarrow l$

else

largest  $\leftarrow r$

If  $r \leq \text{heap-size}[A]$  and  $A[r] > A[\text{largest}]$  then

largest  $\leftarrow r$

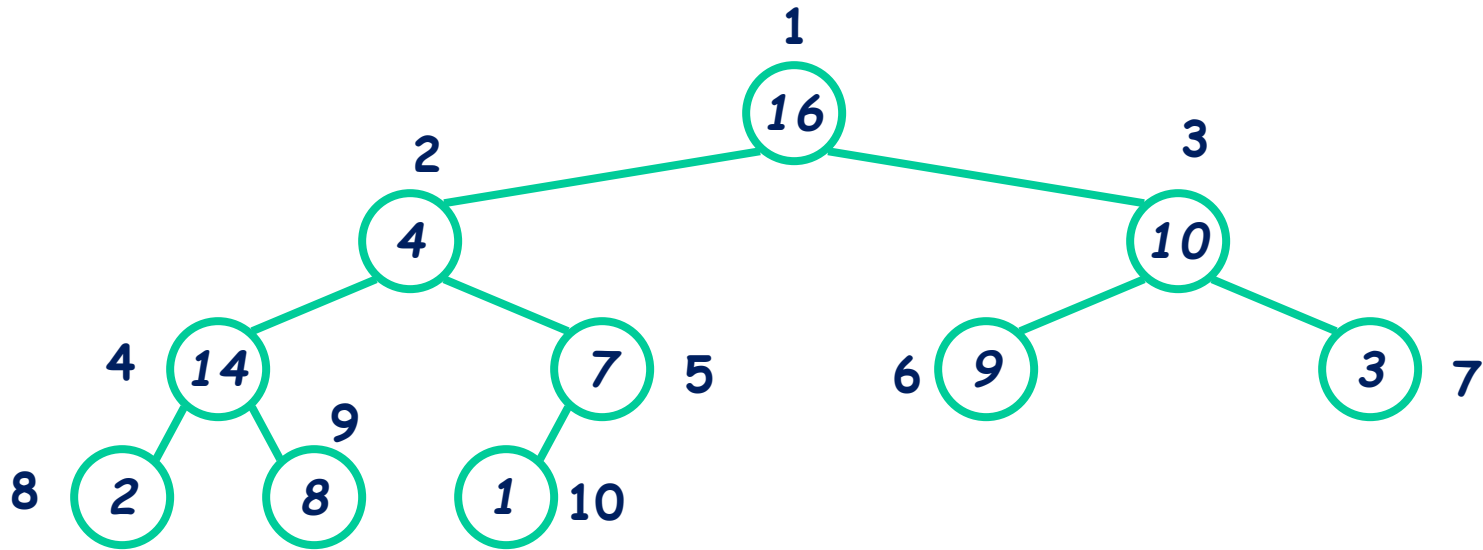
If largest  $\neq i$  then

Exchange  $A[i] \leftrightarrow A[\text{largest}]$

MAX-HEAPIFY(A, largest)

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## MAX-HEAPIFY(A, i) Algorithm: Example



$A =$

16	4	10	14	7	9	3	2	8	1
1	2	3	4	5	6	7	8	9	10

## MAX-HEAPIFY(A,2) Algorithm: Example

$l \leq \text{Heapsize}(A)$ :

$2 \leq 10$ : Yes

And  $A[l] > A[i]$ :

$14 > 4$ : Yes then  
 $\text{largest} = 4$

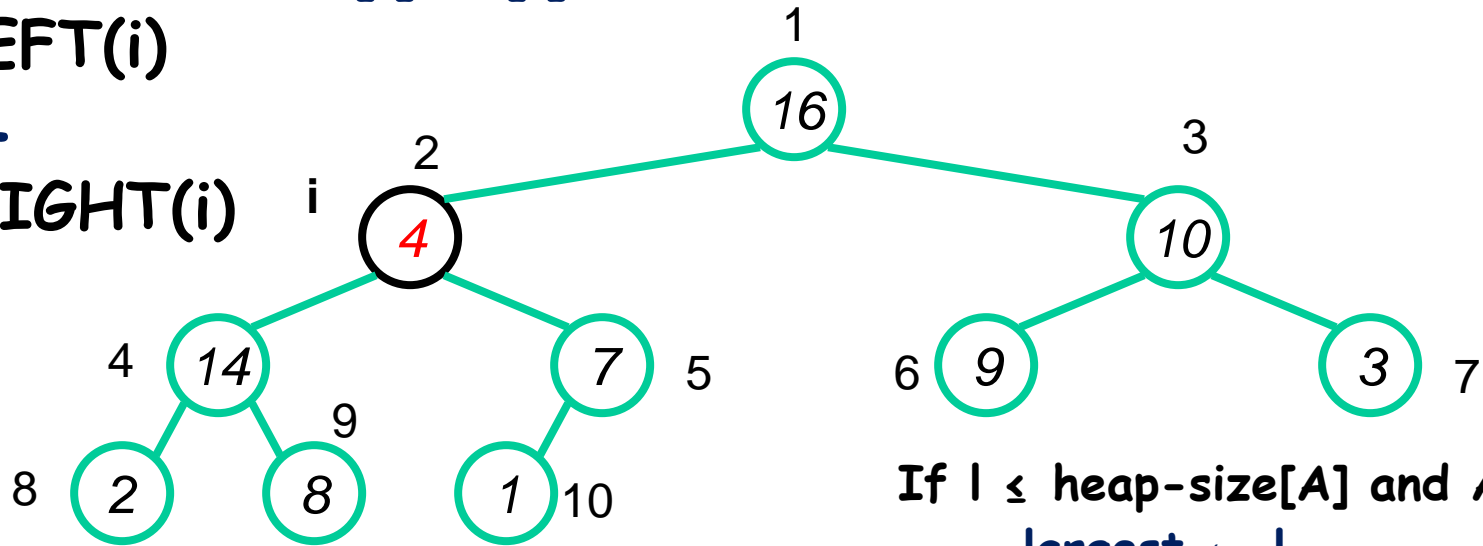
$i = 2$

$l = \text{LEFT}(i)$

$l = 4$

$r = \text{RIGHT}(i)$

$r = 5$

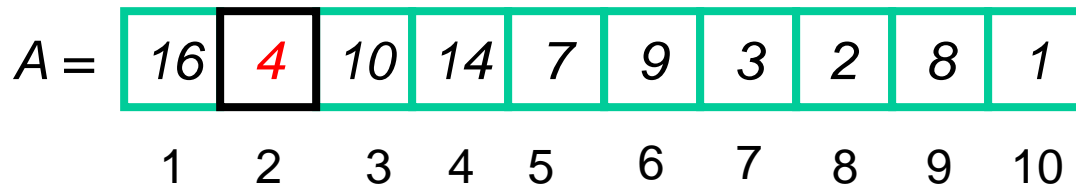


If  $l \leq \text{heap-size}[A]$  and  $A[l] > A[i]$

$\text{largest} \leftarrow l$

else

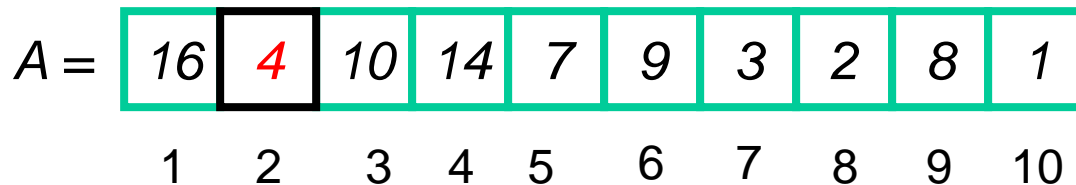
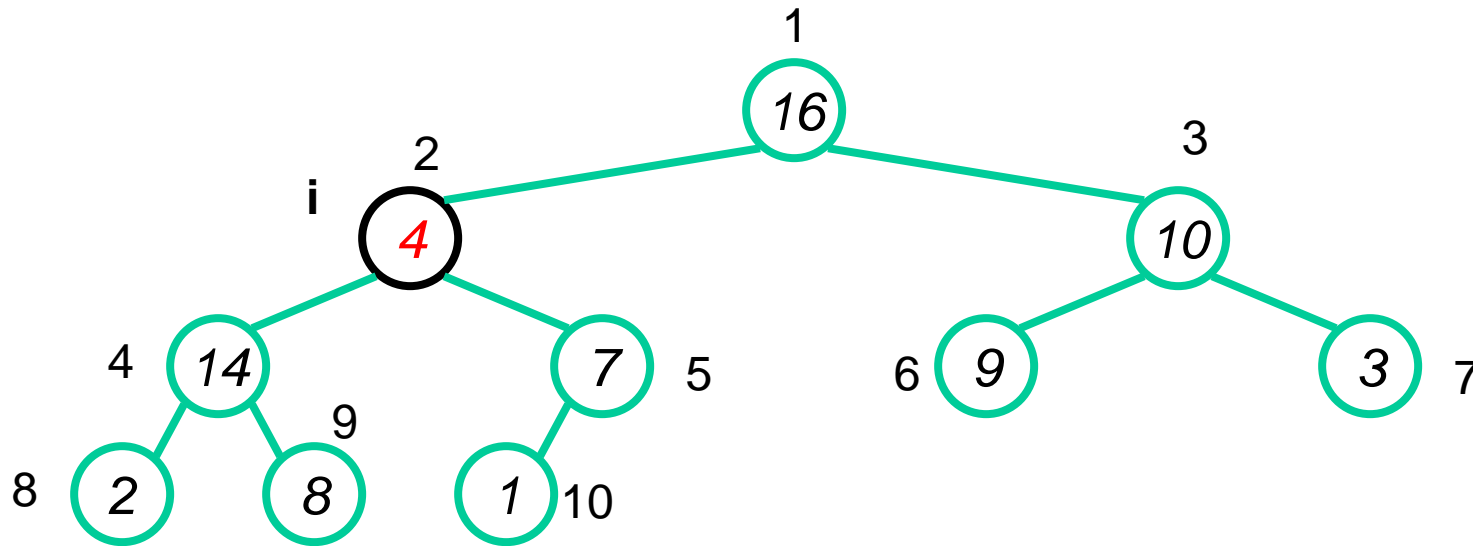
$\text{largest} \leftarrow r$



## MAX-HEAPIFY(A,2) Algorithm: Example

$i=2$   
 $l=4$ ,  
 $r=5$

$r \leq \text{Heapsize}(A)$ :  $7 > 14$ : No  
 $5 \leq 10$ : Yes  
And  $A[r] > A[\text{largest}]$ :



## MAX-HEAPIFY(A,2) Algorithm: Example

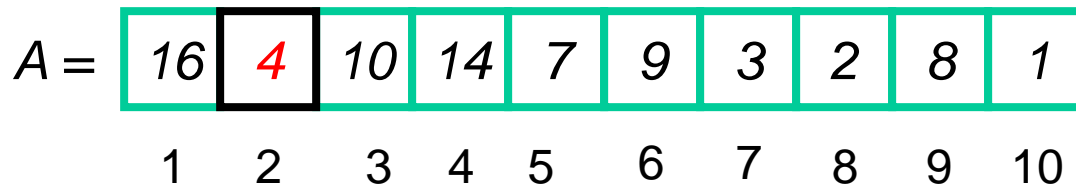
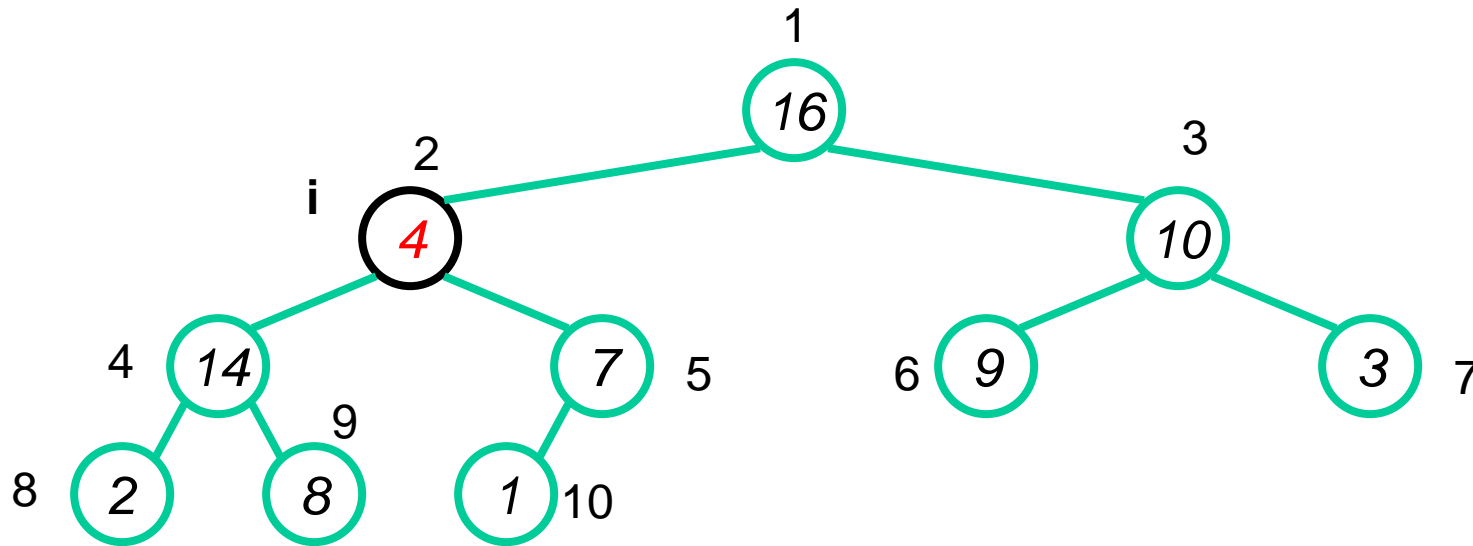
$i=2$   
 $l=4,$   
 $r=5$

largest  $\neq i$ :

$4 \neq 2$ : Yes Then

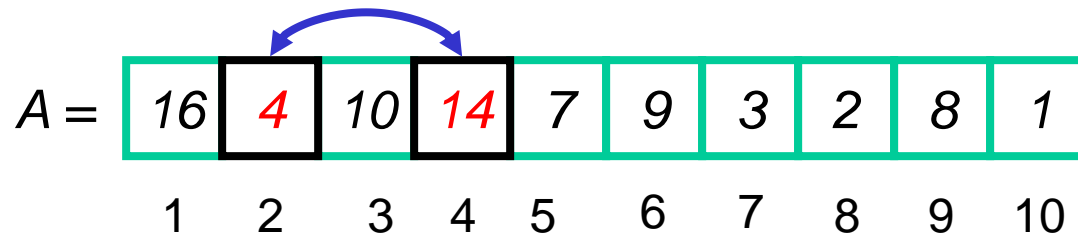
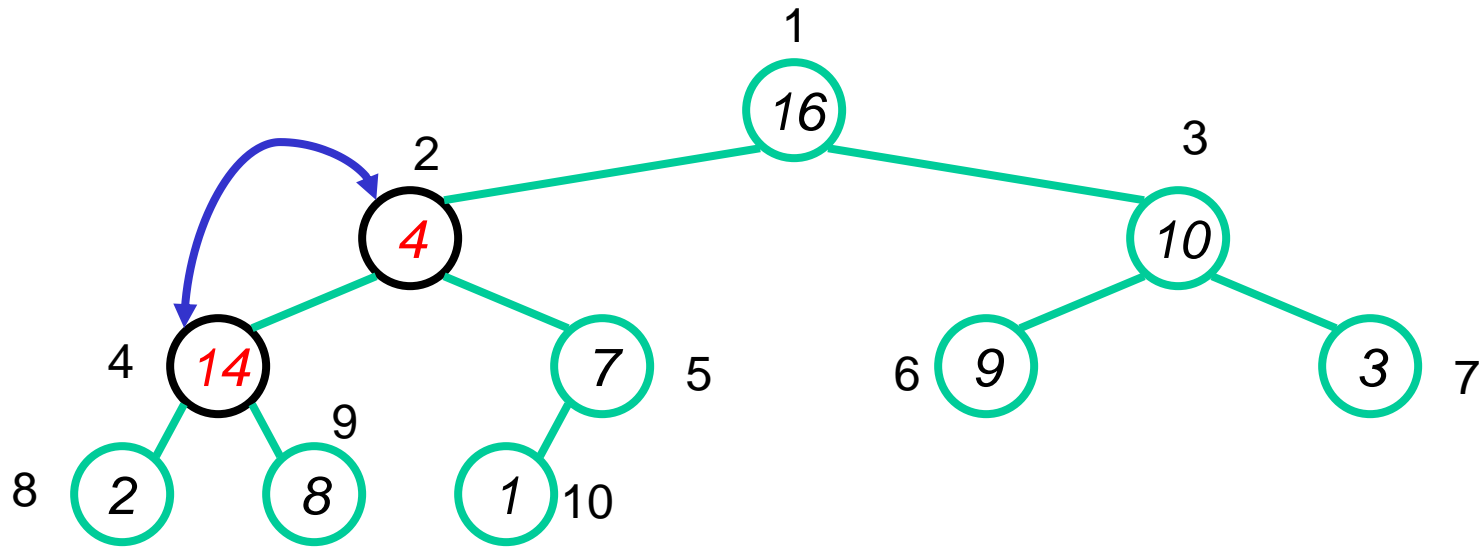
Exchange  $A[2]$  and  $A[4]$

Exchange  $A[i]$  and  $A[\text{largest}]$ :



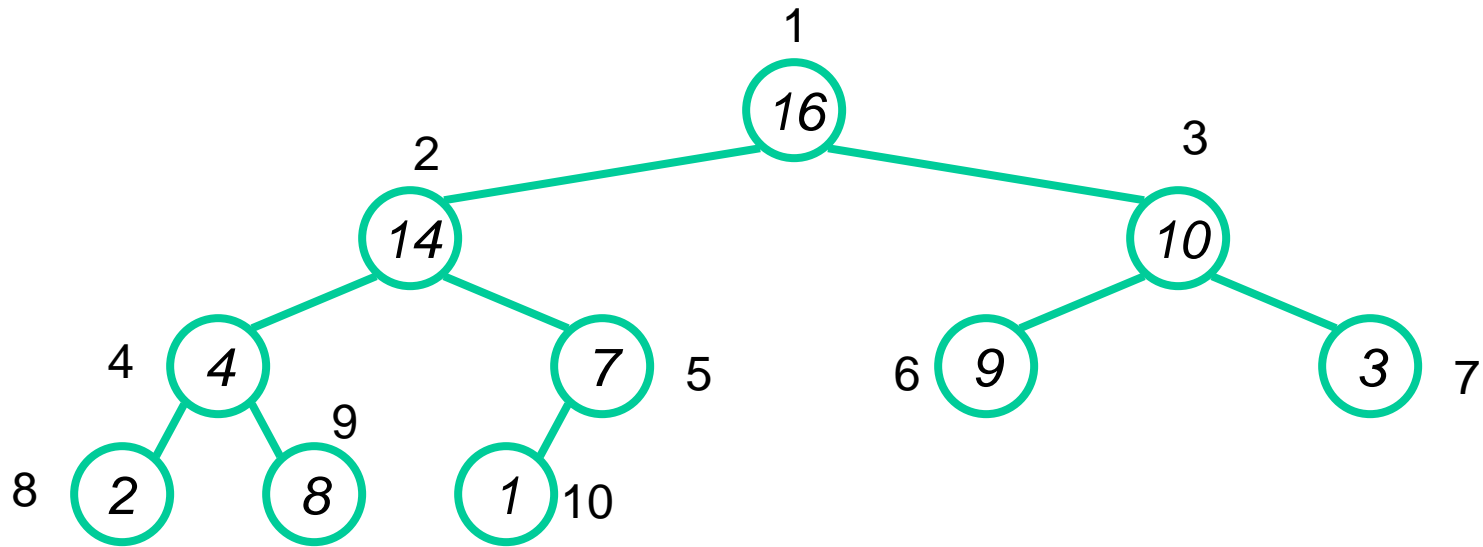


## MAX-HEAPIFY(A,2) Algorithm: Example



# MAX-HEAPIFY(A, largest) Algorithm: Example

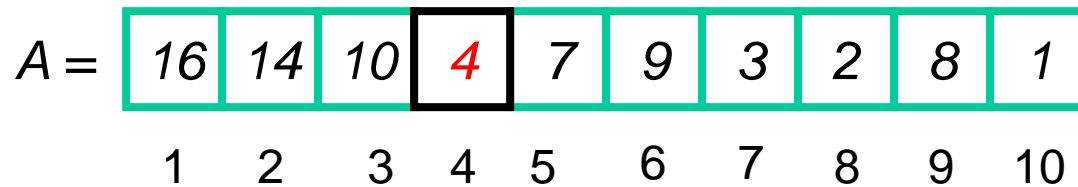
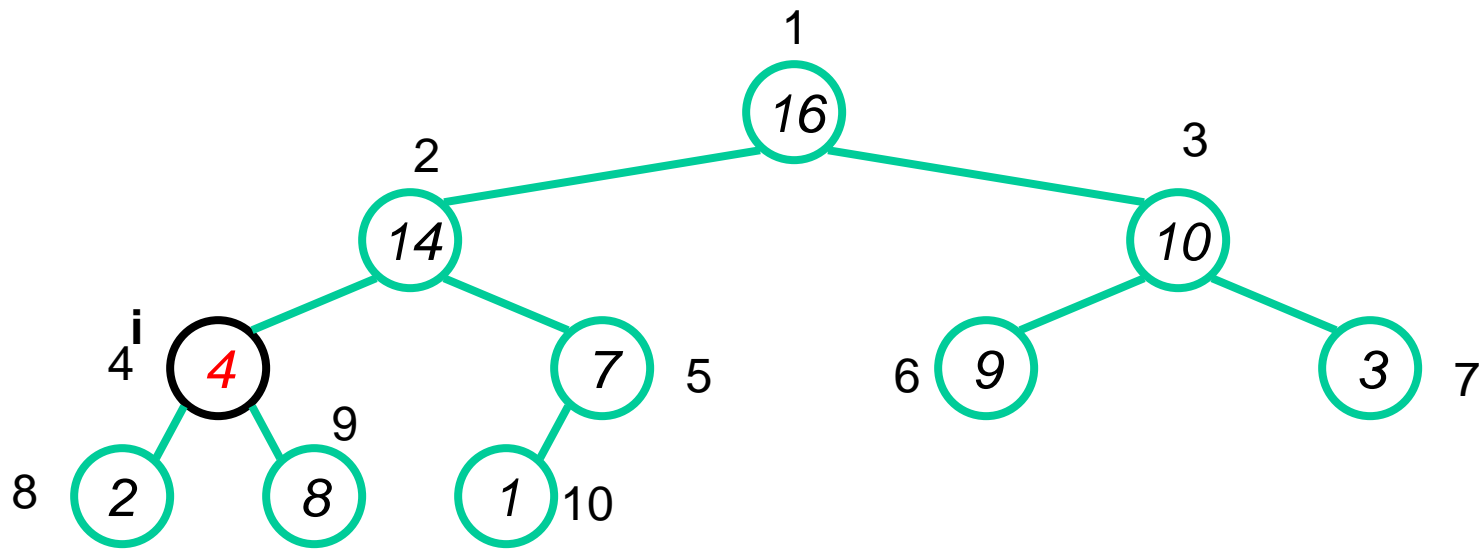
## MAX-HEAPIFY(A, 4)



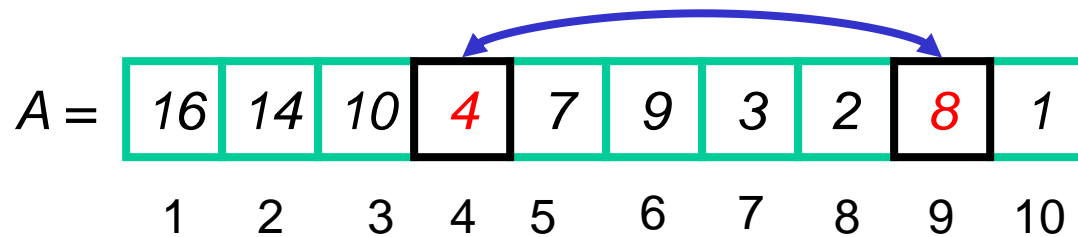
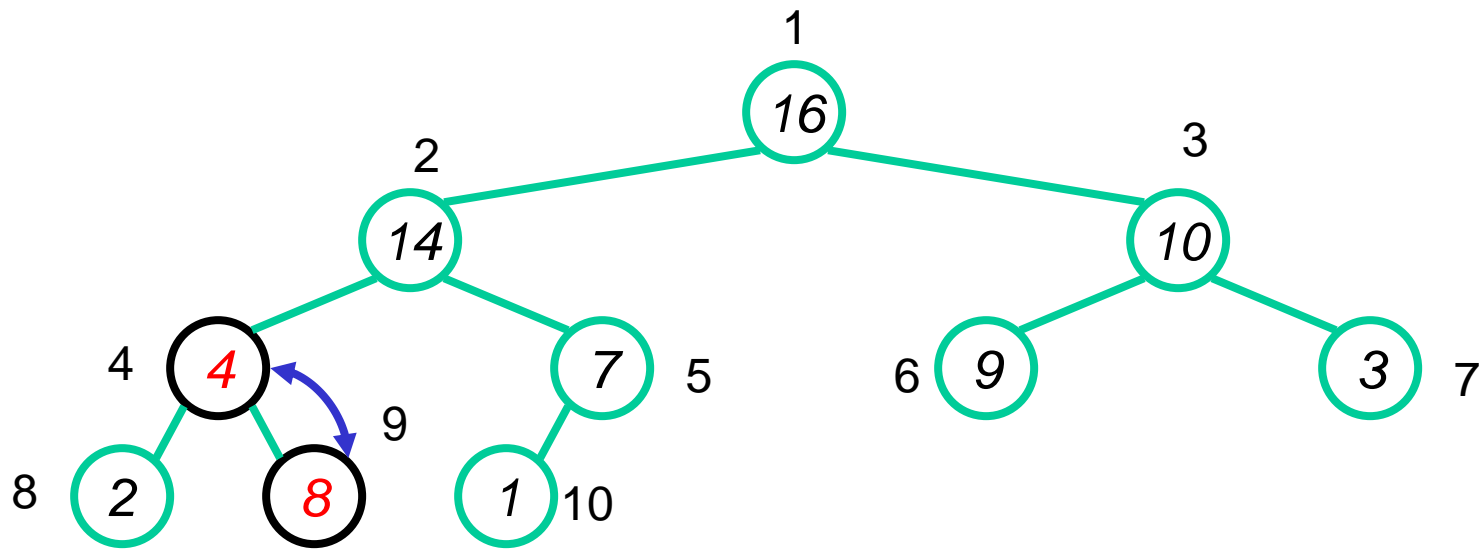
A =

16	14	10	4	7	9	3	2	8	1
1	2	3	4	5	6	7	8	9	10

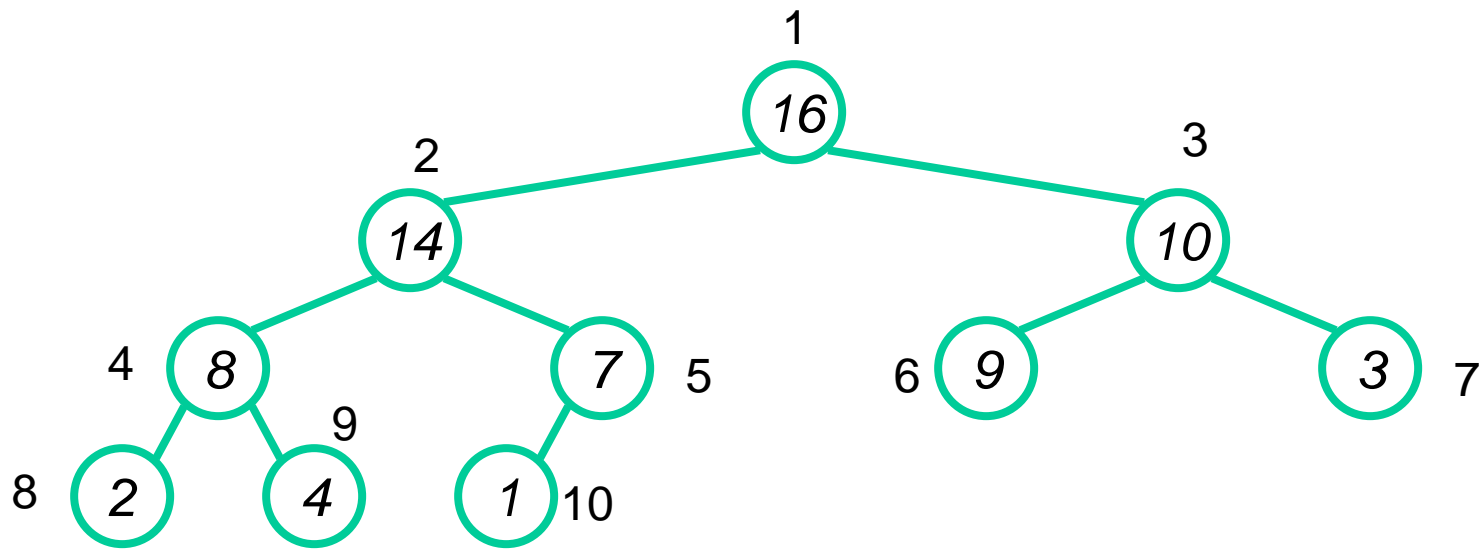
## MAX-HEAPIFY(A, 4)



## MAX-HEAPIFY(A, 4)



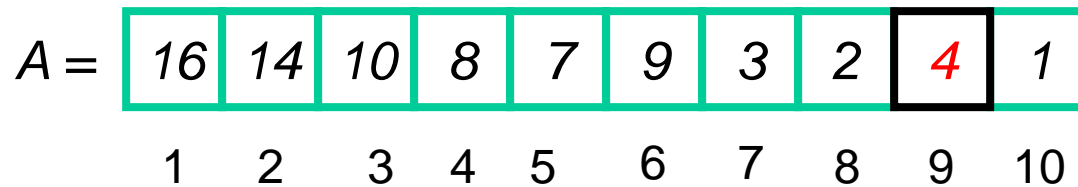
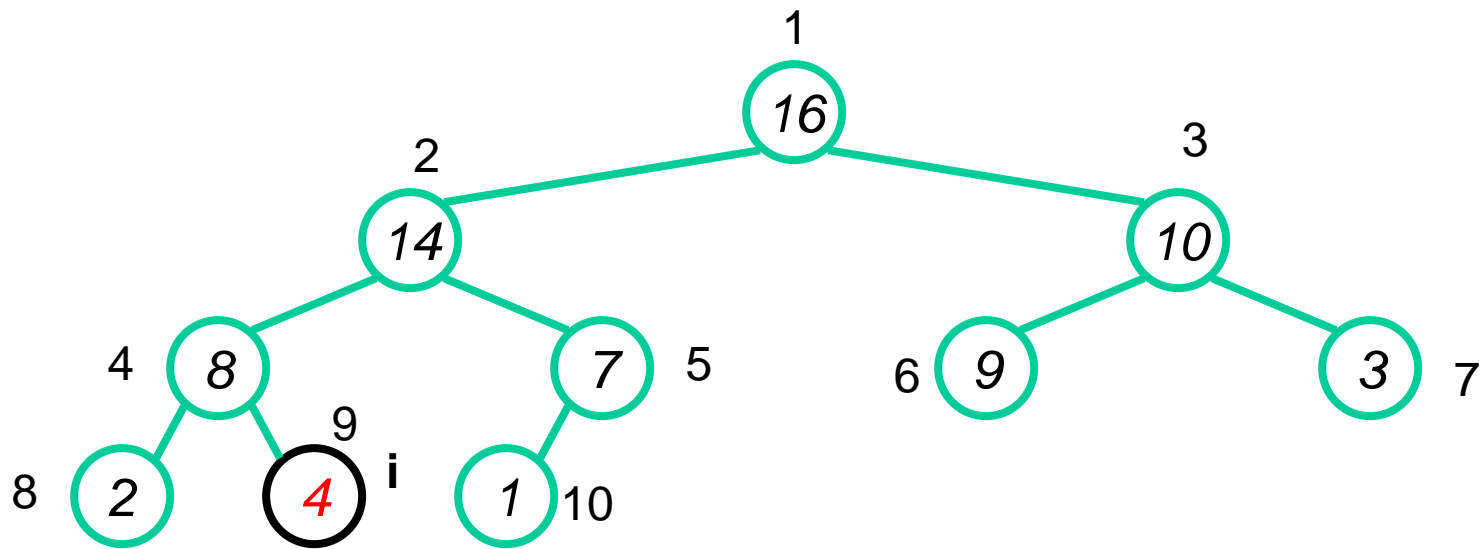
## MAX-HEAPIFY(A, 4)



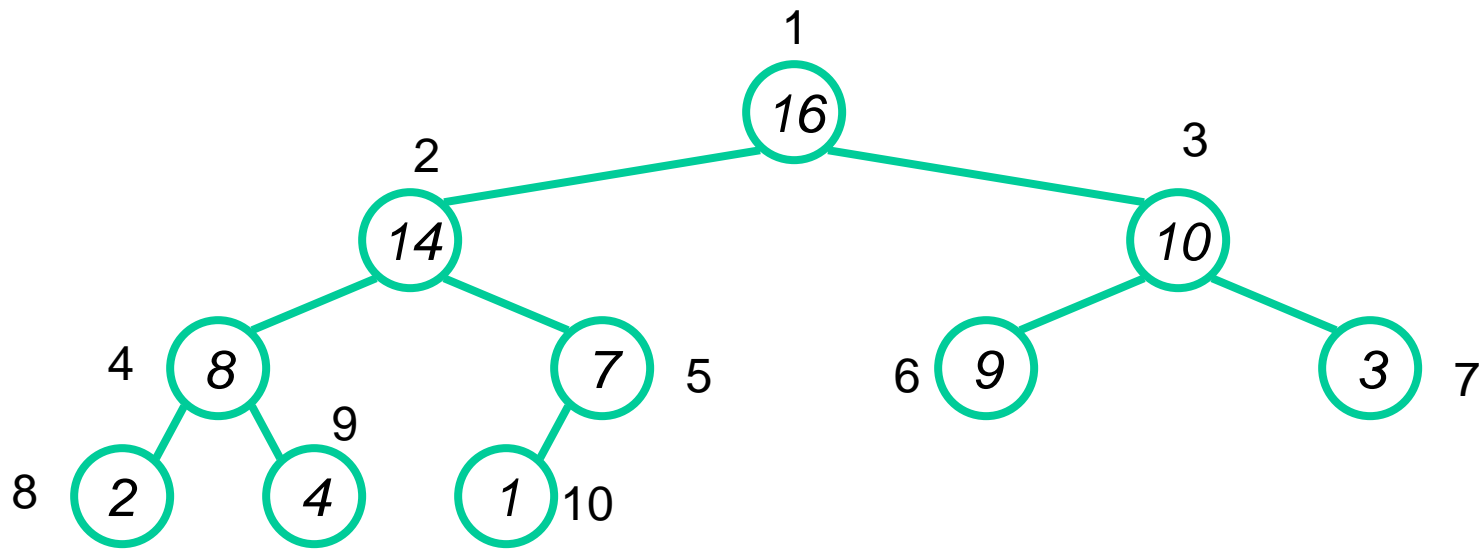
A =

16	14	10	8	7	9	3	2	4	1
1	2	3	4	5	6	7	8	9	10

## MAX-HEAPIFY(A, 9)



## MAX-HEAPIFY(A, 9)



A =

16	14	10	8	7	9	3	2	4	1
1	2	3	4	5	6	7	8	9	10

## BUILD-MAX-HEAP()

### BUILD-MAX-HEAP(A)

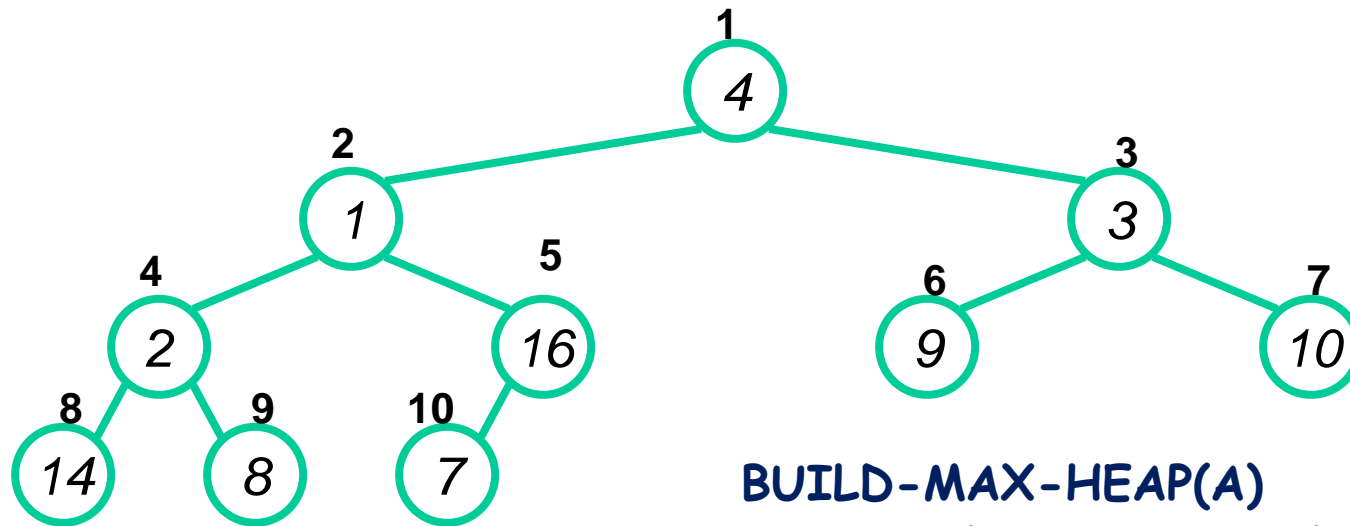
- 1 heap-size[A]  $\leftarrow$  length[A]
- 2 For  $i \leftarrow \lfloor \text{length}[A]/2 \rfloor$  downto 1
- 3     do MAX-HEAPIFY(A, i)



## BUILD-MAX-HEAP() Example

$A = \{4, 1, 3, 2, 16, 9, 10, 14, 8, 7\}$

$n=10, n/2=5$



**BUILD-MAX-HEAP(A)**

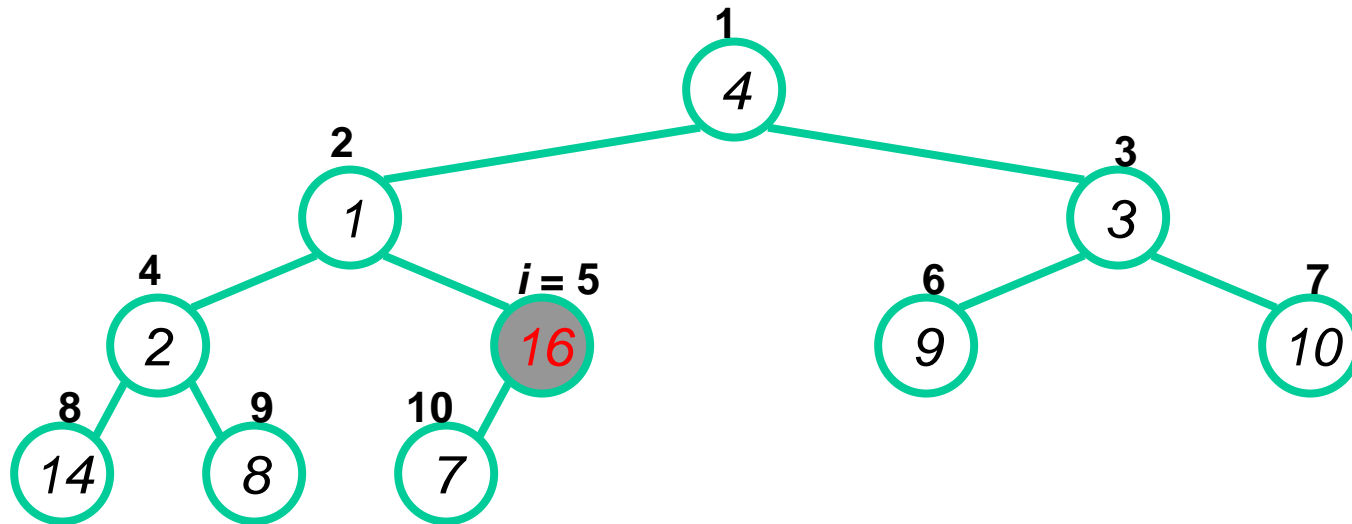
- 1 heap-size[A] ← length[A]
- 2 For  $i \leftarrow \lfloor \text{length}[A]/2 \rfloor$  downto 1
- 3     do MAX-HEAPIFY(A, i)

## BUILD-MAX-HEAP() Example

$A = \{4, 1, 3, 2, 16, 9, 10, 14, 8, 7\}$

for  $i=5$  to 1

do **MAX-HEAPIFY(A,i)**

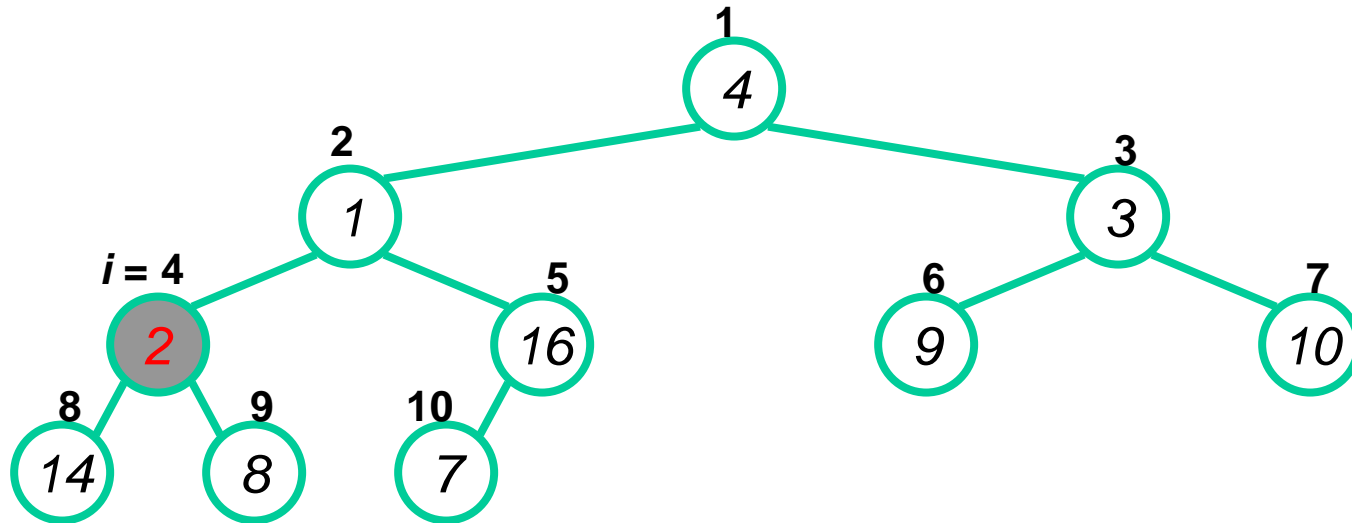


## BUILD-MAX-HEAP() Example

$A = \{4, 1, 3, 2, 16, 9, 10, 14, 8, 7\}$

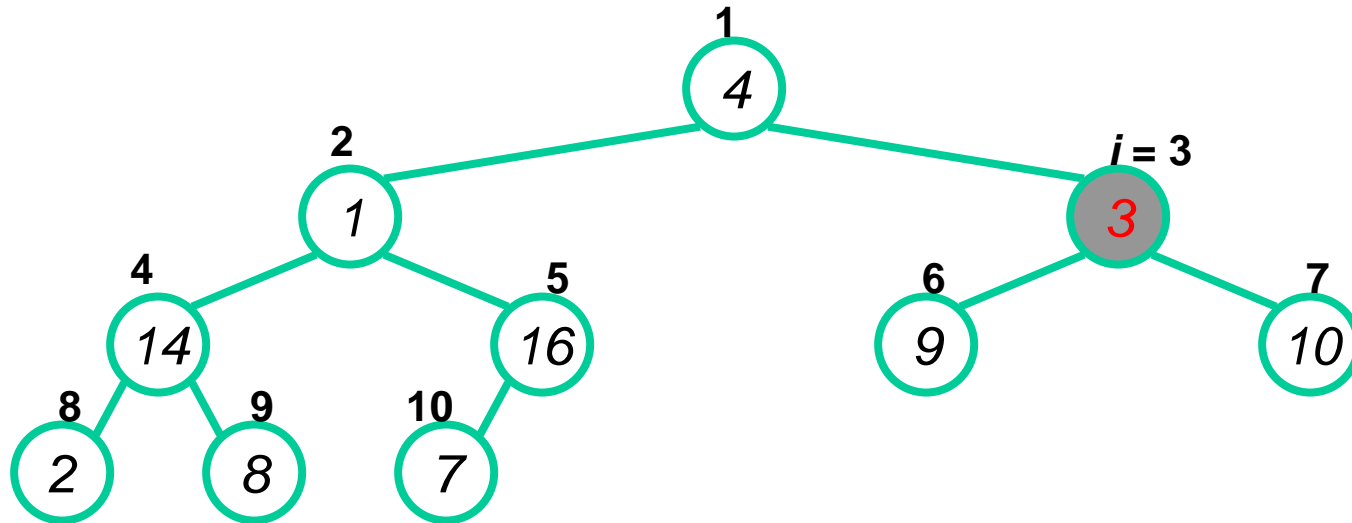
for  $i=5$  to  $1$

do **MAX-HEAPIFY**( $A, i$ )



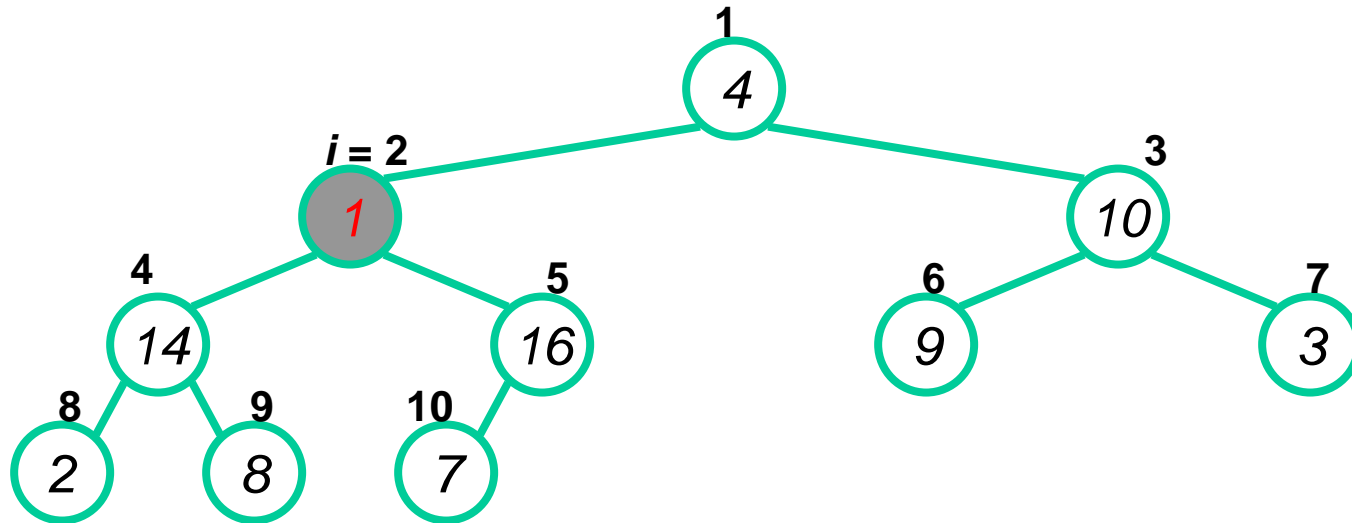
## BUILD-MAX-HEAP() Example

$A = \{4, 1, 3, 14, 16, 9, 10, 2, 8, 7\}$



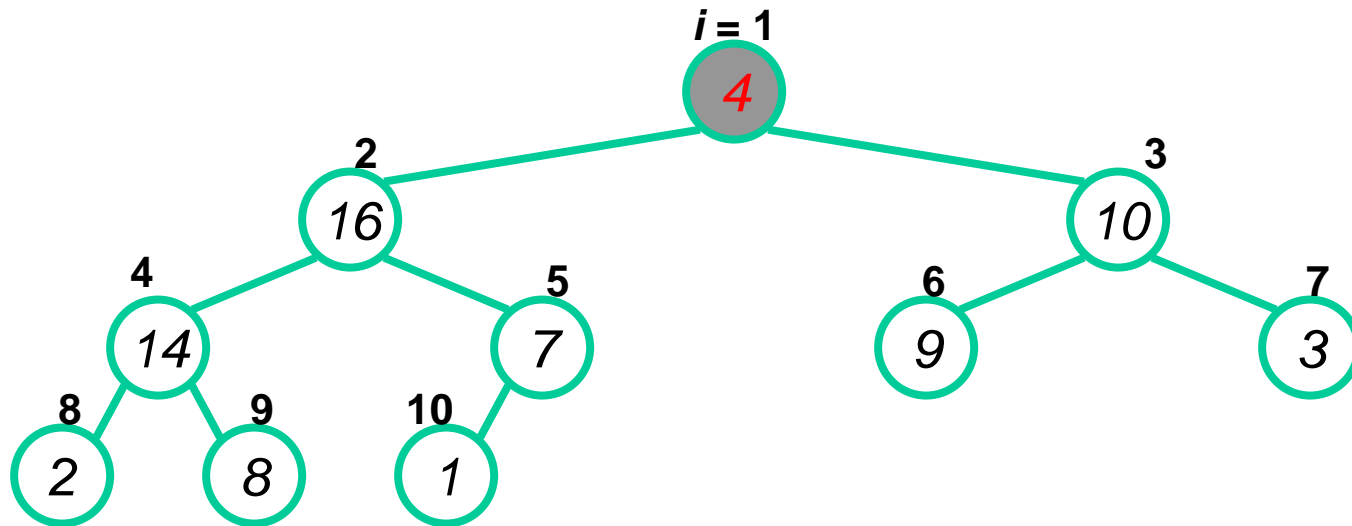
## BUILD-MAX-HEAP() Example

$A = \{4, 1, 10, 14, 16, 9, 3, 2, 8, 7\}$



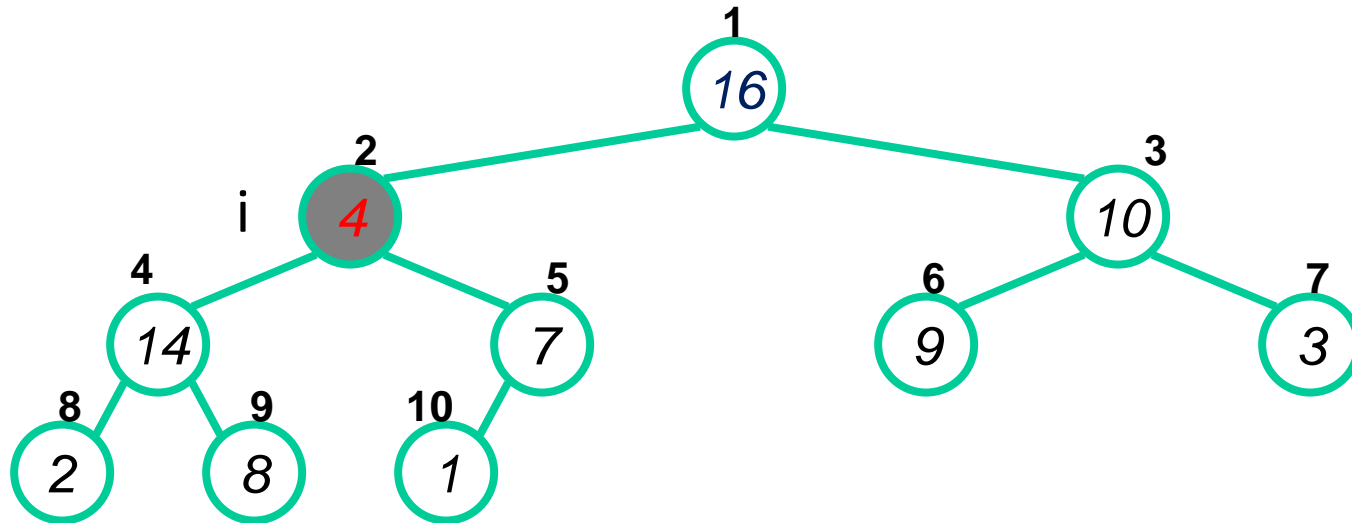
## BUILD-MAX-HEAP() Example

$A = \{4, 16, 10, 14, 7, 9, 3, 2, 8, 1\}$



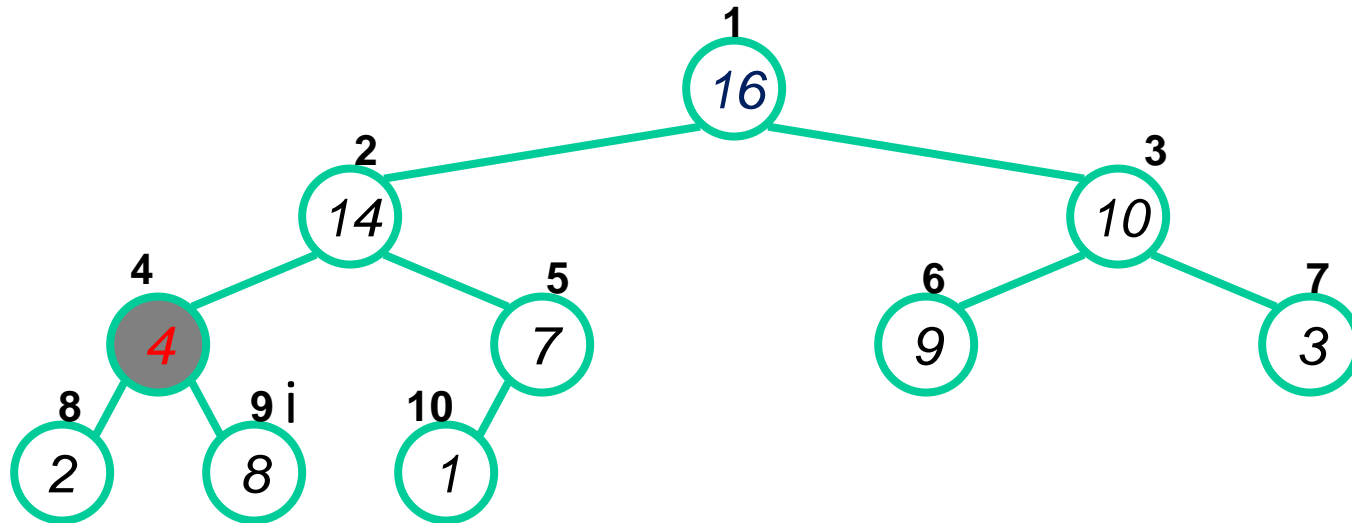
## BUILD-MAX-HEAP() Example

$A = \{16, 4, 10, 14, 7, 9, 3, 2, 8, 1\}$



## BUILD-MAX-HEAP() Example

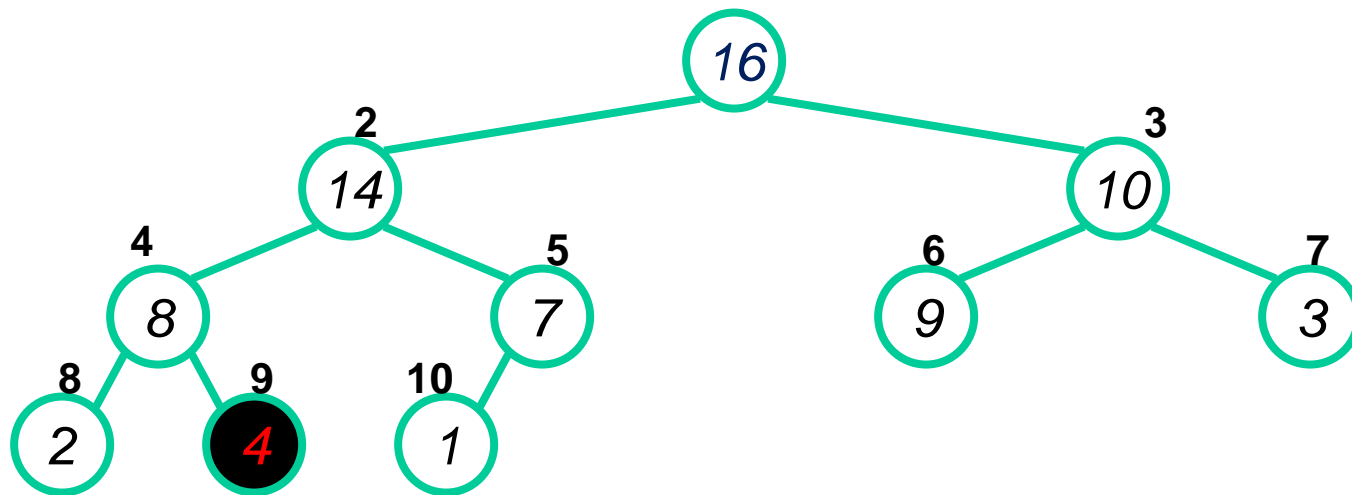
$A = \{16, 14, 10, 4, 7, 9, 3, 2, 8, 1\}$





## BUILD-MAX-HEAP() Example

$A = \{16, 14, 10, 8, 7, 9, 3, 2, 4, 1\}$



## BUILD-MAX-HEAP() Example

$A = \{16, 14, 10, 8, 7, 9, 3, 2, 4, 1\}$

